

# MATH97110: Numerical Methods for Finance

## Topic 2: Path-dependent Options Pricing with Lattice Methods

Imperial College London

2021-2022

# Overview

- ① Introduction to path-dependent options pricing
- ② Forward shooting grid method
- ③ Two specific examples:
  - ▶ Lookback option
  - ▶ Asian option

# Path-dependent options

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- Examples:

Floating strike lookback put option:  $\left( \max_{i=0, \dots, N} S_i - S_N \right)$

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- Suppose the payoff is paid to the option holder at the terminal time  $N$ , then the time-zero fair option price is

$$e^{-rN\Delta t} \mathbb{E}_{\mathbb{Q}}[g(S_0, S_1, \dots, S_N)] \quad (2.1)$$

where  $r$  is the (annualised) interest rate and  $\Delta t$  is the calendar time (in year) corresponding to each period of the tree

# A naive approach to price a path-dependent option

- A direct way to compute the expectation in (2.1) is to loop through all possible realisation of paths
- Example: Pricing a fixed strike Asian call option in a 3-period binomial tree with  $S_0 = 100$ ,  $K = 100$ ,  $u = 2$ ,  $d = 0.5$  and  $q = 0.6$ , and interest rate  $r = 0$

Path scenario	$S_1$	$S_2$	$S_3$	probability	payoff = $(\frac{1}{4} \sum_{i=0}^3 S_i - K)^+$
up up up	200	400	800	$q^3 = 0.216$	275
up up down	200	400	200	$q^2(1-q) = 0.144$	125
up down up	200	100	200	$q^2(1-q) = 0.144$	50
up down down	200	100	50	$q(1-q)^2 = 0.096$	12.5
down up up	50	100	200	$q^2(1-q) = 0.144$	12.5
down up down	50	100	50	$q(1-q)^2 = 0.144$	0
down down up	50	25	50	$q(1-q)^2 = 0.096$	0
down down down	50	25	12.5	$(1-q)^3 = 0.064$	0

and the fair option price is  $e^{-rN\Delta t} (\sum_i \text{probability}_i \times \text{payoff}_i) = 87.6$

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- In an  $N$ -period binomial tree, there are  $2^N$  possible realisations of the stock price path  $(S_0, S_1, \dots, S_N)$  and looping through all possible scenarios is not feasible for large  $N$



# Pricing path-dependent option: introducing an auxiliary process

- For many types of path-dependent options, the payoff function  $g(S_0, S_1, \dots, S_N)$  can be simplified by introducing a new **auxiliary state process**  $F = (F_n)_{n=0, \dots, N}$ , and then the payoff function can be rewritten as  $g(S_N, F_N)$

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- Example: Floating strike lookback put option  $\left( \max_{i=0,1,\dots,N} S_i - S_N \right)$ 
  - ▶ Define  $M_n := \max_{i=0,1,\dots,n} S_i$  which represents the running maximum of the stock price process between time zero and time  $n$ , then the payoff function becomes

$$g(S_N, M_N) = M_N - S_N$$

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- Define  $A_n := \frac{1}{n+1} \sum_{i=0}^n S_i$  which represents the running average of the stock price process between time zero and time  $n$ , then the payoff function becomes

$$g(A_N) = (A_N - K)^+$$

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- Now the path-dependent nature of the payoff disappears, but this comes at the cost of analysing an additional process

# The pricing algorithm for path-dependent options

- If the payoff can be rewritten as  $g(S_N, F_N)$  for some suitably chosen auxiliary process  $F$ , then the fair price of the path-dependent option at time  $n$  becomes

$$V^n := e^{-r(N-n)\Delta t} \mathbb{E}_{\mathbb{Q}} \left[ g(S_N, F_N) \middle| \mathcal{F}_n \right]$$

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- The backward induction algorithm for computing option price remains largely the same as in Prop 1.1 of Topic 1:

$$V^n = \begin{cases} g(S_N, F_N), & n = N; \\ e^{-r\Delta t} \mathbb{E}_{\mathbb{Q}} \left[ V^{n+1} \middle| \mathcal{F}_n \right], & n = 0, 1, \dots, N-1 \end{cases}$$

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- The main challenge is to take into account the random transition of  $F$  from time  $n$  to  $n+1$  as well when calculating  $\mathbb{E}_{\mathbb{Q}} \left[ V^{n+1} \middle| \mathcal{F}_n \right]$

# Pricing path-dependent options in a binomial tree

- Recall that in a binomial tree:

- ▶ We set

$$s_k^n = S_0 u^{n-k} d^k, \quad \text{for } n = 0, 1, \dots, N \text{ and } k = 0, 1, \dots, n$$

as the possible values of stock price at each time point

- ▶ The transition dynamics of  $S$  is known precisely: when  $S_n = s_k^n$  moves to  $s_{k_{new}}^{n+1}$ , then  $k_{new} = k$  or  $k_{new} = k + 1$
- ▶ For a simple non-path-dependent option, if we denote  $V_k^n$  as the time- $n$  option price when the current stock price is  $s_k^n$  then we know

$$V_k^{n+1} = e^{-r\Delta t} [qV_k^{n+1} + (1-q)V_{k+1}^{n+1}]$$



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- For a path-dependent option, the option price in general is a function of current time  $n$ , current stock price level  $S_n$  and current value of  $F_n$

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- For a path-dependent option, the option price in general is a function of current time  $n$ , current stock price level  $S_n$  and current value of  $F_n$
- Let  $V_{k,j}^n$  be the time- $n$  option price when  $S_n = s_k^n$  and  $F_n = f_j^n$ . Then there are **two questions we must address**:
  - ▶ We need to define  $f_j^n$  which represents the grid of values to be taken by the process  $F$
  - ▶ We need to understand how  $j$  the state index of  $F$  moves as  $s_k^n \rightarrow s_{k_{\text{new}}}^{n+1}$ :

$$V_{k,j}^n = e^{-r\Delta t} [qV_{k+1,?}^{n+1} + (1-q)V_{k,?}^{n+1}]$$

# The forward shooting grid method

The pricing strategy of a path-dependent option under a tree model roughly goes as follows:

- 1 Pick a tree model of the stock price
- 2 Pick a suitable auxiliary process  $F = (F_n)$  based on the nature of the path-dependent option
- 3 Specify a grid of values  $f_j^n$  to contain the possible values of the auxiliary variable  $F_n$  at each time point
- 4 Study how the auxiliary variable moves as the stock price changes, and identify how the indexing system should be updated
- 5 Write down the recursive equation for the fair option price (with interpolation adjustment if needed)

We demonstrate how these procedures are implemented in two concrete examples.

# Lookback option

- Lookback option is a derivative product which payoff function is directly linked to the maximum or minimum stock price level over the product's lifetime. The variations include:

Type of lookback option	Payoff function	
	Discrete model	Continuous model
Fixed strike lookback call	$\left( \max_{i=0,1,\dots,N} S_i - K \right)^+$	$\left( \sup_{0 \leq t \leq T} S_t - K \right)^+$
Fixed strike lookback put	$\left( K - \min_{i=0,1,\dots,N} S_i \right)^+$	$\left( K - \inf_{0 \leq t \leq T} S_t \right)^+$
Floating strike lookback call	$\left( S_N - \min_{i=0,1,\dots,N} S_i \right)$	$\left( S_T - \inf_{0 \leq t \leq T} S_t \right)$
Floating strike lookback put	$\left( \max_{i=0,1,\dots,N} S_i - S_N \right)$	$\left( \sup_{0 \leq t \leq T} S_t - S_T \right)$

# Pricing floating strike lookback put with binomial tree

We use an  $N$ -period binomial tree with Cox-Ross-Rubinstein (CRR) parameterisation to price a floating strike lookback put option with payoff

$$\left( \max_{i=0, \dots, N} S_i - S_N \right).$$

Suppose the tree parameters  $(u, d = 1/u, q)$  have been determined accordingly.

## 1 Fix a tree model of stock price:

- ▶ We are using a binomial tree model with CRR specification such that  $ud = 1$ . Then the possible node values of the stock price are

$$s_k^n = S_0 u^{n-k} d^k = S_0 u^{n-2k} \quad \text{for } n = 0, 1, \dots, N \text{ and } k = 0, \dots, n$$

## 2 Pick an auxiliary variable:

- ▶ Since we are working with lookback option involving the maximum value of the stock price path, we choose  $M_n := \max_{i=0, \dots, n} S_i$  representing the running maximum of the stock price between time zero and time  $n$ . The payoff function becomes

$$g(S_N, M_N) := M_N - S_N$$

# Pricing floating strike lookback put with binomial tree (cont')

## 3 Specify a grid of values for the auxiliary variable $M_n$ :

- ▶ Since the initial stock price is  $S_0$  and the largest possible price at time  $n$  is  $S_0 u^n$ , the maximum price recorded between time zero and time  $n$  must take value in the set

$$\{S_0, S_0 u, S_0 u^2, \dots, S_0 u^n\}$$

- ▶ Hence we construct the grid for  $M_n$  as

$$m_j^n = S_0 u^{n-j} \quad \text{for } n = 0, 1, \dots, N \text{ and } j = 0, 1, \dots, n$$

which represents the  $j$ -th possible value of  $M_n$

# Pricing floating strike lookback put with binomial tree (cont')

## ④ Evolution of auxiliary variable as stock price changes:

- ▶ As  $M_{n+1}$  is the running maximum from time zero to time  $n + 1$ , we expect

$$M_{n+1} = \max(M_n, S_{n+1}) \quad (2.2)$$

- ▶ We need to translate (2.2) to our indexing system to understand how  $M_n$  evolves from  $m_j^n$  as  $s_k^n \rightarrow s_{k_{new}}^{n+1}$
- ▶ Using (2.2) and the definition that  $s_k^n = S_0 u^{n-2k}$  and  $m_j^n = S_0 u^{n-j}$ , we have

$$\begin{aligned} m_{j_{new}}^{n+1} &= \max(m_j^n, s_{k_{new}}^{n+1}) \\ \implies S_0 u^{n+1-j_{new}} &= \max(S_0 u^{n-j}, S_0 u^{(n+1)-2k_{new}}) \\ \implies n+1-j_{new} &= \max(n-j, n+1-2k_{new}) \\ \implies j_{new} &= n+1 - \max(n-j, n+1-2k_{new}) \\ &= \min(2k_{new}, j+1) \\ &=: \phi(k_{new}, j) \end{aligned} \quad (2.3)$$

- ▶  $\phi(k_{new}, j)$  is called the **shooting function** which describes the new state of the auxiliary variable  $M$  when stock price moves to a new state  $k_{new}$  and that the current state of  $M$  is  $j$

# Pricing floating strike lookback put with binomial tree (cont')

## 5 Write down the recursive equation for the fair option prices:

- ▶ If  $V_{k,j}^n$  represents the time- $n$  option value when the stock price is  $S_n = s_k^n$  and the running maximum is  $M_n = m_j^n$ , then

$$V_{k,j}^n = e^{-r\Delta t} [qV_{k,\phi(k,j)}^{n+1} + (1-q)V_{k+1,\phi(k+1,j)}^{n+1}]$$

where  $\phi(\cdot, \cdot)$  is given by (2.3)

- ▶ Note that it is not necessary to compute  $V_{k,j}^n$  for all combinations of  $(k,j)$ . Since by definition we expect  $M_n \geq S_n$ , we only need to consider the combinations of  $(k,j)$  such that

$$m_j^n \geq s_k^n \implies S_0 u^{n-j} \geq S_0 u^{n-2k} \implies j \leq 2k$$



# Pricing floating strike lookback put: a summary

The complete algorithm is as follows:

- ① Define  $s_k^n = S_0 u^{n-2k}$  and  $m_j^n = S_0 u^{n-j}$  for  $n = 0, 1, \dots, N$  and  $j, k = 0, 1, \dots, n$
- ② Option price at terminal time  $n = N$  is given by the payoff function:

- ▶ For each  $k = 0, 1, \dots, N$ 
  - ★ For each  $j = 0, 1, \dots, \min(2k, N)$ , compute

$$V_{k,j}^N = g(s_k^N, m_j^N) = m_j^N - s_k^N$$

- ③ Loop backward in time. For  $n = N - 1, N - 2, \dots, 0$ :
  - ▶ For each  $k = 0, 1, \dots, n$ 
    - ★ For each  $j = 0, 1, \dots, \min(2k, n)$

$$V_{k,j}^n = e^{-r\Delta t} [qV_{k,\phi(k,j)}^{n+1} + (1-q)V_{k+1,\phi(k+1,j)}^{n+1}]$$

where the shooting function is  $\phi(k_{\text{new}}, j) := \min(2k_{\text{new}}, j + 1)$  as derived in (2.3) previously

- ④ The required time-zero option value is  $V_{0,0}^0$

## Pricing floating strike lookback put: numerical example

Use a two-period ( $N = 2$ ) CRR binomial tree to price a floating strike lookback put option with maturity of  $T = 1$  year. Other parameters are:  $S_0 = 100$ ,  $r = 1\%$ ,  $\sigma = 20\%$ .

- Compute the CRR tree parameters as

$$\Delta t = \frac{T}{N} = 0.5, \quad u = e^{\sigma\sqrt{\Delta t}} = 1.15, \quad d = \frac{1}{u} = 0.87, \quad q = \frac{e^{r\Delta t} - d}{u - d} = 0.48$$

- Grid of possible stock price values  $s_k^n = S_0 u^{n-k} d^k$  are given by

		$n$		
		0	1	2
$j$	0	100	115.19	132.69
	1		86.81	100
	2			75.36

- Grid of possible running maximum values  $m_j^n = S_0 u^{n-j}$  are given by

		$n$		
		0	1	2
$j$	0	100	115.19	132.69
	1		100	115.19
	2			100

## Pricing floating strike lookback put: numerical example (cont')

- Option values at terminal time  $n = 2$  are given by the payoff function such that  $V_{k,j}^2 = m_j^2 - s_k^2$ . Here all values of  $V_{k,j}^2$  can be computed as

		$k$	0	1	2
		$s_k^2$	132.69	100	75.36
$j$	$m_j^2$				
0	132.69		0	32.69	57.33
1	115.19			15.19	39.83
2	100			0	24.64

Recall it is sufficient to consider the combination of  $(k, j)$  such that  $m_j^2 \geq s_k^2 \implies j \leq 2k$

- Backward induction: the recursive equation of option price is given by

$$V_{k,j}^n = e^{-r\Delta t} [qV_{k,\phi(k,j)}^{n+1} + (1-q)V_{k+1,\phi(k+1,j)}^{n+1}]$$

for  $k = 0, 1, \dots, n$  and  $j = 0, 1, \dots, \min(2k, n)$ . Recall that  $\phi(k_{\text{new}}, j) = \min(2k_{\text{new}}, j + 1)$

- As an example say we want to compute  $V_{0,0}^1$  (i.e. the option value at time  $n = 1$  when  $S_1 = s_0^1 = 115.19$  and  $M_1 = m_0^1 = 115.19$ ), then

$$\phi(0, 0) = \min(0, 1) = 0, \quad \phi(1, 0) = \min(2, 1) = 1$$

and thus

$$\begin{aligned} V_{0,0}^1 &= e^{-r\Delta t} [qV_{0,\phi(0,0)}^2 + (1-q)V_{1,\phi(1,0)}^2] \\ &= e^{-r\Delta t} [qV_{0,0}^2 + (1-q)V_{1,1}^2] = 7.82 \end{aligned}$$

# Pricing floating strike lookback put: numerical example (cont')

- We can compute the values of  $V_{1,0}^1$ , and  $V_{1,1}^1$  similarly where the results are

$j$	$m_j^1$	$k$ $s_k^1$	0	1
			115.19	86.81
0	115.19		7.82	27.80
1	100			12.69

- Finally, the fair option value at time zero is

$$\begin{aligned}V_{0,0}^0 &= e^{-r\Delta t} [qV_{0,\phi(0,0)}^1 + (1-q)V_{1,\phi(1,0)}^1] \\&= e^{-r\Delta t} [qV_{0,0}^1 + (1-q)V_{1,1}^1] \\&= 10.29\end{aligned}$$

# Asian option

- Asian option is a derivative product which payoff function depends on the arithmetic average of the stock prices over the product's lifetime. The variations include:

Type of Asian option	Payoff function	
	Discrete model	Continuous model
Fixed strike Asian call	$\left(\frac{1}{N+1} \sum_{i=0}^N S_i - K\right)^+$	$\left(\frac{1}{T} \int_0^T S_u du - K\right)^+$
Fixed strike Asian put	$\left(K - \frac{1}{N+1} \sum_{i=0}^N S_i\right)^+$	$\left(K - \frac{1}{T} \int_0^T S_u du\right)^+$
Floating strike Asian call	$\left(S_N - \frac{1}{N+1} \sum_{i=0}^N S_i\right)^+$	$\left(S_T - \frac{1}{T} \int_0^T S_u du\right)^+$
Floating strike Asian put	$\left(\frac{1}{N+1} \sum_{i=0}^N S_i - S_N\right)^+$	$\left(\frac{1}{T} \int_0^T S_u du - S_T\right)^+$

# Pricing fixed strike Asian call with binomial tree

We use an  $N$ -period binomial tree with CRR parameterisation to price a fixed strike Asian call option with payoff

$$\left( \frac{1}{N+1} \sum_{i=0}^N S_i - K \right)^+.$$

Suppose the tree parameters  $(u, d = 1/u, q)$  have been determined accordingly.

## 1 Fix a tree model of stock price:

- ▶ Again we are using a binomial tree model with CRR specification with  $ud = 1$ . The possible node values of the stock price are

$$s_k^n = S_0 u^{n-k} d^k = S_0 u^{n-2k} \quad \text{for } n = 0, 1, \dots, N \text{ and } k = 0, \dots, n$$

## 2 Pick an auxiliary variable:

- ▶ Take  $A_n := \frac{1}{n+1} \sum_{i=0}^n S_i$  which represents the running average of the stock price up to time  $n$ . The payoff function becomes

$$g(A_N) := (A_N - K)^+$$

# Pricing fixed strike Asian call with binomial tree (cont')

## 3 Specify a grid of values for the auxiliary variable $A_n$ :

- ▶ It is computational infeasible to construct a grid covering all possible values of  $A_n$  because in general the total number of possible values at time  $n$  is  $2^n$
- ▶ Hence we use a parameterised grid instead:
  - ★ Fix  $\rho > 0$ . Define

$$a_j^n = S_0 u^{\rho j}, \quad -J_{min}^n \leq j \leq J_{max}^n$$

- ★ We are going to choose the positive integers  $J_{min}^n$  and  $J_{max}^n$  just large enough to ensure the  $a_j^n$ 's cover the possible range of  $A_n$
- ★ Note that  $\ln a_j^n$ 's are evenly spaced:  $\ln a_{j+1}^n - \ln a_j^n = \rho \ln u$
- ★  $\rho$  is called the **quantisation parameter**

# Pricing fixed strike Asian call with binomial tree (cont')

## 3 Specify a grid of values for the auxiliary variable $A_n$ (cont'):

- ▶ At  $n = 0$ ,  $S_0 = A_0$  so we can pick  $J_{max}^0 = J_{min}^0 = 0$
- ▶ Observe that

$$A_n = \frac{1}{n+1} \sum_{i=0}^n S_i = \frac{1}{n+1} \left( \sum_{i=0}^{n-1} S_i + S_n \right) = \frac{nA_{n-1} + S_n}{n+1}$$

- ▶ Suppose  $J_{max}^{n-1}$  has been fixed. Then since the largest possible value of  $S_n$  is  $S_0 u^n$ , the largest possible value of  $A_n$  will be

$$a_{max}^n = \frac{nS_0 u^{\rho J_{max}^{n-1}} + S_0 u^n}{n+1}$$

- ▶ We want to find the smallest possible  $J_{max}^n$  such that  $a_{max}^n \leq S_0 u^{\rho J_{max}^n}$ . We hence pick

$$J_{max}^n = \left\lceil \frac{1}{\rho \ln u} \ln \frac{a_{max}^n}{S_0} \right\rceil = \left\lceil \frac{1}{\rho \ln u} \ln \frac{nu^{\rho J_{max}^{n-1}} + u^n}{n+1} \right\rceil \quad (2.4)$$

where  $\lceil \cdot \rceil$  denotes the ceiling of a number. With  $J_{max}^0 = 0$ , (2.4) inductively defines  $J_{max}^n$  for all  $n$

- ▶ With similar logic, we choose

$$J_{min}^n = \left\lceil \frac{-1}{\rho \ln u} \ln \frac{nu^{-\rho J_{min}^{n-1}} + d^n}{n+1} \right\rceil \quad (2.5)$$



# Pricing fixed strike Asian call with binomial tree (cont')

## ❶ Evolution of auxiliary variable as stock price changes:

- ▶ The transition function of  $A_n$  is

$$A_{n+1} = \frac{(n+1)A_n + S_{n+1}}{n+2} \quad (2.6)$$

We need to translate (2.6) to our indexing system to understand how  $A_n$  evolves from  $a_j^n$  as  $s_k^n \rightarrow s_{k_{new}}^{n+1}$

- ▶ Using (2.6) and the definition that  $s_k^n = S_0 u^{n-2k}$  and  $a_j^n = S_0 u^{\rho j}$ , we have

$$\begin{aligned} S_0 u^{\rho j_{new}} &= \frac{(n+1)S_0 u^{\rho j} + S_0 u^{n+1-2k_{new}}}{n+2} \\ \implies j_{new} &= \frac{1}{\rho \ln u} \ln \frac{(n+1)u^{\rho j} + u^{n+1-2k_{new}}}{n+2} \\ &=: \phi(n, k_{new}, j) \end{aligned} \quad (2.7)$$

- ▶ The problem is that  $\phi(n, k_{new}, j)$  may not give a proper integer! This happens when the predetermined grid of  $a_j^{n+1}$  does not contain the projected value of  $A_{n+1}$

# Pricing fixed strike Asian call with binomial tree (cont')

## 5 Write down the recursive equation of the fair option prices:

- ▶ The recursive equation is supposed to be

$$V_{k,j}^n = e^{-r\Delta t} [qV_{k,\phi(n,k,j)}^{n+1} + (1-q)V_{k+1,\phi(n,k+1,j)}^{n+1}]$$

but of course the above is not well-defined if  $\phi(\cdot, \cdot, \cdot)$  is not an integer

- ▶ Recall that  $V_{k_{\text{new}},\phi(n,k_{\text{new}},j)}^{n+1}$  is supposed to give the option value when  $S_{n+1} = s_{k_{\text{new}}}^{n+1}$  and

$$A_{n+1} = a_{\phi(n,k_{\text{new}},j)}^{n+1} = \frac{(n+1)a_j^n + s_{k_{\text{new}}}^{n+1}}{n+2}.$$

We will estimate  $V_{k_{\text{new}},\phi(n,k_{\text{new}},j)}^{n+1}$  by linear interpolation on the grid of  $\ln a_j^{n+1}$

- ▶ Let  $P[\cdot]$  denote the linear interpolation operator such that

$$P[V_{k_{\text{new}},\phi(n,k_{\text{new}},j)}^{n+1}] := \alpha V_{k_{\text{new}},\phi_-(n,k_{\text{new}},j)}^{n+1} + (1-\alpha)V_{k_{\text{new}},\phi_+(n,k_{\text{new}},j)}^{n+1}$$

with

$$\phi_+(n, k_{\text{new}}, j) := \lceil \phi(n, k_{\text{new}}, j) \rceil, \quad \phi_-(n, k_{\text{new}}, j) := \lfloor \phi(n, k_{\text{new}}, j) \rfloor$$

and

$$\alpha := \frac{\ln a_{\phi_+(n,k_{\text{new}},j)}^{n+1} - \ln a_{\phi(n,k_{\text{new}},j)}^{n+1}}{\ln a_{\phi_+(n,k_{\text{new}},j)}^{n+1} - \ln a_{\phi_-(n,k_{\text{new}},j)}^{n+1}} = \frac{\ln a_{\phi_+(n,k_{\text{new}},j)}^{n+1} - \ln a_{\phi(n,k_{\text{new}},j)}^{n+1}}{\rho \ln u}.$$

Here  $\lceil \cdot \rceil$  and  $\lfloor \cdot \rfloor$  denote the ceiling and floor of a number respectively

- ▶ The modified recursion is

$$V_{k,j}^n = e^{-r\Delta t} \left\{ qP[V_{k,\phi(n,k,j)}^{n+1}] + (1-q)P[V_{k+1,\phi(n,k+1,j)}^{n+1}] \right\}$$

## Pricing fixed strike Asian call: numerical example

Use a three-period ( $N = 3$ ) CRR binomial tree to price a fixed strike Asian call option of maturity  $T = 1$  year. Other parameters are:  $S_0 = 100$ ,  $K = 100$ ,  $r = 1\%$ ,  $\sigma = 20\%$ . Use  $\rho = 0.5$  for the grid of the running averages.

- The CRR tree parameters are given by

$$\Delta t = \frac{T}{N} = 0.33, \quad u = e^{\sigma\sqrt{\Delta t}} = 1.12, \quad d = \frac{1}{u} = 0.89, \quad q = \frac{e^{r\Delta t} - d}{u - d} = 0.49$$

- Stock price grid values are given by  $s_k^n = S_0 u^{n-k} d^k$ :

		$n$			
		0	1	2	3
$k$	0	100	112.24	125.98	141.40
	1		89.09	100	112.24
	2			79.38	89.09
	3				70.72

- Use (2.4) and (2.5) to work out the values of  $J_{max}^n$  and  $J_{min}^n$  for the construction of grid for  $A_n$ :

		$n$			
		0	1	2	3
$J_{max}^n$		0	2	3	4
$J_{min}^n$		0	1	2	3

# Pricing fixed strike Asian call: numerical example (cont')

- The grid of  $A_n$  is constructed as  $a_j^n = S_0 u^{\rho j}$  for  $-J_{min}^n \leq j \leq J_{max}^n$

		$n$			
		0	1	2	3
$j$	4				125.98
	3			118.91	118.91
	2		112.24	112.24	112.24
	1		105.94	105.94	105.94
	0	100	100	100	100
	-1		94.39	94.39	94.39
	-2			89.09	89.09
	-3				84.10

- The terminal option price at  $n = 3$  is  $V_{k,j}^3 = (a_j^3 - K)^+$ . Hence for all  $k = 0, 1, 2, 3$  we have

$j$	4	3	2	1	0	-1	-2	-3
$a_j^3$	125.98	118.91	112.24	105.94	100	94.39	89.09	84.10
$V_{k,j}^3$	25.98	18.91	12.24	5.94	0.00	0.00	0.00	0.00

## Pricing fixed strike Asian call: numerical example (cont')

- Suppose we want to compute  $V_{1,0}^2$  which is the option value at time  $n = 2$  when the current stock price is  $s_1^2 = 100$  and the current running average is  $a_0^2 = 100$
- Recall the transition dynamics  $A_{n+1} = \frac{(n+1)A_n + S_{n+1}}{n+2}$ . There are two cases:
  - ▶ If the stock goes up such that  $S_3 = s_1^3 = 112.24$ , then  $A_3 = \frac{3 \times 100 + 112.24}{4} = 103.06$
  - ▶ If the stock goes down such that  $S_3 = s_2^3 = 89.09$ , then  $A_3 = \frac{3 \times 100 + 89.09}{4} = 97.27$
  - ▶ Neither of these values lie on the grid of  $a_j^3$
- The standard recursive equation is supposed to be

$$V_{k,j}^n = e^{-r\Delta t} [qV_{k,\phi(n,k,j)}^{n+1} + (1-q)V_{k+1,\phi(n,k+1,j)}^{n+1}]$$

where  $\phi(n, k_{new}, j)$  is defined in (2.7). We have

$$\phi(2, 1, 0) = 0.52, \quad \phi(2, 2, 0) = -0.48$$

and thus

$$V_{1,0}^2 = e^{-r\Delta t} [qV_{1,0.52}^3 + (1-q)V_{2,-0.48}^3]$$

- We need to use interpolation to approximate  $V_{1,0.52}^3$  and  $V_{2,-0.47}^3$

# Pricing fixed strike Asian call: numerical example (cont')

- To approximate  $V_{1,0.52}^3$ , we look at the nearest available option values:

	$j$		
	1	0.52	0
$a_j^3$	105.94	103.06	100.00
$\ln a_j^3$	4.66	4.64	4.61
$V_{1,j}^3$	5.94	?	0.00

Using linear interpolation in  $\ln a_j^3$ , we compute

$$\begin{aligned}
 P[V_{1,0.52}^3] &= \alpha V_{1,\phi_-(2,1,0)}^3 + (1 - \alpha) V_{1,\phi_+(2,1,0)}^3 \\
 &= \frac{\ln a_1^3 - \ln a_{0.52}^3}{\rho \ln u} V_{1,0}^3 + \left(1 - \frac{\ln a_1^3 - \ln a_{0.52}^3}{\rho \ln u}\right) V_{1,1}^3 \\
 &= \frac{4.66 - 4.64}{0.05} \times 0 + \frac{4.64 - 4.61}{0.05} \times 5.94 = 3.10
 \end{aligned}$$

- By similar interpolation exercise again, we can obtain  $P[V_{2,-0.47}^3] = 0$
- This finally gives  $V_{1,0}^2 \approx e^{-r\Delta t} \{qP[V_{1,0.52}^3] + (1 - q)P[V_{2,-0.47}^3]\} = 1.50$
- Ultimately we can get  $V_{0,0}^0 = 4.81$

# American path-dependent option

- There also exists American style of path-dependent option where the option holder can freely exercise the option at any stopping time  $\tau$  to receive the payoff  $g(S_0, S_1, \dots, S_\tau)$  immediately
- If the payoff function can be converted into  $g(S_\tau, F_\tau)$  for some auxiliary process  $F$ , then the time- $n$  fair value of the option is

$$\sup_{\tau \in \mathcal{T}_{n,N}} \mathbb{E}_{\mathbb{Q}} \left[ e^{-r(\tau-n)\Delta t} g(S_\tau, F_\tau) | \mathcal{F}_n \right]$$

where  $\mathcal{T}_{n,N}$  is the set of stopping times taking values in  $\{n, n+1, \dots, N\}$

- Just like Prop 1.6 in Topic 1, pricing of the American version of the path-dependent option is very easy in a lattice model where we just need to consider the maximum of intrinsic value and continuation value of the option in the backward induction procedure
- E.g. in a binomial tree model the recursive equation will become

$$V_{k,j}^n = \max \left\{ g(s_k^n, f_j^n), e^{-r\Delta t} \left[ qV_{k,\phi(n,k,j)}^{n+1} + (1-q)V_{k+1,\phi(n,k+1,j)}^{n+1} \right] \right\}$$

where  $g(s_k^n, f_j^n)$  is the intrinsic value of the option and  $\phi(\cdot, \cdot, \cdot)$  is some suitably constructed shooting function

- The ability to handle American path-dependent options with ease is the most powerful feature of a lattice method

## Optional reading

- Barraquand, J. & Pudet, T. (1996). Pricing of American Path-dependent Contingent Claims. *Mathematical Finance*, 6(1), 17 - 51.
- Hull, J. & White A. (1993). Efficient Procedures for Valuing European and American Path Dependent Options. *Journal of Derivatives*, 1, 21 - 31.
- Forsyth, P. A., Vetzal, K. R. & Zvan, R. (2002). Convergence of Numerical Methods for Valuing Path-dependent Options Using Interpolation. *Review of Derivatives Research*, 5(3), 273 - 314.