CS 33

Multithreaded Programming IV

Quiz 1

```
void long_running_procedure()
{
   pthread_mutex_lock(&m);
   state = function(state);
   pthread_mutex_unlock(&m);
}
void display(state_t *statep)
{
   pthread_mutex_lock(&m);
   print_state(statep)
   pthread_mutex_unlock(&m);
}
```

long_running_procedure is run by the main thread; display is run by the thread that is handling signals (via sigwait). Is there a potential deadlock resulting from their use of mutexes?

- a) No, since the functions are run by separate threads
- b) Yes, since display is called in response to a signal and thus uses the same stack as does the call to long_running_procedure

Some Thread Gotchas ...

- Exit vs. pthread_exit
- Handling multiple arguments

Worker Threads

```
int main() {
  pthread_t thread[10];
  for (int i=0; i<10; i++)
    pthread_create(&thread[i], 0,
         worker, (void *)i);
  return 0;
}</pre>
```

Better Worker Threads

```
int main() {
  pthread_t thread[10];
  for (int i=0; i<10; i++)
    pthread_create(&thread[i], 0,
         worker, (void *)i);
  pthread_exit(0);
}</pre>
```

Multiple Arguments

```
void relay(int left, int right) {
 pthread t LRthread, RLthread;
  pthread create (&LRthread,
     0,
      copy,
      left, right); // Can't do this ...
  pthread create (&RLthread,
      0,
      copy,
      right, left);
                     // Can't do this ...
```

Multiple Arguments

```
Quiz 2
typedef struct args
                       Does this work?
  int src;
                          a) yes
  int dest;
                          b) no
} args t;
void relay(int left, int right) {
  args t LRargs, RLargs;
  pthread t LRthread, RLthread;
  pthread create (&LRthread, 0, copy, &LRargs);
  pthread create (&RLthread, 0, copy, &RLargs);
  pthread join (LRthread, 0);
  pthread join (RLthread, 0);
```

Multiple Arguments

```
struct 2args {
   int src;
   int dest;
} args;
```

```
Quiz 3

Does this work?

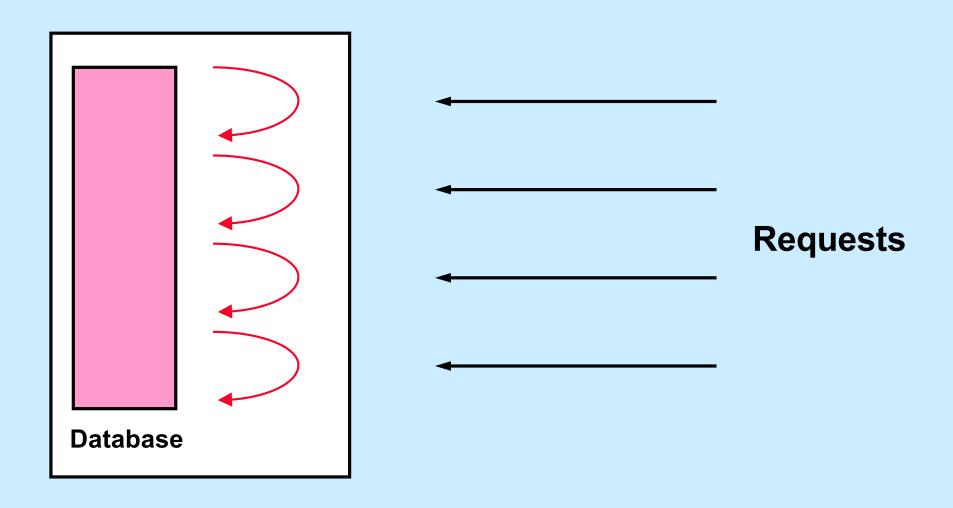
a) yes
b) no
```

```
void relay(int left, int right) {
  pthread_t LRthread, RLthread;
  args.src = left; args.dest = right;
  pthread_create(&LRthread, 0, copy, &args);
  args.src = right; args.dest = left;
  pthread_create(&RLthread, 0, copy, &args);
}
```

Cancellation



Multithreaded Database Server



Sample Code

```
void *thread code(void *arg) {
  node t *head = 0;
  while (1) {
    node t *nodep;
    nodep = (node t *) malloc(sizeof(node t));
    nodep->next = head;
    head = nodep;
    if (read(0, &node->value,
        sizeof(node->value))
                               pthread cancel(thread);
      free (nodep);
      break;
  return head;
```

Quiz 2

```
void *thread code(void *arg) {
      node t *head = 0;
3
      while (1) {
        node t *nodep;
4
        nodep = (node t *) malloc(size
        nodep->next = head;
6
        head = nodep;
        if (read(0, &node->value,
            sizeof(node->value)) == 0) {
          free (nodep);
10
          break;
11
12
13
  return head;
14
```

Where is it safe to terminate a thread within thread_code?

- a) At all lines
- b) At all lines other than 5 and 9
- c) At all lines other than 8
- d) At all lines other than 5, 8, and 9
- e) At no lines

Cancellation Concerns

- Getting cancelled at an inopportune moment
- Cleaning up

Cancellation State

Pending cancel

```
- pthread cancel (thread)
```

Cancels enabled or disabled

```
- int pthread_setcancelstate(
     {PTHREAD_CANCEL_DISABLE
     PTHREAD_CANCEL_ENABLE},
     &oldstate)
```

Asynchronous vs. deferred cancels

```
- int pthread_setcanceltype(
     {PTHREAD_CANCEL_ASYNCHRONOUS,
     PTHREAD_CANCEL_DEFERRED),
     &oldtype)
```

Sample Code – Cancellation Point

```
void *thread code(void *arg) {
  node t *head = 0;
  while (1) {
    node t *nodep;
    nodep = (node t *) malloc(sizeof(node t));
    nodep->next = head;
    head = nodep;
    if (read(0, &node->value,
        sizeof(node->value)) == 0) {
      free (nodep);
      break;
  return head;
```

Cancellation Points

- aio_suspend
- close
- creat
- fcntl (when F_SETLCKW is the command)
- fsync
- mq_receive
- mq_send
- msync
- nanosleep
- open
- pause
- pthread_cond_wait
- pthread_cond_timedwait
- pthread_join

- pthread testcancel
- read
- sem_wait
- sigwait
- sigwaitinfo
- sigsuspend
- sigtimedwait
- sleep
- system
- tcdrain
- wait
- waitpid
- write

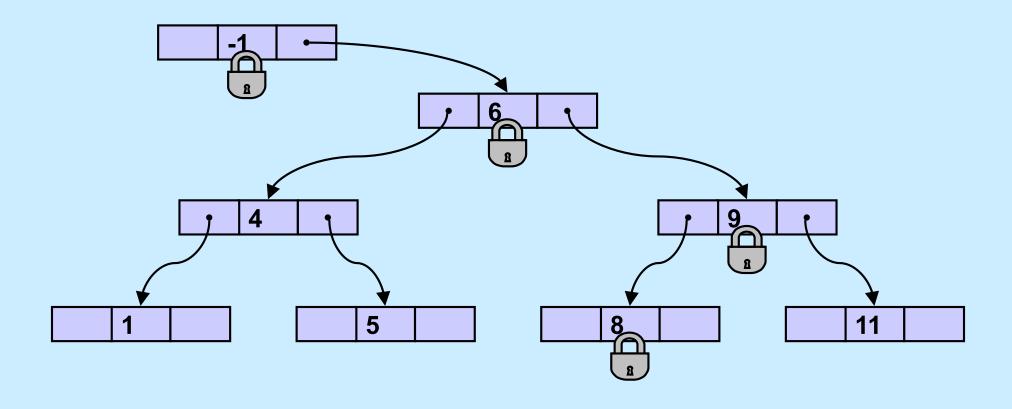
Cleaning Up

- void pthread_cleanup_push((void)(*routine)(void *),void *arg)
- void pthread cleanup pop(int execute)

Sample Code, Revisited

```
void *thread code(void *arg) {
                                          void cleanup(void *arg) {
  node t *head = 0;
                                            node t **headp = arg;
  pthread cleanup push (
                                            while(*headp) {
      cleanup, &head);
                                              node t *nodep = head->next;
                                              free (*headp);
  while (1) {
    node t *nodep;
                                              *headp = nodep;
    nodep = (node t *)
       malloc(sizeof(node t));
    nodep->next = head;
    head = nodep;
    if (read(0, &nodep->value,
        sizeof(nodep->value)) == 0) {
      free (nodep);
      break;
  pthread cleanup pop(0);
  return head;
```

A More Complicated Situation ...



Start/Stop





Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait(&s->queue, &s->mutex);
  pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```

Start/Stop

Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait (&s->queue,
      &s->mutex);
  pthread mutex unlock (&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```



Not a Quiz

You're in charge of designing POSIX threads. Should *pthread_cond_wait* be a cancellation point?

- a) no
- b) yes; cancelled threads must acquire mutex before invoking cleanup handler
- c) yes; but they don't acquire mutex

Cancellation and Conditions

```
pthread_mutex_lock(&m);
pthread_cleanup_push(cleanup_handler, &m);
while(should_wait)
   pthread_cond_wait(&cv, &m);

read(0, buffer, len); // read is a cancellation point
pthread_cleanup_pop(1);
```

Quiz 3

Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  pthread cleanup push (
    cleanup func, cleanup arg);
  while(s->state == stopped)
    pthread cond wait (&s->queue, &s->mutex);
  pthread cleanup pop(1);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast (&s->queue);
  pthread mutex unlock(&s->mutex);
```

What should be used for cleanup_func and cleanup_arg?

- a) pthread_mutex_unlock
 and &s->mutex
- b) that and more
- c) there's no need for a cleanup function

Quiz 4

Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  pthread cleanup push (
    cleanup func, cleanup arg);
  while(s->state == stopped)
    pthread cond wait (&s->queue, &s->mutex);
  pthread cleanup pop(1);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast (&s->queue);
  pthread mutex unlock(&s->mutex);
```

What should be used for cleanup_func and cleanup_arg?

- a) pthread_mutex_unlock
 and &s->mutex
- b) that and more
- c) there's no need for a cleanup function

A Problem ...

In thread 1:

In thread 2:

There's only one errno!

However, somehow it works.

What's done???

A Solution ...

```
#define errno (* errno location())
```

- __errno_location returns an int * that's different for each thread
 - thus each thread has, effectively, its own copy of errno

Process Address Space

Stack, etc. Thread 1

Stack, etc. Thread 2

Stack, etc. Thread 3

Dynamic

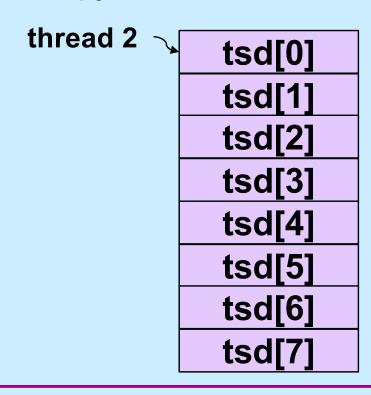
Data

Text

Generalizing

- Thread-specific data (sometimes called thread-local storage)
 - data that's referred to by global variables, but each thread has its own private copy

48 8 4 .	
thread 1	tsd[0]
	tsd[1]
	tsd[2]
	tsd[3]
	tsd[4]
	tsd[5]
	tsd[6]
	tsd[7]



Some Machinery

- pthread_key_create(&key, cleanup_routine)
 - allocates a slot in the TSD arrays
 - provides a function to cleanup when threads terminate
- value = pthread_getspecific(key)
 - fetches from the calling thread's array
- pthread_setspecific(key, value)
 - stores into the calling thread's array

errno (Again)

```
// executed before threads are created
pthread_key_t errno_key;
pthread_key_create(&errno_key, NULL);

// redefine errno to use thread-specific value
#define errno pthread_getspecific(errno_key);

// set current thread's errno
pthread_set_specific(errno_key, (void *)ENOMEM);
```

Beyond POSIX TLS Extensions for ELF and gcc

Thread Local Storage (TLS)

```
__thread int x=6;

// Each thread has its own copy of x,

// each initialized to 6.

// Linker and compiler do the setup.

// May be combined with static or extern.

// Doesn't make sense for local variables!
```

Example: Per-Thread Windows

```
void *tfunc(void * arg) {
                                          getWindow();
typedef struct {
  wcontext t win context;
                                          threadWrite("started");
  int file descriptor;
} win t;
  thread static win t my win;
                                          func2 (...);
void getWindow() {
  my win.win context = ...;
  my win.file decriptor = ...;
                                        void func2(...) {
int threadWrite(char *buf) {
  int status = write to window(
                                          threadWrite(
      &my win, buf);
                                               "important msg");
  return (status);
```

Static Local Storage and Threads

```
char *strtok(char *str, const char *delim) {
    static char *saveptr;

    ... // find next token starting at either
    ... // str or saveptr
    ... // update saveptr

    return(&token);
}
```

Coping

- Use thread local storage
- Allocate storage internally; caller frees it
- Redesign the interface

Thread-Safe Version

Shared Data

Thread 1:

```
printf("goto statement reached");
```

Thread 2:

```
printf("Hello World\n");
```

Printed on display:

go to Hell

Coping

- Wrap library calls with synchronization constructs
- Fix the libraries

Efficiency

- Standard I/O example
 - getc() and putc()
 - » expensive and thread-safe?
 - » cheap and not thread-safe?
 - two versions
 - » getc() and putc()
 - expensive and thread-safe
 - » getc unlocked() and putc unlocked()
 - cheap and not thread-safe
 - made thread-safe with flockfile() and funlockfile()

Efficiency

Naive

```
for (i=0; i<lim; i++)</pre>
 putc(out[i]);
```

Efficient

```
flockfile(stdout);
for(i=0; i<lim; i++)
 putc unlocked(out[i]);
funlockfile(stdout);
```

What's Thread-Safe?

Everything except

asctime()	ecvt()	gethostent()	getutxline()	putc_unlocked()
basename()	encrypt()	getlogin()	gmtime()	putchar_unlocked()
catgets()	endgrent()	getnetbyaddr()	hcreate()	putenv()
crypt()	endpwent()	getnetbyname()	hdestroy()	pututxline()
ctime()	endutxent()	getnetent()	hsearch()	rand()
dbm_clearerr()	fcvt()	getopt()	inet_ntoa()	readdir()
dbm_close()	ftw()	getprotobyname()	l64a()	setenv()
dbm_delete()	gcvt()	getprotobynumber()	lgamma()	setgrent()
dbm_error()	getc_unlocked()	getprotoent()	lgammaf()	setkey()
dbm_fetch()	getchar_unlocked()	getpwent()	lgammal()	setpwent()
dbm_firstkey()	getdate()	getpwnam()	localeconv()	setutxent()
dbm_nextkey()	getenv()	getpwuid()	localtime()	strerror()
dbm_open()	getgrent()	getservbyname()	Irand48()	strtok()
dbm_store()	getgrgid()	getservbyport()	mrand48()	ttyname()
dirname()	getgrnam()	getservent()	nftw()	unsetenv()
dlerror()	gethostbyaddr()	getutxent()	nl_langinfo()	wcstombs()
drand48()	gethostbyname()	getutxid()	ptsname()	wctomb()