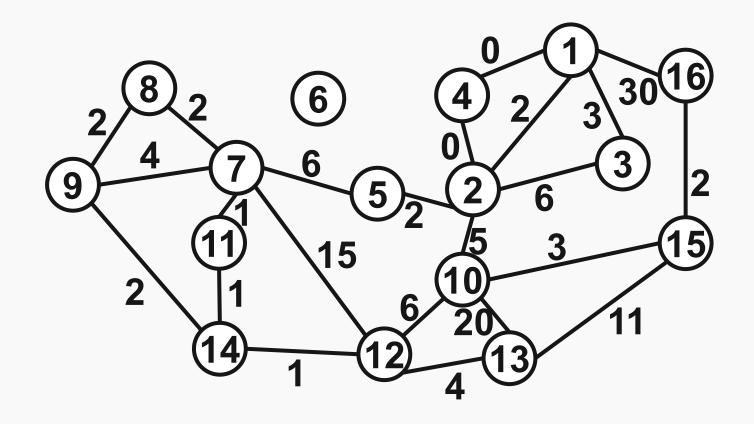




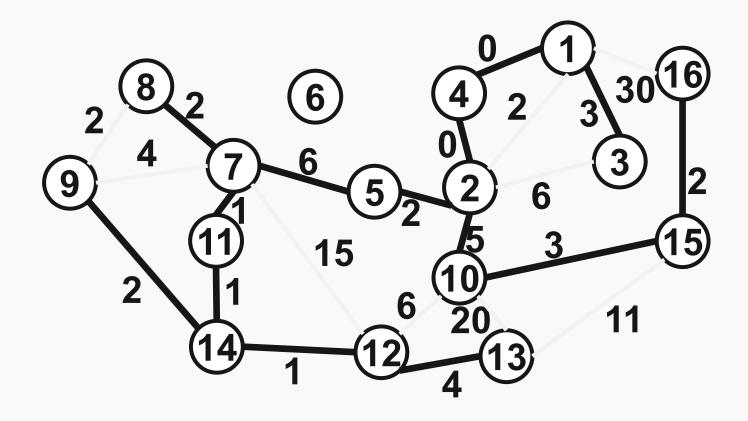


Arbori minimi de acoperire





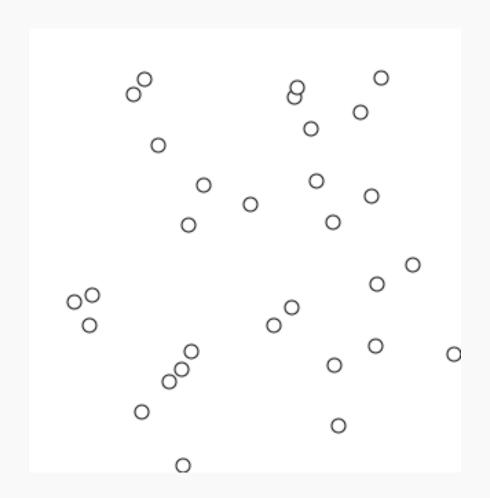
Arbori minimi de acoperire



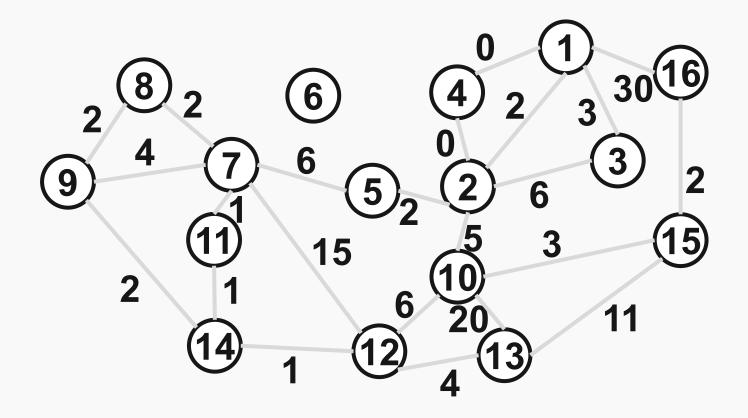


```
tree Kruskal(G) {
    sort(G.E); // sort by weight
    A = \{\};
    for each (node in G.V)
        Make set(node);
    for each ((u, v) in G.E) {
        if (Find set(u) != Find set(v)) {
            A = A \cup \{(u, v)\};
            Union(Find set(u), Find_set(v));
    return A;
```

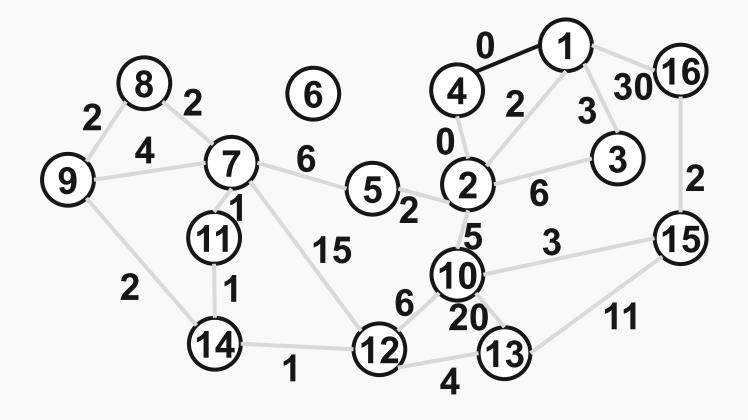




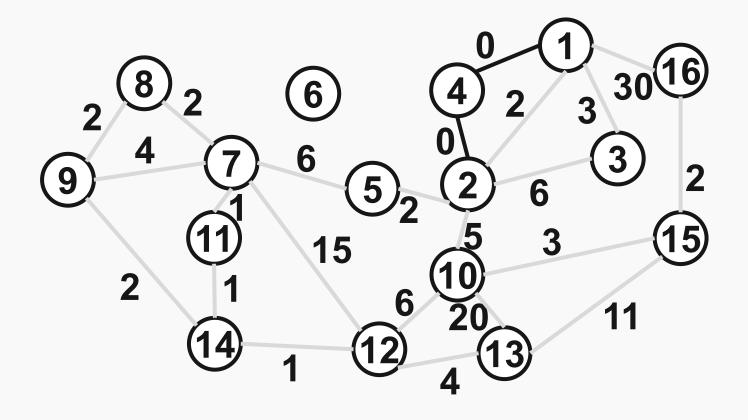




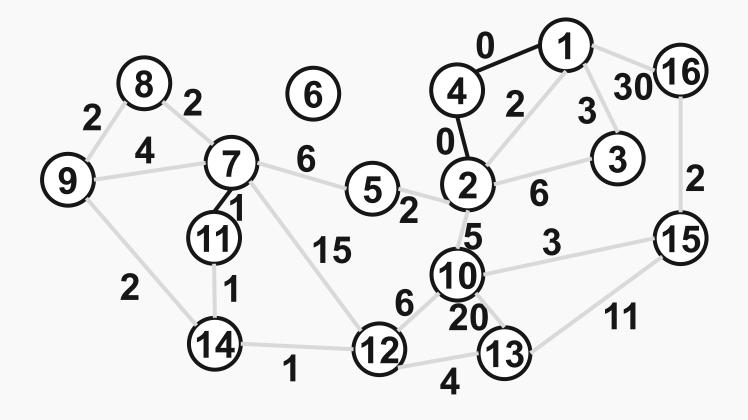




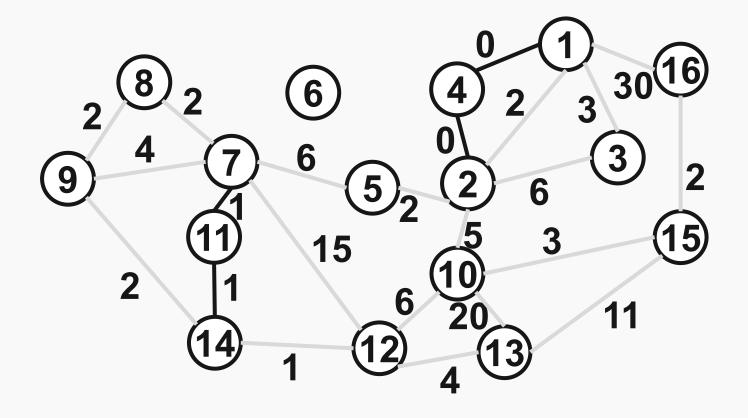




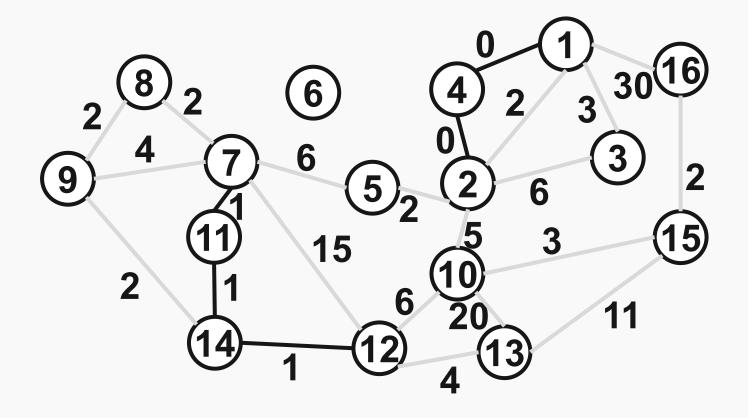




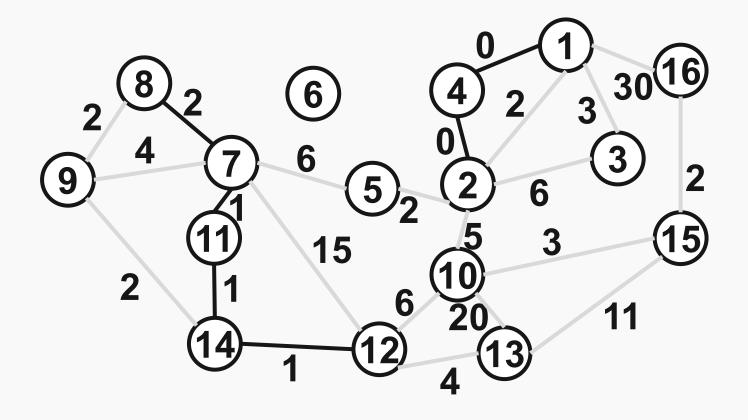




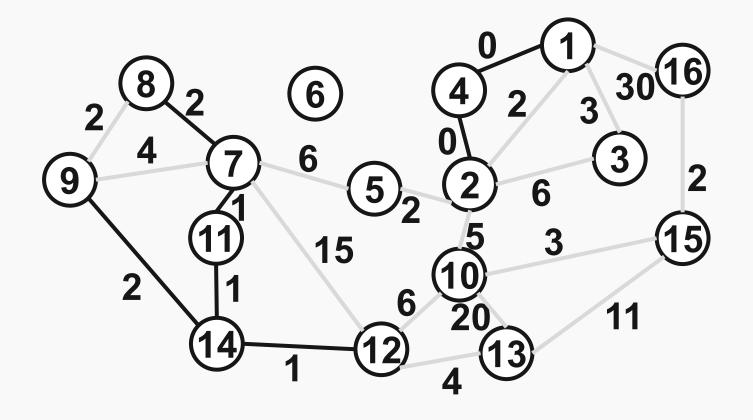




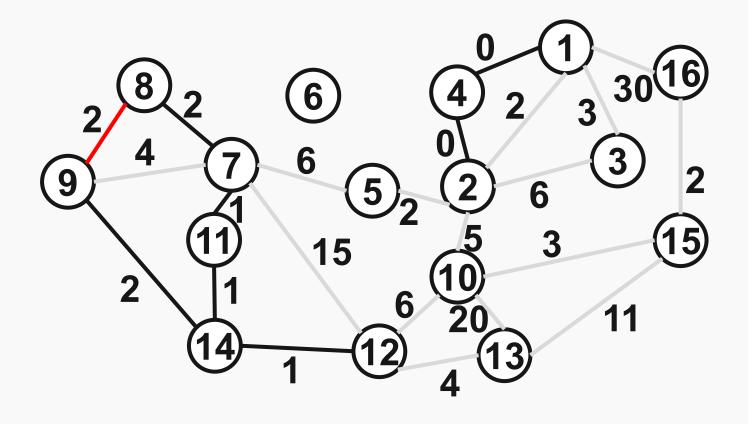




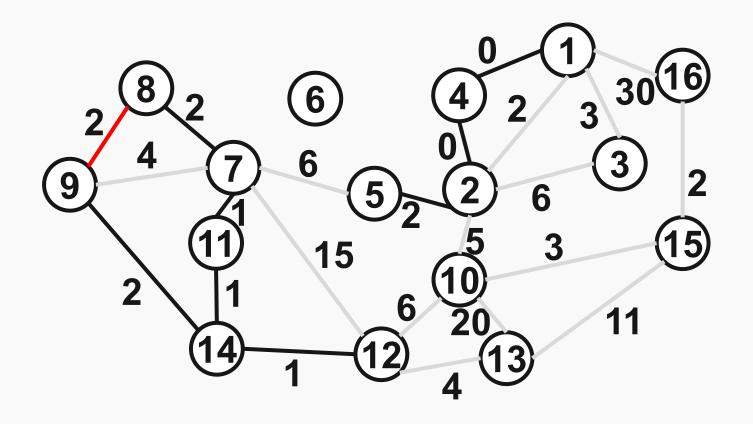




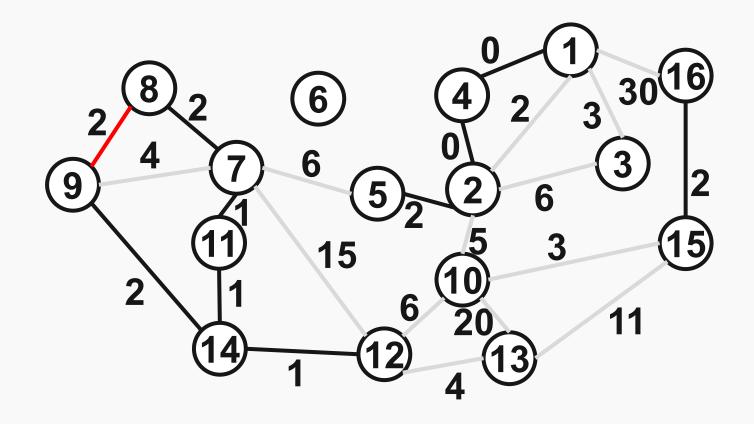




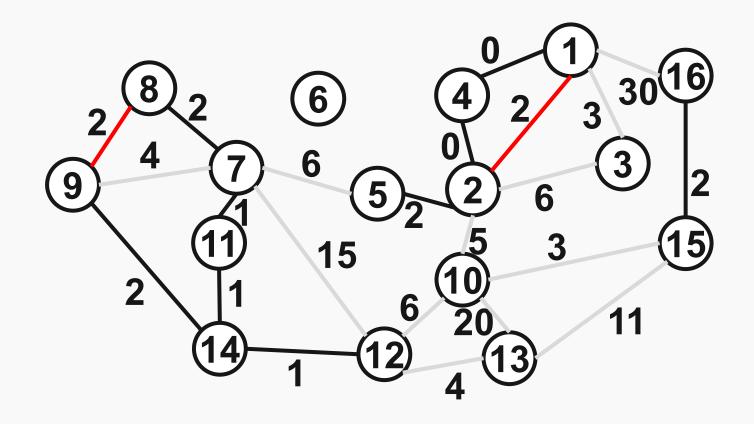




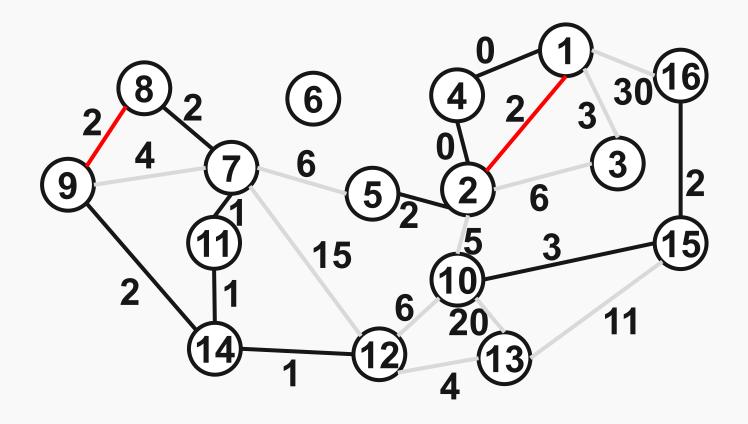




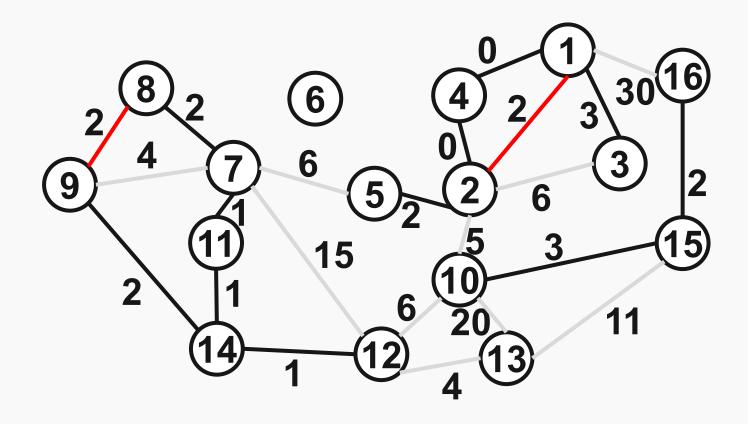




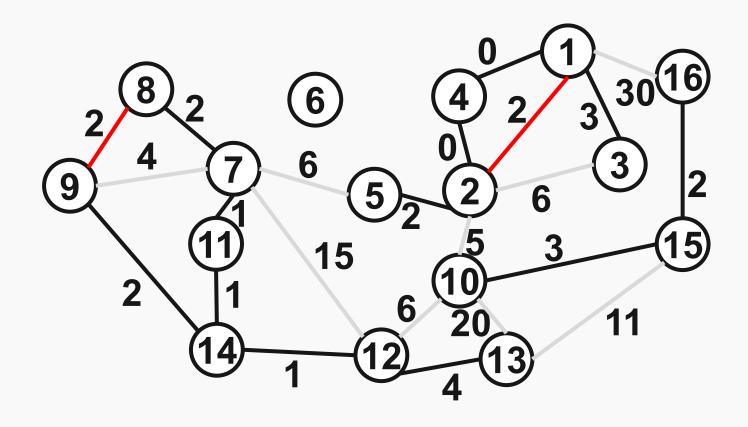




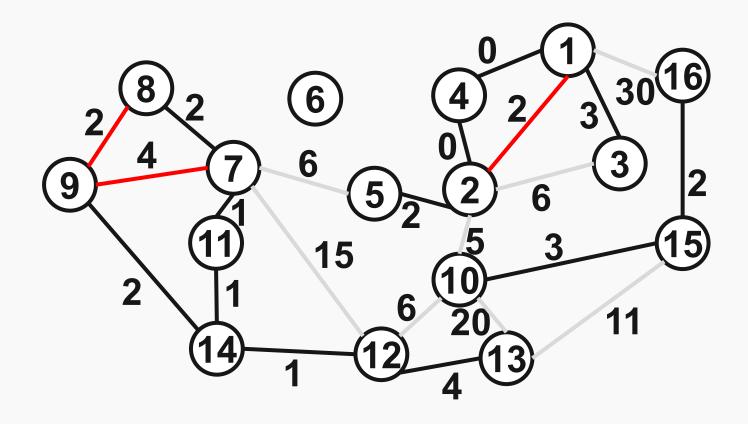




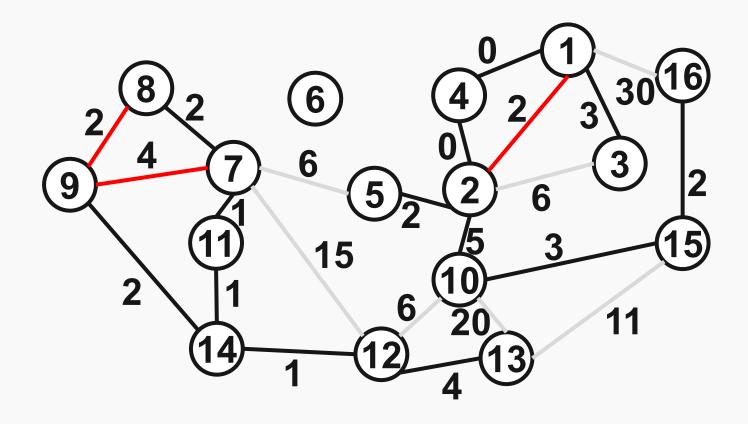




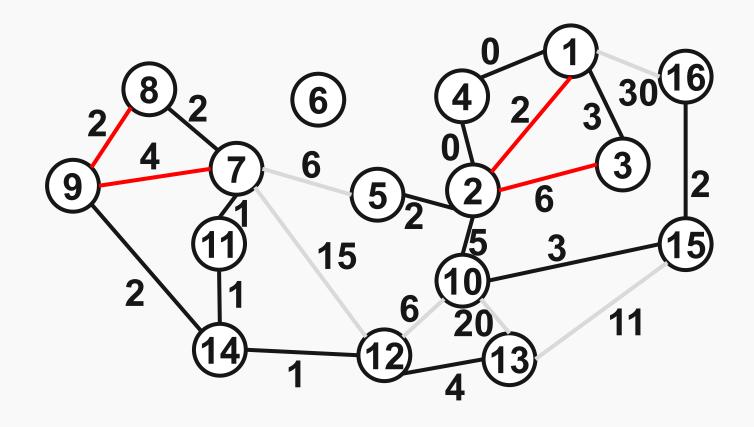




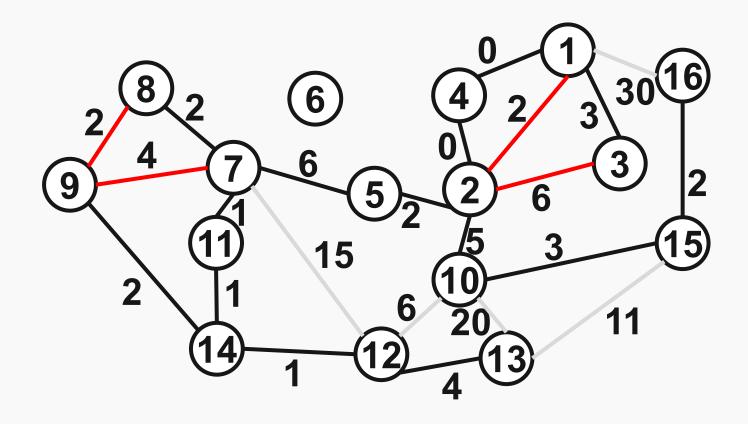




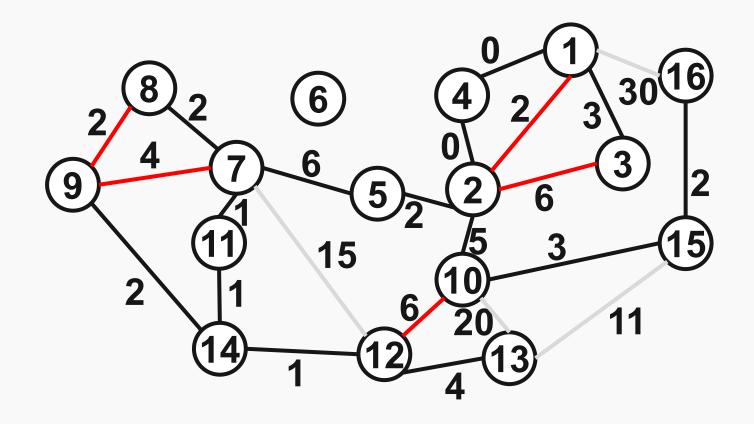




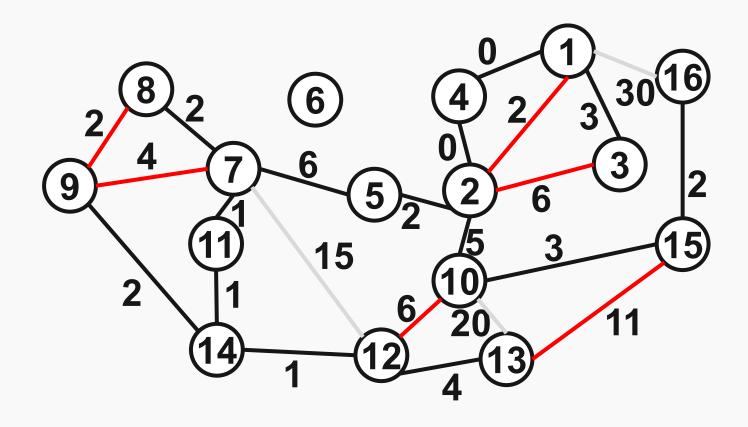




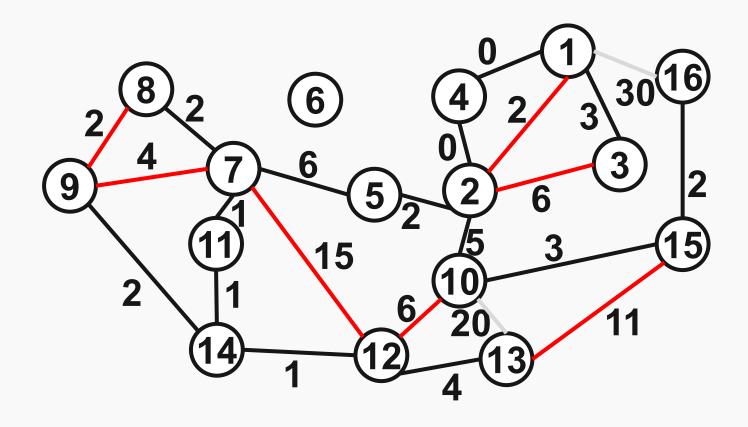




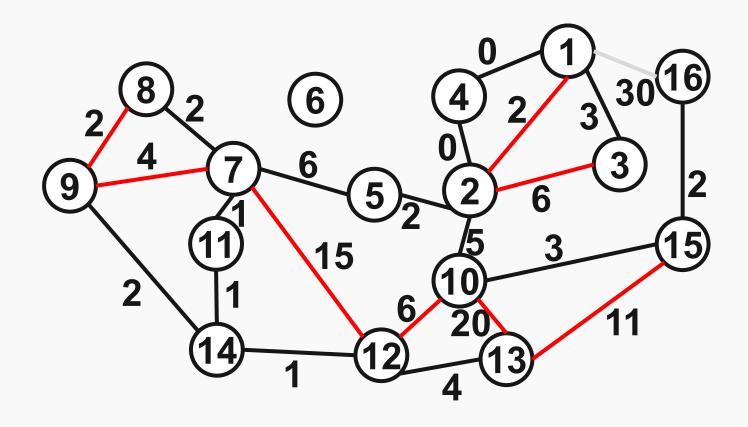




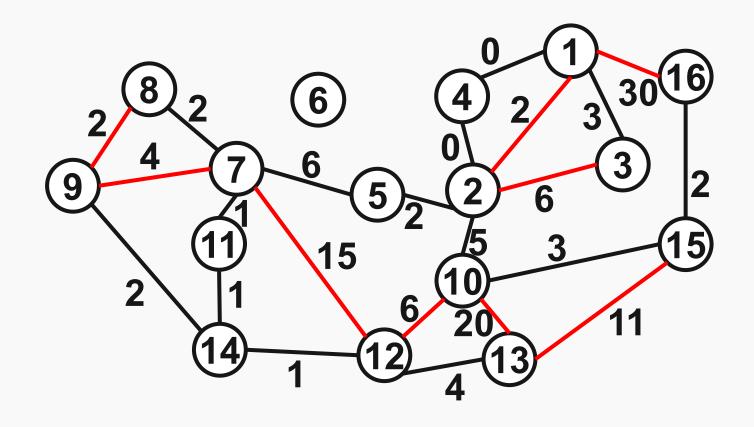




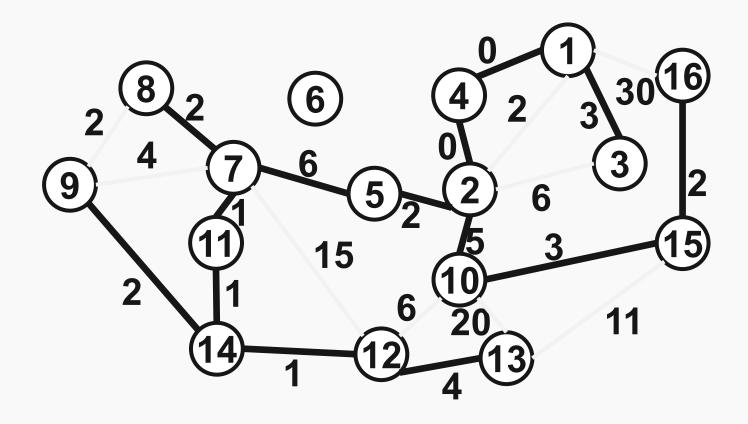














Complexitate?

```
tree Kruskal(G) {
    sort(G.E); // sort by weight
    A = \{\};
    for each (node in G.V)
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    for each ((u, v) in G.E) {
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            A = A \cup \{(u, v)\};
            Union(Find set(u), Find_set(v));
    return A;
```



Complexitate?

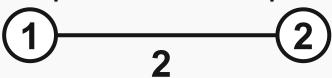


Flux maxim



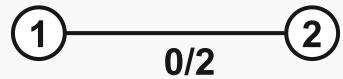
Graf capacitate

În general valoarea de pe muchie reprezintă o distanță.



O distanță mai mare face muchia mai greu de parcurs.

Valoarea muchiei poate reprezenta o capacitate.



- Similar apei/curentului, cu cât capacitatea e mai mare cu atât e mai ușor de parcurs.
- Şoselele au și distanță și capacitate (număr benzi/viteză max)



Flux maxim Algoritmul Ford-Fulkerson

- c capacitate muchie
- f flow muchie. Capacitate folosită
- $c_f(u,v) = c(u,v) f(u,v)$ diferența de capacitate
- G_f graful cu muchii c_f

```
FORD-FULKERSON(G, s, t)

1 for each edge (u, v) \in G.E

2 (u, v).f = 0

3 while there exists a path p from s to t in the residual network G_f

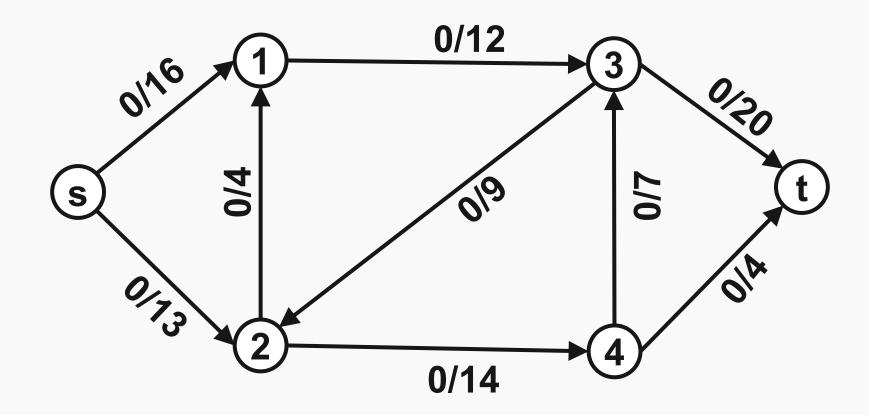
4 c_f(p) = \min\{c_f(u, v) : (u, v) \text{ is in } p\}

5 for each edge (u, v) in p

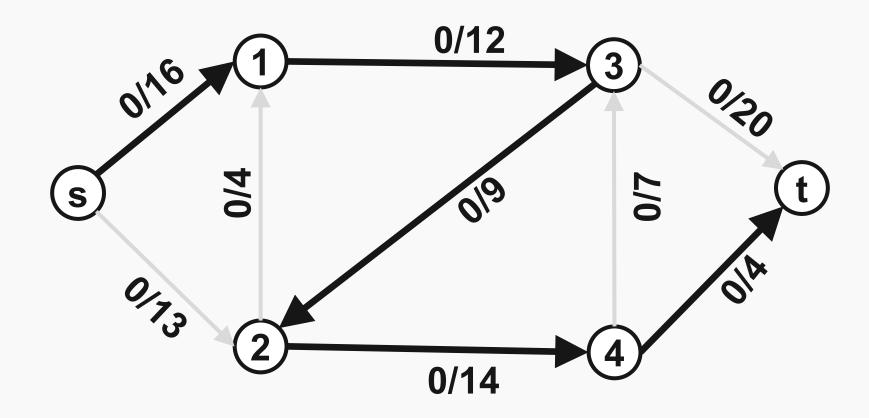
6 (u, v).f = (u, v).f + c_f(p)

7 (v, u).f = (v, u).f - c_f(p)
```

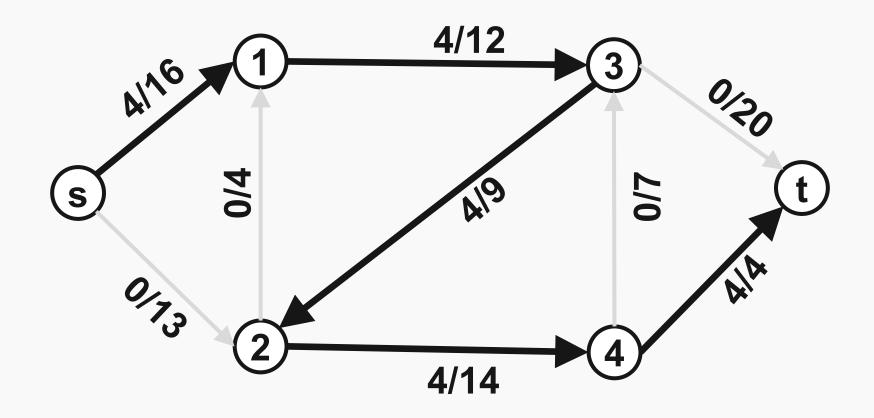




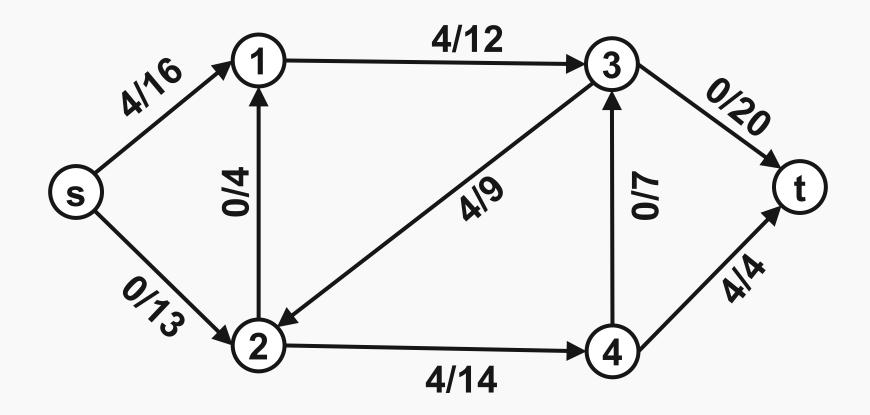




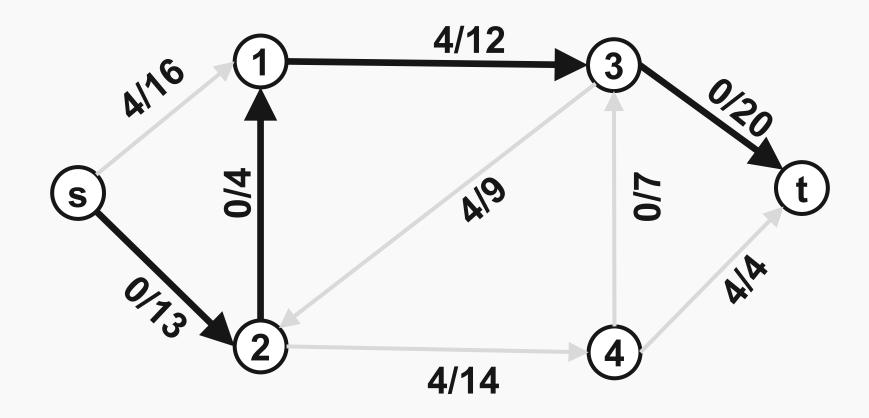




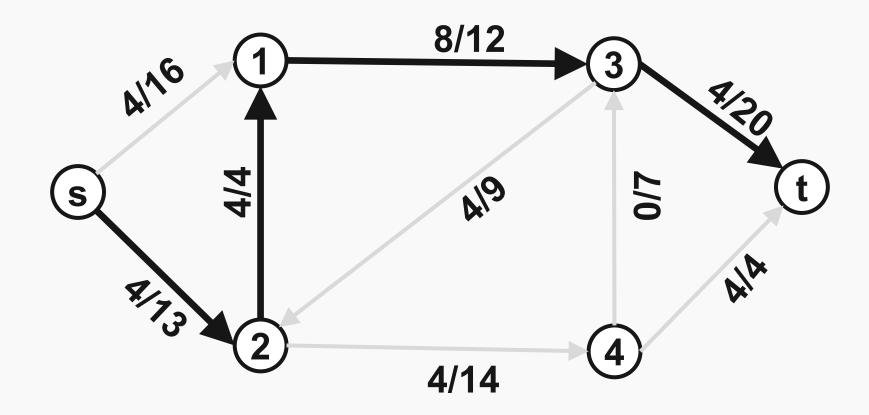




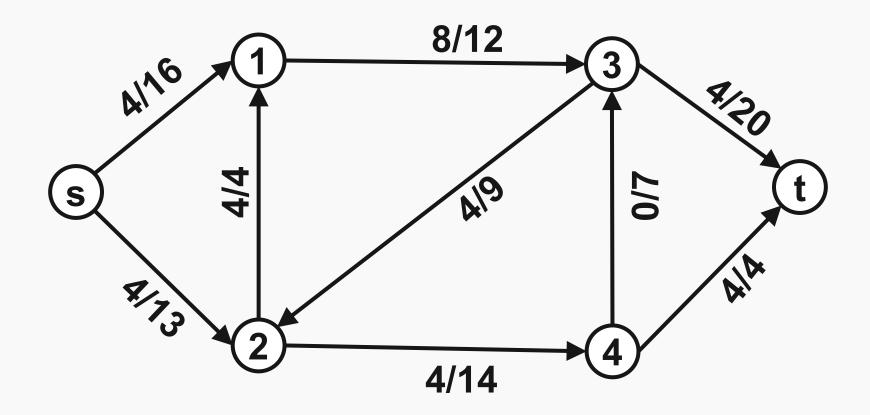




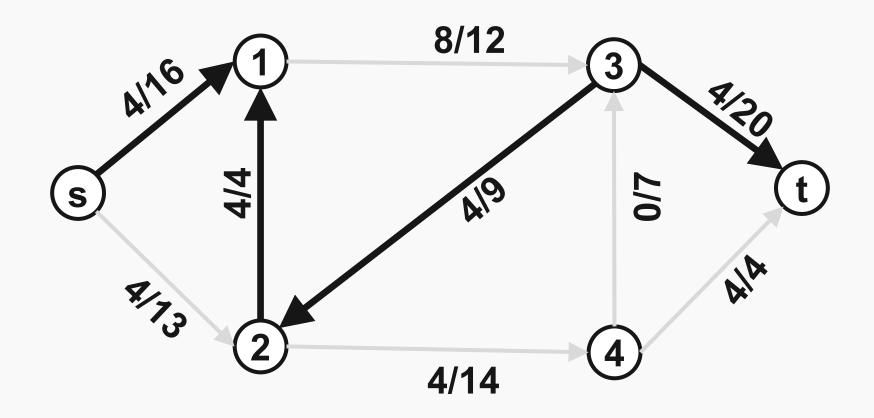




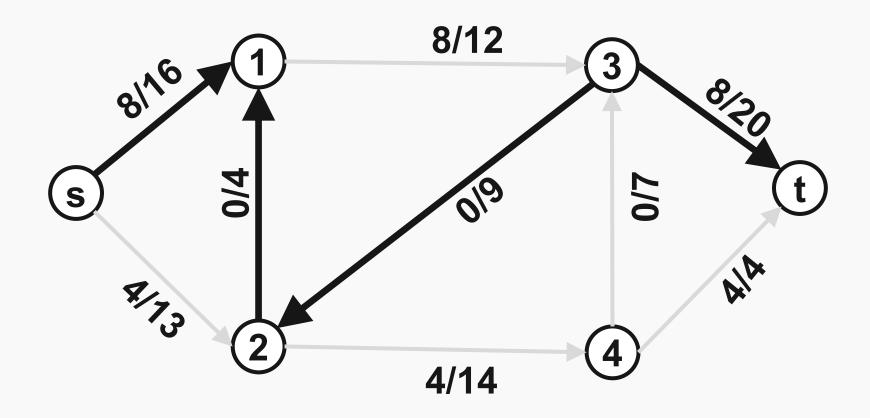




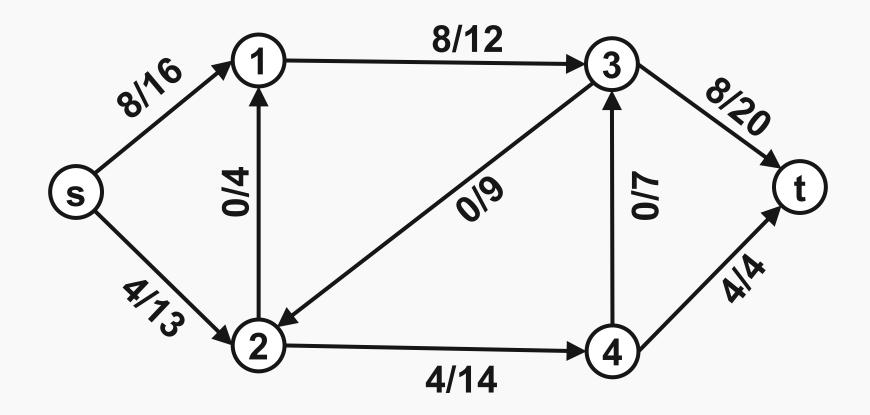




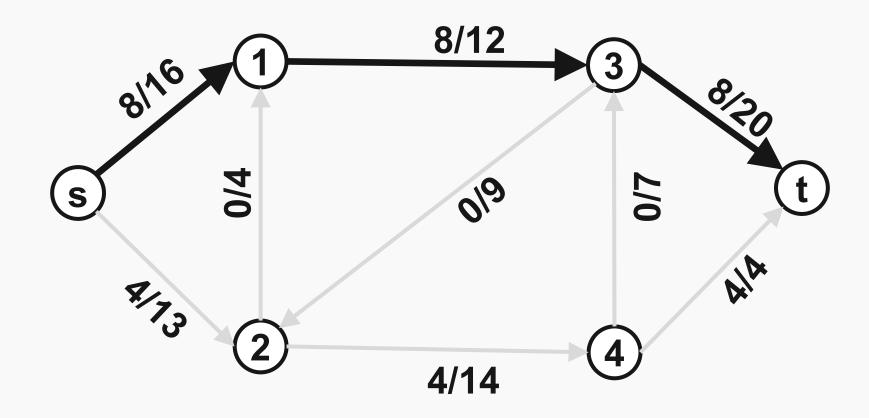




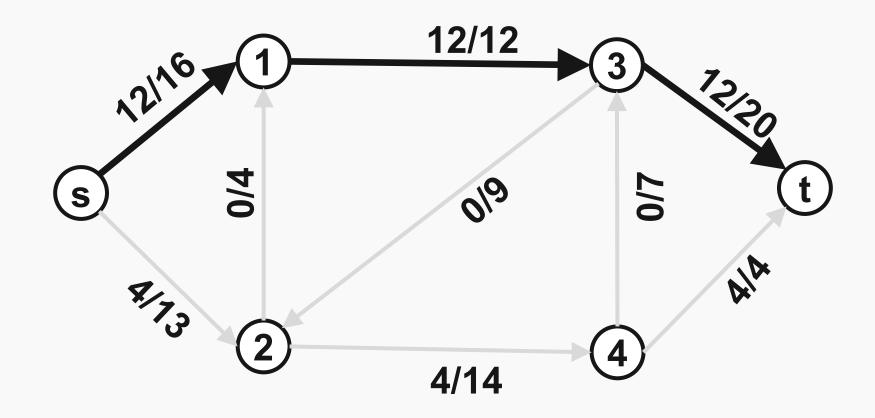




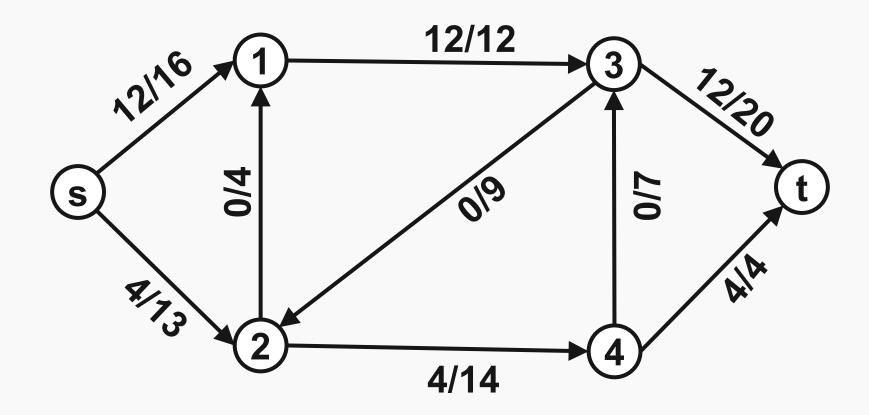




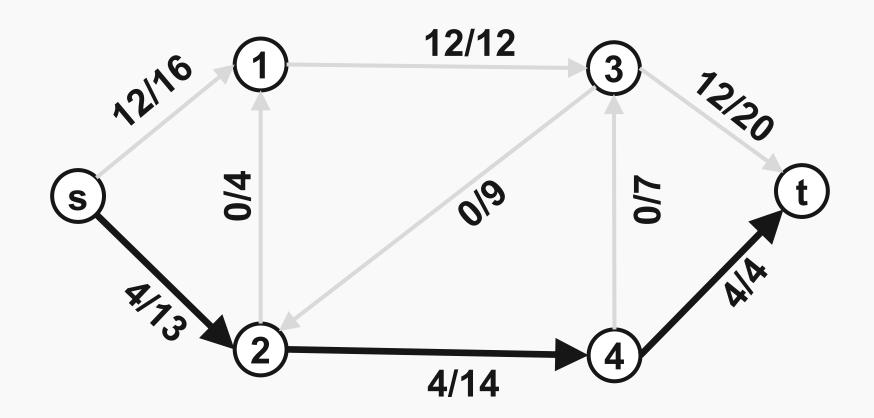




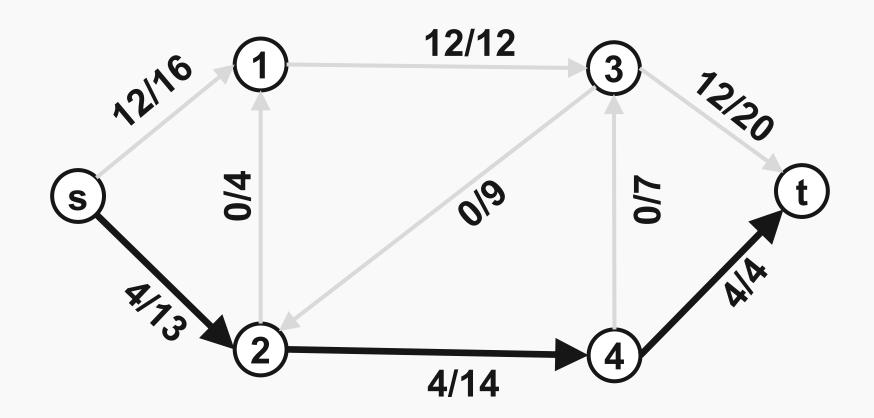




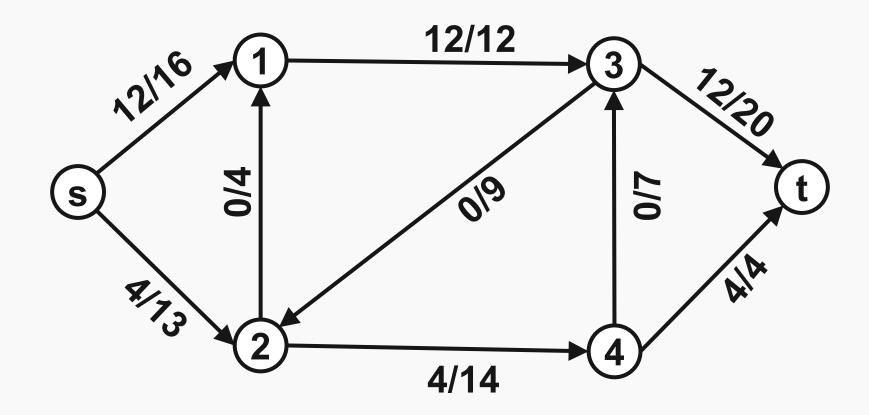




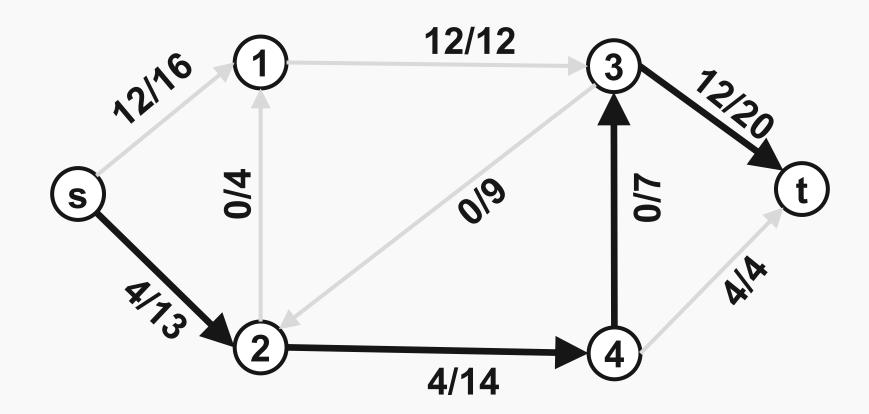




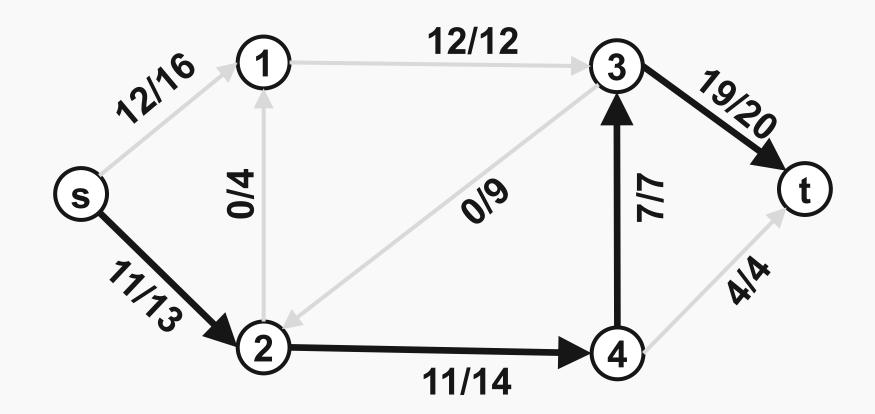




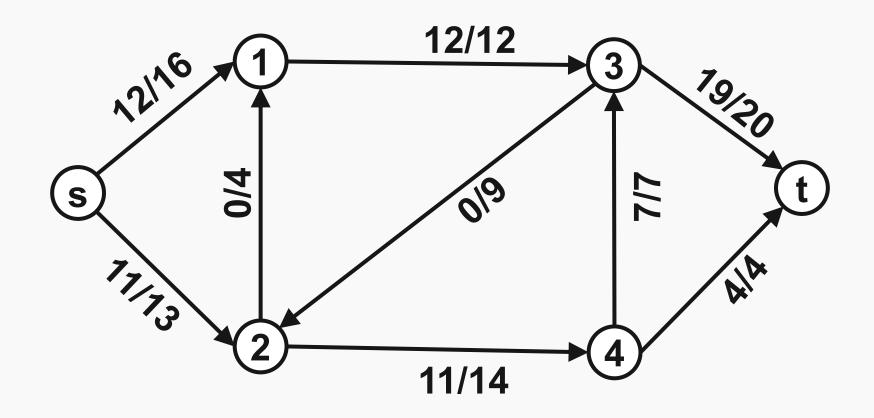














Alte considerente grafuri

Se pot schimba în timp.

LineGraph – Pentru un graf non-direcțional muchiile devin noduri și nodurile muchii.

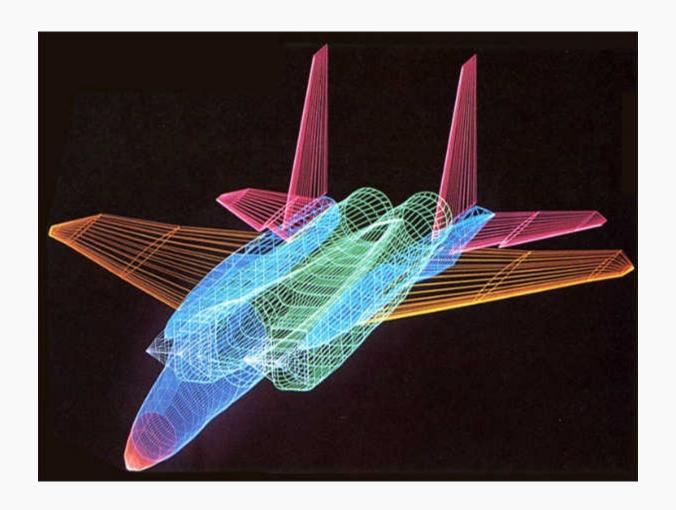


Use case grafuri – Granițe





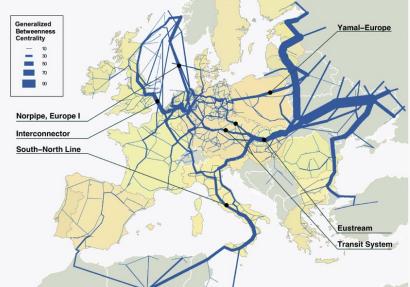
Use case grafuri – Grafică calculator





Use caser grafuri – utilități



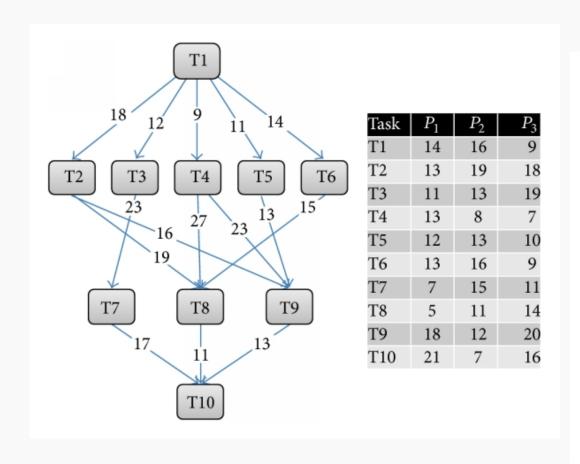


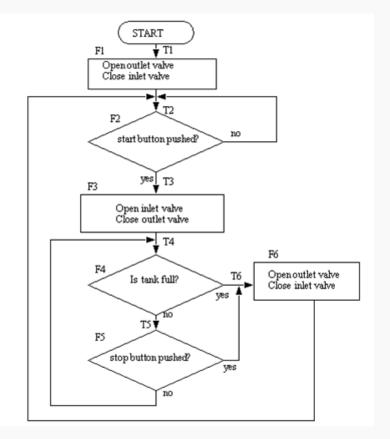






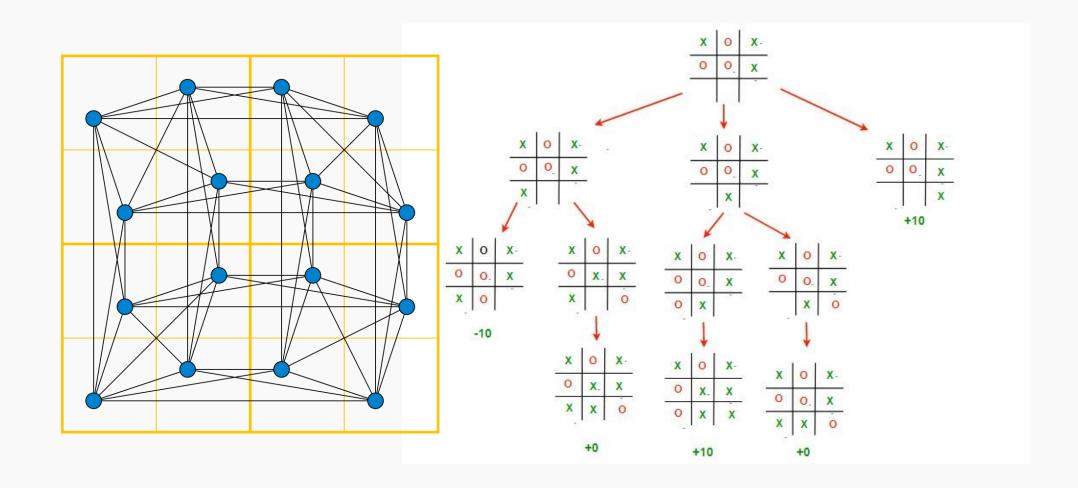
Use case grafuri – Code





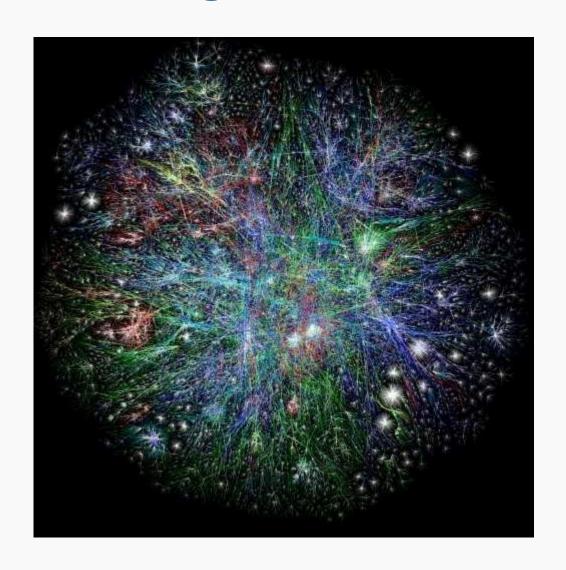


Use case grafuri – Reprezentare Jocuri



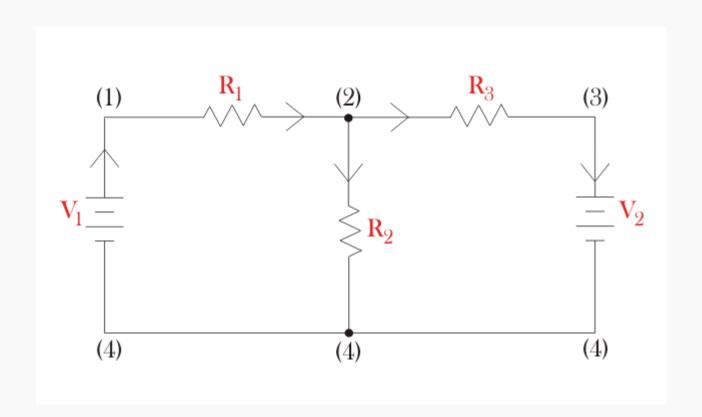


Use case grafuri - Internet



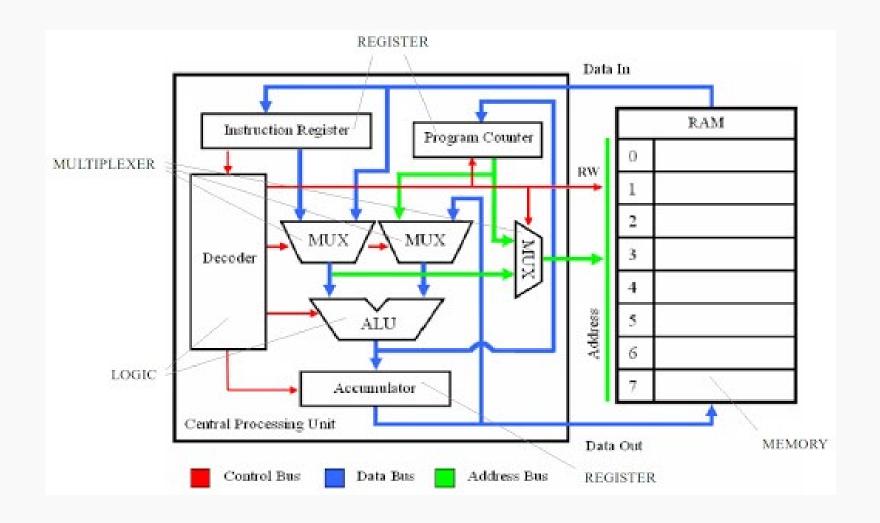


Use case grafuri – Circuite electrice



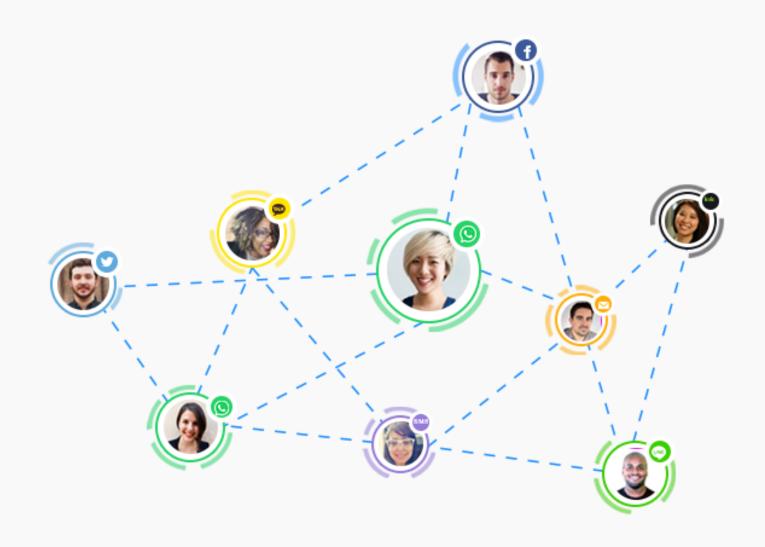


Use case grafuri – Circuite logice



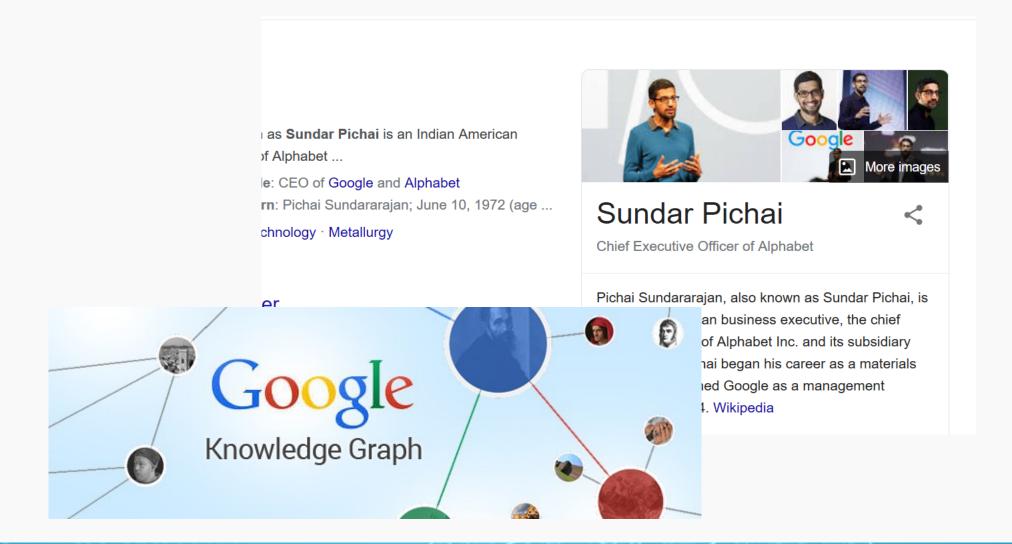


Use case grafuri – Grafuri sociale





Use case-uri – Knowledge graph





Use case-uri – Organigrame

