

$$d_k = \frac{(2n-k)!}{2^{n-k} (k!) (n-k)!} \quad k = 0, 1, 2, \dots, n$$

$B_n(s)$ satisfies the following recursive equation

$$B_n(s) = (2n-1)B_{n-1}(s) + s^2 B_{n-2}(s)$$

with

$$B_0(s) = 1, \quad B_1(s) = s+1$$

(4) cut-off frequency varies with the order of the filter (disadvantage)

$$\omega_c = \omega_0^{1/n}$$

Before
Bilinear
z-transform

IIR Digital Filter Design Methods

Indirect methods

1- Impulse Invariant Method:

Impulse response of the digital filter is the sampled version of the impulse response of the analog filter. Assume that the analog filter has

$$H_A(s) = \sum_{i=1}^n \frac{A_i}{(s+\alpha_i)}$$

$-\alpha_i$: distinct poles of $H_A(s)$.

$$A_i = (s+\alpha_i) H_A(s) \Big|_{s=-\alpha_i}$$

If all the poles are ~~real and~~ distinct then

$$h_A(t) = \mathcal{L}^{-1}[H_A(s)] = \sum_{i=1}^n A_i e^{-\alpha_i t}$$

The impulse response of the digital filter is

$$h_A(t) \big|_{t=KT} = h_D(KT) = \sum_{i=1}^n A_i e^{-\alpha_i KT}$$

The transfer function $H_D(z)$ is

$$\begin{aligned} H_D(z) &= \sum_{k=0}^{\infty} h_D(KT) z^{-k} \\ &= \sum_{k=0}^{\infty} \sum_{i=1}^n A_i e^{-\alpha_i KT} z^{-k} = \sum_{i=1}^n A_i \sum_{k=0}^{\infty} (z^{-1} e^{-\alpha_i T})^k \\ &= \sum_{i=1}^n \frac{A_i}{(1 - e^{-\alpha_i T} z^{-1})} \end{aligned}$$

Comparison with

$$H_A(s) = \sum_{i=1}^n \frac{A_i}{(s + \alpha_i)}$$

gives the mapping as

$$H_D(z) = H_A(s) \big|_{(s + \alpha_i) \rightarrow (1 - z^{-1} e^{-\alpha_i T})}$$

Thus $H_D(z)$ can be obtained from $H_A(s)$ without evaluating $h_A(t)$ or $h_D(KT)$.

For stable analog filter $\alpha_i > 0$ for $\forall i \in [1, n]$

The corresponding digital filter has poles at

$$z = e^{-\alpha_i T} \text{ with } |z| < 1 \text{ when } \alpha_i > 0$$

i.e. the digital filter is also ^{stable} filter.

Using the sampling, the frequency responses are related by

$$H_D(e^{j\Omega}) = \frac{1}{T} \sum_{n=-\infty}^{\infty} H_A\left(\frac{\Omega + 2\pi n}{T}\right) \quad \Omega: \text{digital freq.}$$

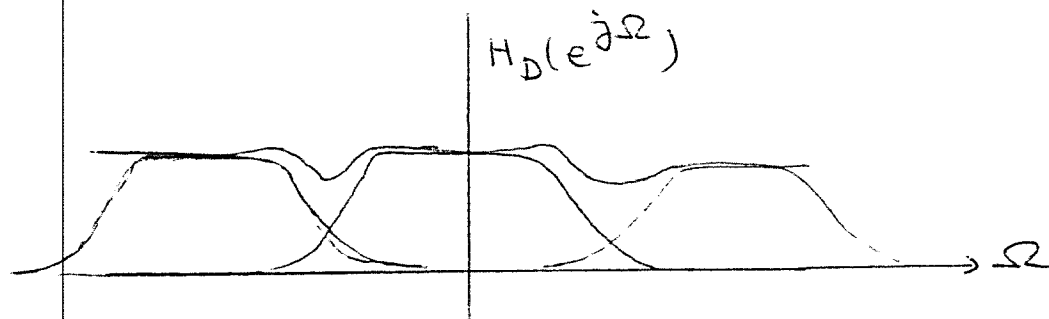
or since $\omega T = \Omega$

$$H_D(\omega) = \frac{1}{T} \sum_{n=-\infty}^{\infty} H_A(\omega + n\omega_s)$$

$\frac{2\pi}{T} = \omega_s$ ω : Analogy freq.

If $H_A(\omega) = 0$ when $|\omega| \geq \frac{\pi}{T} = \frac{\omega_s}{2}$ i.e. Bandlimited
 then $H_D(e^{j\Omega}) = \frac{1}{T} H_A\left(\frac{\Omega}{T}\right)$ i.e. no aliasing

Unfortunately $H_A(\omega)$ is not bandlimited in practical cases and consequently aliasing problem is inevitable



As a result this method is appropriate only for Bandlimit LPF and HPF and BPF would require additional bandlimiting to avoid severe aliasing distortion.

Remarks

1- Let $\alpha_i = \sigma + j\omega$ then

$$z = e^{-\sigma T} e^{-j\omega T}$$

Compare with $z = r e^{-j\Omega} \Rightarrow r = e^{-\sigma T}$ and

$$\boxed{\Omega = \omega T}$$

i.e. linear mapping. (between frequencies)

2- For Complex Conjugate poles one can alternatively use

$$\frac{s + \sigma}{s^2 + 2\sigma s + \sigma^2 + \omega^2} \xrightarrow{\text{mapped}} \frac{1 - \bar{z}^{-1} e^{-\sigma T} \cos \omega T}{1 - 2\bar{z}^{-1} e^{-\sigma T} \cos \omega T + \bar{z}^{-2} e^{-2\sigma T}}$$

$$\frac{\omega}{s^2 + 2\sigma s + \sigma^2 + \omega^2} \xrightarrow{\text{mapped}} \frac{\bar{z}^{-1} e^{-\sigma T} \sin \omega T}{1 - 2\bar{z}^{-1} e^{-\sigma T} \cos \omega T + \bar{z}^{-2} e^{-2\sigma T}}$$

Example (LPP)

Design a digital filter of Butterworth type with

$$n = 3, \quad \omega_s = 15 \omega_c$$

using impulse invariant method.

From the tables we can get the normalized transfer function

$$H_N(s) = \frac{1}{1 + 2s + 2s^2 + s^3} = \frac{1}{(s+1)(s^2+s+1)}$$

$$= \frac{A_1}{(s+1)} + \frac{A_2}{(s+\alpha)} + \frac{A_3}{(s+\alpha^*)}$$

where

$$\alpha = \frac{1}{2} (1 - j\sqrt{3})$$

Direct Methods (after Bilinski 2-)

Least Square Inverse Design

This method leads to a set of linear equations

$\{h_D(n)\}$: $n \in [0, L-1]$ Desired Impulse Response (first L samples)

Filter transfer function

$$H(z) = \frac{b_0}{1 - \sum_{k=1}^N a_k z^{-k}}$$

must find a_k 's

Generalization is proposed by Shanks and Burrus and Parks

The output of the inverse of $H(z)$ must approximate a unit sample ($\delta(n)$) when the input is $h_D(n)$. If $v(n)$ denotes

Let $v(n)$: the output of inverse system with transfer function

$1/H(z)$ then

$$V(z) = \frac{H_D(z)}{H(z)}$$



Thus we can write

$$b_0 v(n) = h_D(n) - \sum_{r=1}^N a_r h_D(n-r)$$

Recall that we require $v(n) \approx \delta(n)$ thus $b_0 = h_D(0)$

and that $v(n)$ be as small as possible for $n > 0$. Thus we

Choose the remaining coeff to minimize

$$E = \sum_{n=1}^{\infty} (v(n))^2$$

$$= \frac{1}{b_0^2} \left[\sum_{n=1}^{\infty} (h_D(n))^2 - 2 \sum_{n=1}^{\infty} h_D(n) \sum_{r=1}^N a_r h_D(n-r) + \sum_{n=1}^{\infty} \left[\sum_{r=1}^N a_r h_D(n-r) \right]^2 \right]$$

$$\frac{\partial E}{\partial a_i} = 0$$

$$i \in [1, N]$$

$$\sum_{r=1}^N a_r \sum_{n=1}^{\infty} h_D(n-r) h_D(n-i) = \sum_{n=1}^{\infty} h_D(n) h_D(n-i)$$

Define the autocorrelation function

$$\Phi(i, r) = \sum_{n=1}^{\infty} h_D(n-r) h_D(n-i)$$

Then the above equation becomes a normal equation

$$\sum_{r=1}^N a_r \Phi(i, r) = \Phi(i, 0) \quad i \in [1, N]$$

These equations can be solved using GE or LU factorization.

~~Toeplitz form~~ The correlation matrix is Toeplitz and efficient procedure known as Levinson algorithm exist for the solution

$$a_1 \Phi(i, 1) + \dots + a_N \Phi(i, N) = \Phi(i, 0)$$

$$\begin{bmatrix} \Phi(1,1) & \dots & \Phi(1,N) \\ \Phi(2,1) & \Phi(2,2) & \dots & \Phi(2,N) \\ \vdots & \vdots & \vdots & \vdots \\ \Phi(N,1) & \dots & \Phi(N,N) \end{bmatrix} \begin{bmatrix} a_1 \\ \vdots \\ a_N \end{bmatrix} = \begin{bmatrix} \Phi(1,0) \\ \vdots \\ \Phi(N,0) \end{bmatrix}$$

then the error introduced

$$E = \sum_{n=1}^{\infty} (v(n))^2 \quad \text{MSE}$$

we have to select a_i 's such that E is minimum.

$$\hat{a} \quad \frac{\partial E}{\partial a_i} = 0 \quad \forall i \in [1, N]$$

$$E = \frac{1}{b_0^2} \sum_{n=1}^{\infty} (h_D(n))^2 - 2 \sum_{n=1}^{\infty} h_D(n) \sum_{r=1}^N a_r h_D(n-r) + \sum_{n=1}^{\infty} \left[\sum_{r=1}^N a_r h_D(n-r) \right]^2$$

$$\therefore \frac{\partial E}{\partial a_i} = -2 \sum_{n=1}^{\infty} h_D(n) h_D(n-i) + 2 \sum_{r=1}^N a_r \sum_{n=1}^{\infty} h_D(n-r) h_D(n-i) = 0$$

Denote

$$p(i, r) \triangleq \sum_{n=1}^{L+r-1} h_D(n-r) h_D(n-i)$$

$$p(i, 0) \triangleq \sum_{n=1}^{L-1} h_D(n) h_D(n-i)$$

then

$$\sum_{r=1}^N a_r p(i, r) = p(i, 0) \quad (*) \quad i \in [1, N]$$

using (*) we can compute a_i 's.

Direct Method

This is direct design method for IIR filters.

(i) Least Square Inverse Method

Solving a set of linear eq^s
 Desired impulse response
 Let $\{h_D(n)\}$: 1st L-point of the impulse response $n \in [0, L-1]$

$$H(z) = \frac{b_0}{1 - \sum_{k=1}^N a_k z^{-k}} = \sum_{n=0}^{\infty} h(n) z^{-n}$$

objective: find a_k 's



$$\begin{aligned} V(z) &= H_D(z) \cdot \frac{1}{H(z)} \\ &= H_D(z) \cdot \left[\frac{1}{b_0} \left(1 - \sum_{k=1}^N a_k z^{-k} \right) \right] \end{aligned}$$

take IZT

$$b_0 v(n) = h_D(n) - \sum_{k=1}^N a_k h_D(n-k)$$

require that $v(n) \approx s(n) \Rightarrow b_0 = h_D(0)$

all $v(n)$, $n > 0$ should be very small.

eg $r=1$

$$p(0,1) = \sum_{n=1}^L h_0(n-1) h_0(n)$$

$$= h_0(0) \cdot h_0(1) + h_0(1) h_0(2) + \dots \\ \dots + h_0(L-1) \cdot h_0(L)$$

system of linear eq^s can be solved simultaneously to get the values of a_i 's.

$$\begin{bmatrix} p(1,1) & p(1,2) & \dots & p(1,N) \\ p(2,1) & p(2,2) & \dots & p(2,N) \\ \vdots & & & \\ p(N,1) & \dots & \dots & p(N,N) \end{bmatrix} \begin{bmatrix} a_1 \\ a_2 \\ \vdots \\ a_N \end{bmatrix} = \begin{bmatrix} p(1,0) \\ p(2,0) \\ \vdots \\ p(N,0) \end{bmatrix}$$

use ^{- Durbin} Levinson Algorithm to solve the above.

(2) Linear Programming Approach:

Given the desired spurred magnitude response for the desired filter $A_d(\omega)$.
Design a filter with $A(\omega)$ as close as possible to $A_d(\omega)$ and obtain the relevant coefficient.

$$A(\omega) = H(\omega) \cdot H(-\omega)$$

Let $H(z) = \sum_{i=0}^M b_i z^{-i}$

$$H(z) = \sum_{j=0}^N a_j z^{-j}$$

$$\begin{aligned} \therefore H(z) H(z^{-1}) &= \left(\sum_{i=0}^M b_i z^{-i} \right) \left(\sum_{i=0}^M b_i z^i \right) \\ &= \left(\sum_{j=0}^N a_j z^{-j} \right) \left(\sum_{j=0}^N a_j z^j \right) \end{aligned}$$

$$= \sum_{i=-M}^M c_i z^{-i}$$

$$= \sum_{j=-N}^N d_j z^j$$

c_i 's and d_j 's are expressed in terms of a_i 's and b_j 's and we have symmetry property

$$c_i = c_{-i}$$

$$d_i = d_{-i}$$

now we can write \tilde{A}^2 as follows;

$$\tilde{A}^2(e^{j\Omega}) = \frac{H(z) H(z^*)}{z = e^{j\Omega}}$$

$$= \frac{c_0 + \sum_{i=1}^M 2c_i \cos i\Omega}{d_0 + \sum_{i=1}^N 2d_i \cos i\Omega}$$

$$\triangleq \frac{N_1(e^{j\Omega})}{D_1(e^{j\Omega})}$$

Given $\tilde{A}_D^2(\Omega)$ find $\tilde{A}^2(\Omega) = \frac{N_1(e^{j\Omega})}{D_1(e^{j\Omega})}$ such that

$$-E(\Omega) \leq \tilde{A}^2(\Omega) - \tilde{A}_D^2(\Omega) \leq E(\Omega) \quad \text{--- tolerance function}$$

$$-E(\Omega) \leq \frac{N_1(e^{j\Omega})}{D_1(e^{j\Omega})} - \tilde{A}_D^2(\Omega) \leq E(\Omega)$$

$E(\Omega)$: tolerance function

$$+ N_1(e^{j\Omega}) - D_1(e^{j\Omega}) [\tilde{A}_D^2(\Omega) + E(\Omega)] \leq 0$$

$$- N_1(e^{j\Omega}) + D_1(e^{j\Omega}) [\tilde{A}_D^2(\Omega) - E(\Omega)] \leq 0$$

$$N_1(e^{j\Omega}) \neq 0$$

$$D_1(e^{j\Omega}) \neq 0$$

now introduce a dummy variable s to both equations.

$$-N_1(e^{j\omega}) - D_1(e^{j\omega}) [A_2^2(\omega) + \epsilon(\omega)] - s \leq 0 \quad (*)$$

$$-N_1(e^{j\omega}) + D_1(e^{j\omega}) [A_2^2(\omega) - \epsilon(\omega)] - s \leq 0 \quad (*)$$

solve $(*)$ eq's using linear programming techniques.

if $s = 0$ you have solⁿ
 $s > 0$ you can still make $s = 0$
 $s < 0$ you have to redesign with compromises $A_2^2(\omega)$.