

Arhitecturi Paralele Introducere

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Curs susținut în parteneriat cu Prof. Florin Pop







Regulament

- Laboratoarele sunt obligatorii
- Temele se rezolvă individual şi vor fi verificate anti-plagiat



Punctaje

6p Teme

Minim 3p pentru promovare

2p Examen parțial

Minim 1p pentru promovare

2p Examen final

Minim 1p pentru promovare



Obiective

Dezvoltarea abilităților pentru:

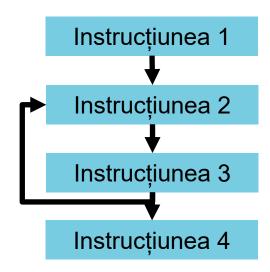
- Proiectarea și implementarea aplicațiilor multi-thread
- Proiectarea și implementarea aplicațiilor distribuite
- Depanarea unor aplicații multi-thread sau distribuite
- Demonstrarea corectitudinii şi scalabilităţii unui program multi-thread sau distribuit
- Modelarea complexității unui algoritm multi-thread sau distribuit
- Recunoașterea soluțiilor clasice de tip multi-thread sau distribuit în probleme reale



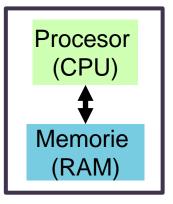
Calcul Paralel vs Distribuit vs Secvențial



Calcul Secvențial

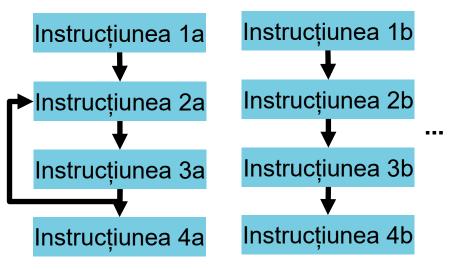


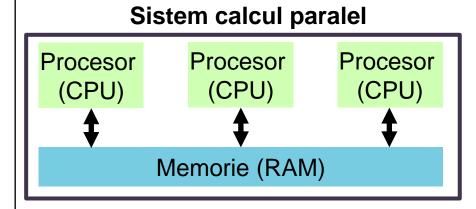
Sistem calcul secvențial





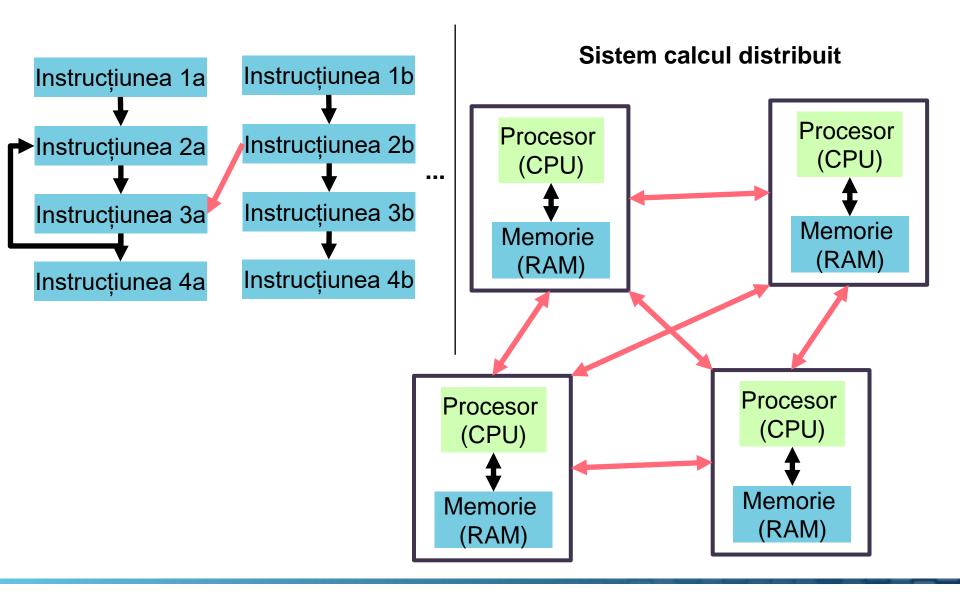
Calcul Paralel







Calcul Distribuit





Resurse fizice

Procesor – multi-core – 48 core-uri(64 la AMD)



Cluster



Grid/Cloud





Supercomputers

Rank	System	Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Supercomputer Fugaku - Supercomputer Fugaku, A64FX 48C 2.2GHz, Tofu interconnect D, Fujitsu RIKEN Center for Computational Science Japan	7,299,072	415,530.0	513,854.7	28,335
2	Summit - IBM Power System AC922, IBM POWER9 22C 3.07GHz, NVIDIA Volta GV100, Dual-rail Mellanox EDR Infiniband, IBM DOE/SC/Oak Ridge National Laboratory United States	2,414,592	148,600.0	200,794.9	10,096
3	Sierra - IBM Power System AC922, IBM POWER9 22C 3.1GHz, NVIDIA Volta GV100, Dual-rail Mellanox EDR Infiniband, IBM / NVIDIA / Mellanox DOE/NNSA/LLNL United States	1,572,480	94,640.0	125,712.0	7,438
4	Sunway TaihuLight - Sunway MPP, Sunway SW26010 260C 1.45GHz, Sunway, NRCPC National Supercomputing Center in Wuxi China	10,649,600	93,014.6	125,435.9	15,371
5	Tianhe-2A - TH-IVB-FEP Cluster, Intel Xeon E5-2692v2 12C 2.2GHz, TH Express-2, Matrix-2000, NUDT National Super Computer Center in Guangzhou China	4,981,760	61,444.5	100,678.7	18,482



Fugaku



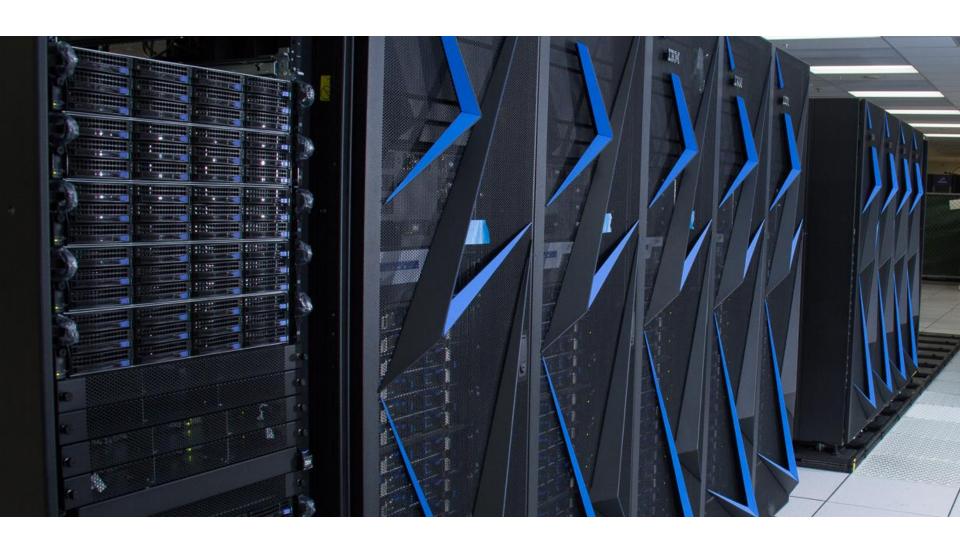


Summit





Sierra



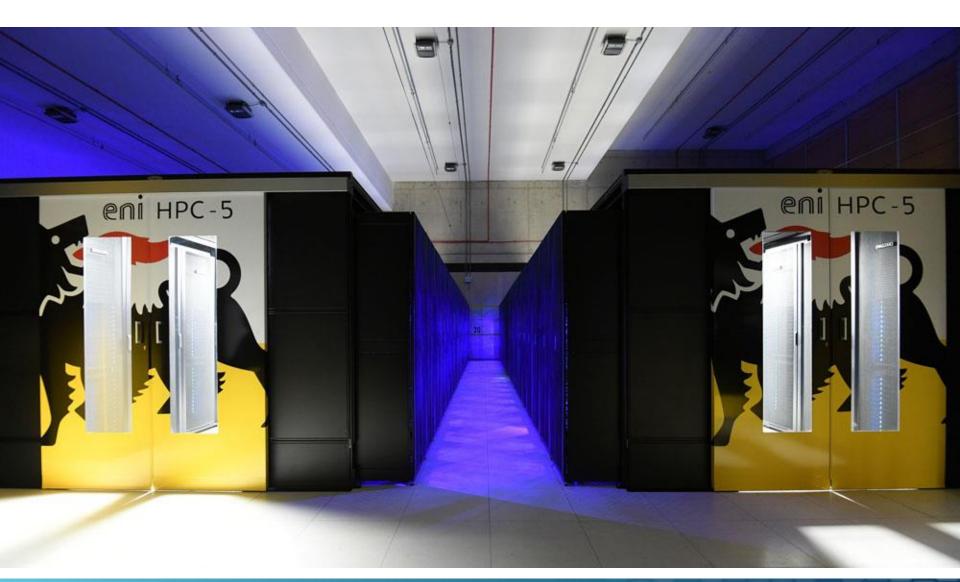


Sunway TaihuLight





HPC - 5







Computing power

24-hour average: 27.702 PetaFLOPS.

Active: 87,060 volunteers, 918,047 computers.



De ce calcul paralel și distribuit?

- Timp de execuție mai scurt
- Permite abordarea problemelor de dimensiuni mari
- Accesul resurselor aflate la distanță
- Reducerea costurilor
- Toleranța la defecte
- Ascunderea timpilor de așteptare
- Redundanţa
- Scalabilitatea
- Scăderea timpului de răspuns
- Securitate



Limitele programării secvențiale?

Cramming More Components onto Integrated Circuits

GORDON E. MOORE, LIFE FELLOW, IEEE

With unit cost falling as the number of components per circuit rises, by 1975 economics may dictate squeezing as many as 65 000 components on a single silicon chip.

The future of integrated electronics is the future of electronics itself. The advantages of integration will bring about a proliferation of electronics, pushing this science into many new areas.

Integrated circuits will lead to such wonders as home computers—or at least terminals connected to a central computer—automatic controls for automobiles, and personal portable communications equipment. The electronic wristwatch needs only a display to be feasible today.

But the biggest potential lies in the production of large systems. In telephone communications, integrated circuits Each approach each borrowed to believe the way ovarious approach

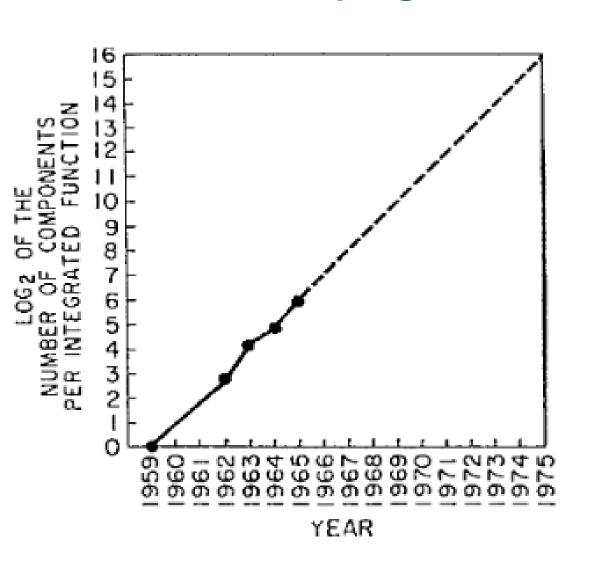
The advocates already using the resistors by apply conductor substraupon films are de

attachment of active semiconductor devices to the passive film arrays.

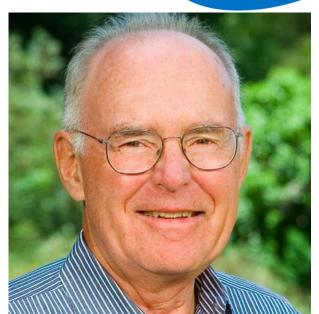
Both approaches have worked well and are being used in equipment today.



Limitele programării secvențiale?







cofounded Intel



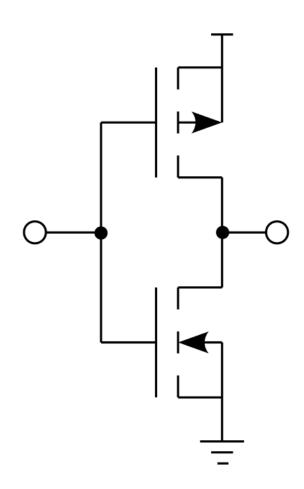
Limitele programării secvențiale

- Viteza de transmisie
 - Maxim c viteza luminii
- Miniaturizare
 - Tranzistor de mărimea unui atom
- Economic
 - Costuri enorme pentru cercetare şi proiectarea unui nou timp de procesor



Limitele programării secvențiale

CMOS





- Apar tot mai multe tehnologii distribuite
 - Blockchain; Peer-to-Peer
- Chiar și un procesor de ceas are mai multe core-uri
 - LG Watch Sport MSM8909w Processor
 - Quad-Core
- Suport în noile IDE-uri
 - Eclipse; Visual Studio
- Număr mare de aplicații distribuite
 - Dropbox; Spark; <u>Boinc</u>



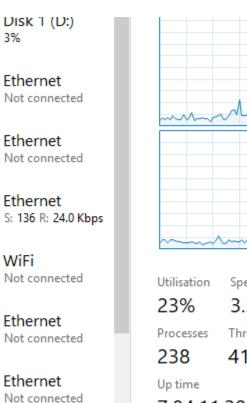


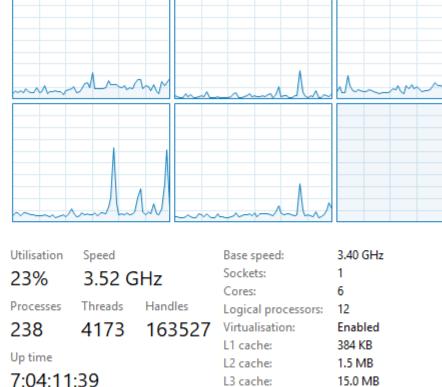
EuroHPC Leading the way in European Supercomputing

The EuroHPC Joint Undertaking is a joint initiative between the EU, European countries and private partners to develop a World Class Supercomputing Ecosystem in Europe.



Aproape toate aplicațiile au mai multe thread-uri





etails | 🔕 Open Resource Monitor

DI..... LAL DANI



Processes Performance	App his	tory Start-up	Users Details	Services			
Name	PID	Status	Username	CPU	Memory (p	Threads	Description
explorer.exe	4948	Running	cristian.chi	. 00	154,348 K	260	Windows Explorer
I System	4	Running	SYSTEM	00	20 K	240	NT Kernel & System
Uropbox.exe	7900	Running	cristian.chi	. 00	161,068 K	148	Dropbox
NVIDIA Web Helper.	21388	Running	cristian.chi	. 00	29,024 K	96	NVIDIA Web Helper Service
Origin.exe	13304	Running	cristian.chi	. 00	100,008 K	96	Origin
nvcontainer.exe	5752	Running	SYSTEM	00	30,000 K	86	NVIDIA Container
🥑 firefox.exe	28740	Running	cristian.chi	. 00	424,460 K	85	Firefox
nvcontainer.exe	8964	Running	NETWORK	. 00	11,180 K	83	NVIDIA Container
🥑 firefox.exe	10140	Running	cristian.chi	. 00	770,132 K	73	Firefox
S Skype.exe	9156	Running	cristian.chi	. 00	225,100 K	66	Skype
POWERPNT.EXE	27828	Running	cristian.chi	. 10	183,364 K	64	Microsoft PowerPoint
🥑 firefox.exe	28840	Running	cristian.chi	. 00	890,532 K	64	Firefox
CorsairLink4.Service.	15076	Running	SYSTEM	01	40,780 K	58	Corsair LINK 4 Service
🥑 firefox.exe	23840	Running	cristian.chi	. 00	506,796 K	57	Firefox
🥑 firefox.exe	22076	Running	cristian.chi	. 00	602,072 K	56	Firefox
BitTorrent.exe	1168	Running	cristian.chi	. 00	67,916 K	54	BitTorrent
■ MsMpEng.exe	29124	Running	SYSTEM	00	114,688 K	50	Antimalware Service Executable
■ SearchUl.exe	10212	Suspended	cristian.chi	. 00	102,652 K	49	Search and Cortana application
FortiTray.exe	20760	Running	cristian.chi	. 00	5,240 K	48	FortiClient System Tray Controller
OVRServer_x64.exe	7968	Running	cristian.chi	. 00	53,356 K	46	OVRServer_x64.exe 676007-public SC:67765493906
Steam.exe	13576	Running	cristian.chi	. 00	309,492 K	44	Steam Client Bootstrapper
googledrivesync.exe	10116	Running	cristian.chi	. 00	178,840 K	42	googledrivesync.exe



Taxonomia Flynn

Some Computer Organizations and Their Effectiveness

MICHAEL J. FLYNN, MEMBER, IEEE

Abstract—A hierarchical model of computer organizations is developed, based on a tree model using request/service type resources as nodes. Two aspects of the model are distinguished: logical and physical.

General parallel- or multiple-stream organizations are examined as to type and effectiveness—especially regarding intrinsic logical difficulties.

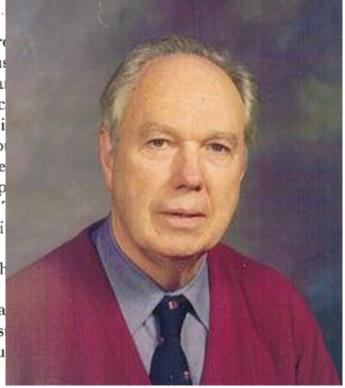
The overlapped simplex processor (SISD) is limited by data dependencies. Branching has a particularly degenerative effect.

The parallel processors [single-instruction stream-multiple-data stream (SIMD)] are analyzed. In particular, a nesting type explanation is offered for Minsky's conjecture—the performance of a parallel processor increases as $\log M$ instead of M (the number of data stream processors).

Multiprocessors (MIMD) are subjected to a saturation syndrome based on general communications lockout. Simplified queuing models indicate that saturation develops when the fraction of task time spent locked out (L/E) approaches 1/n, where n is the number of processors. Resources sharing in multiprocessors can be used to avoid

more "macro particular us must be shar more signific

- 1) There is limiting resonant est will either tions will appreciations. Computer with cerned with potential, whisiderations.
- 2) We ma sets. It is ass set of instru





Taxonomia Flynn

SISD

- Single Instruction Stream, Single Data Stream
- Calculatorul "clasic" one-core

SIMD

- Single Instruction Stream, Multiple Data Streams
- Suportul SSE; procesoare GPU

MISD

- Multiple Instruction Streams, Multiple Data Streams
- Sisteme specializate

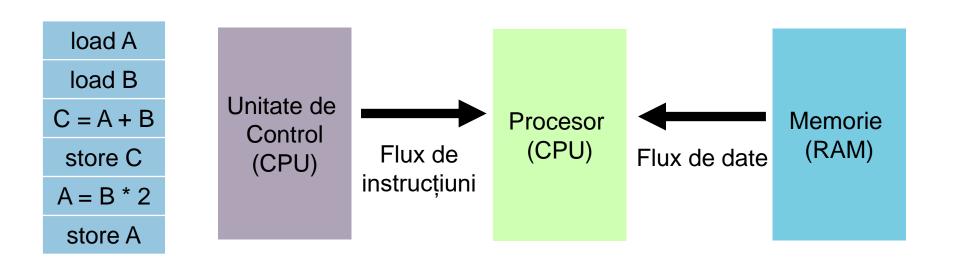
MIMD

- Multiple Instruction Streams, Multiple Data Streams
- Procesoare actuale (ce aveţi acasă şi în buzunar)



SISD

Model clasic Arhitectura von Neumann





SISD

Model clasic Arhitectura von Neumann

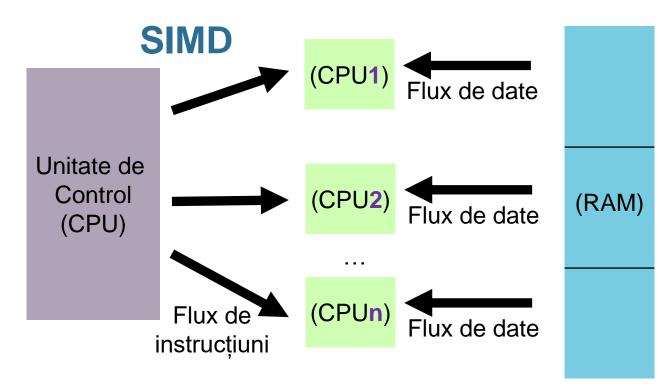
First Draft of a Report on the EDVAC

by

John von Neumann







load A[1]	load A[2]		load A[n]
load B[1]	load B[2]		load B[n]
C[1] = A[1] + B[1]	C[2] = A[2] + B[2]		$C[\mathbf{n}] = A[\mathbf{n}] + B[\mathbf{n}]$
store C[1]	store C[2]	•	store C[n]
A[1] = B[1] * 2	A[2] = B[2] * 2		$A[\mathbf{n}] = B[\mathbf{n}] * 2$
store A[1]	store A[2]		store A[n]



SIMD – Memorie Partajată

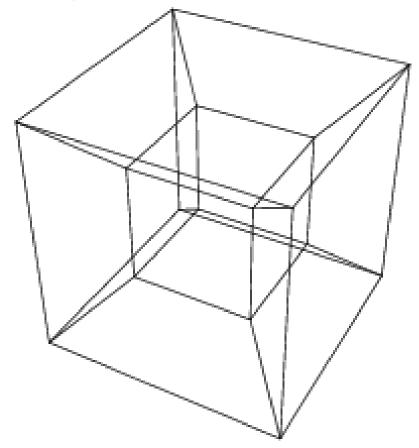
- Shared memory
 - Parallel Random Access Memory
 - PRAM

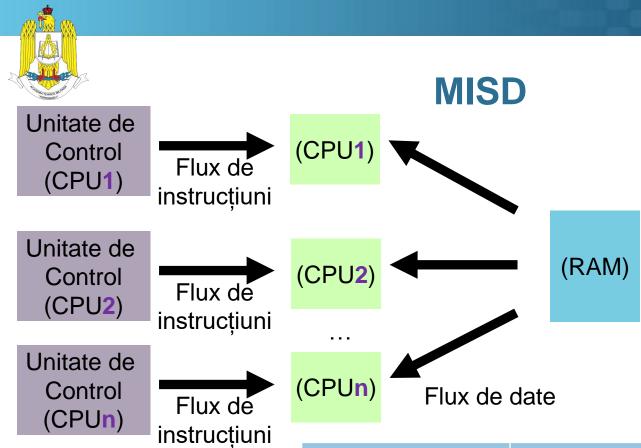
- EREW Exclusive Read Exclusive Write
- CREW Concurrent Read Exclusive Write -- cel mai des întâlnit
- ERCW Exclusive Read Concurrent Write
- CRCW Concurrent Read Concurrent Write
- O variabilă poate fi citiă într-un pas în model CR dar în log(N) în model ER.



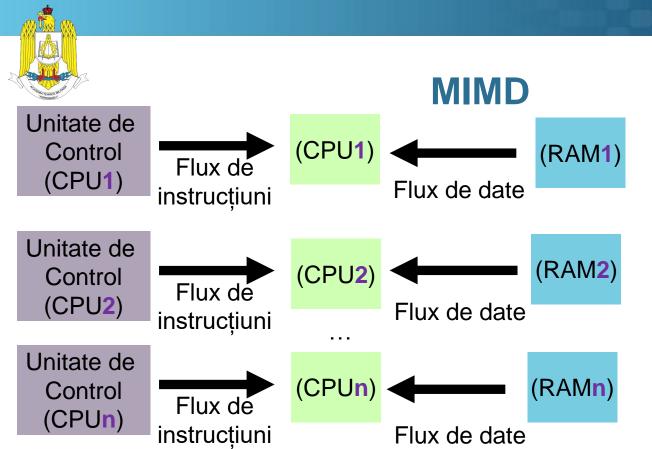
Rețele de configurare

- Topologii
 - Tablou
 - Arbore
 - Cub
 - Hipercub
- Depinde de
 - Aplicație
 - Performanțe dorite
 - Număr procesoare disponibile
- Exemple: IBM 9000, Cray C90, Fujitsu VP





load A[1]	load A[1]		load A[1]
load B[1]	load B[1]	•	load B[1]
C[1] = A[1] + B[1]	C[1] = A[1] / B[1]	•	C[1] = A[1] * B[1]
store C[1]	store C[1]	•	store C[1]
A[1] = B[1] * 2	load D[1]		C[1] = A[1] + B[1]
store A[1]	A[1] = D[1]		store C[1]

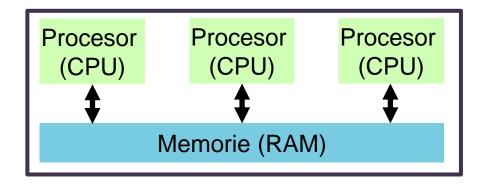


load A	load Z		A = B + C
D = A	store H	•	store A
load C	H=H*100	•	G[1] = G[2] + G[3]
load E	H++	•	store G[1]
load F	store H		G[4] = G[5] + G[6]



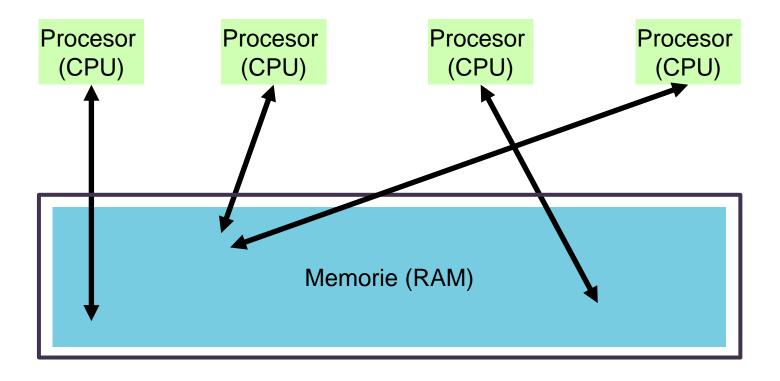
Memorie Partajată

Uniform Memory Access (UMA)





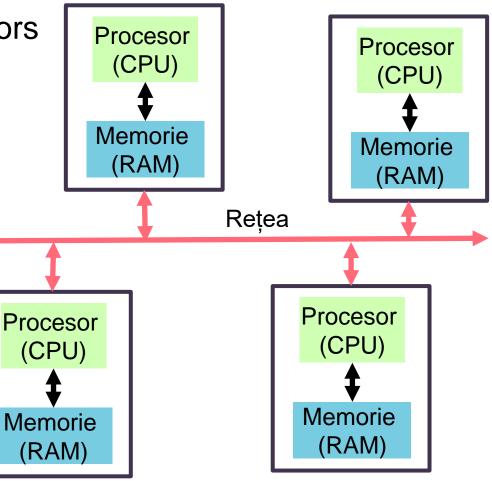
Memorie Partajată - Acces





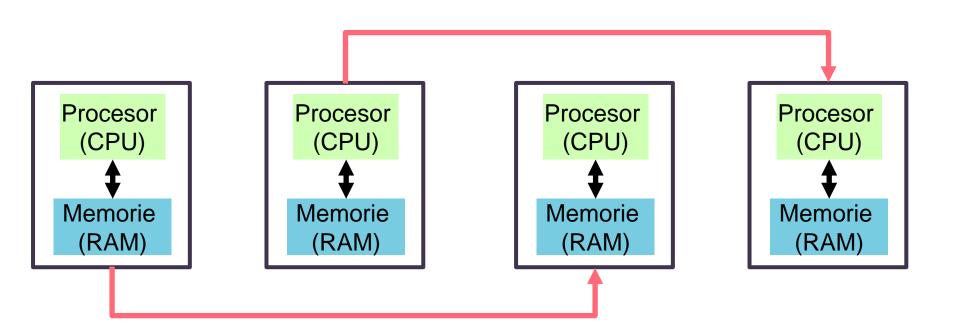
Memorie Distribuită

- Non-Uniform Memory Access (NUMA)
- Massively Parallel Processors
- Network of workstations



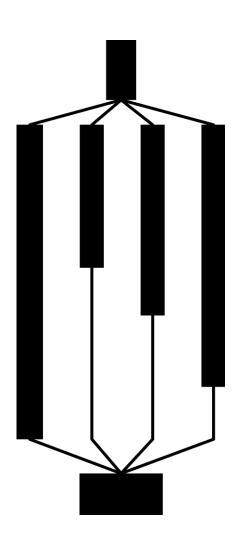


Memorie Distribuită - Acces

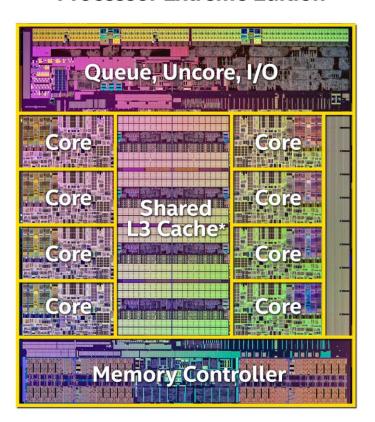




Threads vs cores



New 8-Core Intel® Core™ i7 Processor Extreme Edition



Intel® Core™ i7-5960X Processor Extreme Edition Transistor count: 2.6 Billion Die size: 17.6mm x 20.2mm

* 20MB of cache is shared across all 8 cores



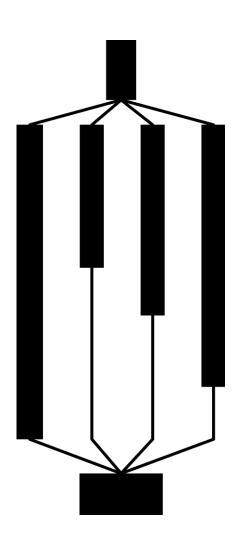


Hyperthreading – the confusion

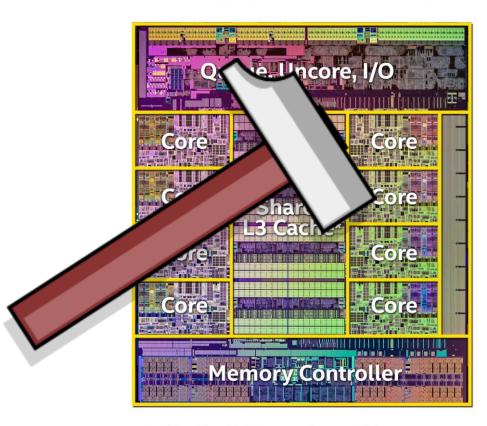




Hyperthreading – the confusion



New 8-Core Intel® Core™ i7 Processor Extreme Edition



Intel® Core™ i7-5960X Processor Extreme Edition Transistor count: 2.6 Billion Die size: 17.6mm x 20.2mm



^{* 20}MB of cache is shared across all 8 cores









Task 1

























Thread 2

Thread 3

Thread 4

Thread 1





Thread 1

Thread 2

Thread 4

Thread 3





Thread 1

Thread 2

Thread 4 Thread 3





Thread 2

Thread 3

Thread 4

Thread 1





Thread 1

Thread 4

Thread 2 Thread 3





Folosesc același principiu, ba chiar și aceleași sisteme!







Deci care sunt diferențele?

Un proces are mai multe thread-uri

Thread-urile unui proces împart

memoria







Poţi avea multi-tasking pe mai multe core-uri?

Poți avea multi-thread-ing pe un singur core?







Poţi avea multi-tasking pe mai multe core-uri? **DA**

Poţi avea multi-thread-ing pe un singur core? **DA**





