Discrete Mathematics #3

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8.3

Problem 29

$$\begin{split} &f(n) = af(\frac{n}{b}) + cn^d \\ &f(\frac{n}{b}) = af(\frac{n}{b^2}) + c(\frac{n}{b})^d \\ &f(n) = a(af(\frac{n}{b^2}) + c(\frac{n}{b})^d) + cn^d \\ &= a^2f(\frac{n}{b^2}) + ac(\frac{n}{b})^d + cn^d \\ &f(\frac{n}{b^2}) = af(\frac{n}{b^3}) + c(\frac{n}{b^2})^d \\ &f(n) = a^3f(\frac{n}{b^3}) + a^2c(\frac{n}{b^2})^d + ac(\frac{n}{b})^d + cn^d \\ & \therefore f(n) = a^kf(1) + c\sum_{i=0}^{k-1}a^i(\frac{n}{b^i})^d \\ &a = b^d, a^i(\frac{n}{b^i})^d = n^d \\ &f(n) = a^kf(1) + cn^d\sum_{i=0}^{k-1}1 \\ &= a^kf(1) + cn^dk \\ &n = b^k, k = \log_b n \\ &a^k = a^{\log_b n} = n^{\log_b a}, \log_b a = d \end{split}$$

 $f(n) = n^d f(1) + cn^d \log_b n$

$$f(n) = n^d f(1) + c n^d \log_b n$$

$$n^d f(1) = O(n^d)$$

$$cn^d \log_b n = cn^d \tfrac{\log n}{\log b} = \tfrac{c}{\log b} n^d \log n = O(n^d \log n)$$

 n^d 보다 $n^d \log n$ 이 더 빠르게 증가하므로

$$\therefore f(n) = O(n^d \log n)$$

Problem 29와 마찬가지로

$$\begin{split} &f(n) = a^k f(1) + c \sum_{i=0}^{k-1} a^i (\frac{n}{b^i})^d \\ &f(n) = a^k f(1) + c n^d \sum_{i=0}^{k-1} (ab^{-d})^i \\ &= a^k f(1) + c n^d \frac{1 - (ab^{-d})^k}{1 - ab^{-d}} \\ &n = b^k, k = \log_b n \\ &f(n) = a^k f(1) + c n^d \frac{1 - (ab^{-d})^{\log_b n}}{1 - ab^{-d}} \\ &= n^{\log_b a} f(1) + \frac{c n^d}{1 - ab^{-d}} - \frac{c n^d (ab^{-d})^{\log_b n}}{1 - ab^{-d}} \\ &= n^{\log_b a} f(1) + n^d \frac{cb^d}{b^d - a} + \frac{cb^d n^d (ab^{-d})^{\log_b n}}{a - b^d} \\ &= n^{\log_b a} f(1) + n^d \frac{cb^d}{b^d - a} + \frac{cb^d n^d n^{\log_b a}^{\log_b n}}{a - b^d} \\ &= n^{\log_b a} f(1) + n^d \frac{cb^d}{b^d - a} + \frac{cb^d n^{d + \log_b a}^{\log_b n}}{a - b^d} \\ &= n^{\log_b a} f(1) + n^d \frac{cb^d}{b^d - a} + \frac{cb^d n^{d + \log_b a}^{\log_b n}}{a - b^d} \\ &= n^{\log_b a} (f(1) + \frac{cb^d}{b^d - a}) + n^d (\frac{cb^d}{b^d - a}) \\ &C_1 = \frac{b^d c}{b^d - a}, C_2 = f(1) + \frac{b^d c}{a - b^d}$$
 라고 할 때,
$$f(n) = C_1 n^d + C_2 n^{\log_b a} \end{split}$$

$$f(n) = C_1 n^d + C_2 n^{\log_b a}$$

$$a < b^d, \log_b a < d$$

$$n^d$$
이 $n^{\log_b a}$ 보다 빠르게 증가

$$\therefore f(n) = O(n^d)$$

Problem 32의 과정에서 $a>b^d$ 인 경우로 생각하면

 $n^{\log_b a}$ 이 n^d 보다 빠르게 증가한다는 사실을 알 수 있다

$$\therefore f(n) = O(n^{\log_b a})$$