



Lecture 2

summary of Java SE & Network Programming – section 2

presentation

DAD – Distributed Applications Development

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Cristian Toma – Business Card



Cristian Toma

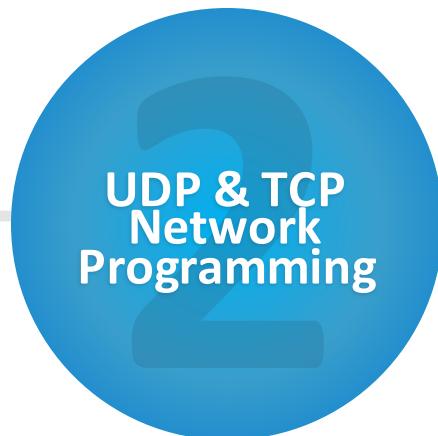
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Agenda for Lecture 2



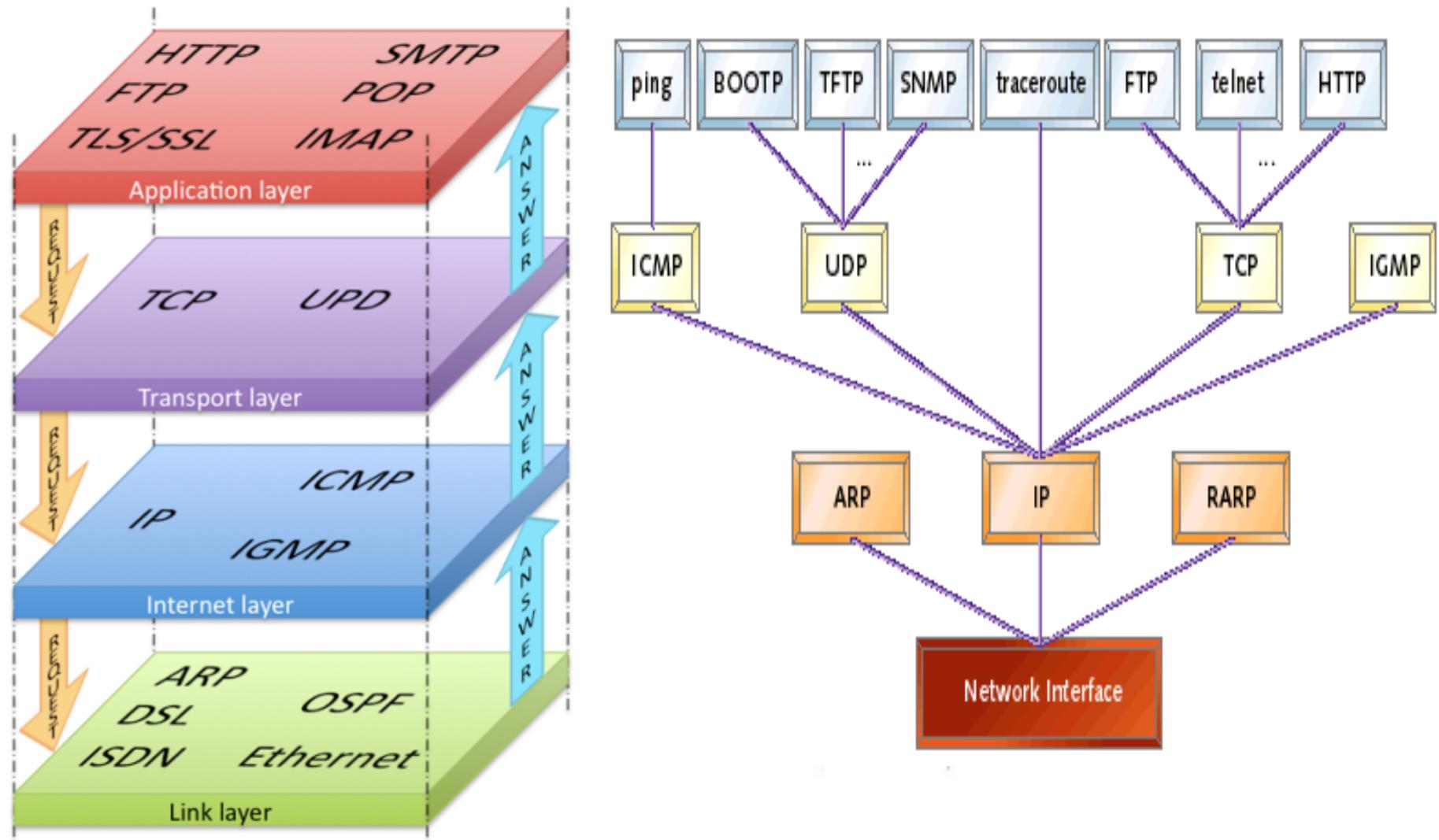


Networking IP, UDP and TCP programming, TCP/IP state machine, SNMP, SMTP

Networking Recapitulation

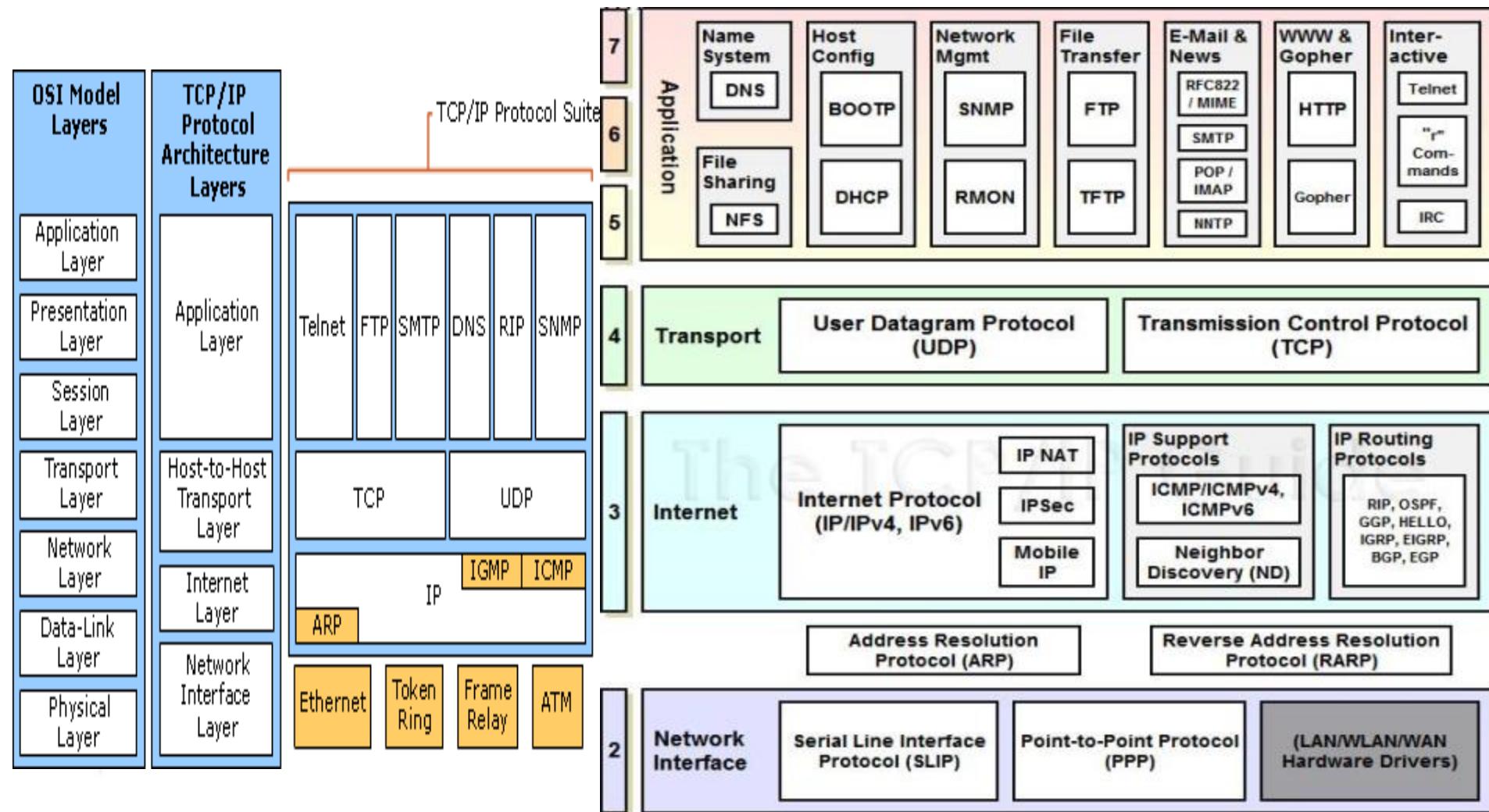
1. Networking TCP/IP Stack

HOW TCP/IP Works:



1. Networking TCP/IP Stack

TCP/IP Stack Model:



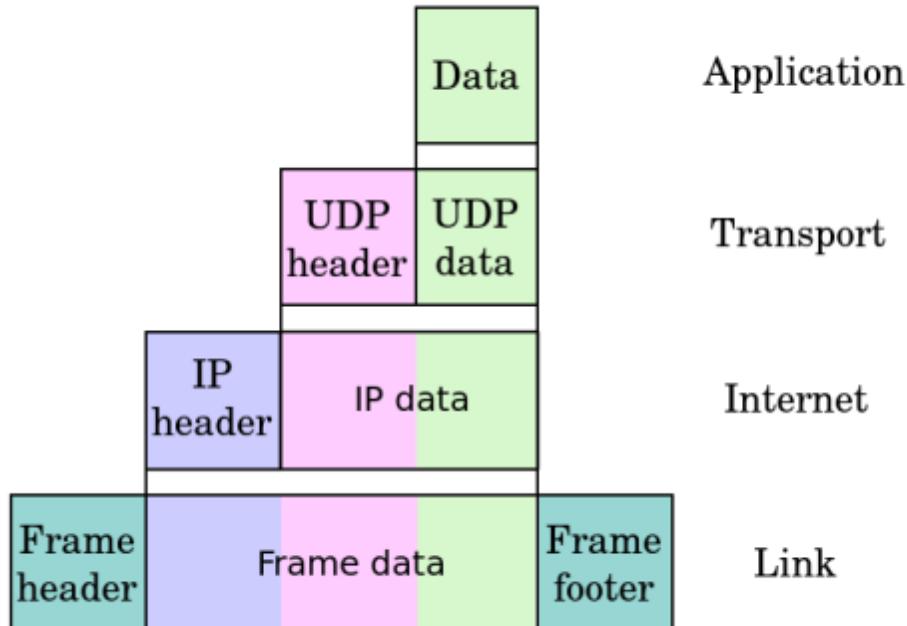
1. Networking TCP/IP Stack

ISO/OSI Model vs. TCP/IP:

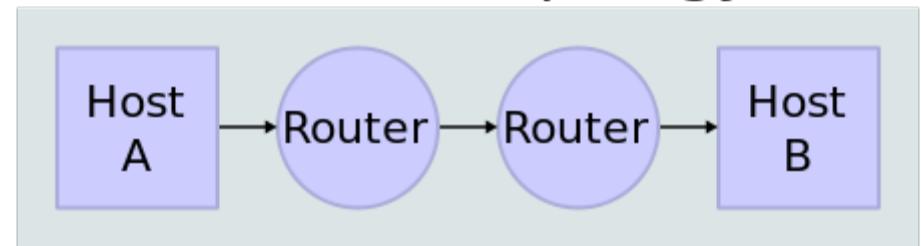
TCP/IP DoD Model			OSI Model
Application Layer (Services Layers 5,6,7) PDU: Data	HTTP: port 80 HTTPS/TLS/SSL: port 443 NNTP: port 119 FTP: port 21, 20 Telnet: port 23 SSH: port 22 POP3: port 110 IMAP4: port 143 SMTP: port 25	DNS: port 53 TFTP: port 69 DHCP/BootP: port 67,68 SNMP: port 162, 161 NTP: port 123 Syslog: port 514	Application Layer (7) Scribe. APIs, network services Serves the King/User
Transport Layer (Host to Host Layer 4) PDU: Segments	TCP: protocol 6	UDP: protocol 17	Presentation Layer (6) Translator. Reformats, encrypts/de-crypts, compress/de-compress
Internet Layer (Network Layer 3) PDU: Packets	IP	IP	Session Layer (5) Negotiator. Establishes, manages and ends sessions.
Network Acces Layer 1 & 2 PDU: Frame	Ethernet, PPP Frame Relay MAC addresses, ARP	Ethernet, PPP Frame Relay MAC addresses, ARP	Transport Layer (4) Middle Manager. Segment ID/Assembly Network Layer (3) Mail Room Guy. IP Addressing/Routing
Network Access Layer 1 & 2 PDU: Bits or Data Stream	Electrons, RF or Light	Electrons, RF or Light	Data-Link Layer (2) Envelope Stuffer. Organizes bits into frames Physical Layer (1) The Truck. Movement of bits.

1. Networking TCP/IP Stack

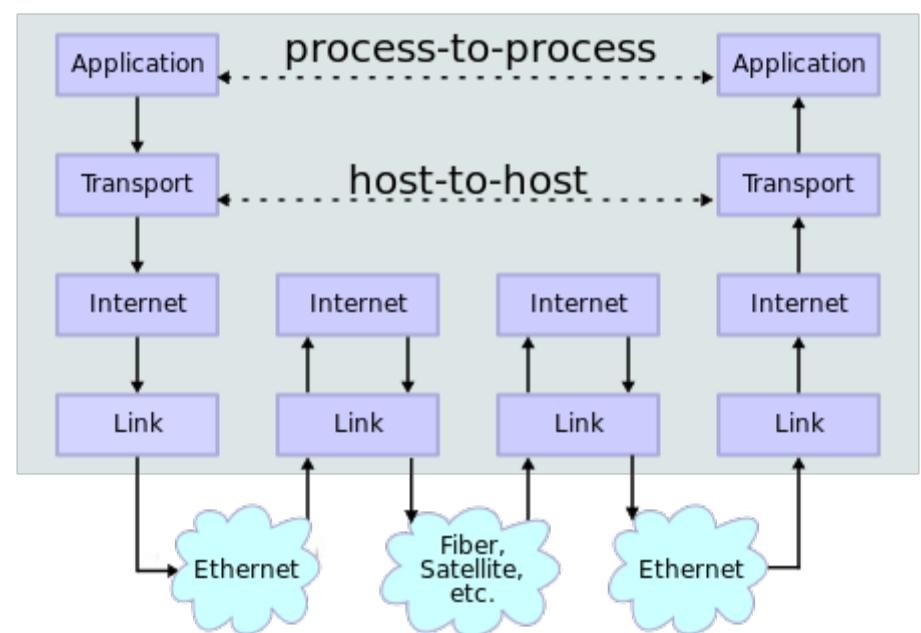
TCP/IP Message Encapsulation:



Network Topology

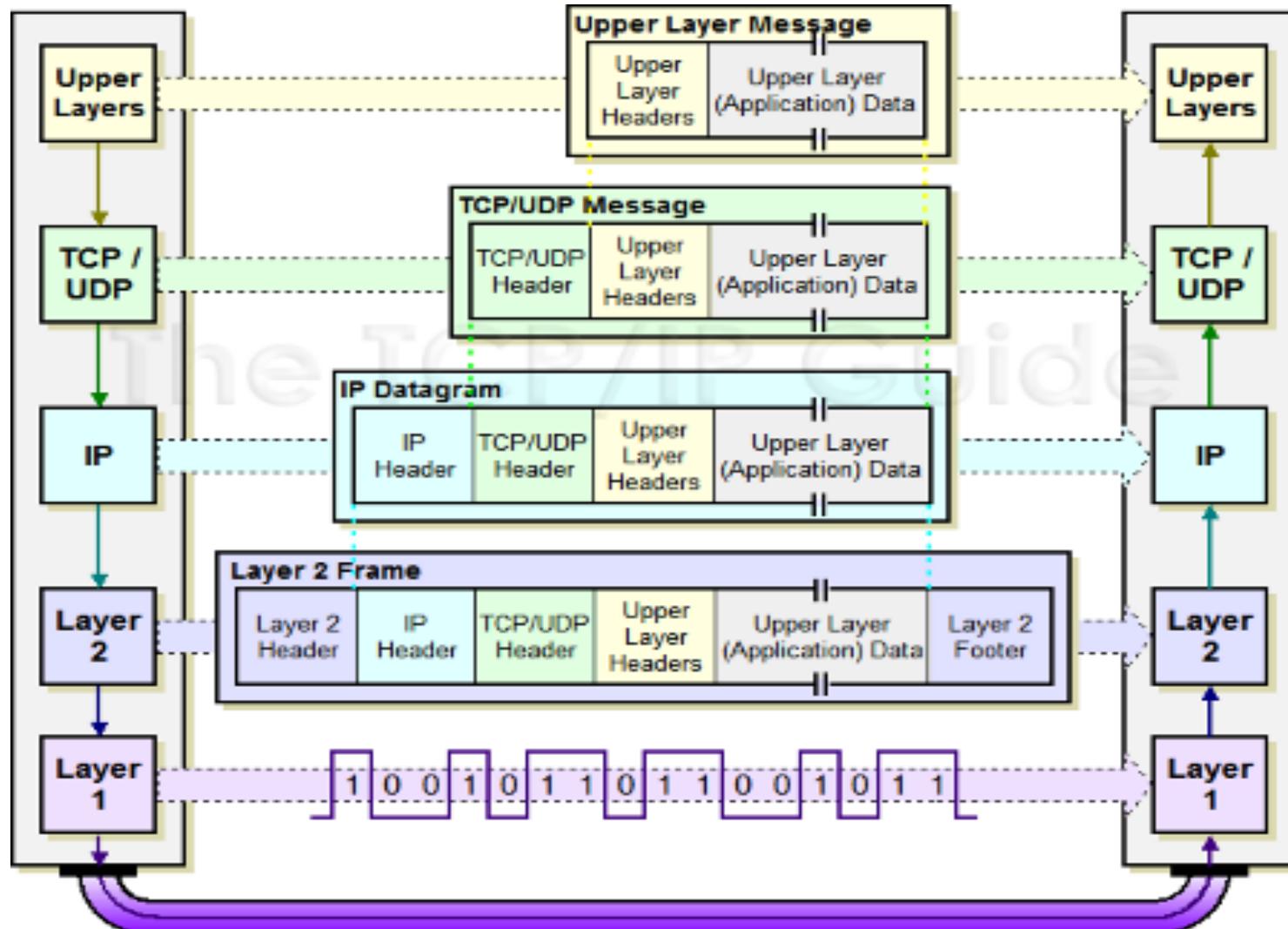


Data Flow



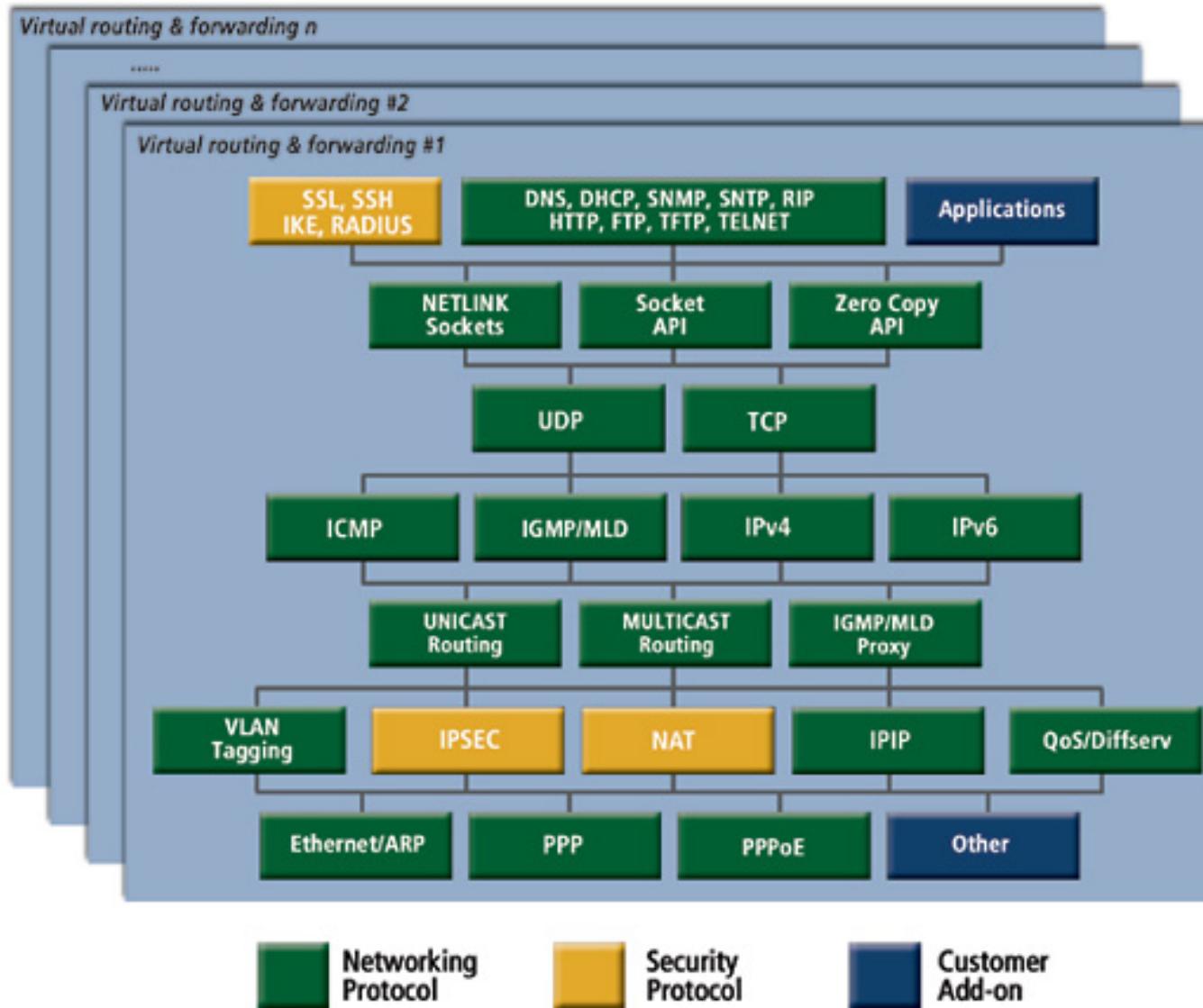
1. Networking TCP/IP Stack

TCP/IP Message Flow:



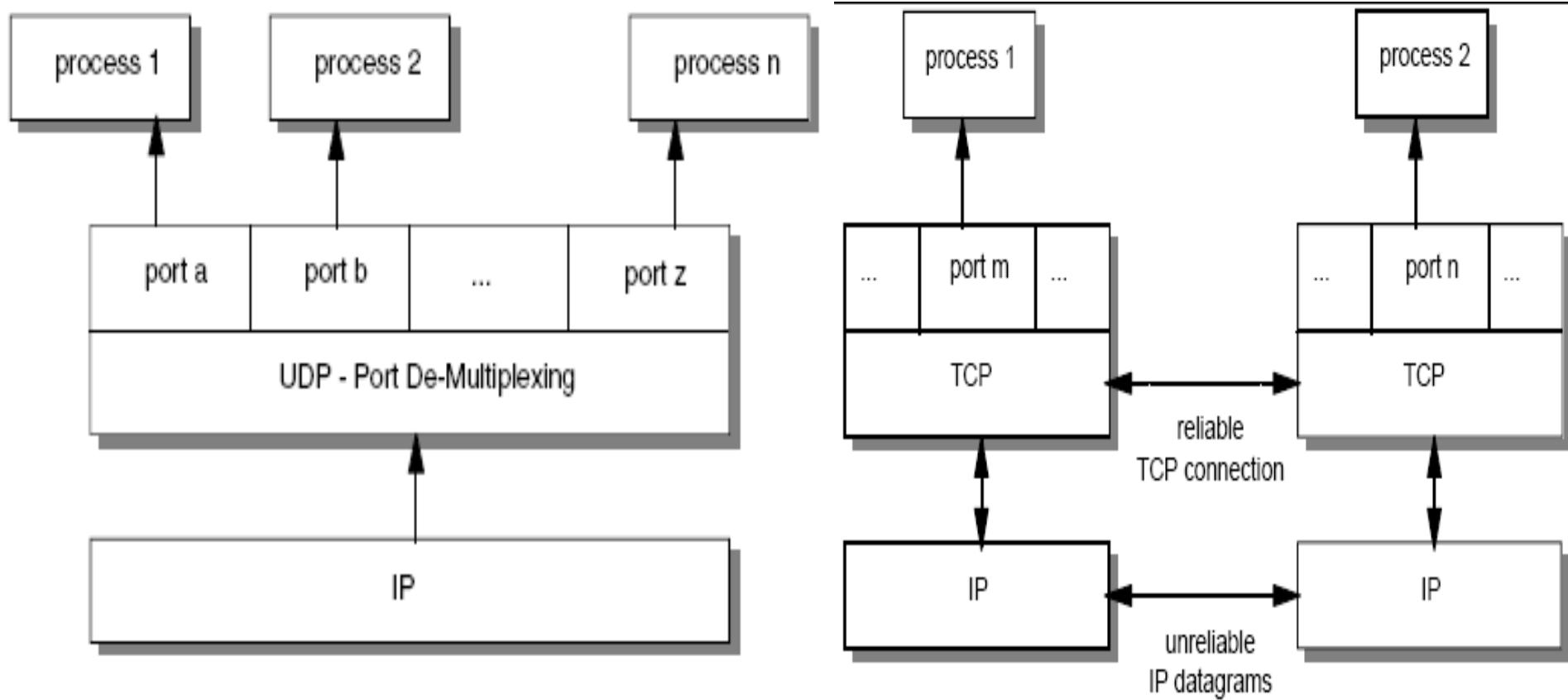
1. Networking TCP/IP Stack

TCP/IP Virtual LAN/Networks: http://www.ghs.com/products/comm_tcp-ip.html



1. Networking TCP/IP Stack

TCP/IP App/Port Multiplexing:



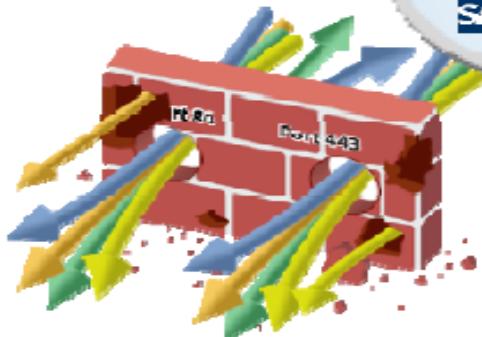
1. Networking TCP/IP Stack

What is running on port 80?



Applications Have Changed – Firewalls Have Not

- The gateway at the trust border is the right place to enforce policy control
 - Sees all traffic
 - Defines trust boundary

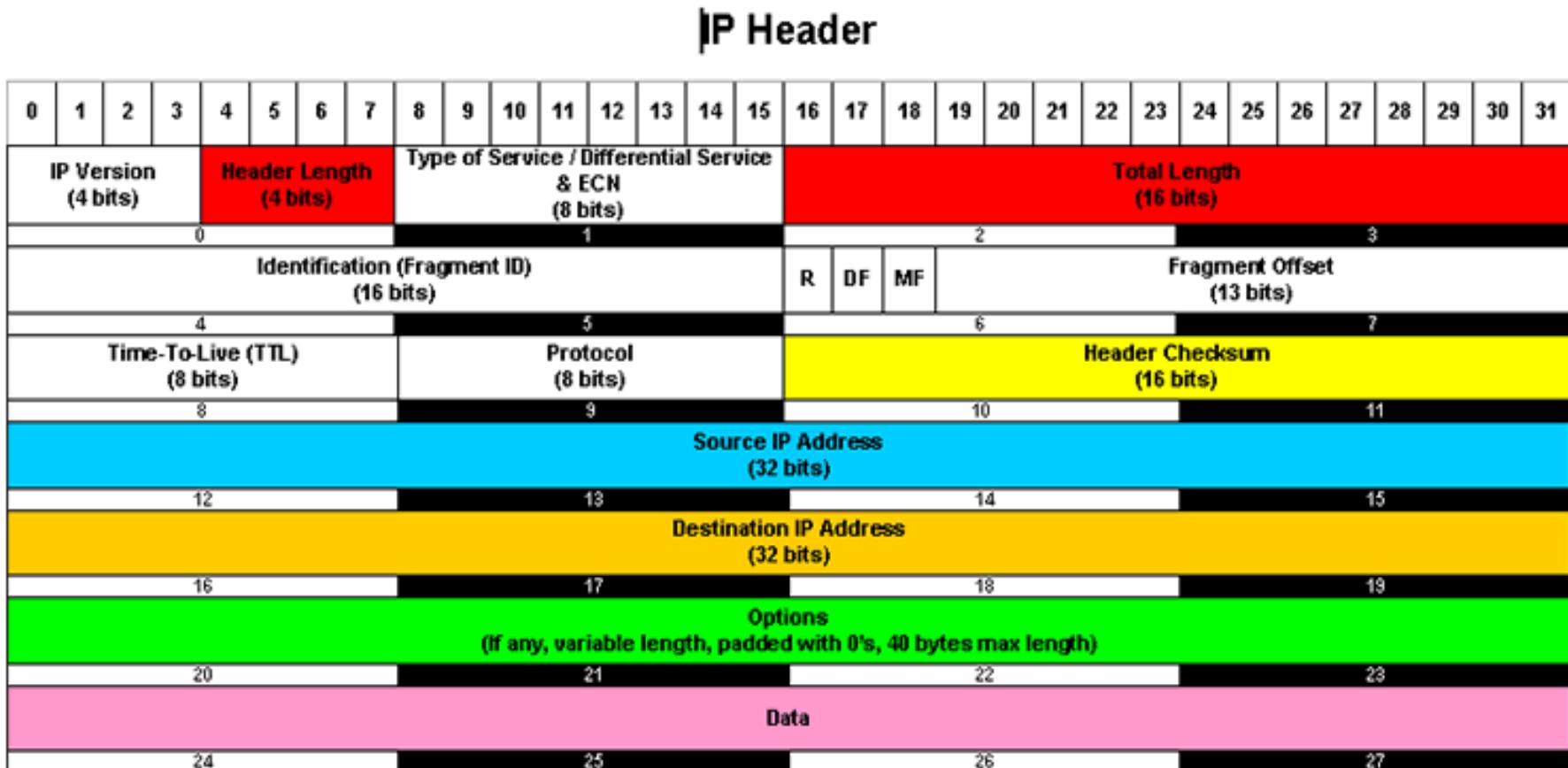


- BUT... Applications Have Changed
 - Ports ≠ Applications
 - IP Addresses ≠ Users
 - Packets ≠ Content

Need to Restore Visibility and Control in the Firewall

1. Networking TCP/IP Stack

IP Header - RFC 791 – Submasking + Routing + NAT:



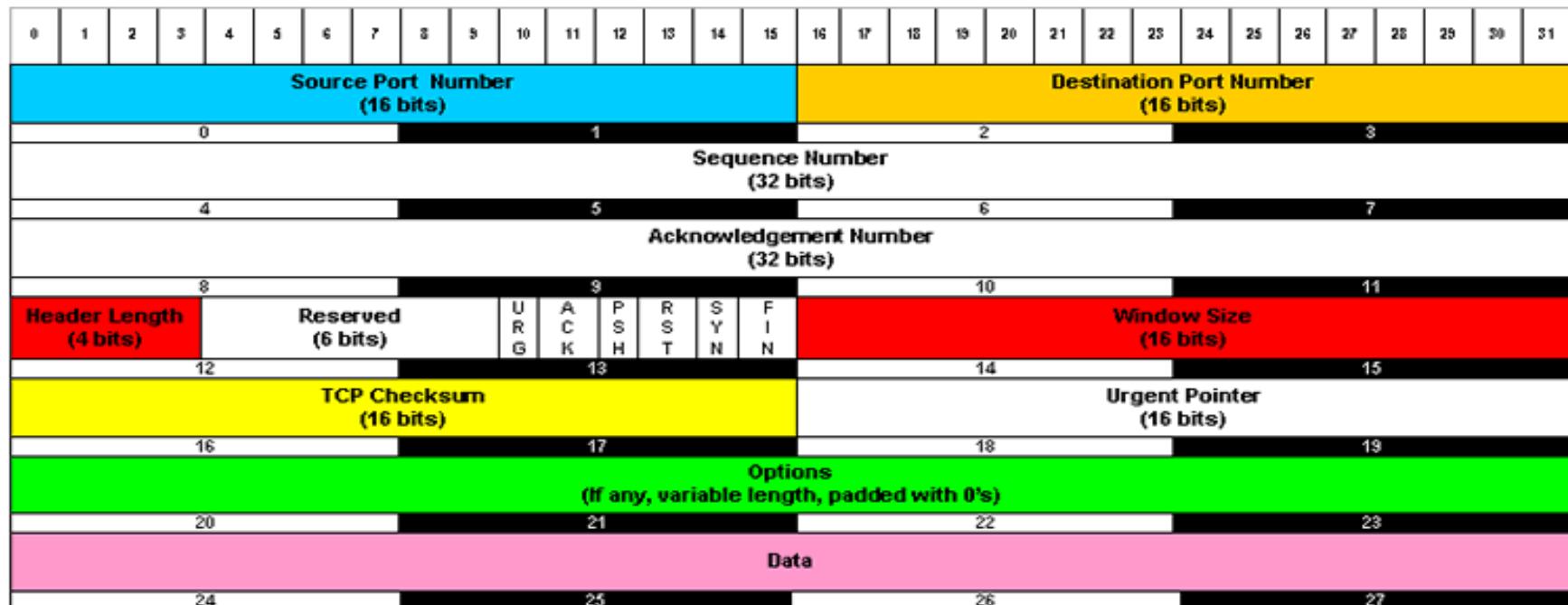
1. Networking TCP/IP Stack

UDP Header – RFC 768 – Connection-less vs. TCP Header – RFC 793 – Connection-oriented

UDP Header

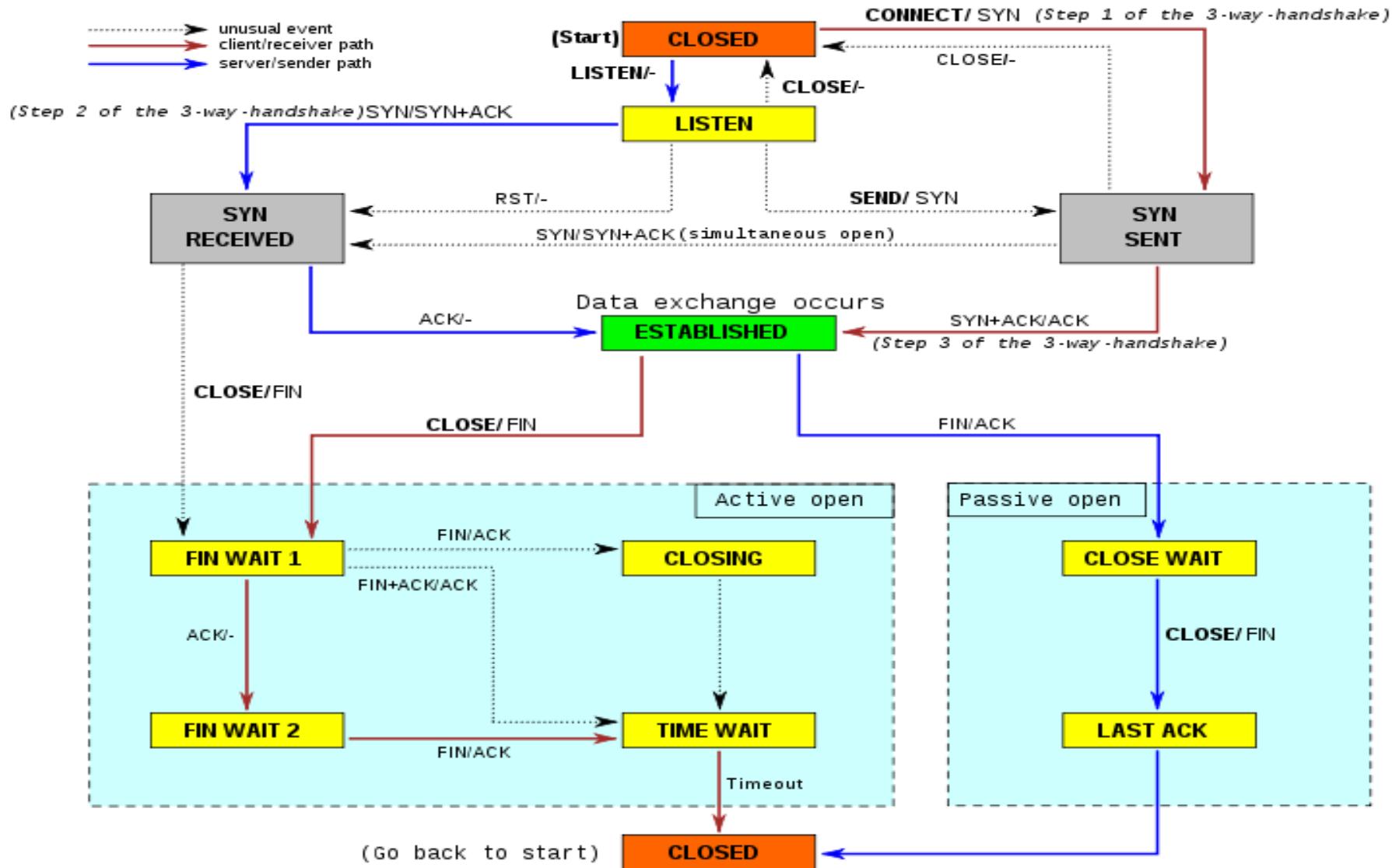


TCP Header



1. Networking TCP/IP Stack

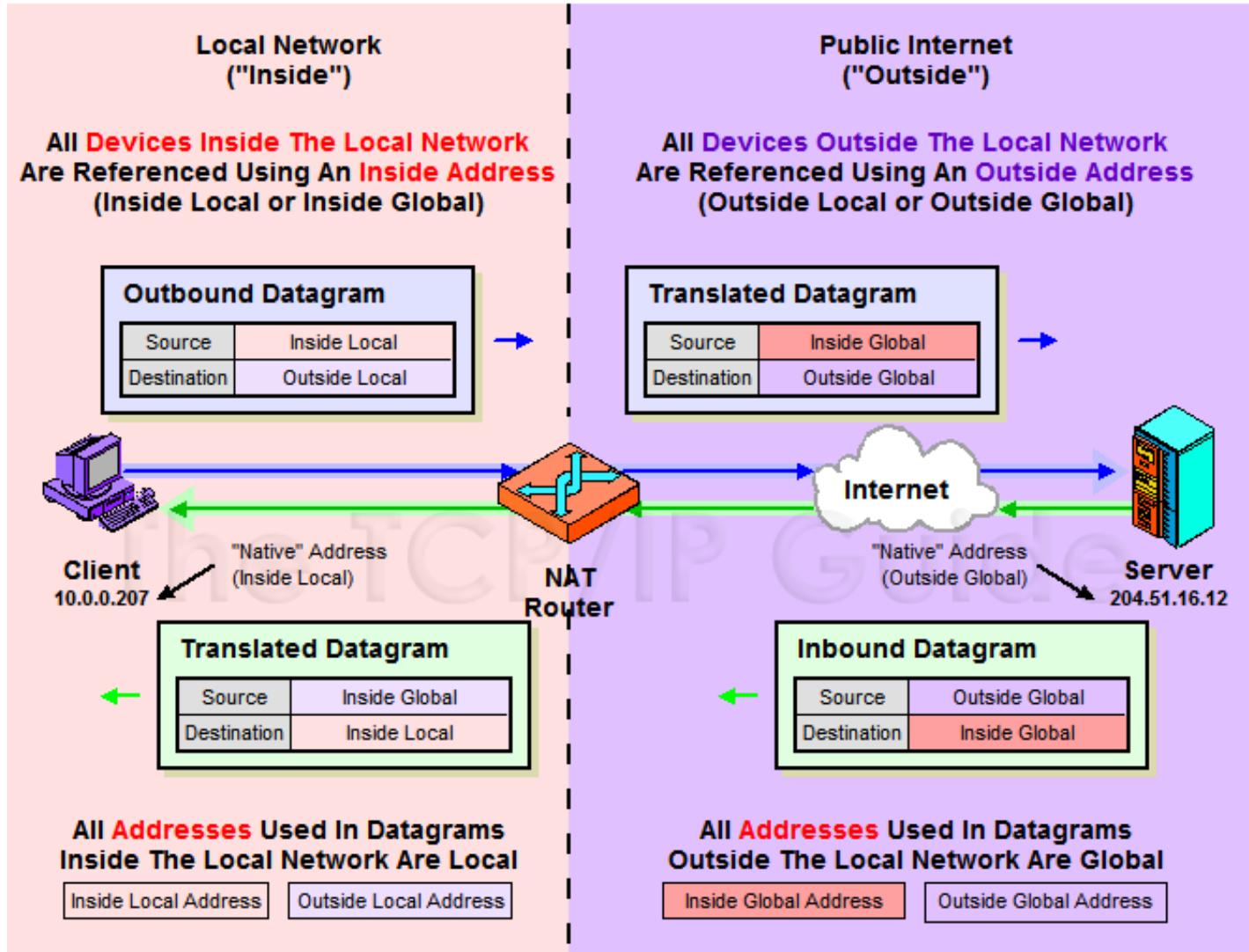
TCP State Machine – RFC 793:



1. Networking TCP/IP Stack

TCP/IP NAT & PAT:

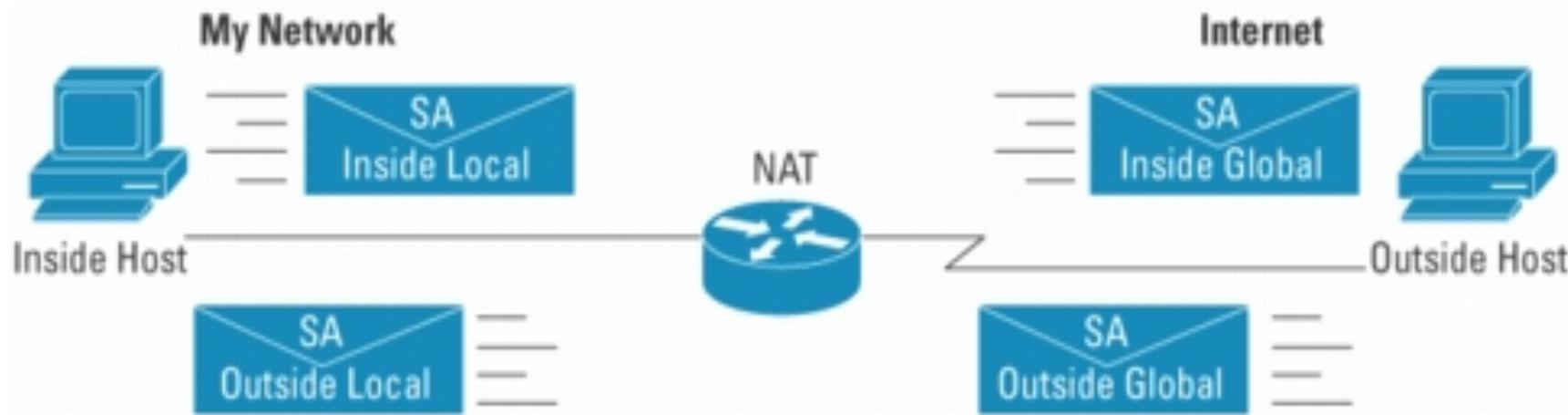
http://www.tcpipguide.com/free/t_IPNATAddressTerminology-3.htm



1. Networking TCP/IP Stack

TCP/IP NAT:

NAT Concepts



An IP address is either local or global

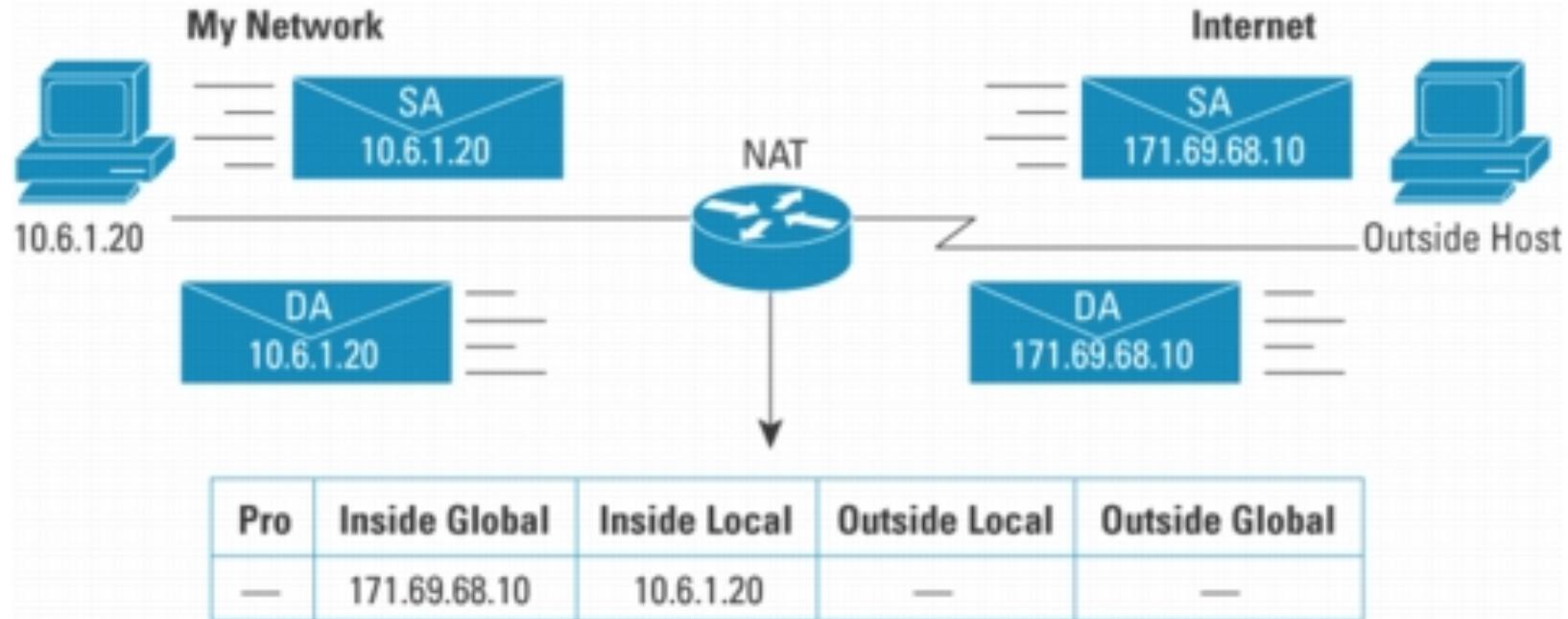
Local IP addresses are seen in the inside network

Global IP addresses are seen in the Outside network

1. Networking TCP/IP Stack

TCP/IP NAT:

NAT Terminology « Inside Addressing »



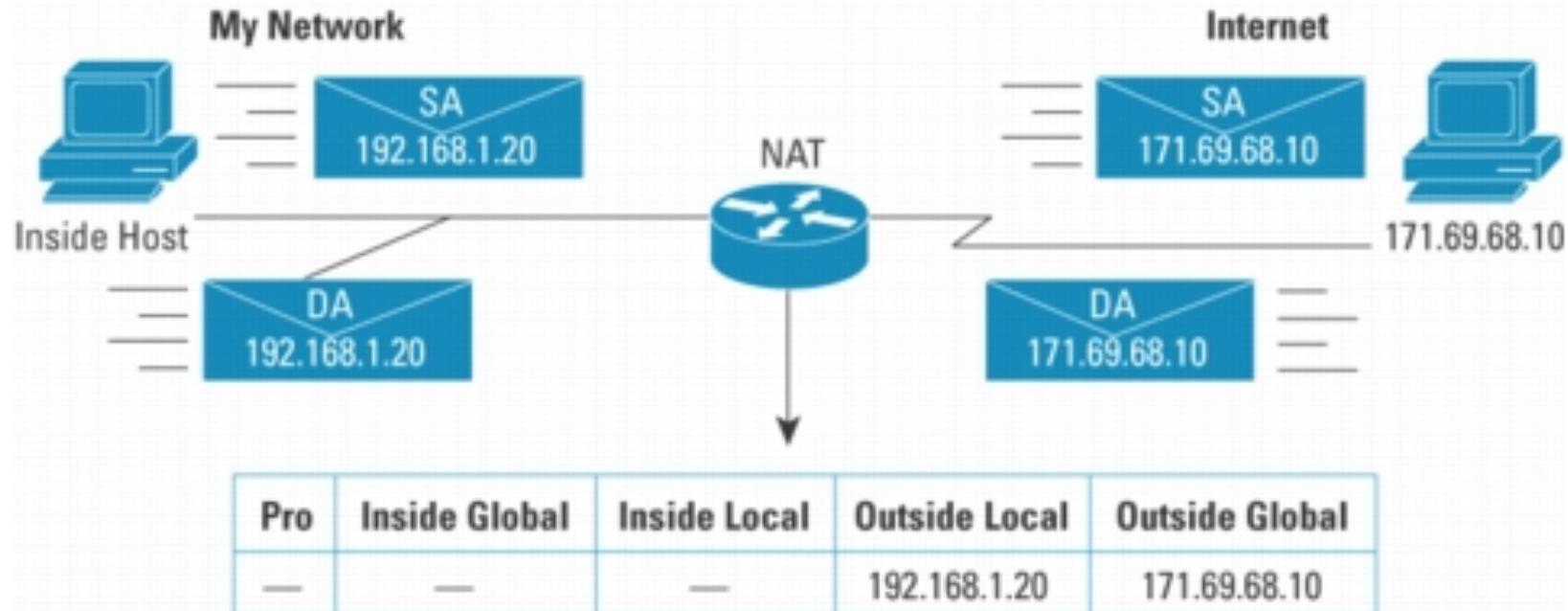
10.6.1.20 is inside local address

171.69.68.10 is inside global address

1. Networking TCP/IP Stack

TCP/IP NAT:

NAT Terminology "Outside Addressing"



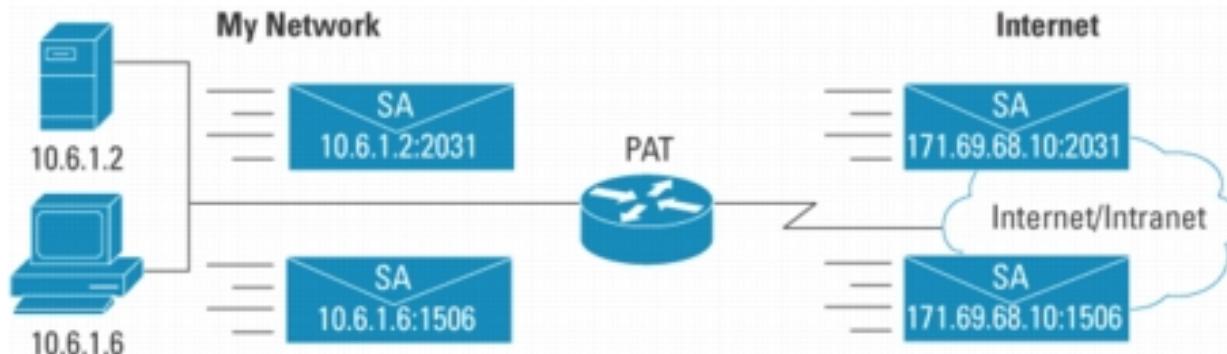
192.168.1.20 is inside local address

171.69.68.10 is inside global address

1. Networking TCP/IP Stack

TCP/IP PAT:

Basic Concepts of PAT



Port Address Translation (PAT) extends NAT from "1 to 1" to "many-to-1" by associating the source port with each flow

Figure 5

Unique Source Port per Translation Entry

Pro	Inside Global	Inside Local	Outside Local	Outside Global
tcp	171.69.68.5:1405	10.6.15.2:1405	204.71.200.69:80	204.71.200.69:80

PAT (Port Address Translation) includes ports in addition to IP addresses

Many-to-one translation

Maps multiple IP addresses to 1 or a few IP addresses

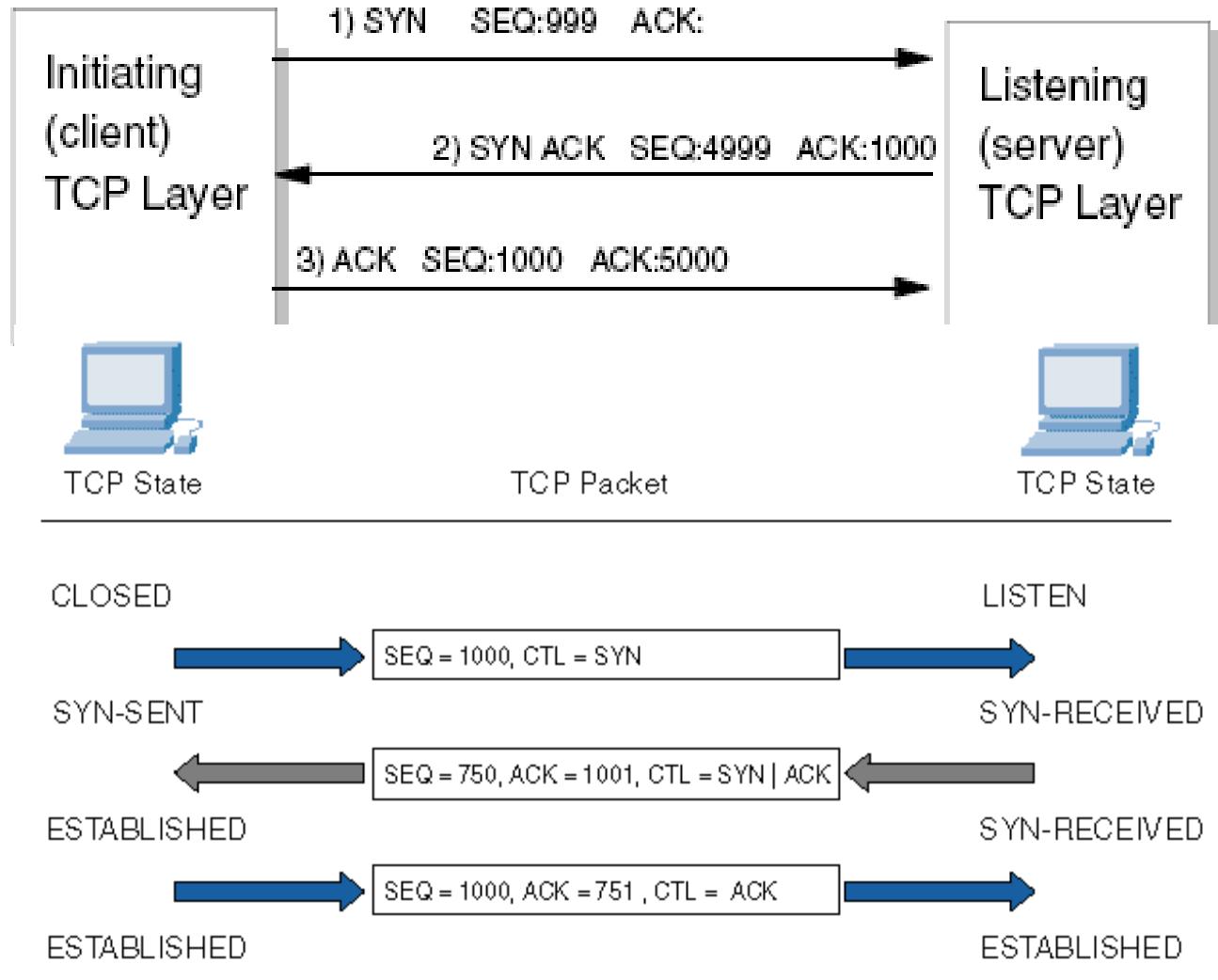
Unique source port number identifies each session

Conserves registered IP addresses

Also called NAPT in IETF documents

1.2 TCP/IP Networking Programming

TCP Handshake:



Section Conclusion

Fact: **DAD needs Networking**

In few **samples** it is simple to remember: UDP and TCP programming is useful for HTC ... didactical samples: SNMP over UDP for monitoring, SMTP/IMAP4/POP3 over TCP for e-mail notification, FTP over TCP for file and data transfer, HTTP over TCP for web & web services, and in general TCP for RPC/RMI, CORBA, JMS, SOAP, P2P-JXTA – for distributed computing and systems.

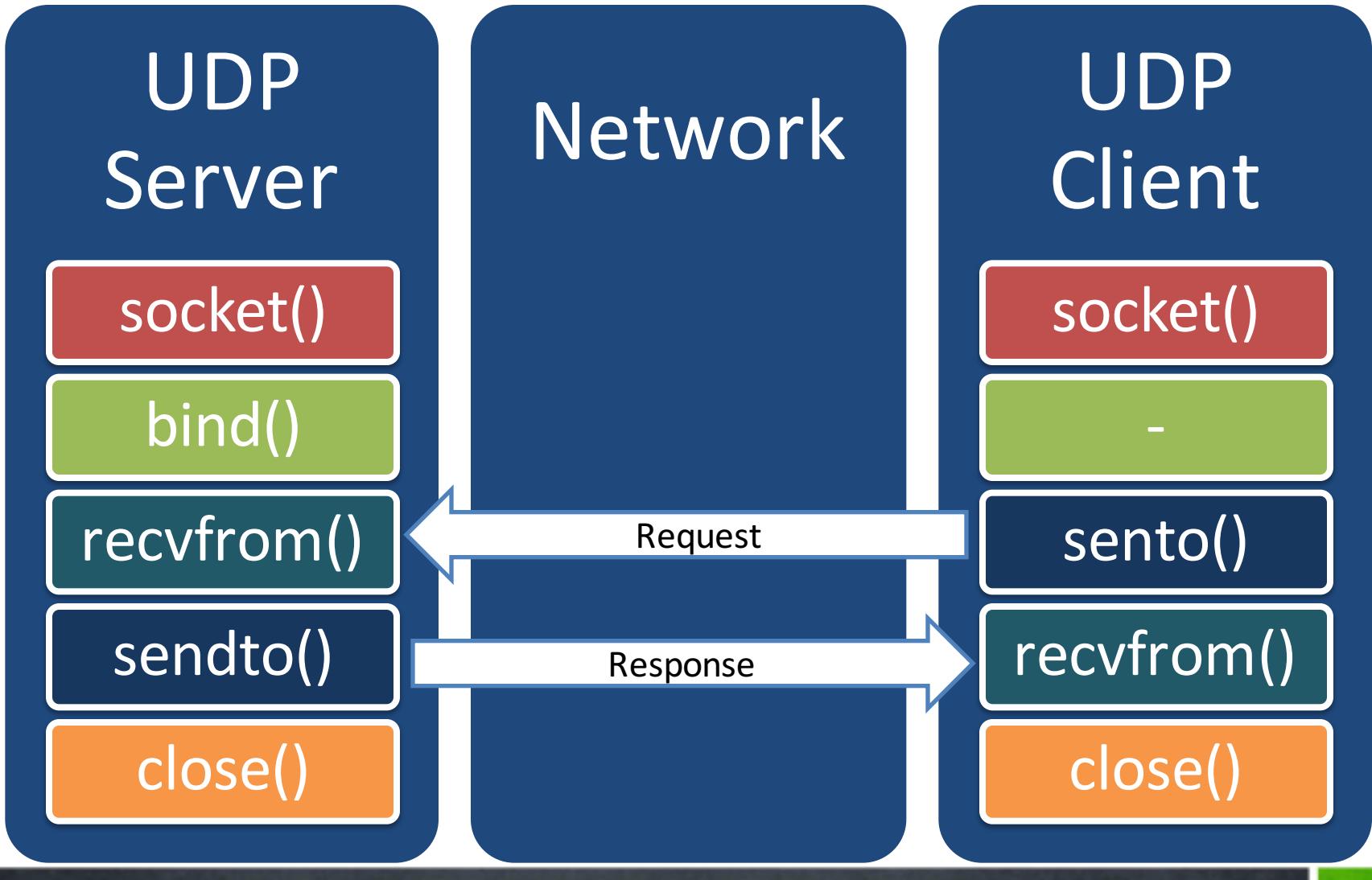




UDP Client-Server programming, SNMP client sample, TCP Client-Server programming, SMTP client sample

Network UDP & TCP Programming

2.1 TCP/IP Networking Programming – UDP Programming – Socket Primitives:



2.1 TCP/IP Networking Programming – UDP Programming – Socket Primitives:

```
package eu.ase.net.udp;
import java.io.*;
import java.net.*;
public class UDPServer {
    public static void main(String[] args) {
        // get a datagram socket
        DatagramSocket socket = null;
        byte[] bufResp = null;
        byte[] bufRecv = null;
        try {
            socket = new DatagramSocket(778); //it is correct because this constructor executes "bind"
            while(true) {
                bufRecv = new byte[256];
                // receive request
                DatagramPacket packet = new DatagramPacket(bufRecv, bufRecv.length);
                socket.receive(packet);

                // figure out response
                String respString = new String("OK");
                bufResp = respString.getBytes();

                // send the response to the client at "address" and "port"
                InetAddress address = packet.getAddress();
                int port = packet.getPort();
                packet = new DatagramPacket(bufResp, bufResp.length, address, port);
                socket.send(packet);
            }
        } catch(IOException ioe) {
```

2.1 TCP/IP Networking Programming – UDP Programming – Socket Primitives:

```
package eu.ase.net.udp;
import java.io.*;
import java.net.*;
public class UDPCClient {
    public static void main(String[] args) throws IOException {
        // get a datagram socket
        DatagramSocket socket = new DatagramSocket();

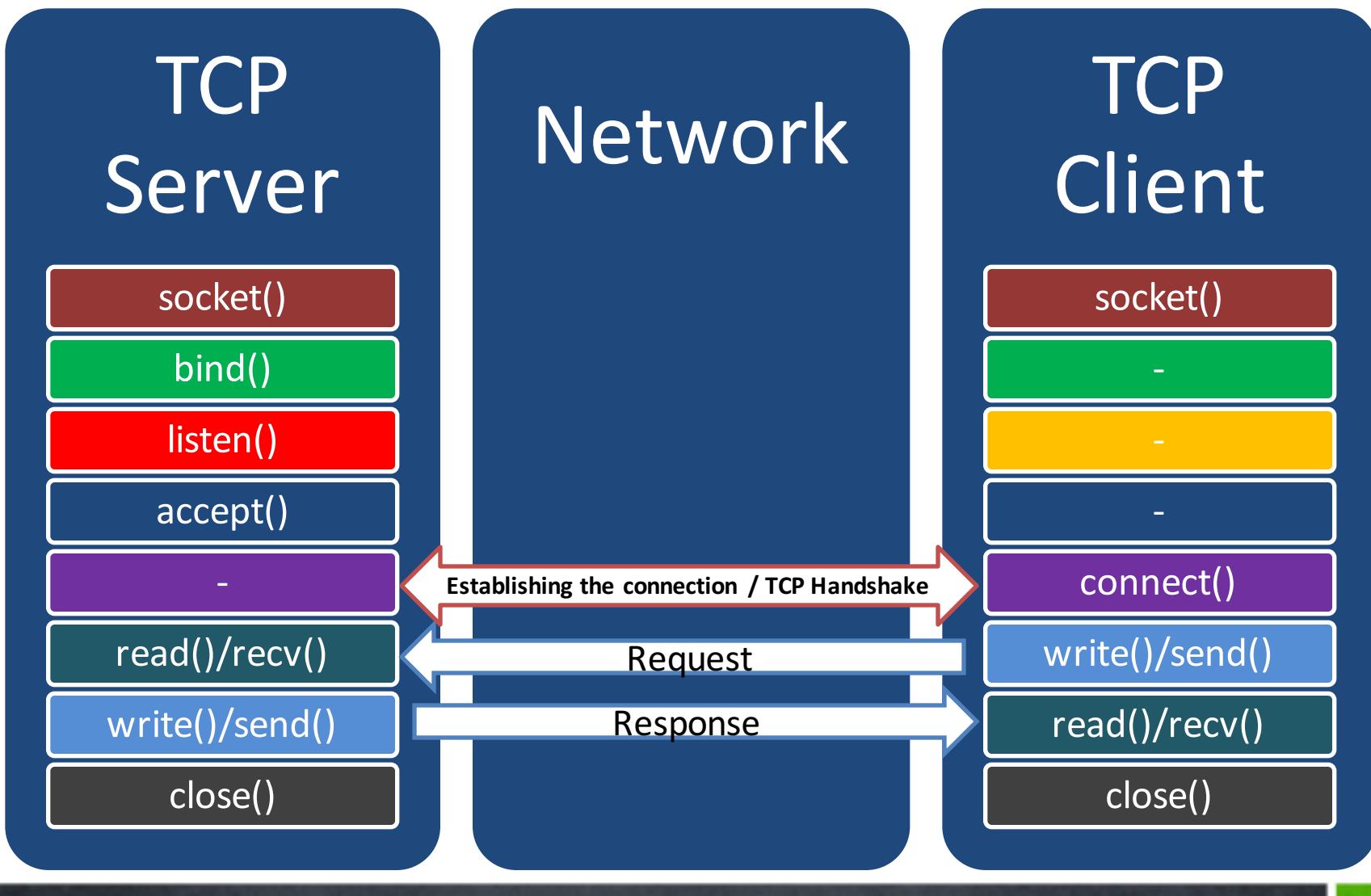
        // send request
        byte[] buf = new byte[256];
        InetAddress address = InetAddress.getByName("127.0.0.1");
        DatagramPacket packet = new DatagramPacket(buf, buf.length, address, 778);
        socket.send(packet);

        // get response
        byte[] bufResp = new byte[256];
        packet = new DatagramPacket(bufResp, bufResp.length);
        socket.receive(packet);

        // display response
        String received = new String(packet.getData());
        System.out.print("Client de la server: " + received);

        // close socket
        socket.close();
    }
}
```

2.2 TCP/IP Networking Programming – TCP Programming – Socket Primitives:



2.2 TCP/IP Networking Programming – TCP Programming – Socket Primitives:



Server

```
ServerSocket serverSocket = null;  
Socket clientSocket = null;  
  
boolean listening = true;  
  
OutputStream os = null; PrintWriter out = null;  
InputStream is = null; BufferedReader in = null;  
String inputLine = null, outputLine = null;
```

```
//SEVERSOCKET = SOCKET+BIND+LISTEN  
serverSocket = new ServerSocket(4801);  
clientSocket = serverSocket.accept(); //ACCEPT
```

```
//STABILIREA CONEXIUNII  
is = clientSocket.getInputStream();  
in = new BufferedReader(new InputStreamReader(is));  
  
os = clientSocket.getOutputStream();  
out = new PrintWriter(os, true);
```

```
while ((inputLine = in.readLine()) != null) {  
    System.out.println(inputLine);  
    outputLine = new String("OK");  
    out.println(outputLine);  
    out.flush();  
    if (outputLine.compareTo("La revedere!") == 0) {break;}  
}
```



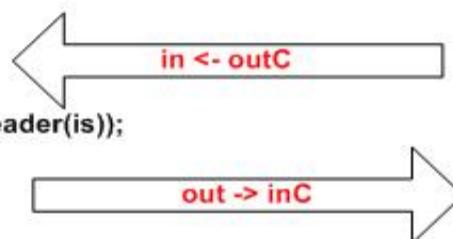
Client

```
Socket clientSocket = null;  
PrintWriter outC = null;  
BufferedReader inC = null;
```

```
clientSocket = new Socket(args[0],  
Integer.parseInt(args[1])); //SOCKET
```

```
//STABILIREA CONEXIUNII  
//CONNECT = OUT2SERVER + INfromSERVER
```

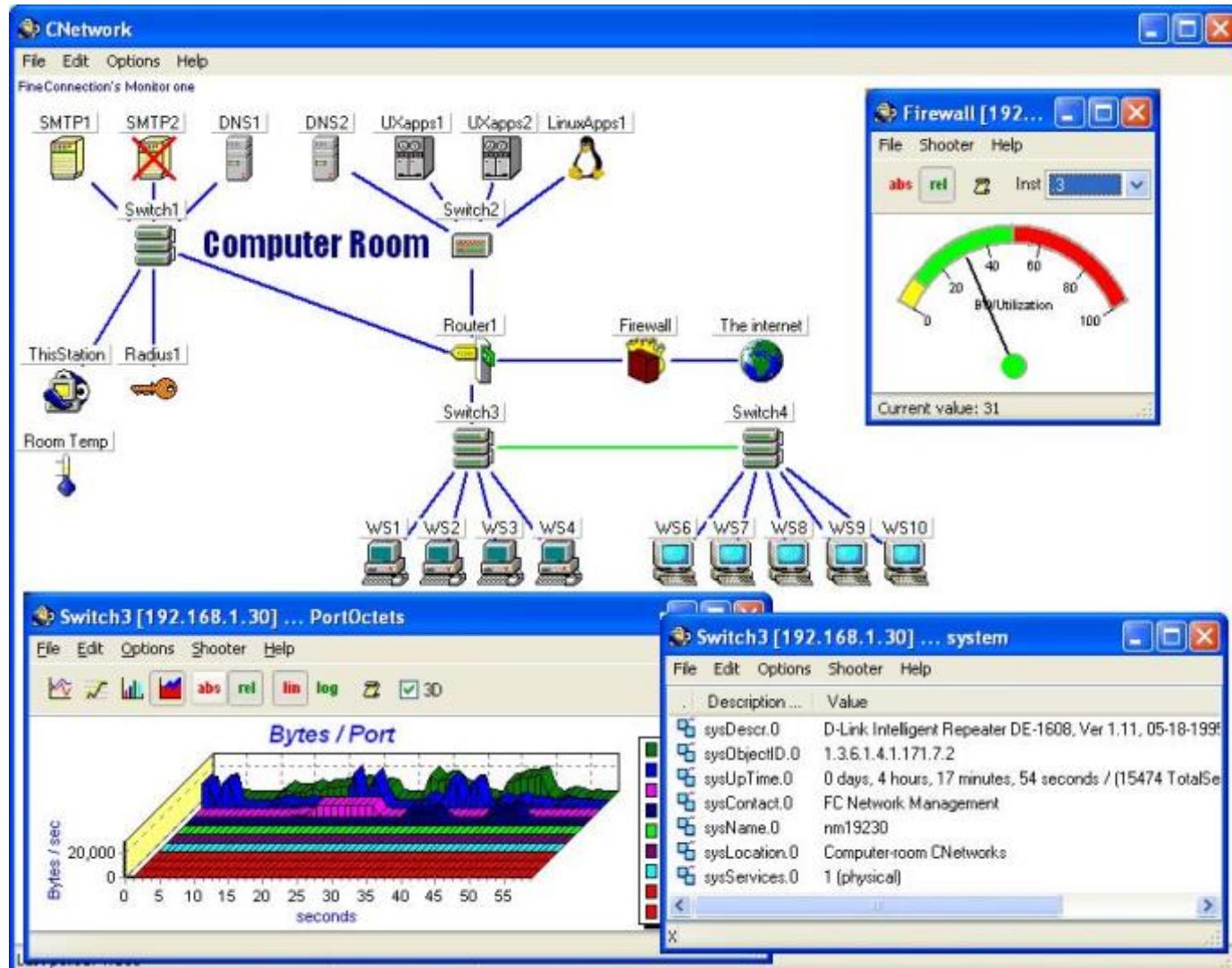
```
//OUT2SERVER  
outC = new PrintWriter(clientSocket.getOutputStream(), true);  
  
//INfromSERVER  
inC = new BufferedReader(new  
InputStreamReader(clientSocket.getInputStream()));
```



```
String lin = "";  
outC.println("As vrea sa ma conectez."); //SEND  
lin = inC.readLine(); //RECV  
System.out.println("Sever: " + lin);
```

2.3 SNMP – RFC 1157

What is the necessity for SNMP – Simple? Network Management Protocol?



2.3 SNMP – RFC 1157

ASN.1 DER and BER:

RFC1157-SNMP DEFINITIONS ::= BEGIN

```
IMPORTS
    ObjectName, ObjectSyntax, NetworkAddress, IpAddress, TimeTicks
    FROM RFC1155-SMI;

-- top-level message

Message ::= SEQUENCE {
    version          -- version-1 for this RFC
        INTEGER {
            version-1(0)                                -- PDUS
        },
    community        -- community name
        OCTET STRING,
    data             -- e.g., PDUs if trivial
        ANY           -- authentication is being used
}

-- protocol data units

PDUs ::= CHOICE {
    get-request      GetRequest-PDU,
    get-next-request GetNextRequest-PDU,
    get-response     GetResponse-PDU,
    set-request      SetRequest-PDU,
    trap             Trap-PDU
}

GetRequest-PDU ::= [0] IMPLICIT PDU
GetNextRequest-PDU ::= [1] IMPLICIT PDU
GetResponse-PDU ::= [2] IMPLICIT PDU
SetRequest-PDU ::= [3] IMPLICIT PDU

PDU ::= SEQUENCE {
    request-id      INTEGER,
    error-status     -- sometimes ignored
        INTEGER {
            noError(0),
            tooBig(1),
            noSuchName(2),
            badValue(3),
        }
}
```

2.3 SNMP – RFC 1157

ASN.1 DER and BER:

Some of the **data types** used in ASN.1 are presented in the following listing:

Data Type	Sort	Tag Number	Meaning
BOOLEAN	Primitive	0x01	Boolean value: yes/no
INTEGER	Primitive	0x02	Negative and positive integers
BIT STRING	Primitive	0x03	Bit sequence
OCTET STRING	Primitive	0x04	Byte sequence (1 byte=8 bits=1octet)
NULL	Primitive	0x05	A null value
OBJECT IDENTIFIER	Primitive	0x06	An object identifier, which is a sequence of integer components that identify an object such as an algorithm or attribute type
PrintableString	Primitive	0x13	An arbitrary string of printable characters
T61String	Primitive	0x14	An arbitrary string of T.61 (eight-bit) characters
IA5String	Primitive	0x16	An arbitrary string of IA5 (ASCII) characters
UTCTime	Primitive	0x17	A "coordinated universal time" or Greenwich Mean Time (GMT) value

2.3 SNMP – RFC 1157

ASN.1 DER and BER:

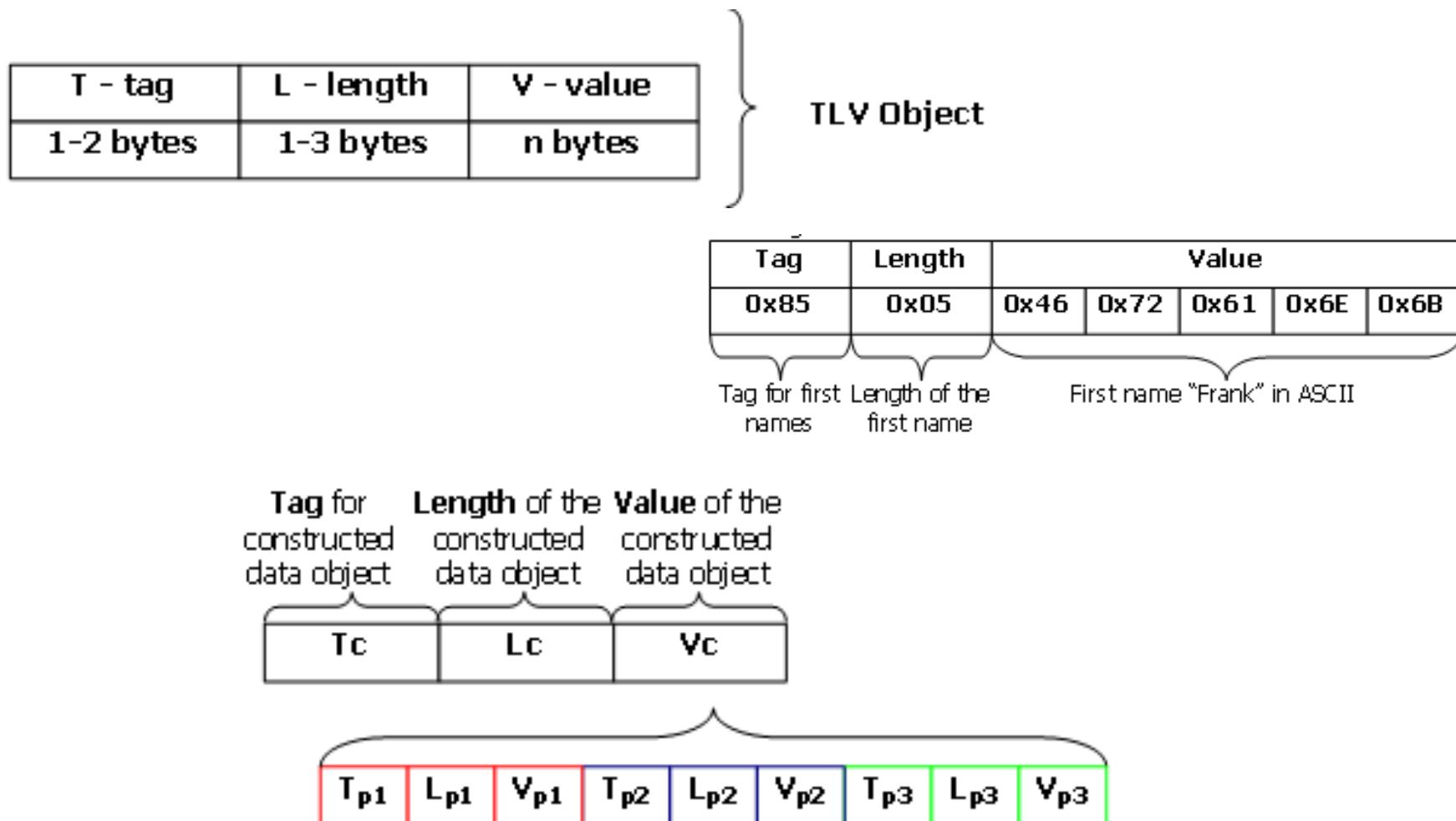
Some of **structured types** defined in ASN.1 are presented in the following listing:

Data Type	Sort	Tag Number	Meaning
SEQUENCE	Constructed	0x30	An ordered collection of one or more types
SEQUENCE OF	Constructed	0x10	An ordered collection of zero or more occurrences of a given type
SET	Constructed	0x31	An unordered collection of one or more types
SET OF	Constructed	0x11	An unordered collection of zero or more occurrences of a given type

A0, A1, ...	Constructed	0xAz	Where z = 0..F in hex and it represents the z-th element in SEQUENCE data type
-------------	-------------	------	--

2.3 SNMP – RFC 1157

ASN.1 TLV:



2.3 SNMP – RFC 1157

ASN.1 DER and BER:

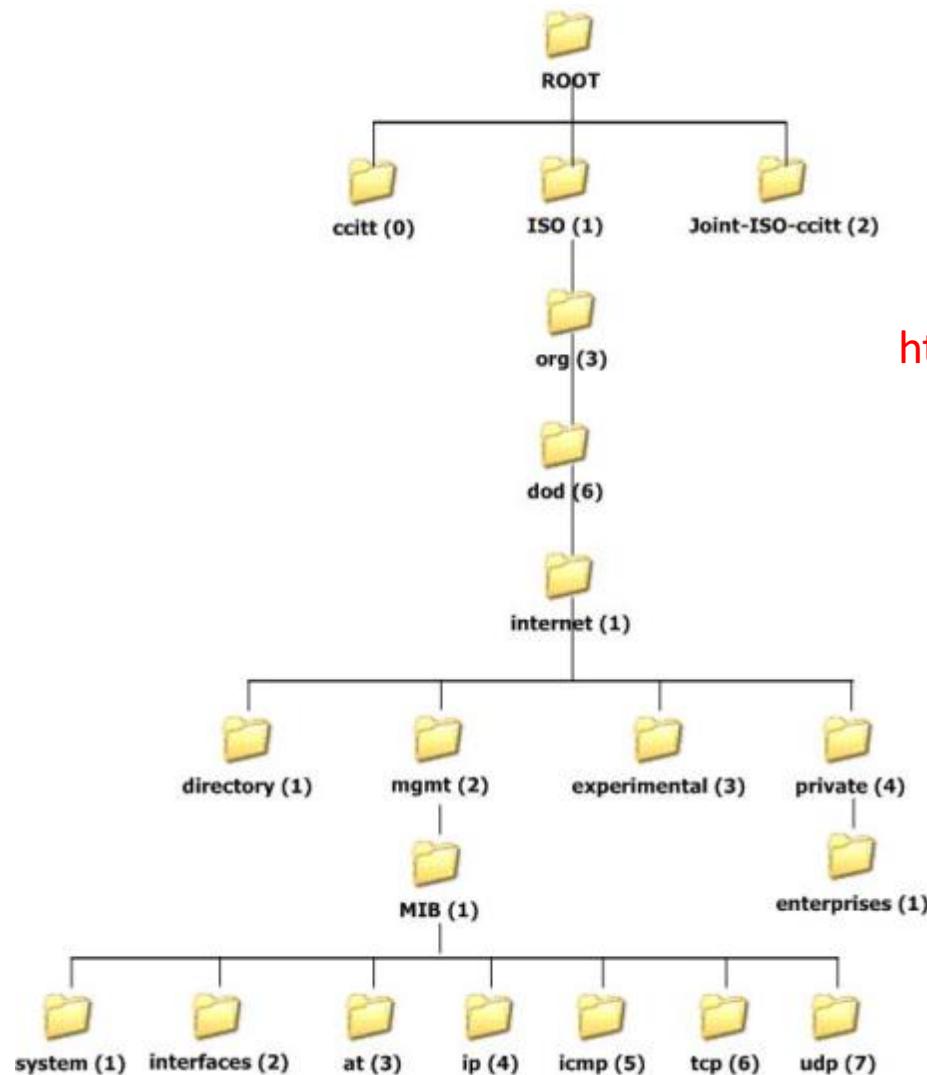
T - tag								L - length															
b7	b6	b5	b4	b3	b2	b1	b0	Meaning	b7	b6	b5	b4	b3	b2	b1	b0	Byte 1	Byte 2	Byte 3	Meaning			
0	0	Universal class	0	0 - 127	One byte is needed			
0	1	Application class	0x81	128-255			Two bytes are needed			
1	0	Context-specific class	The value of tag code in range 31-127								0x82	256-65535		Three bytes are needed			
1	1	Private class															
...	...	0	Primitive data object															
...	...	1	Constructed data obj.															
...	Y	Y	Y	Y	Y	Tag code (0-30)															
...	1	1	1	1	1	Pointer to the following byte (byte 2), which specifies the tag code															

Diagram illustrating the structure of ASN.1 DER/BER bytes:

- Byte 1:** Contains the T-tag (Universal, Application, Context-specific, Private) and L-length (0-127).
- Byte 2:** Contains the tag code (0-30) or a pointer to the following bytes (1-127).
- Bytes 3,4,5:** Contain the length of the data (128-65535).

2.3 SNMP – RFC 1157

ASN.1 OID:



<http://www.oid-info.com/cgi-bin/display>

2.3 SNMP – RFC 1157

ASN.1 OID:

1.2.840.113549.1.1.5

({iso(1) member-body(2) us(840) rsadsi(113549)
pkcs(1) pkcs-1(1) sha1-with-rsa-signature(5)})

sample of OID is 1.2.840.113549.1.1.5 ({iso(1) member-body(2) us(840) rsadsi(113549) pkcs(1) pkcs-1(1) sha1-with-rsa-signature(5)}) for signature obtained from SHA-1 digest on the message and RSA algorithm applied to the digest. This encoding will be at byte level: 0x06 0x09 0x2A 0x86 0x48 0x86 0xF7 0x0D 0x01 0x01 0x05. The first byte shows that here there is a OID – OBJECT IDENTIFIER field. The OID has 9 bytes length because of second byte in array. Because the first bit in length field is not set the length field has only one byte. The length field has value 9 which means the OID structure has the payload data in 9 bytes.

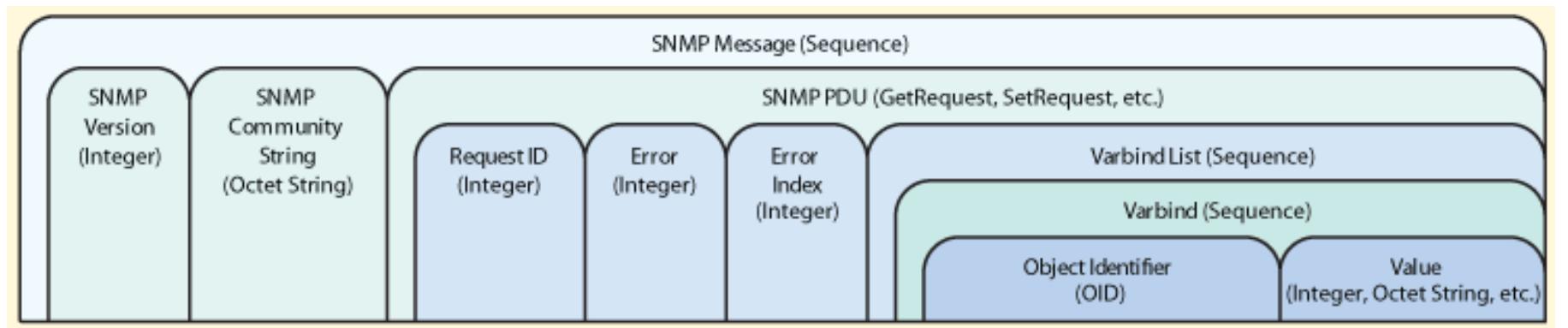
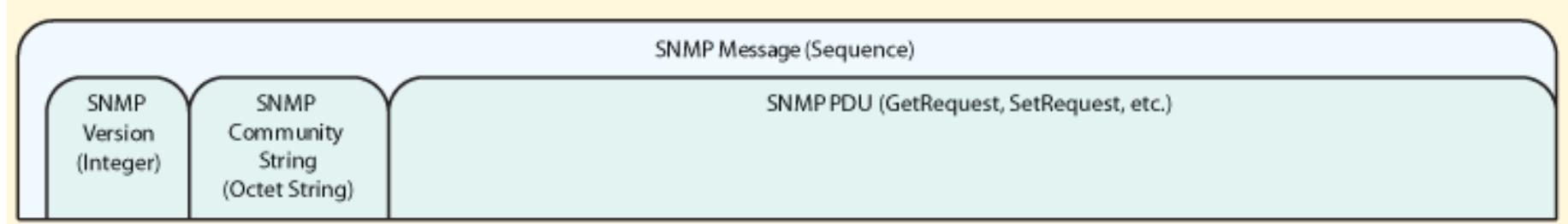
2.3 SNMP – RFC 1157

ASN.1 OID:

According to BER, the first two numbers of any OID (x.y) are encoded as one value using the formula $(40*x)+y$. The first two numbers in an OID are here 1.2. Therefore, the first two numbers of an OID are encoded as 42 or 0x2A, because $(40*1)+2 = 42$. After the first two numbers are encoded, the subsequent numbers in the OID are each encoded as a byte. However, a special rule is required for large numbers because one byte (eight bits) can only represent a number from 0-255. This is the case for 840 and 113549. For 840 from the OID, the first bit of the first byte should be set. The number occupies enough number of bytes till the last byte of representation is not having the first bit set. In 840 case 0x48 has the first bit NOT set. So, the formula is the last hex digit from the first byte multiplied with $2^7 = 128$ (1 bit is for multiple bytes representation) and added with the value from the second byte as long as the second byte has a value less than 0x80. In 840 case, the formula is: $0x06*128 + 0x48$ (from 0x86 0x48)= $768 + 72 = 840$. In case of 113549 we choose the bytes 0x86 0xF7 0x0D because 0x0D is the last byte with first bit (sign bit) not set. The formula for 113549 is:

$0x06*2^{14} + 0x77*2^7 + 0x0d*2^0 = 6*16384 + 119*128 + 13*1 = 98304 + 15232 + 13 = 113549$. For the remaining encoding {...pkcs(1) pkcs-1(1) sha1-with-rsa-signature(5)}, we will have only one byte for each number: 0x01 0x01 0x05.

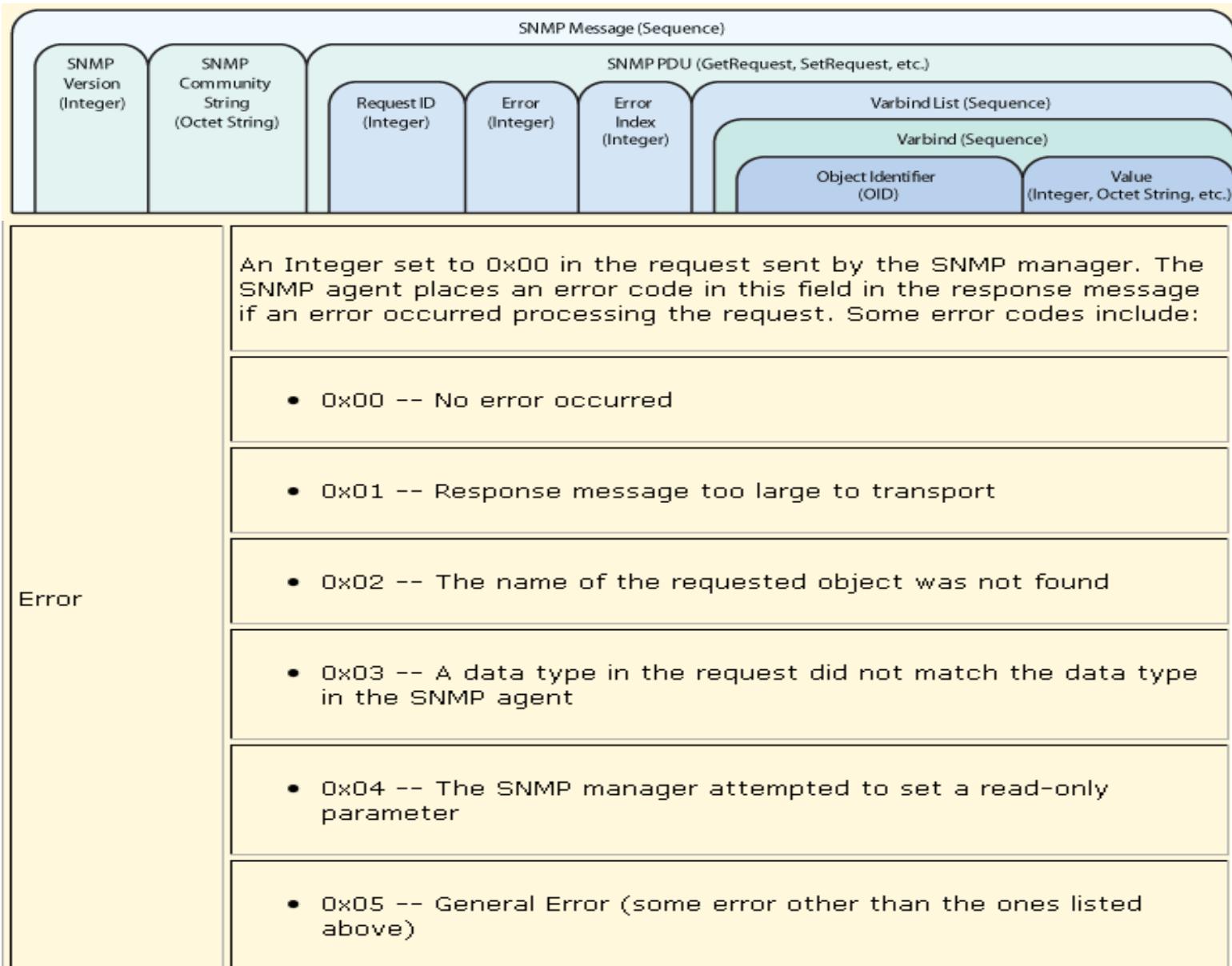
2.3 SNMP – RFC 1157



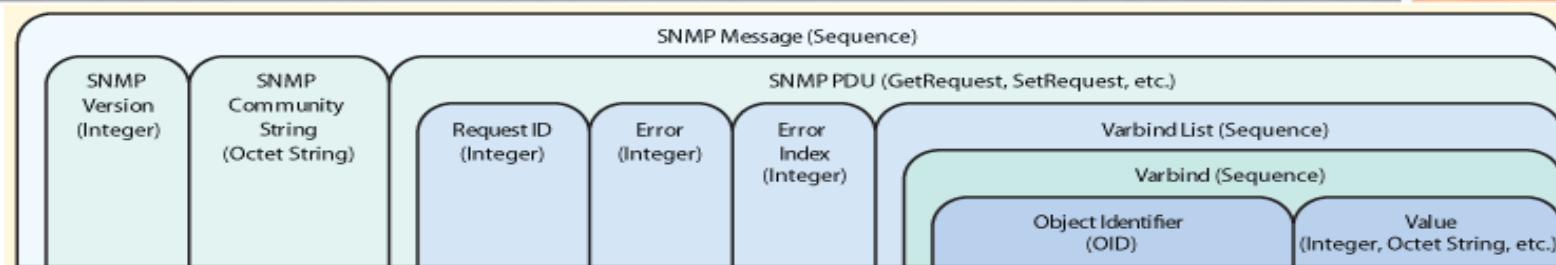
2.3 SNMP – RFC 1157

SNMP Message (Sequence)						
Field	Description					
SNMP Version (Integer)	An Integer that identifies the version of SNMP. SNMPv1 = 0					
SNMP Community String (Octet String)	An Octet String that may contain a string used to add security to SNMP devices.					
SNMP PDU (GetRequest, SetRequest, etc.)	A Sequence representing the entire SNMP message consisting of the SNMP version, Community String, and SNMP PDU.					
Request ID (Integer)	An Integer that identifies a particular SNMP request. This index is echoed back in the response from the SNMP agent, allowing the SNMP manager to match an incoming response to the appropriate request.					
Error (Integer)						
Error Index (Integer)						
Varbind List (Sequence)						
Varbind (Sequence)						
Object Identifier (OID)						
Value (Integer, Octet String, etc.)						

2.3 SNMP – RFC 1157



2.3 SNMP – RFC 1157



Error Index	If an Error occurs, the Error Index holds a pointer to the Object that caused the error, otherwise the Error Index is 0x00.
Varbind List	A Sequence of Varbinds.
Varbind	A Sequence of two fields, an Object ID and the value for/from that Object ID.
Object Identifier	An Object Identifier that points to a particular parameter in the SNMP agent.
Value	<p>SetRequest PDU -- Value is applied to the specified OID of the SNMP agent.</p> <p>GetRequest PDU -- Value is a Null that acts as a placeholder for the return data.</p> <p>GetResponse PDU -- The returned Value from the specified OID of the SNMP agent.</p>

2.3 SNMP – RFC 1157

SNMP Version Type = Integer Length = 1 Value = 0	SNMP Community String Type = Octet String Length = 7 Value = private
30 2C	02 01 00 04 07 70 72 69 76 61 74 65

SNMP Message Type = Sequence, Length = 44

SNMP PDU Type = GetRequest, Length = 30

Request ID Type = Integer Length = 1 Value = 1	Error Type = Integer Length = 1 Value = 0	Error Index Type = Integer Length = 1 Value = 0		Varbind List Type = Sequence, Length = 19		
A0 1E	02 01 01	02 01 00	02 01 00	30 13	30 11 06 0D 2B 06 01 04 01 94 78 01 02 07 03 02 00	Object Identifier Type = Object Identifier Length = 13 Value = 1.3.6.1.4.1.2680.1.2.7.3.2.0 Value Type = Null Length = 0 05 00

2.4 SMTP – RFC 2821

D.1 A Typical SMTP Transaction Scenario

This SMTP example shows mail sent by Smith at host bar.com, to Jones, Green, and Brown at host foo.com. Here we assume that host bar.com contacts host foo.com directly. The mail is accepted for Jones and Brown. Green does not have a mailbox at host foo.com.

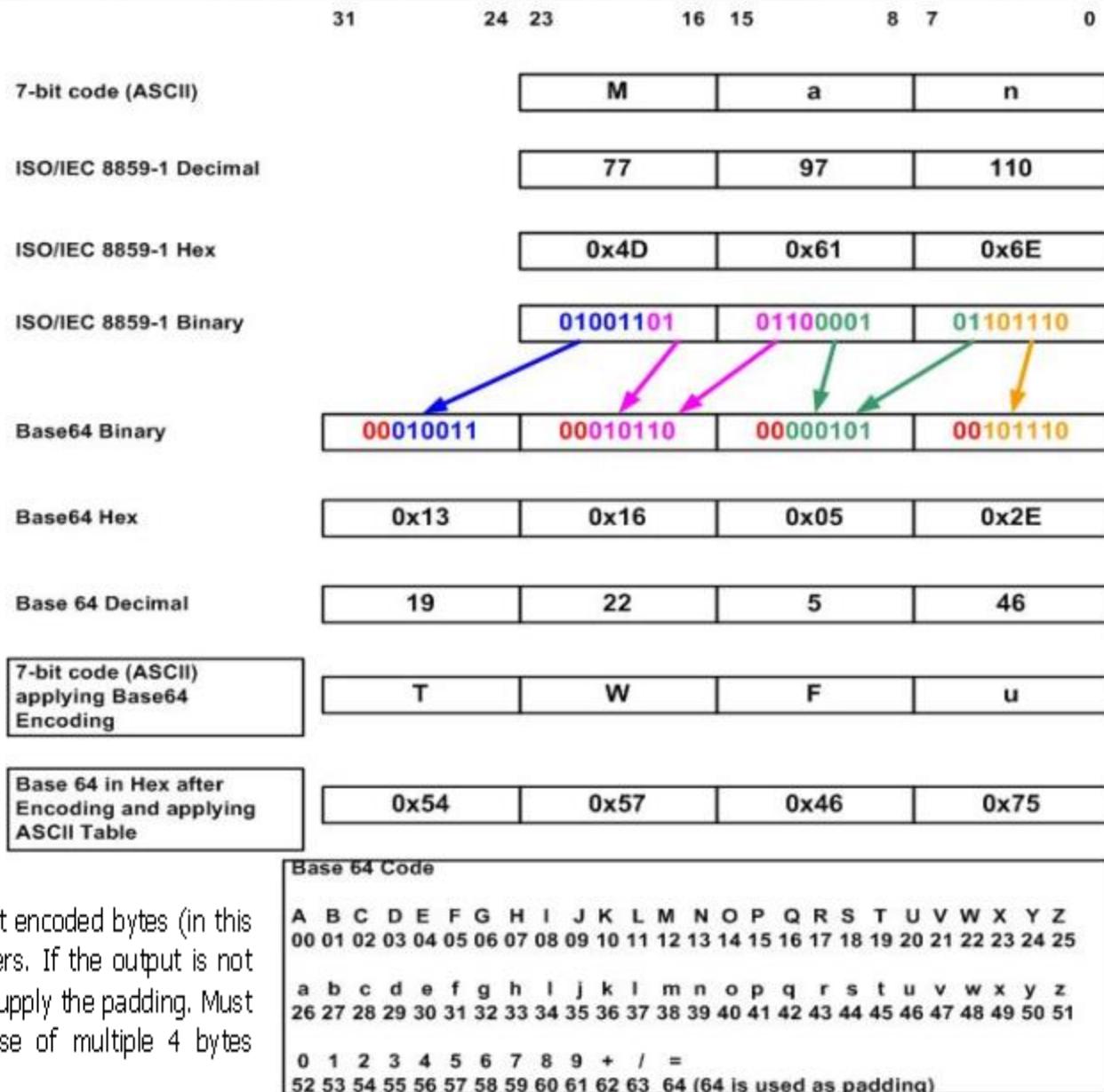
```
S: 220 foo.com Simple Mail Transfer Service Ready
C: EHLO bar.com
S: 250-foo.com greets bar.com
S: 250-8BITMIME
S: 250-SIZE
S: 250-DSN
S: 250 HELP
C: MAIL FROM:<Smith@bar.com>
S: 250 OK
C: RCPT TO:<Jones@foo.com>
S: 250 OK
C: RCPT TO:<Green@foo.com>
S: 550 No such user here
C: RCPT TO:<Brown@foo.com>

S: 250 OK
C: DATA
S: 354 Start mail input; end with <CRLF>. <CRLF>
C: Blah blah blah...
C: ...etc. etc. etc.
C: .
S: 250 OK
C: QUIT
S: 221 foo.com Service closing transmission channel
```

2.4 SMTP – RFC 2821

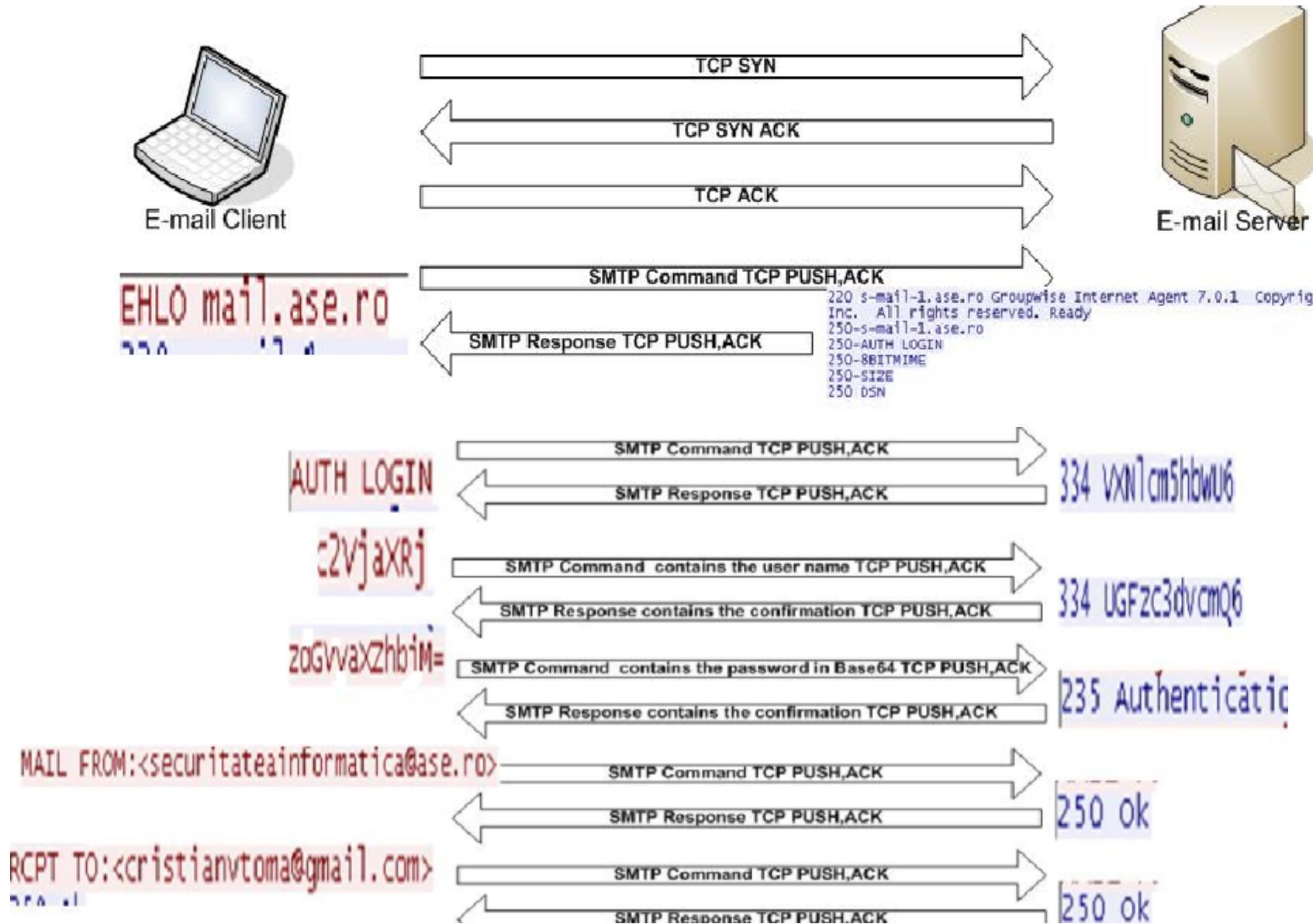
Base64 encoding is used in practice usually for transport over the network and heterogeneous environments binary code such as pictures or executable code. The techniques is very simple: to transform each 3 bytes values into 4 bytes value in order to avoid to obtain values greater then 127 per byte.

For instance, if the scope is to encode the word “Man” into Base64 encoding then it is encoded as “TWFu”. Encoded in ASCII (in ISO 8859-1, one value per byte), M, a, n are stored as the bytes 77 (0x4D), 97 (0x61), 110 (0x6E), which are 01001101, 01100001, 01101110 in base 2.



As this example illustrates, the encoding converts 3 not encoded bytes (in this case, ASCII characters) into 4 encoded ASCII characters. If the output is not multiple of 4 bytes then the '=' sign is put in order to supply the padding. Must be considered the padding with = (64 value) in case of multiple 4 bytes number.

2.4 SMTP – RFC 2821



2.4 SMTP – RFC 2821

DATA

SMTP Command TCP PUSH,ACK

SMTP Response TCP PUSH,ACK

354 Enter mail, end with "." on a line by itself

Date: Sun Jun 01 15:24:37 BST 2008
Subject: SECITC 2008 Account Registration
To: cristianvtoma@gmail.com
From: <secitc@ase.ro>

SMTP Command TCP PUSH,ACK

SMTP Response TCP PUSH,ACK

250 ok

Thank you for registering at SECITC 2008 - 1
Information Technology and Communication.
Please activate your account by clicking the
. <https://www.secitc.eu:443/registration?jID>
. username:
. password:
. Best regards,
SECITC 2008 Team

quit

SMTP Command TCP PUSH,ACK

SMTP Response TCP PUSH,ACK

221 s-mail-1.ase.ro Closing transmission channel

TCP FIN ACK

TCP ACK

TCP FIN ACK

TCP ACK

2.4 SMTP – RFC 2821

```
Frame 4 (72 bytes on wire, 72 bytes captured)
Ethernet II, Src: IntelCor_cd:71:7f (00:13:e8:cd:71:7f), Dst: Giga-Byt_4a:6c:47 (00:14:85:4a:6c:47)
Internet Protocol, Src: 192.168.0.5 (192.168.0.5), Dst: 193.226.34.82 (193.226.34.82)
Transmission Control Protocol, Src Port: 1414 (1414), Dst Port: smtp (25), Seq: 1, Ack: 1, Len: 18
Simple Mail Transfer Protocol
Command: EHLO mail.ase.ro\r\n
```

0000	00	14	85	4a	6c	47	00	13	e8	cd	71	7f	08	00	45	00	...JIG.. .q...E.
0010	00	3a	15	c2	40	00	80	06	40	1a	c0	a8	00	05	c1	e2	.:.@... @.....
0020	22	52	05	86	00	19	e0	63	16	8d	76	06	eb	7e	50	18	"R....c ..v...~P.
0030	b3	a6	33	65	00	00	45	48	4c	4f	20	6d	61	69	6c	2e	..3e..EH LO mail.
0040	61	73	65	2e	72	6f	0d	0a									ase.ro..

simple Mail Transfer Protocol

```
Response: 220 s-mail-1.ase.ro Groupwise Internet Agent 7.0.1 Copyright (c)
Response code: 220
Response parameter: s-mail-1.ase.ro Groupwise Internet Agent 7.0.1 Copy
```

0000	00	13	e8	cd	71	7f	00	14	85	4a	6c	47	08	00	45	00q... .JIG..E.
0010	00	9e	de	88	40	00	31	06	c5	ef	c1	e2	22	52	c0	a8	.:.@.1."R..
0020	00	05	00	19	05	86	76	06	eb	7e	e0	63	16	9f	50	18v. ..~.C..P.
0030	16	d0	29	e8	00	00	32	32	30	20	73	2d	6d	61	69	6c	..)...)22 0 s-mail
0040	2d	31	2e	61	73	65	2e	72	6f	20	47	72	6f	75	70	57	-1.ase.r o Groupw
0050	69	73	65	20	49	6e	74	65	72	6e	65	74	20	41	67	65	ise Inte rnet Age
0060	6e	74	20	37	2e	30	2e	31	20	20	43	6f	70	79	72	69	nt 7.0.1 Copyri
0070	67	68	74	20	28	63	29	20	31	39	39	33	2d	32	30	30	ght (c) 1993-200
0080	36	20	4e	6f	76	65	6c	6c	2c	20	49	6e	63	2e	20	20	6 Novell , Inc.
0090	41	6c	6c	20	72	69	67	68	74	73	20	72	65	73	65	72	All righ ts reser
00a0	76	65	64	2e	20	52	65	61	64	79	0d	0a					ved. Rea dy..

Java Network Programming for easy sharing

Section Conclusions

**Java Network Programming uses for UDP:
DatagramSocket and DatagramPacket classes on
both server and client side.**

**Java Network Programming uses for TCP:
ServerSocket and Socket classes on server side.
Only Socket class on client side.**

**For both server and client, it is necessary to create
byte/char Input (socket.getInputStream()) and
Output (socket.getOutputStream()) streams between
the Random Access Memory – RAM and the network
communications channel.**

**As practical approach, also developed in the laboratory
activities, there is a need to implement “raw” SNMP
request, but third party libraries – *snmp4j.org* or
product – *Nagios, Cacti, Zabbix* may be used.**

**As practical approach, there is a need to implement
“raw” SMTP in order to send e-mails, but
Sun/Oracle or third party libraries may be used –
*java.mail.**.**



Network Programming & Java Sockets

Communicate & Exchange Ideas





Questions & Answers!

But wait...

There's More!





Thanks!



DAD – Distributed Application Development
End of Lecture 2 – Summary of Java SE & Network
Programming – section 2

