

# Query languages with structural and analytic properties

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## 1 MATLANG syntax and semantics

We assume that we have a supply of matrix variables. The definition of an instance  $I$  on MATLANG is a function defined on a nonempty set  $var(I) = \{A, B, M, C, \dots\}$ , that assigns a concrete matrix to each element (matrix *name*) of  $var(I)$ .

Every expression  $e$  is a matrix, either a matrix of  $var(I)$  (*base* matrix, if you will) or a result of an operation over matrices.

The syntax of MATLANG expressions is defined by the following grammar. Every sentence is an expression itself.

$$\begin{aligned}
 e &= M && \text{(matrix variable)} \\
 \text{let } M = e_1 \text{ in } e_2 &&& \text{(local binding)} \\
 e^* &&& \text{(conjugate transpose)} \\
 \mathbf{1}(e) &&& \text{(one-vector)} \\
 \text{diag}(e) &&& \text{(diagonalization of a vector)} \\
 e_1 \cdot e_2 &&& \text{(matrix multiplication)} \\
 \text{apply}[f](e_1, \dots, e_n) &&& \text{(pointwise application of } f)
 \end{aligned}$$

The operations used in the semantics of the language are defined over complex numbers.

- **Transpose:** if  $A$  is a matrix then  $A^*$  is its conjugate transpose.
- **One-vector:** if  $A$  is a  $n \times m$  matrix then  $\mathbf{1}(A)$  is the  $n \times 1$  column vector full of ones.
- **Diag:** if  $v$  is a  $m \times 1$  column vector then  $\text{diag}(v)$  is the matrix

$$\begin{bmatrix}
 v_1 & 0 & 0 & \dots & 0 \\
 0 & v_2 & 0 & \dots & 0 \\
 \dots & \dots & \dots & \dots & \dots \\
 0 & 0 & 0 & \dots & v_m
 \end{bmatrix}$$

- **Matrix multiplication:** if  $A$  is a  $n \times m$  matrix and  $B$  is a  $m \times p$  matrix then  $A \cdot B$  is the  $n \times p$  matrix with  $(AB)_{ij} = \sum_{k=1}^n A_{ik}B_{kj}$ .
- **Pointwise application:** if  $A^{(1)}, \dots, A^{(n)}$  are  $m \times p$  matrices, then  $\text{apply}[f](A^{(1)}, \dots, A^{(n)})$  is the  $m \times p$  matrix  $C$  where  $C_{ij} = f(A_{ij}^{(1)}, \dots, A_{ij}^{(n)})$ .

The formal semantics have a set of rules for an expression  $e$  to be valid on an instance  $I$ , this is,  $e$  succesfully evaluates to a matrix  $A$  on the instance  $I$ . This success is denoted as  $e(I) = A$ . Here  $I[M := A]$  denotes the instance that is equal to  $I$  except that maps  $M$  to the matrix  $A$ .

Expression	Condition for validity
(let $M = e_1$ in $e_2$ )( $I$ ) = $B$	$e_1(I) = A, e_2(I[M := A]) = B$
$e^*(I) = A^*$	$e(I) = A$
$\mathbf{1}(e)(I) = \mathbf{1}(A)$	$e(I) = A$
$\text{diag}(e)(I) = \text{diag}(A)$	$e(I) = A, A$ is a column vector
$e_1 \cdot e_2(I) = A \cdot B$	$\# \text{ columns of } A = \# \text{ rows of } B$
$\text{apply}[f](e_1, \dots, e_n)(I) = \text{apply}[f](A_1, \dots, A_n)$	$\forall k, e_k(I) = A$ and all $A_k$ have the same dimention

For example,

$$\text{let } N = \mathbf{1}(M)^* \text{ in } \text{apply}[c](\mathbf{1}(N)),$$

is an expression, where  $c$  denotes the constant function  $c : x \rightarrow c$ . The result is a  $1 \times 1$  matrix with  $c$  as its entry. This expression is equivalent to  $\text{apply}[c](\mathbf{1}(\mathbf{1}(M)^*))$ .

An example of what can be computed in MATLANG is the mean of a vector:

$$\begin{aligned} &\text{let } N = \mathbf{1}(v)^* \cdot \mathbf{1}(v) \text{ in} \\ &\text{let } S = v^* \cdot \mathbf{1}(v) \text{ in} \\ &\text{let } R = \text{apply}[\div](S, N) \text{ in } R. \end{aligned}$$

Here,  $N$  is the  $1 \times 1$  matrix with the dimension of  $v$  as its entry.  $S$  computes the sum of all the entries of  $v$ . Finally, in  $R$  we store the result of the sum divided by the dimension.

One thing that is worth keeping in mind is that pointwise function application is powerful. With the expression  $\text{apply}[f](\cdot)$  one can compute other elemental operations over matrices that can be studied separately, such as:

- Scalar multiplication: we compute  $c \cdot A$  as

$$\begin{aligned} &\text{let } C = \text{apply}[c](\mathbf{1}(\mathbf{1}(A)^*)) \text{ in} \\ &\text{let } M = \mathbf{1}(A) \cdot C \cdot \mathbf{1}(A^*)^* \text{ in } \text{apply}[\times](M, A). \end{aligned}$$

- Addition: we compute  $A + B$  as

$$\text{let } \text{apply}[+](A, B).$$

- Trace: let

$$m(x, y) = \begin{cases} x & \text{if } x - y > 0 \\ 0 & \text{if } x - y \leq 0 \end{cases}$$

Then we compute the trace of  $A$ ,  $\text{tr}(A)$ , as

$$\begin{aligned} &\text{let } I = \text{diag}(\mathbf{1}(A)) \text{ in} \\ &\text{let } B = \text{apply}[-](A, I) \text{ in} \\ &\text{let } C = \text{apply}[m](A, B) \text{ in} \\ &\text{let } T = \text{apply}[+](A, I) \text{ in } \mathbf{1}(A)^* \cdot T \cdot \mathbf{1}(A) \end{aligned}$$

## 2 Adding canonical vectors to MATLANG

One thing that we cannot do in MATLANG is to obtain a specific entry of a matrix. This entry is expected to be a  $1 \times 1$  matrix. We can do this by adding the standard unit vectors  $e_j$  where

$$e_i = \begin{bmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{bmatrix} \rightarrow i\text{-th position}$$

We know show some examples of what can we express with this new feature. For illustrative reasons, we assume that all the dimensions are well suited for the corresponding operation.

- Get  $A_{ij}$  with  $e_i^* \cdot A \cdot e_j$ .
- The expression  $e_i \cdot e_j^*$  is the matrix that has a 1 in the position  $i, j$  and zero everywhere else.
- Given a vector  $v$ , the expression  $v \cdot e_i^*$  is the matrix

$$\begin{bmatrix} 0 & \cdots & v_1 & \cdots & 0 \\ 0 & \cdots & v_2 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & v_n & \cdots & 0 \end{bmatrix}$$

- Replace column  $j$  of  $A$  with zeros:  $A(I - e_j \cdot e_j^*)$ .
- Replace column  $j$  of  $A$  with a vector  $v$ :  $A(I - e_j \cdot e_j^*) + v \cdot e_j^*$ .

Note that  $I = \text{diag}(\mathbf{1}(A))$  and the sum of matrices can be implemented as  $\text{apply}[+](A, B)$ .

### 3 Connection with logic

We show the expressive power of MATLANG. Let RRA be the class of algebras of binary relations with the operations all, identity, set difference, converse and relational composition. Also, let  $\text{FO}_b^3$  denote first-order logic with three variables, equality and infinitely many binary relation symbols. This is,  $\text{FO}_b^3$  are  $\text{FO}^3$  graph queries.

It is known that the logic captured by RRA is  $\text{FO}_b^3$ . Now, we show that RRA can be interpreted into MATLANG, thus the expressive power of MATLANG is at least the same as  $\text{FO}_b^3$ .

Let  $U$  be a nonempty finite set with  $n$  elements (and thus has an enumeration  $u_1, \dots, u_n$ ). Let  $\mathcal{A}^U \in \text{RRA}$ , this is

$$\mathcal{A}^U = \langle A, \cup, -, \circ, {}^{-1}, I \rangle,$$

where

- $A \subseteq \mathcal{P}(U \times U)$ , this is,  $\forall R \in A. R \subseteq U \times U$ . So  $A$  is a set of binary relations over  $U$ .
- $\cup, -, \circ, ^{-1}$  denote the operations union, set difference, relational composition and converse, respectively. All of them defined over binary relations.

- $I$  is the constant relation symbol that denotes the set  $\{(u, u) : u \in U\}$ .

Let  $R \in A$ . Define  $M^R$  a matrix such that

$$M_{ij}^R = \begin{cases} 1 & \text{if } (u_i, u_j) \in R \\ 0 & \text{if } (u_i, u_j) \notin R \end{cases}$$

Note that the MATLANG instance of  $\mathcal{A}$ , has  $|A|$  matrices with dimensions  $|U| \times |U|$ . Thus binary relations are represented as adjacency matrices.

We now show how to express the RRA operations in MATLANG.

- **All:** let  $R \in A$  be any relation. We express  $U \times U$  as  $M^{U \times U} = \mathbf{1}(M^R) \cdot \mathbf{1}(M^R)^*$ .
- **Identity:** let  $R \in A$  be any relation. Then we express the constant relation  $I$  as  $M^I = \text{diag}(\mathbf{1}(M^R))$ .
- **Union:** let  $R, S \in A$ . Then  $M^{R \cup S} = \text{apply}[x \vee y](M^R, M^S)$ .
- **Set difference:** let  $R, S \in A$ . Then  $M^{R-S} = \text{apply}[x \vee \neg y](M^R, M^S)$ .
- **Converse:** let  $R \in A$ . Then we express  $R^{-1} = \{(u, v) \in U \times U : (v, u) \in R\}$  as  $M^{R^{-1}} = \mathbf{t}(M^R)$ .
- **Relational composition:** let  $R, S \in A$ , then  $M^{R \circ S} = \text{apply}[x > 0](M^R \cdot M^S)$ . Where

$$x > 0 = \begin{cases} 1 & \text{if } x > 0 \\ 0 & \text{if } x \leq 0 \end{cases}$$

Thus RRA can be interpreted into MATLANG, and hence so does  $\text{FO}_b^3$ , this is,  $\text{FO}^3$  graph queries.

## 4 Boolean MATLANG

## 5 Expressions as functions

Another way of looking at expressions in MATLANG is as functions between matrix spaces, this is

$$e : (M_1, \dots, M_k) \rightarrow M,$$

where  $M, M_1, \dots, M_k$  are matrix spaces, i.e.,  $M, M_1, \dots, M_k \in \{\mathcal{M}^{n \times m} : n, m \in \mathbb{N}^+\}$ .

Some examples:

- If  $A$  is  $n \times m$  then

$$\begin{aligned} e(A) = t(A) : \mathcal{M}_{n \times m} &\rightarrow \mathcal{M}_{m \times n} \\ A &\rightarrow A^* \end{aligned}$$

- If  $A$  is  $n \times m$  then

$$e(A) = \mathbf{1}(A) : \mathcal{M}_{n \times m} \rightarrow \mathcal{M}_{n \times 1}$$

$$A \rightarrow n \text{ times} \left\{ \begin{bmatrix} 1 \\ \vdots \\ 1 \\ \vdots \\ 1 \end{bmatrix} \right.$$

- If  $v$  is  $n \times 1$  then

$$e(v) = \text{diag}(v) : \mathcal{M}_{n \times 1} \rightarrow \mathcal{M}_{n \times n}$$

$$v \rightarrow \begin{bmatrix} v_1 & 0 & 0 & \dots & 0 \\ 0 & v_2 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & v_n \end{bmatrix}$$

- If  $A$  is  $n \times m$  and  $B$  is  $m \times p$  then

$$e(A, B) = A \cdot B : \mathcal{M}_{n \times m} \times \mathcal{M}_{m \times p} \rightarrow \mathcal{M}_{n \times p}$$

$$(A, B) \rightarrow A \cdot B$$

- If  $A^{(1)}, \dots, A^{(n)}$  are  $m \times p$  matrices then

$$e(A^{(1)}, \dots, A^{(n)}) = \text{apply}[f](A^{(1)}, \dots, A^{(n)})$$

has domains

$$\mathcal{M}_{m \times p}^n \rightarrow \mathcal{M}_{m \times p}$$

$$(A^{(1)}, \dots, A^{(n)}) \rightarrow C : C_{ij} = f(A_{ij}^{(1)}, \dots, A_{ij}^{(n)}).$$

We can start to analyze if this functions are increasing, decreasing, boolean, etc. Also, we can study the effects of disturbances on the input in the output of these functions.

## 6 Core of Matlab and R

### MATLAB

The basic operations of MATLAB over matrices are:

- **mldivide**(A, B): returns  $x$  such that  $Ax = B$ .
- **descomposition**(A): returns a decomposition or factorization  $LU, LDL, QR$ , Cholesky, etc.
- **inv**(A): returns  $A^{-1}$ .
- **multiplication**: compute  $A \cdot B$ .

- **transpose(A)**: returns  $A^T$ .
- **conjugate transpose(A)**: returns  $A'$ .
- **matrix power(A, k)**: returns  $A^k$ .
- **eigen(A)**: returns the eigenvectors and the eigenvectors matrices of  $A$ .
- **funm(A, f)**: returns matrix  $B$  with elements  $b_{ij} = f(a_{ij})$ .
- **crossprod(a, b)**: vectorial product, returns  $c$  such that  $c \perp a, b$ .
- **dotprod(a,b)**: returns  $a \cdot b$ .
- **diag(v)**:  $v$  vector. Returns the following matrix:

$$\begin{bmatrix} v_1 & 0 & 0 & \dots & 0 \\ 0 & v_2 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & v_n \end{bmatrix}$$

- **diag(A)**: given matrix  $A$ , it returns

$$\begin{bmatrix} a_{11} \\ a_{22} \\ \vdots \\ a_{nn} \end{bmatrix}$$

- **det(A)**: returns the determinant of  $A$ .
- **zeros(n, m)**: returns a  $n \times m$  matrix full of zeros.
- **ones(n, m)**: returns a  $n \times m$  matrix full of ones.
- **A[i, j]**: you can get  $A_{ij}$ .

## R

The basic operations of the language R over matrices are:

- **A%%B**: matrix multiplication.
- **A\*B**: pointwise multiplication.
- **t(A)**: transpose.
- **diag(v)**: returns the matrix

$$\begin{bmatrix} v_1 & 0 & 0 & \dots & 0 \\ 0 & v_2 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & v_n \end{bmatrix}$$

- **diag(A)**: Returns the vector

$$\begin{bmatrix} a_{11} \\ a_{22} \\ \vdots \\ a_{nn} \end{bmatrix}$$

- **diag(k)**:  $k$  scalar. It creates the  $k \times k$  identity matrix.
- **matrix(k, n, m)**: returns the  $n \times m$  matrix, where every entry is equal to  $k$ .
- **solve(A, b)**: returns  $x$  such that  $Ax = b$ .
- **solve(A)**: returns  $A^{-1}$ .
- **det(A)**: determinant of  $A$ .
- **y<-eigen(A)**: stores the eigenvalues of  $A$  in `y$val` and the eigenvectors in `y$vec`.
- **y<-svd(A)**: it computes and stores the following:
  - `y$d`: vector of the singular values of  $A$ .
  - `y$u`: matrix of the left singular vectors of  $A$ .
  - `y$v`: matrix of the right singular vectors of  $A$ .
- **R<-chol(A)**: Cholesky factorization,  $R'R = A$ .
- **y<-qr(A)**:  $QR$  decomposition, stored in `y$qr`.
- **cbind(A,B, v, ...)**: joins matrices and vector horizontally, returns a matrix.
- **rbind(A,B, v, ...)**: joins matrices and vector vertically, returns a matrix.
- **rowMeans(A)**: returns the vector of the averages over the rows of  $A$ .
- **colMeans(A)**: returns the vector of the averages over the columns of  $A$ .
- **rowSums(A)**: returns the vector of the sums over the rows of  $A$ .
- **colSums(A)**: returns the vector of the sums over the columns of  $A$ .
- **outer(A, B, f)**: applies  $f(\cdot, \cdot)$ . Returns matrix  $C$  of entries  $c_{ij} = f(a_{ij}, b_{ij})$ .
- **A[i, j]**: you can get  $A_{ij}$ .