**CS 33** 

**Multithreaded Programming VII** 

#### **Implementing Mutexes**

#### Strategy

- make the usual case (no waiting) very fast
- can afford to take more time for the other case (waiting for the mutex)

#### **Futexes**

- Safe, efficient kernel conditional queueing in Linux
- All operations performed atomically

» otherwise return

```
- futex wake(futex t *futex)
```

» wake up one thread from futex's wait queue, if there are any waiting threads

## **Ancillary Functions**

```
• int atomic inc(int *val)

    add 1 to *val, return its original value

• int atomic dec(int *val)

    subtract 1 from *val, return its original value

• int CAS(int *ptr, int old, int new) {
      int tmp = *ptr;
      if (*ptr == old)
          *ptr = new;
      return tmp;
```

#### Attempt 1

```
void lock(futex_t *futex) {
  int c;
  while ((c = atomic_inc(&futex->val)) != 0)
    futex_wait(futex, c+1);
}

void unlock(futex_t *futex) {
  futex->val = 0;
  futex_wake(futex);
}
```

#### Attempt 2

```
void lock(futex_t *futex) {
  int c;
  if ((c = CAS(&futex->val, 0, 1) != 0)
    do {
      if (c == 2 || (CAS(&futex->val, 1, 2) != 0))
        futex wait(futex, 2);
    while ((c = CAS(\&futex->val, 0, 2)) != 0))
void unlock(futex t *futex) {
  if (atomic dec(&futex->val) != 1) {
    futex->val = 0;
    futex wake(futex);
```

# **Memory Allocation**

Multiple threads

One heap

**Bottleneck?** 

#### **Solution 1**

- Divvy up the heap among the threads
  - each thread has its own heap
  - no mutexes required
  - no bottleneck
- How much heap does each thread get?

#### **Solution 2**

- Multiple "arenas"
  - each with its own mutex
  - thread allocates from the first one it can find whose mutex was unlocked
    - » if none, then creates new one
  - deallocations go back to original arena

#### Solution 3

- Global heap plus per-thread heaps
  - threads pull storage from global heap
  - freed storage goes to per-thread heap
    - » unless things are imbalanced
      - then thread moves storage back to global heap
  - mutex on only the global heap
- What if one thread allocates and another frees storage?

## Malloc/Free Implementations

- ptmalloc
  - based on solution 2
  - in glibc (i.e., used by default)
- tcmalloc
  - based on solution 3
  - from Google
- Which is best?

## **Test Program**

```
const unsigned int N=64, nthreads=32, iters=10000000;
int main() {
  void *tfunc(void *);
 pthread t thread[nthreads];
  for (int i=0; i<nthreads; i++) {
    pthread create(&thread[i], 0, tfunc, (void *)i);
    pthread detach(thread[i]);
 pthread exit(0);
void *tfunc(void *arg) {
  long i;
  for (i=0; i<iters; i++) {
    long *p = (long *) malloc(sizeof(long) * ((i%N) +1));
    free(p);
  return 0;
```

#### **Not a Quiz**

Which is fastest?

- a) glibc (i.e., standard Linux)
- b) Google

#### Compiling It ...

```
% gcc -o ptalloc alloc.c -lpthread
% gcc -o tcalloc alloc.c -lpthread -ltcmalloc
```

## Running It (2014) ...

```
$ time ./ptalloc
real 0m5.142s
user 0m20.501s
sys 0m0.024s
$ time ./tcalloc
real 0m1.889s
user 0m7.492s
sys 0m0.008s
```

# Running It (2023) ...

```
$ time ./ptalloc
real 0m0.666s
user 0m5.815s
sys 0m0.004s
$ time ./tcalloc
real 0m0.496s
user 0m4.197s
sys 0m0.008s
```

## What's Going On (2014)?

```
$ strace -c -f ./ptalloc
% time seconds usecs/call calls errors syscall
100.00 0.040002 13 3007 520 futex
$ strace -c -f ./tcalloc
% time seconds usecs/call calls errors syscall
 0.00 0.000000
                0 59 13 futex
```

# What's Going On (2023)?

```
$ strace -c -f ./ptalloc
% time seconds usecs/call calls errors syscall
                       5 1 futex
 0.02 0.000001
$ strace -c -f ./tcalloc
% time seconds usecs/call calls errors syscall
0.26 0.000006 0 23 3 futex
```

```
#define N 64
#define npairs 16
#define allocsPerIter 1024
const long iters = 8*1024*1024/allocsPerIter;
#define BufSize 10240
typedef struct buffer {
  int *buf[BufSize];
  unsigned int nextin;
  unsigned int nextout;
  sem t empty;
  sem t occupied;
 pthread t pthread;
 pthread t cthread;
} buffer t;
```

```
int main() {
  long i;
 buffer t b[npairs];
  for (i=0; i<npairs; i++) {
   b[i].nextin = 0;
   b[i].nextout = 0;
    sem init(&b[i].empty, 0, BufSize/allocsPerIter);
    sem init(&b[i].occupied, 0, 0);
   pthread create(&b[i].pthread, 0, prod, &b[i]);
   pthread create(&b[i].cthread, 0, cons, &b[i]);
  for (i=0; i<npairs; i++) {
   pthread join(b[i].pthread, 0);
   pthread join(b[i].cthread, 0);
  return 0;
```

```
void *prod(void *arg) {
  long i, j;
  buffer t *b = (buffer t *) arg;
  for (i = 0; i<iters; i++) {
    sem wait(&b->empty);
    for (j = 0; j<allocsPerIter; j++) {</pre>
      b->buf[b->nextin] = malloc(sizeof(int)*((j%N)+1));
      if (++b->nextin >= BufSize)
       b->nextin = 0:
    sem post(&b->occupied);
  return 0;
```

```
void *cons(void *arg) {
  long i, j;
  buffer t *b = (buffer t *) arg;
  for (i = 0; i<iters; i++) {
    sem wait(&b->occupied);
    for (j = 0; j<allocsPerIter; j++) {</pre>
      free(b->buf[b->nextout]);
      if (++b->nextout >= BufSize)
       b->nextout = 0;
    sem post(&b->empty);
  return 0;
```

## Running It (2014) ...

```
$ time ./ptalloc2
real 0m1.087s
user 0m3.744s
sys 0m0.204s
$ time ./tcalloc2
real 0m3.535s
user 0m11.361s
sys 0m2.112s
```

## Running It (2023) ...

```
$ time ./ptalloc2
real 0m0.365s
user 0m1.378s
sys 0m0.536s
$ time ./tcalloc2
real 0m8.019s
user 1m1.348s
sys 0m7.161s
```

# What's Going On (2014)?

```
$ strace -c -f ./ptalloc2
% time seconds usecs/call calls errors syscall
93.04 8.246196 117 70173 20775 futex
$ strace -c -f ./tcalloc2
% time seconds usecs/call calls errors syscall
99.92 47.796676 153 311012 7244 futex
```

# What's Going On (2023)?

```
$ strace -c -f ./ptalloc2
% time seconds usecs/call calls errors syscall
98.48 42.883196 79 539757 179723 futex
$ strace -c -f ./tcalloc2
% time seconds usecs/call calls errors syscall
99.99 346.746205 146 2372684 44547 futex
```

#### You'll Soon Finish CS 33 ...

- You might
  - celebrate



- take another systems course
  - » 320
  - » 1380
  - » 1660
  - » 1670
  - » 1680



become a 33 TA



#### **Systems Courses Next Semester**

- CS 320 (Intro to Software Engineering)
  - you've mastered low-level systems programming
  - now do things at a higher level
  - learn software-engineering techniques using Java, XML, etc.
- CS 1380 (Distributed Systems)
  - you now know how things work on one computer
  - what if you've got lots of computers?
  - some may have crashed, others may have been taken over by your worst (and smartest) enemy
- CS 1660/1620/2660 (Computer Systems Security)
  - liked buffer?
  - you'll really like 1660
- CS 1670/1690/2670 (Operating Systems)
  - still mystified about what the OS does?
  - write your own!

# The End

Well, not quite ...

Database is due on 12/15

**Happy Coding and Happy Holidays!**