

# **Arhitecturi Paralele Pthread + Primitive sincronizare**

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Elemente preluate din cursul Prof. Ciprian Dobre





Bit level

Instruction level

Task parallelism

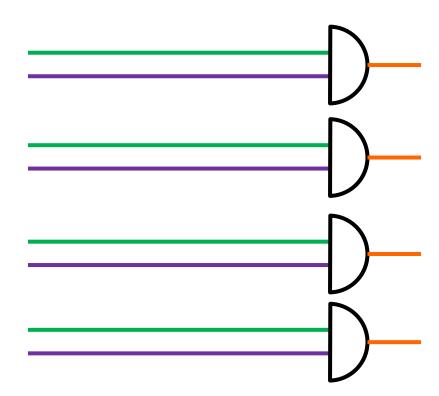


Bit level

$$\circ$$
 C = A & B

Instruction level

Task parallelism





Bit level

Instruction level

Task parallelism

$$C = A + B$$

Adunarea a doi vectori



Bit level

Instruction level

Task parallelism

- Multi-Tasking (pot comunica şi procesele)
- Multi-Threading



Bit level

Instruction level

Task parallelism

- Multi-Tasking (pot comunica şi procesele)
- Multi-Threading

Cum pot comunica două procese?



# Notații pseudocod

co S1 || S2 || ... || Sn oc  
Ex.1:  

$$x=0; y=0;$$
  
co  $x=x+1 || y=y+1$  oc  
 $z=x+y;$ 



# Notații pseudocod

co [cuantificator] {Sj}

Ex. 2:

**co**  $[j=1 \text{ to } n] \{a[j]=0; b[j]=0; \}$ 



## **POSIX** threads

```
pthread_t thread;
pthread_create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
pthread_join(thread, NULL);
```



# **Compilare pthread**

gcc -o executabil cod.c -lpthread -lrt

#include<pthread.h>

#include<semaphore.h>



```
pthread_t thread;
pthread_create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
               Acest element reprezintă thread-ul.
               Poate fi considerat handle
```



```
pthread_t thread;
pthread_create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
          Aici am putea să facem recomandări
          sistemului de operare. Gen să
          folosească anumite core-uri.
```



```
pthread_t thread;
pthread_create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
        Când se crează thread-ul, acesta va
        începe la această funcție.
```



```
pthread_t thread;
pthread_create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
        Aşa trimitem date thread-ului
```



```
pthread_t thread;
pthread_create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
         Astfel se pot extrage
         informații din thread
```





```
int main(int argc, char **argv)
 int i, P=2;
 pthread_t tid[P];
 for(i = 0; i < P; i++) {
   pthread_create(&(tid[i]), NULL, threadFunction, NULL);
  for(i = 0; i < P; i++) {
    pthread_join(tid[i], NULL);
  return 0;
```



```
int main(int argc, char **argv)
 int i, P=2;
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  return 0;
```



void\* threadFunction(void \*arg)

```
int main(int argc, char **argv)
                                               a = a + 1;
                                               return NULL;
 int i, P=2;
 pthread_t tid[P];
 for(i = 0; i < P; i++) {
   pthread_create(&(tid[i]), NULL, threadFunction, NULL);
  for(i = 0; i < P; i++) 
    pthread_join(tid[i], NULL);
  return 0;
```



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                                               return NULL;
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    pthread_join(tid[i], NULL);
  return 0;
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void\* threadFunction(void \*arg)

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  return 0;
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```
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                                                 a = a + 1;
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 int i, P=2;
 pthread_t tid[P];
 for(i = 0; i < P; i++) {
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  for(i = 0; i < P; i++) 
    pthread_join(tid[i], NULL);
                                            void* threadFunction(void *arg)
                                              a = a + 1;
  return 0;
                                              return NULL;
```



```
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  for(i = 0; i < P; i++) {
    pthread_join(tid[i], NULL);
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  for(i = 0; i < P; i++) {
    pthread_join(tid[i], NULL);
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                                              a = a + 1;
  return 0;
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int main(int argc, char **argv)
                                                 a = a + 1;
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 pthread_t tid[P];
 for(i = 0; i < P; i++) {
   pthread_create(&(tid[i]), NULL, threadFunction, NULL);
  for(i = 0; i < P; i++) 
    pthread_join(tid[i], NULL);
                                            void* threadFunction(void *arg)
                                              a = a + 1;
  return 0;
                                              return NULL;
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                               a = a + 1;
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 int I, P=2;
 pthread_t tid[P];
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  for(i = 0; i < P; i++) {
    pthread_join(tid[i], NULL);
  return 0;
```





Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

What is the value of a?



$$a = 0$$

Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

What is the value of a?



$$a = 0$$

Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

What is the value of a?

4



$$a = 0$$

Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

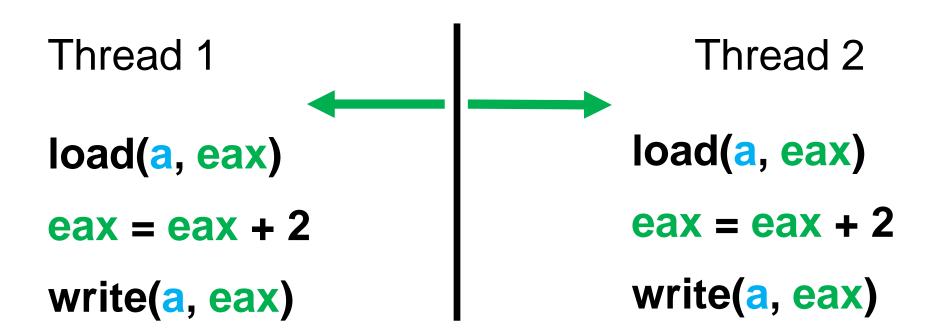
What is the value of a?

4 **AND** 2



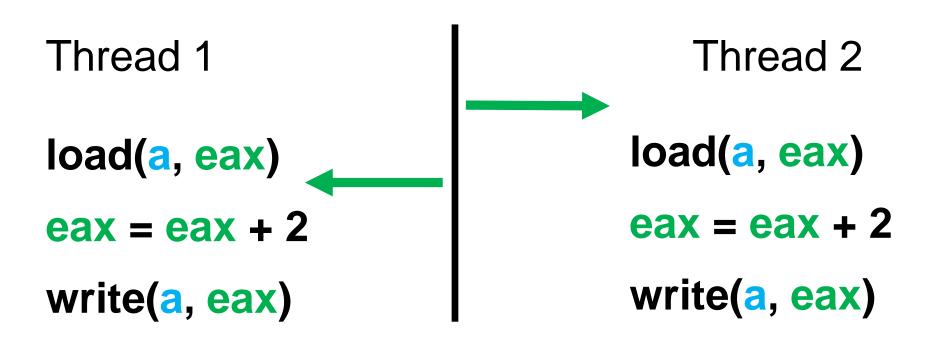


$$a = 0$$





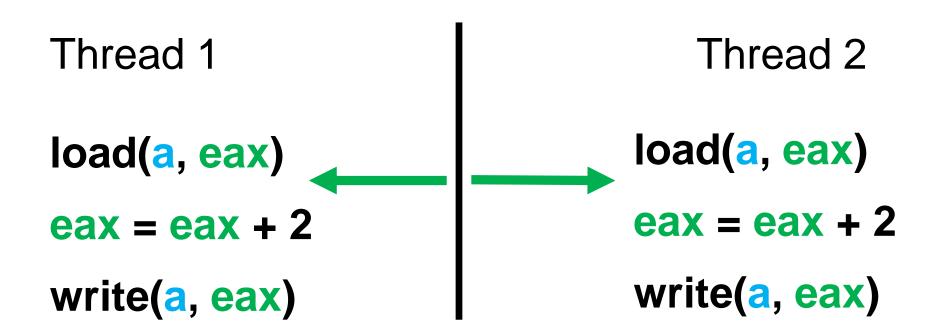
$$a = 0$$



$$eax = 0$$



$$a = 0$$

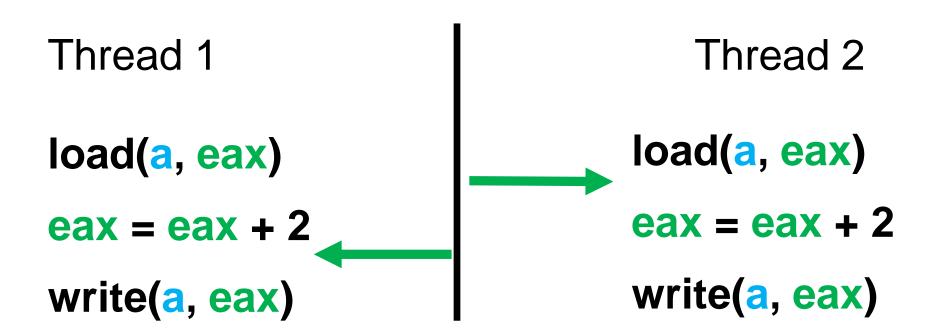


$$eax = 0$$

$$eax = 0$$



$$a = 0$$



$$eax = 2$$

$$eax = 0$$



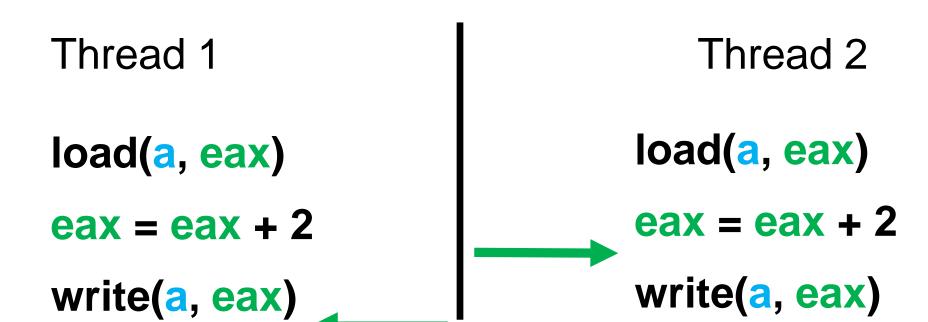
$$a = 0$$

$$eax = 2$$

$$eax = 2$$



$$a = 2$$



$$eax = 2$$

$$eax = 2$$



$$a = 2$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

eax = 2

Thread 2

load(a, eax)

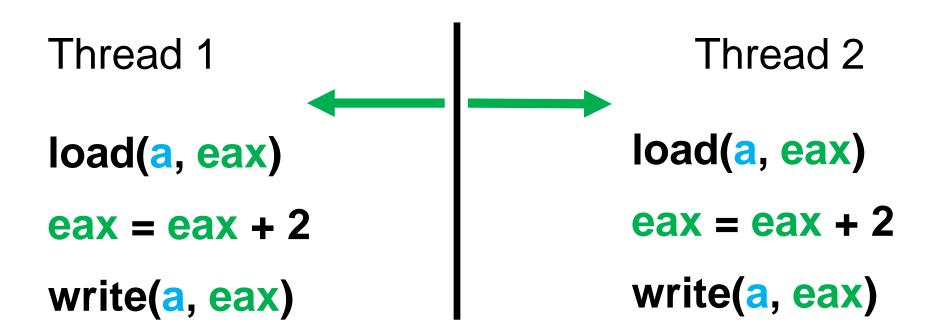
eax = eax + 2

write(a, eax)



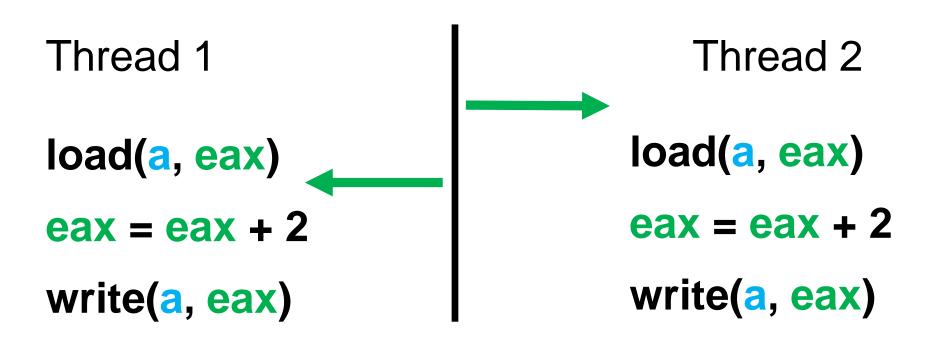


$$a = 0$$





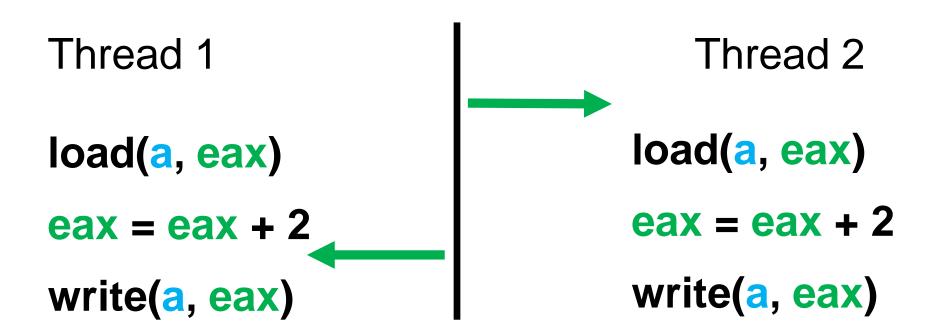
$$a = 0$$



$$eax = 0$$



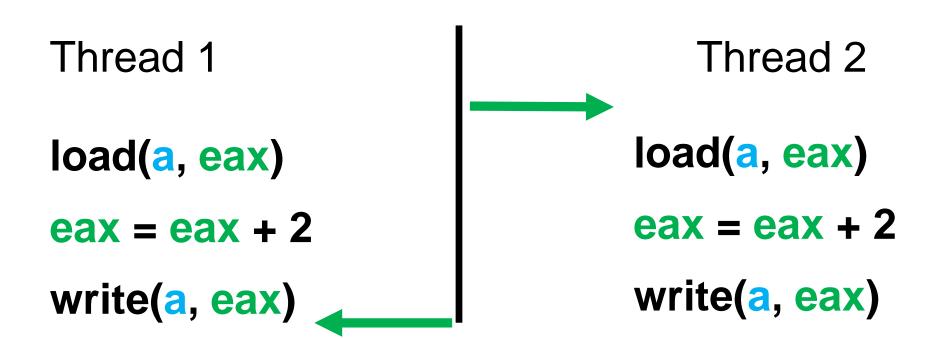
$$a = 0$$



$$eax = 2$$



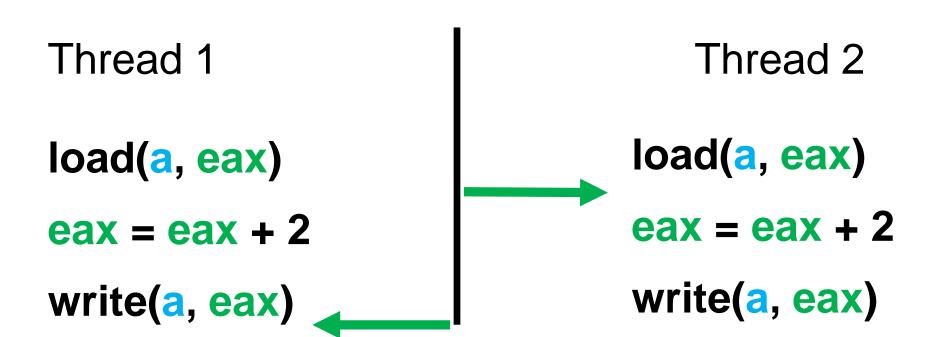
$$a = 2$$



$$eax = 2$$



$$a = 2$$

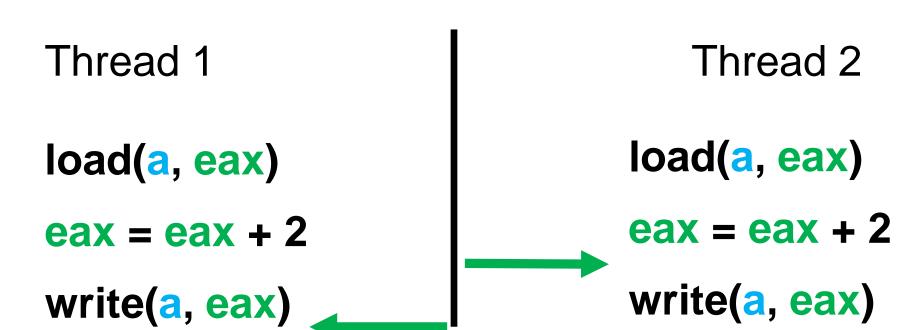


$$eax = 2$$

$$eax = 2$$



$$a = 2$$



$$eax = 2$$

$$eax = 4$$



$$a = 4$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

eax = 2

Thread 2

load(a, eax)

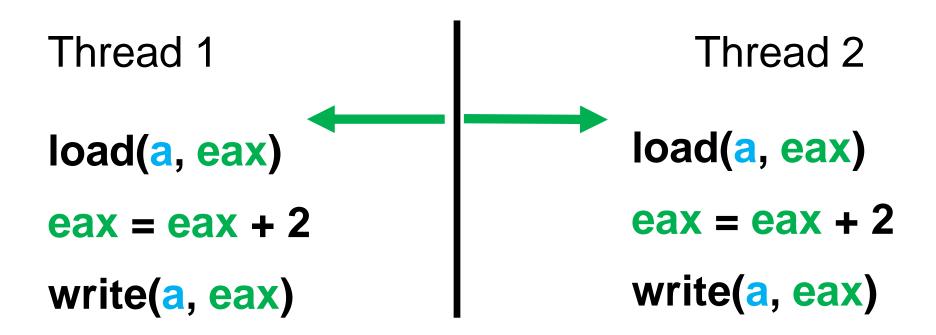
eax = eax + 2

write(a, eax)



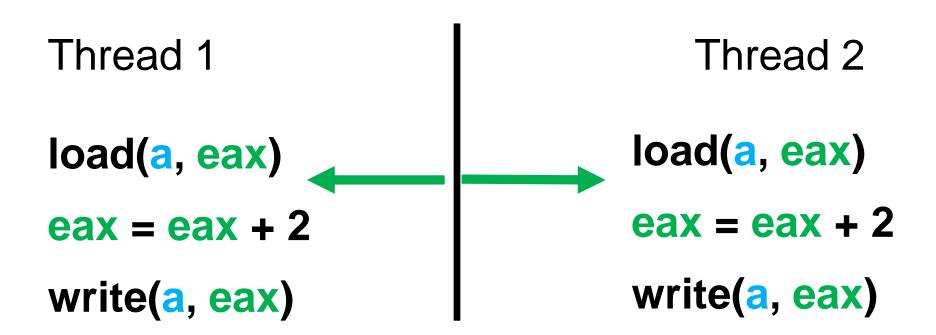


$$a = 0$$





$$a = 0$$

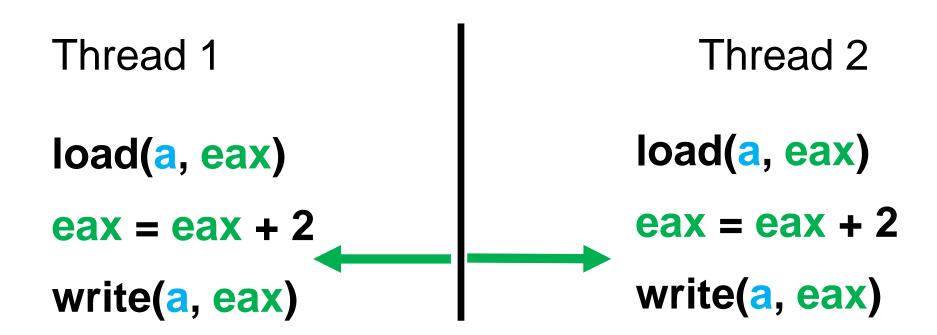


$$eax = 0$$

$$eax = 0$$



$$a = 0$$



$$eax = 2$$

$$eax = 2$$



$$a = 2$$

load(a, eax)

eax = eax + 2

write(a, eax)

eax = 2

Thread 2

load(a, eax)

eax = eax + 2

write(a, eax)



$$a = 2$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

eax = 2

Thread 2

load(a, eax)

eax = eax + 2

write(a, eax)



$$a = 2$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

eax = 2

Thread 2

load(a, eax)

eax = eax + 2

write(a, eax)

CREW





# Synchronization primitives

- Atomics
- Semaphore
  - -Binary semaphore (Mutex)
  - Critical section
- Barrier



## **Atomics**

■ Fie variabilelor de 64 biţi pe un procesor 64 biţi C = A + B

load(A, eax)

load(B, ebx)

eax = eax + ebx

write(C, eax)



#### **Atomics**

■ Fie variabilelor de 64 biţi pe un procesor 32 biţi C = A + B

```
load(A[0], eax)
load(B[0], ebx)
eax = eax + ebx
write(C[0], eax)
load(A[1], eax)
load(B[1], ebx)
eax = eax + ebx
write(C[1], eax)
```



#### **Atomics**

Fie variabilelor de 64 biţi pe un procesor 32 biţi C = A + B

```
load(A[0], eax)
load(B[0], ebx)
eax = eax + ebx
write(C[0], eax)
load(A[1], eax)
load(B[1], ebx)
eax = eax + ebx
write(C[1], eax)
```

Putem avea doar jumătate de C modificat



#### **Atomics**

■ Fie variabilelor de 64 biţi pe un procesor 32 biţi C = A + B

load(A[0], eax)

load(B[0], ebx)

eax = eax + ebx

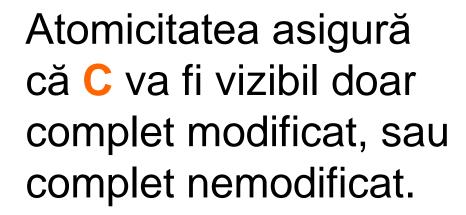
write(C[0], eax)

load(A[1], eax)

load(B[1], ebx)

eax = eax + ebx

write(C[1], eax)





### **Mutual exclusion - mutex**



#### Mutual exclusion - Dekker's solution

EWD35 - 1

EWD35.htm

About the sequentiality of process descriptions.

Over de sequentialiteit van procesbeschrijvingen.

Het is niet ongebruikelijk, wanneer een spreker zijn een inleiding. Omdat ik mij hier mogelijk richt tot een g minder vertrouwd is met de problematiek, die ik wil aansn die ik zal moeten gebruiken, wilde ik in dit geval ter in inleidingen houden, nl. een om de achtergrond van de prob een tweede, om U een gevoel te geven voor de aard van de wij tegen het lijf zullen lopen.



Theodorus (Dirk) J. Dekker



# Mutual exclusion - Dijsktra's Solution

#### Solution of a Problem in Concurrent Programming Control

E. W. Dijkstra Technological University, Eindhoven, The Netherlands

A number of mainly independent sequential-cyclic processes with restricted means of communication with each other can be made in such a way that at any moment one and only one of them is engaged in the "critical section" of its cycle.

#### Introduction

Given in this paper is a solution to a problem for which, to the knowledge of the author, has been an open question since at least 1962, irrespective of the solvability. The paper consists of three parts: the problem, the solution, and the proof. Although the setting of the problem might seem somewhat academic at first, the author trusts that anyone familiar with the logical problems that arise in

computer can only request one one-way message at a time.

And only this w
this problem is f

#### The Solution

The common "Boo

The integer will only be set it by the others. It well outside the mentioned set to

The program

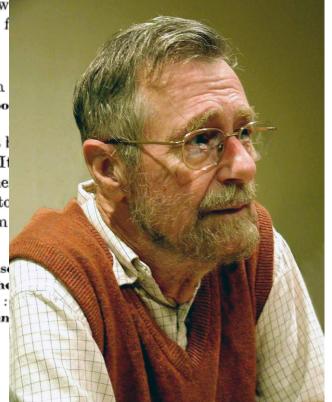
#### "integer j;

Li0: b[i] := fals $Li1: if k \neq i the$ 

Li2: begin c[i]: Li3: if b[k] then

go to Lil

else





#### **Dekker's Solution**

```
wants_to_enter[0] = true
while wants_to_enter[1] {
      if turn != 0 {
            wants_to_enter[0] = false
            while turn != 0 { // busy wait }
            wants_to_enter[0] = true
// critical section ...
turn = 1
wants_to_enter[0] = false
```



## **Dijsktra's Solution**

```
b[i] = fals
while(sw[i]) {
          sw[i] = F
          if (k!=i) {
                     c[i] = true
                     if(b[k])
                                k = i
                     sw[i] = T
          } else {
          c[i] = false
          for(j=0;j<N;j++)
                     if(j!=i && !c[j])
                                sw[i] = T
//critical
b[i] = true
c[i] = true
```

lock()

unlock()



#### **Peterson's Solution**

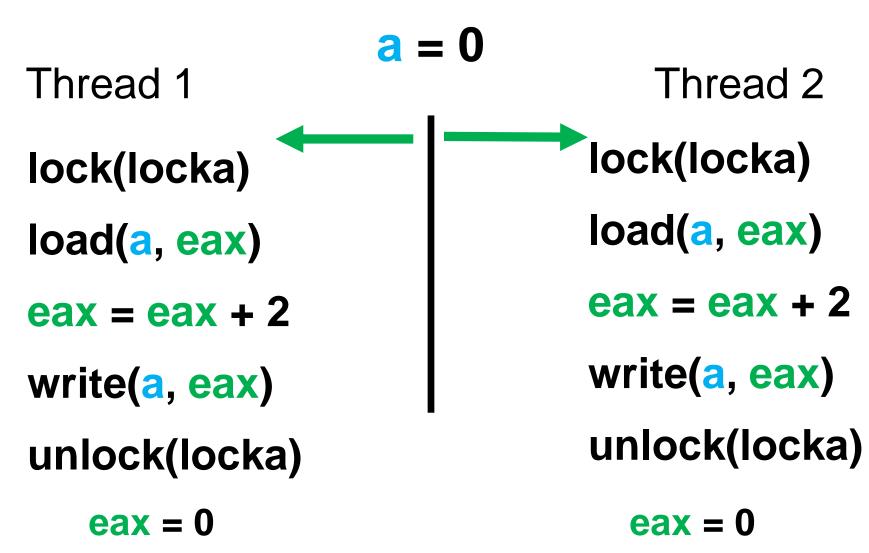
```
flag[0] = true;
P0_gate: turn = 1;
while (flag[1] == && turn == 1) { // busy wait }
// critical section ...
flag[0] = false;
```



#### **Hardware assisted Solution**

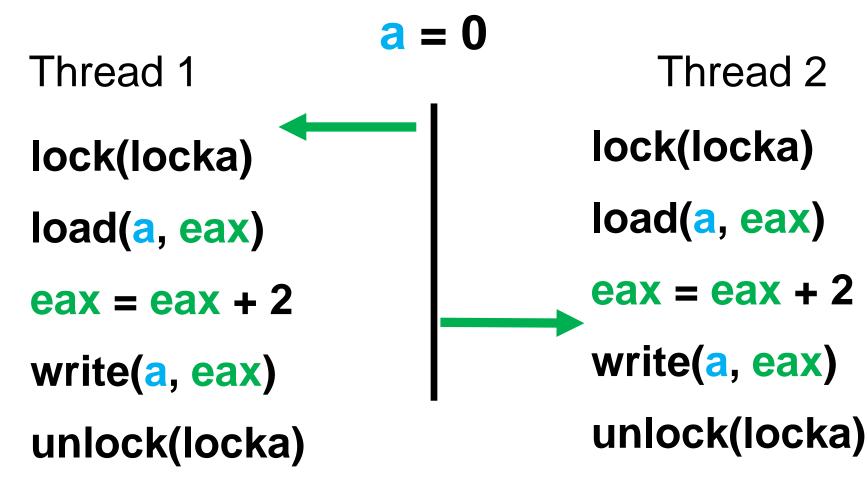
```
while (test_and_set(lock));
// critical section
lock = 0
```







eax = 2







$$a = 0$$

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 0

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)



$$a = 2$$
Thread 1

Thread 2

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load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 0

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)





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unlock(locka)

eax = 0

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

 $\mathbf{a} = \mathbf{0}$ 



Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)





$$a = 0$$

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)





lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

# a = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)





$$a = 2$$

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)





$$a = 4$$

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)





$$a = 4$$

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

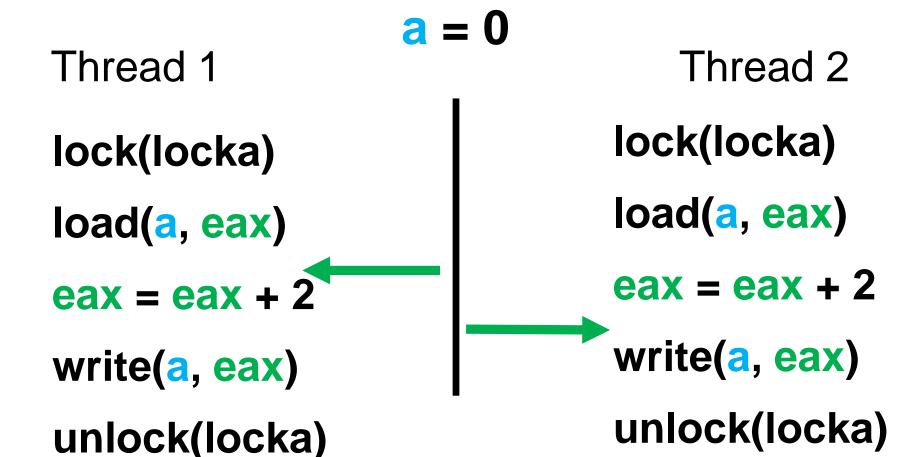
load(a, eax)

eax = eax + 2

write(a, eax)

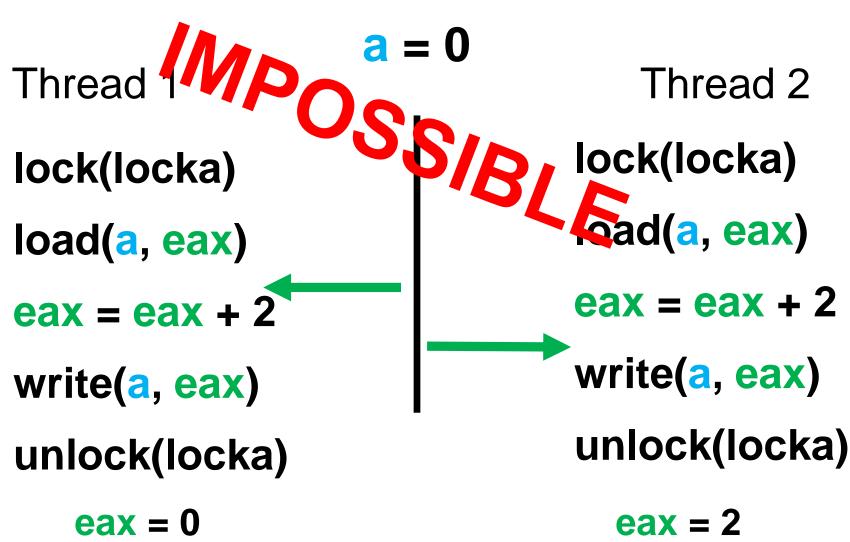
unlock(locka)





$$eax = 2$$







# ÎN MAIN

Înainte de a porni thread-urile

```
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL);
```



```
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL);
```

Poate fi folosit să anunțe că acest mutex e împărțit mai multor procese



```
pthread_mutex_lock(&mutex);
load(a, eax)
eax = eax + 2
write(a, eax)
pthread_mutex_unlock(&mutex);
```



# ÎN MAIN

După ce au terminat thread-urile

pthread\_mutex\_destroy(&mutex);





# ÎN MAIN

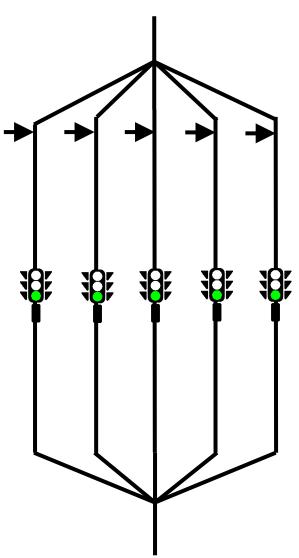
Înainte de a porni thread-urile

```
sem_t semaphore;
int semaphore_value= 4;
sem_init(& semaphore, 0, semaphore_value);
```



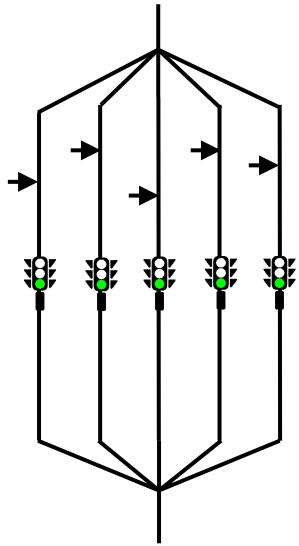


P() sau Proberen - Dijsktra



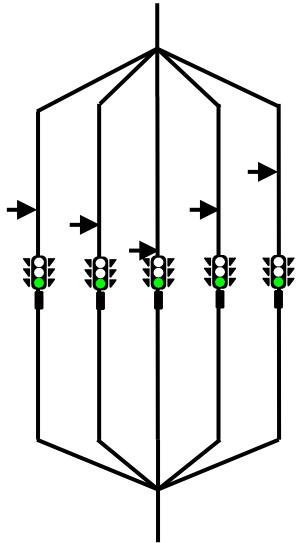






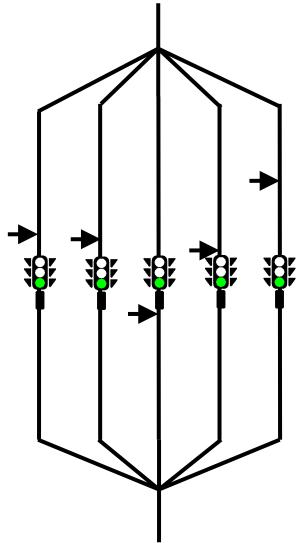








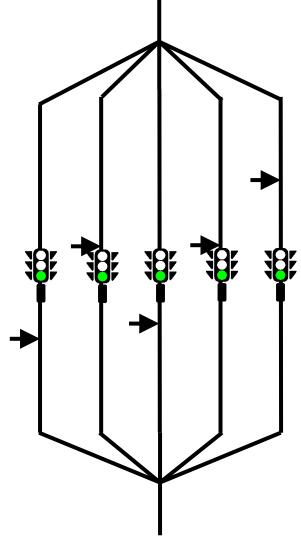






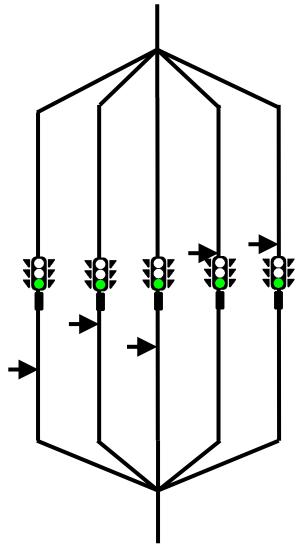
Nu contează că doua thread-uri au ajuns simultan la semafor, acesta este protejat, la fel ca un mutex.

sem\_wait(&semaphore);



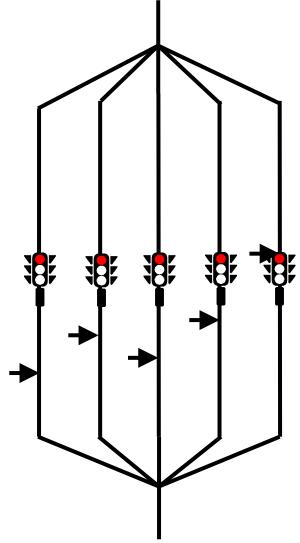








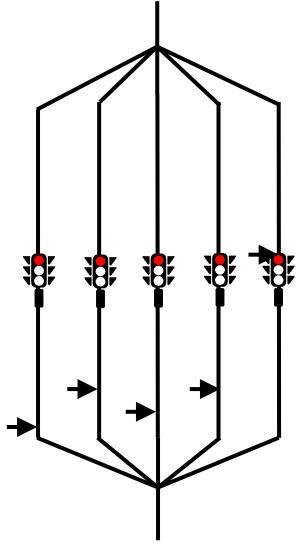






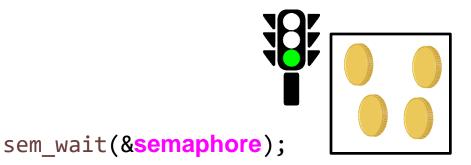
# **Semaphore**

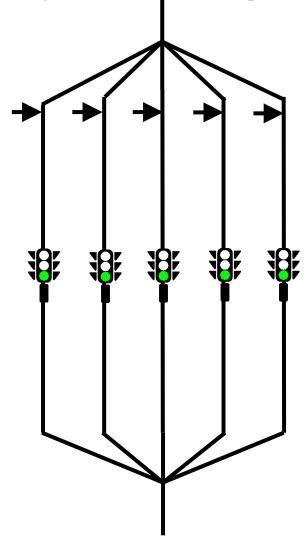




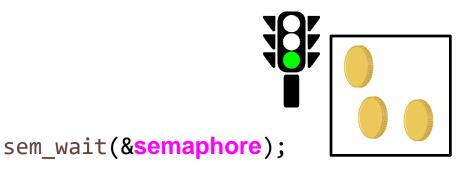


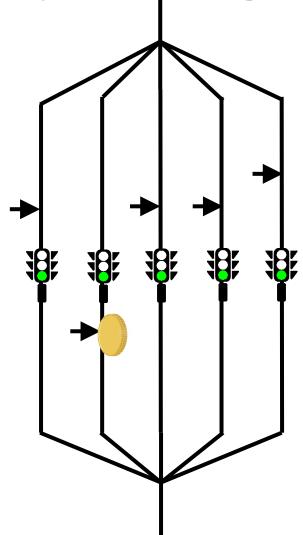
Un semafor are un set de token-uri.



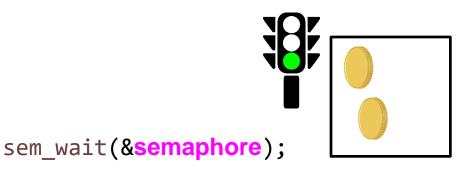


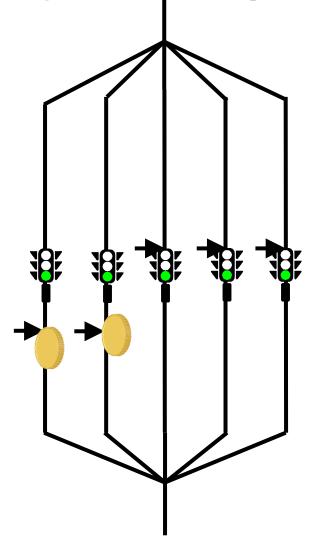




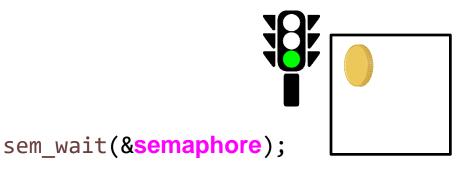


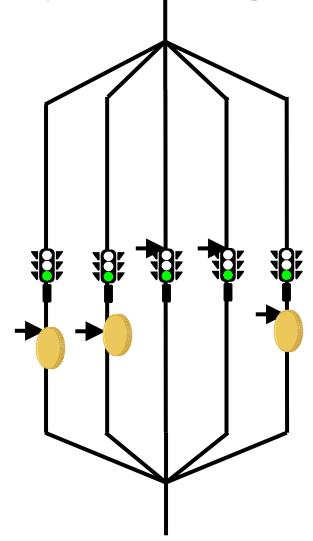






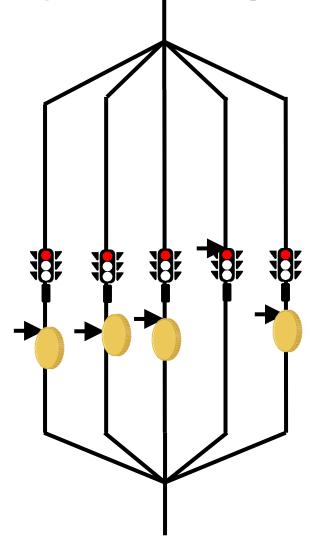














**Semaphore – Signaling** sem\_wait(&semaphore); sem\_post(&semaphore); sem\_post(&semaphore); V() sau Verogen - Dijsktra













# **Semaphore**

# ÎN MAIN

După ce au terminat thread-urile

sem\_destroy(& semaphore);







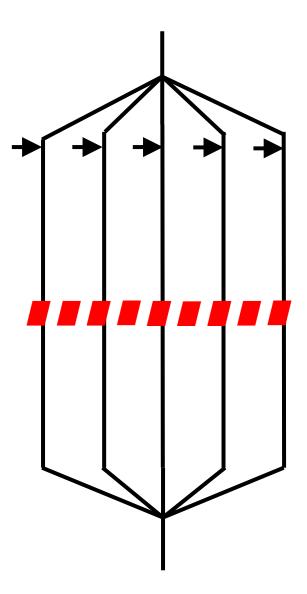
# ÎN MAIN

Înainte de a porni thread-urile

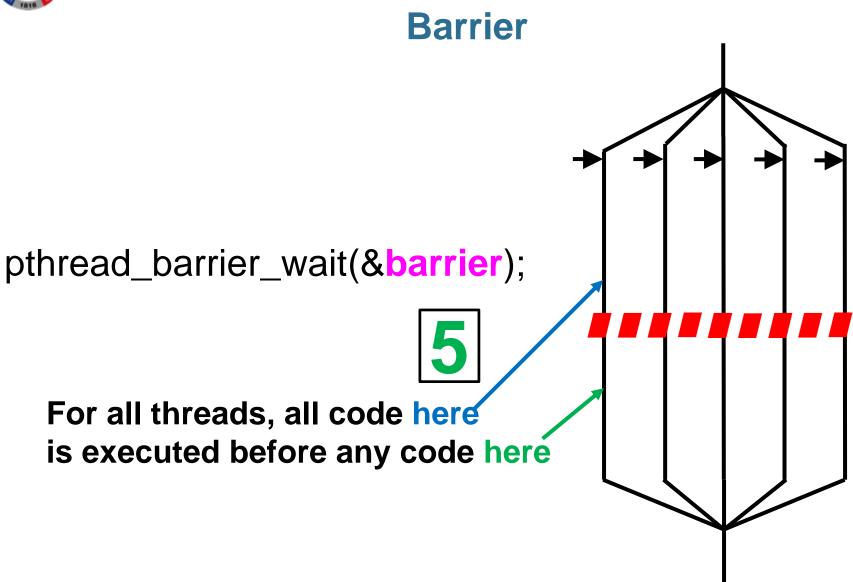
```
pthread_barrier_t barrier;
int num_threads = 5;
pthread_barrier_init(&barrier, NULL, num_threads);
```





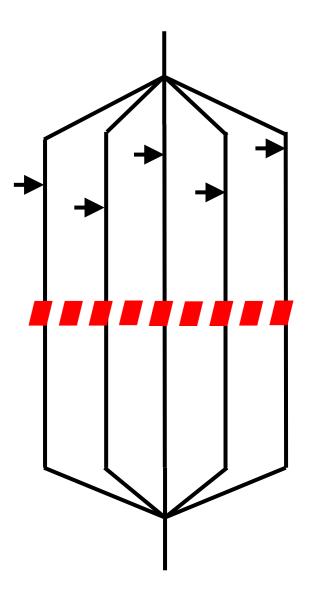






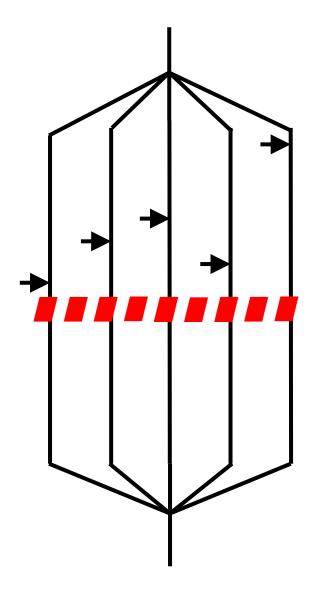






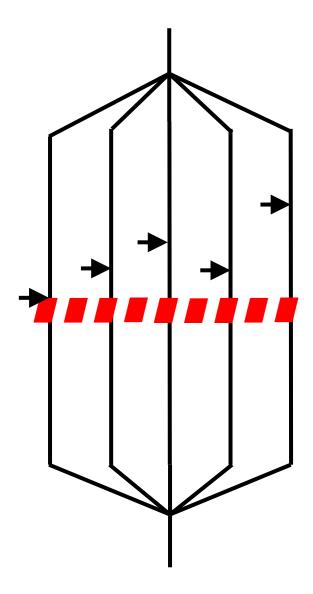






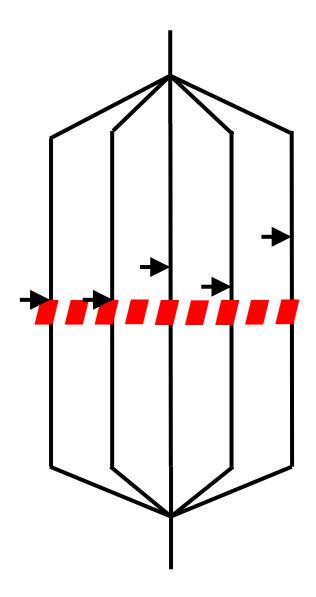






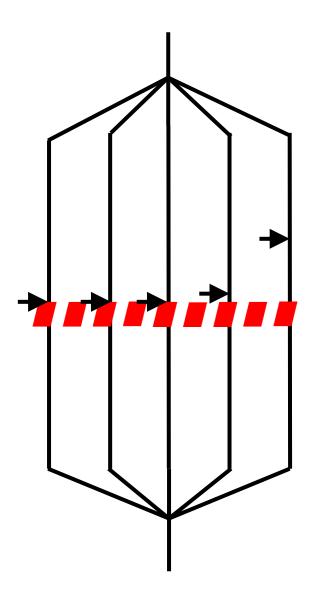






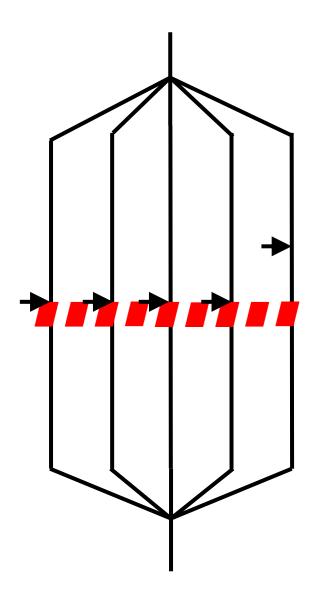






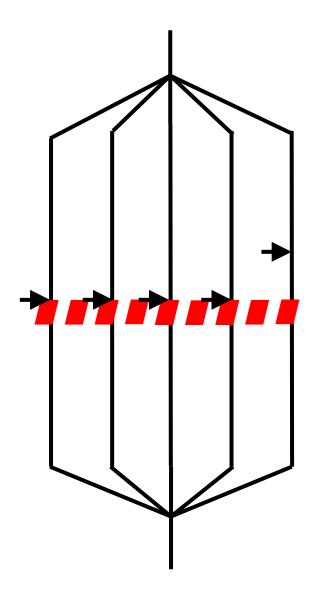






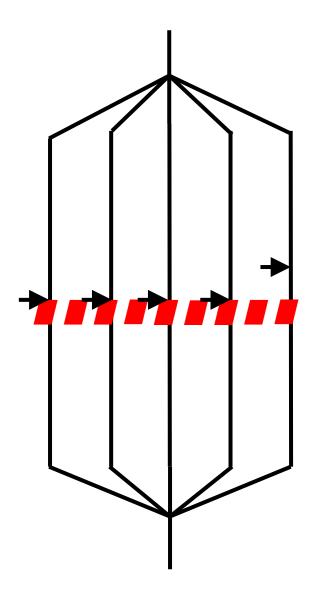






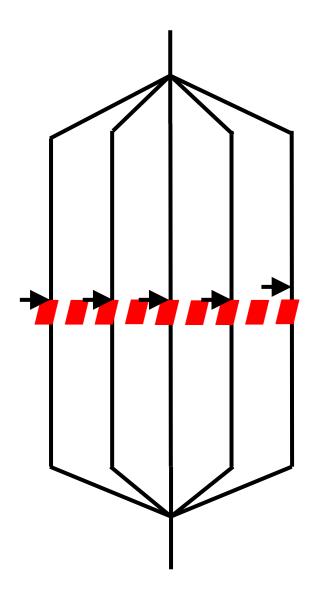






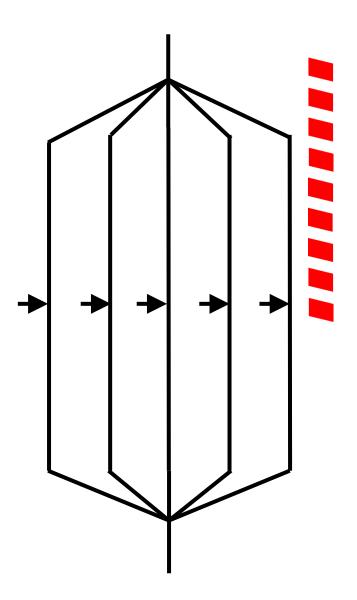






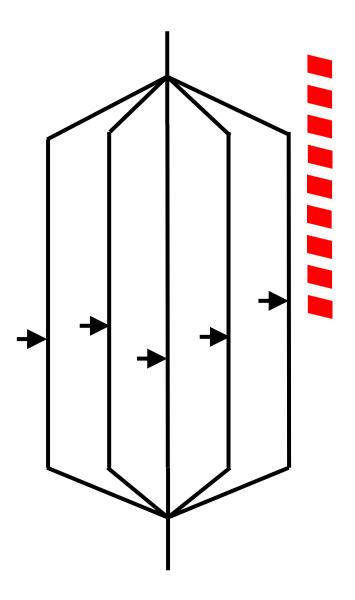












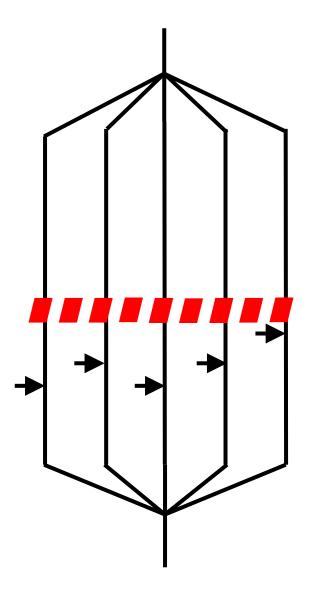


Cum știe o barieră când să se reseteze?

O soluție ar fi: Reusable Barrier in The Little Book of Semaphores By Allen B. Downey









# ÎN MAIN

După ce au terminat thread-urile

pthread\_barrier\_destroy(&barrier);





# Some problems can not be parallelized

Calculating the hash of a hash of a hash ... of a string.

Deep First Search

Huffman decoding

Outer loops of most simulations

P complete problems



# Splitting a problem to make it parallel

Sunt o serie de probleme care sunt extrem de ușor paralelizabile.

**Embarrassingly parallel** 



### **Embarrassingly parallel problems**

Multiplicare unui vector cu un scalar

\* 3

27 18 27 12 6 21 18 15 18 3



### **Embarrassingly parallel problems**

Toate calculele pot fi efectuate în același timp

\* 3

27 18 27 12 6 21 18 15 18 3



# **Embarrassingly parallel problems**

Câte elemente sunt?





Câte elemente sunt?





Câte elemente sunt? N

969427656 ...



Dar câte elemente de procesare?





Dar câte elemente de procesare?

9 6 9 4 2 7 6 5 6 ...



Dar câte thread-uri?





Dar câte thread-uri?



În majoritatea cazurilor obținem performanță maximă când numărul de thread-uri este egal cu numărul de elemente de procesare, sau core-uri.



Cum este P față de N?



P << N

969427656 ....

1



Caz concret: P = 2 Cum împărțim?

969427656 ....

Florin Pop – Cristian Chilipirea



Caz concret: P = 2 Cum împărțim?

1

Thread 1



Caz concret: P = 2 Cum împărțim?

1

Thread 1



Caz concret: P = 2 Cum împărțim?

1

Thread 1



Caz concret: P = 2 Cum împărțim? Putem și random

Thread 1



Caz concret: P = 2 Cum împărțim?

1

Thread 1



Caz concret: P = 2 Cum împărțim?

Este utilă?

Thread 1



Caz concret: P = 2 Cum împărțim?

Ce ne dorim?

Thread 1



Ce ne dorim?

Thread 1

Thread 2

Aproximativ același număr elemente



# **Aproximativ N/P elemente**

Thread 1



# **Aproximativ N/P elemente**

# Dacă N nu se divide perfect la P?

Thread 1



# Aproximativ N/P elemente Dacă N nu se divide perfect la P?

1

6

4 2 7

9 6 9

Thread 1

8

4 9 2

5 6 3



# floor(N/P) elemente floor(15/2) = 7

1

6

4 2 7

9 6 9

Thread 1

8

4 9 2

5 6 3



ceil(N/P) elemente ceil(15/2) = 8

6 5

4 2 7

9 6 9

Thread 1

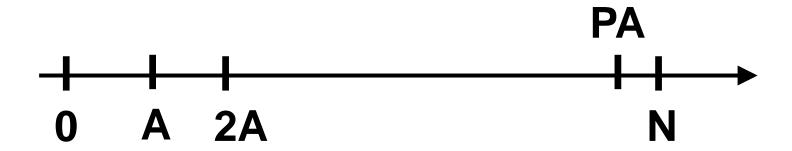
8 1

4 9 2

6 3

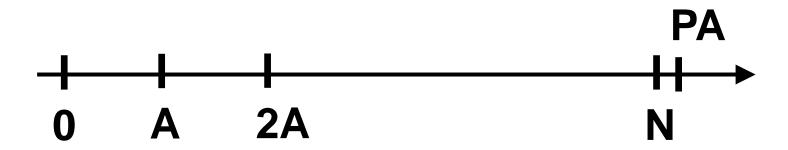


$$A = floor(N/P)$$





$$A = ceil(N/P)$$





#### Formule elegante:

Tid este identificator de thread, are valori de la 0 la P

