CAP Theoram [Brewer's Theoram]

C= Consistency A = Availability All nodes see same data all the him Each request reueves a response (recent or not) l = Partitition Tolerance. > system continues to operate despite n/w partitions of conno factures by nodes Mongoob rpbms for a distributed Velti ia Redis Apache NBASE System, It is surpossible to Cassandrabb DynamoBB achieve all 3(c Ap) Couch OB simultaneousy. During Paritition Polerance LNetwork falure], either Consistency on Availability can be provide by a system (DB), ACID

Atomicily Consistency & solution Durabily

properly & ensure reliable fransaction processing

Txno

Atomicchy: All or Nothing of any part of ten fails the entine pransachon is solled back.

- Aco. 18 g no changes are applied.

80 (Not updated) -> holled back 4× n Account A 100 -> sent 100 Consistency: bottoms rules brings des prom one valid state eg: - ve salance (Rule Violation)
Sum of total Lemains same Frolation: Concurrent Txus are isolated RCA) (RCB) (RCA). Costome withdrawing from same ashome withdrawing acount, isolation 11 sehedule ensures that tothe Selval

Selval

Shedule

Schedule

Lost is always consistent

Durability:

All changes in DR should be

permanent even of system

permanent pail

Neplication

SQL] In don't se e eachother BASE Properties

Basically Available
Soft State Eventually consistent

Basially Available:

Stem - soil out not updated on welsite during high traffic

System Granulees availability. DR will respond to all regrests even of data sent is stale on i vico implete

estations

1x DD

Py Nom

Rephia

2min eventral

update

he that I mainte

Even wo Enpert from outside, the state of system charges. Due to the eventual updates across distributed system

Eventvally: System will eventvally consistent become consistent given enough fine. for updalt it well make huis (replicas) But eventually become consistent