

# A 2D Nearest-Neighbor Quantum Architecture for Factoring

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March 25, 2013

This document responds to comments by Referee 2, which were received on November 30, 2012. These comments are quoted and responded to below. We thank the Referee for the constructive feedback and suggestions.

## 1 General Comments

The paper "A 2D Nearest-Neighbor Quantum Architecture for Factoring", by Pham and Svore, contains new ideas and represents a step forward in our understanding of how to implement arithmetic on a quantum computer. The authors do two things that separate their paper from previous analyses of Shor's algorithm: they consider a 2D network, and they choose to optimize depth rather than width. While it is still too early to say which physical constraints on quantum computers will be most restrictive, the authors explore an important new part of the space of algorithms. This paper is appropriate for publication in QIC.

There are a few issues that should be addressed before publication.

The introduction does not state the main result. At the very least, there should be a statement of the asymptotic behavior of the circuit.

It is even possible that most (or all) of Section 8 should be moved earlier, including figure 11. The reader should not have to read to the end to find out the punch line.

*We have reworded the introduction to reflect the main contribution of our work.*

The introduction to Section 4 is confusing. Is the pair  $(u, v)$  a CSE number only when it arises from this construction? Is  $u_{n-1} = 0$  a convention within the paper or part of the definition? The example in Figure 2 might suggest that  $u_i = v_i = 1$  is not permitted, when in fact it is. The authors need to clarify standard definitions versus their conventions. A reference to pre-quantum literature on carry-save addition would help.

*The sum of any two numbers  $(u+v)$  can be interpreted as a CSE number. The point of the definition  $a+b+c = u+v$  is to associate a given  $u+v$  with a particular  $a+b+c$ , but this association is non-unique. However, given the truth table for parity ( $u$ ) and majority ( $v$ ), this is a one-to-one mapping between  $a, b, c$  and  $u, v$ . By that truth table,  $u_{n-1} = 0$ .*

*In the related work section, we have included a reference to Wallace Trees and the 3-2 adder. These are the classical CSA techniques.*

The authors imply that the final output of their exponentiation circuit is left in CSE form, which is not a unique representation of the answer. This could mess up the next step of Shor's algorithm, in which states with the same answer collapse. The authors need to either (a) explain why this is not a problem, or (b) explicitly say that they convert the answer to standard form. (Note that, at worst, this conversion can be done in log depth, so the asymptotic analysis of the algorithm is not affected.)

*Thank you for catching this mistake. Resources have been added at the end to upper bound the conversion of the final CSE output to a conventional number.*

## 2 Minor comments

Page 1: "best-known" should be "best known". ("best-known" means "most famous", which is not what is intended.)

*Corrected.*

Top of Section 2.1: Should be "following Van Meter and Itoh [Van Meter and Itoh (2005)]" or "following [Van Meter and Itoh (2005)]".

*Corrected.*

Page 2: “where each qubit has four neighbors” is misleading since it implies a torus rather than a bounded planar region, and since the authors are about to change four to six. Simply “where there is an extra...” would suffice.

*Corrected to “planar” neighbors.*

Last sentence of Section 2: “there is no known way”. Do the authors mean (a) no one has found a way, (b) there is provably no way, or (c) there is provably no way for a particular family of circuits? The at-large citation to Rosenbaum is unhelpful. One way to clarify would be to cite a specific result from Rosenbaum.

*This sentence has been removed. Recent private communication indicates a constructive way to unentangle the source qubit.*

Bottom of page 6: “The circuit operations out-of-place and produces two garbage qubits”. No. An out-of-place circuit leaves its input intact; this circuit overwrites  $a_i$ , so it operates in place. Calling  $b_i$  and  $c_i$  “garbage” seems not quite right, since if they were not present the circuit would not be reversible. One could design an out-of-place version, but in this case cleaning up  $b_i$  and  $c_i$  would be straightforward and there would be no garbage.

*Sentence has been corrected.*

Page 7: The phrase “n-bit” appears twice in one sentence, once as an adjective and once as a noun, meaning two different things. This is not technically ambiguous, but it is awkward.

*Corrected.*

Page 7, before (6): The term “significance” has not been defined and seems not to be used elsewhere. The sentence could be rewritten.

*We added the definition of significance in Section 4.*

Page 7, proof of lemma 1: extraneous “.” in “Our”.

*Corrected.*

Pages 8 and 9: The argument that  $v'''_{n+2} = 0$  is convoluted, working backward from the conclusion to a true statement (and misusing the word “implies” in the process), and contains at least one mistake. It would be cleaner to work forward. For example: since  $u'_{(n)} + v'_{(n+1)} = u_{(n)} + v_{(n)} \leq 2^{n+1}$ , the bits  $u'_n$  and  $v'_{n+1}$  cannot both be 1. But  $u''_{n+1} = v'_{n+1}$  and  $v''_{n+1} = u'_n v'_n$ , so  $u''_{n+1}$  and  $v''_{n+1}$  cannot both be 1, and hence  $v'''_{n+2} = 0$ .

*We thank the referee for his shorter proof and for pointing out the mistake. We have incorporated this change into the paper.*