**CS 33** 

Files Part 3

## **Setting File Permissions**

```
#include <sys/types.h>
#include <sys/stat.h>
int chmod(const char *path, mode_t mode)
```

- sets the file permissions of the given file to those specified in *mode*
- only the owner of a file and the superuser may change its permissions
- nine combinable possibilities for mode
   (read/write/execute for user, group, and others)

```
» S_IRUSR (0400), S_IWUSR (0200), S_IXUSR (0100)
» S_IRGRP (040), S_IWGRP (020), S_IXGRP (010)
» S IROTH (04), S IWOTH (02), S IXOTH (01)
```

#### **Umask**

 Standard programs create files with "maximum needed permissions" as mode

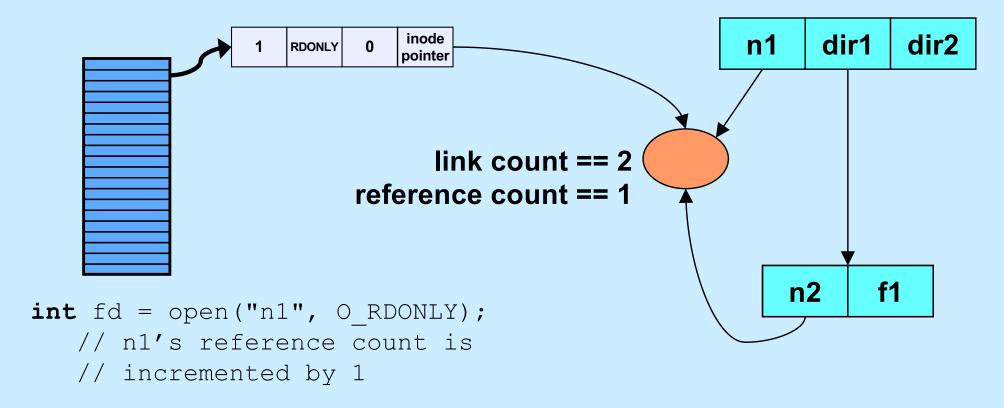
– compilers: 0777

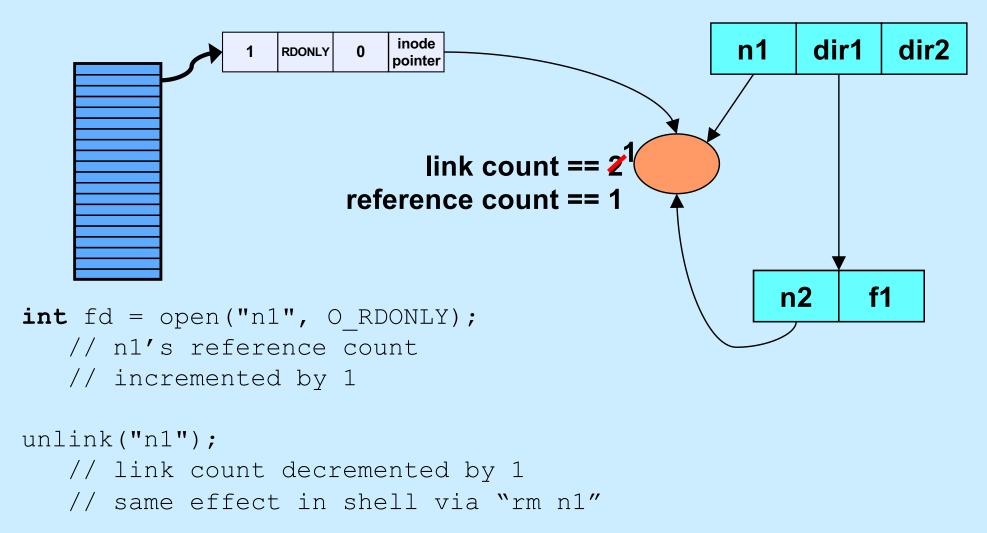
– editors: 0666

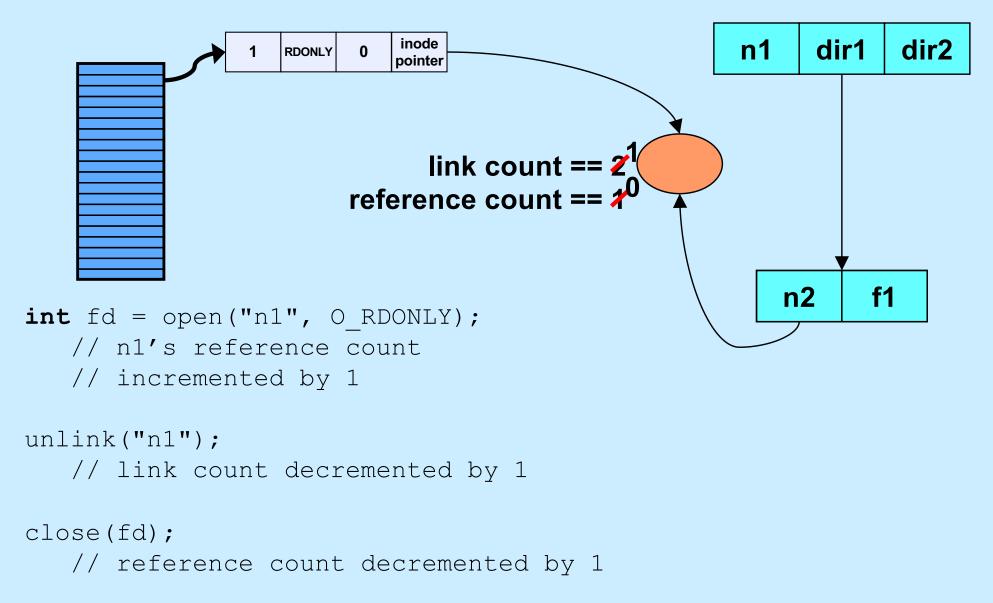
- Per-process parameter, umask, used to turn off undesired permission bits
  - e.g., turn off all permissions for others, write permission for group: set umask to 027
    - **»** compilers: permissions =  $0777 \& \sim (027) = 0750$
    - > editors: permissions = 0666 &  $\sim$ (027) = 0640
  - set with umask system call or (usually) shell command

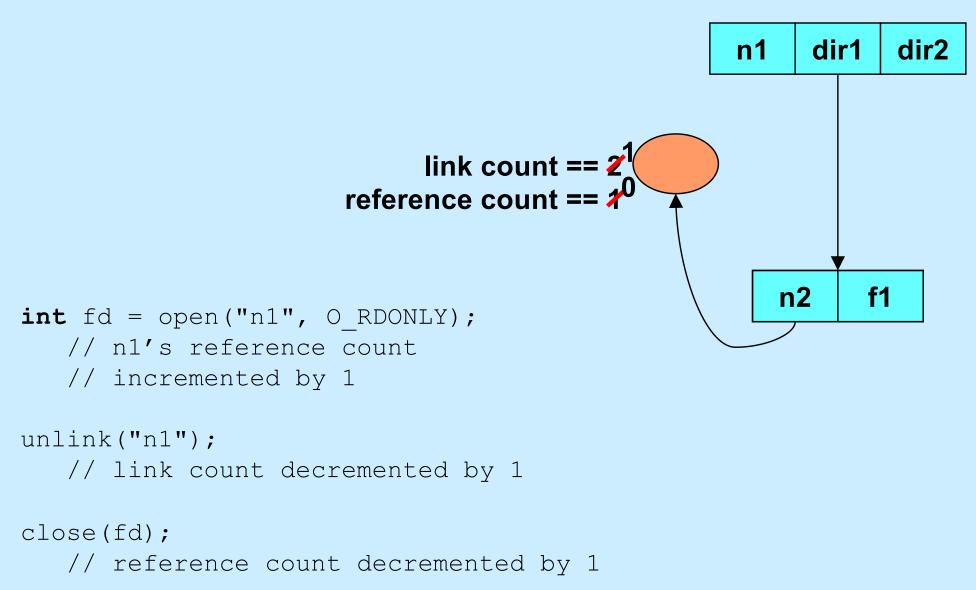
## **Creating a File**

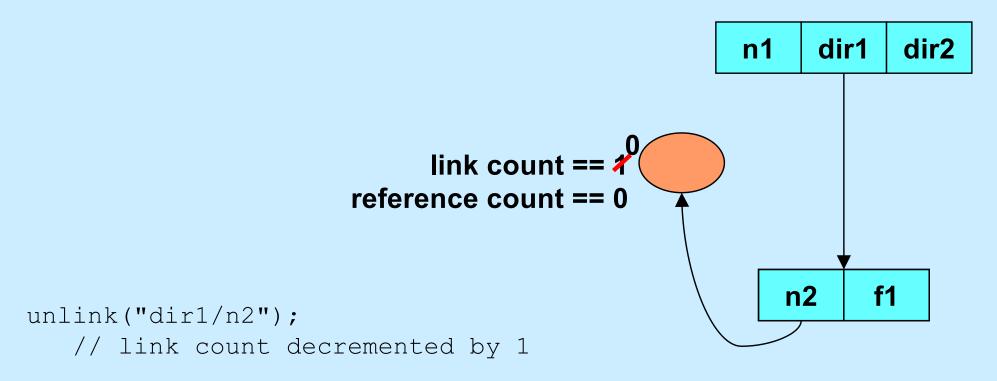
- Use either open or creat
  - open(const char \*pathname, int flags, mode\_t mode)
    - » flags must include O\_CREAT
  - creat(const char \*pathname, mode\_t mode)
    - » open is preferred
- The mode parameter helps specify the permissions of the newly created file
  - permissions = mode & ~umask











#### Quiz 1

```
int main() {
  int fd = open("file", O_RDWR|O_CREAT, 0666);
  unlink("file");
  PutStuffInFile(fd);
  GetStuffFromFile(fd);
  return 0;
}
```

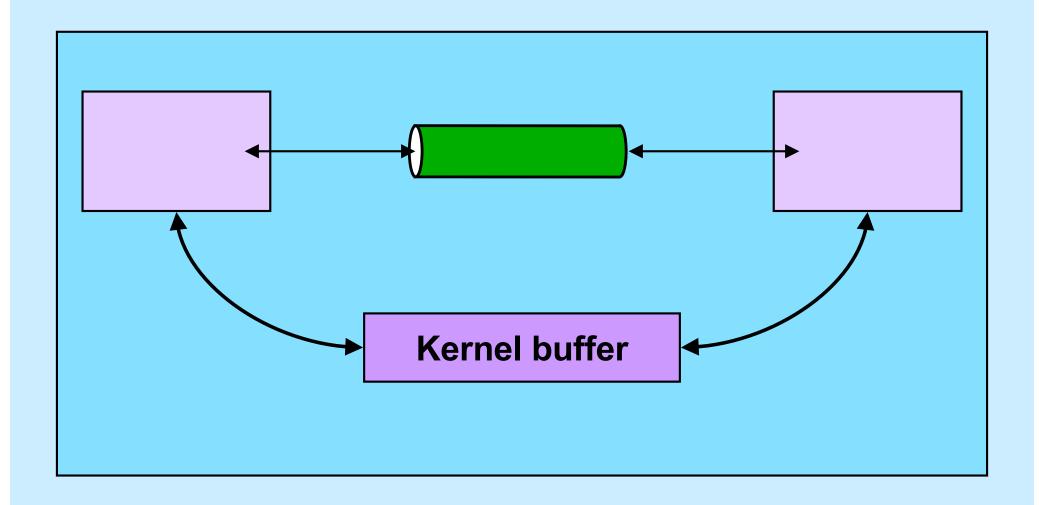
Assume that *PutStuffInFile* writes to the given file, and *GetStuffFromFile* reads from the file.

- a) This program is doomed to failure, since the file is deleted before it's used
- b) Because the file is used after the unlink call, it won't be deleted
- c) The file will be deleted when the program terminates

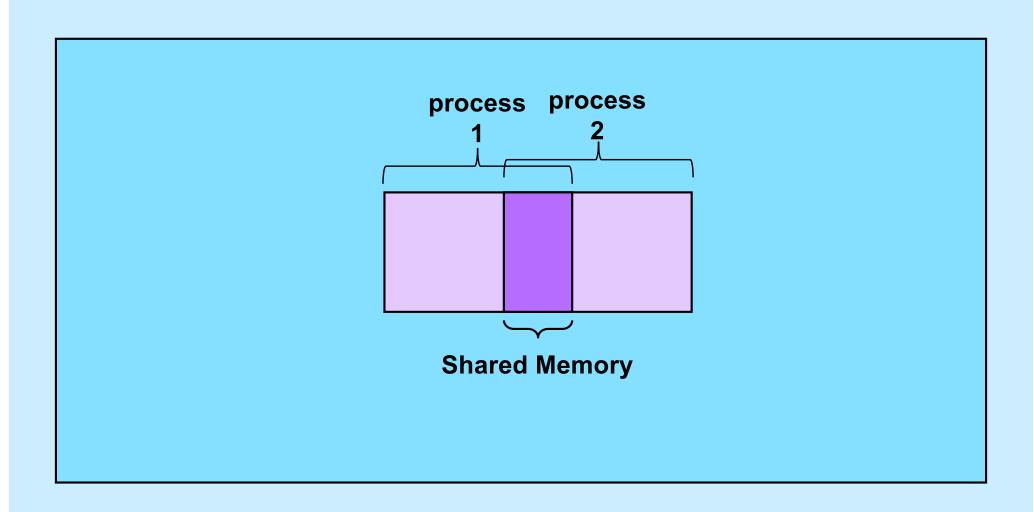
# Interprocess Communication (IPC): Pipes



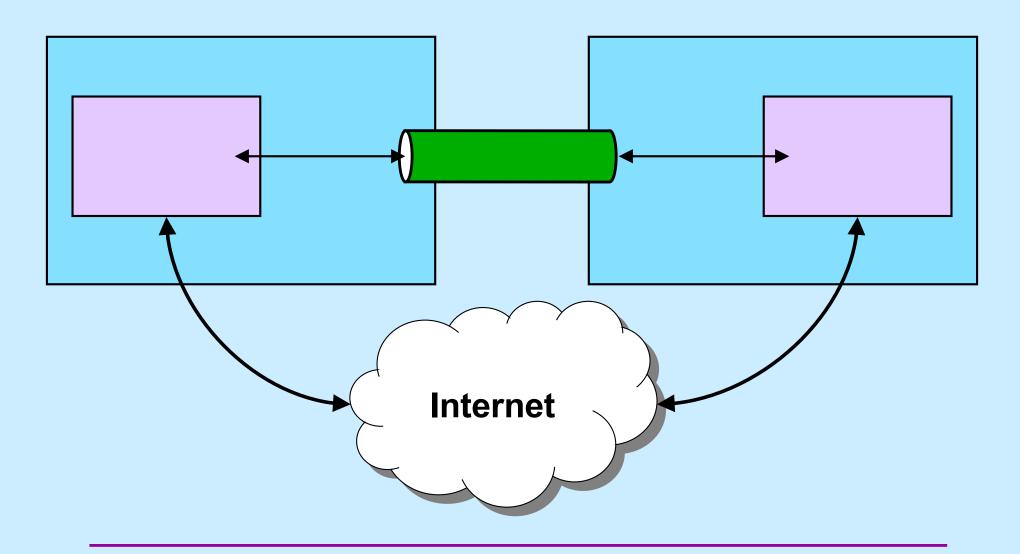
# Interprocess Communication: Same Machine I



# **Interprocess Communication: Same Machine II**

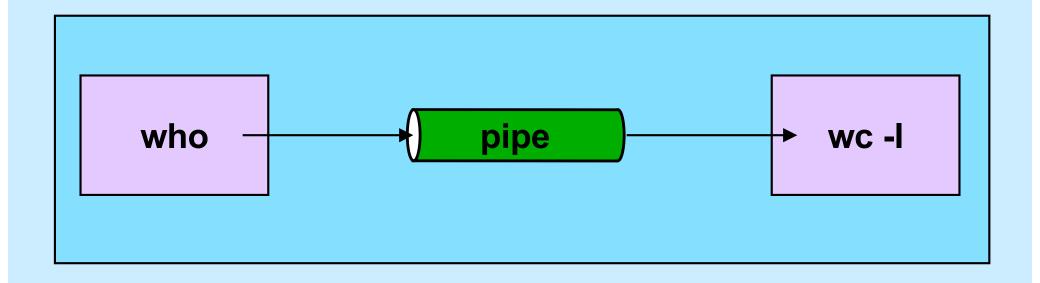


# **Interprocess Communication: Different Machines**



# **Pipes**

\$cslab2e who | wc -1



## **Using Pipes in C**

\$cslab2e who | wc -1 **int** fd[2]; fd[1] pipe **→** fd[0] pipe(fd); **if** (fork() == 0) { close(fd[0]); close(1);dup(fd[1]); close(fd[1]); execl("/usr/bin/who", "who", 0); // who sends output to pipe **if** (fork() == 0) { close(fd[1]); close(0);dup(fd[0]); close(fd[0]); execl("/usr/bin/wc", "wc", "-1", 0); // wc's input is from pipe close(fd[1]); close(fd[0]); // ...

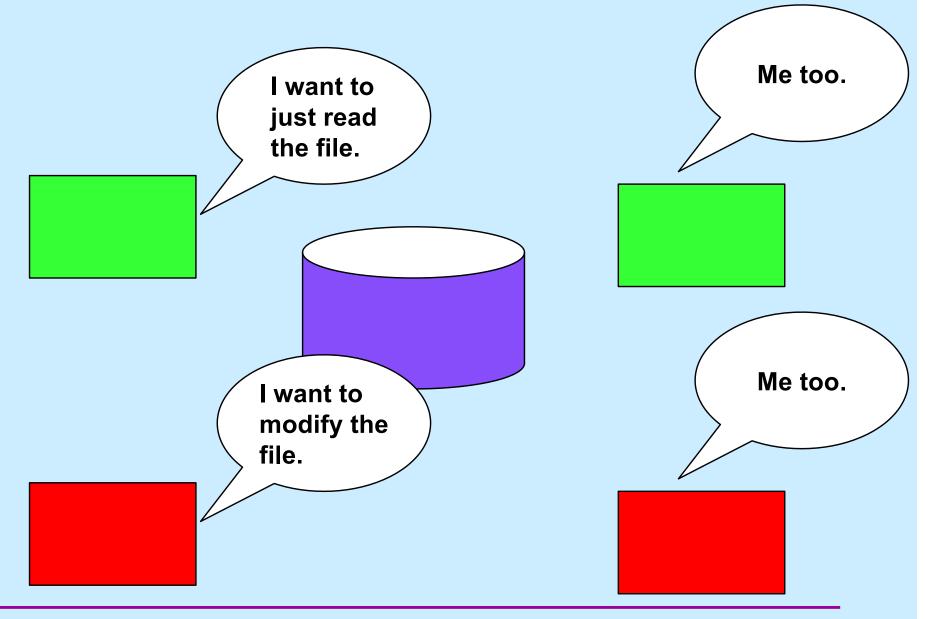
## **Sharing Files**

- You're doing a project with a partner
- You code it as one 15,000-line file
  - the first 7,500 lines are yours
  - the second 7,500 lines are your partner's
- You edit the file, changing 6,000 lines
  - it's now 5am
- Your partner completes her changes at 5:01am
- At 5:02am you look at the file
  - your partner's changes are there
  - yours are not

#### Lessons

- Never work with a partner
- Use more than one file
- Read up on git
- Use an editor and file system that support file locking

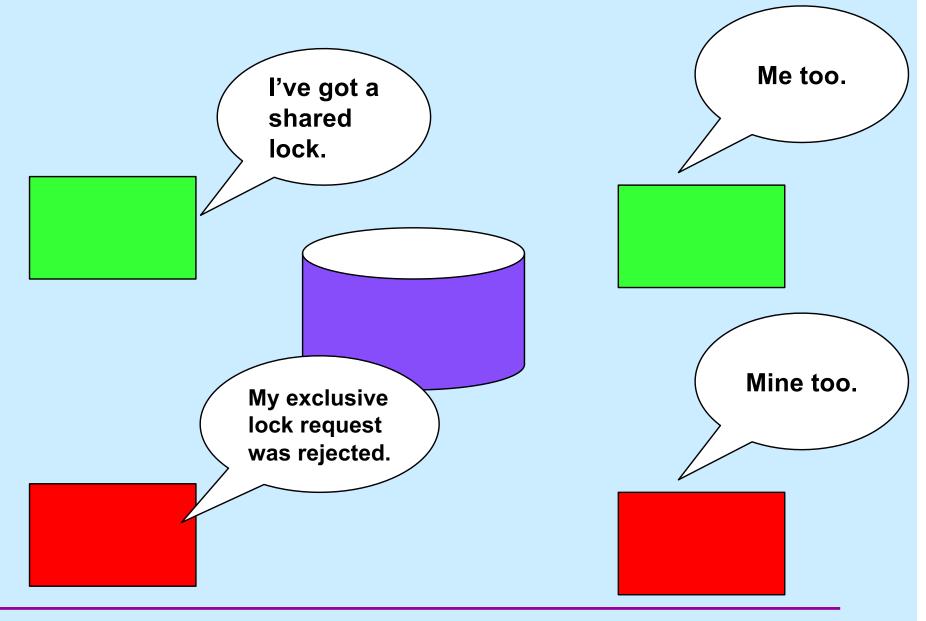
### What We Want ...



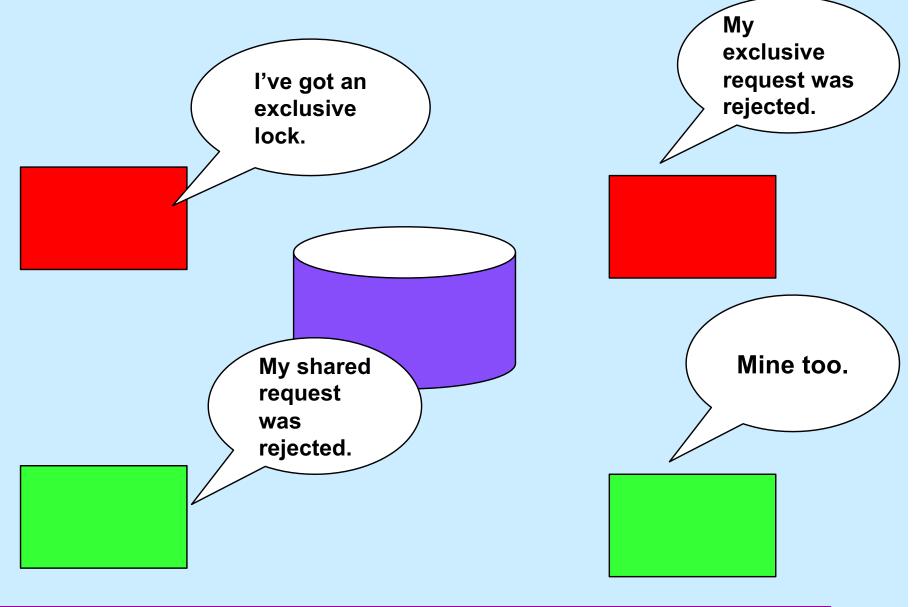
## **Types of Locks**

- Shared (readers) locks
  - any number may have them at same time
  - may not be held when an exclusive lock is held by others
- Exclusive (writers) locks
  - only one at a time
  - may not be held when any other lock is held by others

#### What We Want ...



### What We Want ...



## **Locking Files**

- Early Unix didn't support file locking
- How did people survive?

```
- open ("file.lck", O RDWR | O CREAT | O EXCL, 0666);
```

- » operation fails if *file.lck* exists, succeeds (and creates file.lck) otherwise
- » requires cooperative programs

## Locking Files (continued)

- How it's done in "modern" Unix
  - "advisory locks" may be placed on files
    - » may request shared (readers) or exclusive (writers) lock
      - fcntl system call
    - » either succeeds or fails
    - » open, read, write always work, regardless of locks
    - » a lock applies to a specified range of bytes, not necessarily to the whole file
    - » requires cooperative programs
  - "mandatory locks" supported as a per-file option
    - » set along with permission bits
    - » if set, file can't be used unless process possesses appropriate locks

### Locking Files (still continued)

#### How to:

```
struct flock fl;
fl.l type = F RDLCK; // read lock
// fl.l type = F WRLCK; // write lock
// fl.l type = F UNLCK; // unlock
fl.1 whence = SEEK SET; // starting where
              // offset
fl.1 start = 0;
            // how much? (0 = whole file)
fl.1 len = 0;
fd = open("file", O RDWR);
if (fcntl(fd, F SETLK, &fl) == -1)
 if ((errno == EACCES) || (errno == EAGAIN))
   // didn't get lock
 else
   // something else is wrong
else
 // got the lock!
```

#### Quiz 2

- Your program currently has a shared lock on a portion of a file. It would like to "upgrade" the lock to be an exclusive lock. Would there be any problems with adding an option to fcntl that would allow the holder of a shared lock to wait until it's possible to upgrade to an exclusive lock, then do the upgrade?
  - a) at least one major problem
  - b) either no problems whatsoever or some easy-to-deal-with problems

## **Shell 1: Artisanal Coding**

```
while ((line = get a line()) != 0) {
    tokens = parse line(line);
    for (int i=0; i < ntokens; i++) {
        if (strcmp(tokens[i], ">") == 0) {
            // handle output redirection
        } else if (strcmp(tokens[i], "<") == 0) {</pre>
            // handle input redirection
        } else if (strcmp(tokens[i], "&") == 0) {
            // handle "no wait"
        } ... else {
           // handle other cases
    if (fork() == 0) {
        // . . .
       execv(...);
    // ...
```

## **Shell 1: Non-Artisanal Coding (1)**

```
while ((line = get_a_line()) != 0) {
    tokens = parse_line(line);
    for (int i=0; i < ntokens; i++) {
        // handle "normal" case
    }
    if (fork() == 0) {
        // ...
        execv(...);
    }
    // ...
}</pre>
```

## **Shell 1: Non-Artisanal Coding (2)**

```
next line: while ((line = get a line()) != 0) {
   tokens = parse line(line);
   for (int i=0; i < ntokens; i++) {
       if (redirection symbol(token[i])) {
           // ...
           if (fork() == 0) {
              // . . .
             execv(...); whoops!
           // ...
           goto next line;
       // handle "normal" case
   if (fork() == 0) {
       // ... (whoops!)
       execv(...);
   // ...
```

## **Shell 1: Non-Artisanal Coding (3)**

```
next line: while ((line = get a line()) != 0) {
    tokens = parse line(line);
    for (int i=0; i < ntokens; i++) {
        if (redirection symbol(token[i])) {
            // ...
            if (fork() == 0) {
               // . . .
              execv(...);
            // ... deal with &
            goto next line;
        // handle "normal" case
    if (fork() == 0) {
       // ...
        execv(...);
    // ... also deal with & here!
```

## **Shell 1: Non-Artisanal Coding (Worse)**

```
next line: while ((line = get a line()) != 0) {
tokens = parse line(line);
for (int i=0; i < ntokens; i++) {
if (redirection symbol(token[i])) {
// ...
if (fork() == 0) {
// . . .
execv(...);
// ... deal with &
goto next line;
// handle "normal" case
if (fork() == 0) {
// ...
execv(...);
// ... also deal with & here!
```

## **Artisanal Programming**

- Factor your code!
  - A; FE | B; FE | C; FE = (A | B | C); FE
- Format as you write!
  - don't run the formatter only just before handing it in
  - your code should always be well formatted
- If you have a tough time understanding your code, you'll have a tougher time debugging it and TAs will have an even tougher time helping you

#### It's Your Code

- Be proud of it!
  - it not only works; it shows skillful artisanship
- It's not enough to merely work
  - others have to understand it
    - » (not to mention you ...)
  - you (and others) have to maintain it
    - » shell 2 is coming soon!