CS 33

Machine Programming (5)

Arguments and Local Variables (C Code)

```
int mainfunc() {
  long array[3] =
      {2,117,-6};
  long sum =
      ASum(array, 3);
    ...
  return sum;
}
```

- Local variables usually allocated on stack
- Arguments to functions pushed onto stack

```
long ASum(long *a,
    unsigned long size) {
    long i, sum = 0;
    for (i=0; i<size; i++)
        sum += a[i];
    return sum;
}</pre>
```

 Local variables may be put in registers (and thus not on stack)

Arguments and Local Variables (1)

```
mainfunc:
  pushq %rbp
                             # save old %rbp
                             # set %rbp to point to stack frame
   mova %rsp, %rbp
   subq $32, %rsp
                             # alloc. space for locals (array and sum)
   movq \$2, -32(\$rbp) # initialize array[0]
   movq $117, -24(%rbp) # initialize array[1]
   movq \$-6, -16(\$rbp) # initialize array[2]
                             # push arg 2
   pusha $3
   leaq -32(%rbp), %rax
                       # array address is put in %rax
                             # push arg 1
   pushq %rax
   call ASum
   addq $16, %rsp
                            # pop args
   movq %rax, -8(%rbp)
                             # copy return value to sum
   addq $32, %rsp
                             # pop locals
```

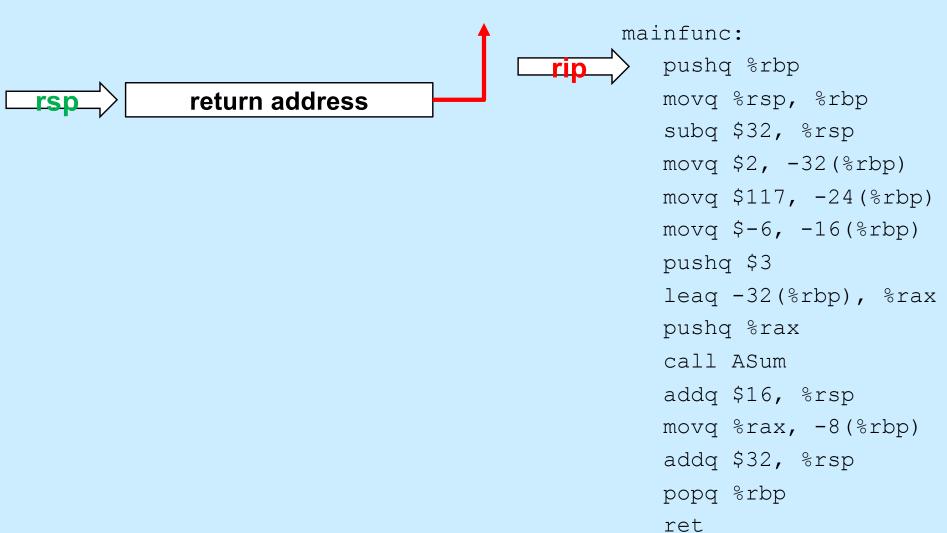
popq %rbp

pop and restore old %rbp

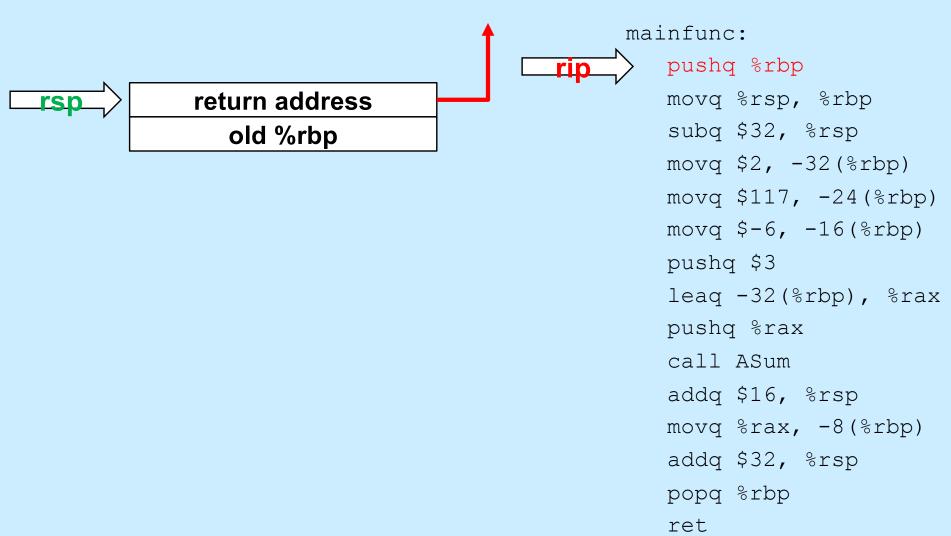
Arguments and Local Variables (2)

```
ASum:
   pushq %rbp
                              # save old %rbp
   movq %rsp, %rbp
                              # set %rbp to point to stack frame
                              # i in %rcx
   movq $0, %rcx
   movq $0, %rax
                              # sum in %rax
   movq 16(%rbp), %rdx
                              # copy arg 1 (array) into %rdx
loop:
   cmpq 24(%rbp), %rcx # i < size?</pre>
   jge done
   addq (%rdx, %rcx, 8), %rax # sum += a[i]
   incq %rcx
                              # i++
   ja loop
done:
                              # pop and restore %rbp
   popq %rbp
   ret
```

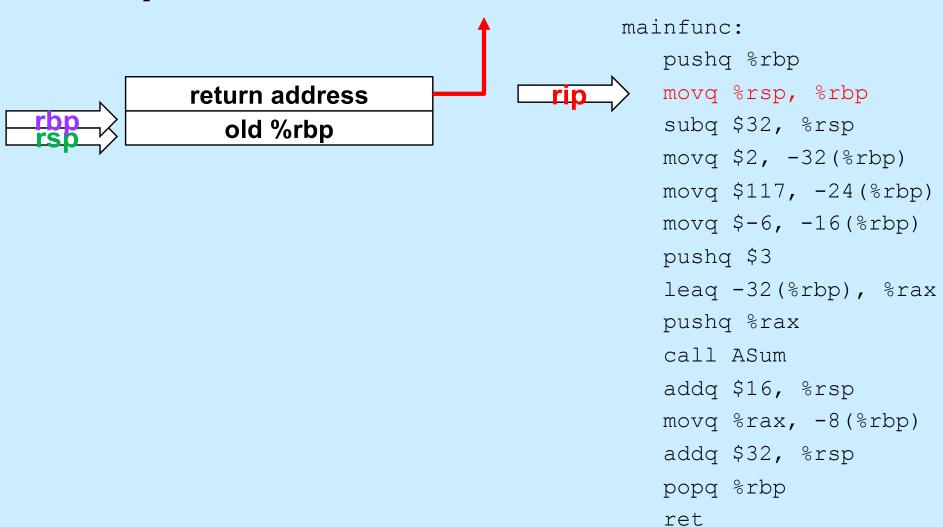
Enter mainfunc



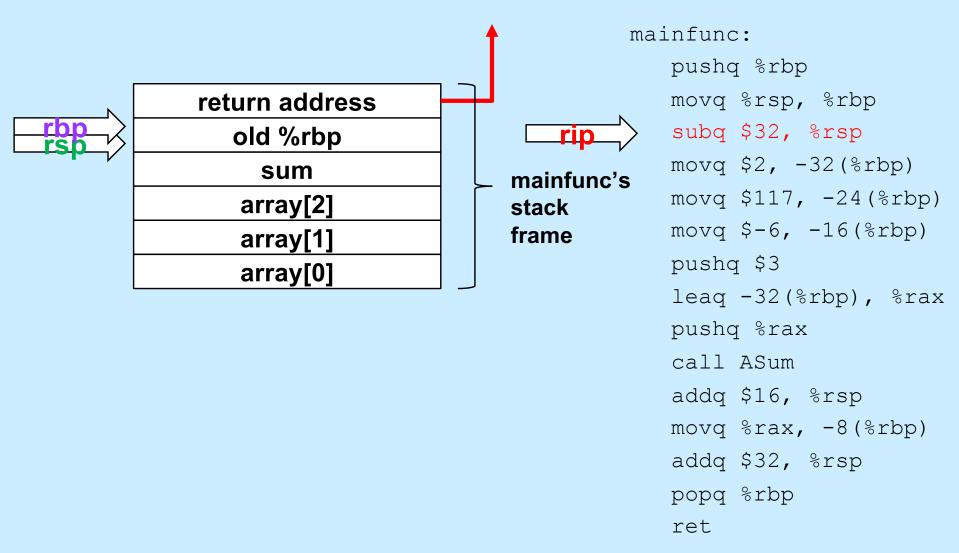
Enter mainfunc



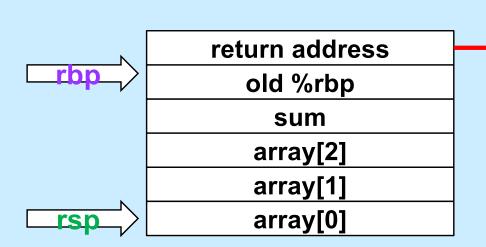
Setup Frame

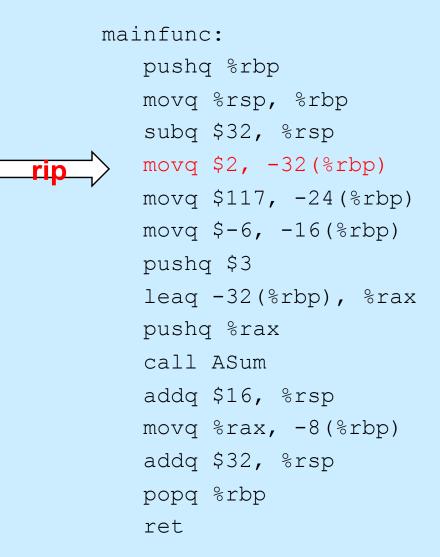


Allocate Local Variables

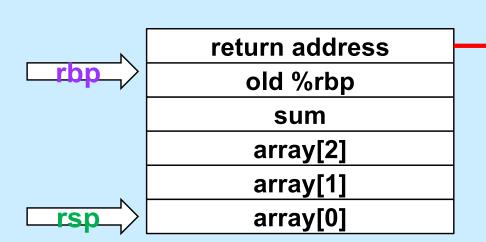


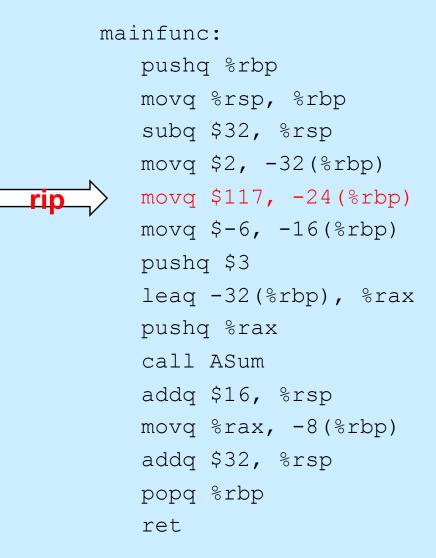
Initialize Local Array



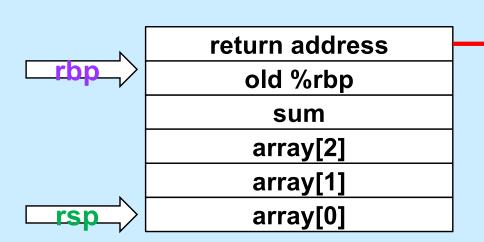


Initialize Local Array



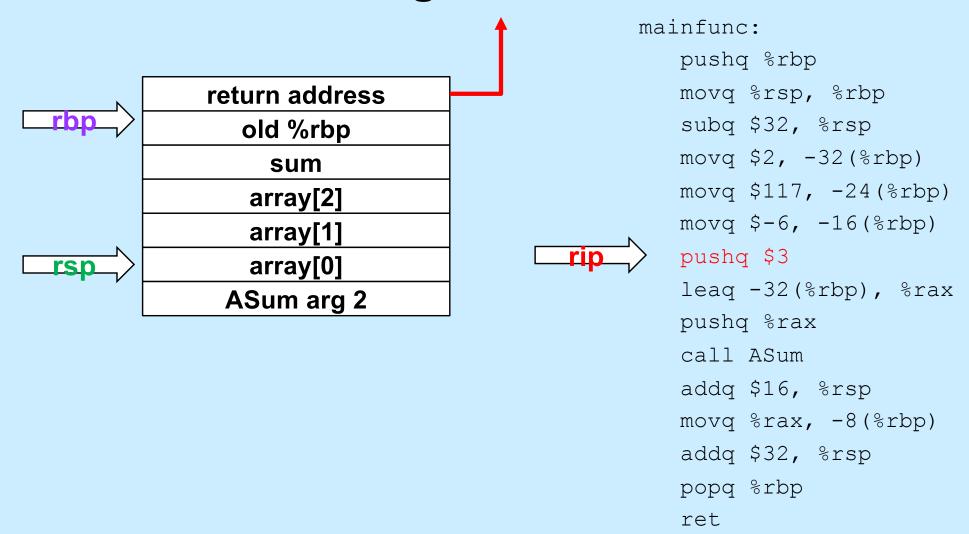


Initialize Local Array

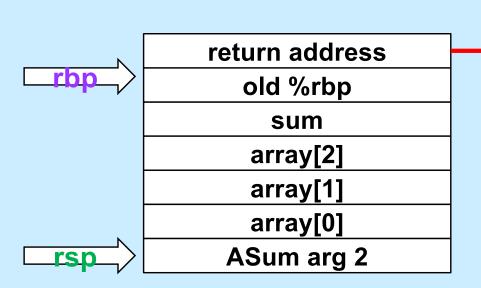


```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movq $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pushq $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Push Second Argument

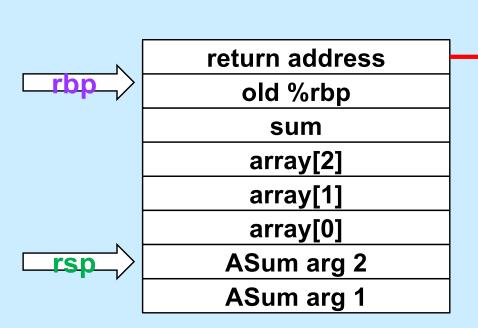


Get Array Address



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq \$-6, -16(%rbp)
   pushq $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

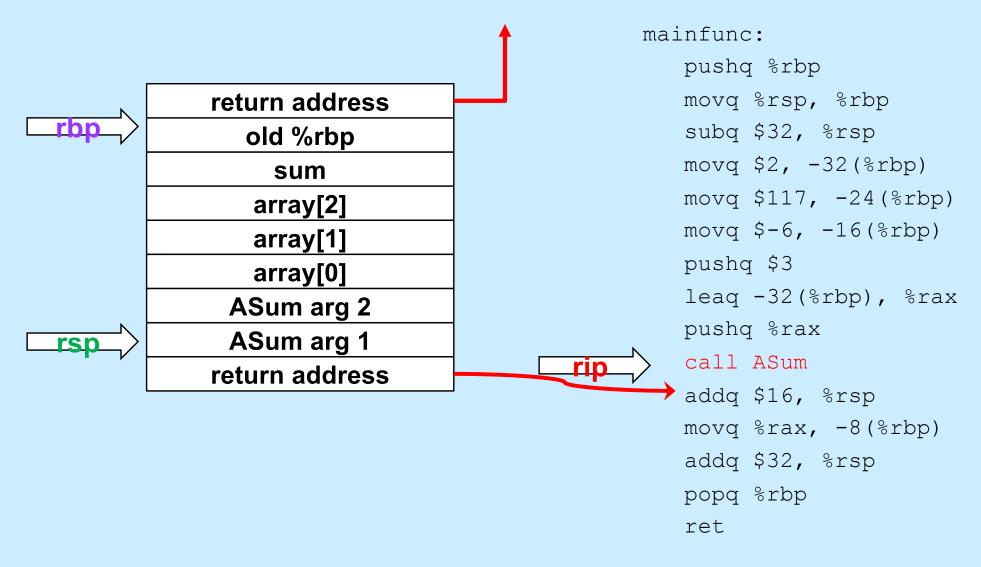
Push First Argument



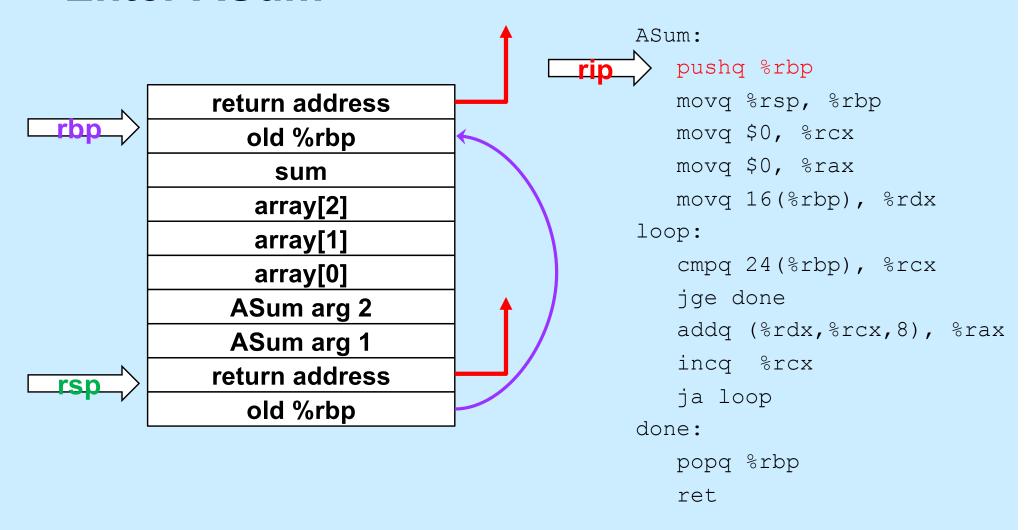
mainfunc: pushq %rbp movq %rsp, %rbp subq \$32, %rsp movq \$2, -32(%rbp) movg \$117, -24(%rbp) movq \$-6, -16(%rbp) pusha \$3 leaq -32(%rbp), %rax pushq %rax call ASum addq \$16, %rsp movq %rax, -8(%rbp) addq \$32, %rsp popq %rbp

ret

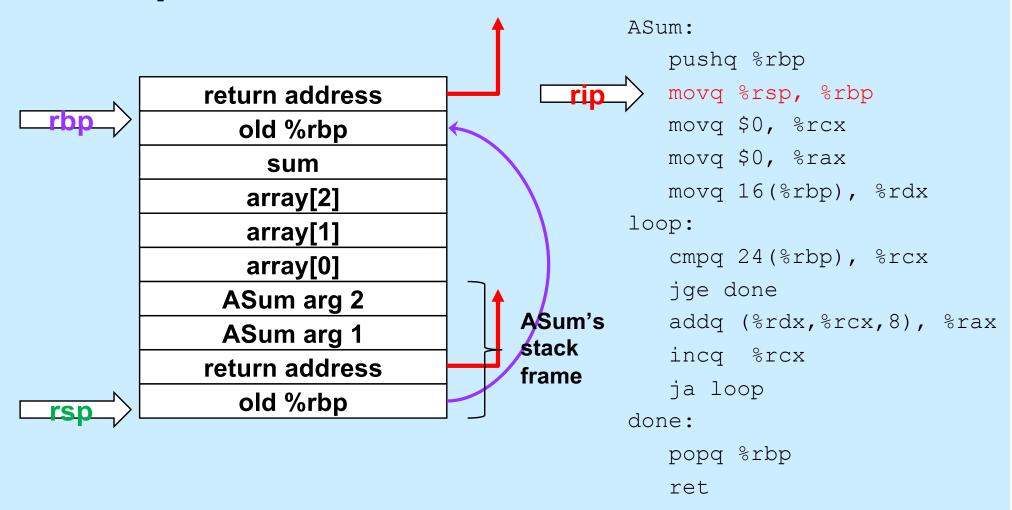
Call ASum



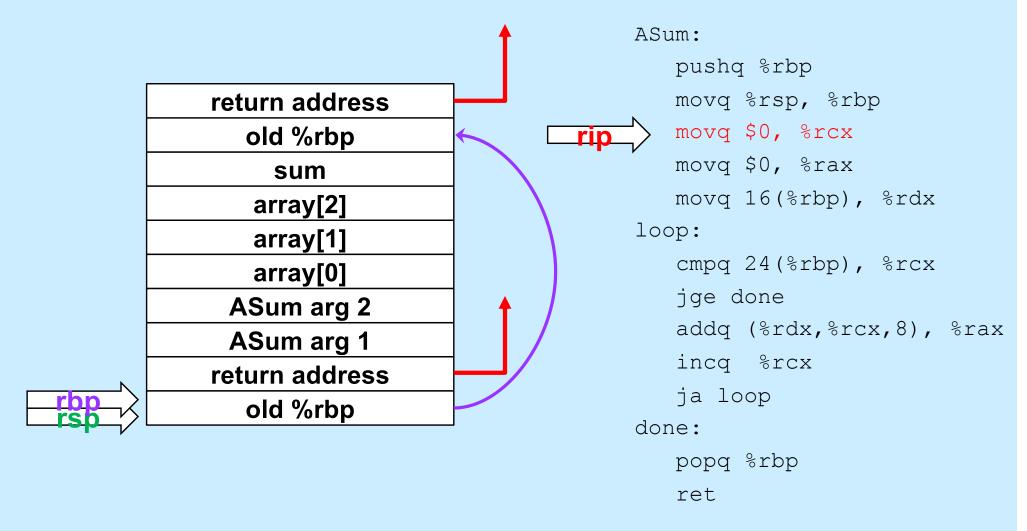
Enter ASum



Setup Frame



Execute the Function



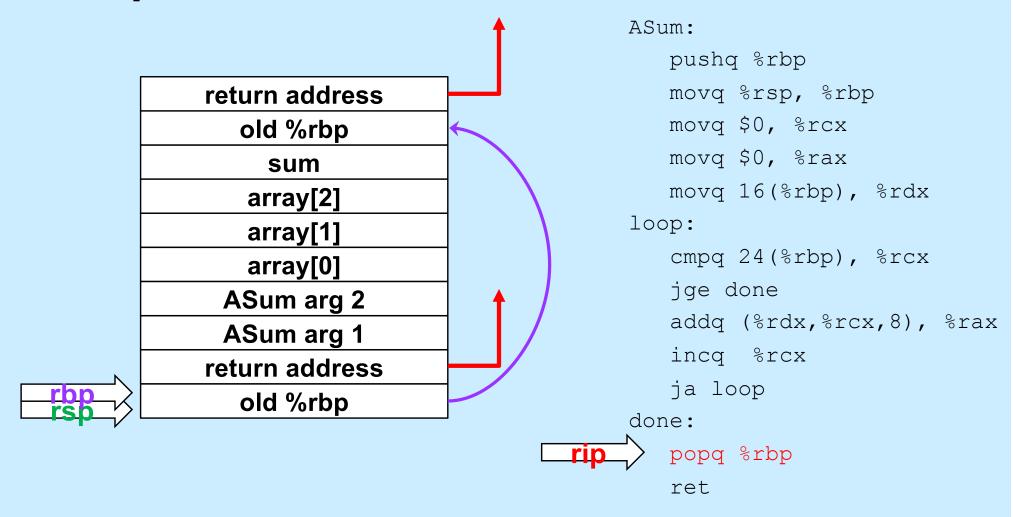
Quiz 1

What's at 16(%rbp) (after the second instruction is executed)?

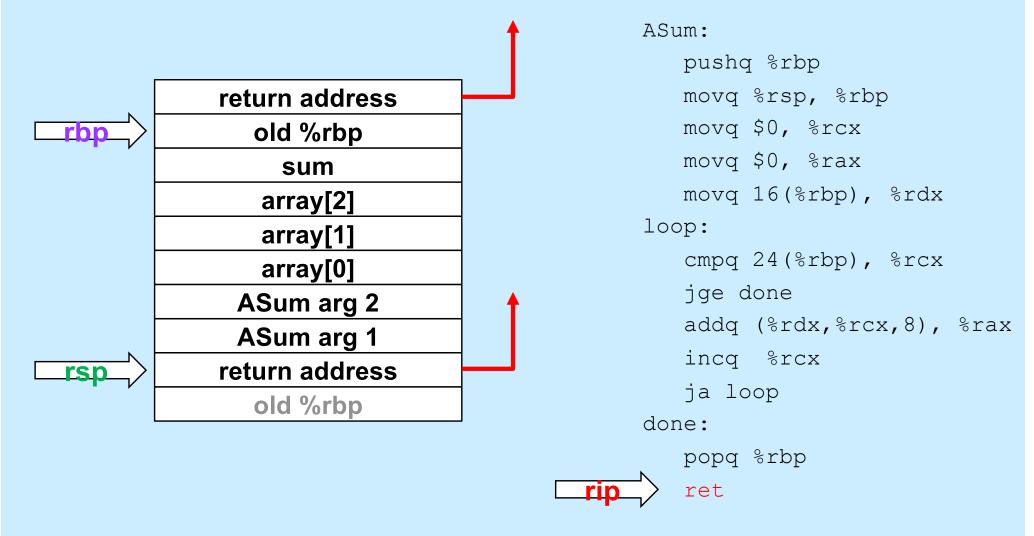
- a) a local variable
- b) the first argument to ASum
- c) the second argument to ASum
- d) something else

```
ASum:
   pushq %rbp
   movq %rsp, %rbp
   movq $0, %rcx
   movq $0, %rax
   movq 16(%rbp), %rdx
loop:
   cmpq 24(%rbp), %rcx
   jge done
   addq (%rdx,%rcx,8), %rax
   incq %rcx
   ja loop
done:
   popq %rbp
   ret
```

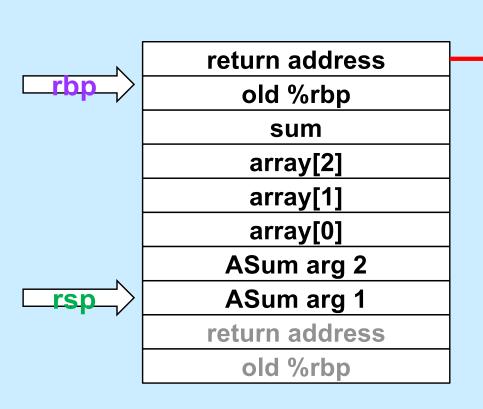
Prepare to Return



Return

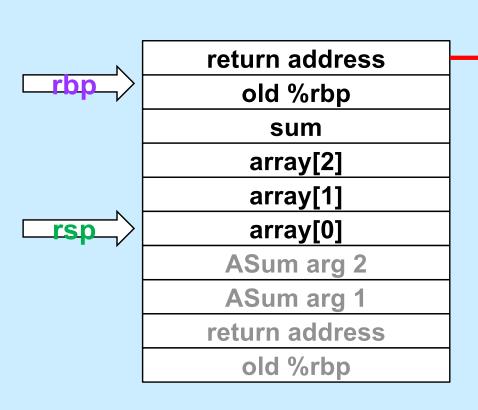


Pop Arguments



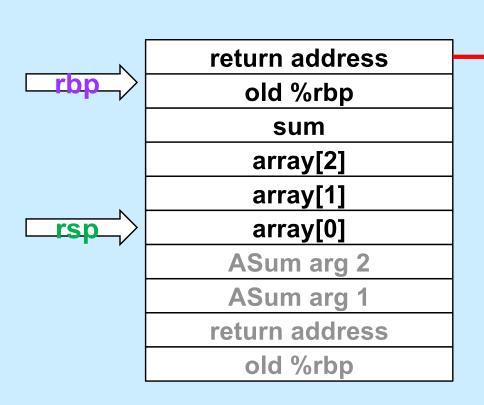
```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Save Return Value



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq \$-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

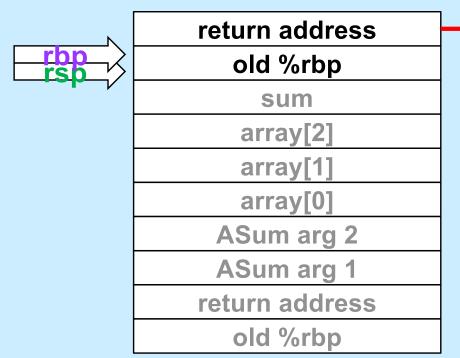
Pop Local Variables



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
```

ret

Prepare to Return



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movg $-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Return

return address
old %rbp
sum
array[2]
array[1]
array[0]
ASum arg 2
ASum arg 1
return address
old %rbp

mainfunc: pushq %rbp movq %rsp, %rbp subq \$32, %rsp movq \$2, -32(%rbp) movq \$117, -24(%rbp) movg \$-6, -16(%rbp)pushq \$3 leaq -32(%rbp), %rax pushq %rax call ASum addq \$16, %rsp movq %rax, -8(%rbp) addq \$32, %rsp popq %rbp

Using Registers

- ASum modifies registers:
 - %rsp
 - %rbp
 - %rcx
 - %rax
 - %rdx
- Suppose its caller uses these registers

```
movq $33, %rcx
movq $167, %rdx
pushq $6
pushq array
call ASum
    # assumes unmodified %rcx and %rdx
addq $16, %rsp
addq %rax,%rcx  # %rcx was modified!
addq %rdx, %rcx  # %rdx was modified!
```

```
ASum:
   pushq %rbp
   movq %rsp, %rbp
   movq $0, %rcx
   movq $0, %rax
   movq 16(%rbp), %rdx
loop:
   cmpq 24(%rbp), %rcx
   jge done
   addq (%rdx,%rcx,8), %rax
   incq %rcx
   ja loop
done:
   popq %rbp
   ret
```

Register Values Across Function Calls

- ASum modifies registers:
 - %rsp
 - %rbp
 - %rcx
 - %rax
 - %rdx
- May the caller of ASum depend on its registers being the same on return?
 - ASum saves and restores %rbp and makes no net changes to %rsp
 - » their values are unmodified on return to its caller
 - %rax, %rcx, and %rdx are not saved and restored
 - » their values might be different on return

```
ASum:
   pushq %rbp
   movq %rsp, %rbp
   movq $0, %rcx
   movq $0, %rax
   movq 16(%rbp), %rdx
loop:
   cmpq 24(%rbp), %rcx
   jge done
   addq (%rdx,%rcx,8), %rax
   incq %rcx
   ja loop
done:
   popq %rbp
   ret
```

Register-Saving Conventions

Caller-save registers

 if the caller wants their values to be the same on return from function calls, it must save and restore them

```
pushq %rcx
call func
popq %rcx
```

Callee-save registers

 if the callee wants to use these registers, it must first save them, then restore their values before returning

```
func:
    pushq %rbx
    movq $6, %rbx
    ...
    popq %rbx
```

x86-64 General-Purpose Registers: Usage Conventions

%rax	Return value
%rbx	Callee saved
%rcx	Caller saved
%rdx	Caller saved
%rsi	Caller saved
%rdi	Caller saved
%rsp	Stack pointer
%rbp	Base pointer

%r8	Caller saved
%r9	Caller saved
%r10	Caller saved
%r11	Caller Saved
%r12	Callee saved
%r13	Callee saved
%r14	Callee saved
%r15	Callee saved

Passing Arguments in Registers

Observations

- accessing registers is much faster than accessing primary memory
 - » if arguments were in registers rather than on the stack, speed would increase
- most functions have just a few arguments

Actions

- change calling conventions so that the first six arguments are passed in registers
 - » in caller-save registers
- any additional arguments are pushed on the stack

Why Bother with a Base Pointer?

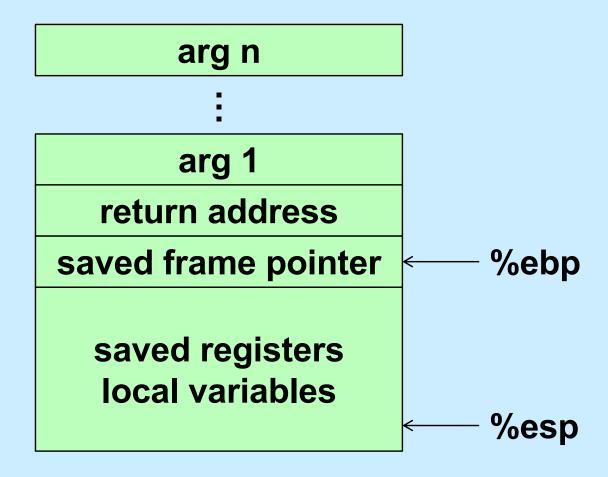
- It (%rbp) points to the beginning of the stack frame
 - making it easy for people to figure out where things are in the frame
 - but people don't execute the code ...
- The stack pointer always points somewhere within the stack frame
 - it moves about, but the compiler knows where it is pointing
 - » a local variable might be at 8(%rsp) for one instruction, but at 16(%rsp) for a subsequent one
 - » tough for people, but easy for the compiler
- Thus the base pointer is superfluous
 - it can be used as a general-purpose register

x86-64 General-Purpose Registers: Updated Usage Conventions

%rax	Return value
%rbx	Callee saved
%rcx	Argument #4
%rdx	Argument #3
%rsi	Argument #2
%rdi	Argument #1
%rsp	Stack pointer
%rbp	Callee saved

%r8	Argument #5
8r9	Argument #6
%r10	Caller saved
%r11	Caller Saved
%r12	Callee saved
%r13	Callee saved
%r14	Callee saved
%r15	Callee saved

The IA32 Stack Frame



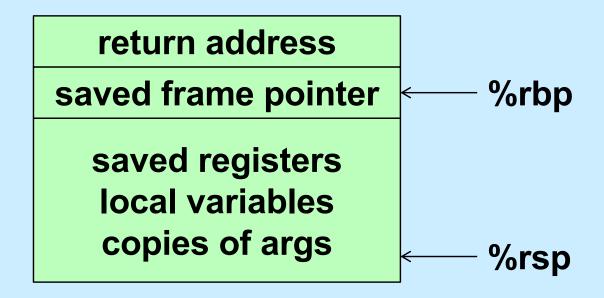
The x86-64 Stack Frame

return address

saved registers local variables

%rsp

The -O0 x86-64 Stack Frame (Buffer)



Summary

- What's pushed on the stack
 - return address
 - saved registers
 - » caller-saved by the caller
 - » callee-saved by the callee
 - local variables
 - function parameters
 - » those too large to be in registers (structs)
 - » those beyond the six that we have registers for
 - large return values (structs)
 - » caller allocates space on stack
 - » callee copies return value to that space

Quiz 2

Suppose function A is compiled using the convention that %rbp is used as the base pointer, pointing to the beginning of the stack frame. Function B is compiled using the convention that there's no need for a base pointer. Will there be any problems if A calls B or if B calls A?

- a) Neither case will work
- b) A calling B works, but B calling A doesn't
- c) B calling A works, but A calling B doesn't
- d) Both work

Exploiting the Stack

Buffer-Overflow Attacks

String Library Code

Implementation of Unix function gets()

```
/* Get string from stdin */
char *gets(char *dest)
{
   int c = getchar();
   char *p = dest;
   while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
   }
   *p = '\0';
   return dest;
}
```

- no way to specify limit on number of characters to read
- Similar problems with other library functions
 - strcpy, strcat: copy strings of arbitrary length
 - scanf, fscanf, sscanf, when given %s conversion specification

Vulnerable Buffer Code

```
/* Echo Line */
void echo()
{
    char buf[4];    /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```
int main() {
    echo();

return 0;
}
```

```
unix>./echo
123
123
```

```
unix>./echo
123456789ABCDEF01234567
123456789ABCDEF01234567
```

```
unix>./echo
123456789ABCDEF012345678
Segmentation Fault
```

Buffer-Overflow Disassembly

echo:

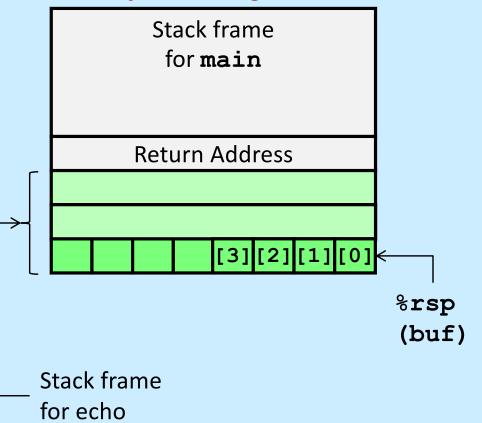
```
000000000040054c <echo>:
 40054c:
               48 83 ec 18
                                sub
                                      $0x18,%rsp
 400550:
               48 89 e7
                                      %rsp,%rdi
                               mov
 400553:
               e8 d8 fe ff ff
                               callq
                                      400430 <gets@plt>
 400558:
               48 89 e7
                                      %rsp,%rdi
                               mov
 40055b:
               e8 b0 fe ff ff
                               callq
                                      400410 <puts@plt>
 400560:
               48 83 c4 18
                               add
                                      $0x18,%rsp
 400564:
               c3
                                retq
```

main:

```
0000000000400565 <main>:
 400565:
               48 83 ec 08
                               sub
                                      $0x8,%rsp
 400569:
               b8 00 00 00 00
                                      $0x0, %eax
                               mov
 40056e:
               e8 d9 ff ff ff
                                      40054c <echo>
                               callq
 400573:
               b8 00 00 00 00
                                      $0x0, %eax
                               mov
              48 83 c4 08
 400578:
                               add
                                      $0x8,%rsp
 40057c:
               c3
                               retq
```

Buffer-Overflow Stack

Before call to gets



```
/* Echo Line */
void echo()
{
   char buf[4];  /* Too small! */
   gets(buf);
   puts(buf);
}
```

```
echo:

subq $24, %rsp

movq %rsp, %rdi

call gets

movq %rsp, %rdi

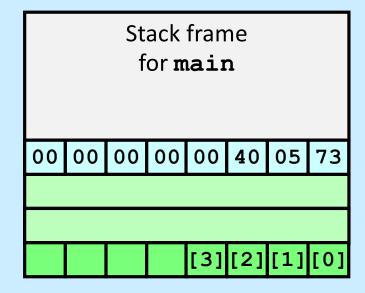
call puts

addq $24, %rsp

ret
```

Buffer Overflow Stack Example

unix> gdb echo
(gdb) break echo
Breakpoint 1 at 0x40054c
(gdb) run
Breakpoint 1, 0x000000000040054c in echo ()
(gdb) print /x \$rsp
\$1 = 0x7fffffffe988
(gdb) print /x *(unsigned *)\$rsp
\$2 = 0x400573



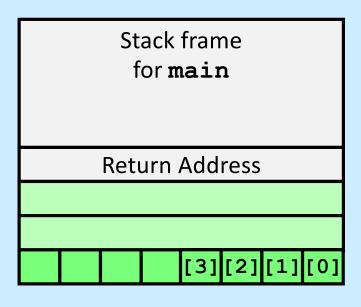
40056e: e8 d9 ff ff ff callq 40054c <echo>

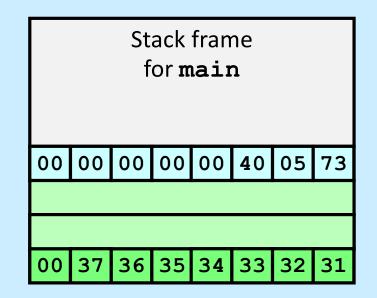
400573: b8 00 00 00 mov \$0x0, %eax

Buffer Overflow Example #1

Before call to gets

Input 1234567





Overflow buf, but no problem

40056e: e8 d9 ff ff ff callq 40054c <echo>

400573: b8 00 00 00 mov \$0x0, %eax

Buffer Overflow Example #2

Before call to gets

Stack frame for main Return Address [3][2][1][0]

Input 123456789ABCDEF01234567

	Stack frame for main									
00	00	00	00	00	40	05	73			
00	37	36	35	34	33	32	31			
30	46	45	44	43	42	41	39			
38	37	36	35	34	33	32	31			

Still no problem

40056e: e8 d9 ff ff ff callq 40054c <echo>

400573: b8 00 00 00 mov \$0x0, %eax

Buffer Overflow Example #3

Before call to gets

Stack frame for main Return Address [3][2][1][0]

Input 123456789ABCDEF012345678

	Stack frame for main									
Ì	00	00	00	00	00	40	05	00		
ĺ	38	37	36	35	34	33	32	31		
	30	46	45	44	43	42	41	39		
	38	37	36	35	34	33	32	31		

Return address corrupted

40056e: e8 d9 ff ff ff callq 40054c <echo>

400573: b8 00 00 00 00 mov \$0x0, %eax

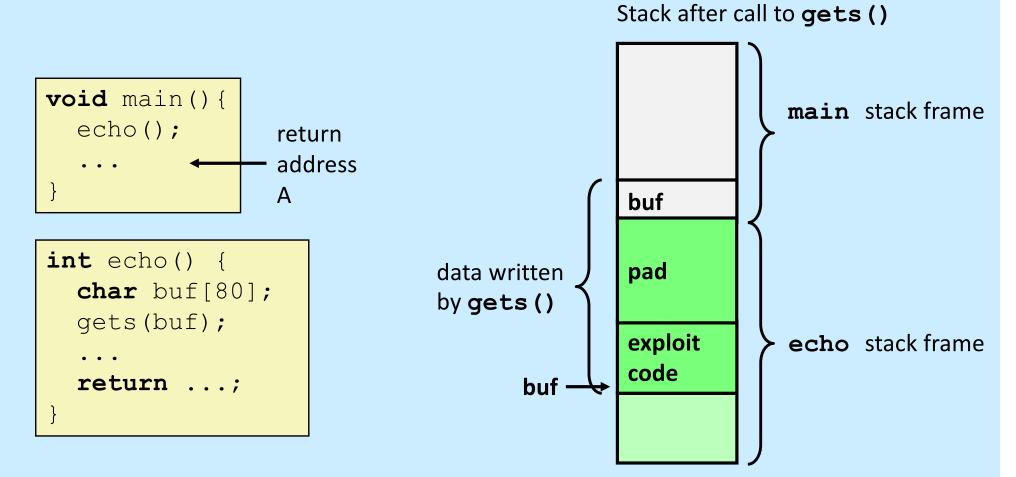
Avoiding Overflow Vulnerability

```
/* Echo Line */
void echo()
{
   char buf[4]; /* Way too small! */
   fgets(buf, 4, stdin);
   puts(buf);
}
```

Use library functions that limit string lengths

- fgets instead of gets
- strncpy instead of strcpy
- don't use scanf with %s conversion specification
 - » use fgets to read the string
 - » or use %ns where n is a suitable integer

Malicious Use of Buffer Overflow



- Input string contains byte representation of executable code
- Overwrite return address A with address of buffer buf
- When echo() executes ret, will jump to exploit code

```
int main() {
        char buf[80];
        gets (buf);
        puts(buf);
        return 0;
main:
  subq $88, %rsp # grow stack
 movq %rsp, %rdi # setup arg
  call gets
 movq %rsp, %rdi # setup arg
  call puts
 movl $0, %eax # set return value
  addq $88, %rsp # pop stack
  ret
```

previous frame return address **Exploit**

Crafting the Exploit ...

- Code + padding
 - 96 bytes long
 - » 88 bytes for buf
 - » 8 bytes for return address

Code (in C):

previous frame

return address

buf (88 bytes)

Quiz 3

The exploit code will be read into memory starting at location 0x7ffffffe948. What value should be put into the return-address portion of the stack frame?

previous frame

0x7ffffffe9a0 return address

buf (88 bytes)

- a)
- 0x7fffffffe9a0
- 0x7fffffffe948
- it doesn't matter what value goes there

0x7fffffffe948