CS 33

Multithreaded Programming II

Problem

```
pthread_create(&thread, 0, start, 0);
...

void *start(void *arg) {
  long BigArray[128*1024*1024];
  ...
  return 0;
}
```

Thread Attributes

```
pthread t thread;
pthread attr t thr attr;
pthread attr init(&thr attr);
/* establish some attributes */
pthread create (&thread, &thr attr, startroutine, arg);
```

Stack Size

```
pthread t thread;
pthread attr t thr attr;
pthread attr init(&thr attr);
pthread attr setstacksize(&thr attr, 130*1024*1024);
pthread create (&thread, &thr attr, startroutine, arg);
```

Execution Speed

Mutual Exclusion



Threads and Mutual Exclusion

Thread 1:

x = x+1; /* movl x,%eax incr %eax movl %eax,x */

Thread 2:

```
x = x+1;
/*
  movl x,%eax
  incr %eax
  movl %eax,x
  */
```

Quiz 1

Suppose gcc produces the following code. Will it still be the case that x's value might not be incremented by 2?

- a) yes
- b) no

Thread 1:

x = x+1; /* incr x */

Thread 2:

POSIX Threads Mutual Exclusion

```
pthread mutex t m =
     PTHREAD MUTEX INITIALIZER;
     // shared by both threads
int x; // ditto
 pthread mutex lock(&m);
 x = x+1;
 pthread mutex unlock (&m);
```

Correct Usage

```
pthread mutex lock(&m);
                                // in thread 1
// critical section
                               pthread mutex lock (&m);
                               // critical section
pthread mutex unlock(&m);
                               return;
                                   in thread 2
                               pthread mutex unlock (&m);
```

A Queue

```
void enqueue (node t *item) {
    pthread mutex lock(&mutex);
    item->next = NULL;
    if (tail == NULL) {
        head = item;
        tail = item;
    } else {
        tail->next = item;
    pthread mutex unlock(&mutex);
                              tail
```

head

```
node t *dequeue() {
    node t *ret;
    pthread mutex lock(&mutex);
    if (head == NULL) {
        ret = NULL;
    } else {
        ret = head;
        head = head->next;
        if (head == NULL)
            tail = NULL;
    pthread mutex unlock(&mutex);
    return ret;
```

Removing a Freelist Block

```
void pull_from_freelist(fblock_t *fbp) {
    ...
fbp->blink->flink = fbp->flink;
fbp->flink->blink = fbp->blink;
    ...
}
```

Parallelizing It

- Coarse grained
 - one mutex for the heap
 - threads lock the mutex before doing any operation
 - unlock it afterwards
 - only one thread at a time

- Fine grained
 - one mutex for each block
 - threads lockmutexes of onlythe blocks theyare using
 - multiple threads at a time

Removing a Freelist Block: Coarse Grained

```
void pull_from_freelist(fblock_t *fbp) {
    pthread_mutex_lock(&heap_mutex);
    ...
    fbp->blink->flink = fbp->flink;
    fbp->flink->blink = fbp->blink;
    ...
    pthread_mutex_unlock(&heap_mutex);
}
```

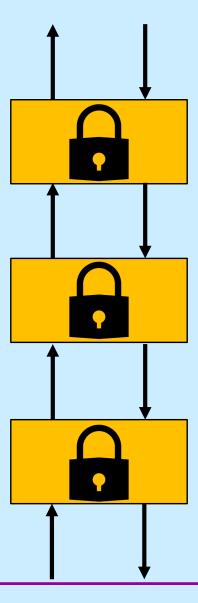
Removing a Freelist Block: Fine Grained (1)

```
void pull_from_freelist(fblock_t *fbp) {
    pthread_mutex_lock(&fpp->mutex);
    ...
    fbp->blink->flink = fbp->flink;
    fbp->flink->blink = fbp->blink;
    ...
    pthread_mutex_unlock(&fpp->mutex);
}
```

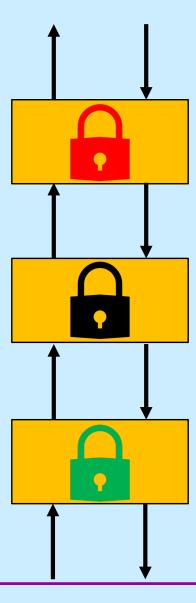
Removing a Freelist Block: Fine Grained (2)

```
void pull from freelist(fblock t *fbp) {
    pthread mutex lock(&fpp->mutex);
    pthread mutex lock(&fpp->blink->mutex);
    fbp->blink->flink = fbp->flink;
    pthread mutex lock(&fpp->flink->mutex);
    fbp->flink->blink = fbp->blink;
    pthread mutex unlock(&fpp->blink->mutex);
    pthread mutex unlock(&fpp->flink->mutex);
    pthread mutex unlock(&fpp->mutex);
```

Multiple Pulls



Multiple Pulls

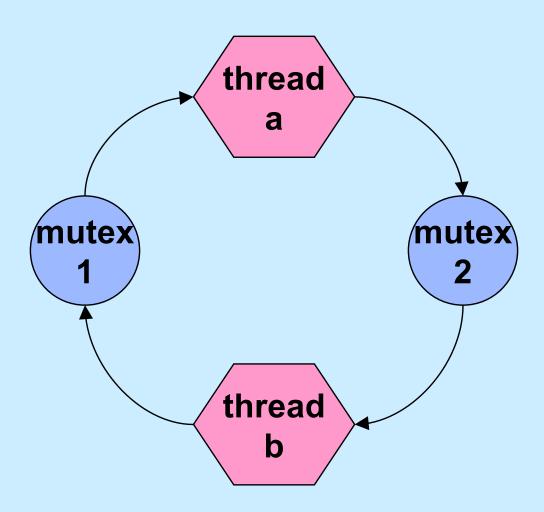


Taking Multiple Locks

```
func1() {
 pthread mutex lock(&m1);
 /* use object 1 */
 pthread mutex lock(&m2); pthread mutex lock(&m1);
```

```
func2() {
                pthread mutex lock(&m2);
                      /* use object 2 */
pthread mutex unlock(&m2); pthread mutex unlock(&m1);
pthread mutex unlock(&m1); pthread mutex unlock(&m2);
```

Preventing Deadlock

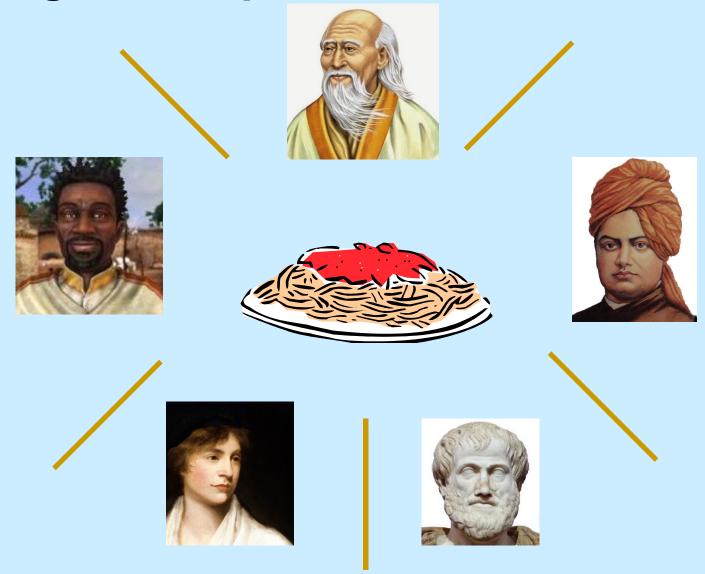


Taking Multiple Locks, Safely

```
proc1() {
 pthread mutex lock(&m1);
pthread mutex lock(&m1);
  /* use object 1 */
 pthread mutex lock(&m2);
pthread mutex lock(&m2);
```

```
proc2() {
                     /* use object 1 */
pthread mutex unlock(&m2); pthread mutex unlock(&m2);
pthread mutex unlock(&m1); pthread mutex unlock(&m1);
```

Dining Philosophers Problem



Practical Issues with Mutexes

- Used a lot in multithreaded programs
 - speed is really important
 - » shouldn't slow things down much in the success case
 - checking for errors slows things down (a lot)
 - » thus errors aren't checked by default

Set Up

Stupid (i.e., Common) Mistakes ...

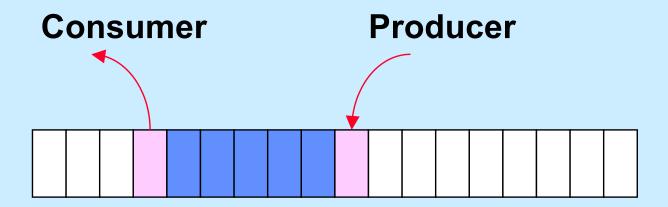
```
pthread_mutex_lock(&m1);
pthread_mutex_lock(&m1);
  // really meant to lock m2 ...

pthread_mutex_lock(&m1);
  ...
pthread_mutex_unlock(&m2);
  // really meant to unlock m1 ...
```

Runtime Error Checking

```
pthread mutexattr t err chk attr;
pthread mutexattr init(&err chk attr);
pthread mutexattr settype (&err chk attr,
       PTHREAD MUTEX ERRORCHECK);
pthread mutex t mut1;
pthread mutex init(&mut1, &err chk attr);
pthread mutex lock(&mut1);
if (pthread mutex lock(&mut1) == EDEADLK)
  fprintf(stderr, "error caught at runtime\n");
if (pthread mutex unlock(&mut2) == EPERM)
  fprintf(stderr, "another error: you didn't lock it!\n");
```

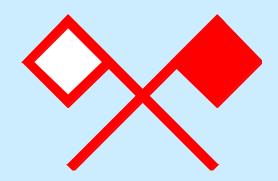
Producer-Consumer Problem



Guarded Commands

```
when (guard) [
 /*
    once the guard is true, execute this
    code atomically
   * /
```

Semaphores



• P(S) operation:

V(S) operation:

$$[S = S + 1;]$$

Quiz 2

```
semaphore S = 1;
int count = 0;
void func() {
  P(S);
  count++;
  count--;
  V(S);
```

The function func is called concurrently by n threads. What's the maximum value that count will take on?

- a) 1
- b) 2
- c) n
- d) indeterminate

```
P(S) operation:
```

```
when (S > 0) [
S = S - 1;
]
```

• V(S) operation:

$$[S = S + 1;]$$

Producer/Consumer with Semaphores

```
Semaphore empty = BSIZE;
             Semaphore occupied = 0;
             int nextin = 0:
             int nextout = 0;
char item;
 P(empty);
                             P(occupied);
 buf[nextin] = item;
 if (++nextin >= BSIZE)
                             item = buf[nextout];
   nextin = 0;
                             if (++nextout >= BSIZE)
 V (occupied);
                               nextout = 0;
                             V(empty);
                             return item;
```

POSIX Semaphores

```
#include <semaphore.h>
int sem_init(sem_t *semaphore, int pshared, int init);
int sem_destroy(sem_t *semaphore);
int sem_wait(sem_t *semaphore);
    /* P operation */
int sem_trywait(sem_t *semaphore);
    /* conditional P operation */
int sem_post(sem_t *semaphore);
    /* V operation */
```

Producer-Consumer with POSIX Semaphores

```
sem init(&empty, 0, BSIZE);
              sem init(&occupied, 0, 0);
             int nextin = 0;
             int nextout = 0;
void produce(char item) {          char consume() {
                                char item;
  sem wait(&empty);
                                sem wait(&occupied);
 buf[nextin] = item;
                                item = buf[nextout];
                                if (++nextout >= BSIZE)
  if (++nextin >= BSIZE)
   nextin = 0;
                                  nextout = 0;
  sem post (&occupied);
                                sem post(&empty);
                                return item;
```

Quiz 3

Does the POSIX version of the producerconsumer solution work with multiple producers and consumers?

- a) Yes
- b) No, but it can be made to work by using mutexes to make sure that only one thread is executing the producer code at a time and only one thread is executing the consumer code at a time
- c) It can't easily be made to work





```
void wait_for_start(state_t *s);

void start(state_t *s);

void stop(state t *s);
```





```
void wait for start(state_t *s) {
  if (s->state == stopped)
    sleep();
void start(state t *s) {
  state = started;
  wakeup all();
void stop(state t *s) {
  state = stopped;
```





```
void wait for start(state t *s) {
 pthread mutex lock(&s->mutex);
  if (s->state == stopped) {
    pthread mutex unlock(&s->mutex);
    sleep();
  else pthread mutex unlock(&s->mutex);
void start(state t *s) {
 pthread mutex lock(&s->mutex);
  state = started;
  wakeup_all();
 pthread mutex unlock(&s->mutex);
```







```
void wait for start(state_t *s) {
 pthread mutex lock(&s->mutex);
  if (s->state == stopped) {
    sleep();
 pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  state = started;
  wakeup all();
 pthread mutex unlock(&s->mutex);
```





```
void wait for start(state t *s) {
 pthread mutex lock(&s->mutex);
 while(s->state == stopped)
    pthread cond wait(&s->queue, &s->mutex);
 pthread mutex unlock(&s->mutex);
void start(state t *s) {
 pthread mutex lock(&s->mutex);
  s->state = started;
 pthread cond broadcast (&s->queue);
 pthread mutex unlock(&s->mutex);
```

Condition Variables

```
when (quard) [
                                    pthread mutex lock(&mutex);
                                    while(!quard)
  statement 1;
                                      pthread cond wait (
  statement n;
                                          &cond var, &mutex);
                                    statement 1:
                                    statement n;
                                    pthread mutex unlock (&mutex);
// code modifying the guard:
                                    pthread mutex lock(&mutex);
                                    // code modifying the guard:
                                    pthread cond broadcast (
                                        &cond var);
                                    pthread mutex unlock(&mutex);
```

Set Up

PC with Condition Variables (1)

```
typedef struct buffer {
   pthread_mutex_t m;
   pthread_cond_t more_space;
   pthread_cond_t more_items;
   int next_in;
   int next_out;
   int empty;
   char buf[BSIZE];
}
```

PC with Condition Variables (2)

```
void produce(buffer t *b,
    char item) {
                                    char item;
  pthread mutex lock(&b->m);
  while (!(b->empty > 0))
   pthread cond wait (
       &b->more space, &b->m); item = b->buf[b->nextout];
  b->buf[b->nextin] = item;
  if (++(b->nextin) == BSIZE)
    b->nextin = 0;
  b->empty--;
  pthread cond signal (
     &b->more items);
  pthread mutex unlock(&b->m);
```

```
char consume(buffer t *b) {
  pthread mutex lock(&b->m);
  while (!(b->empty < BSIZE))</pre>
   pthread cond wait (
       &b->more items, &b->m);
  if (++(b->nextout) == BSIZE)
    b->nextout = 0:
  b->empty++;
  pthread cond signal (
     &b->more space);
  pthread mutex unlock(&b->m);
  return item;
```