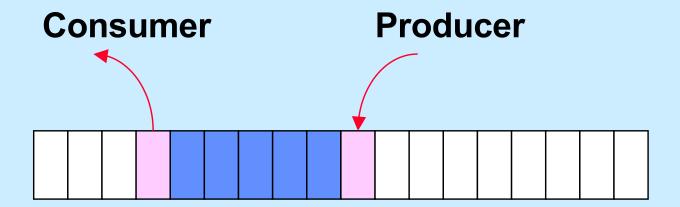
CS 33

Multithreaded Programming III

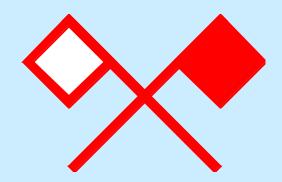
Producer-Consumer Problem



Guarded Commands

```
when (guard) [
 /*
    once the guard is true, execute this
    code atomically
   * /
```

Semaphores



P(S) operation:

V(S) operation:

$$[S = S + 1;]$$

Quiz 1

```
semaphore S = 1;
int count = 0;
void func() {
  P(S);
  count++;
  count--;
  V(S);
```

The function func is called concurrently by n threads. What's the maximum value that count will take on?

- a) 1
- b) 2
- c) n
- d) indeterminate
- P(S) operation:

```
when (S > 0) [
S = S - 1;
]
```

• V(S) operation:

$$[S = S + 1;]$$

Producer/Consumer with Semaphores

```
Semaphore empty = BSIZE;
             Semaphore occupied = 0;
             int nextin = 0:
             int nextout = 0;
P(empty);
                             char item;
                             P(occupied);
 buf[nextin] = item;
 if (++nextin >= BSIZE)
                             item = buf[nextout];
   nextin = 0;
                             if (++nextout >= BSIZE)
                               nextout = 0;
 V (occupied);
                             V(empty);
                             return item;
```

POSIX Semaphores

```
#include <semaphore.h>
int sem_init(sem_t *semaphore, int pshared, int init);
int sem_destroy(sem_t *semaphore);
int sem_wait(sem_t *semaphore);
    /* P operation */
int sem_trywait(sem_t *semaphore);
    /* conditional P operation */
int sem_post(sem_t *semaphore);
    /* V operation */
```

Producer-Consumer with POSIX Semaphores

```
sem init(&empty, 0, BSIZE);
             sem init(&occupied, 0, 0);
             int nextin = 0;
             int nextout = 0;
void produce(char item) {         char consume() {
                                char item;
  sem wait(&empty);
                                 sem wait (&occupied);
                                item = buf[nextout];
 buf[nextin] = item;
  if (++nextin >= BSIZE)
                                if (++nextout >= BSIZE)
   nextin = 0;
                                  nextout = 0;
  sem post(&occupied);
                                 sem post(&empty);
                                return item;
```





```
void wait_for_start(state_t *s);

void start(state_t *s);

void stop(state t *s);
```





```
void wait for start(state t *s) {
  if (s->state == stopped)
    sleep();
void start(state t *s) {
  state = started;
  wakeup all();
void stop(state t *s) {
  state = stopped;
```



```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  if (s->state == stopped) {
    pthread mutex unlock(&s->mutex);
    sleep();
  else pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  state = started;
  wakeup all();
  pthread mutex unlock(&s->mutex);
```



```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  if (s->state == stopped) {
    sleep();
  pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  state = started;
  wakeup all();
  pthread mutex unlock(&s->mutex);
```





```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait(&s->queue, &s->mutex);
  pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```

Condition Variables

```
when (guard) [
                                    pthread mutex lock(&mutex);
  statement 1;
                                    while(!quard)
                                      pthread cond wait (
                                         &cond var, &mutex);
  statement n;
                                    statement 1;
                                    statement n;
                                    pthread mutex unlock (&mutex);
// code modifying the guard:
                                    pthread mutex lock(&mutex);
                                    // code modifying the guard:
                                    pthread cond broadcast (
                                         &cond var);
                                    pthread mutex unlock(&mutex);
```

Set Up

PC with Condition Variables (1)

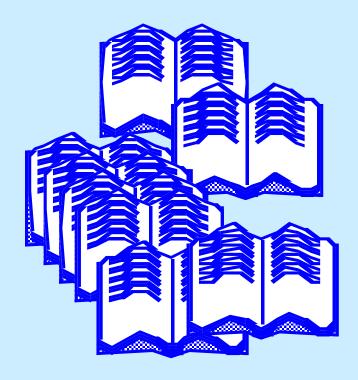
```
typedef struct buffer {
   pthread_mutex_t m;
   pthread_cond_t more_space;
   pthread_cond_t more_items;
   int next_in;
   int next_out;
   int empty;
   char buf[BSIZE];
}
```

PC with Condition Variables (2)

```
void produce(buffer_t *b,
    char item) {
  pthread mutex lock(&b->m);
  while (!(b->empty > 0))
   pthread cond wait (
       \&b->more space, \&b->m);
  b->buf[b->nextin] = item;
  if (++(b->nextin) == BSIZE)
    b->nextin = 0;
  b->empty--;
  pthread cond signal (
     &b->more items);
  pthread mutex unlock(&b->m);
```

```
char consume(buffer t *b) {
  char item;
  pthread mutex lock(&b->m);
  while (!(b->empty < BSIZE))</pre>
   pthread cond wait (
       &b->more items, &b->m);
  item = b->buf[b->nextout];
  if (++(b->nextout) == BSIZE)
    b->nextout = 0;
  b->empty++;
  pthread cond signal (
     &b->more space);
  pthread mutex unlock(&b->m);
  return item;
```

Readers-Writers Problem





Pseudocode

```
reader() {
    when (writers == 0) [
      readers++;
    ]

    /* read */
    [readers--;]
}
```

Pseudocode with Assertions

```
writer() {
 when ((writers == 0) &&
     (readers == 0)) [
   writers++;
  assert((readers == 0) &&
     (writers == 1));
  /* write */
  [writers--;]
```

Solution with POSIX Threads

```
reader() {
 pthread mutex lock(&m);
 while (!(writers == 0))
    pthread cond wait (
        &readersQ, &m);
  readers++;
 pthread mutex unlock(&m);
  /* read */
 pthread mutex lock(&m);
 if (--readers == 0)
    pthread cond signal (
        &writers();
 pthread mutex unlock(&m);
```

```
writer() {
  pthread mutex lock(&m);
  while(!((readers == 0) &&
      (writers == 0))
   pthread cond wait (
       &writersQ, &m);
  writers++;
  pthread mutex unlock(&m);
  /* write */
  pthread mutex lock(&m);
  writers--;
  pthread cond signal (
       &writersO);
  pthread cond broadcast (
     &readers();
  pthread mutex unlock (&m);
```

New Pseudocode

```
writer() {
reader() {
 when (writers == 0) [
                               [writers++;]
                               when ((readers == 0) &&
   readers++;
                                  (active writers == 0)) [
                                active writers++;
 /* read */
                               /* write */
  [readers--;]
                               [writers--;
                               active writers--;]
```

Improved Reader

```
reader() {
 pthread mutex lock(&m);
                                  pthread mutex lock(&m);
 while (!(writers == 0)) {
                                  if (--readers == 0)
   pthread cond wait (
                                   pthread cond signal (
        &readersQ, &m);
                                         &writersQ);
 readers++;
                                  pthread mutex unlock (&m);
 pthread mutex unlock(&m);
 /* read */
```

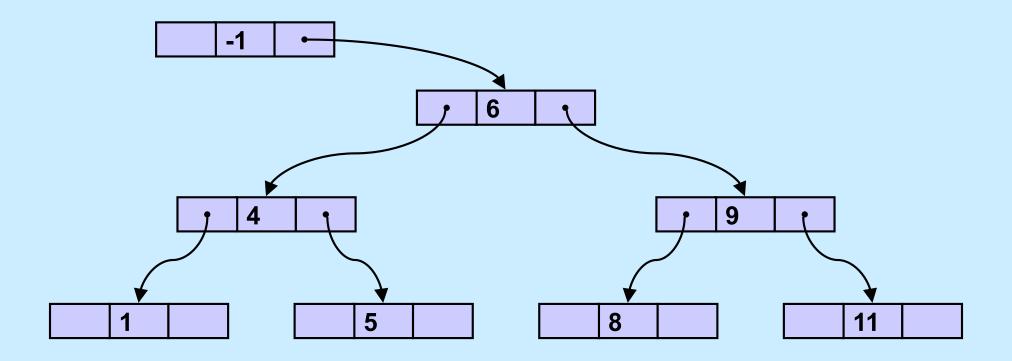
Improved Writer

```
writer() {
  pthread mutex lock(&m);
                                        pthread mutex lock(&m);
                                        writers--;
  writers++;
                                          active writers--;
  while (!((readers == 0) &&
                                        if (writers)
        (active writers == 0))) {
                                         pthread cond signal (
   pthread cond wait (
                                              &writers();
         &writersO, &m);
                                        else
                                         pthread cond broadcast (
  active writers++;
                                               &readersO);
  pthread mutex unlock (&m);
                                        pthread mutex unlock (&m);
  /* write */
```

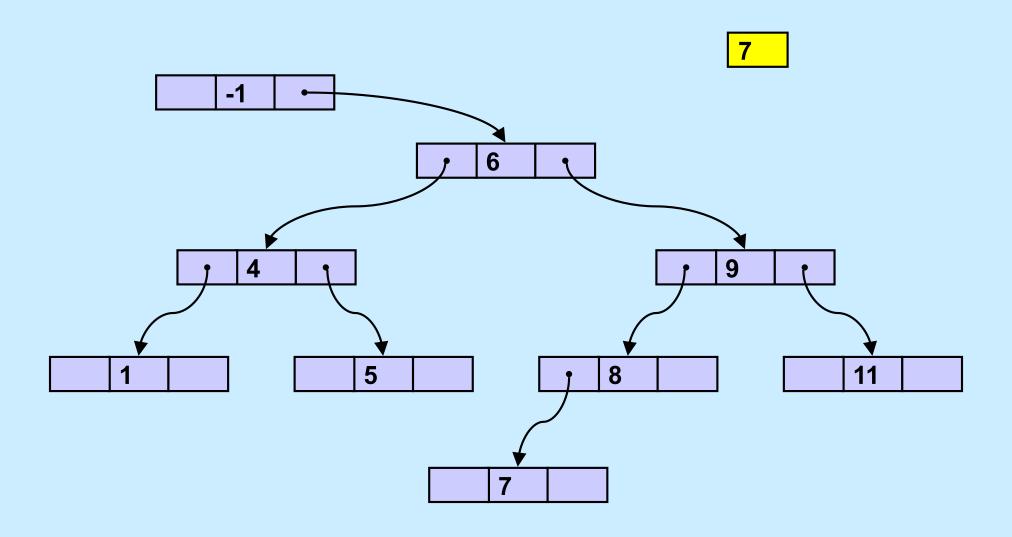
New, From POSIX!

```
int pthread rwlock init (pthread rwlock t *lock,
      pthread rwlockattr t *att);
int pthread rwlock destroy(pthread rwlock t *lock);
int pthread rwlock rdlock (pthread rwlock t *lock);
int pthread rwlock wrlock(pthread rwlock t *lock);
int pthread rwlock tryrdlock(pthread rwlock t *lock);
int pthread rwlock trywrlock(pthread rwlock t *lock);
int pthread timedrwlock rdlock (pthread rwlock t *lock,
      struct timespec *ts);
int pthread timedrwlock wrlock (pthread rwlock t *lock,
      struct timespec *ts);
int pthread rwlock unlock(pthread rwlock t *lock);
```

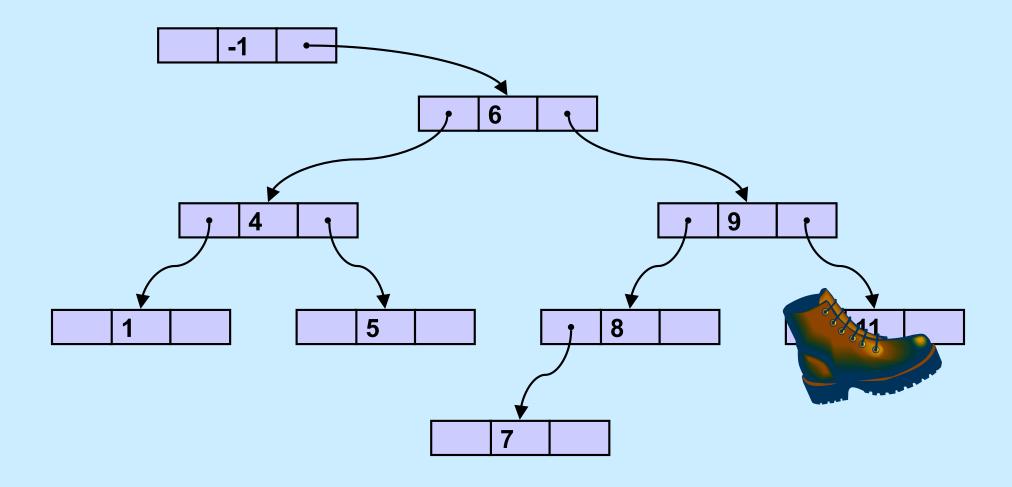
Binary Search Tree



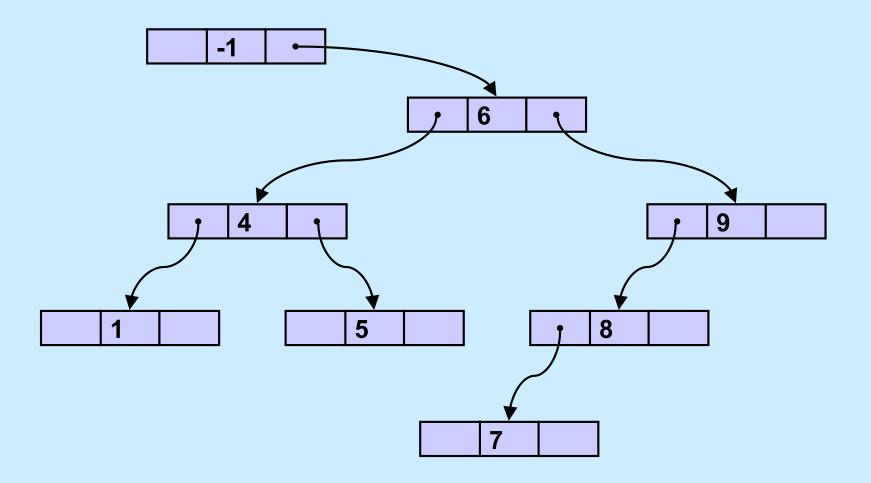
Binary Search Tree: Insertion



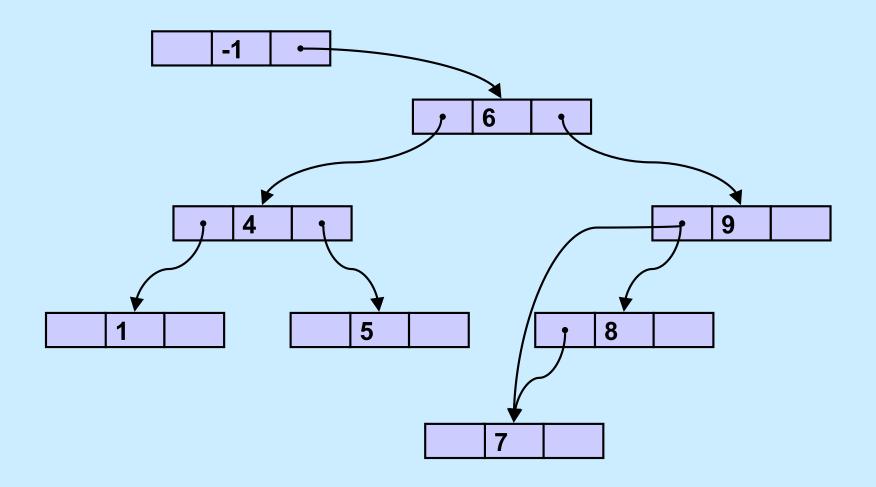
Binary Search Tree: Deletion of Leaf



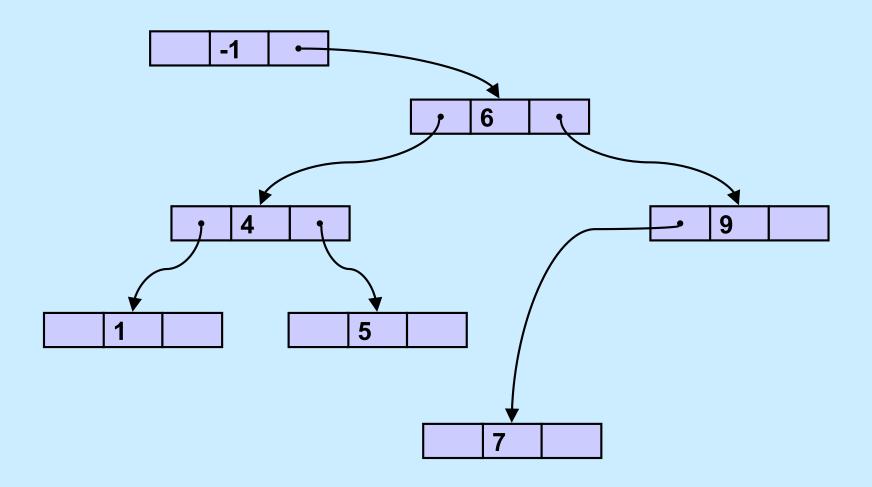
Binary Search Tree: Deletion of Leaf



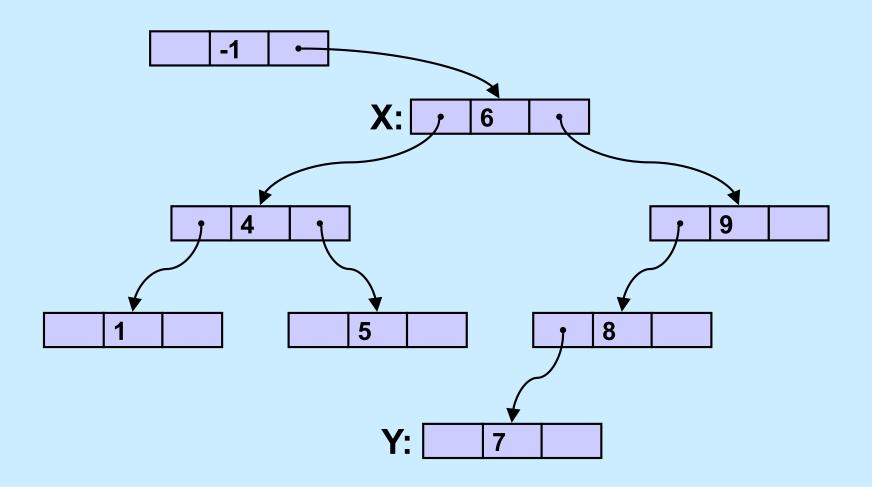
Binary Search Tree: Deletion of Node with One Child



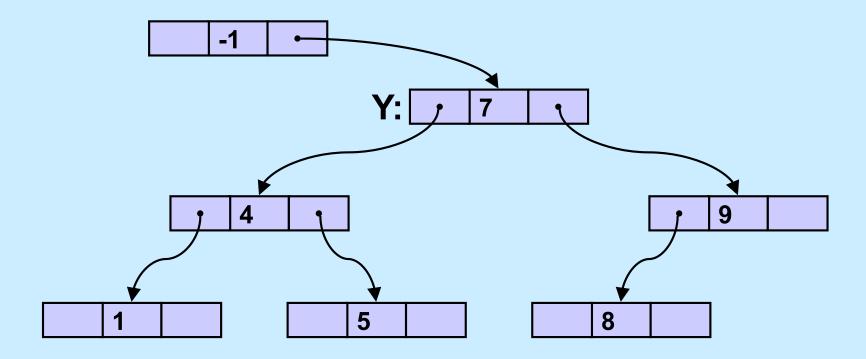
Binary Search Tree: Deletion of Node with One Child



Binary Search Tree: Deletion of Node with Two Children



Binary Search Tree: Deletion of Node with Two Children



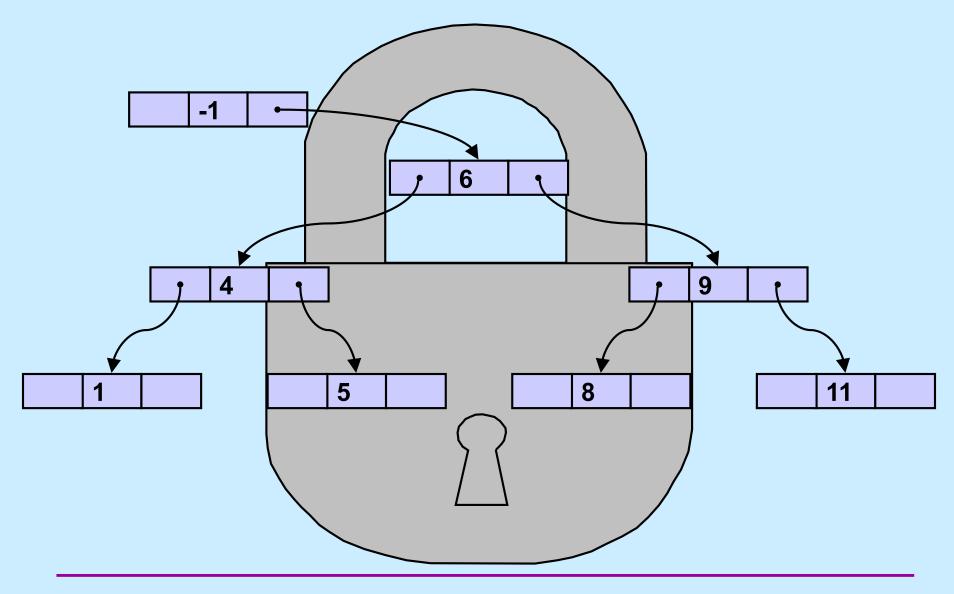
C Code: Search

```
Node *search(int key,
                                        } else {
    Node *parent, Node **parentp) {
                                          if ((next = parent->rchild)
  Node *next;
                                              == 0)
  Node *result;
                                            result = 0;
  if (key < parent->key) {
                                          } else {
    if ((next = parent->lchild)
                                            if (key == next->key) {
        == 0)
                                              result = next;
      result = 0;
                                            } else {
    } else {
                                              result = search(key,
      if (key == next->key) {
                                                  next, parentpp);
        result = next;
                                              return result;
      } else {
        result = search(key,
            next, parentpp);
                                        if (parentpp != 0)
        return result;
                                          *parentpp = parent;
                                        return result;
```

C Code: Add

```
int add(int key) {
 Node *parent, *target, *newnode;
  if ((target = search(key, &head, &parent)) != 0) {
    return 0;
 newnode = malloc(sizeof(Node));
 newnode->key = key;
 newnode->lchild = newnode->rchild = 0;
  if (name < parent->name)
    parent->lchild = newnode;
 else
    parent->rchild = newnode;
  return 1;
```

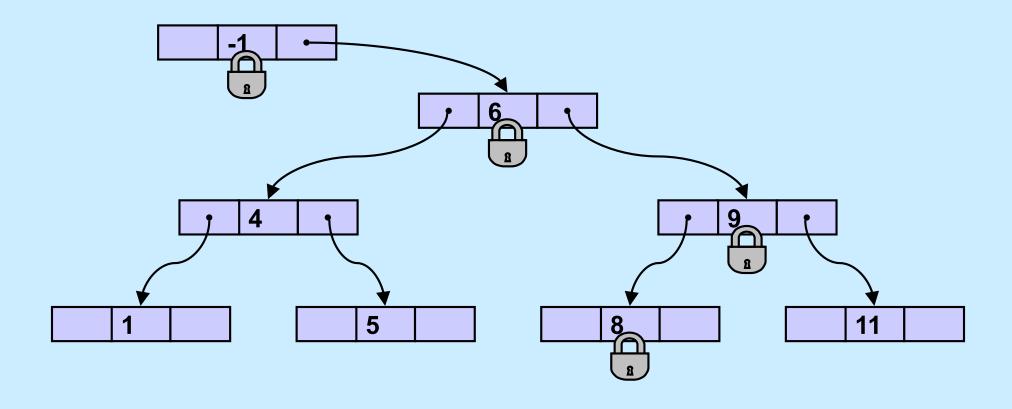
Binary Search Treewith Coarse-Grained Synchronization



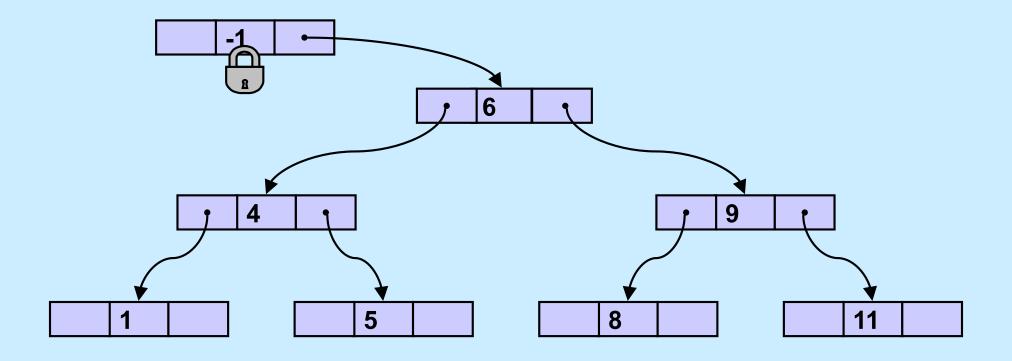
C Code: Add with Coarse-Grained Synchronization

```
int add(int key) {
 Node *parent, *target, *newnode;
 pthread rwlock wrlock (&tree lock);
 if ((target = search(key, &head, &parent)) != 0) {
   pthread rwlock unlock (&tree lock);
   return 0:
 newnode = malloc(sizeof(Node));
 newnode -> key = key;
 newnode->lchild = newnode->rchild = 0;
 if (name < parent->name)
   parent->lchild = newnode;
 else
   parent->rchild = newnode;
 pthread rwlock unlock (&tree lock);
 return 1:
```

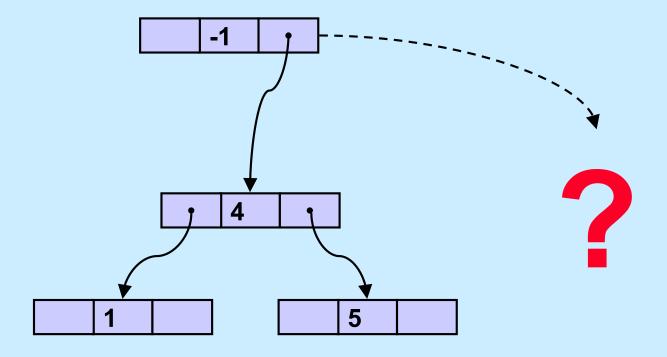
Binary Search Treewith Fine-Grained Synchronization I



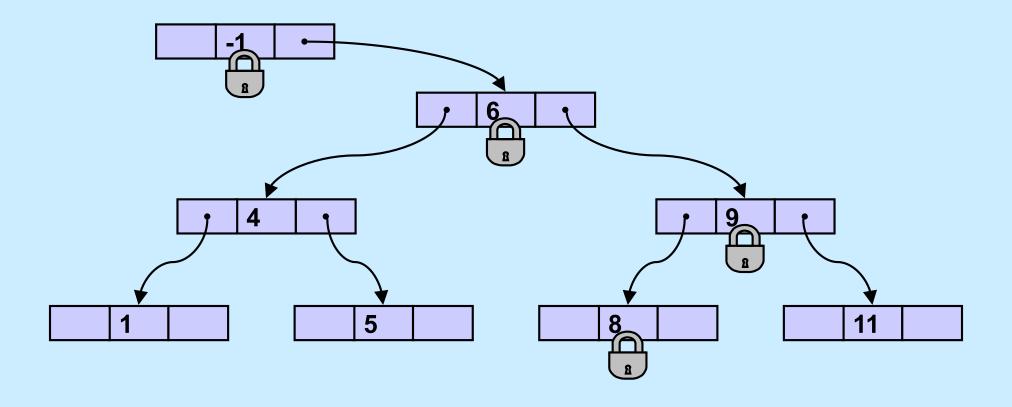
Binary Search Treewith Fine-Grained Synchronization II



Binary Search Treewith Fine-Grained Synchronization III



Doing It Right ...



C Code: Fine-Grained Search I

```
enum locktype {l read, l write};
                                         } else {
                                              lock(lt, &next->lock);
#define lock(lt, lk) ((lt) == l read)?
                                              if (key == next->key) {
      pthread rwlock rdlock(lk):
                                                result = next;
      pthread rwlock wrlock(lk)
                                              } else {
                                                pthread rwlock unlock (
Node *search(int key,
                                                     &parent->lock);
    Node *parent, Node **parentp,
                                                result = search(key,
    enum locktype lt) {
                                                    next, parentpp, lt);
   // parent is locked on entry
                                                 return result;
 Node *next;
 Node *result;
  if (key < parent->key) {
    if ((next = parent->lchild)
        == 0)
      result = 0;
```

C Code: Fine-Grained Search II

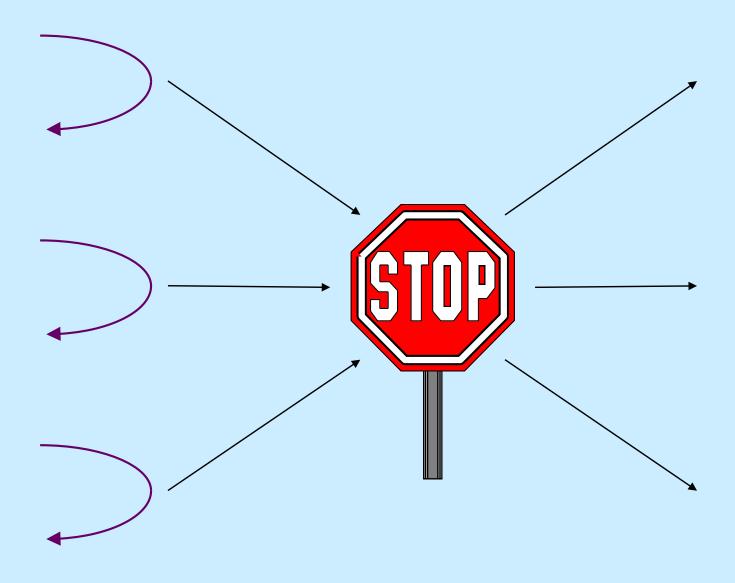
```
} else {
                                           } else {
  if ((next = parent->rchild)
                                            pthread rwlock unlock (
      == 0)
                                                 &parent->lock);
     result = 0;
                                             result = search(key,
   } else {
                                                 next, parentpp, lt);
     lock(lt, &next->lock);
                                            return result;
     if (key == next->key) {
       result = next;
                                      if (parentpp != 0) {
                                        // parent remains locked
                                        *parentpp = parent;
                                      } else
                                        pthread rwlock unlock (
                                             &parent->lock);
                                      return result;
```

C Code: Add with Fine-Grained Synchronization I

C Code: Add with Fine-Grained Synchronization II

```
newnode = malloc(sizeof(Node));
newnode->key = key;
newnode->lchild = newnode->rchild = 0;
pthread_rwlock_init(&newnode->lock, 0);
if (name < parent->name)
   parent->lchild = newnode;
else
   parent->rchild = newnode;
pthread_rwlock_unlock(&parent->lock);
return 1;
```

Barriers



A Solution?

```
pthread mutex lock(&m);
if (++count == number) {
  pthread cond broadcast (&cond var);
} else while (!(count == number)) {
  pthread cond wait (&cond var, &m);
pthread mutex unlock (&m);
```

How About This?

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
   count = 0;
} else while (!(count == number)) {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

And This ...

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
   count = 0;
} else {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

Quiz 2

Does it work?

a) definitely

Barrier in POSIX Threads

```
pthread mutex lock(&m);
if (++count < number) {</pre>
  int my generation = generation;
  while (my generation == generation) {
    pthread cond wait(&waitQ, &m);
 else {
  count = 0;
  generation++;
  pthread cond broadcast (&waitQ);
pthread mutex unlock (&m);
```

More From POSIX!

```
int pthread barrier init(pthread barrier t *barrier,
      pthread barrierattr t *attr,
      unsigned int count);
int pthread barrier destroy(
      pthread barrier t *barrier);
int pthread barrier wait(
      pthread_barrier_t *barrier);
```

Why cond wait is Weird ...

```
pthread cond wait(pthread_cond_t *c, pthread_mutex_t *m) {
   pthread mutex unlock(m);
   sem wait(c->sem);
   pthread mutex lock(m);
pthread cond signal(pthread_cond_t *c) {
   sem post(c->sem);
```