**CS 33** 

**Multithreaded Programming VI** 

#### A Problem ...

In thread 1:

```
if ((ret = open(path,
    O_RDWR) == -1) {
   if (errno == EINTR) {
     ...
}
```

In thread 2:

#### There's only one errno!

However, somehow it works.

What's done???

#### A Solution ...

```
#define errno (* errno location())
```

- \_\_errno\_location returns an int \* that's different for each thread
  - thus each thread has, effectively, its own copy of errno

## **Process Address Space**

Stack, etc. Thread 1

Stack, etc. Thread 2

Stack, etc. Thread 3

Dynamic

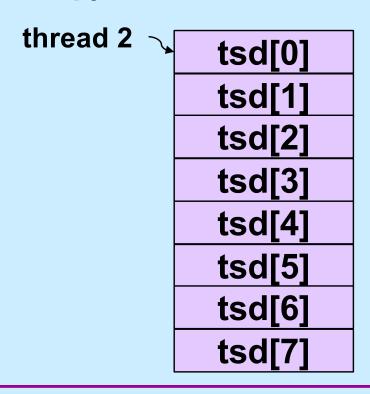
Data

Text

# Generalizing

- Thread-specific data (sometimes called thread-local storage)
  - data that's referred to by global variables, but each thread has its own private copy

41 14 .			
thread 1	tsd[0]		
	tsd[1]		
	tsd[2]		
	tsd[3]		
	tsd[4]		
	tsd[5]		
	tsd[6]		
	4 15-23		



tsd[7

# **Some Machinery**

- pthread\_key\_create(&key, cleanup\_routine)
  - allocates a slot in the TSD arrays
  - provides a function to cleanup when threads terminate
- value = pthread getspecific(key)
  - fetches from the calling thread's array
- pthread setspecific(key, value)
  - stores into the calling thread's array

# errno (Again)

```
// executed before threads are created
pthread_key_t errno_key;
pthread_key_create(&errno_key, NULL);

// redefine errno to use thread-specific value
#define errno pthread_getspecific(errno_key);

// set current thread's errno
pthread_set_specific(errno_key, (void *)ENOMEM);
```

# **Beyond POSIX TLS Extensions for ELF and gcc**

Thread Local Storage (TLS)

```
__thread int x=6;

/*

  * Each thread has its own copy of x,

  * each initialized to 6.

  * Linker and compiler do the setup.

  * May be combined with static or extern.

  * Doesn't make sense for non-static

  * local variables!

  */
```

# **Example: Per-Thread Windows**

```
void *tfunc(void * arg) {
                                          getWindow();
typedef struct {
  wcontext t win context;
                                          threadWrite("started");
  int file descriptor;
} win t;
  thread static win t my win;
                                          func2 (...);
void getWindow() {
  my win.win context = ...;
  my win.file decriptor = ...;
                                        void func2(...) {
int threadWrite(char *buf) {
  int status = write to window(
                                          threadWrite(
      &my win, buf);
                                               "important msg");
  return (status);
```

# Static Local Storage and Threads

```
char *strtok(char *str, const char *delim) {
    static char *saveptr;

    ... // find next token starting at either
    ... // str or saveptr
    ... // update saveptr

    return(&token);
}
```

# Coping

- Use thread local storage
- Allocate storage internally; caller frees it
- Redesign the interface

#### **Thread-Safe Version**

#### **Shared Data**

Thread 1:

```
printf("goto statement reached");
```

Thread 2:

```
printf("Hello World\n");
```

Printed on display:

go to Hell

# Coping

- Wrap library calls with synchronization constructs
- Fix the libraries

# **Efficiency**

- Standard I/O example
  - getc() and putc()
    - » expensive and thread-safe?
    - » cheap and not thread-safe?
  - two versions
    - » getc() and putc()
      - expensive and thread-safe
    - » getc unlocked() and putc unlocked()
      - cheap and not thread-safe
      - made thread-safe with flockfile() and funlockfile()

# **Efficiency**

#### Naive

```
for (i=0; i<lim; i++)</pre>
 putc(out[i]);
```

#### Efficient

```
flockfile(stdout);
for (i=0; i<lim; i++)</pre>
 putc unlocked(out[i]);
funlockfile(stdout);
```

#### What's Thread-Safe?

#### Everything except

asctime()	ecvt()	gethostent()	getutxline()	putc_unlocked()
basename()	encrypt()	getlogin()	gmtime()	putchar_unlocked()
catgets()	endgrent()	getnetbyaddr()	hcreate()	putenv()
crypt()	endpwent()	getnetbyname()	hdestroy()	pututxline()
ctime()	endutxent()	getnetent()	hsearch()	rand()
dbm_clearerr()	fcvt()	getopt()	inet_ntoa()	readdir()
dbm_close()	ftw()	getprotobyname()	I64a()	setenv()
dbm_delete()	gcvt()	getprotobynumber()	lgamma()	setgrent()
dbm_error()	getc_unlocked()	getprotoent()	lgammaf()	setkey()
dbm_fetch()	getchar_unlocked()	getpwent()	lgammal()	setpwent()
dbm_firstkey()	getdate()	getpwnam()	localeconv()	setutxent()
dbm_nextkey()	getenv()	getpwuid()	localtime()	strerror()
dbm_open()	getgrent()	getservbyname()	Irand48()	strtok()
dbm_store()	getgrgid()	getservbyport()	mrand48()	ttyname()
dirname()	getgrnam()	getservent()	nftw()	unsetenv()
dlerror()	gethostbyaddr()	getutxent()	nl_langinfo()	wcstombs()
drand48()	gethostbyname()	getutxid()	ptsname()	wctomb()

# Concurrency

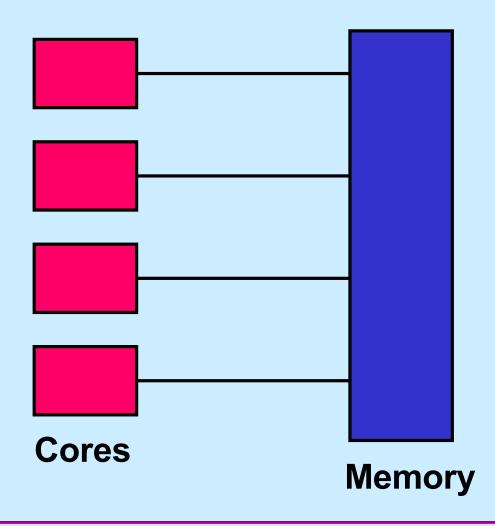
#### Real

- many things happen at once
- multiple threads running on multiple cores

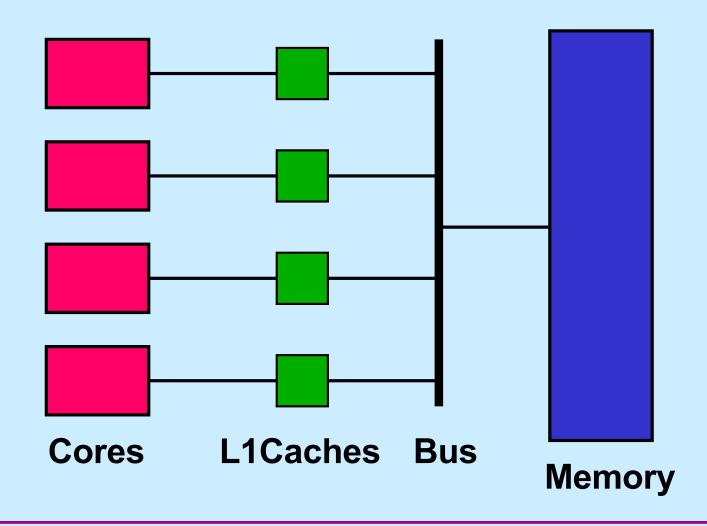
#### Simulated

- things appear to happen at once
- a single core is multiplexed among multiple threads
  - » time slicing

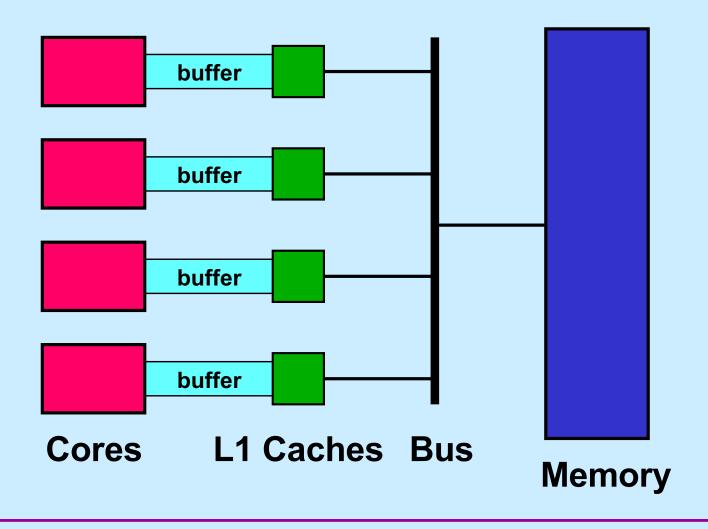
# Multi-Core Processor: Simple View



## Multi-Core Processor: More Realistic View



## Multi-Core Processor: Even More Realistic



# **Concurrent Reading and Writing**

#### Thread 1:

#### Thread 2:

```
i = shared_counter; shared_counter++;
```

#### **Mutual Exclusion w/o Mutexes**

```
void peterson(long me) {
                            // shared
 static long loser;
 static long active[2] = \{0, 0\}; // shared
 long other = 1 - me;
                            // private
 active[me] = 1;
 loser = me;
 while (loser == me && active[other])
 // critical section
 active[me] = 0;
```

#### Quiz 1

```
void peterson(long me) {
                             // shared
 static long loser;
 static long active[2] = \{0, 0\}; // shared
 long other = 1 - me;
                            // private
 active[me] = 1;
 loser = me;
 while (loser == me && active[other])
 // critical section
                       This works on sunlab
 active[me] = 0;
                       machines.
                       a) true
                       b) false
```

# **Busy-Waiting Producer/Consumer**

```
char item;
 while(in - out == BSIZE)
                         while (in - out == 0)
 buf[in%BSIZE] = item;
                         item = buf[out%BSIZE];
 in++;
                         out++;
                         return (item);
```

#### Quiz 2

```
void producer(char item) {
                           char consumer() {
                                  char item;
 while(in - out == BSIZE)
                                  while (in - out == 0)
 buf[in%BSIZE] = item;
                                  item = buf[out%BSIZE];
  in++;
                                  out++;
       This works on sunlab
                                  return(item);
       machines.
       a) true
       b) false
```

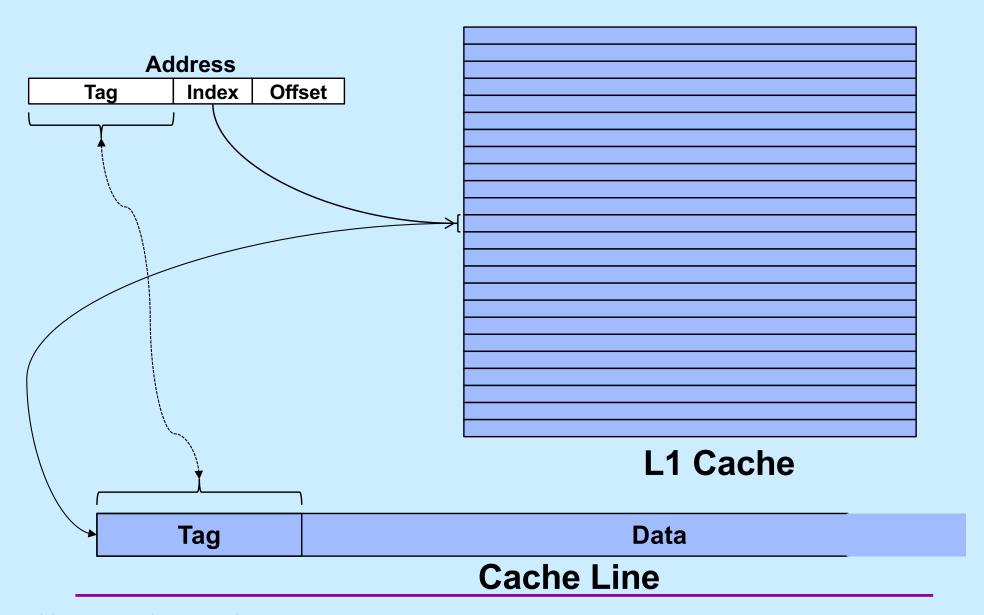
# Coping

- Don't rely on shared memory for synchronization
- Use the synchronization primitives

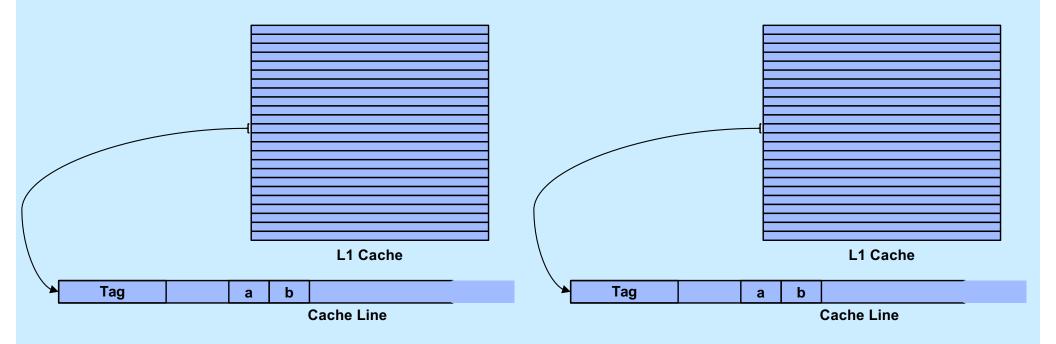
#### Which Runs Faster?

```
volatile int a, b;
                                    volatile int a,
                                      padding[128], b;
void *thread1(void *arg) {
                                    void *thread1(void *arg) {
  int i;
                                      int i;
  for (i=0; i<reps; i++) {
                                      for (i=0; i<reps; i++) {
    a = 1;
                                        a = 1;
void *thread2(void *arg) {
                                    void *thread2(void *arg) {
  int i;
                                      int i;
  for (i=0; i<reps; i++) {</pre>
                                      for (i=0; i<reps; i++) {</pre>
   b = 1;
                                        b = 1;
```

#### **Cache Lines**



# **False Sharing**



# **Implementing Mutexes**

#### Strategy

- make the usual case (no waiting) very fast
- can afford to take more time for the other case (waiting for the mutex)

#### **Futexes**

- Safe, efficient kernel conditional queueing in Linux
- All operations performed atomically

- » otherwise return
- futex wake(futex t \*futex)
  - » wake up one thread from futex's wait queue, if there are any waiting threads

# **Ancillary Functions**

```
• int atomic inc(int *val)

    add 1 to *val, return its original value

• int atomic dec(int *val)

    subtract 1 from *val, return its original value

• int CAS(int *ptr, int old, int new) {
      int tmp = *ptr;
      if (*ptr == old)
          *ptr = new;
      return tmp;
```

# Attempt 1

```
void lock(futex t *futex) {
  int c;
  while ((c = atomic inc(&futex->val)) != 0)
    futex wait(futex, c+1);
void unlock(futex_t *futex) {
  futex->val = 0;
  futex wake(futex);
```

#### Quiz 3

```
void lock(futex t *futex) {
  int c;
  while ((c = atomic inc(&futex->val)) != 0)
    futex wait(futex, c+1);
void unlock(futex t *futex)
  futex->val = 0;
  futex wake (futex);
```

Which of the following won't happen if the futex's value is zero and three threads call lock at the same time?

- a) one will return immediately, two will call futex\_wait.
- b) even though unlock is called appropriately, one thread will never return from futex\_wait.
- c) threads might return from futex\_wait immediately, because the futex's value is not equal to c+1.

## Attempt 2

```
void lock(futex_t *futex) {
  int c;
  if ((c = CAS(&futex->val, 0, 1) != 0)
    do {
      if (c == 2 || (CAS(&futex->val, 1, 2) != 0))
        futex wait(futex, 2);
    while ((c = CAS(\&futex->val, 0, 2)) != 0))
void unlock(futex t *futex) {
  if (atomic dec(&futex->val) != 1) {
    futex->val = 0;
    futex wake(futex);
```

# **Memory Allocation**

Multiple threads

One heap

**Bottleneck?** 

#### **Solution 1**

- Divvy up the heap among the threads
  - each thread has its own heap
  - no mutexes required
  - no bottleneck
- How much heap does each thread get?

#### Solution 2

- Multiple "arenas"
  - each with its own mutex
  - thread allocates from the first one it can find whose mutex was unlocked
    - » if none, then creates new one
  - deallocations go back to original arena

#### **Solution 3**

- Global heap plus per-thread heaps
  - threads pull storage from global heap
  - freed storage goes to per-thread heap
    - » unless things are imbalanced
      - then thread moves storage back to global heap
  - mutex on only the global heap
- What if one thread allocates and another frees storage?

### Malloc/Free Implementations

- ptmalloc
  - based on solution 2
  - in glibc (i.e., used by default)
- tcmalloc
  - based on solution 3
  - from Google
- Which is best?

### **Test Program**

```
const unsigned int N=64, nthreads=32, iters=10000000;
int main() {
  void *tfunc(void *);
 pthread t thread[nthreads];
  for (int i=0; i<nthreads; i++) {
    pthread create(&thread[i], 0, tfunc, (void *)i);
    pthread detach(thread[i]);
 pthread exit(0);
void *tfunc(void *arg) {
  long i;
  for (i=0; i<iters; i++) {
    long *p = (long *) malloc(sizeof(long) * ((i%N) +1));
    free(p);
  return 0;
```

#### Compiling It ...

% gcc -o ptalloc alloc.cc -lpthread
% gcc -o tcalloc alloc.cc -lpthread -ltcmalloc

### Running It (2014) ...

```
$ time ./ptalloc
real 0m5.142s
user 0m20.501s
sys 0m0.024s
$ time ./tcalloc
real 0m1.889s
user 0m7.492s
sys 0m0.008s
```

# Running It (2022) ...

```
$ time ./ptalloc real 0m1.156s user 0m3.456s sys 0m0.004s $ time ./tcalloc real 0m0.876s user 0m3.460s sys 0m0.004s
```

### What's Going On (2014)?

```
$ strace -c -f ./ptalloc
% time seconds usecs/call calls errors syscall
100.00 0.040002 13 3007 520 futex
$ strace -c -f ./tcalloc
% time seconds usecs/call calls errors syscall
 0.00 0.000000
                0 59 13 futex
```

### What's Going On (2022)?

```
$ strace -c -f ./ptalloc
% time seconds usecs/call calls errors syscall
31.23 0.019968 416 48 6 futex
$ strace -c -f ./tcalloc
% time seconds usecs/call calls errors syscall
0.00 0.000000 0 42 3 futex
```

```
#define N 64
#define npairs 16
#define allocsPerIter 1024
const long iters = 8*1024*1024/allocsPerIter;
#define BufSize 10240
typedef struct buffer {
  int *buf[BufSize];
  unsigned int nextin;
  unsigned int nextout;
  sem t empty;
  sem t occupied;
 pthread t pthread;
 pthread t cthread;
} buffer t;
```

```
int main() {
  long i;
 buffer t b[npairs];
  for (i=0; i<npairs; i++) {
   b[i].nextin = 0;
   b[i].nextout = 0;
    sem init(&b[i].empty, 0, BufSize/allocsPerIter);
    sem init(&b[i].occupied, 0, 0);
   pthread create(&b[i].pthread, 0, prod, &b[i]);
   pthread create(&b[i].cthread, 0, cons, &b[i]);
  for (i=0; i<npairs; i++) {
   pthread join(b[i].pthread, 0);
   pthread join(b[i].cthread, 0);
  return 0;
```

```
void *prod(void *arg) {
  long i, j;
  buffer t *b = (buffer t *) arg;
  for (i = 0; i<iters; i++) {
    sem wait(&b->empty);
    for (j = 0; j<allocsPerIter; j++) {</pre>
      b->buf[b->nextin] = malloc(sizeof(int)*((j%N)+1));
      if (++b->nextin >= BufSize)
       b->nextin = 0:
    sem post(&b->occupied);
  return 0;
```

```
void *cons(void *arg) {
  long i, j;
  buffer t *b = (buffer t *) arg;
  for (i = 0; i<iters; i++) {
    sem wait(&b->occupied);
    for (j = 0; j<allocsPerIter; j++) {</pre>
      free(b->buf[b->nextout]);
      if (++b->nextout >= BufSize)
       b->nextout = 0;
    sem post(&b->empty);
  return 0;
```

### Running It (2014) ...

```
$ time ./ptalloc2
real 0m1.087s
user 0m3.744s
sys 0m0.204s
$ time ./tcalloc2
real 0m3.535s
user 0m11.361s
sys 0m2.112s
```

### Running It (2022) ...

```
$ time ./ptalloc2
real 0m0.367s
user 0m1.187s
sys 0m0.179s
$ time ./tcalloc2
real 0m0.426s
user 0m1.211s
sys 0m0.290s
```

# What's Going On (2014)?

# What's Going On (2022)?