${\rm Lab}~04$ Designing the Controller of the CPU

	Ryan Cruz	Zachary Davis
Category	ryan.cruz25@uga.edu	zachdav@uga.edu
Pre-lab	50	50
In-lab Module & Testbench Design	50	50
In-lab Testbench Sim. & Analysis	50	50
In-lab FPGA Synthesis & Analysis	50	50
Lab Report Writing	50	50

February 8, 2018

Contents

1	Lab Purpose	3
2	Implementation Details	3
	2.1 Part 1	3
	2.2 Part 2	7
	2.3 Part 3	8
3	Experimental Results	11
4	Significance	12
5	Comments/Suggestions	12

1 Lab Purpose

In this lab we design the controller for the CPU, covering all possible states and transitions layed out by the control signal table we completed for prelab (Figure 2.). Like always, all test cases will be tested with a test bench and corresponding waveforms, shown in the Expiremental Results section.

2 Implementation Details

2.1 Part 1

Prelab There were 3 parts to the prelab: Designing the FSM for the controller, the full control signal table, and the corresponding boolean equations for those signals.

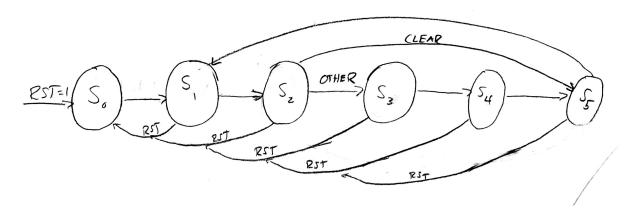


Figure 1: FSM for the controller of the CPU. Created using information from the lecture slides

Transition	LD_D	CL_D	LD_IR	CL_IR	LD_AC	CL_AC	ADD/SUB	R/W	MEM_EN	D/PC	LD_PC	PC_CNT	CL_PC
Any> S0	0	1	0	1	0	1	0	0	0	0	0	0	1
S0> S1	0	0	0	0	0	0	0	0	1	1	0	0	0
S1> S2	0	0	1	0	0	0	0	0	0	1	0	1	0
S2> S5	0	0	0	0	0	1	0	0	0	1	0	0	0
S2> S3	0	0	0	0	0	0	0	0	1	1	0	0	0
S3> S4	1	0	0	0	0	0	0	0	0	1	0	0	0
S4> S5 Add	0	0	0	0	1 - 1	0	0	0	0	1	0	1	0
S4> S5 Sub	0	0	0	0	1	0	7-1	0	0	1	0	1	0
S4> S5 Store	0	0	0	0	0	0	0	1	1	0	0	1	0
S4> S5 BNZ (Z=1)	0	0	0	0	0	0	0	0	0	0	1	0	0
S4> S5 BNZ (Z=0)	0	0	0	0	0	0	0	0	0	0	0	1	0
S5> S1	0	0	0	0	0	0	0	0	01	1	08	0	0

Figure 2: Control Signals. Filled in using states from the lecture slides using one-hot encoding, left any not-on state 0.

```
LD_D = S3

CL_D = RST

LD_IR = S,

CL_IR = RST

LD_A(= Sy . (ADD + SUB)

ADDORSUB = Sy . STORE

R|W = Sy . STORE

D|PC = [RST + Sy . (BNZ + STORE)]'

LD_PC = Sy . BNZ . ZERO

P(-(NT = S1 + Sy . [(APD + SUB) + BNZ . ZERO' + STORE]

CL_PC = RST

MEM_ENC = S1 + Sz + Sy . STORE
```

Figure 3: Boolean Equations for figure 2. control signals

```
//control.v
'timescale 1s/1s
module control(CLK,CLR,RESET,S0,S1,S2,S3,S4,S5);
        input CLK;
        input CLR,RESET;
        output S0,S1,S2,S3,S4,S5;
                     S0, next_S0,
        reg
                         S1, next_S1,
                         S2, next_S2,
                         S3, next_S3,
                         S4, next_S4,
                         S5, next_S5;
                     STATEO, next_STATEO,
        reg
                         STATE1, next_STATE1,
                         STATE2, next_STATE2,
                         STATE3, next_STATE3,
                         STATE4, next_STATE4,
                         STATE5, next_STATE5;
        always @(posedge CLK)
                begin
                         STATEO = next_STATEO;
                         STATE1 = next_STATE1;
                         STATE2 = next_STATE2;
                         STATE3 = next_STATE3;
                         STATE4 = next_STATE4;
                         STATE5 = next_STATE5;
                         S0 = next_S0;
                         S1 = next_S1;
                         S2 = next_S2;
                         S3 = next_S3;
                         S4 = next_S4;
                         S5 = next_S5;
                 end
        always @ (CLR or RESET or STATEO or STATE1 or STATE2 or STATE3
         or STATE4 or STATE5)
                 begin
                         if ( RESET )
                                 next_STATE0=1;
                         else
                                 next_STATE0=0;
```

```
if ( "RESET & STATEO | "RESET & STATE5 )
        next_STATE1=1;
else
        next_STATE1=0;
if ( ~RESET & STATE1 )
        next_STATE2=1;
else
        next_STATE2=0;
if ( "RESET & "CLR & STATE2 )
        next_STATE3=1;
else
        next_STATE3=0;
if ( "RESET & STATE3 )
        next_STATE4=1;
else
        next_STATE4=0;
if ( "RESET & CLR & STATE2 | "RESET & STATE4 )
        next_STATE5=1;
else
        next_STATE5=0;
if ( RESET )
        next_S0=1;
else
        next_S0=0;
if ( "RESET & STATEO | "RESET & STATE5 )
        next_S1=1;
else
        next_S1=0;
if ( ~RESET & STATE1 )
        next_S2=1;
else
        next_S2=0;
if ( "RESET & "CLR & STATE2 )
        next_S3=1;
else
        next_S3=0;
if ( "RESET & STATE3 )
        next_S4=1;
```

Runs through every case assigning states for each possible transition using simple if/else statements. Named each state and transition using variables in verilog.

2.2 Part 2

1-Hot Encoding For this part of the lab we needed to determine how we were going to encode the instruction register and the lab report recommended 1-hot encoding. We have five instructions to encode and an eight bit instruction register so rather than count up by ones we count up in powers of two. This reduces the number of labels sufficently, but we only need five and with counting by powers of two there are eight combinations. This is far easier to construct in a schematic and far more human readible.

Instruction	Encoded Value (Decimal)	Encoded Value (Binary)
ADD	1	00000001
SUB	2	00000010
CLR	4	00000100
BNZ	8	00001000
STR	16	00010000
	ADD SUB CLR BNZ	SUB 2 CLR 4 BNZ 8

2.3 Part 3

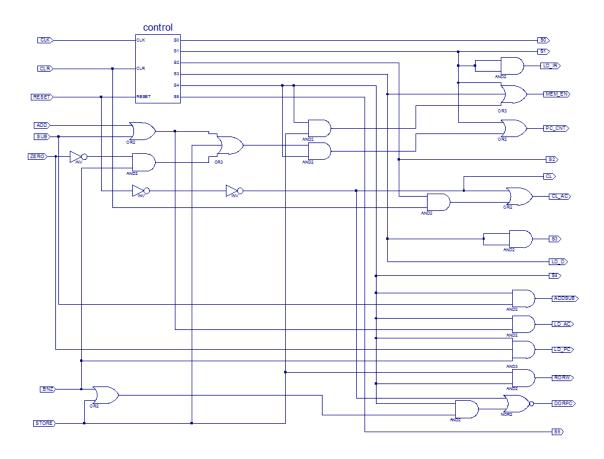


Figure 4: Shematic for the controller using a symbol for the control module we created in part 1.

```
//controller_tb.v
'timescale 1ns/1ps

module controller_tbw_tb_0;

reg ADD = 1'b0;
reg BNZ = 1'b0;
reg CLK = 1'b0;
reg CLR = 1'b0;
reg RESET = 1'b0;
reg RESET = 1'b0;
reg STORE = 1'b0;
reg SUB = 1'b0;
reg ZERO = 1'b0;
```

```
wire DORPC;
wire LD_AC;
wire LD_D;
wire LD_IR;
wire LD_PC;
wire MEM_EN;
wire PC_CNT;
wire RORW;
wire SO;
wire S1;
wire S2;
wire S3;
wire S4;
wire S5;
initial // Clock process for CLK
        begin
                forever
                begin
                         CLK = 1'b0;
                         #100;
                         CLK = 1'b1;
                         #100;
                end
        end
controller_sch UUT(
        .ADD(ADD),
        .BNZ(BNZ),
        .CLK(CLK),
        .CLR(CLR),
        .RESET(RESET),
        .STORE(STORE),
        .SUB(SUB),
        .ZERO(ZERO),
        .ADDSUB(ADDSUB),
        .CL(CL),
        .CL_AC(CL_AC),
        .DORPC(DORPC),
        .LD_AC(LD_AC),
        .LD_D(LD_D),
        .LD_IR(LD_IR),
        .LD_PC(LD_PC),
        .MEM_EN(MEM_EN),
        .PC_CNT(PC_CNT),
        .RORW(RORW),
        .SO(SO),
```

```
.S1(S1),
     .S2(S2),
     .S3(S3),
     .S4(S4),
     .S5(S5));
initial
     begin
          // ----- Current Time: 85ns
          #85;
          RESET = 1'b1;
          // -----
          // ----- Current Time: 285ns
          #200;
          RESET = 1'b0;
          // -----
          // ----- Current Time: 485ns
          #200;
          CLR = 1'b1;
          // -----
          // ----- Current Time: 1085ns
          #600;
          ADD = 1'b1;
          CLR = 1'b0;
          // -----
          // ----- Current Time: 2085ns
          #1000;
          ADD = 1'b0;
          SUB = 1'b1;
          // -----
          // ----- Current Time: 3085ns
          #1000;
          STORE = 1'b1;
          SUB = 1'b0;
          // -----
          // ----- Current Time: 4085ns
          #1000;
          BNZ = 1'b1;
          STORE = 1'b0;
          // -----
```

3 Experimental Results

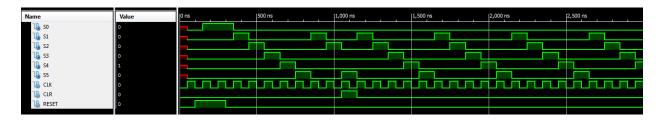


Figure 5: Test bench for the control module. The uniform cascading waves show the simple transitions from 0-5, and at the bottom you can see that when clear is active, S5 occurs, and when reset is active, S0 occurs.

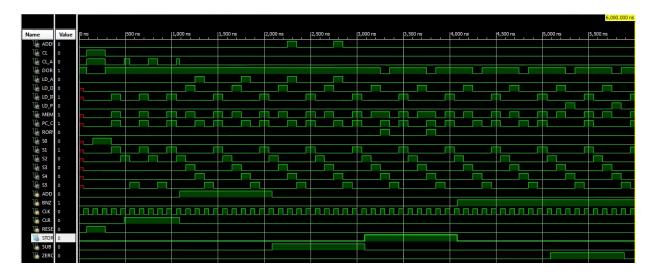


Figure 6: Test bench for the controller schematic. Cycles clock every 100 ns and changes all states to cover all cases. The inputs line the bottom, and the outputs exist on top. The cascading S0-S5 waves that look similar to the first test bench are exactly that, more direct state transitions. Then you can see things like CL_A activating when CLR and S2 are 1, CL on RESET, RORW on STORE and S4, and the rest of the things that occur in the table.

4 Significance

Another piece of the full CPU puzzle, the controller allows us to take in any of the control signals and provide the correct instructions to the CPU. With this controller completed, we can move on to connect this portion to the datapath to complete and assemble the toy processor, which will be able to run a given program.

5 Comments/Suggestions