EEPROM File System Design Notes

Overview

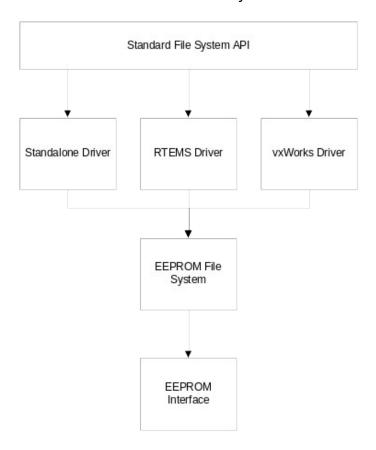
The EEPROM File System is a custom file system designed for data stored in EEPROM. The EEPROM File System is a slot based file system where each slot is a fixed size contiguous region of memory allocated for a single file. Note that each slot can be a different size. The size of each slot is determined when the file system is created and is based on the size of the file to be stored in the slot. Additional free space can be to the end of each slot to allow room for the file to grow in size if necessary and new files can be created at runtime.

Design Goals

- Easy on-orbit maintenance. Since each file is located in a known fixed contiguous region of memory, it is easy to patch or reload the file from the ground.
- Support for the standard file system api. Looks like any other file system to application software.
- Minimize the number of writes to EEPROM.
- Simple implementation.
- Low memory/system overhead.
- The EEFS shall store a single file in a fixed continuous region of memory. This is opposed to a block device where a single file may be broken up into multiple non-continuous memory blocks
- The EEFS shall support a standard file system api.
- Patching
- Disadvantages to this approach vs a block device

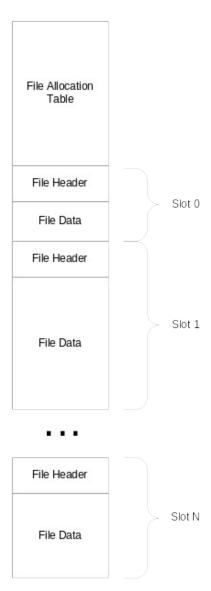
Layered Architecture

The EEPROM File System is split into four layers. At the top most layer the EEPROM File System supports the standard file system api, which makes it look like any other file system to application software. Each implementation of the file system has an OS Specific Driver layer that maps the standard file system api function calls to EEPROM File System specific function calls. The EEPROM File System layer contains the low level EEPROM File System software, and finally the EEPROM Interface layer includes functions that talk directly to the EEPROM hardware.



File System Structure

The EEPROM File System structure is outlined below. The file system contains a File Allocation Table followed by slots for each file in the file system. Each slot contains a File Header and File Data.



File Allocation Table:

The File Allocation Table defines where in memory each slot starts as well as the maximum slot size for each file. This table is a fixed size regardless of how many files actually reside in the file system and never changes unless a new file is added to the file system using the EEFS_LibCreat() function or the whole file system is reloaded. The maximum number of files that can be added to the file system is determined at compile time by the EEFS_MAX_FILES define. It is important to choose this number carefully since a code patch would be required to change it. The file offsets for each file defined in the File Allocation Table are relative offsets

from the beginning of the file system and point to the start of the File Header. Since they are relative offsets the file system in not tied to any physical address. In fact the exact same file system image can be burned into multiple locations of EEPROM and then mounted as different volumes. The size of each slot is specified in bytes and does NOT include the File Header. So each slot actually occupies (File Header Size + Slot Size) bytes of EEPROM.

```
typedef struct
    uint32
                                         Crc;
    uint32
                                         Magic;
    uint32
                                         Version;
                                         FreeMemoryOffset;
    uint32
    uint32
                                         FreeMemorySize;
                                         NumberOfFiles;
    uint32
} EEFS_FileAllocationTableHeader_t;
typedef struct
    uint32
                                         FileHeaderOffset;
    uint32
                                         MaxFileSize;
} EEFS_FileAllocationTableEntry_t ;
typedef struct
    EEFS_FileAllocationTableHeader_t
                                         Header:
                                         File[EEFS_MAX_FILES];
    EEFS_FileAllocationTableEntry_t
} EEFS_FileAllocationTable_t;
```

File Header:

Each slot in the file system starts with a File Header. The File Header contains information about the file contained in the slot such as its filename. The File Header is initialized to all 0's if the slot is unused or the file has been deleted.

File Data:

File Data starts immediately following the File Header and may or may not use all of the available space in the slot.

Inode Table

The Inode Table is a ram table that is used by the file system api to access the file system and is similar in structure to the File Allocation Table. The Inode table is initialized when the function EEFS_LibInitFS() is called and once the Inode table is initialized the File Allocation Table is no longer used. One important difference between the File Allocation Table and the Inode Table is that the Inode Table

contains physical address pointers to the start of each file instead of relative offsets. Note also that the Inode table does not cache any information about the file in ram. This means that the file could be patched or reloaded to EEPROM without the need to patch the Inode Table, i.e. the file updates are available to the file system immediately. A disadvantage to this approach is that each File Header must be read from EEPROM when searching the file system for a specific file, for example when a file is opened.

```
typedef struct
                                        *FileHeaderPointer;
    void
    uint32
                                         MaxFileSize;
} EEFS_InodeTableEntry_t;
typedef struct
    uint32
                                         BaseAddress:
    void
                                        *FreeMemoryPointer;
    uint32
                                         FreeMemorySize;
    uint32
                                         NumberOfFiles;
    EEFS_InodeTableEntry_t
                                         File[EEFS_MAX_FILES];
} EEFS_InodeTable_t;
```

File Descriptor Table

The File Descriptor Table manages all File Descriptors for the EEPROM File System. There is only one File Descriptor Table that is shared by all EEPROM File System volumes. The maximum number of files that can be open at one time is determined at compile time by the EEFS_MAX_OPEN_FILES define.

```
typedef struct
    uint32
                                         InUse;
    uint32
                                         Mode:
    void
                                        *FileHeaderPointer;
    void
                                        *FileDataPointer:
    uint32
                                         ByteOffset;
    uint32
                                         FileSize;
    uint32
                                         MaxFileSize:
    EEFS_InodeTable_t
                                         *InodeTable;
    uint32
                                         InodeIndex;
} EEFS_FileDescriptor_t;
EEFS_FileDescriptor_t
                                         EEFS_FileDescriptorTable[EEFS_MAX_OPEN_FILES];
```

Directory Descriptor Table

The Directory Descriptor Table manages the Directory Descriptor for the EEPROM File System. There is currently only one Directory Descriptor that is shared by all EEPROM File System volumes.

```
typedef struct
    uint32
                                         InUse;
                                         InodeIndex:
    uint32
    EEFS_InodeTable_t
                                         *InodeTable;
} EEFS_DirectoryDescriptor_t;
typedef struct
    uint32
                                         InodeIndex;
                                         Filename[EEFS_MAX_FILENAME_SIZE];
    char
    uint32
                                         InUse:
                                        *FileHeaderPointer;
    void
    uint32
                                         MaxFileSize;
} EEFS_DirectoryEntry_t;
```

EEPROM Access

The EEPROM File System software never directly reads or writes to EEPROM, instead it uses implementation specific EEPROM interface functions. Since not all EEPROM is memory mapped, some EEPROM implementations may require implementation specific functions for accessing EEPROM. The implementation specific EEPROM interface functions are defined as macros in the file eefs_macros.h. By default these macros are defined to use memcpy. Note also that the EEPROM interface functions are protected from shared access by the EEPROM File System however there is nothing that prevents other processes from calling the EEPROM interface functions from outside of the EEPROM File System.

```
#define EEFS_LibEEPROMWrite(Dest, Src, Length) memcpy(Dest, Src, Length)
#define EEFS_LibEEPROMRead(Dest, Src, Length) memcpy(Dest, Src, Length)
#define EEFS_LibEEPROMFlush
```

EEPROM Write Protection

The EEPROM File System can be write protected. This feature is implemented through an implementation specific interface function that returns the write protection state of the file system. This interface function is defined as a macro in the file eefs_macros.h. If the file system is read-only then the macro will be defined to TRUE. If the file system is always write enabled then the macro will be defined to FALSE. If the eeprom has an external write protection interface then a custom function can be called to determine the write protect status. By default this macro is defined to FALSE which means that the file system is not write protected.

```
#define EEFS_LibIsWriteProtected EEPROM_IsWriteProtected(InodeTable->BaseAddress)
```

Mutual Exclusion

Mutual exclusion is implemented by the functions EEFS_LibLock and EEFS_LibUnlock. Since the EEPROM File System is not intended to be used very often and to keep things simple it was decided to implement a single locking mechanism that is shared by all EEPROM File System volumes. This locking

mechanism simply locks the shared resource at the start of each function and unlocks the shared resource at the end of each function. It is recommended that semaphores be used as the locking mechanism vs disabling interrupts. Note that since the shared resource is locked for all lower level functions, lower level functions should not be called recursively. The implementation of the EEFS_LibLock and EEFS LibUnlock functions are defined as macros in the file eefs macros.h.

```
#define EEFS_LIB_LOCK semTake(EEFS_semId, WAIT_FOREVER);
#define EEFS_LIB_UNLOCK semGive(EEFS_semId);
```

Time Stamps

Time Stamps are implemented by the function EEFS_LibTime. Time stamps are based on the standard library time_t. The implementation of the EEFS_LibTime function is defined as a macro in the file eefs_macros.h. By default this macro is defined to use the standard library time function.

File Attributes

The only File Attribute supported by the EEPROM File System is the Read Only file attribute. This Attribute can be set on a per file basis when the file system is created or it can be changed at run time by calling the function EEFS_LibSetFileAttributes(EEFS_InodeTable_t *InodeTable, char *Filename, uint32 Attributes).

Directories

The EEPROM File System is a flat file system therefore it only supports a single top level directory and does not support sub directories. If you want to group files separately then they should be placed in different volumes. For example /EEFSO Apps for apps, /EEFSO Tables for tables etc...

CRC's

The EEPROM File System includes a crc in the File Allocation Table and a crc in the File Header for each file. Currently the crc included in the File Allocation Table is calculated across the entire file system, including unused space (i.e. MaxEepromSize), and is currently only used by bootstrap code to verify the integrity of the file system at boot time. This crc is NOT automatically updated by the file system when files are modified, so this must be done manually. The crc included in each File Header is currently not used by the file system.

Rewriting Existing Files

Existing files can be updated or completely rewritten simply by opening an existing file for write access. The only restriction is that existing files are limited in size to the max file size for that slot. If the file grows larger than the max file size then file write calls will return 0 bytes written.

Creating New Files

New files can be created at runtime by opening a file for write access that does not already exist with the O_CREAT flag or by calling the creat function . When a new file is created a new slot is added to the end of the file system that initially has a max file size equal to all available unused space in the file system. This is done because we won't know the final size of the file until it is closed. When the file is closed the slot is resized to the actual size of the file plus a fixed amount of spare and the additional unused space is returned to the file system. The fixed amount of spare is determined at compile time by the EEFS_DEFAULT_CREAT_SPARE_BYTES define. New files can be added to the file system up to EEFS_MAX_FILES or all available space in the file system is used. Note also that only one new file can be created at a time.

Deleting Files

The file system supports deleting files however it is not recommended. Deleting a file renders the slot containing the file unusable. Note that because of this design the file system will not reclaim space from a deleted file.

Patching Files

Existing files can be easily patched outside of the file system api when necessary. In some cases only the file data that is changing needs to be loaded, however if the file size or the file name is changing then the file header will also have to be updated. File metadata is not cached in ram so file updates only need to occur to eeprom, no other patches are necessary. The address of each slot can be found by looking at the file system map file created by the geneepronfs tool. See the section on Building a File System Image for more information.

Micro EEPROM File System

The Micro version of the EEPROM file system allows bootstrap code access to files in an EEPROM File System. The full implementation of the EEPROM File System is too large to be used in bootstrap code so a simple single function version was developed that returns the starting address of a file given its filename. The bootstrap code can then boot the system from a kernel image that is contained in an EEPROM File System. This software is designed to use very little memory, and is independent of the file system size (i.e. maximum number of files). This means that the bootstrap image will NOT have to be updated whenever the size of the EEPROM File System changes.

Building a File System Image

EEPROM File System images are created using the geneepromfs tool. This command line tool reads an input file that describes the files that will be included in the file system and outputs an EEPROM File System image ready to be burned into EEPROM.

Available Options

```
Build a EEPROM File System Image.
  Options:
  -e, --endian=big or little
                                     set the output encoding (big)
  -s, --eeprom_size=SIZE
                                     set the size of the target eeprom (2 Mb)
  -t, --time=TIME
                                     set the file timestamps to a fixed value
  -f, --fill_eeprom
-v, --verbose
                                     fill unused eeprom with 0's
                                     print the name of each file added to the
                                       file system
  -m, --map=FILENAME
                                     output a file system memory map
  -V, --version
-h, --help
                                     output version information and exit
                                     output usage information and exit
  The INPUT_FILE is a formatted text file that specifies the files to be added
    to the file system. Each entry in the INPUT_FILE contains the following
    fields separated by a comma:
   1. Input Filename: The path and name of the file to add to the file system
2. EEFS Filename: The name of the file in the eeprom file system. Note the
         EEFS Filename can be different from the original Input Filename
    3. Spare Bytes: The number of spare bytes to add to the end of the file.
         Note also that the max size of the file is rounded up to the nearest
         4 byte boundary.
    4. Attributes: The file attributes, EEFS_ATTRIBUTE_NONE or EEFS_ATTRIBUTE_READONLY.
    Each entry must end with a semicolon.
    Comments can be added to the file by preceding the comment with an
      exclamation point.
    Example:
     Input Filename EEFS Filename Spare Bytes Attributes
      /../images/cfe-core.slf, file1.slf,
                                                    100,
                                                                 EEFS_ATTRIBUTE_NONE;
```

Sample Input File

! Input Filename	EEFS Filename	Spare Bytes	Attributes
startupA.scr,	startupA.scr,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
startupB.scr,	startupB.scr,	128,	EEFS_ATTRIBUTE_NONE;
cfe_es_startup.scr,	cfe_es_startup.scr,	128,	EEFS_ATTRIBUTE_NONE;
cfe-core.o,	cfe-core.o,	2048,	EEFS_ATTRIBUTE_NONE;
cdh_lib.o.gz,	cdh_lib.o.gz,	2048,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cfs_lib.o.gz,	cfs_lib.o.gz,	2048,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
sw.o.gz,	sw.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
sw_a_netwtbl.tbl,	sw_a_netwtbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
sw_a_sca_rttbl.tbl,	sw_a_sca_rttbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
sw_a_scb_rttbl.tbl,	sw_a_scb_rttbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
sw_a_transtbl.tbl,	sw_a_transtbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
sw_b_netwtbl.tbl,	sw_b_netwtbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
sw_b_sca_rttbl.tbl,	sw_b_sca_rttbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
sw_b_scb_rttbl.tbl,	sw_b_scb_rttbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
sw_b_transtbl.tbl,	sw_b_transtbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
sw_sbccnfgtbl.tbl,	sw_sbccnfgtbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
sw_scommacfgtbl.tbl,	sw_scommacfgtbl.tbl		EEFS_ATTRIBUTE_NONE;
sw_scommbcfgtbl.tbl,	sw_scommbcfgtbl.tbl,	•	EEFS_ATTRIBUTE_NONE;
Xb.o.gz,	Xb.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
xb_cchntblin.tbl,	xb_cchntblin.tbl,	128,	EEFS_ATTRIBUTE_NONE;
xb_cchntblsc.tbl, xb_cdsctblin.tbl,	<pre>xb_cchntblsc.tbl, xb_cdsctblin.tbl,</pre>	128, 128,	EEFS_ATTRIBUTE_NONE;
xb_cdsctblin.tbt, xb_cdsctblsc.tbl,	xb_cdsctblsc.tbl,	128,	EEFS_ATTRIBUTE_NONE; EEFS_ATTRIBUTE_NONE;
xb_ttbl.tbl,	xb_ttbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
Ci.O.gz,	ci.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
ci_fec_keytbl.tbl,	ci_fec_keytbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
gnc.o.gz,	gnc.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
cf.o.gz,	cf.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
cf_cfgtable.tbl,	cf_cfgtable.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cs.o.gz,	cs.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
cs_eepromtbl.tbl,	cs_eepromtbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg.o.gz,	cg.o.gz,	2048,	EEFS_ATTRIBUTE_NONE;
cg_mtb_tbl.tbl,	cg_mtb_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cg_mtb_b_tbl.tbl,	cg_mtb_b_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cg_rw_tbl.tbl,	cg_rw_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cg_rw_f0_tbl.tbl,	cg_rw_f0_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cg_rw_f1_tbl.tbl,	cg_rw_f1_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg_rw_f2_tbl.tbl,	cg_rw_f2_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg_rw_f3_tbl.tbl,	cg_rw_f3_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg_rw_f4_tbl.tbl,	cg_rw_f4_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg_sa_tbl.tbl,	cg_sa_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg_hga_tbl.tbl,	cg_hga_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cg_prop_tbl.tbl,	cg_prop_tbl.tbl, cl.o.gz,	128, 2048,	<pre>EEFS_ATTRIBUTE_NONE; EEFS_ATTRIBUTE_NONE;</pre>
cl.o.gz, cl_rnmode_tbl.tbl,	cl_rnmode_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_spmode_tbl.tbl,	cl_spmode_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_gspmode_tbl.tbl,	cl_gspmode_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_msmode_tbl.tbl,	cl_msmode_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_slew_tbl.tbl,	cl_slew_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_dhmode_tbl.tbl,	cl_dhmode_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_dvmode_tbl.tbl,	cl_dvmode_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_modemgr_tbl.tbl,	cl_modemgr_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cl_hga_tbl.tbl,	cl_hga_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cl_sa_tbl.tbl,	cl_sa_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
cl_system_tbl.tbl,	cl_system_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
cl_target_tbl.tbl,	cl_target_tbl.tbl,	128,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
ds.o.gz,	ds.o.gz,	2048,	<pre>EEFS_ATTRIBUTE_NONE;</pre>
ds_file_tbl.tbl,	ds_file_tbl.tbl,	128,	EEFS_ATTRIBUTE_NONE;
DS_LEO_filter_tbl.tbl,	DS_LEO_filter_tbl.tl		EEFS_ATTRIBUTE_NONE;
DS_DV_filter_tbl.tbl,	DS_DV_filter_tbl.tb		EEFS_ATTRIBUTE_NONE;
DS_Nom_filter_tbl.tbl,	DS_Nom_filter_tbl.tl		EEFS_ATTRIBUTE_NONE;
DS_Safe_filter_tbl.tbl,	DS_Safe_filter_tbl.		EEFS_ATTRIBUTE_NONE;
di.o.gz, di_mtb_b_tbl.tbl,	di.o.gz, di_mtb_b_tbl.tbl,	2048, 128	EEFS_ATTRIBUTE_NONE;
u		128,	EEFS_ATTRIBUTE_NONE;

Sample Map File

Offset 0	Size 4120	Sectior FAT	n Slot	Filename	File Si	ize	Spare	Max Siz	ze Crc	Attrib	utes	
4120	64	Header	Θ									
4184	288	Data	0	startupA.scr	160	128	288	0xFFFF9	1818	0		
4472	64	Header		ocar capition	100	120	200	OMITTE	7010	Ü		
4536	288	Data	1	startupB.scr	160	128	288	0xFFFF9	9888	0		
4824	64	Header										
4888	4112	Data	2	cfe_es_startup	3984	128	4112	0x0000!	5A51	0		
9000	64	Header	3									
9064	333056		3	cfe-core.o	331007	2049	333056	0x00007	713B	0		
342120	64	Header										
342184	26480	Data	4	cdh_lib.o.gz 24429		2051	26480	0xFFFFE	39B0	0		
368664 368728	64 3580	Header Data	5 5	cfs_lib.o.gz	2051	3580	0xFFFFB62B 0					
372308		Header		013_t1b.0.g2	1529	2031	3300	OXFFFF	0020	U		
372372		Data	6	sw.o.gz 21206	2050	23256	0xFFFF8	BACC	Θ			
395628		Header	7	5 g= =====			•		-			
395692		Data	7	sw_a_netwtbl.tbl		180	128	308	0xFFFFI	F6E1	0	
396000	64	Header	8									
396064	396	Data	8	sw_a_sca_rttbl	.tbl	266	130	396	0xFFFFI	0221	0	
396460	64	Header										
396524	396	Data	9	sw_a_scb_rttbl	.tbl	266	130	396	0xFFFFI	32E4	0	
396920	64	Header			461	F00	100	704	00000	25.70	•	
396984 397708	724	Data	10	sw_a_transtbl.	ומז	596	128	724	0×00002	2576	0	
397772		Header Data	11 11	sw_b_netwtbl.t	h1	180	128	308	0×0000	SEOC	0	
398080	64	Header		3W_D_HETWEDT. C	D L	100	120	300	0.00000	3390	U	
398144	396	Data	12	sw_b_sca_rttbl	.tbl	266	130	396	0xFFFF	AADB	0	
398540	64	Header										
398604	396	Data	13	sw_b_scb_rttbl	.tbl	266	130	396	0xFFFF	CA1E	0	
399000	64	Header	14									
399064	724	Data	14	sw_b_transtbl.	tbl	596	128	724	0xFFFF	ABBD	0	
399788	64	Header									_	
399852	332	Data	15	sw_sbccnfgtbl.tbl		204	128	332	0x00002	2A7C	0	
400184	64	Header				204	100	222	٥٧٢٢٢٢	-745	0	
400248 400580	332 64	Data Header	16 17	sw_scommacfgtb		204	128	332	0xFFFFI	=/4F	0	
400644	332	Data	17 17	sw_scommbcfgtb	l thl	204	128	332	0×0000	398F	0	
400976	64	Header		on_500iiiiib01 geb		201	120 002		0,0000	3001	Ü	
401040	30592	Data	18	xb.o.gz 28541 2051		30592	0xFFFFFDE1		0			
431632	64	Header	19	J								
431696	1284	Data	19	xb_cchntblin.t	bl	1156	128	1284	0x00002	2737	0	
432980	64	Header										
433044	1284	Data	20	xb_cchntblsc.t	bl	1156	128	1284	0xFFFFI	BADC	0	
434328	64	Header			L 1	74.0	100	0.4.4	0.,0000	2500	0	
434392		Data	21 22	xb_cdsctblin.t	DL	716	128	844	0×0000	3F3C	0	
435236 435300	64 844	Header Data	22	xb cdsctblsc.t	h1	716	128	844	0×0000	17C/	0	
436144		Header		AD_CUSCIDISC.I	D L	710	120	044	0,0000.	1704	U	
436208		Data	23	xb_ttbl.tbl	1216	128	1344	0x00006	318	0		
437552		Header										
437616		Data	24	ci.o.gz 14447	2049	16496	0xFFFF0	696	0			
454112	64	Header	25									
454176		Data	25	ci_fec_keytbl.	tbl	8308	128	8436	0x00003	3BA1	0	
462612		Header			0.4.0.0							
462676		Data	26	gnc.o.gz	3196	2048	5244	0xFFFFE	388C	0		
467920 467984		Header Data	27 27	cf.o.gz 78290	2050	80340	0xFFFF0	2050	0			
548324		Header		C1.0.y2 76290	2030	00340	UXFFFF	,639	U			
548388		Data	28	cf_cfgtable.tbl		4100	128	4228	0×00002	2769	0	
552616		Header		55. gtab to . tb t					5555521.55			
552680		Data	29	cs.o.gz 20165	2051	22216	0×00002	DFF	Θ			
574896		Header	30	-								
574960		Data	30	cs_eepromtbl.t	308	128	436	0xFFFFI	343C	0		
575396		Header		00 0 == 07043	2054	20000	0.45555	104	0			
575460	29692	Data	31	cg.o.gz 27641	2051	29692	0xFFFFE	3484	0			

605152 64 Header 32 605216 492 Data 32 cg_mtb_tbl.tbl 364 128 492 0x00007A7F 0