

# Parallel Algorithms

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# Contents

- Parallel Programming: Introduction
- Parallel application design:
  - Parallel programming paradigms
  - Parallel algorithms

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# Introduction

- Performance is the main reason of parallel/distributed processing.
- When a user develops a parallel application, he/she expects reaching certain performance improvements.
- These improvements allow reducing execution time, increasing the problem size, maximizing the throughput, etc.

# Introduction

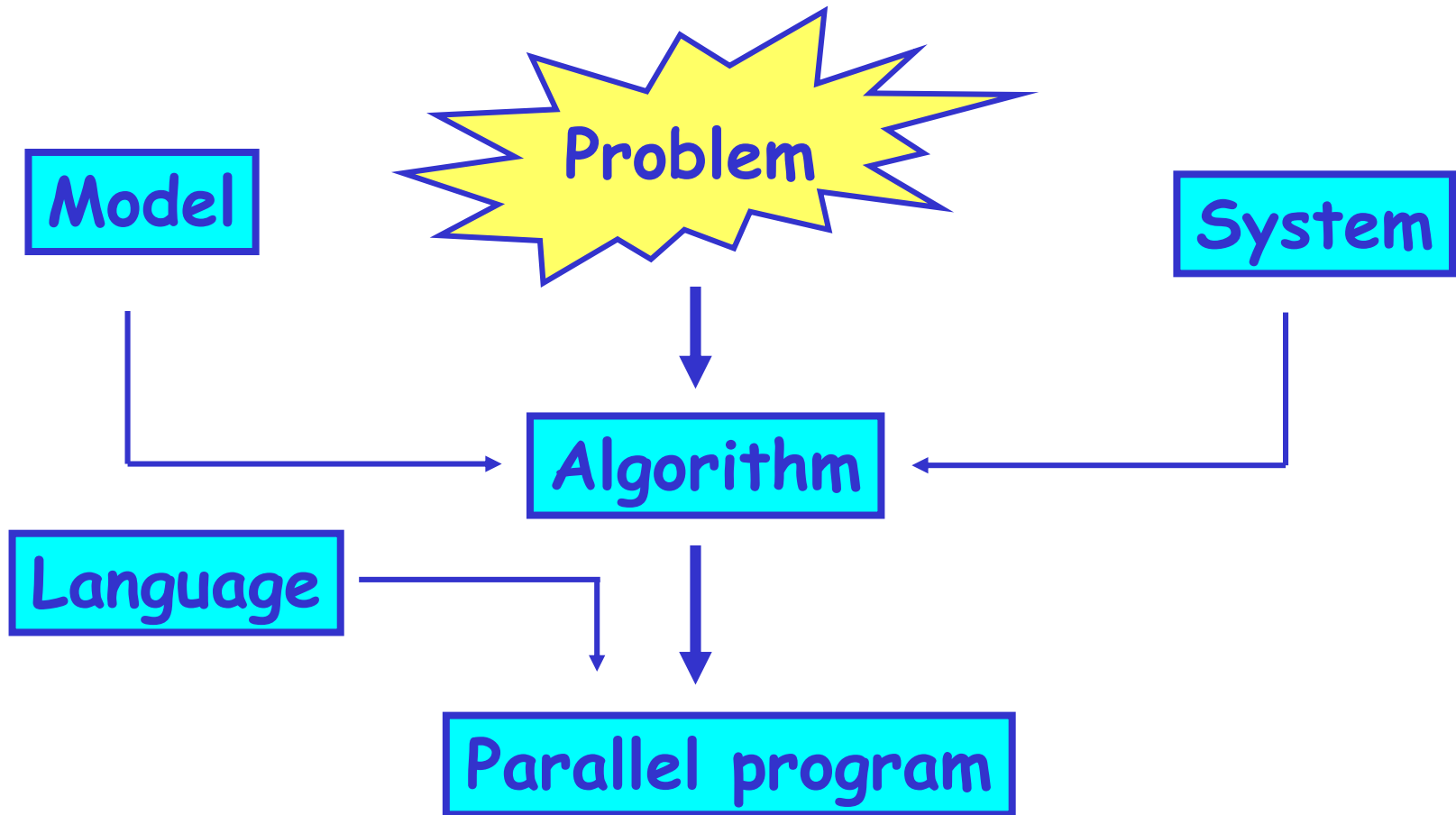
- Parallel Application design is strongly connected to the system where the application will run and to the programming model used.
- Application, system and programming model are closely related.

# Parallel programming models

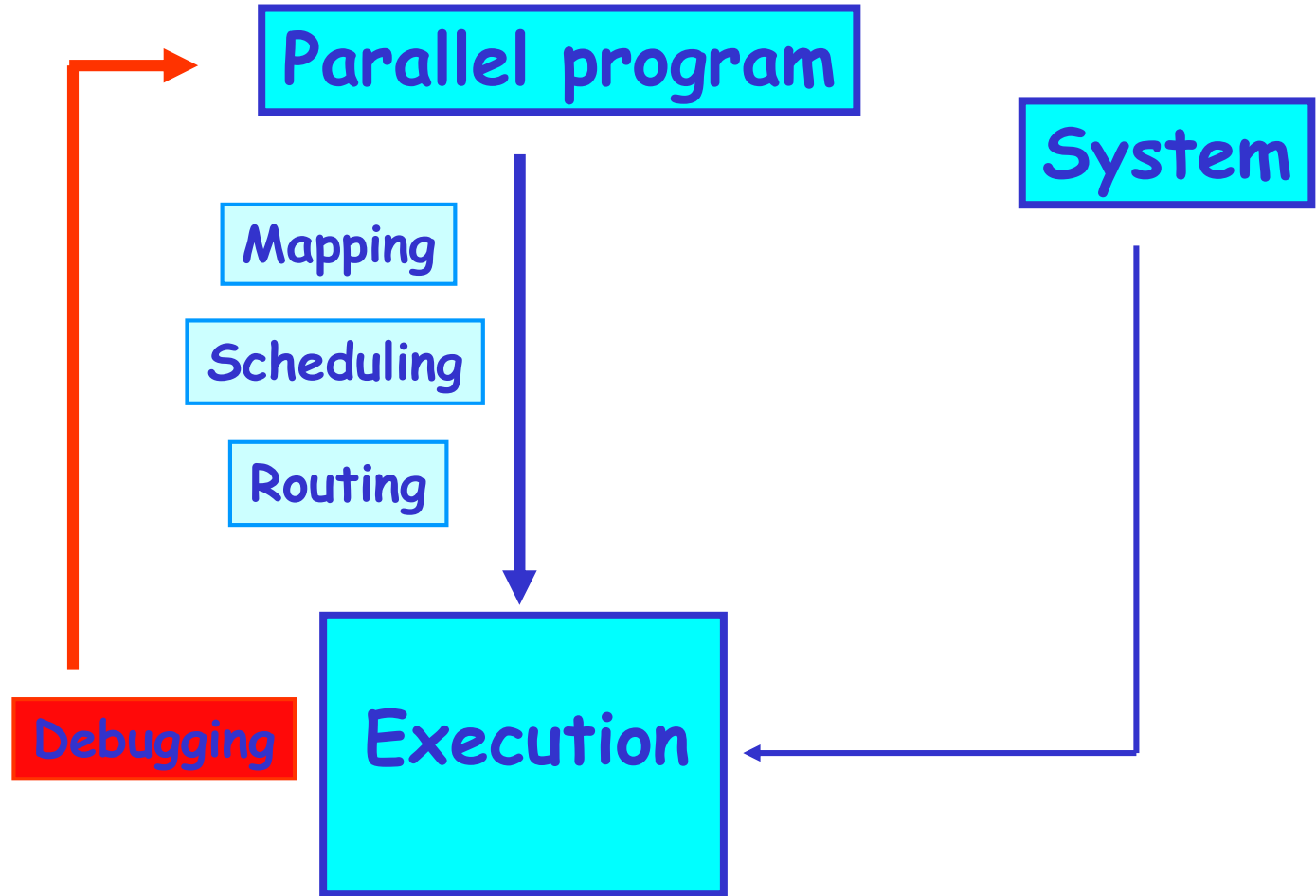
## > Programming models:

- Automatic parallelism extraction
- Message passing
- Shared memory: "Threads"

# Parallel programming life cycle

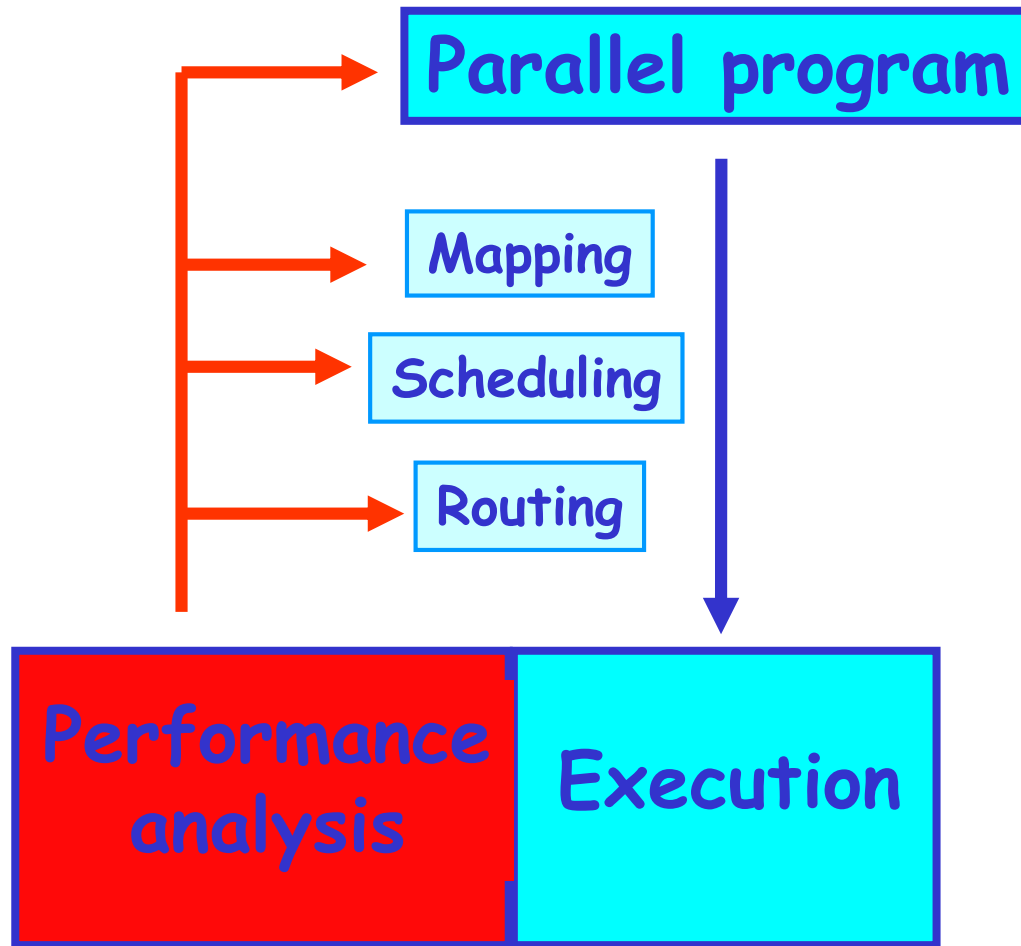


# Parallel programming life cycle





# Parallel programming life cycle



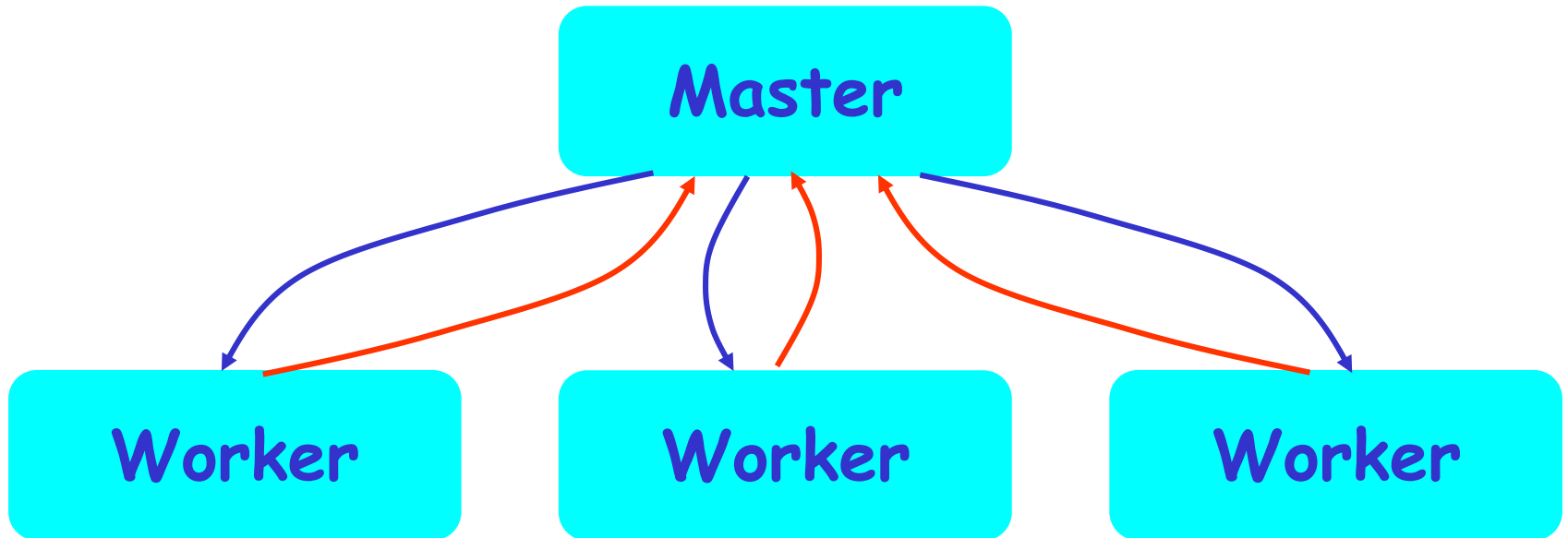
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# Programming paradigms

- There are many algorithms in the literature that can be used to solve many different problems.
- Most algorithms have been designed for a particular system architecture.
- However, there are certain basic programming paradigms that can be used to solve many problems.

# Master/Worker



# Master/Worker

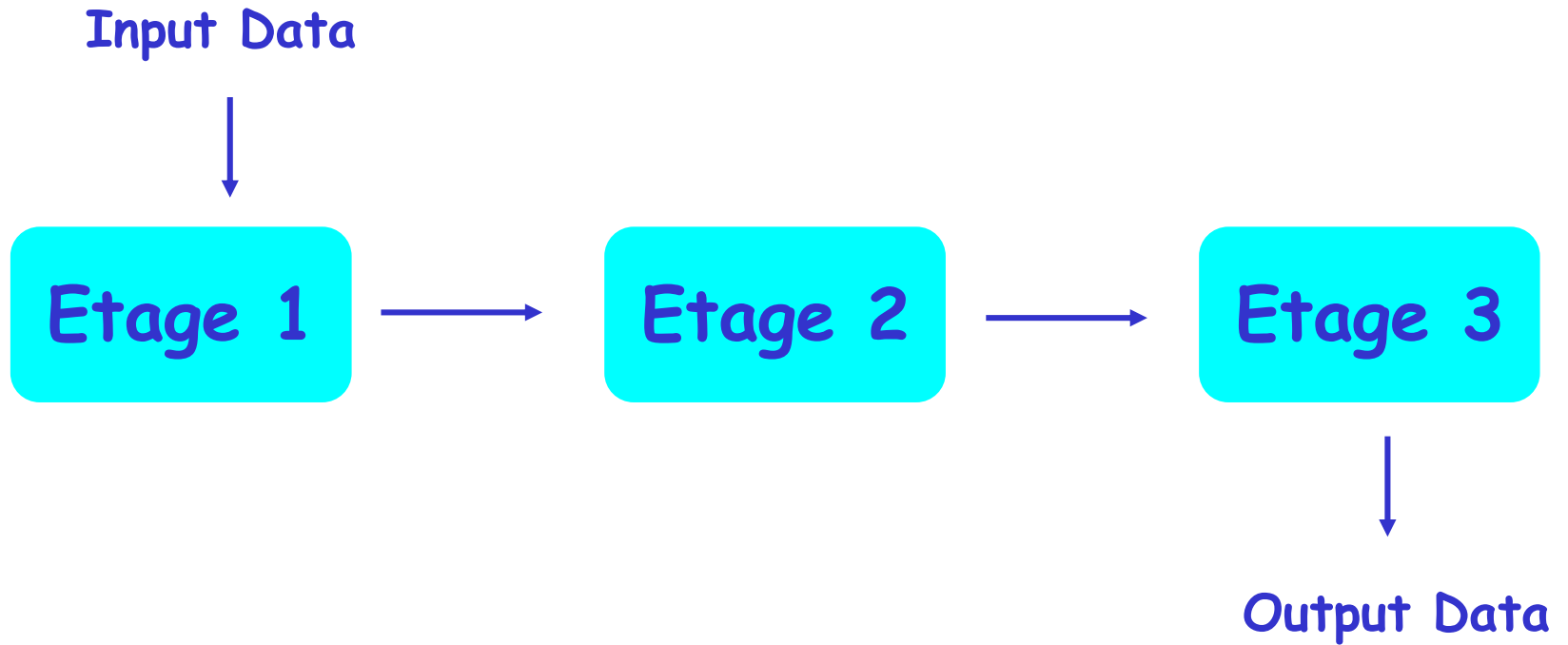
## ➤ Iteration dependence

- Dependence between iterations: Master process requires the results of “all” the workers” to generate the next set of data.
- Independent data input: Data continuously arrives to the Master process and it does not require the results of the previous set of data.

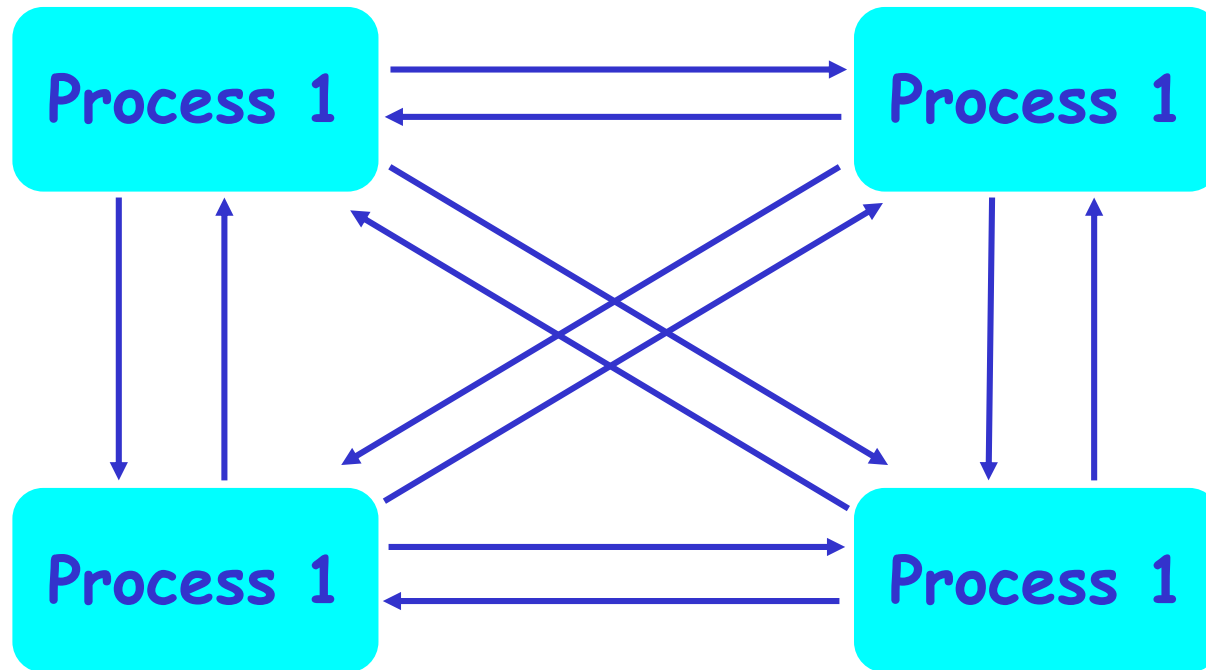
# Master/Worker

- Data partitioning and distribution
  - Distribute all data among the Workers:  
When Master has prepared all the data, it sends them to the Worker processes according to certain strategy.
  - Distribute data on-demand: Master process distributes a subset of data to the Workers and then these workers ask for more data when they finish the previous chunk.

# Pipeline

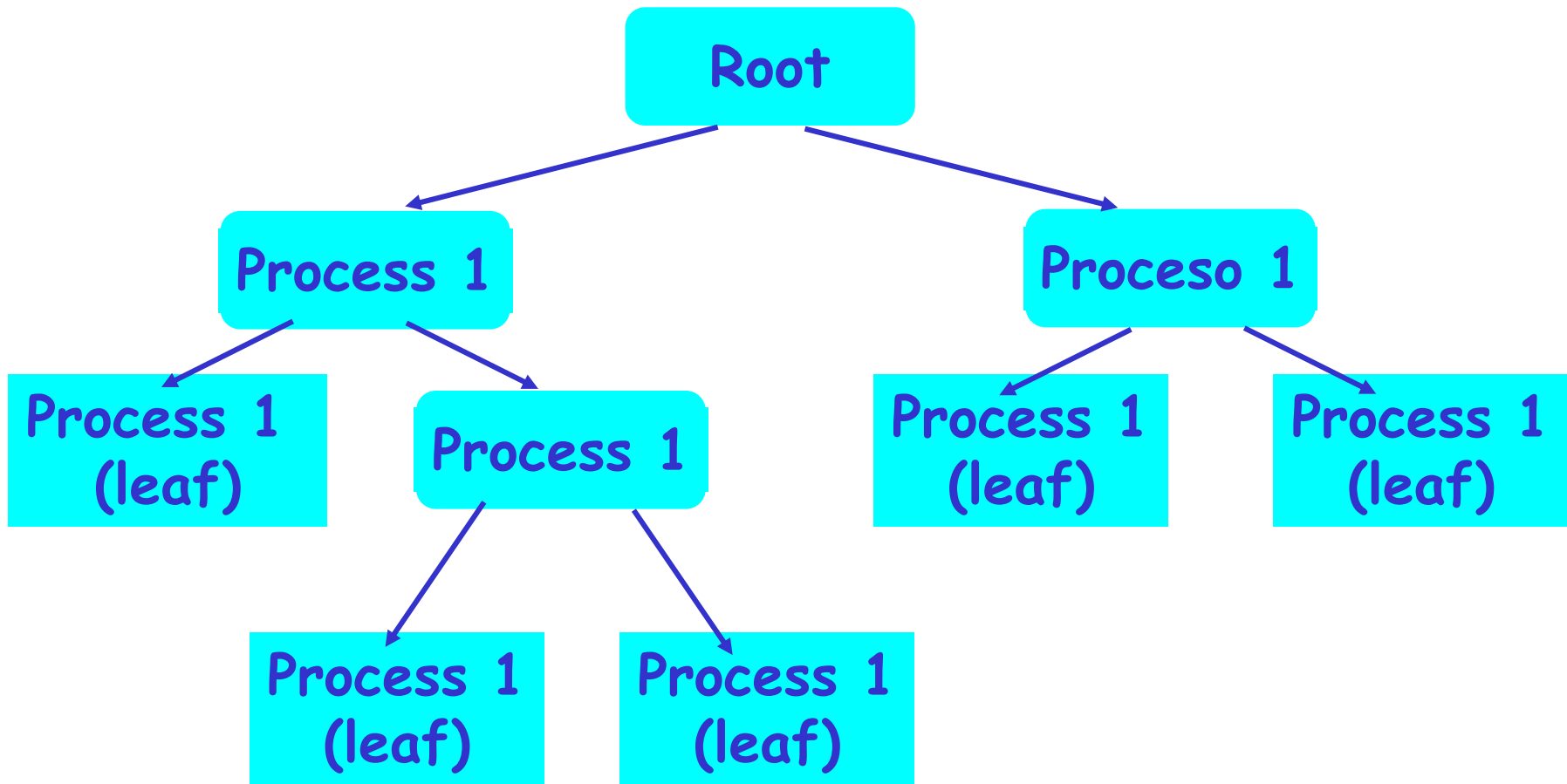


# Single Program Multiple Data SPMD





# Divide and Conquer



# Divide and Conquer

- All processes, but the root, are exactly equal. Each process:
  - Receive a subset of data.
  - If it can process them, it process them.
  - If it cannot process them, it creates a certain number of child processes and distribute the subset of data.

# Programming paradigms

- Usually, these paradigms simplify the programmers task.
- The programmer just develops one copy of the code and, when it is executed, multiple instances are generated.
- Communication patterns are always the same for each paradigm.

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# Matrix multiplication

|          |          |          |          |
|----------|----------|----------|----------|
| $A_{11}$ | $A_{12}$ | $A_{13}$ | $A_{14}$ |
|          |          |          |          |
|          |          |          |          |
|          |          |          |          |

|          |  |  |  |
|----------|--|--|--|
| $B_{11}$ |  |  |  |
| $B_{21}$ |  |  |  |
| $B_{31}$ |  |  |  |
| $B_{41}$ |  |  |  |

|          |  |  |  |
|----------|--|--|--|
| $C_{11}$ |  |  |  |
|          |  |  |  |
|          |  |  |  |
|          |  |  |  |


$$C_{ij} = \sum_{\forall k} A_{ik} * B_{kj}$$

# Matrix multiplication

- There is no dependence among the calculations required to determine the values of the resulting matrix
- All  $C_{ij}$  terms can be calculated in parallel.
- The set of multiplications that must be done to evaluate one term of the resulting matrix are also independent. Therefore, each  $C_{ij}$  calculation can be parallelised.

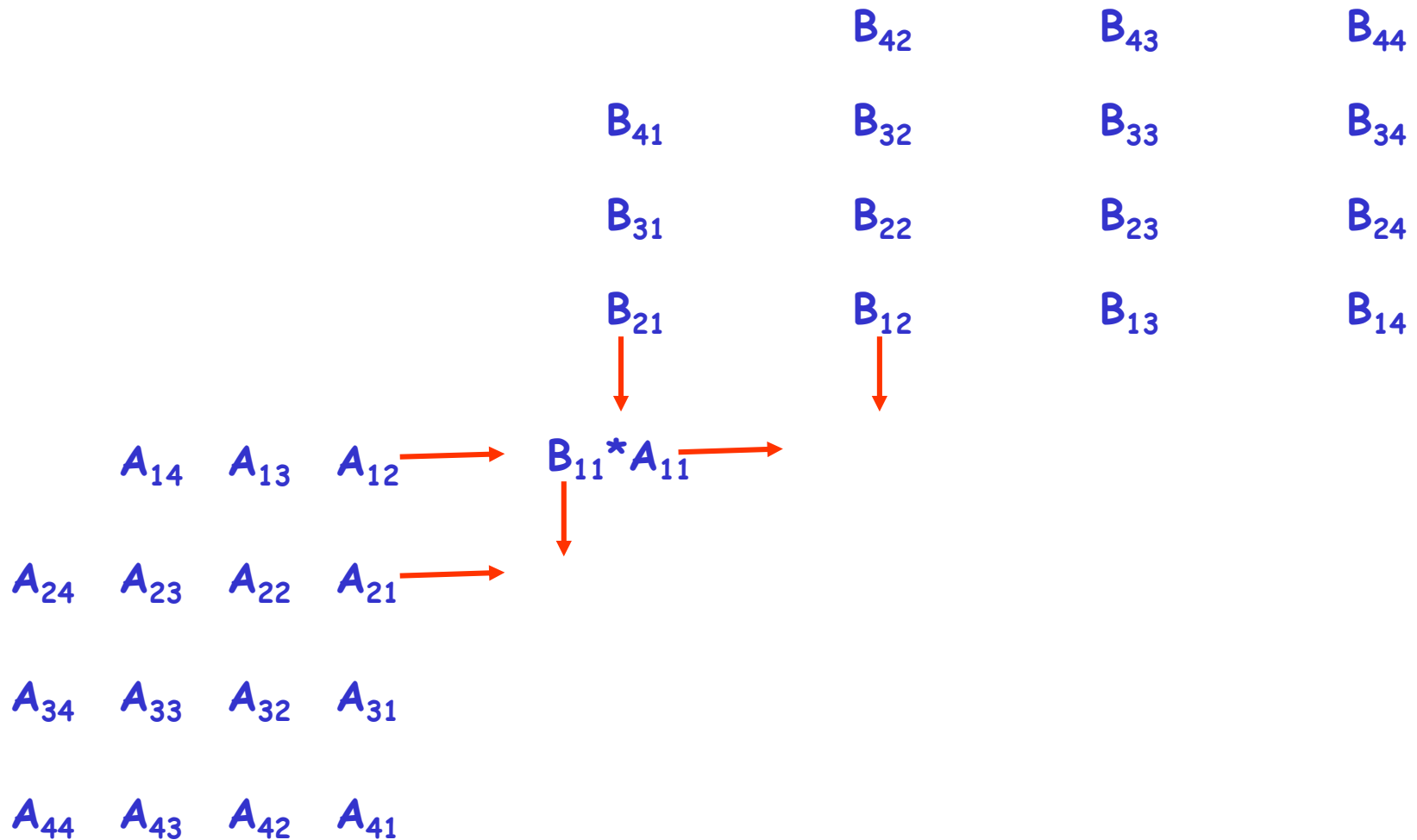
# Matrix multiplication

|          |          |          |          |
|----------|----------|----------|----------|
| $B_{41}$ | $B_{42}$ | $B_{43}$ | $B_{44}$ |
| $B_{31}$ | $B_{32}$ | $B_{33}$ | $B_{34}$ |
| $B_{21}$ | $B_{22}$ | $B_{23}$ | $B_{24}$ |
| $B_{11}$ | $B_{12}$ | $B_{13}$ | $B_{14}$ |



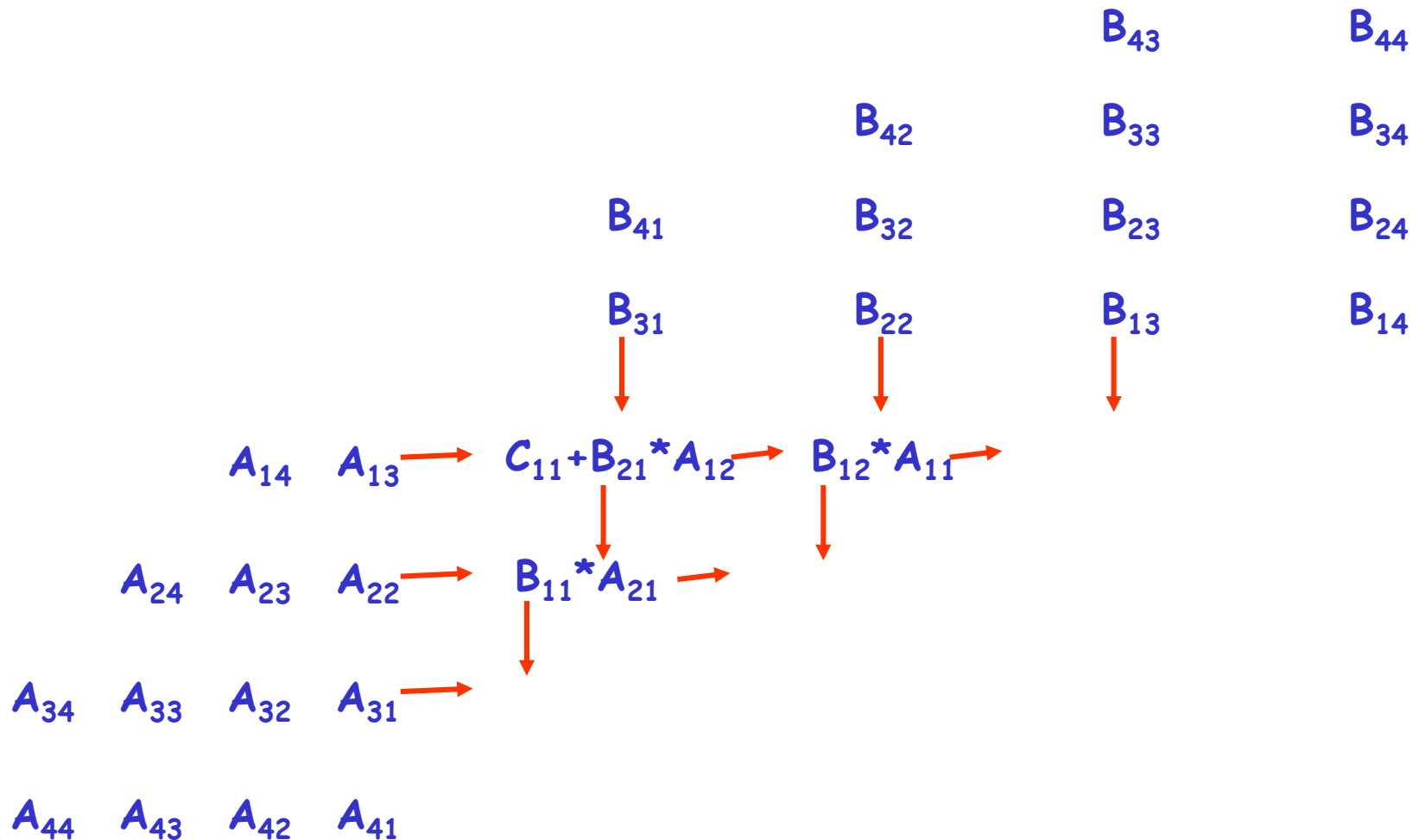
|          |          |          |          |   |
|----------|----------|----------|----------|---|
| $A_{14}$ | $A_{13}$ | $A_{12}$ | $A_{11}$ |  |
| $A_{24}$ | $A_{23}$ | $A_{22}$ | $A_{21}$ |   |
| $A_{34}$ | $A_{33}$ | $A_{32}$ | $A_{31}$ |   |
| $A_{44}$ | $A_{43}$ | $A_{42}$ | $A_{41}$ |   |

# Matrix multiplication

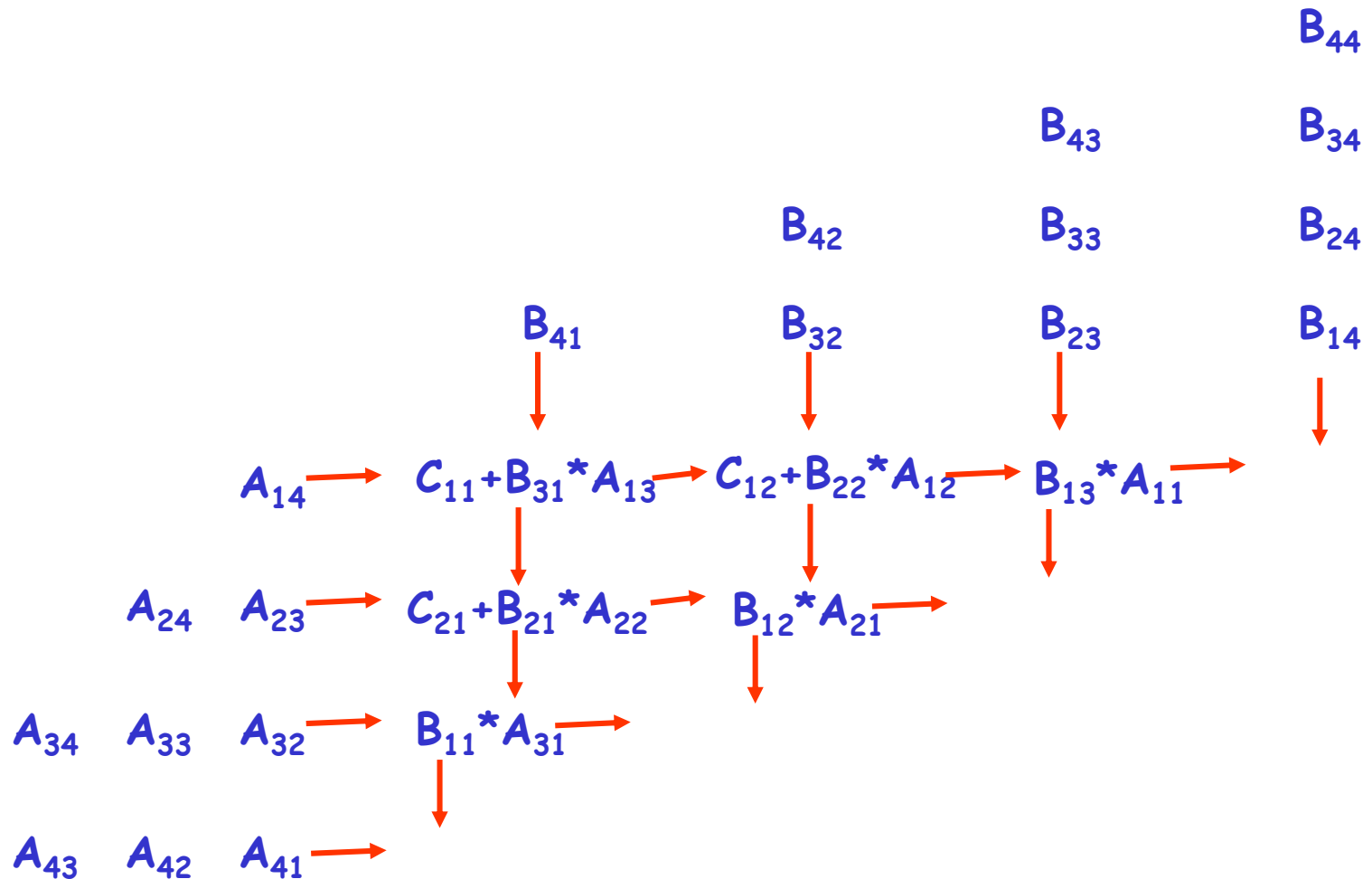




# Matrix multiplication



# Matrix multiplication



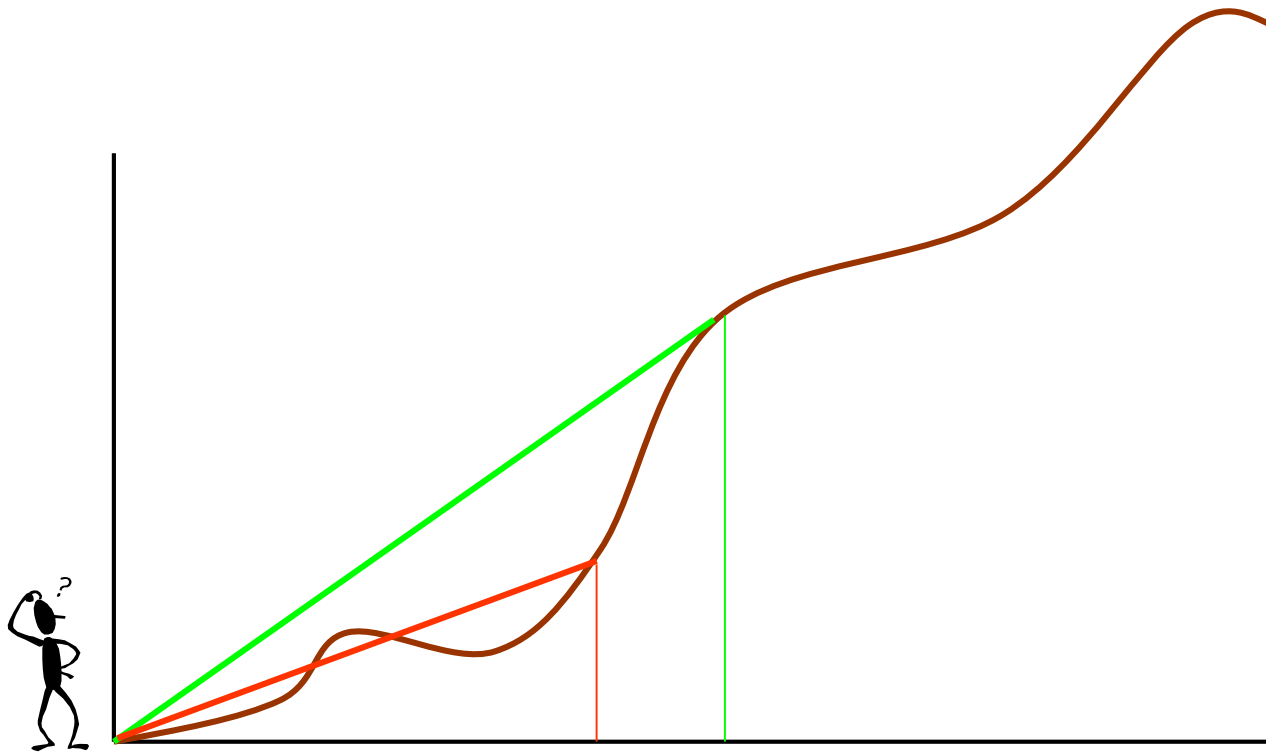
# Thinkings about the algorithm

- The algorithm works properly on a system with a mesh interconnectin network.
- On other systems (PC clusters on ethernet) the performance degrades substancialy.

# Other considerations

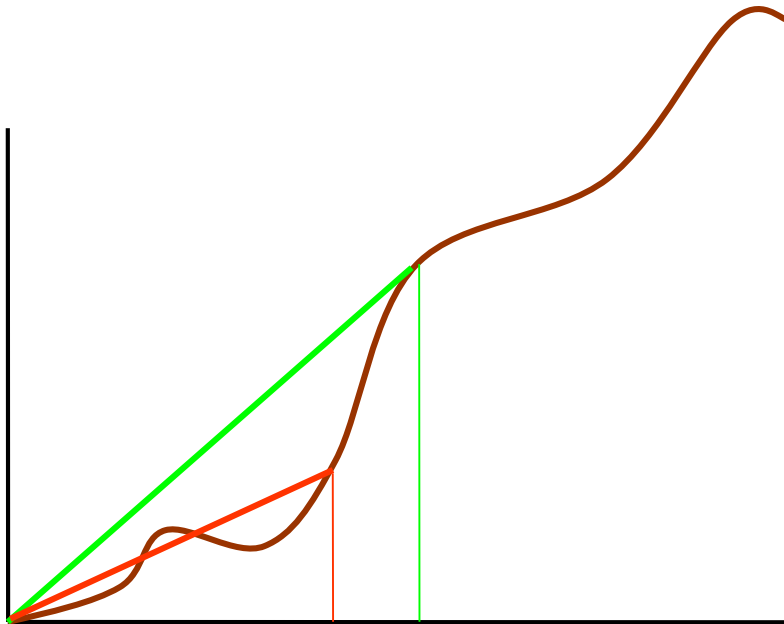
- Besides the topology of the interconnection network, the computation and communication capabilities will also affect the performance of the algorithm.
- It is necessary to determine the granularity of the tasks/data to balance computation and communication.

# Which points are visible?



# Which points are visible?

$$\alpha_i = \arctg \frac{y_i}{x_i}$$



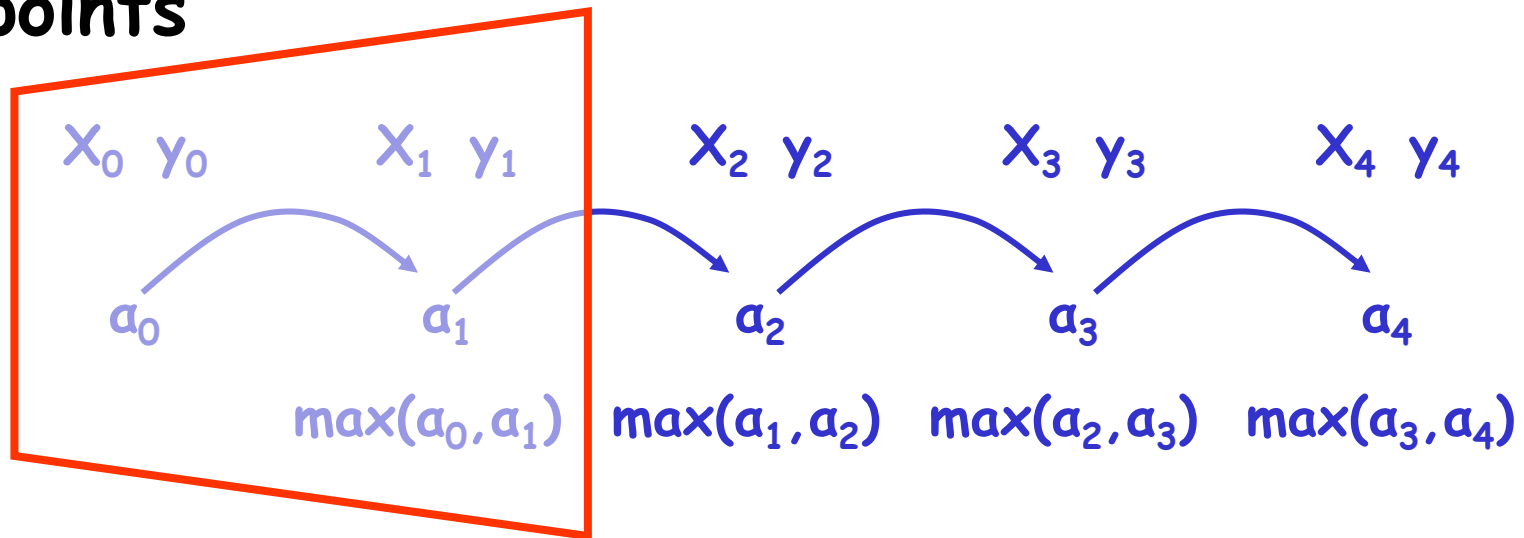
Visibility condition

$$\alpha_i \geq \alpha_k \forall k < i$$

$$\max(\alpha_k)_{k=0}^i = \alpha_i$$

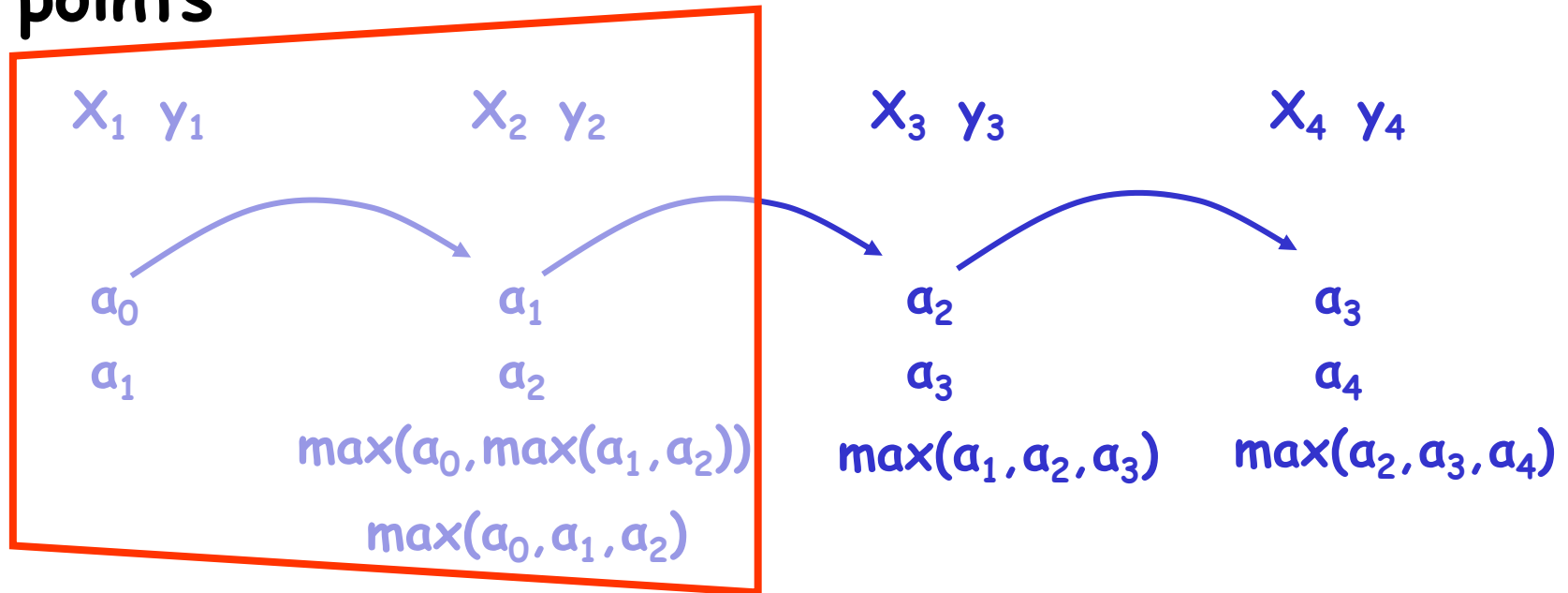
# A parallel algorithm

- To determine if one point is visible, it is necessary to know the angle of all previous points



# A parallel algorithm

- > To determine if one point is visible, it is necessary to know the angle of all previous points





# A parallel algorithm

- › To determine if one point is visible, it is necessary to know the angle of all previous points

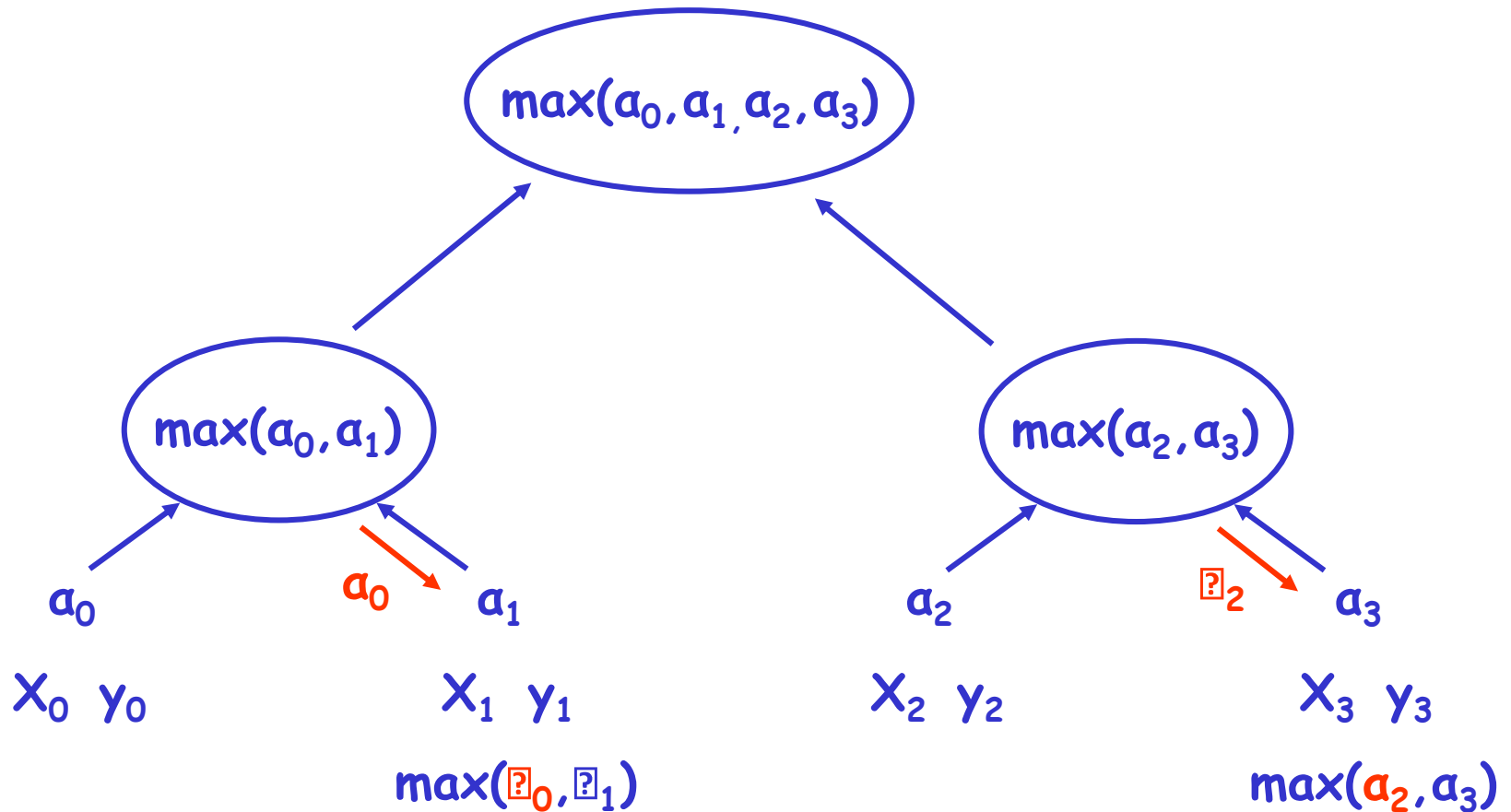
$$x_i \ y_i$$

$$\max(a_0, a_1, \dots, a_i)$$

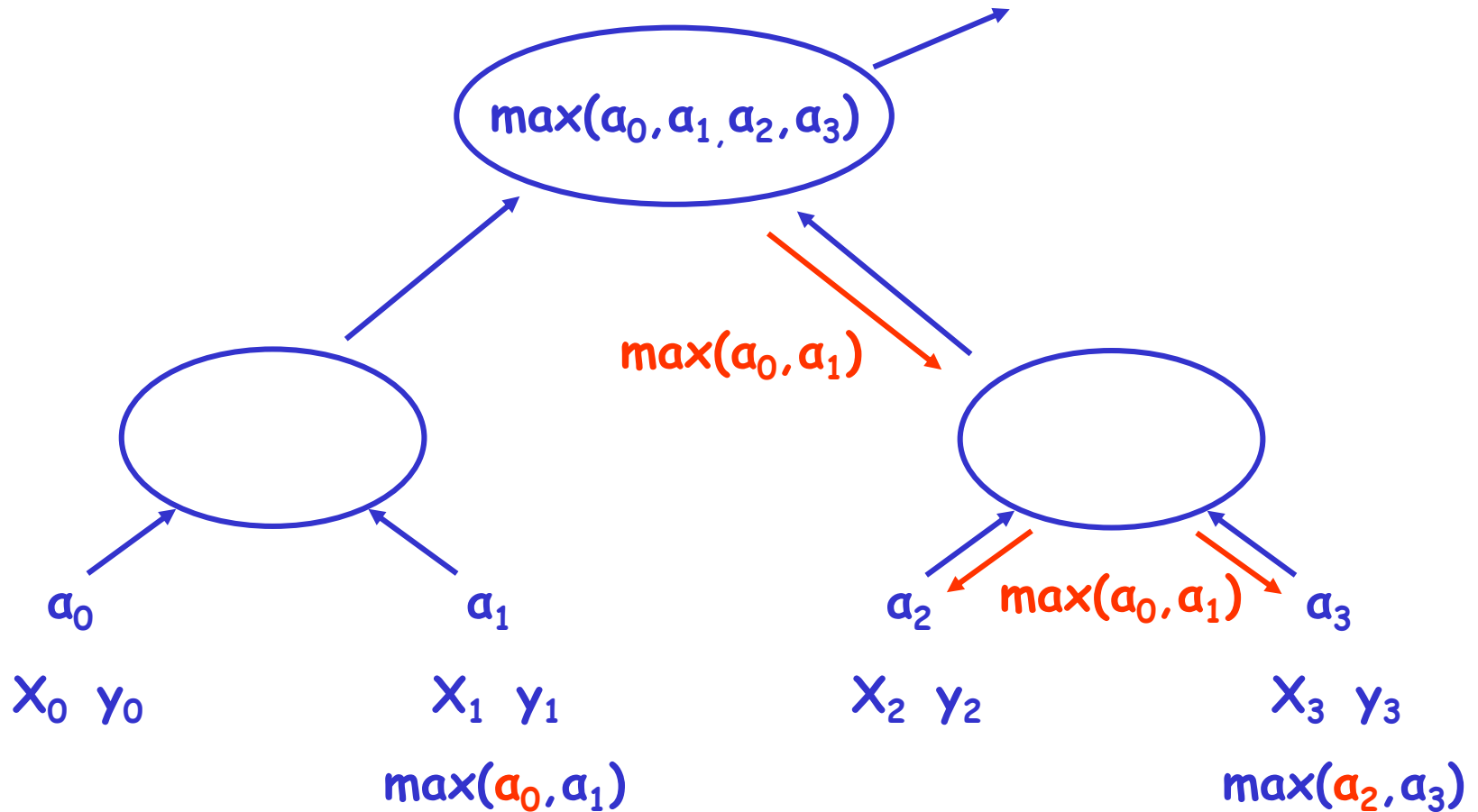
- › If there  $n$  points, we need  $(n-1)$  steps.

$$O(n)$$

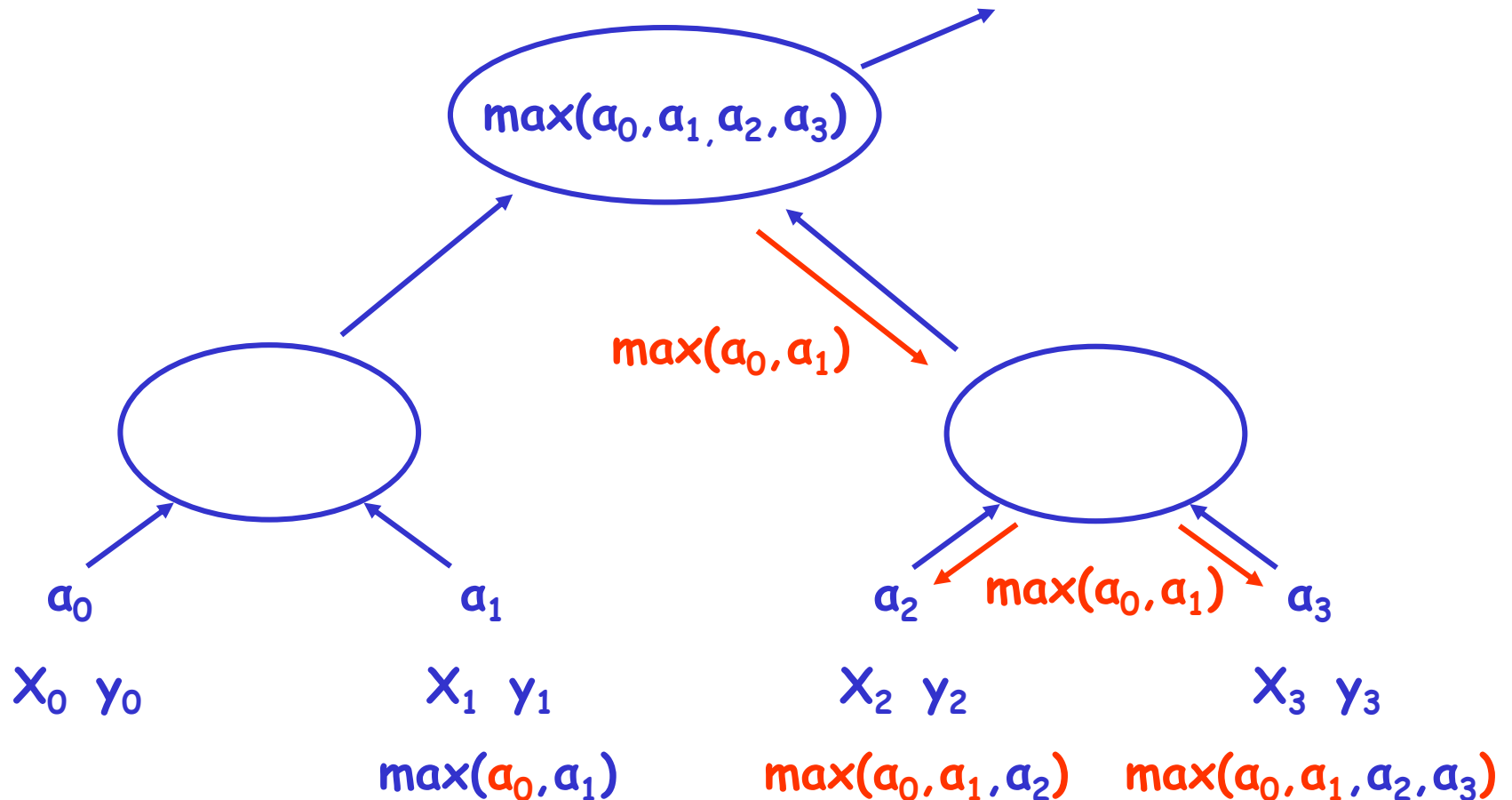
# Parallel Prefix algorithm



# Parallel Prefix algorithm



# Parallel Prefix algorithm



# Parallel Prefix algorithm

- The complexity of this algorithm is logarithmic.

$$O(\log_2 n)$$

- When the number of points increases, the improvement in the execution time is significant.