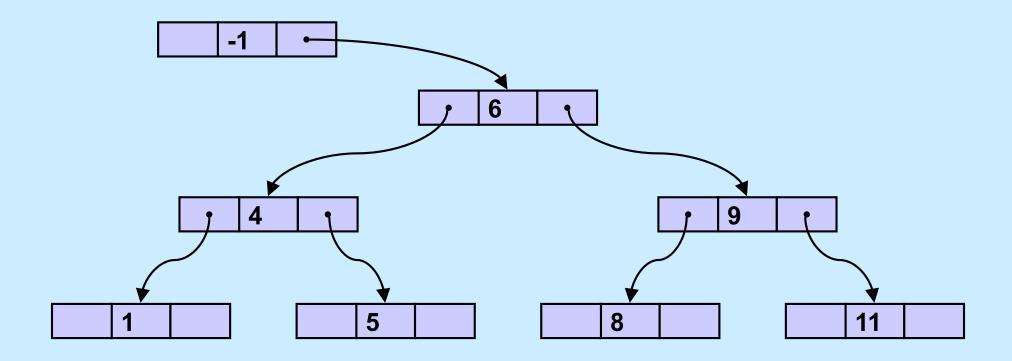
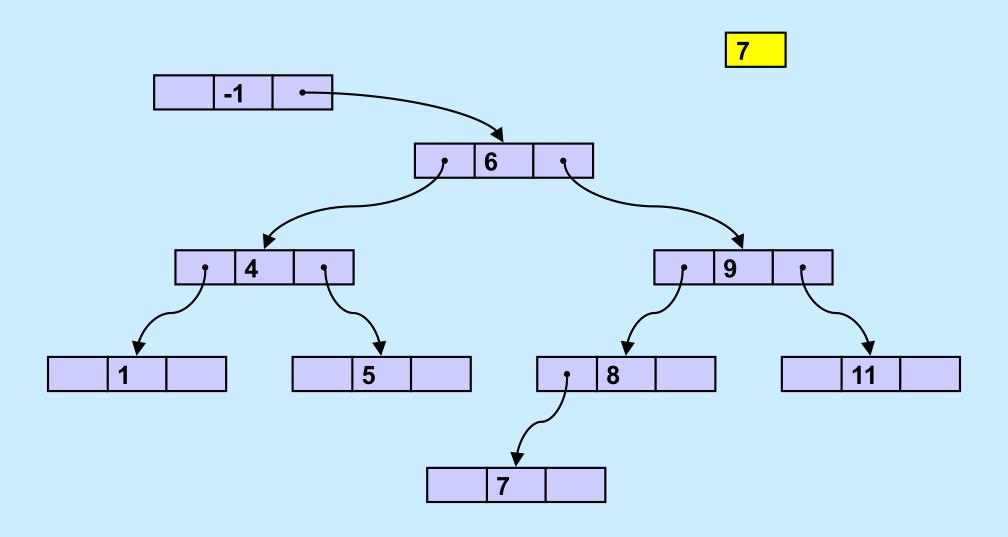
CS 33

Multithreaded Programming IV

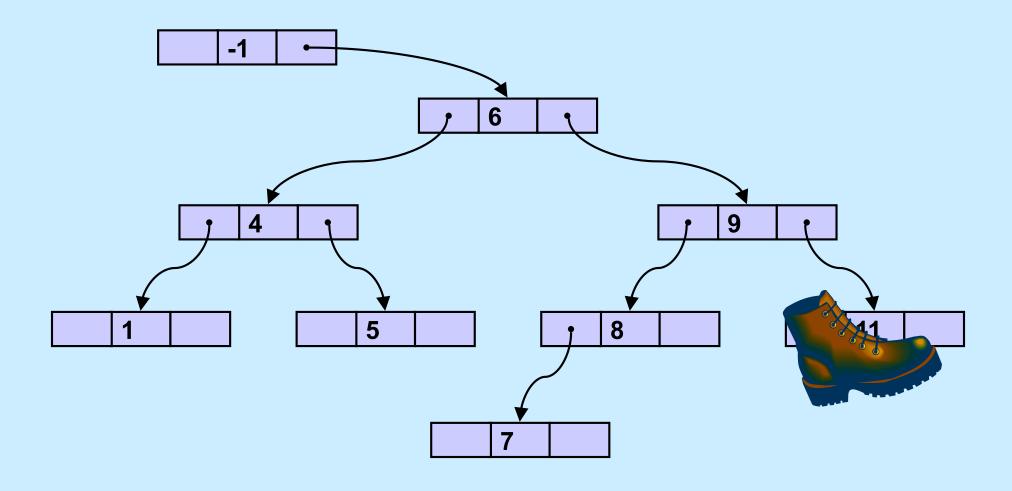
Binary Search Tree



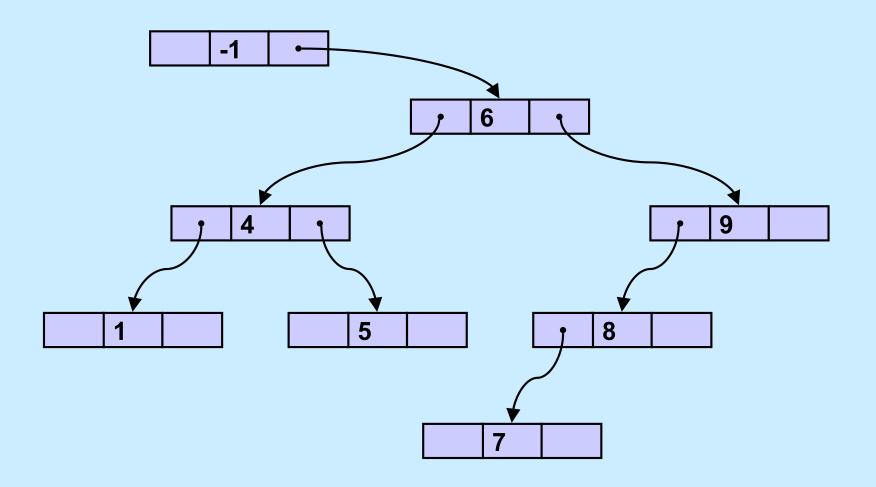
Binary Search Tree: Insertion



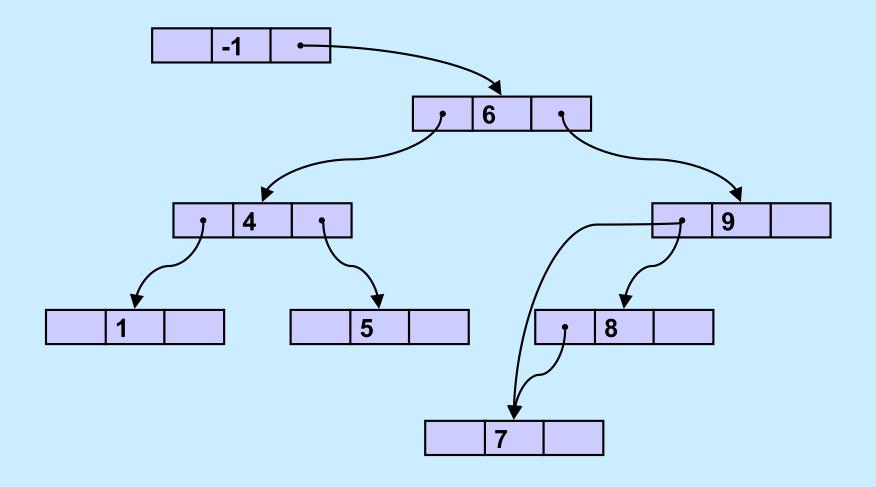
Binary Search Tree: Deletion of Leaf



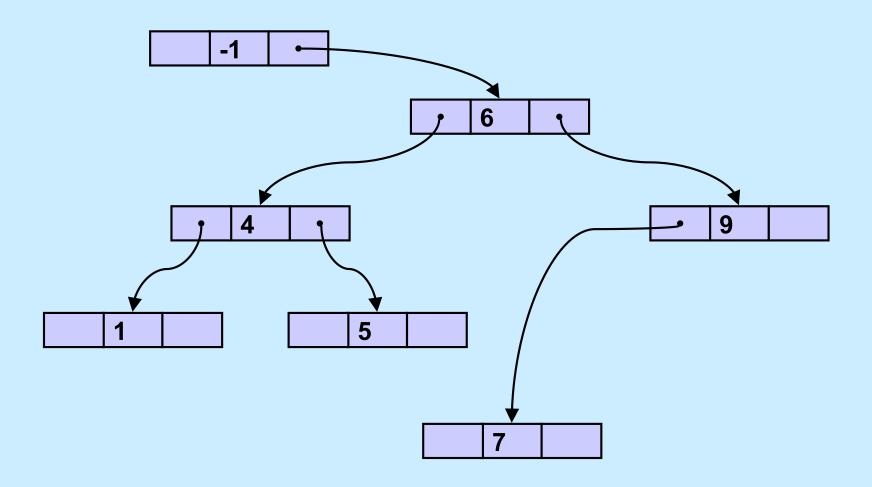
Binary Search Tree: Deletion of Leaf



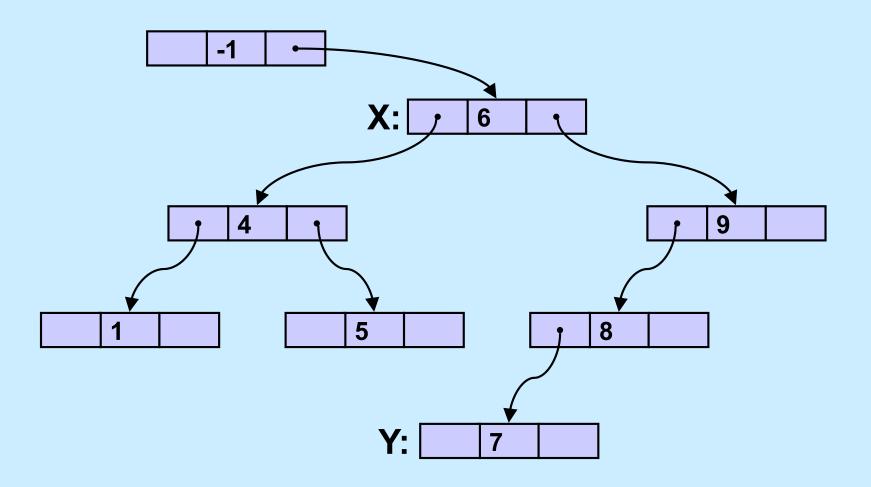
Binary Search Tree: Deletion of Node with One Child



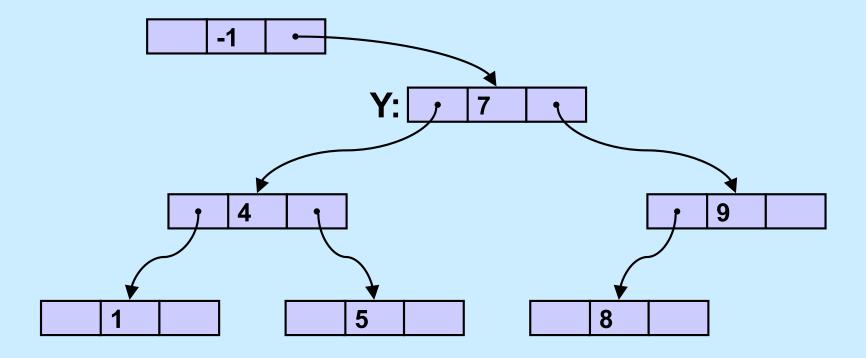
Binary Search Tree: Deletion of Node with One Child



Binary Search Tree: Deletion of Node with Two Children



Binary Search Tree: Deletion of Node with Two Children



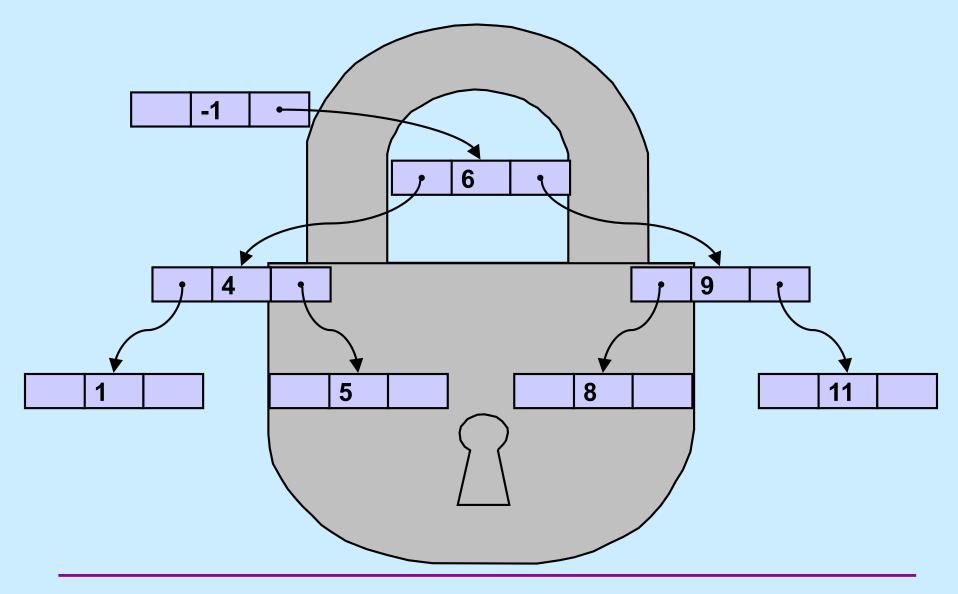
C Code: Search

```
Node *search(int key,
                                        } else {
    Node *parent, Node **parentp) {
                                          if ((next = parent->rchild)
  Node *next;
                                              == 0)
  Node *result;
                                            result = 0;
  if (key < parent->key) {
                                          } else {
    if ((next = parent->lchild)
                                            if (key == next->key) {
        == 0)
                                              result = next;
      result = 0;
                                            } else {
    } else {
                                              result = search(key,
      if (key == next->key) {
                                                  next, parentpp);
        result = next;
                                              return result;
      } else {
        result = search(key,
            next, parentpp);
                                        if (parentpp != 0)
        return result;
                                          *parentpp = parent;
                                        return result;
```

C Code: Add

```
int add(int key) {
 Node *parent, *target, *newnode;
  if ((target = search(key, &head, &parent)) != 0) {
    return 0;
 newnode = malloc(sizeof(Node));
 newnode->key = key;
 newnode->lchild = newnode->rchild = 0;
  if (name < parent->name)
    parent->lchild = newnode;
 else
    parent->rchild = newnode;
  return 1;
```

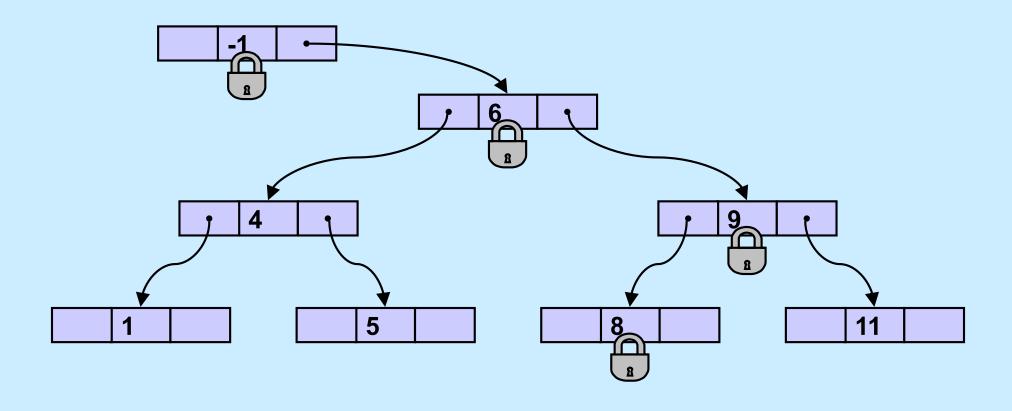
Binary Search Treewith Coarse-Grained Synchronization



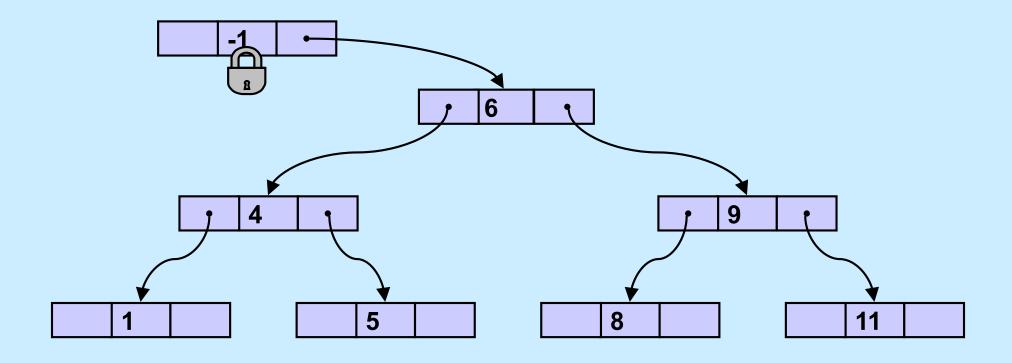
C Code: Add with Coarse-Grained Synchronization

```
int add(int key) {
 Node *parent, *target, *newnode;
 pthread rwlock wrlock (&tree lock);
 if ((target = search(key, &head, &parent)) != 0) {
   pthread rwlock unlock (&tree lock);
   return 0:
 newnode = malloc(sizeof(Node));
 newnode -> key = key;
 newnode->lchild = newnode->rchild = 0;
 if (name < parent->name)
   parent->lchild = newnode;
 else
   parent->rchild = newnode;
 pthread rwlock unlock (&tree lock);
 return 1:
```

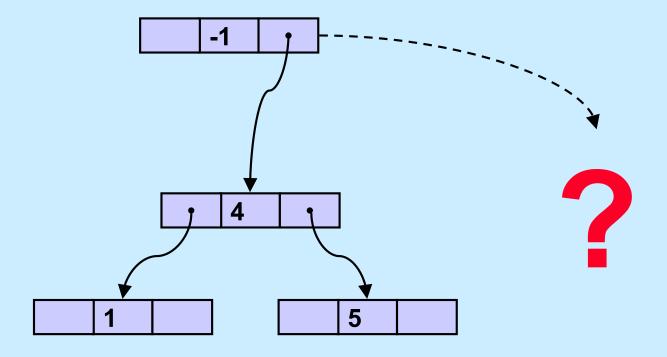
Binary Search Treewith Fine-Grained Synchronization I



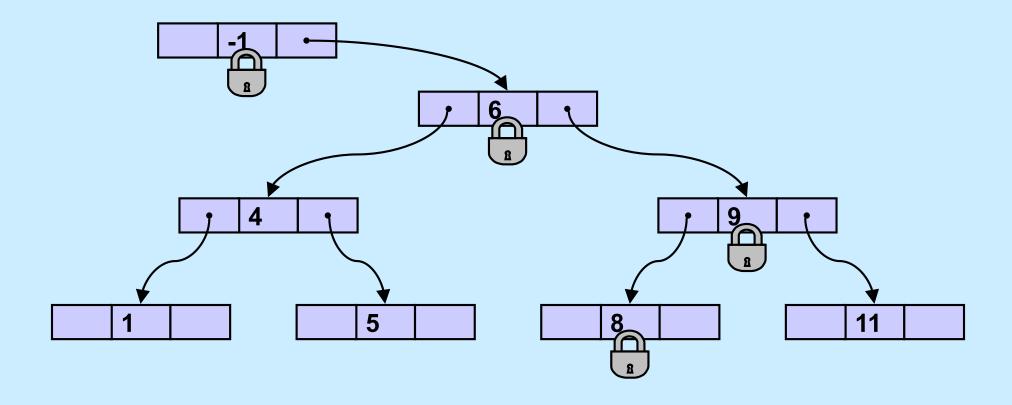
Binary Search Treewith Fine-Grained Synchronization II



Binary Search Treewith Fine-Grained Synchronization III



Doing It Right ...



C Code: Fine-Grained Search I

```
enum locktype {l read, l write};
                                         } else {
                                              lock(lt, &next->lock);
#define lock(lt, lk) ((lt) == l read)?
                                              if (key == next->key) {
      pthread rwlock rdlock(lk):
                                                result = next;
      pthread rwlock wrlock(lk)
                                              } else {
                                                pthread rwlock unlock (
Node *search(int key,
                                                     &parent->lock);
    Node *parent, Node **parentp,
                                                result = search(key,
    enum locktype lt) {
                                                    next, parentpp, lt);
   // parent is locked on entry
                                                 return result;
 Node *next;
 Node *result;
  if (key < parent->key) {
    if ((next = parent->lchild)
        == 0)
      result = 0;
```

C Code: Fine-Grained Search II

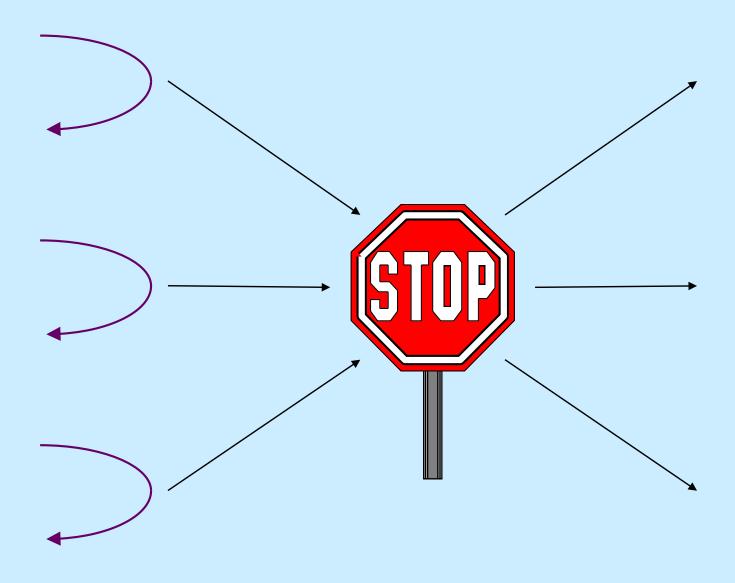
```
} else {
                                           } else {
  if ((next = parent->rchild)
                                            pthread rwlock unlock (
      == 0)
                                                 &parent->lock);
     result = 0;
                                             result = search(key,
   } else {
                                                 next, parentpp, lt);
     lock(lt, &next->lock);
                                            return result;
     if (key == next->key) {
       result = next;
                                      if (parentpp != 0) {
                                        // parent remains locked
                                        *parentpp = parent;
                                      } else
                                        pthread rwlock unlock (
                                             &parent->lock);
                                      return result;
```

C Code: Add with Fine-Grained Synchronization I

C Code: Add with Fine-Grained Synchronization II

```
newnode = malloc(sizeof(Node));
newnode->key = key;
newnode->lchild = newnode->rchild = 0;
pthread_rwlock_init(&newnode->lock, 0);
if (name < parent->name)
   parent->lchild = newnode;
else
   parent->rchild = newnode;
pthread_rwlock_unlock(&parent->lock);
return 1;
```

Barriers



A Solution?

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
} else while (!(count == number)) {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

How About This?

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
   count = 0;
} else while (!(count == number)) {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

And This ...

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
   count = 0;
} else {
   pthread_cond_wait(&cond_var, &m);
}
pthread mutex unlock(&m);
```

- a) definitely
- b) probably

Barrier in POSIX Threads

```
pthread mutex lock(&m);
if (++count < number) {</pre>
  int my generation = generation;
  while (my generation == generation) {
    pthread cond wait(&waitQ, &m);
 else {
  count = 0;
  generation++;
  pthread cond broadcast (&waitQ);
pthread mutex unlock (&m);
```

More From POSIX!

Why cond_wait is Weird ...

```
pthread_cond_wait(pthread_cond_t *c, pthread_mutex_t *m) {
    pthread_mutex_unlock(m);
    sem_wait(c->sem);
    pthread_mutex_lock(m);
}

pthread_cond_signal(pthread_cond_t *c) {
    sem_post(c->sem);
}
```

Deviations

Signals



VS.



- Cancellation
 - tamed lightning

Signals



- who gets them?
- who needs them?



– how do you respond to them?

Dealing with Signals

- Per-thread signal masks
- Per-process signal vectors
- One delivery per signal

Signals and Threads

```
int pthread kill(pthread t thread, int signo);
```

thread equivalent of kill

```
int pthread sigmask (int how,
     const sigset t *newmask,
     sigset t oldmask);
```

thread equivalent of sigprocmask

Asynchronous Signals (1)

```
int main() {
  void handler(int);
   signal(SIGINT, handler);
void handler(int sig) {
```

Asynchronous Signals (2)

```
int main() {
                                 void handler(int sig) {
 void handler(int);
                                   ... // deal with signal
  signal(SIGINT, handler);
                                   printf("equally important "
  ... // complicated program
                                      "message: %s\n", message);
 printf("important message:
     "%s\n", message);
  ... // more program
```

Quiz 2

```
int main() {
  void handler(int);
  signal(SIGINT, handler);
  ... // complicated program
 pthread mutex lock(&mut);
 printf("important message: "
     "%s\n", message);
 pthread mutex unlock (&mut);
  ... // more program
```

```
void handler(int sig) {
  ... // deal with signal
 pthread mutex lock(&mut);
  printf("equally important "
     "message: %s\n", message);
 pthread mutex unlock(&mut);
```

Does this work? a) yes

Synchronizing Asynchrony

```
computation state t state;
sigset t set;
int main() {
  pthread_t thread;
  sigemptyset(&set);
  sigaddset(&set, SIGINT);
  pthread sigmask (SIG BLOCK,
   &set, 0);
  pthread create (&thread, 0,
   monitor, 0);
  long running procedure();
```

```
void *monitor(void *dummy) {
  int sig;
  while (1) {
    sigwait(&set, &sig);
    display(&state);
  }
  return(0);
}
```

Cancellation



Sample Code

```
void *thread code(void *arg) {
  node t *head = 0;
  while (1) {
    node t *nodep;
    nodep = (node t *) malloc(sizeof(node t));
    if (read(0, &node->value,
        sizeof(node->value)) == 0) {
      free (nodep);
      break;
                               pthread cancel(thread);
    nodep->next = head;
    head = nodep;
  return head;
```

Cancellation Concerns

- Getting cancelled at an inopportune moment
- Cleaning up

Cancellation State

Pending cancel

```
- pthread cancel (thread)
```

Cancels enabled or disabled

```
- int pthread_setcancelstate(
     {PTHREAD_CANCEL_DISABLE
     PTHREAD_CANCEL_ENABLE},
     &oldstate)
```

Asynchronous vs. deferred cancels

```
- int pthread_setcanceltype(
     {PTHREAD_CANCEL_ASYNCHRONOUS,
     PTHREAD_CANCEL_DEFERRED),
     &oldtype)
```

Cancellation Points

- aio_suspend
- close
- creat
- fcntl (when F_SETLCKW is the command)
- fsync
- mq_receive
- mq_send
- msync
- nanosleep
- open
- pause
- pthread_cond_wait
- pthread_cond_timedwait
- pthread_join

- pthread_testcancel
- read
- sem_wait
- sigwait
- sigwaitinfo
- sigsuspend
- sigtimedwait
- sleep
- system
- tcdrain
- wait
- waitpid
- write

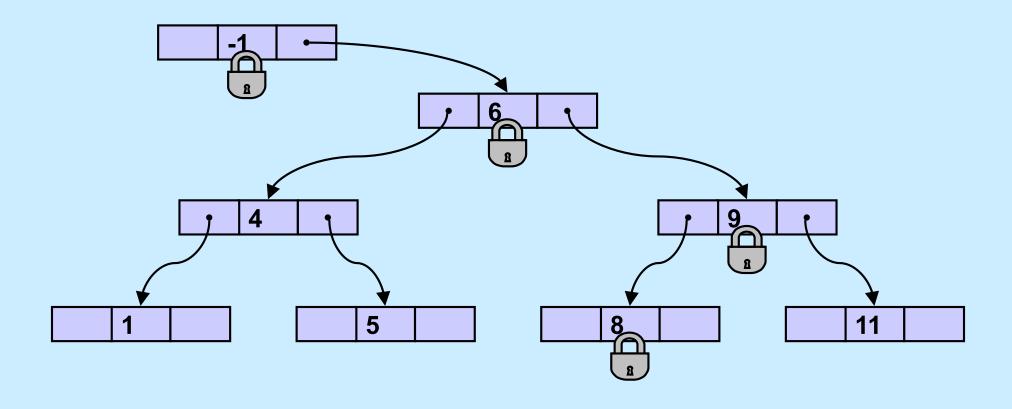
Cleaning Up

- void pthread_cleanup_push((void)(*routine)(void *),void *arg)
- void pthread_cleanup_pop(int execute)

Sample Code, Revisited

```
void *thread code(void *arg) {
                                          void cleanup(void *arg) {
  node t *head = 0;
                                            node t **headp = arg;
  pthread cleanup push (
                                            while(*headp) {
      cleanup, &head);
                                              node t *nodep = head->next;
    while (1) {
                                              free (*headp);
      node t *nodep;
                                              *headp = nodep;
      nodep = (node t *)
      malloc(sizeof(node t));
      if (read(0, &node->value,
          sizeof(node->value)) == 0) {
        free (nodep);
        break;
      nodep->next = head;
      head = nodep;
  pthread cleanup pop(0);
  return head;
```

A More Complicated Situation ...



Start/Stop





Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait(&s->queue, &s->mutex);
  pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```

Start/Stop

Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait (&s->queue,
      &s->mutex);
  pthread mutex unlock (&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```



Quiz 3

You're in charge of designing POSIX threads. Should *pthread_cond_wait* be a cancellation point?

- a) no
- b) yes; cancelled threads must acquire mutex before invoking cleanup handler
- c) yes; but they don't acquire mutex

Start/Stop



Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  pthread cleanup push (
    pthread mutex unlock, &s);
  while(s->state == stopped)
    pthread cond wait (&s->queue, &s->mutex);
  pthread cleanup pop(1);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast (&s->queue);
  pthread mutex unlock(&s->mutex);
```

Cancellation and Conditions

```
pthread_mutex_lock(&m);
pthread_cleanup_push(pthread_mutex_unlock, &m);
while(should_wait)
   pthread_cond_wait(&cv, &m);

// ... (code perhaps containing other cancellation points)
pthread_cleanup_pop(1);
```