

Introduction



ΜΥΥ 801
ΔΙΚΤΥΑ ΥΠΟΛΟΓΙΣΤΩΝ II

Liaskos Christos
UoI CSE

Course Details

Instructor: Christos Liaskos
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Class Hours:

Thu: 16:00-19:00 (Theory - Αμφ. 1)
Fri: 14:00-16:00 (Labs)

Contact:

via email (office hours TBA)

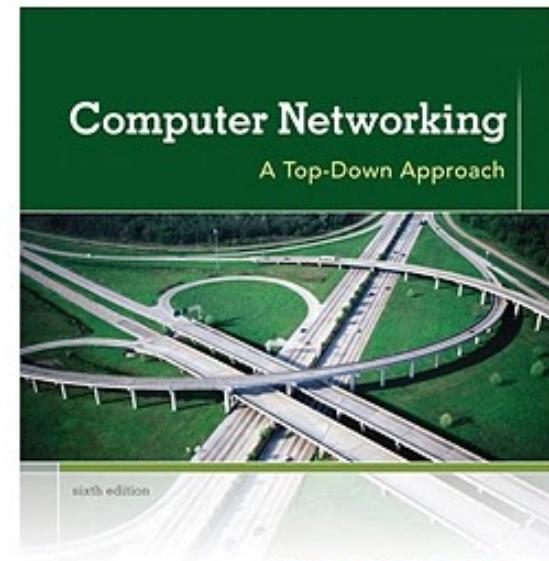
Outline

- Course Info
- Introduction
- Some basic concepts
- History of popular wireless systems

Textbook

Slides used/adapted from the original material by the book authors.

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*Computer
Networking: A Top
Down Approach*
7th edition
Jim Kurose, Keith Ross
Addison-Wesley
March 2012

Textbook outline

-  [Chapter 1 Computer Networks and the Internet](#)

-  Chapter 2 Application Layer
How do applications send data to each other?
-  Chapter 3 Transport Layer
How does the network multiplexes the data?
-  Chapter 4 The Network Layer
How does the network perform routing?
-  Chapter 5 The Link Layer:
Links, Access Networks,
and LANs
How are frames delivered in a LAN?
-  Chapter 6 Wireless and Mobile Networks
-  Chapter 7 Multimedia Networking
-  Chapter 8 Security in Computer Networks

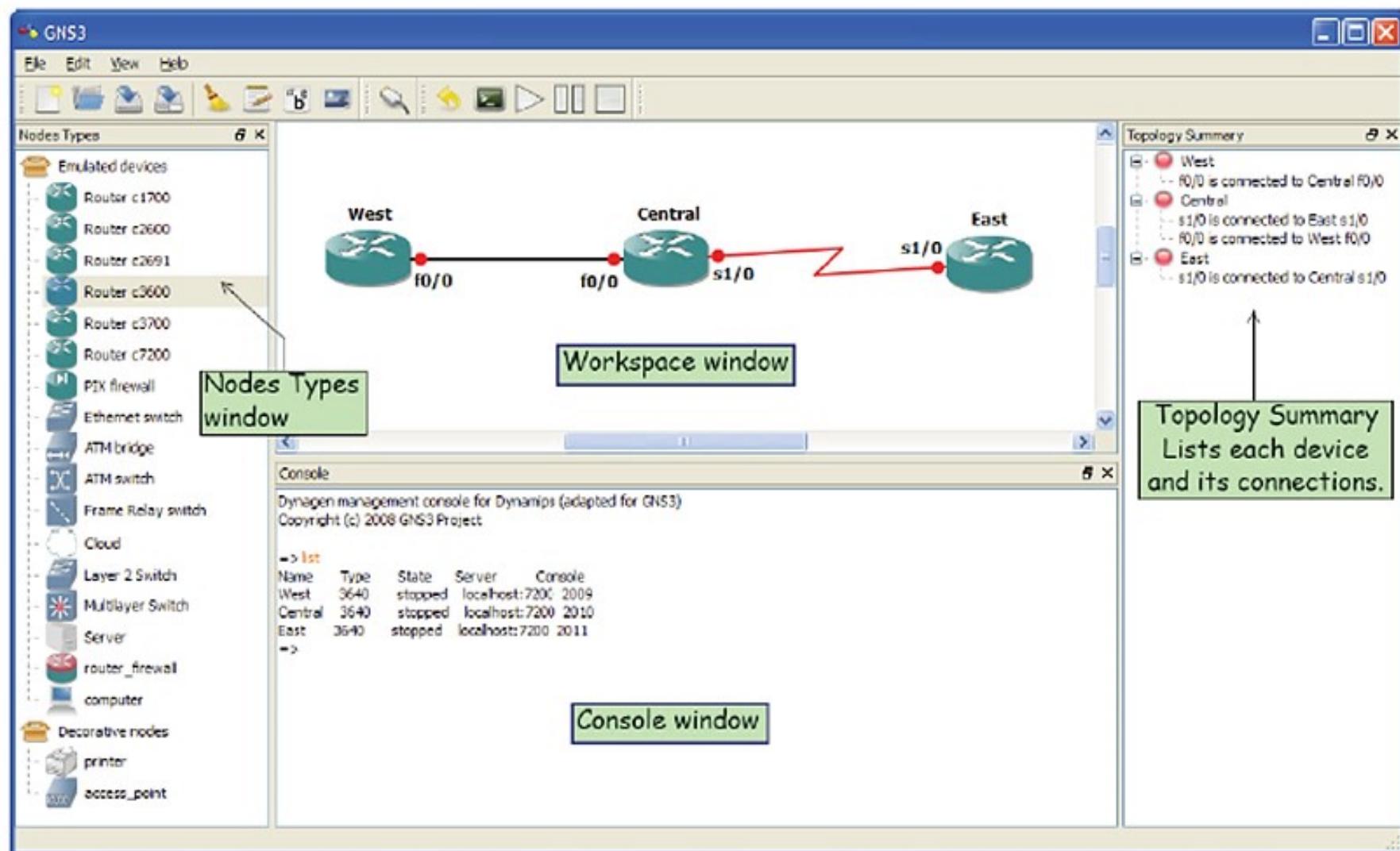
-  Chapter 9 Network Management

We will see some recent DDoS attack types
We will survey a few things about SDN

Grading

- There will be one final course exam
 - ❖ The usual stuff..
- There will be 1 lab exam.
 - ❖ It will count as fixed part of the overall grade.
 - E.g.,
 - If you pass with honors, you get 2 pts, the rest 8 pts via the final exam.
 - If you fail, you go to the exams with a -2pt handicap (max 8).
 - ❖ Lab exercises: based on GNS3
 - Graphical Network Simulator-3.
 - Really cool, realistic network simulator.
 - We will see several networking aspects.
 - You will be examined in much simpler things.

GNS3



GNS3

- Everything we need will be available at the CSE Labs.
- You are encouraged to use this tool on your own as well!
 - ❖ E.g. at your home PC
 - ❖ You will need a decent 8-core PC / 8 GB ram.
- Lots and lots of educational material, tutorials, YouTube channels on how to install, use, etc.

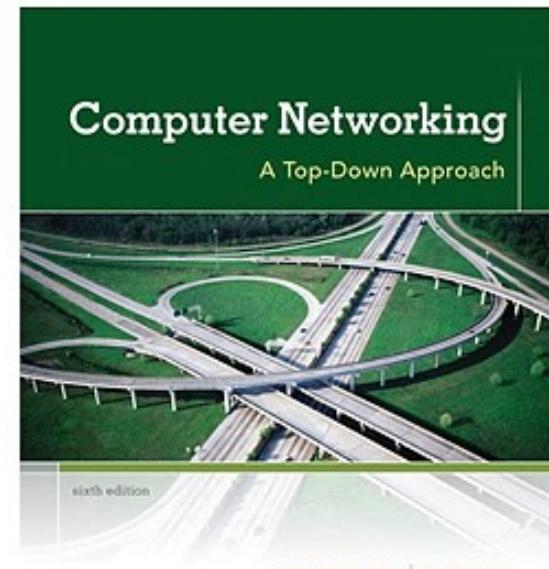
Course/Lab Announcements

- Via my web-site
 - ❖ <https://www.cse.uoi.gr/~cliaskos/>

Chapter I Introduction

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Chapter 1: Introduction

Our goal:

- get "feel" and terminology
- more depth, detail later in course
- approach:
 - ❖ use Internet as example

Overview:

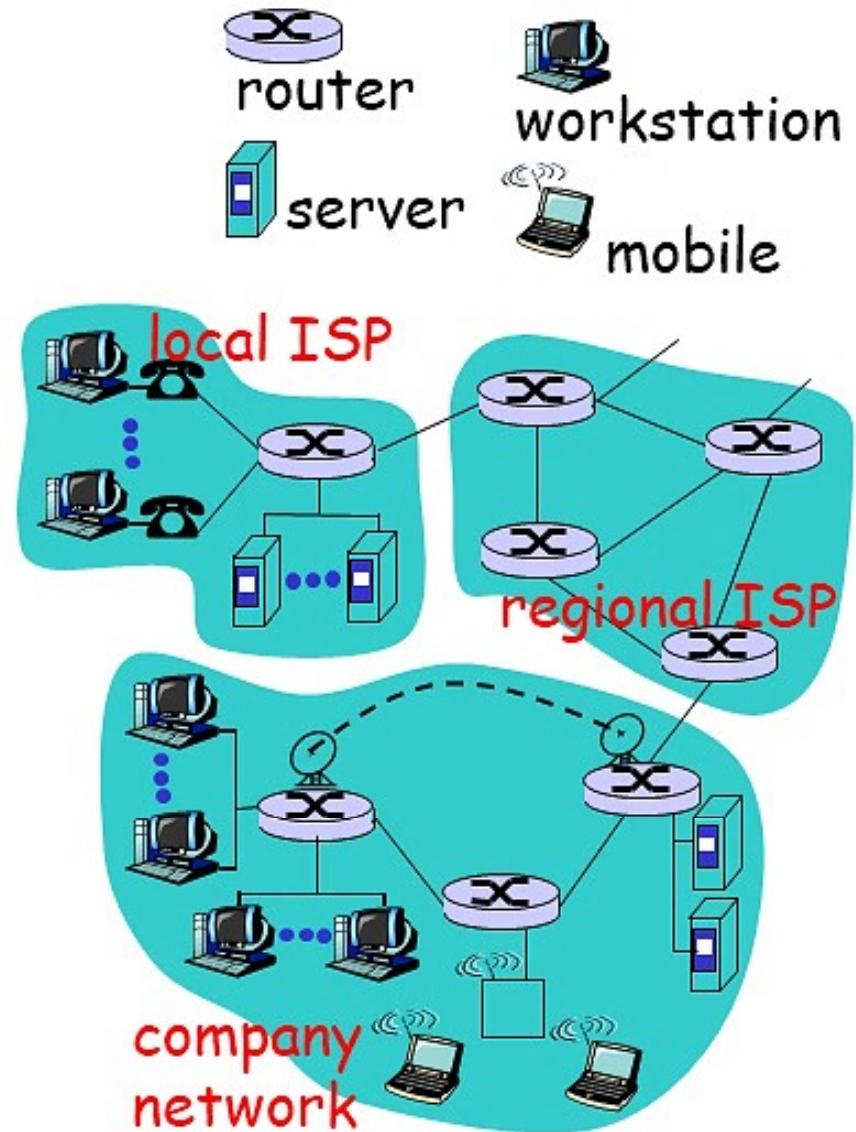
- what's the Internet
- what's a protocol?
- network edge
- network core
- access net, physical media
- Internet/ISP structure
- performance: loss, delay
- protocol layers, service models
- network modeling

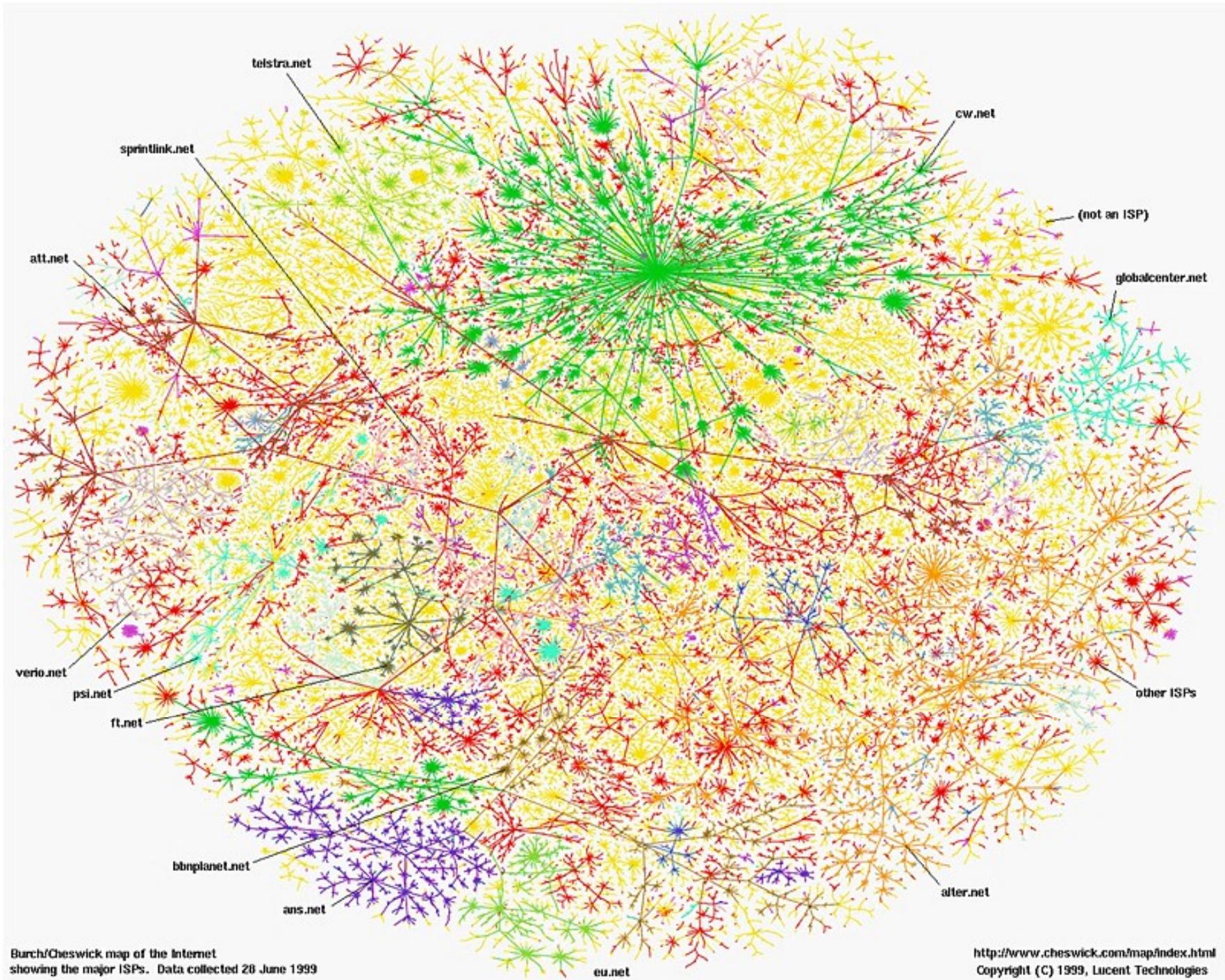
Chapter 1: roadmap

- 1.1 What *is* the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History

What's the Internet: "nuts and bolts" view

- ❑ millions of connected computing devices: *hosts* = *end systems*
- ❑ running *network apps*
- ❑ *communication links*
 - ❖ fiber, copper, radio, satellite
 - ❖ transmission rate = *bandwidth*
- ❑ *routers*: forward packets (chunks of data)





Burch/Cheswick map of the Internet
showing the major ISPs. Data collected 28 June 1999

<http://www.cheswick.com/map/index.html>
Copyright (C) 1999, Lucent Technologies

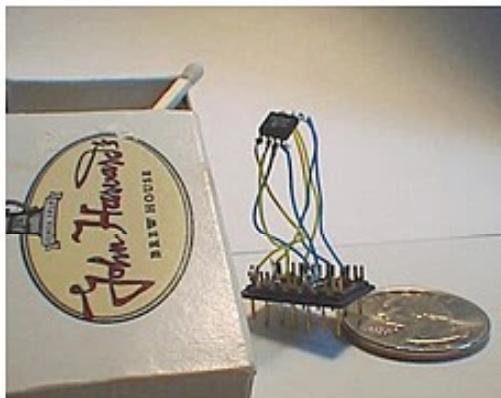
"Cool" old internet appliances



IP picture frame
<http://www.ceiva.com/>



Web-enabled toaster +
weather forecaster

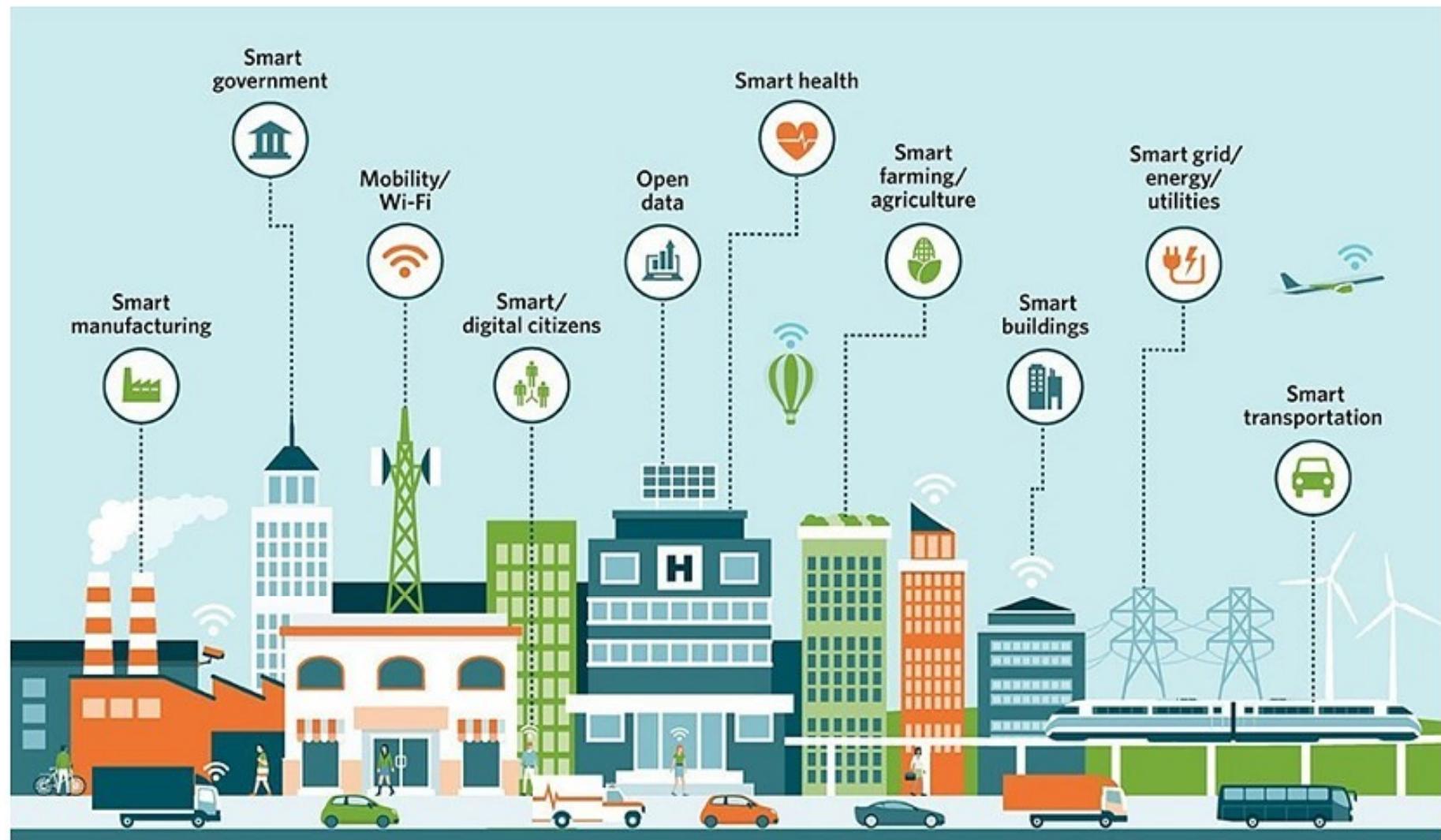


World's smallest web server
<http://www-ccs.cs.umass.edu/~shri/iPic.html>



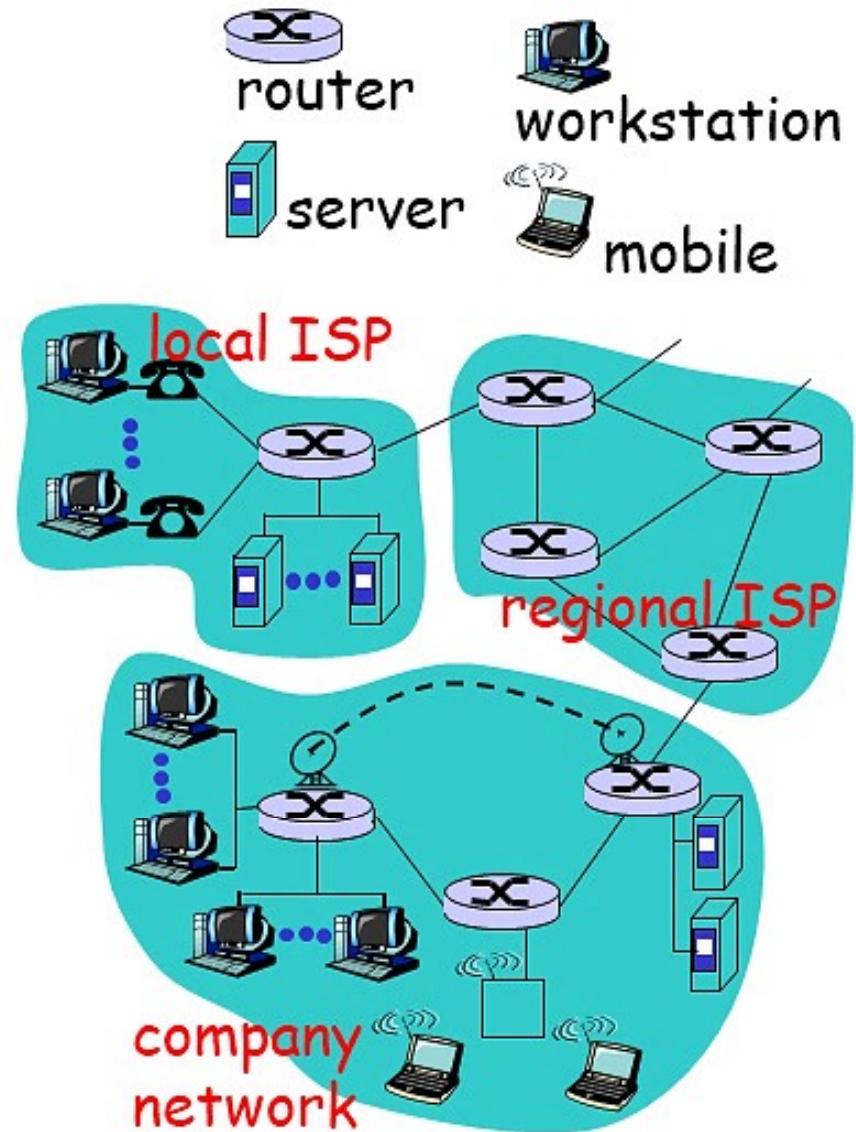
Internet phones

Towards IoE



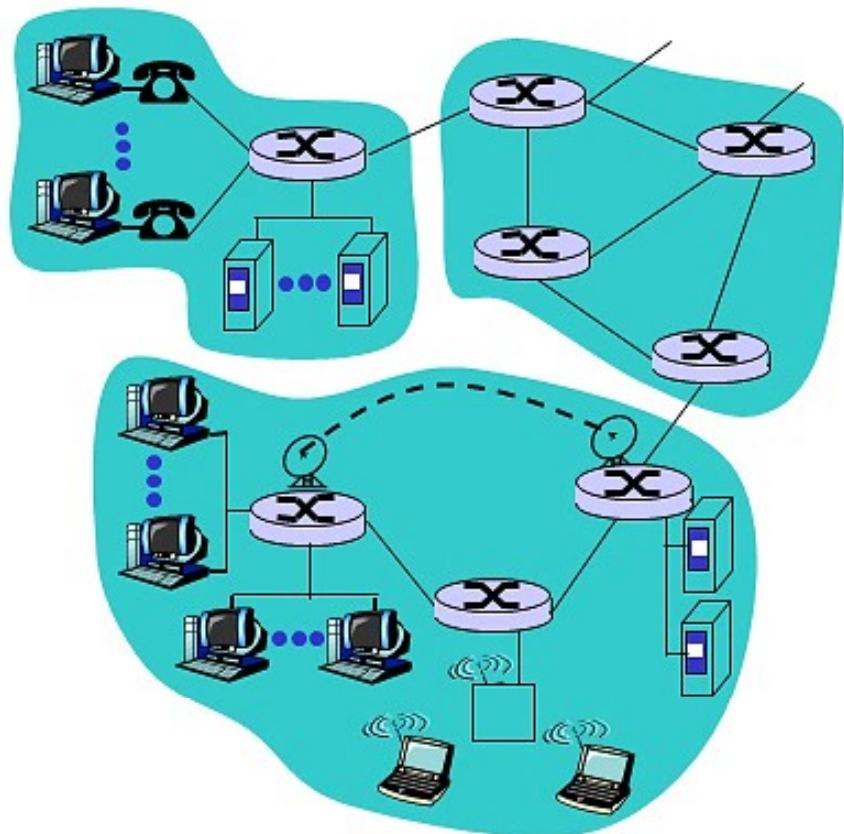
What's the Internet: "nuts and bolts" view

- **protocols** control sending, receiving of msgs
 - ❖ e.g., TCP, IP, HTTP, FTP, PPP
- **Internet: "network of networks"**
 - ❖ loosely hierarchical
 - ❖ public Internet versus private intranet
- **Internet standards**
 - ❖ RFC: Request for comments
 - ❖ IETF: Internet Engineering Task Force



What's the Internet: a service view

- communication *infrastructure* enables distributed applications:
 - ❖ Web, email, games, e-commerce, file sharing
- communication services provided to apps:
 - ❖ Connectionless unreliable
 - ❖ connection-oriented reliable



What's a protocol?

human protocols:

- "what's the time?"
- "I have a question"
- introductions

... specific msgs sent

... specific actions taken
when msgs received,
or other events

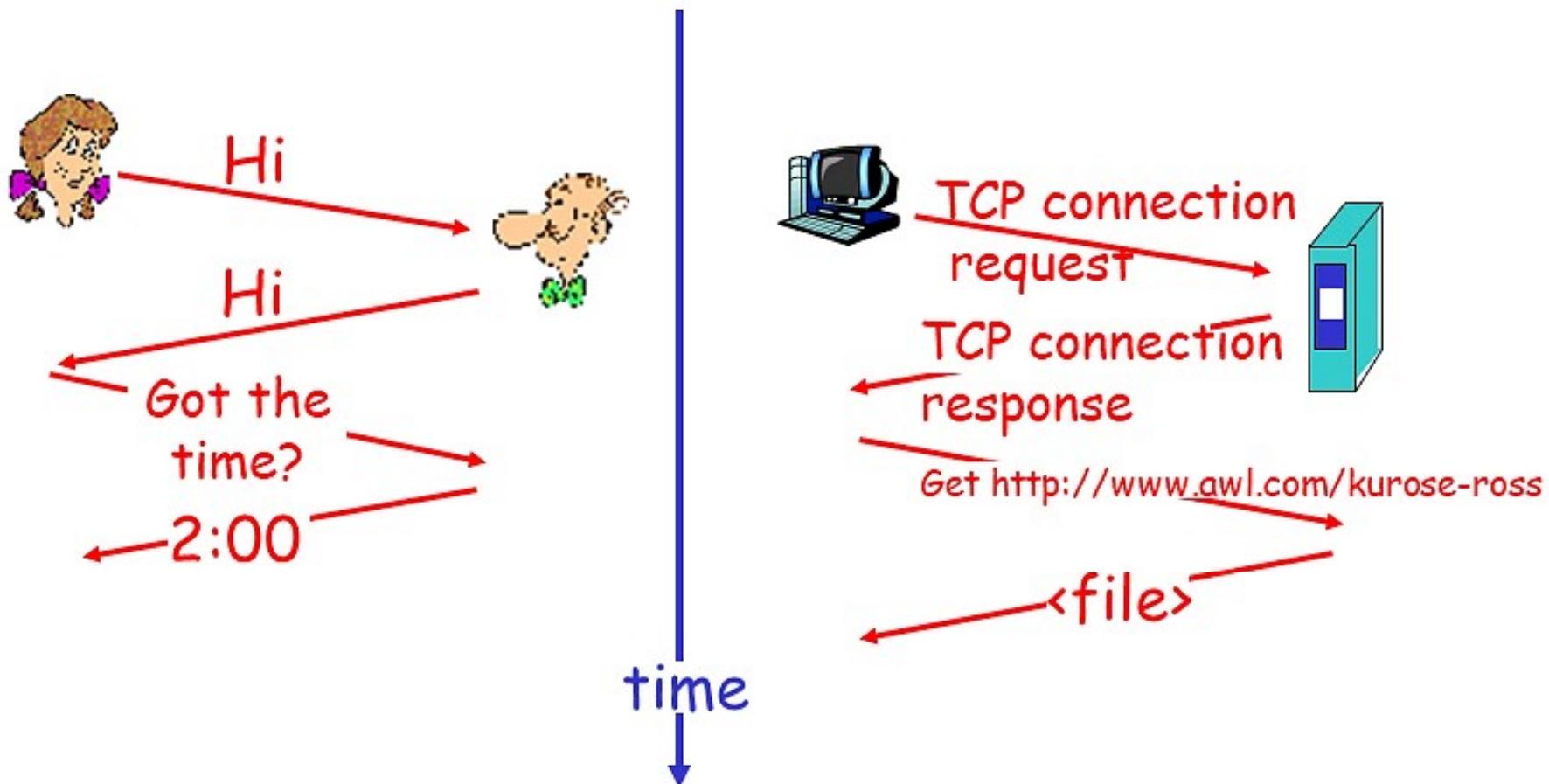
network protocols:

- machines rather than humans
- all communication activity in Internet governed by protocols

*protocols define format,
order of msgs sent and
received among network
entities, and actions
taken on msg
transmission, receipt*

What's a protocol?

a human protocol and a computer network protocol:



Q: Other human protocols?

Chapter 1: roadmap

1.1 What is the Internet?

1.2 Network edge

1.3 Network core

1.4 Network access and physical media

1.5 Internet structure and ISPs

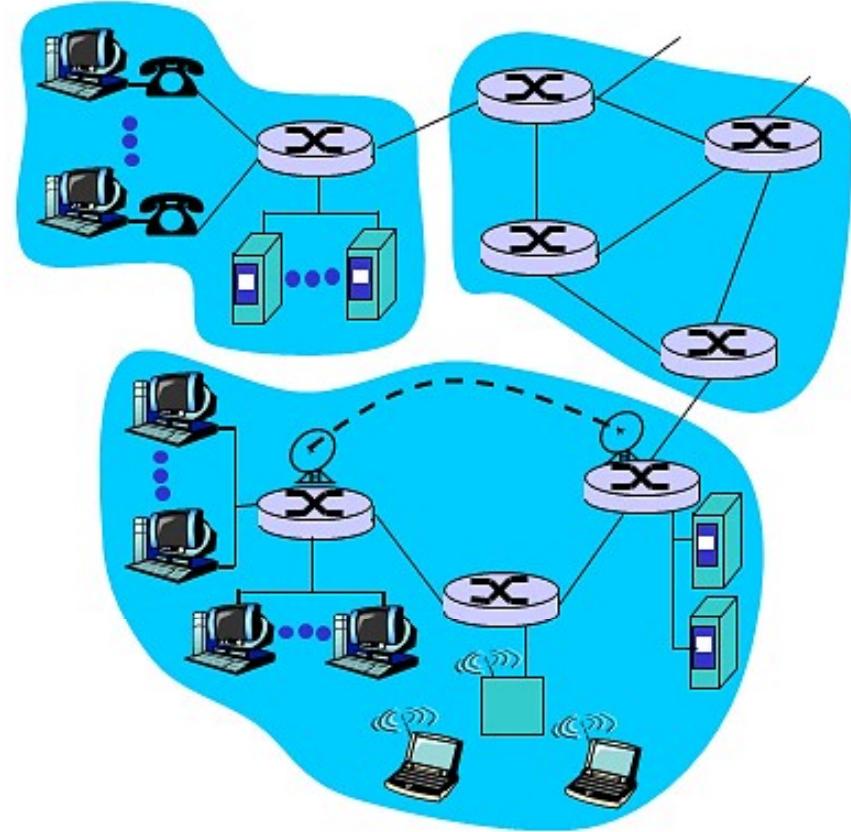
1.6 Delay & loss in packet-switched networks

1.7 Protocol layers, service models

1.8 History

A closer look at network structure:

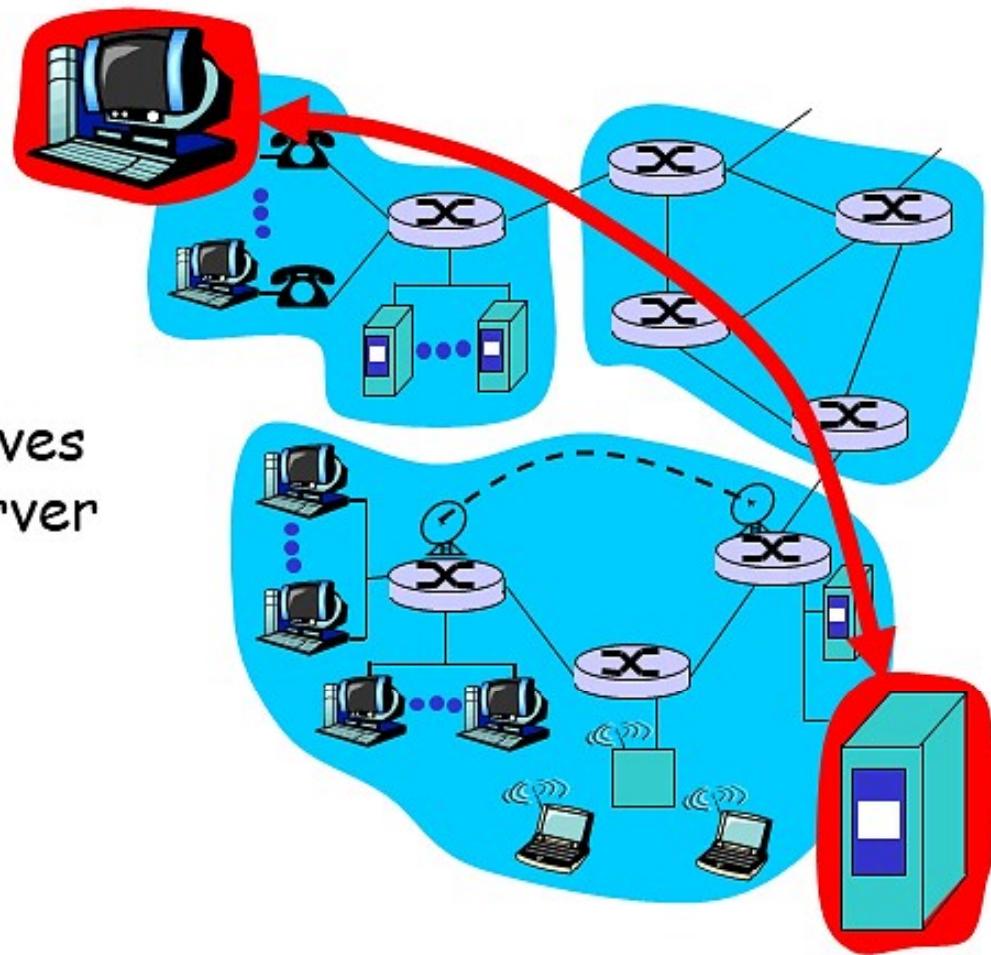
- **network edge:**
applications and hosts
- **network core:**
 - ❖ routers
 - ❖ network of networks
- **access networks, physical media:**
communication links



The network edge:

□ end systems (hosts):

- ❖ run application programs
- ❖ e.g. Web, email
- ❖ at "edge of network"



□ client/server model

- ❖ client host requests, receives service from always-on server
- ❖ e.g. Web browser/server; email client/server

□ peer-peer model:

- ❖ minimal (or no) use of dedicated servers
- ❖ e.g. Skype, BitTorrent

Network edge: connection-oriented service

- Goal:** data transfer between end systems
- **handshaking:** setup (prepare for) data transfer ahead of time
 - ❖ Hello, hello back human protocol
 - ❖ *set up "state"* in two communicating hosts
 - **TCP - Transmission Control Protocol**
 - ❖ Internet's connection-oriented service

TCP service [RFC 793]

- *reliable, in-order byte-stream data transfer*
 - ❖ loss: acknowledgements and retransmissions
- **flow control:**
 - ❖ sender won't overwhelm receiver
- **congestion control:**
 - ❖ senders "slow down sending rate" when network congested

Network edge: connectionless service

- Goal: data transfer between end systems
- ❖ same as before!
 - UDP** - User Datagram Protocol [RFC 768]:
 - ❖ connectionless
 - ❖ unreliable data transfer
 - ❖ no flow control
 - ❖ no congestion control

App's using TCP:

- HTTP (Web), FTP (file transfer), Telnet (remote login), SMTP (email)

App's using UDP:

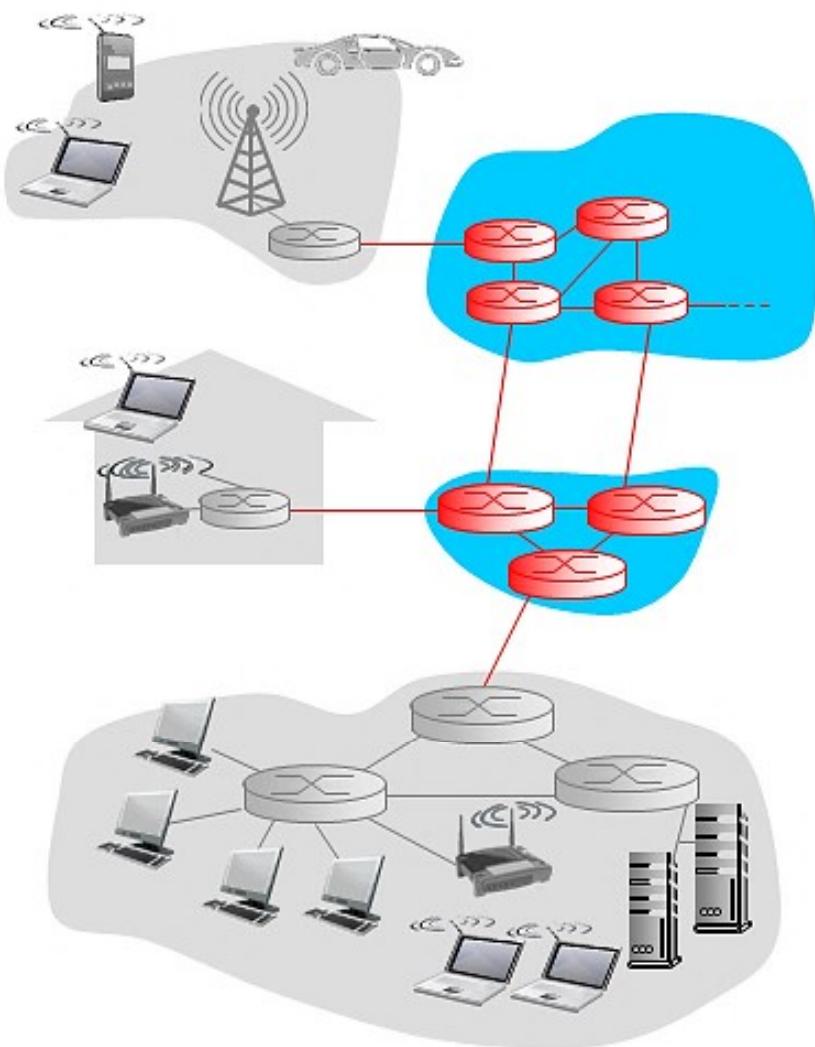
- streaming media, teleconferencing, DNS, Internet telephony

Chapter 1: roadmap

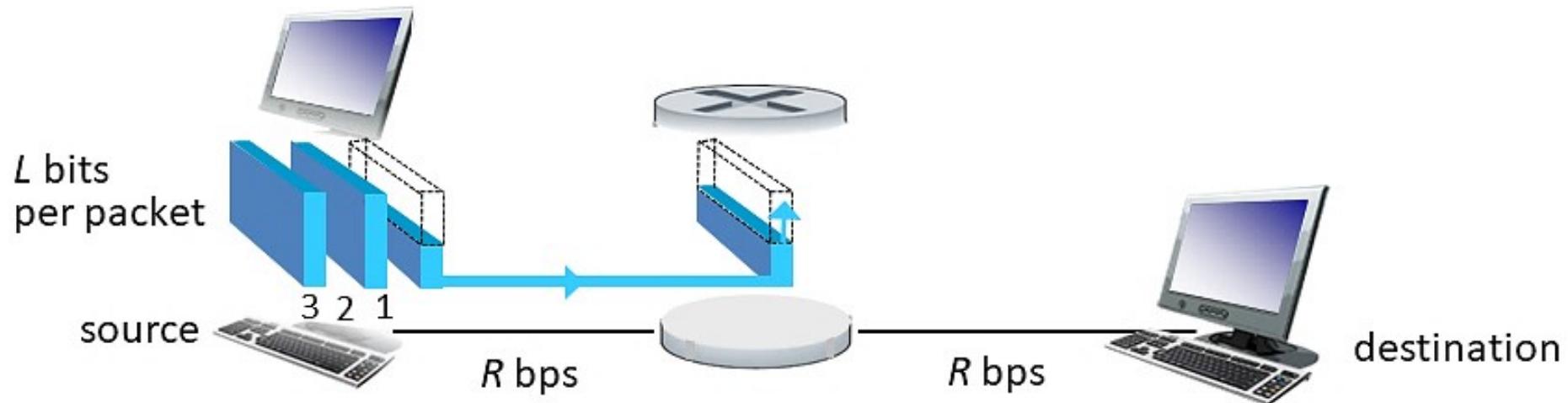
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The network core

- ❖ mesh of interconnected routers
- ❖ **packet-switching: hosts break application-layer messages into *packets***
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



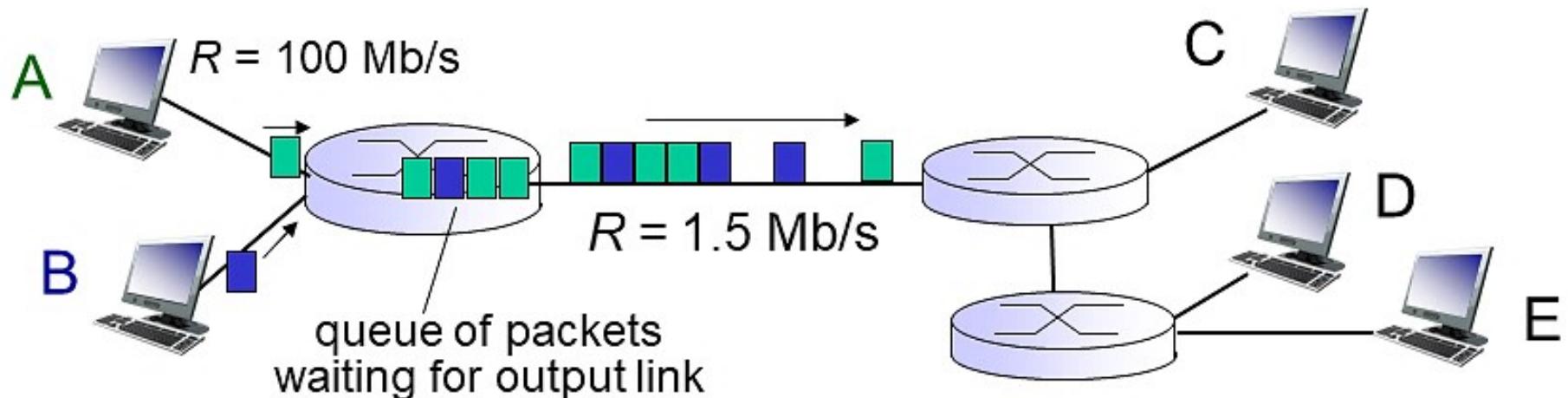
Packet-switching: store-and-forward



- ❖ takes L/R seconds to transmit (push out) L -bit packet into link at R bps
- ❖ ***store and forward:*** entire packet must arrive at router before it can be transmitted on next link
- ❖ end-end delay = $2L/R$ (assuming zero propagation delay)

- one-hop numerical example:*
- $L = 7.5 \text{ Mbits}$
 - $R = 1.5 \text{ Mbps}$
 - one-hop transmission delay = 5 sec
- } more on delay shortly ...

Packet Switching: queueing delay, loss



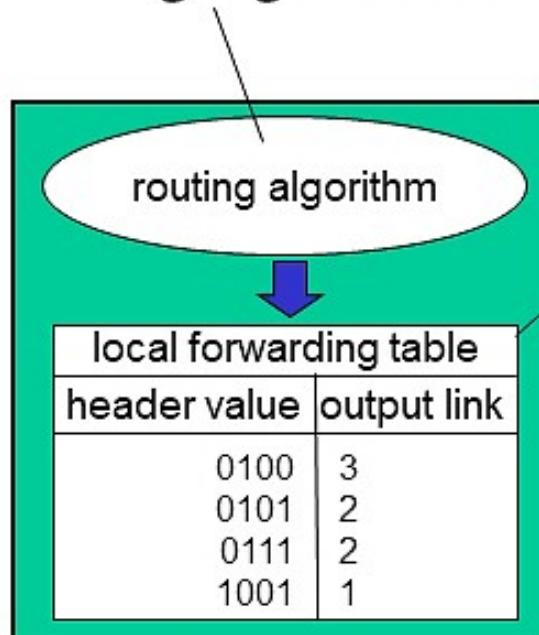
queueing and loss:

- ❖ If arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
 - packets will queue, wait to be transmitted on link
 - packets can be dropped (lost) if memory (buffer) fills up

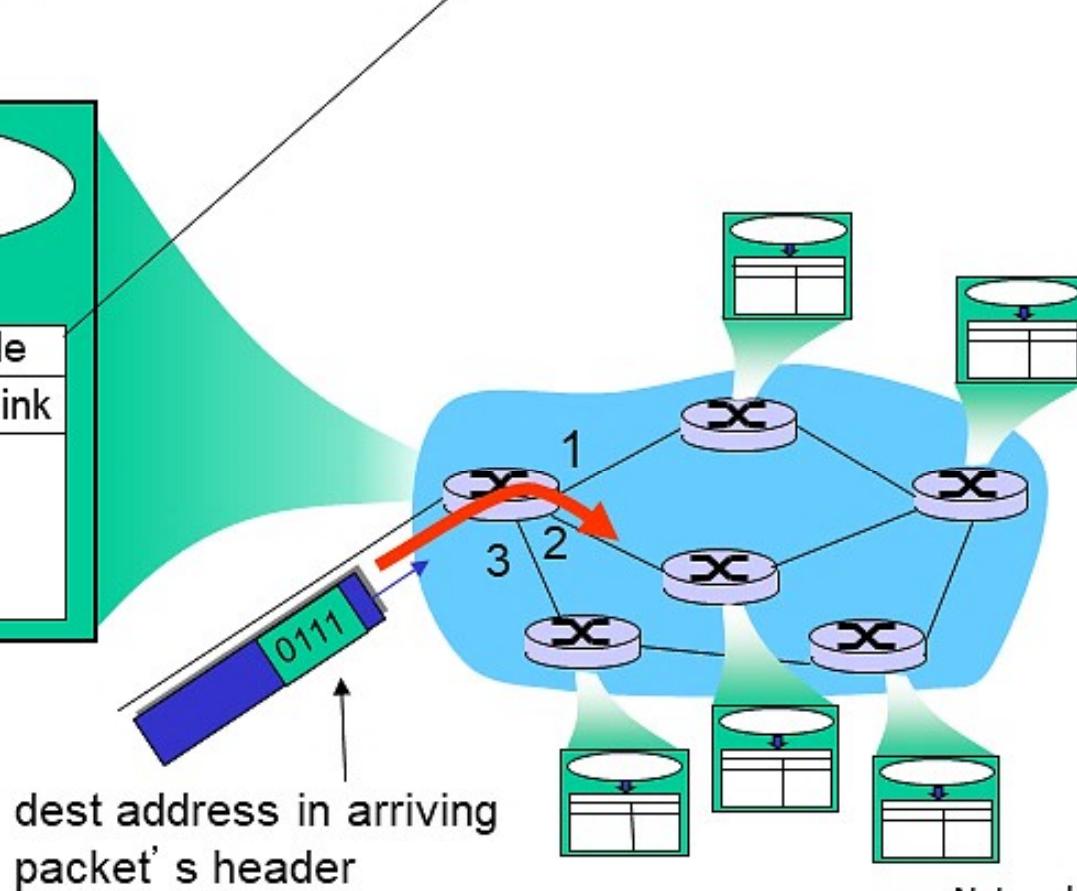
Two key network-core functions

routing: determines source-destination route taken by packets

- *routing algorithms*



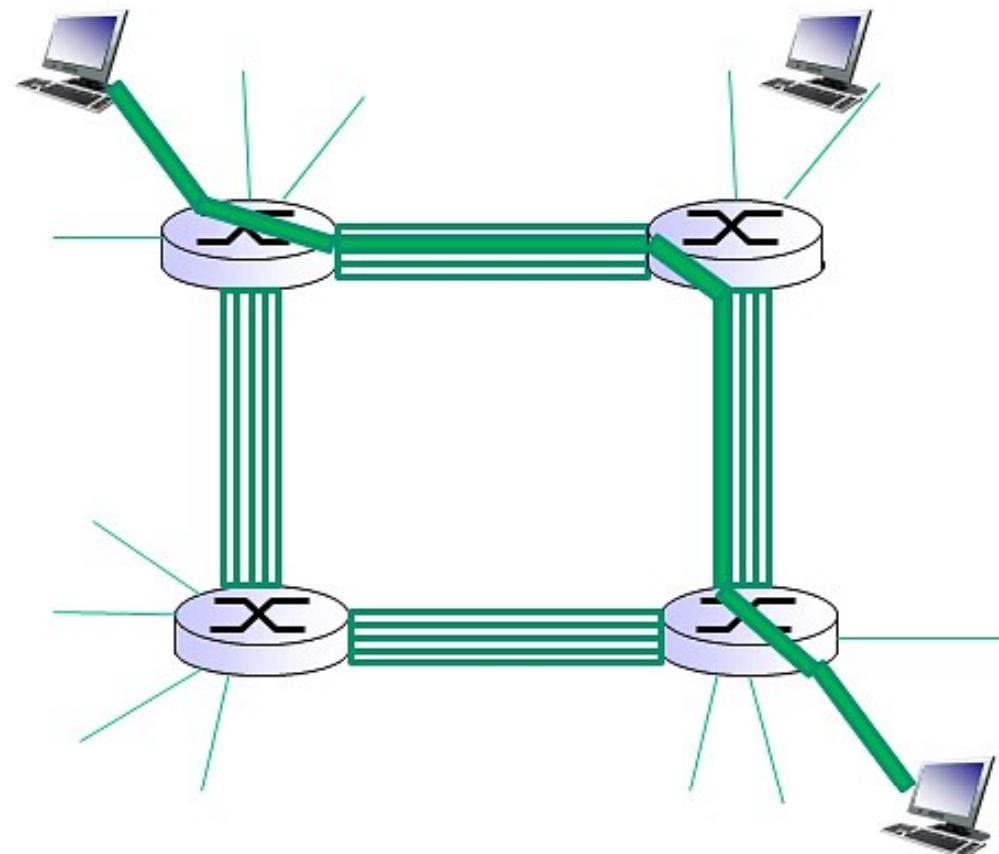
forwarding: move packets from router's input to appropriate router output



Alternative core: circuit switching

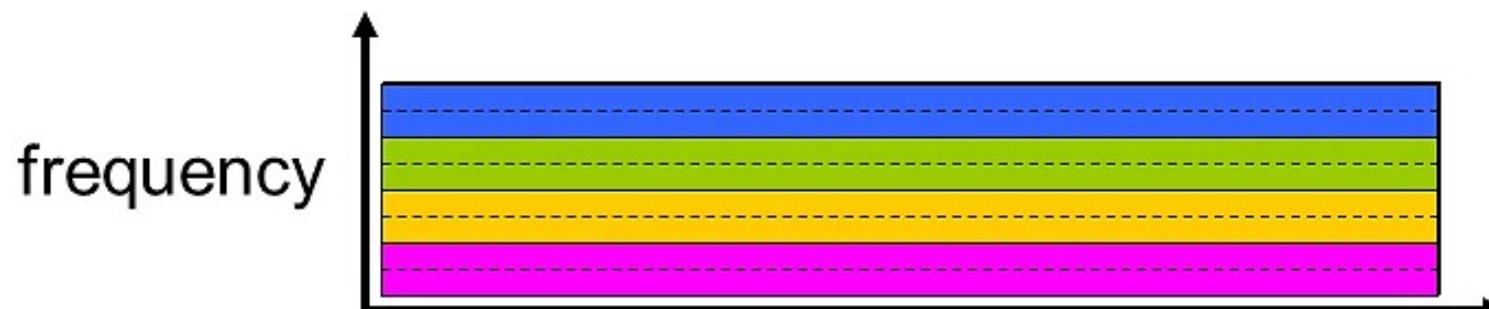
end-end resources allocated
to, reserved for “call”
between source & dest:

- ❖ In diagram, each link has four circuits.
 - call gets 2nd circuit in top link and 1st circuit in right link.
- ❖ dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- ❖ circuit segment idle if not used by call (*no sharing*)
- ❖ Commonly used in traditional telephone networks



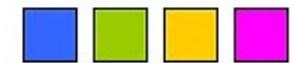
Circuit switching: FDM versus TDM

FDM

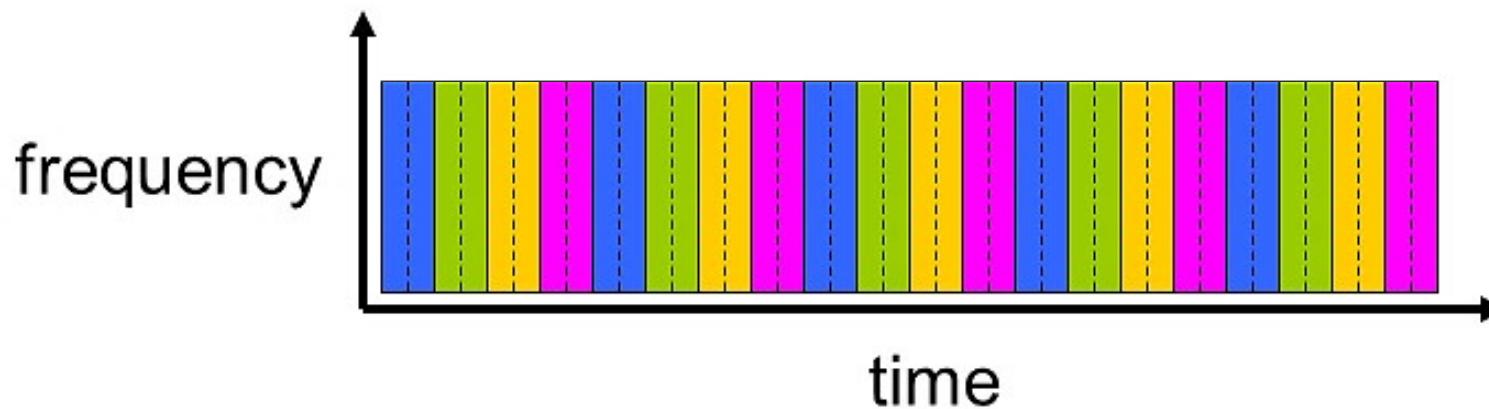


Example:

4 users



TDM

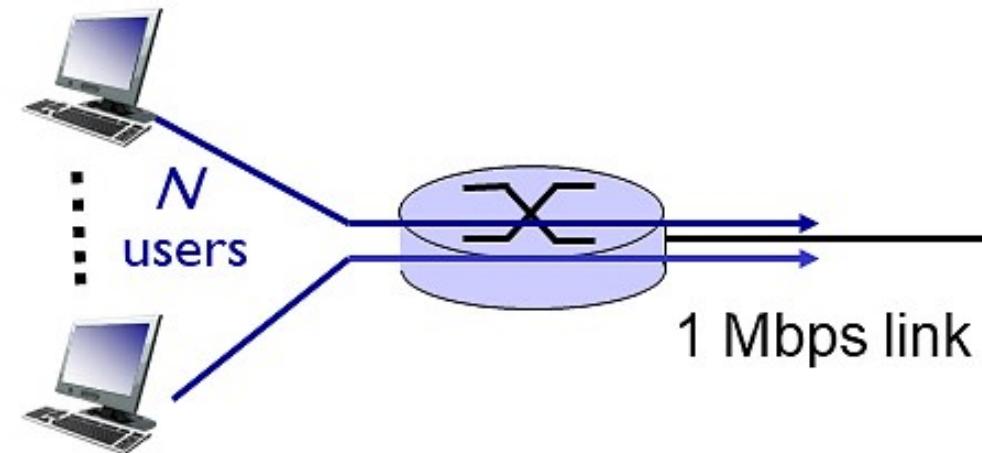


Packet switching versus circuit switching

packet switching allows more users to use network!

example:

- 1 Mb/s link
- each user:
 - 100 kb/s when “active”
 - active 10% of time



❖ *circuit-switching:*

- 10 users

❖ *packet switching:*

- with 35 users, probability > 10 active at same time is less than .0004 *

Q: how did we get value 0.0004?

Q: what happens if > 35 users ?

* Check out the online interactive exercises for more examples

Packet switching versus circuit switching

is packet switching a “slam dunk winner?”

- ❖ great for bursty data
 - resource sharing
 - simpler, no call setup
- ❖ **excessive congestion possible:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- ❖ **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 7)

Q: human analogies of reserved resources (circuit switching)
versus on-demand allocation (packet-switching)?

Chapter 1: roadmap

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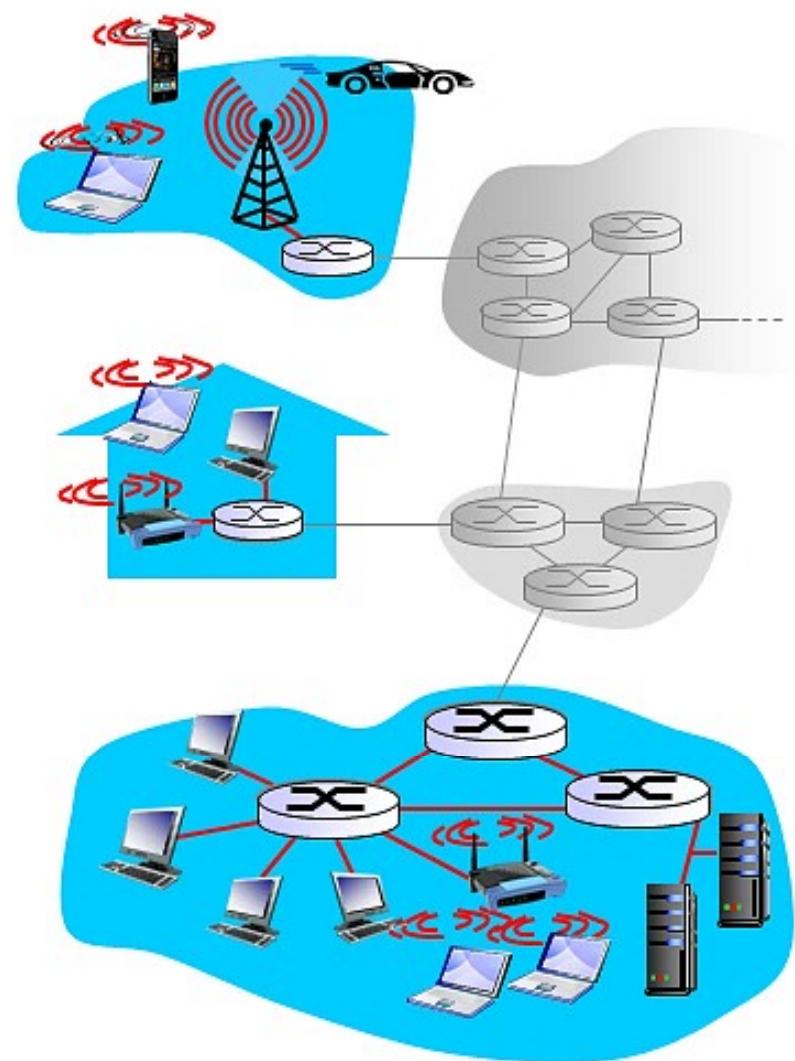
Access networks and physical media

Q: How to connect end systems to edge router?

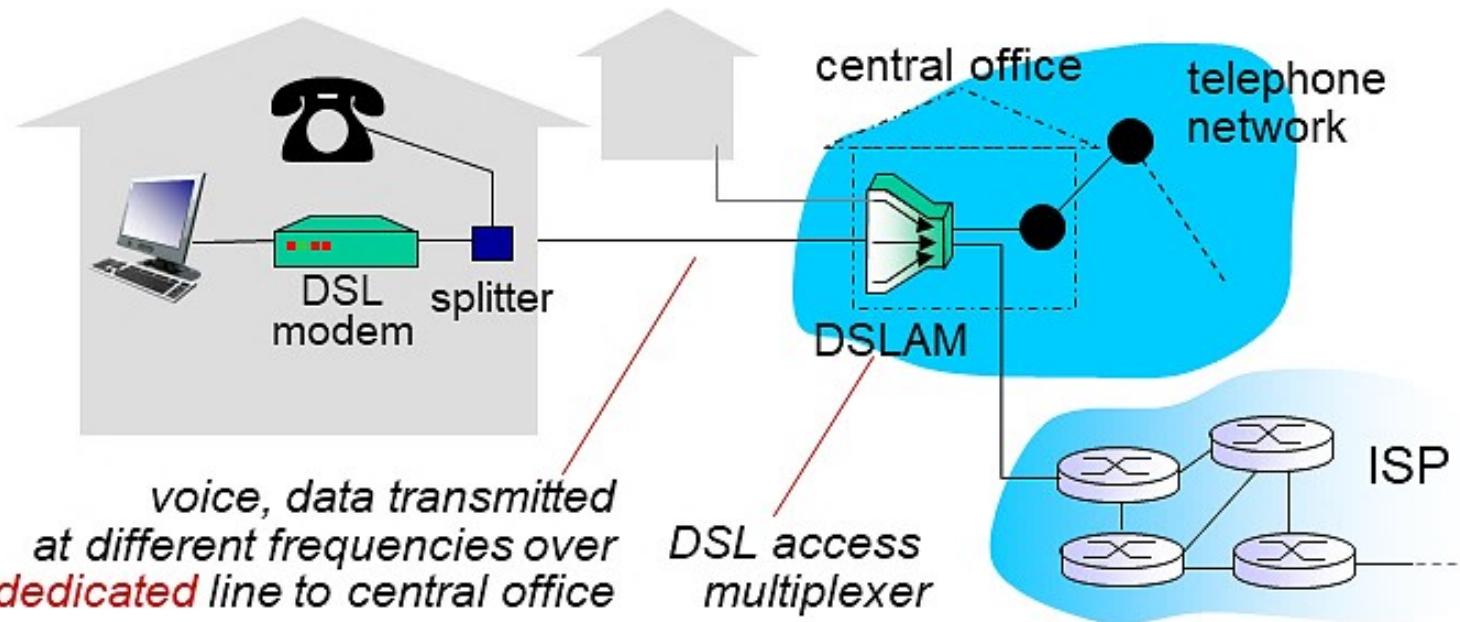
- ❖ residential access nets
- ❖ institutional access networks (school, company)
- ❖ mobile access networks

keep in mind:

- ❖ bandwidth (bits per second) of access network?
- ❖ shared or dedicated?

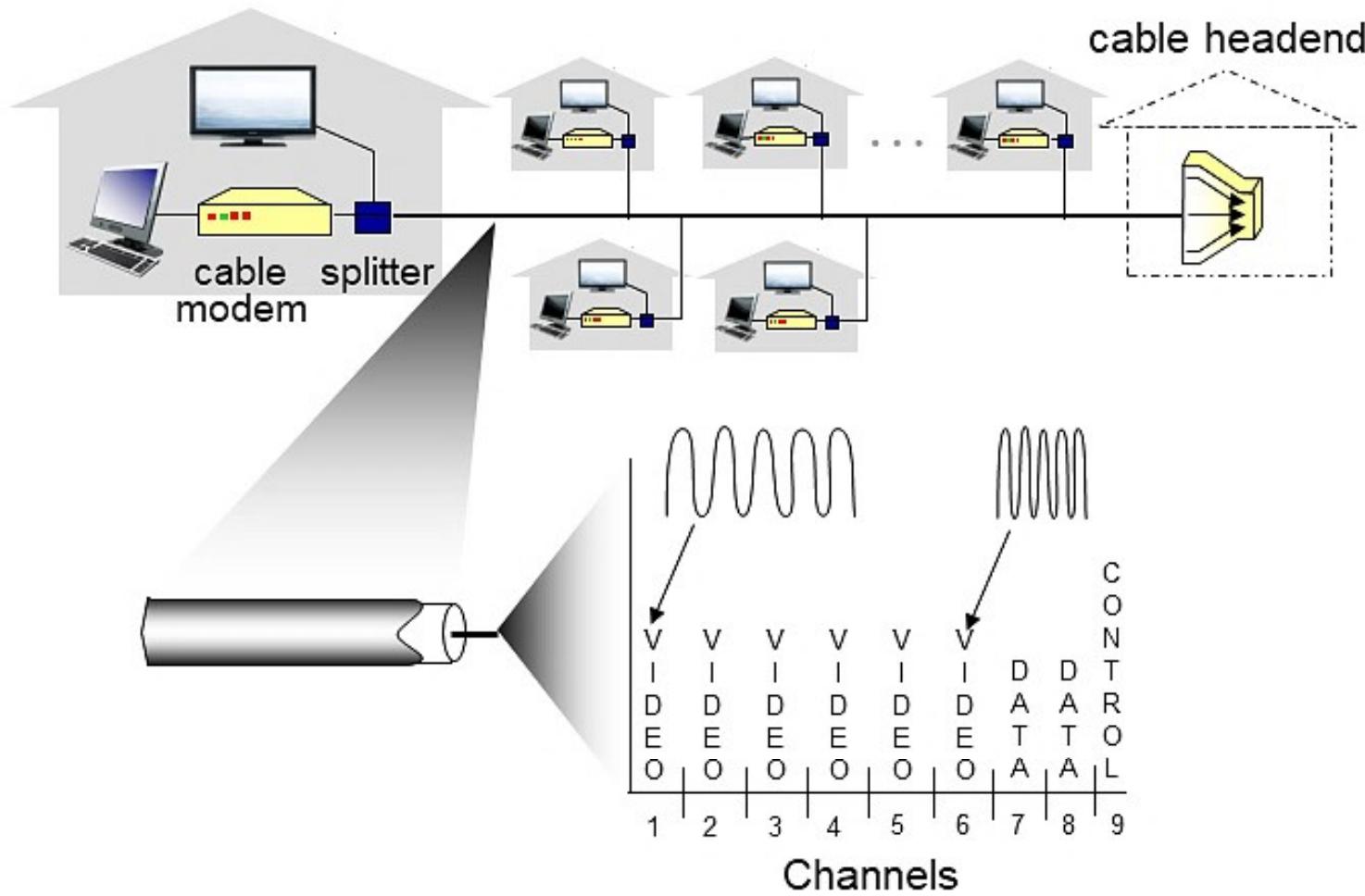


Access net: digital subscriber line (DSL)



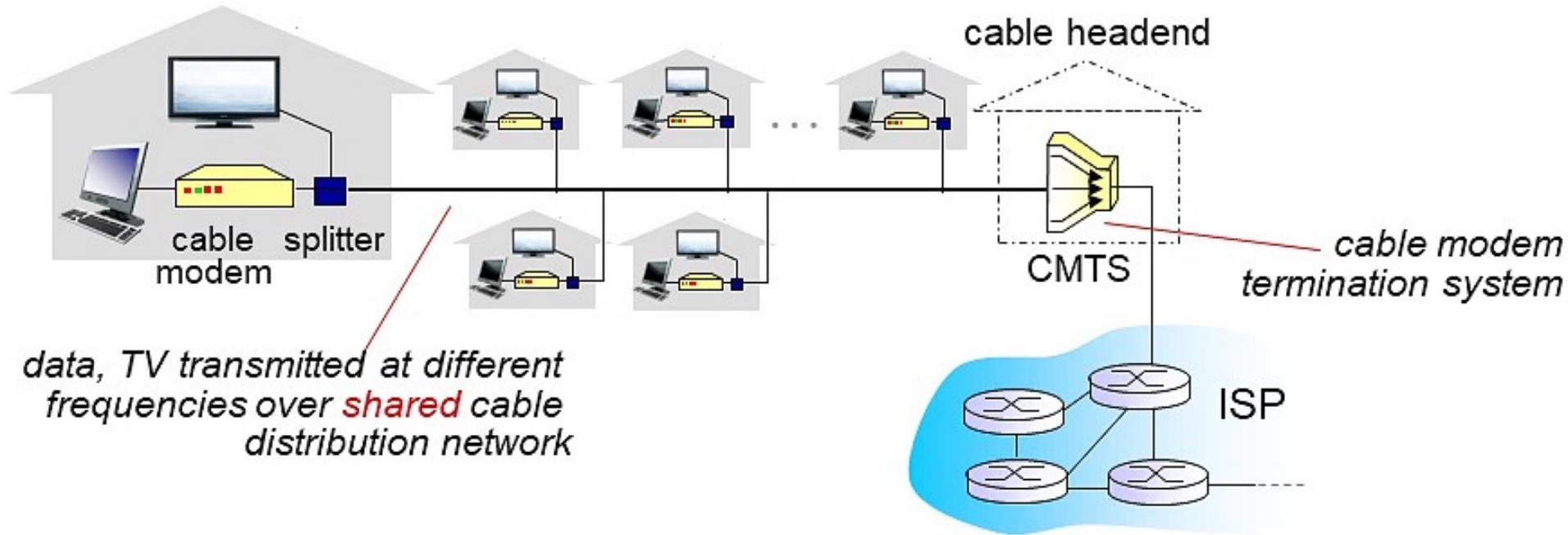
- ❖ use *existing* telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- ❖ < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)
- ❖ < 24 Mbps downstream transmission rate (typically < 10 Mbps)

Access net: cable network



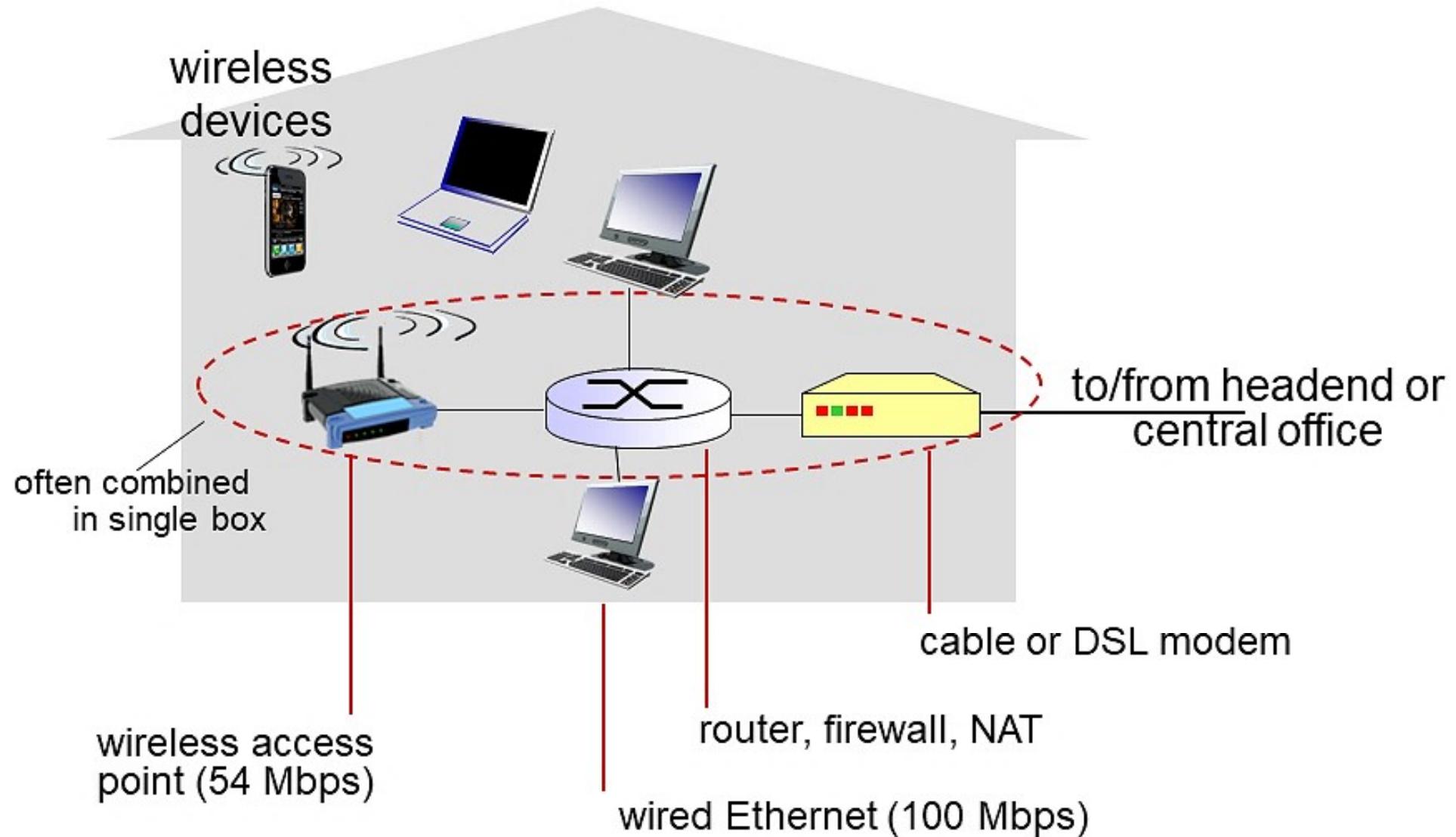
frequency division multiplexing: different channels transmitted in different frequency bands

Access net: cable network

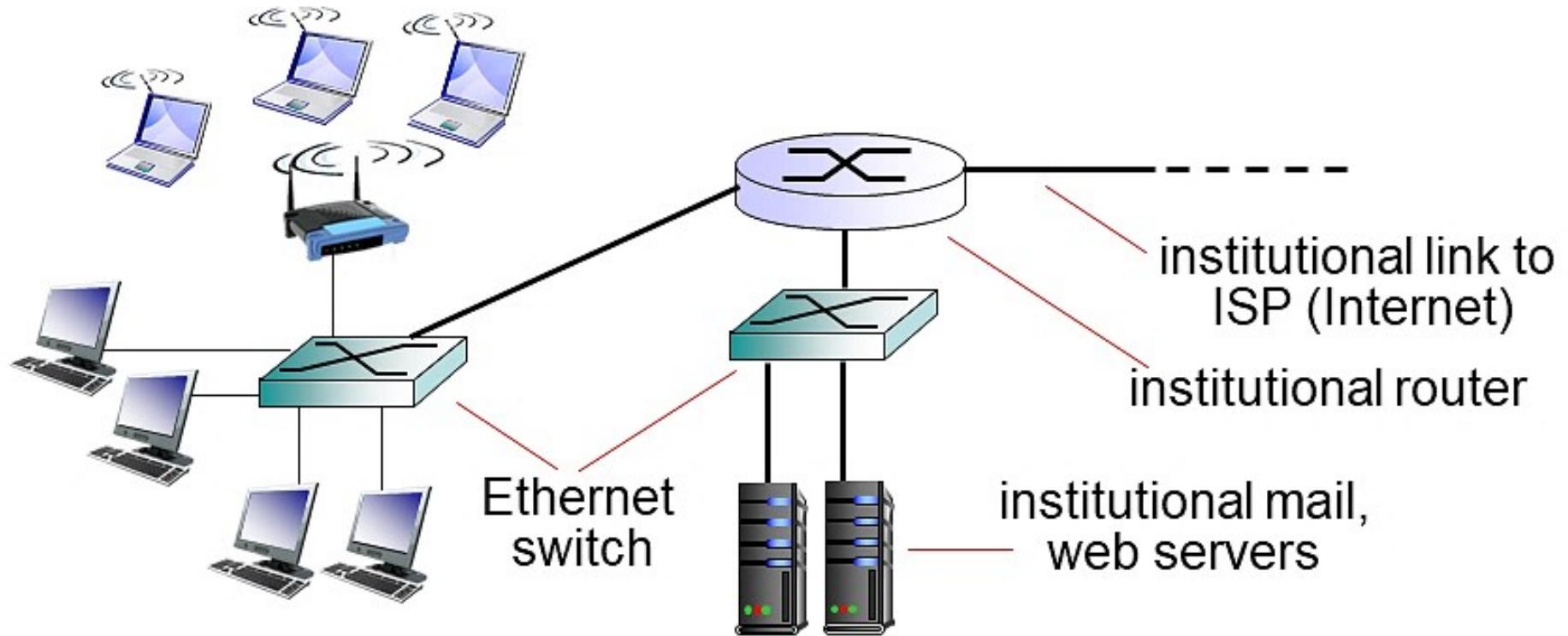


- ❖ HFC: hybrid fiber coax
 - asymmetric: up to 30Mbps downstream transmission rate, 2 Mbps upstream transmission rate
- ❖ network of cable, fiber attaches homes to ISP router
 - homes **share access network** to cable headend
 - unlike DSL, which has dedicated access to central office

Access net: home network



Enterprise access networks (Ethernet)



- ❖ typically used in companies, universities, etc
- ❖ 10 Mbps, 100Mbps, 1 Gbps, 10Gbps transmission rates
- ❖ today, end systems typically connect into Ethernet switch

Wireless access networks

- ❖ shared *wireless* access network connects end system to router
 - via base station aka “access point”

wireless LANs:

- within building (100 ft)
- 802.11b/g (WiFi): 11, 54 Mbps transmission rate



wide-area wireless access

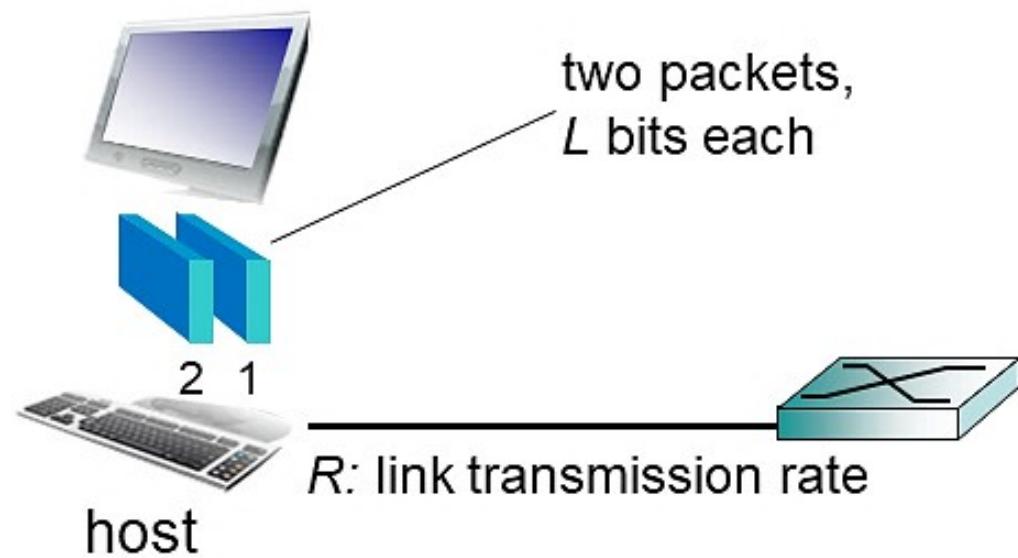
- provided by telco (cellular) operator, 10's km
- between 1 and 10 Mbps
- 3G, 4G: LTE



Host: sends *packets* of data

host sending function:

- ❖ takes application message
- ❖ breaks into smaller chunks, known as *packets*, of length L bits
- ❖ transmits packet into access network at *transmission rate R*
 - link transmission rate, aka link *capacity*, aka *link bandwidth*



$$\text{packet transmission delay} = \frac{\text{time needed to transmit } L\text{-bit packet into link}}{R \text{ (bits/sec)}}$$

Physical media

- ❖ **bit:** propagates between transmitter/receiver pairs
- ❖ **physical link:** what lies between transmitter & receiver
- ❖ **guided media:**
 - signals propagate in solid media: copper, fiber, coax
- ❖ **unguided media:**
 - signals propagate freely, e.g., radio

twisted pair (TP)

- ❖ two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps



Physical media: coax, fiber

coaxial cable:

- ❖ two concentric copper conductors
- ❖ bidirectional
- ❖ broadband:
 - multiple channels on cable
 - HFC



fiber optic cable:

- ❖ glass fiber carrying light pulses, each pulse a bit
- ❖ high-speed operation:
 - high-speed point-to-point transmission (e.g., 10's-100's Gbps transmission rate)
- ❖ low error rate:
 - repeaters spaced far apart
 - immune to electromagnetic noise



Physical media: radio

- ❖ signal carried in electromagnetic spectrum
- ❖ no physical “wire”
- ❖ bidirectional
- ❖ propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

radio link types:

- ❖ terrestrial microwave
 - e.g. up to 45 Mbps channels
- ❖ LAN (e.g., WiFi)
 - 11Mbps, 54 Mbps
- ❖ wide-area (e.g., cellular)
 - 3G cellular: ~ few Mbps
- ❖ satellite
 - Kbps to 45Mbps channel (or multiple smaller channels)
 - 270 msec end-end delay
 - geosynchronous versus low altitude

Chapter 1: roadmap

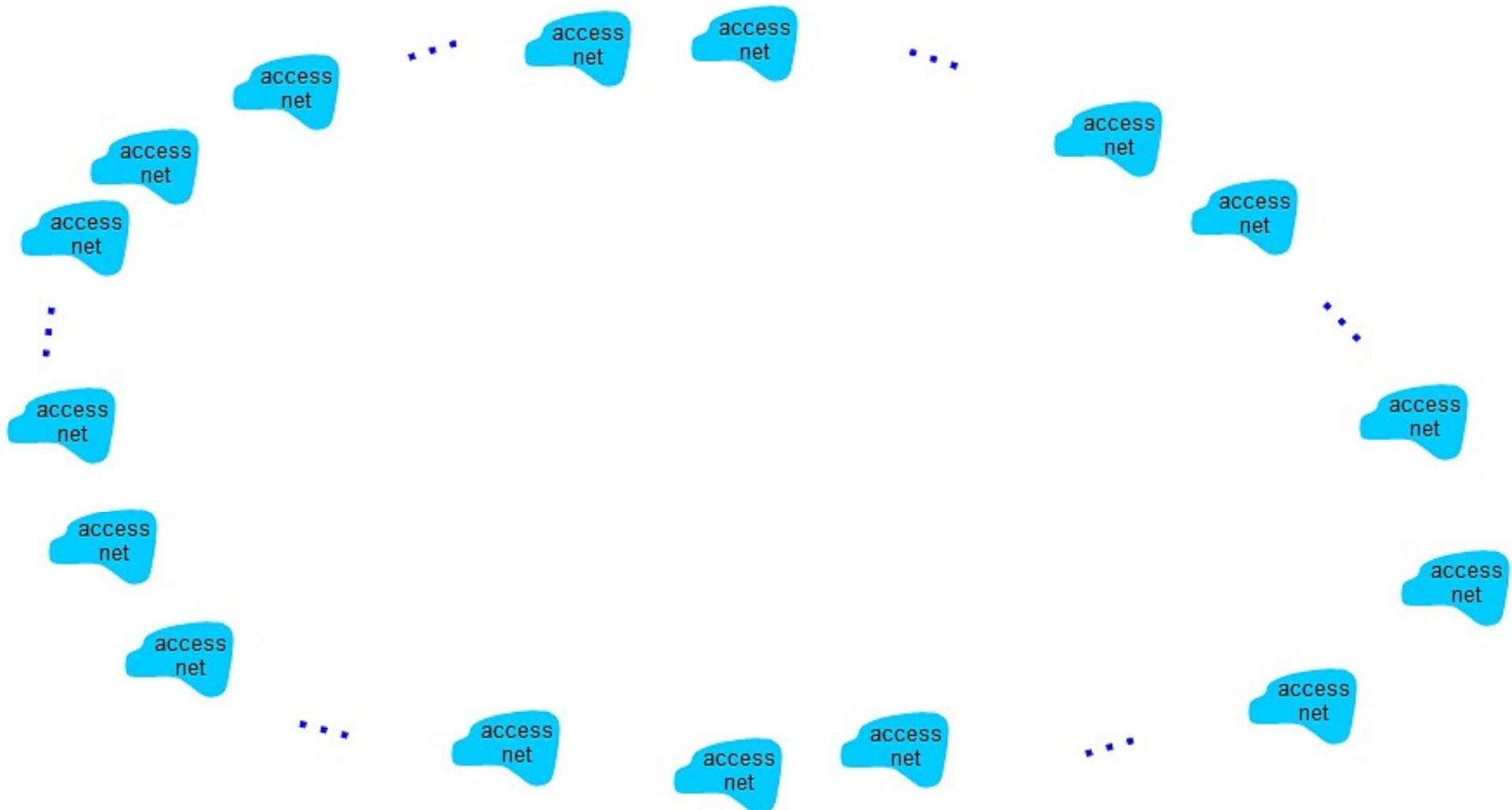
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Internet structure: network of networks

- ❖ End systems connect to Internet via **access ISPs** (Internet Service Providers)
 - Residential, company and university ISPs
- ❖ Access ISPs in turn must be interconnected.
 - ❖ So that any two hosts can send packets to each other
- ❖ Resulting network of networks is very complex
 - ❖ Evolution was driven by **economics** and **national policies**
- ❖ Let's take a stepwise approach to describe current Internet structure

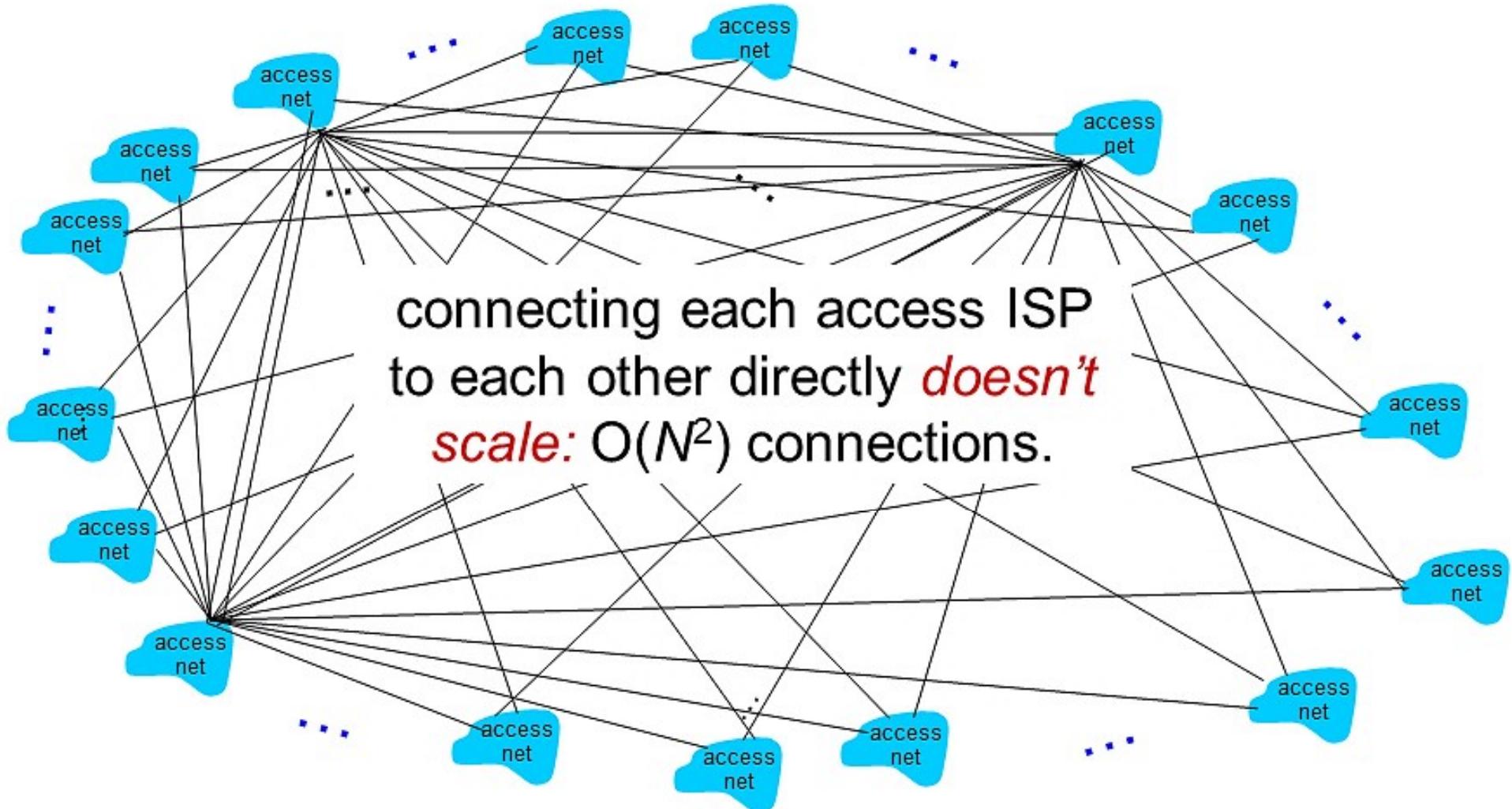
Internet structure: network of networks

Question: given *millions* of access ISPs, how to connect them together?



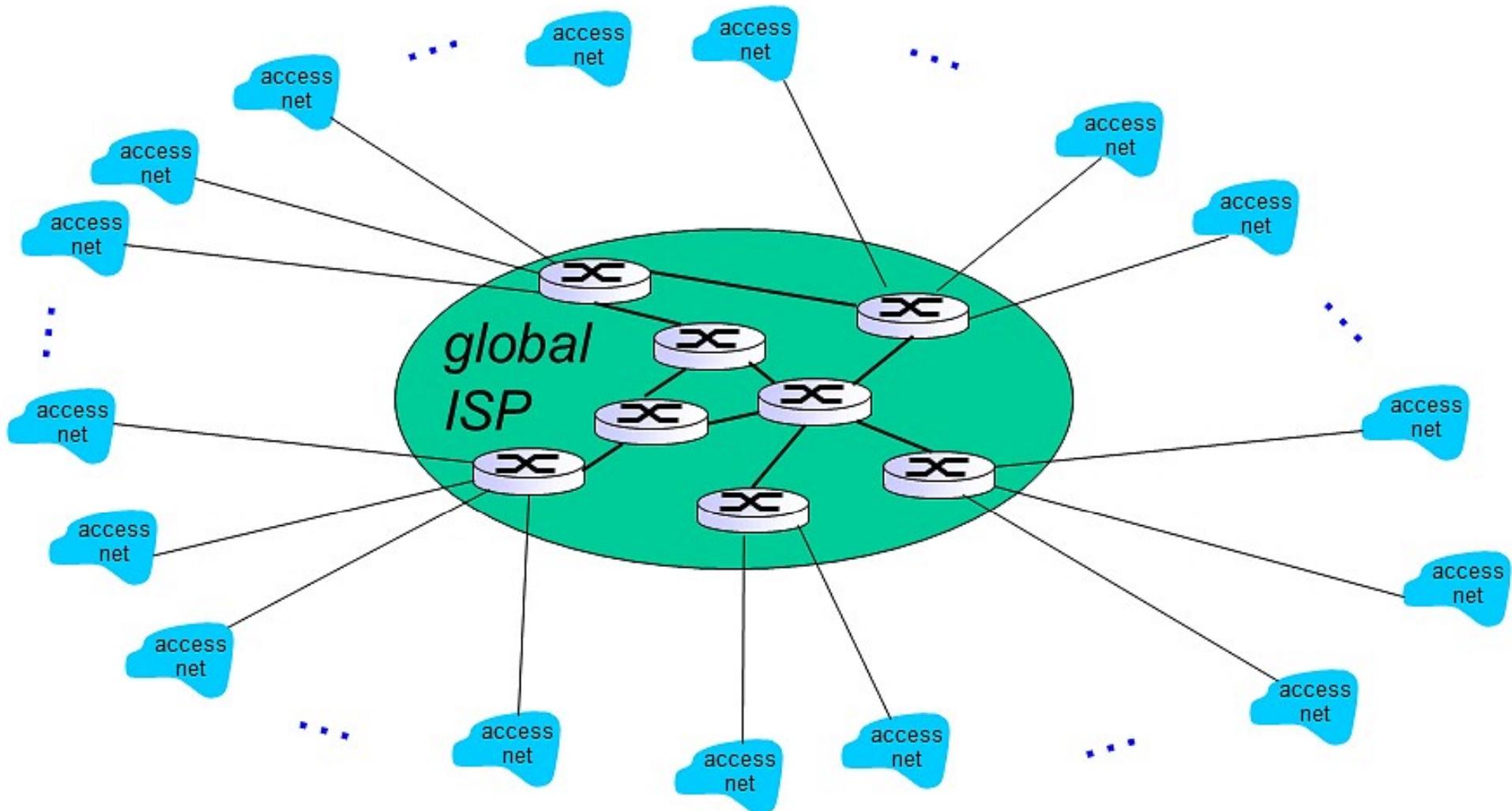
Internet structure: network of networks

Option: connect each access ISP to every other access ISP?



Internet structure: network of networks

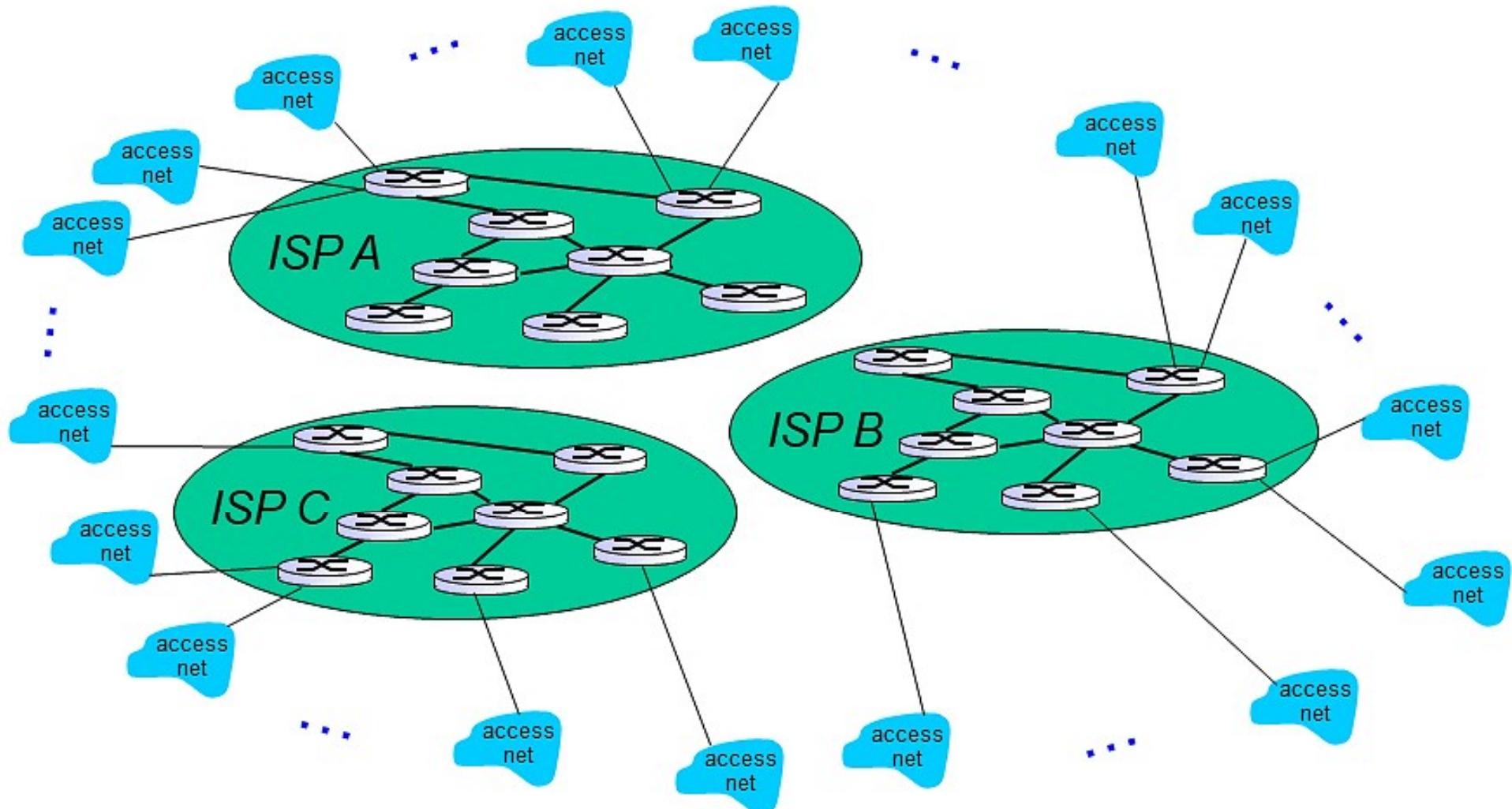
*Option: connect each access ISP to a global transit ISP?
Customer and provider ISPs have economic agreement.*



Internet structure: network of networks

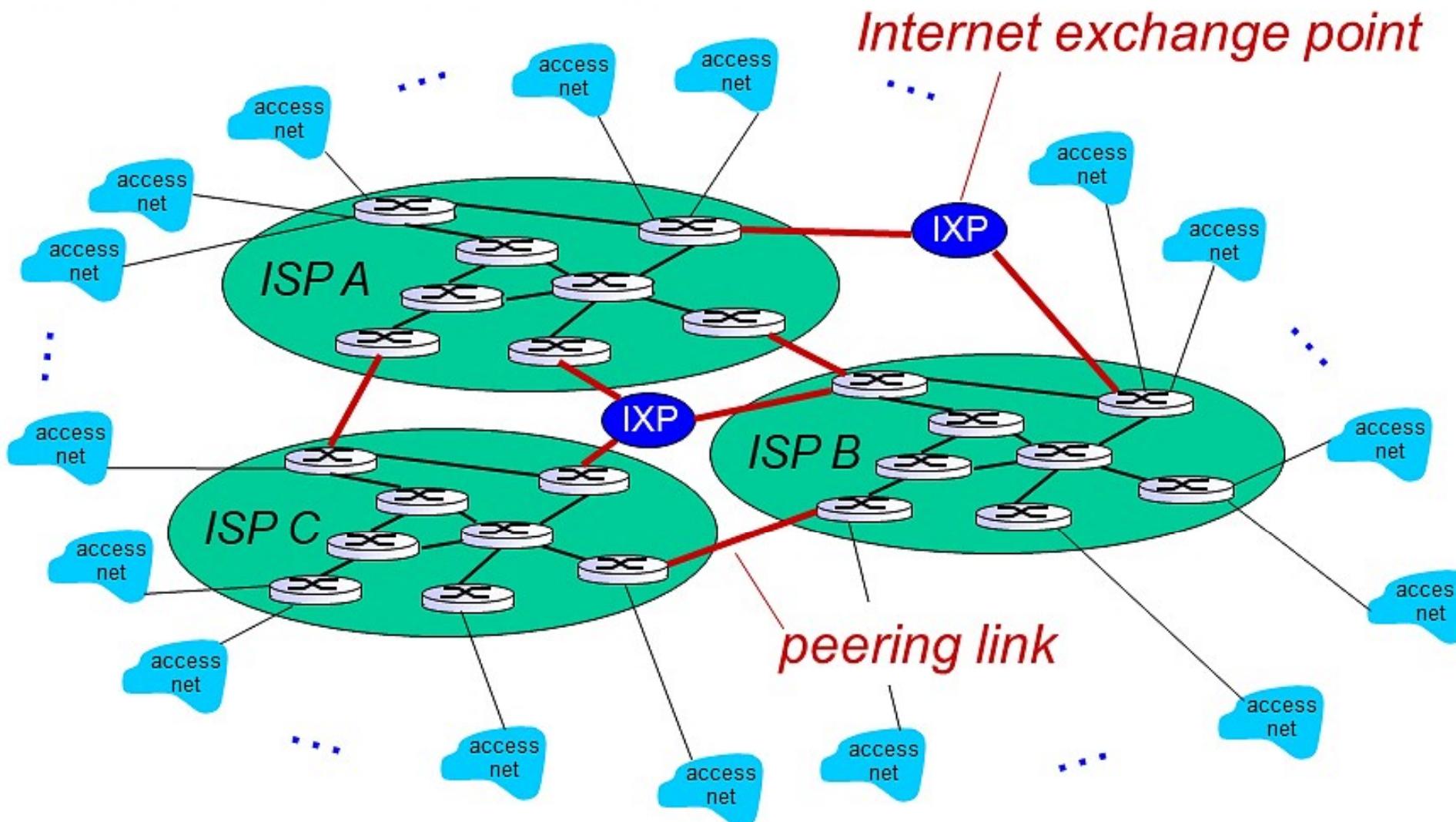
But if one global ISP is viable business, there will be competitors

....



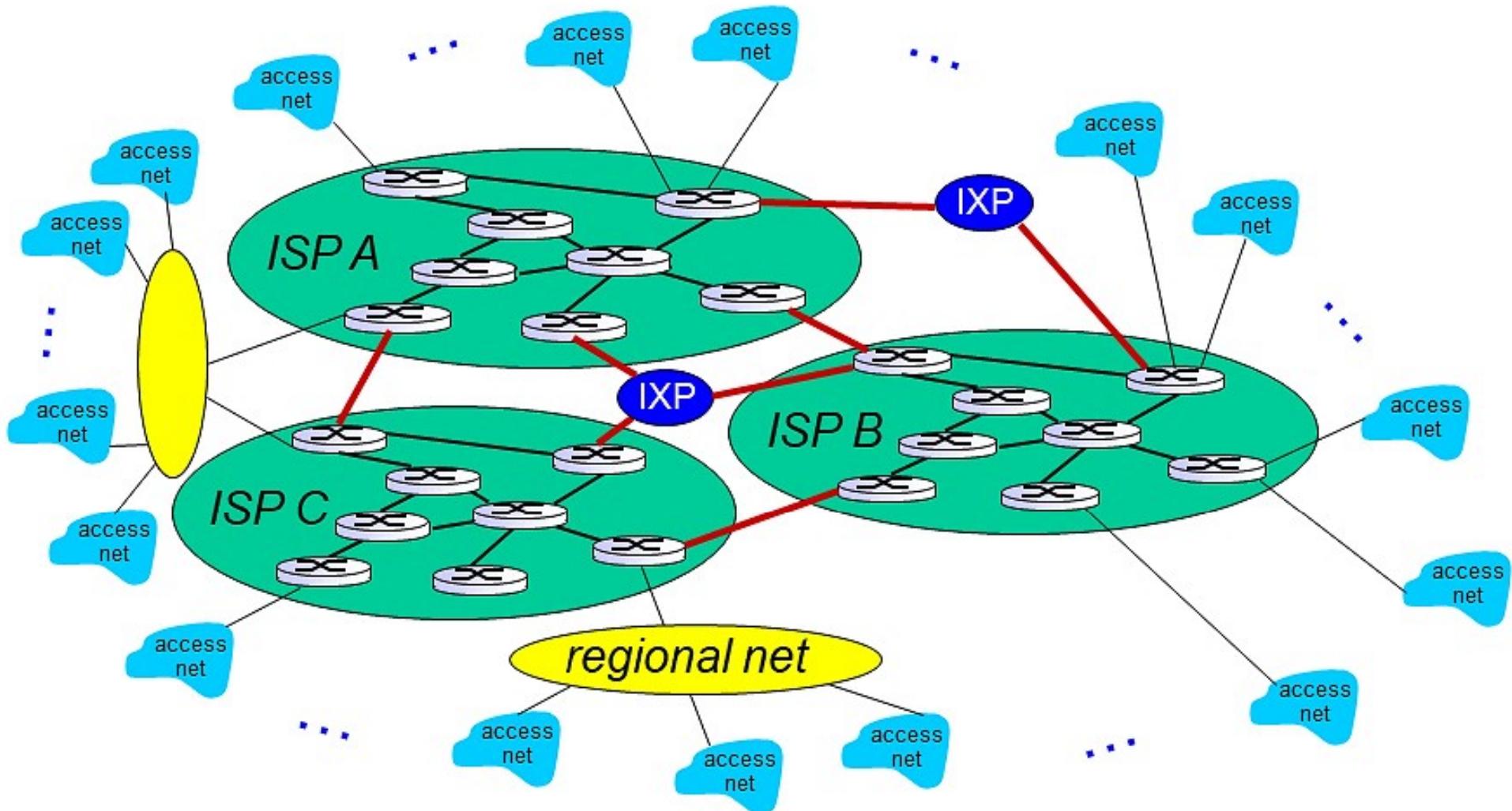
Internet structure: network of networks

But if one global ISP is viable business, there will be competitors
.... which must be interconnected



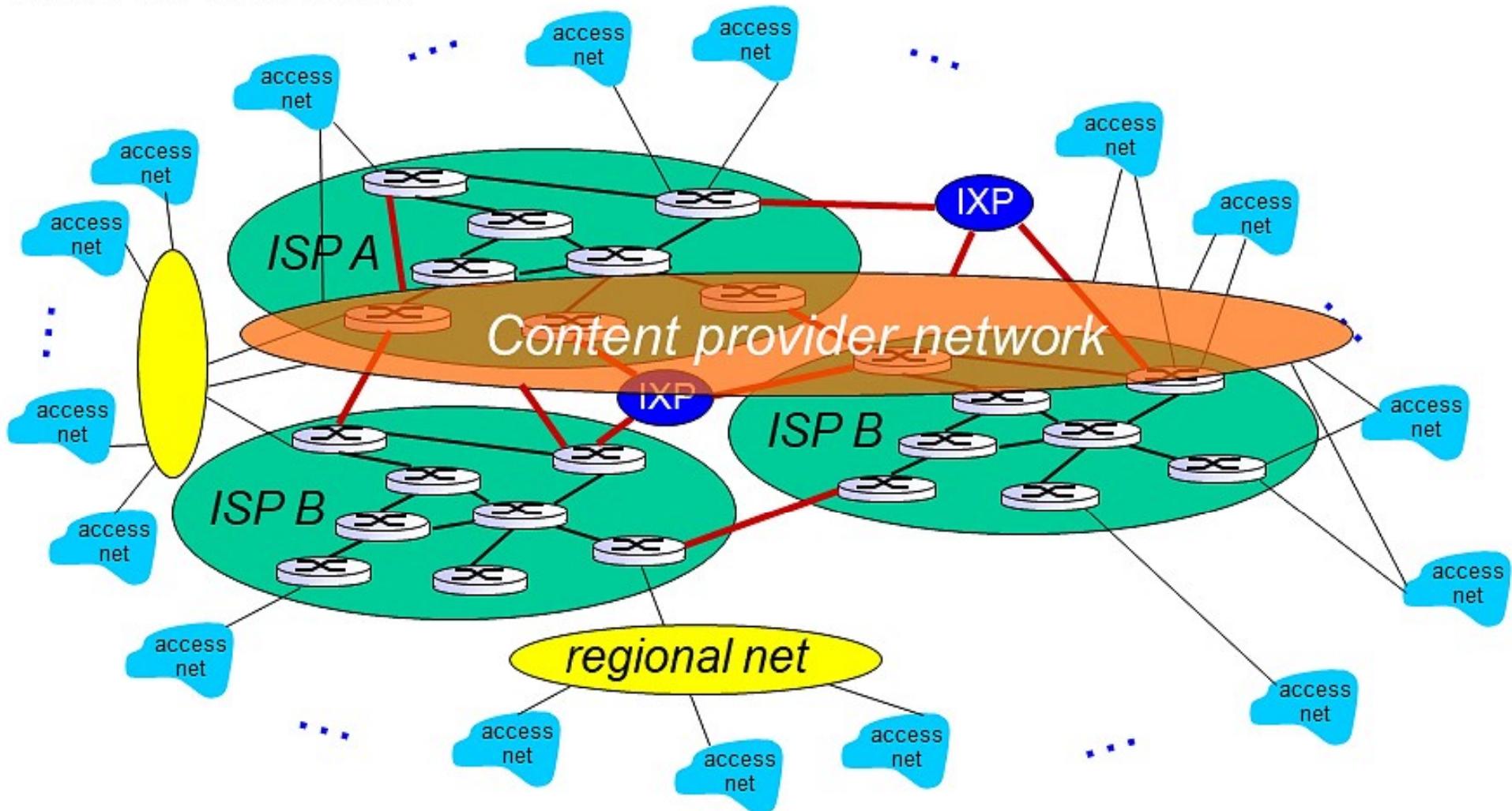
Internet structure: network of networks

... and regional networks may arise to connect access nets to ISPS

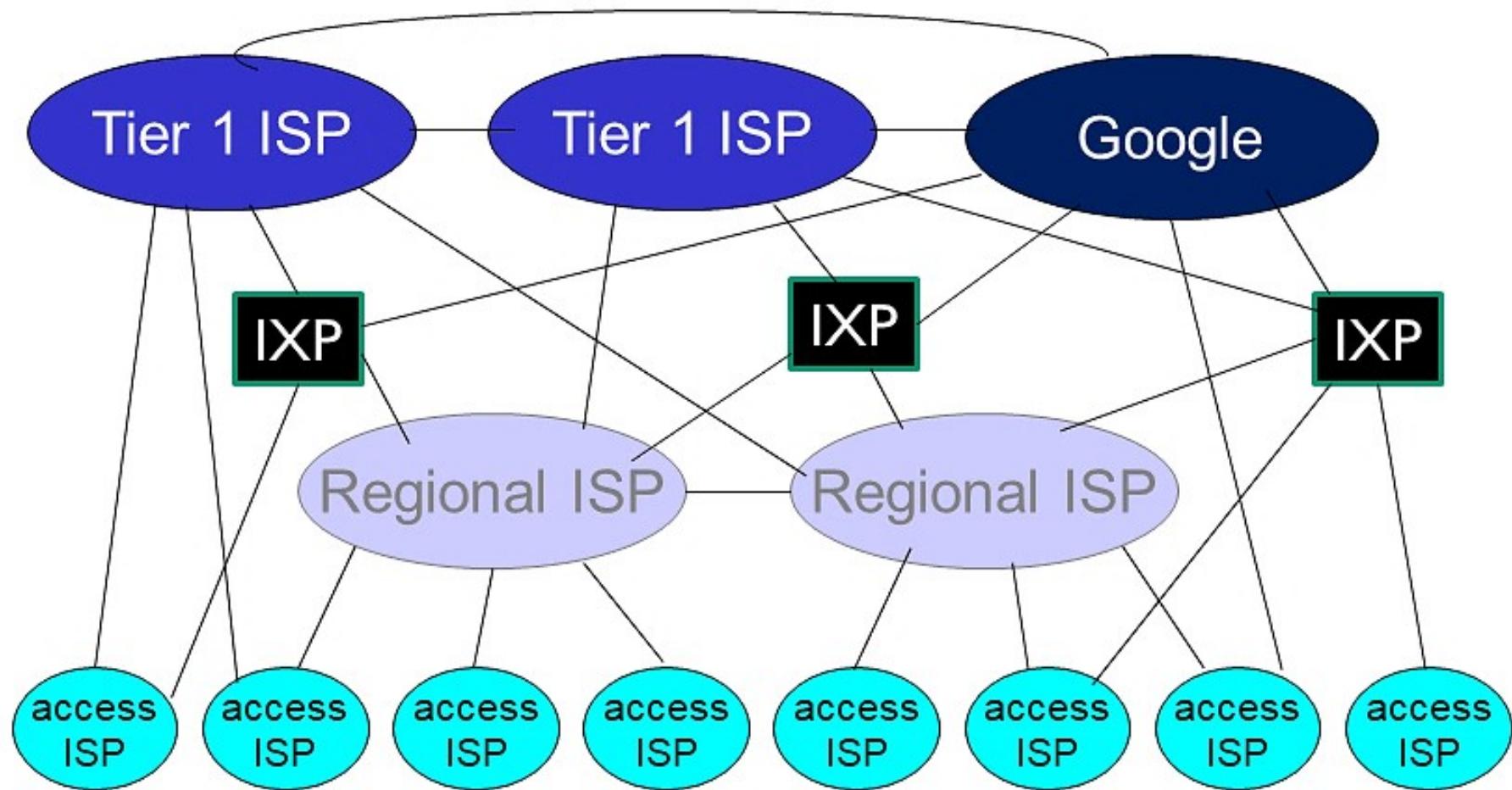


Internet structure: network of networks

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users

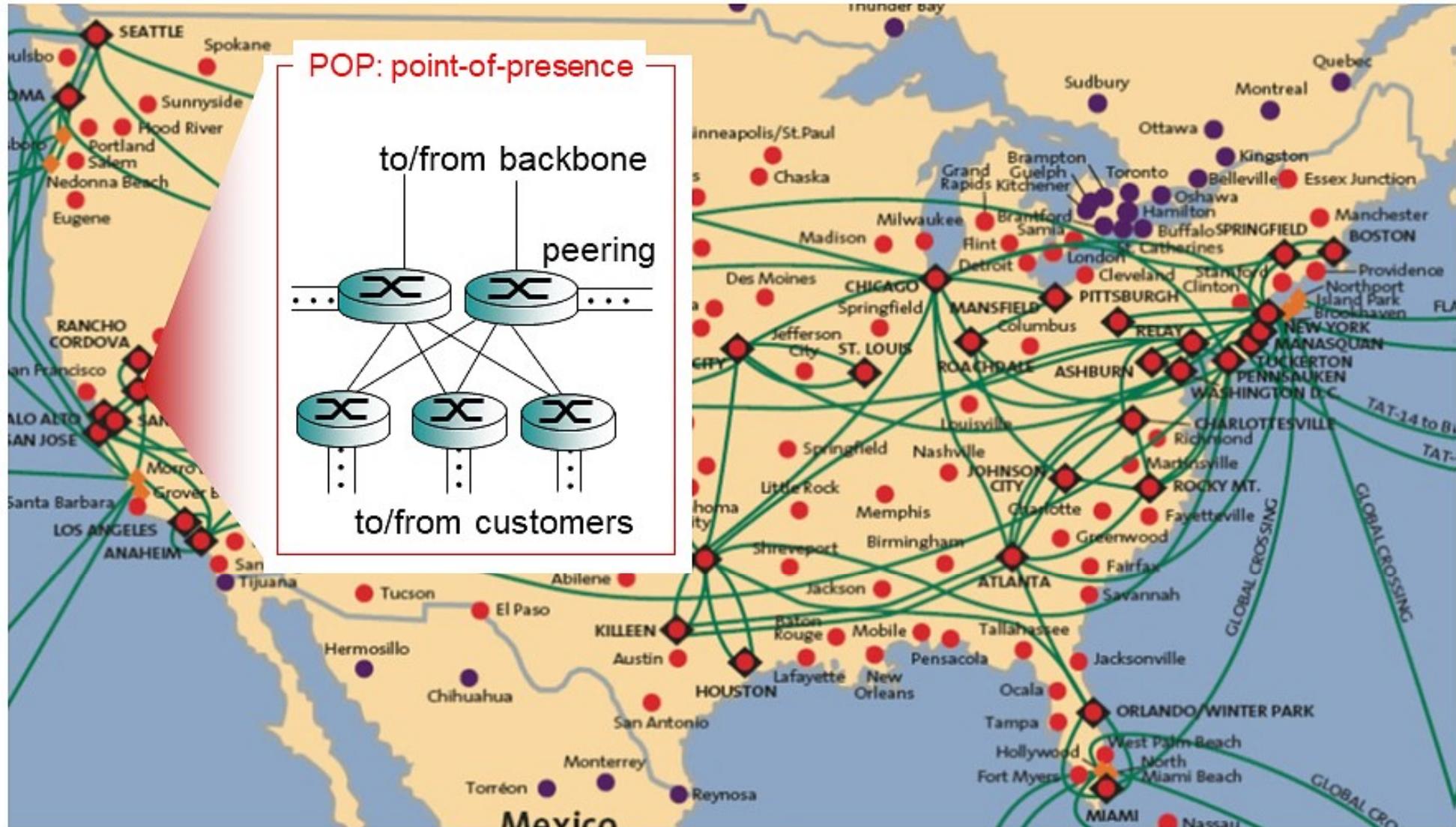


Internet structure: network of networks

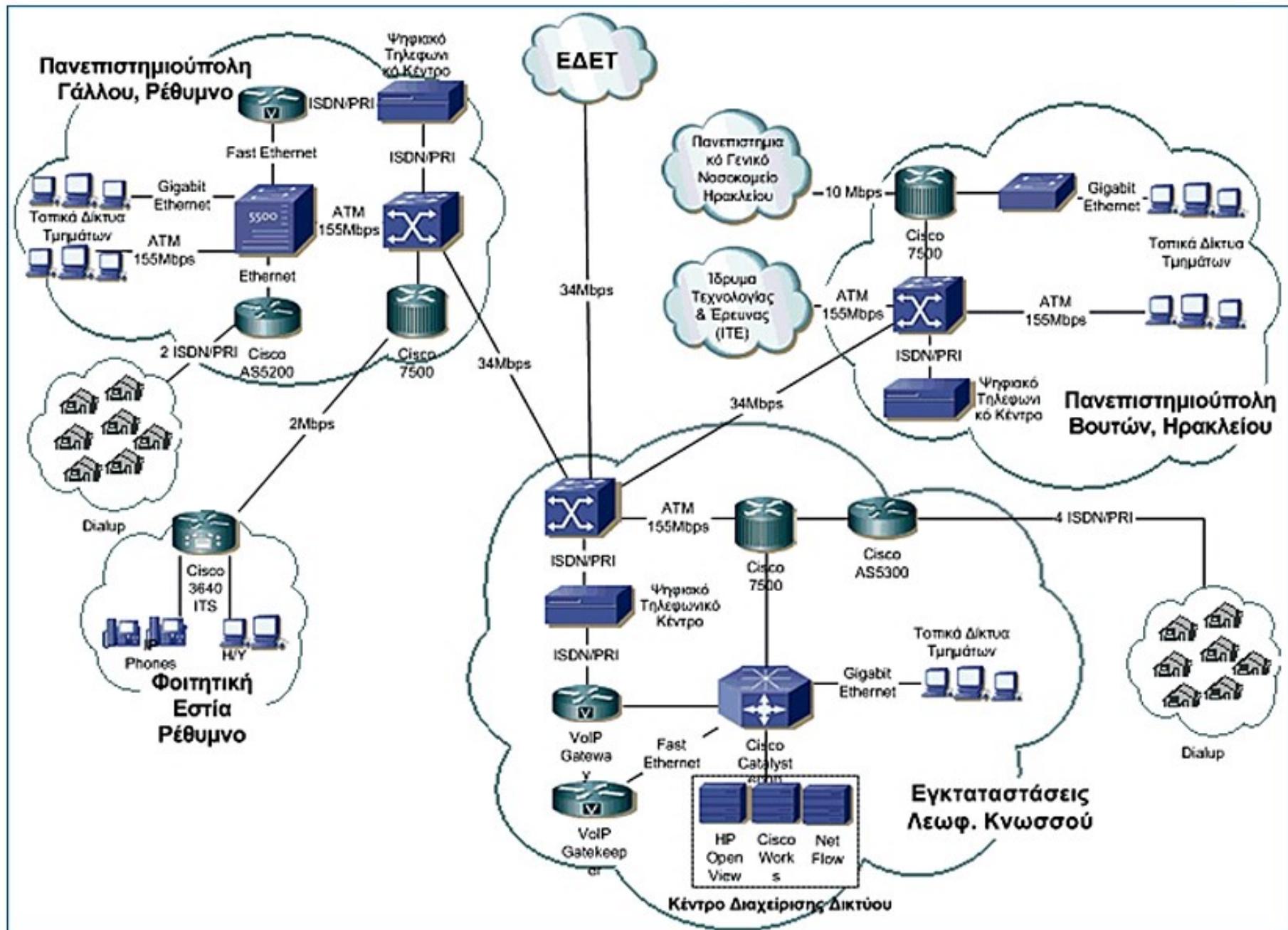


- ❖ at center: small # of well-connected large networks
 - “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - content provider network (e.g, Google): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

Tier-I ISP: e.g., Sprint



Univ of Crete Campus Network



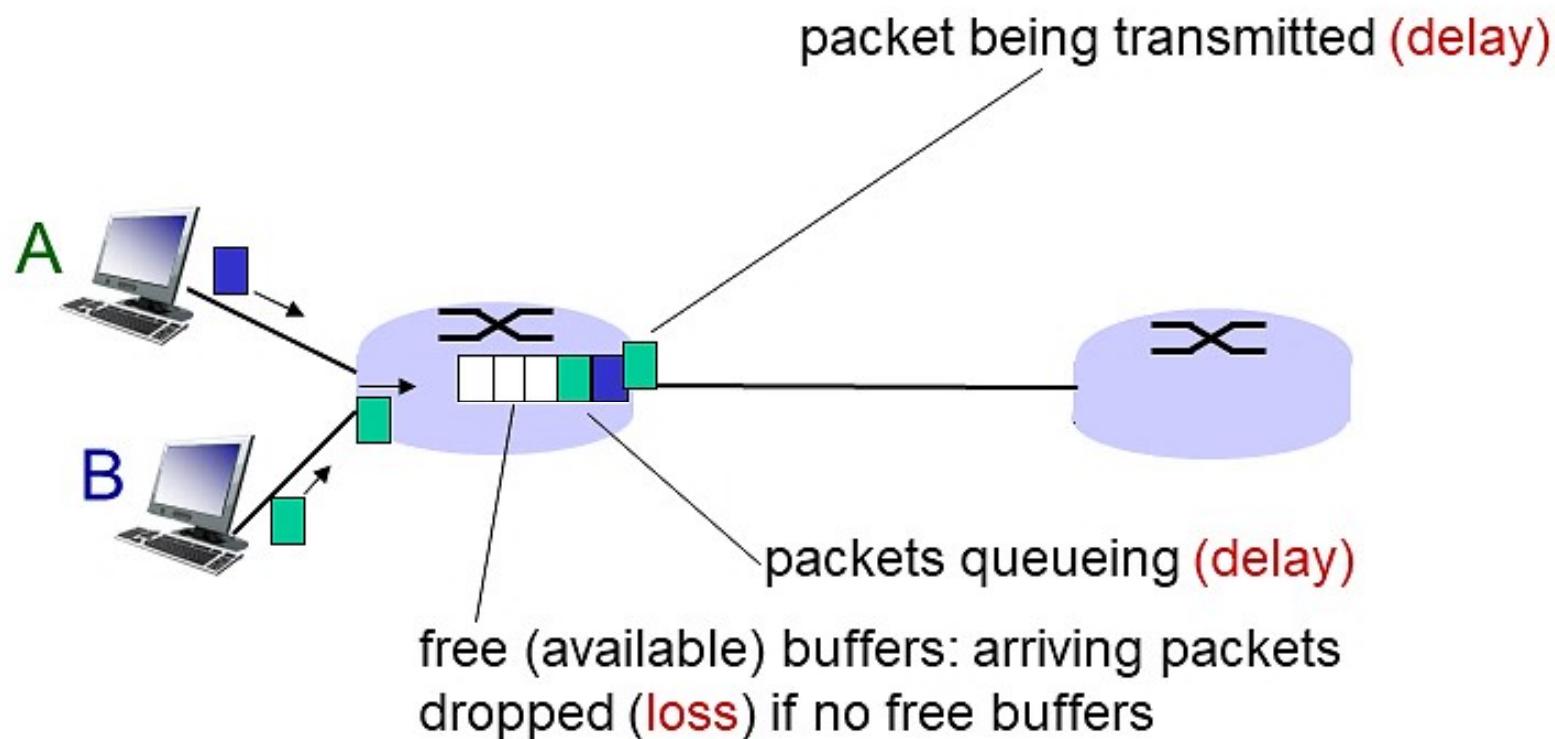
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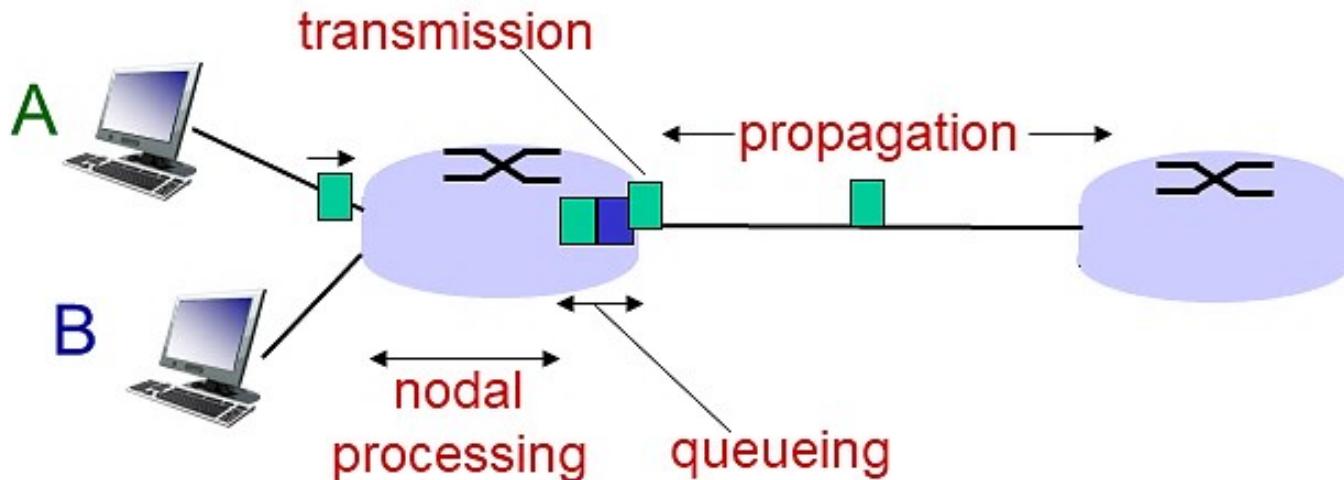
How do loss and delay occur?

packets *queue* in router buffers

- ❖ packet arrival rate to link (temporarily) exceeds output link capacity
- ❖ packets queue, wait for turn



Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

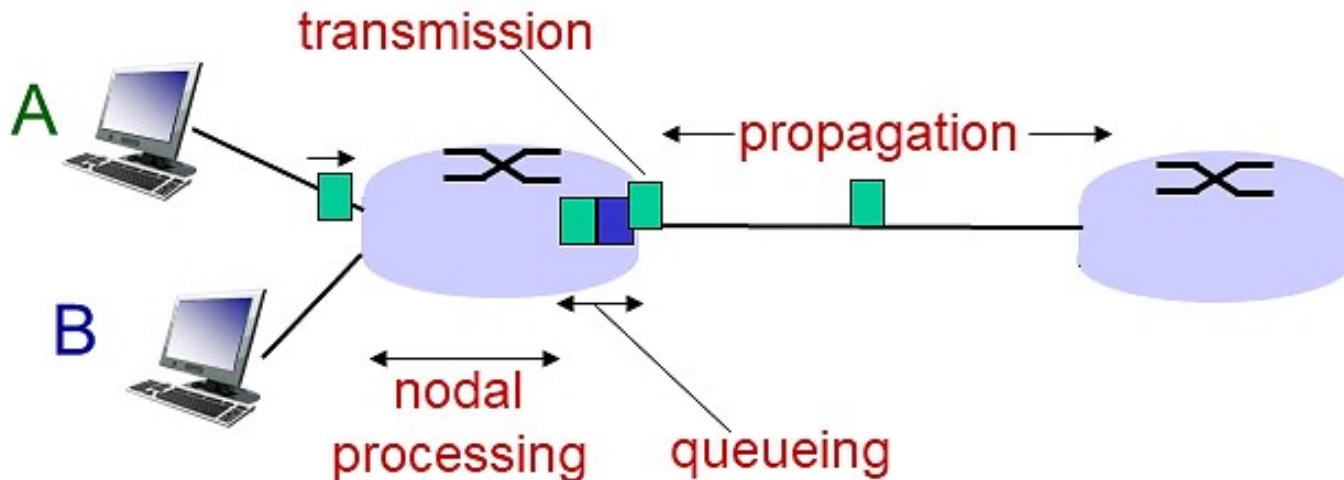
d_{proc} : nodal processing

- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L : packet length (bits)
- R : link *bandwidth (bps)*
- $d_{\text{trans}} = L/R$

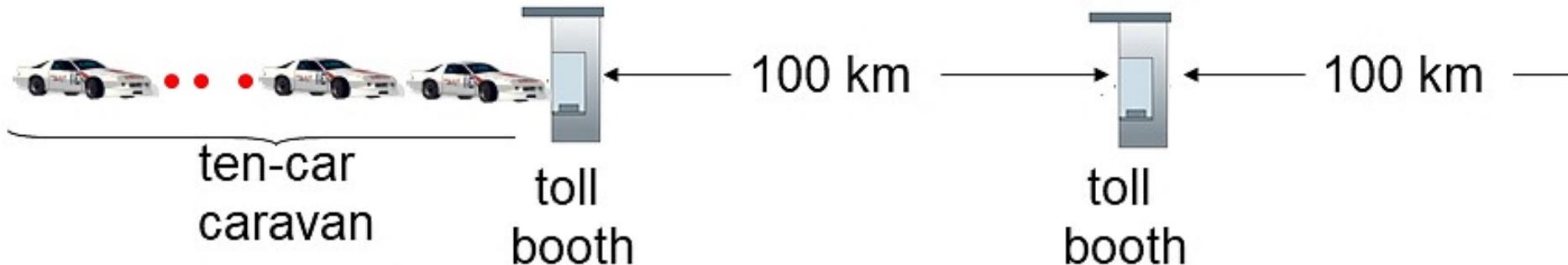
d_{trans} and d_{prop}
very different

d_{prop} : propagation delay:

- d : length of physical link
- s : propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- $d_{\text{prop}} = d/s$

* Check out the Java applet for an interactive animation on trans vs. prop delay

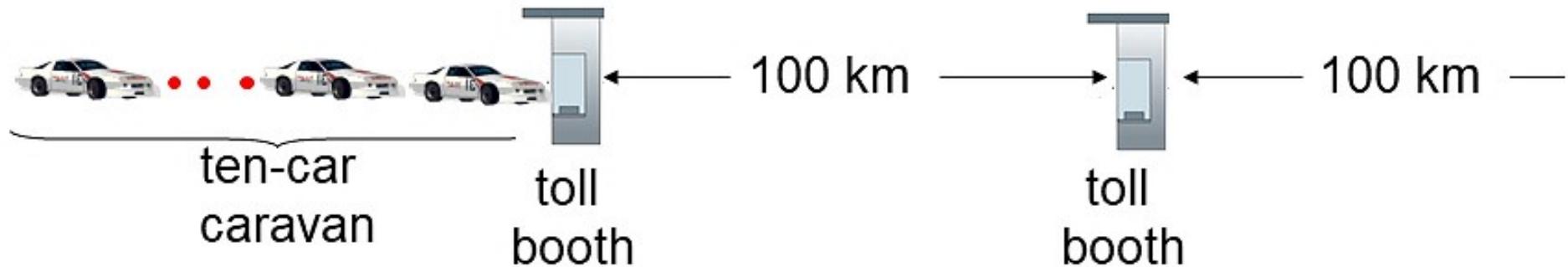
Caravan analogy



- ❖ cars “propagate” at 100 km/hr
- ❖ toll booth takes 12 sec to service car (bit transmission time)
- ❖ car~bit; caravan ~ packet
- ❖ **Q: How long until caravan is lined up before 2nd toll booth?**

- time to “push” entire caravan through toll booth onto highway = $12*10 = 120$ sec
- time for last car to propagate from 1st to 2nd toll both: $100\text{km}/(100\text{km/hr}) = 1$ hr
- **A: 62 minutes**

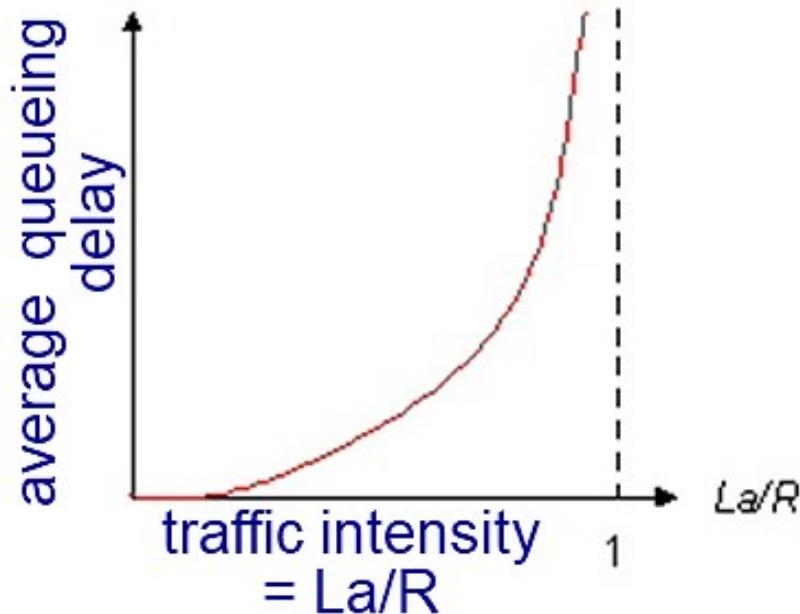
Caravan analogy (more)



- ❖ suppose cars now “propagate” at 1000 km/hr
- ❖ and suppose toll booth now takes one min to service a car
- ❖ **Q:** Will cars arrive to 2nd booth before all cars serviced at first booth?
 - **A: Yes!** after 7 min, 1st car arrives at second booth; three cars still at 1st booth.

Queueing delay (revisited)

- ❖ R : link bandwidth (bps)
- ❖ L : packet length (bits)
- ❖ a : average packet arrival rate



- ❖ $La/R \sim 0$: avg. queueing delay small
- ❖ $La/R \rightarrow 1$: avg. queueing delay large
- ❖ $La/R > 1$: more “work” arriving than can be serviced, average delay infinite!



$La/R \sim 0$

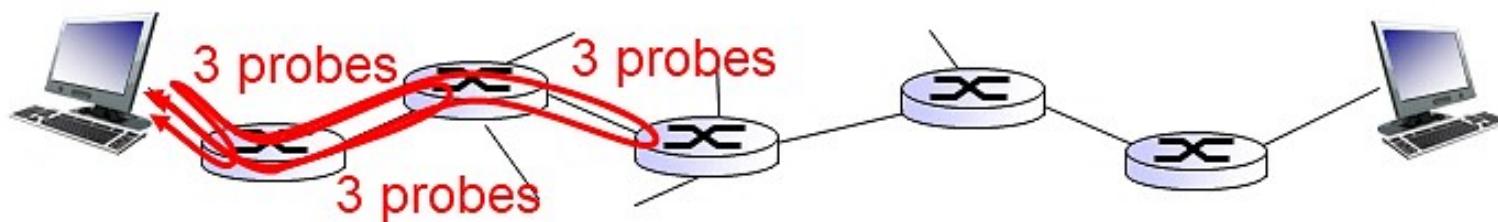


$La/R \rightarrow 1$

* Check out the Java applet for an interactive animation on queuing and loss

“Real” Internet delays and routes

- ❖ what do “real” Internet delay & loss look like?
- ❖ **traceroute** program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



“Real” Internet delays, routes

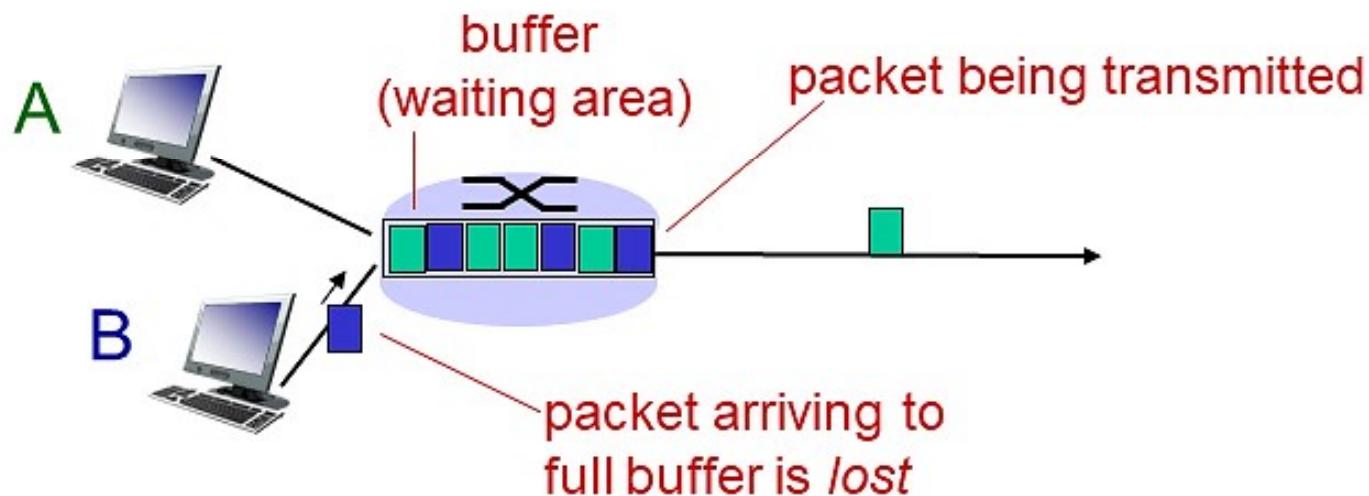
traceroute: gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from gaia.cs.umass.edu to cs-gw.cs.umass.edu						
1	cs-gw (128.119.240.254)	1 ms	1 ms	2 ms		
2	border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145)	1 ms	1 ms	2 ms		
3	cht-vbns.gw.umass.edu (128.119.3.130)	6 ms	5 ms	5 ms		
4	jn1-at1-0-0-19.wor.vbns.net (204.147.132.129)	16 ms	11 ms	13 ms		
5	jn1-so7-0-0-0.wae.vbns.net (204.147.136.136)	21 ms	18 ms	18 ms		
6	abilene-vbns.abilene.ucaid.edu (198.32.11.9)	22 ms	18 ms	22 ms		
7	nycm-wash.abilene.ucaid.edu (198.32.8.46)	22 ms	22 ms	22 ms		
8	62.40.103.253 (62.40.103.253)	104 ms	109 ms	106 ms		
9	de2-1.de1.de.geant.net (62.40.96.129)	109 ms	102 ms	104 ms		
10	de.fr1.fr.geant.net (62.40.96.50)	113 ms	121 ms	114 ms		
11	renater-gw.fr1.fr.geant.net (62.40.103.54)	112 ms	114 ms	112 ms		
12	nio-n2.cssi.renater.fr (193.51.206.13)	111 ms	114 ms	116 ms		
13	nice.cssi.renater.fr (195.220.98.102)	123 ms	125 ms	124 ms		
14	r3t2-nice.cssi.renater.fr (195.220.98.110)	126 ms	126 ms	124 ms		
15	eurecom-valbonne.r3t2.ft.net (193.48.50.54)	135 ms	128 ms	133 ms		
16	194.214.211.25 (194.214.211.25)	126 ms	128 ms	126 ms		
17	***					
18	***	*	means no response (probe lost, router not replying)			
19	fantasia.eurecom.fr (193.55.113.142)	132 ms	128 ms	136 ms		

* Do some traceroutes from exotic countries at www.traceroute.org

Packet loss

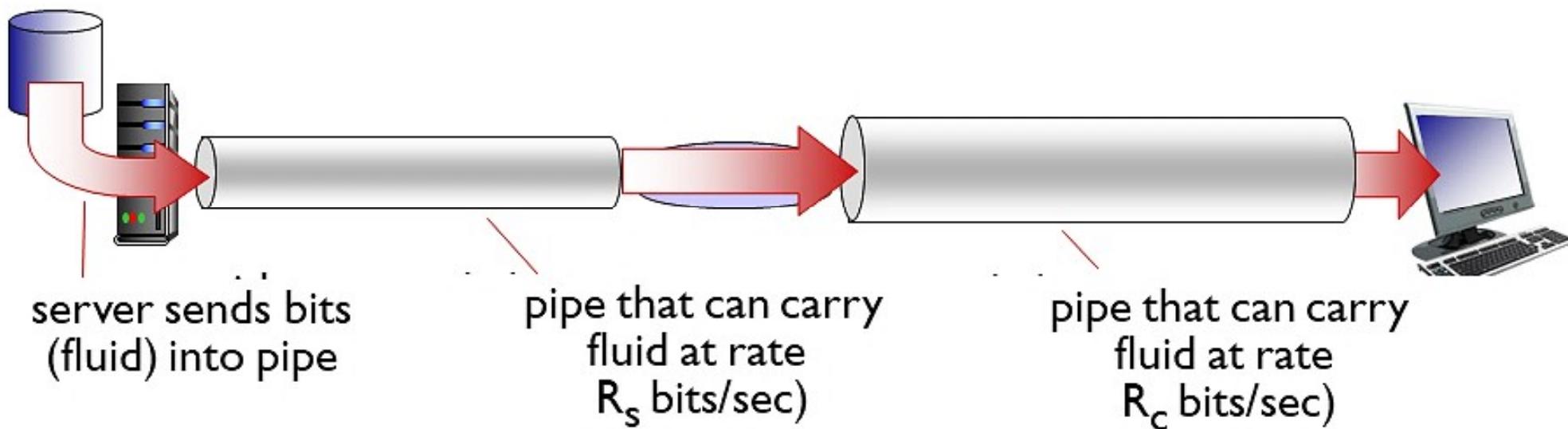
- ❖ queue (aka buffer) preceding link in buffer has finite capacity
- ❖ packet arriving to full queue dropped (aka lost)
- ❖ lost packet may be retransmitted by previous node, by source end system, or not at all



* Check out the Java applet for an interactive animation on queuing and loss

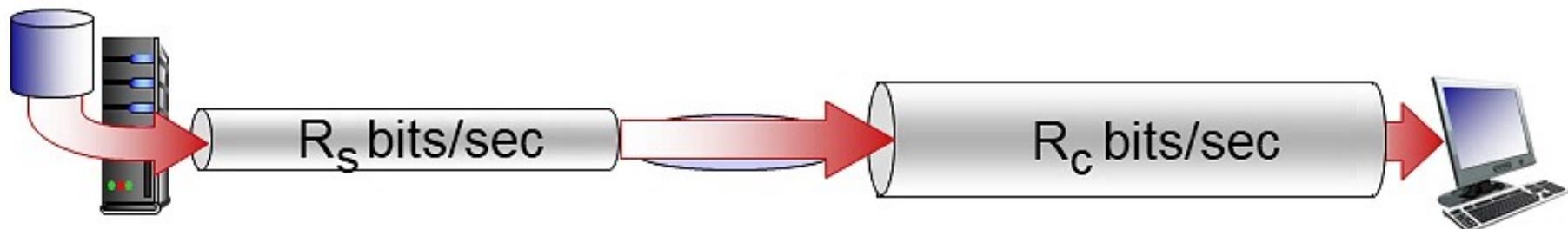
Throughput

- ❖ *throughput*: rate (bits/time unit) at which bits transferred between sender/receiver
 - *instantaneous*: rate at given point in time
 - *average*: rate over longer period of time

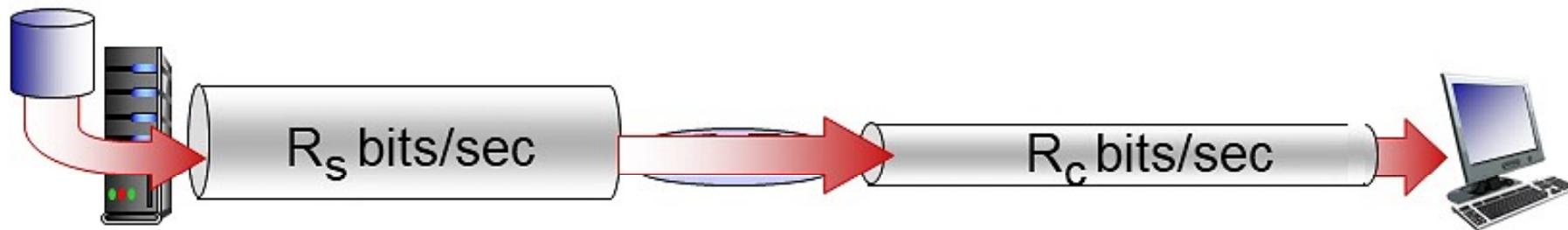


Throughput (more)

- ❖ $R_s < R_c$ What is average end-end throughput?



- ❖ $R_s > R_c$ What is average end-end throughput?

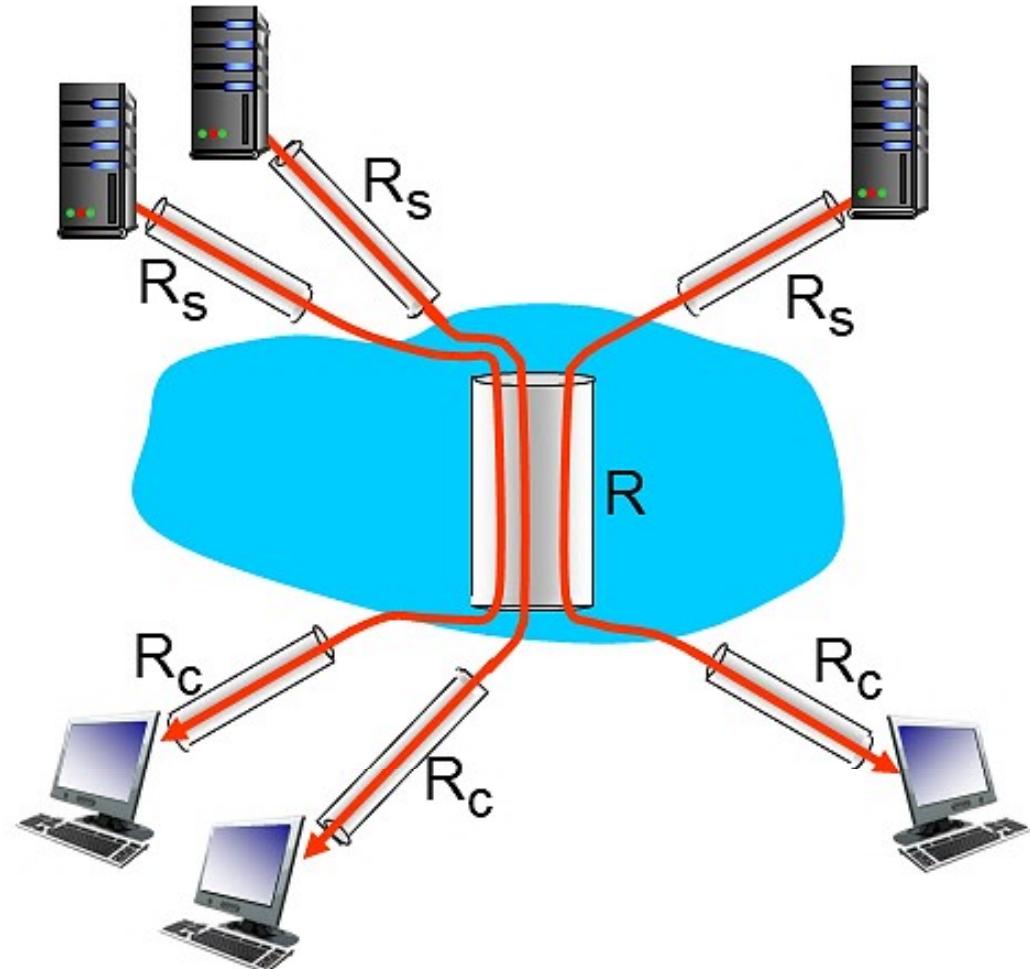


bottleneck

link on *end*-end path that constrains end-end throughput

Throughput: Internet scenario

- ❖ per-connection end-end throughput:
 $\min(R_c, R_s, R/10)$
- ❖ in practice: R_c or R_s is often bottleneck



10 connections (fairly) share
backbone bottleneck link R bits/sec

Chapter 1: roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History

Protocol “layers”

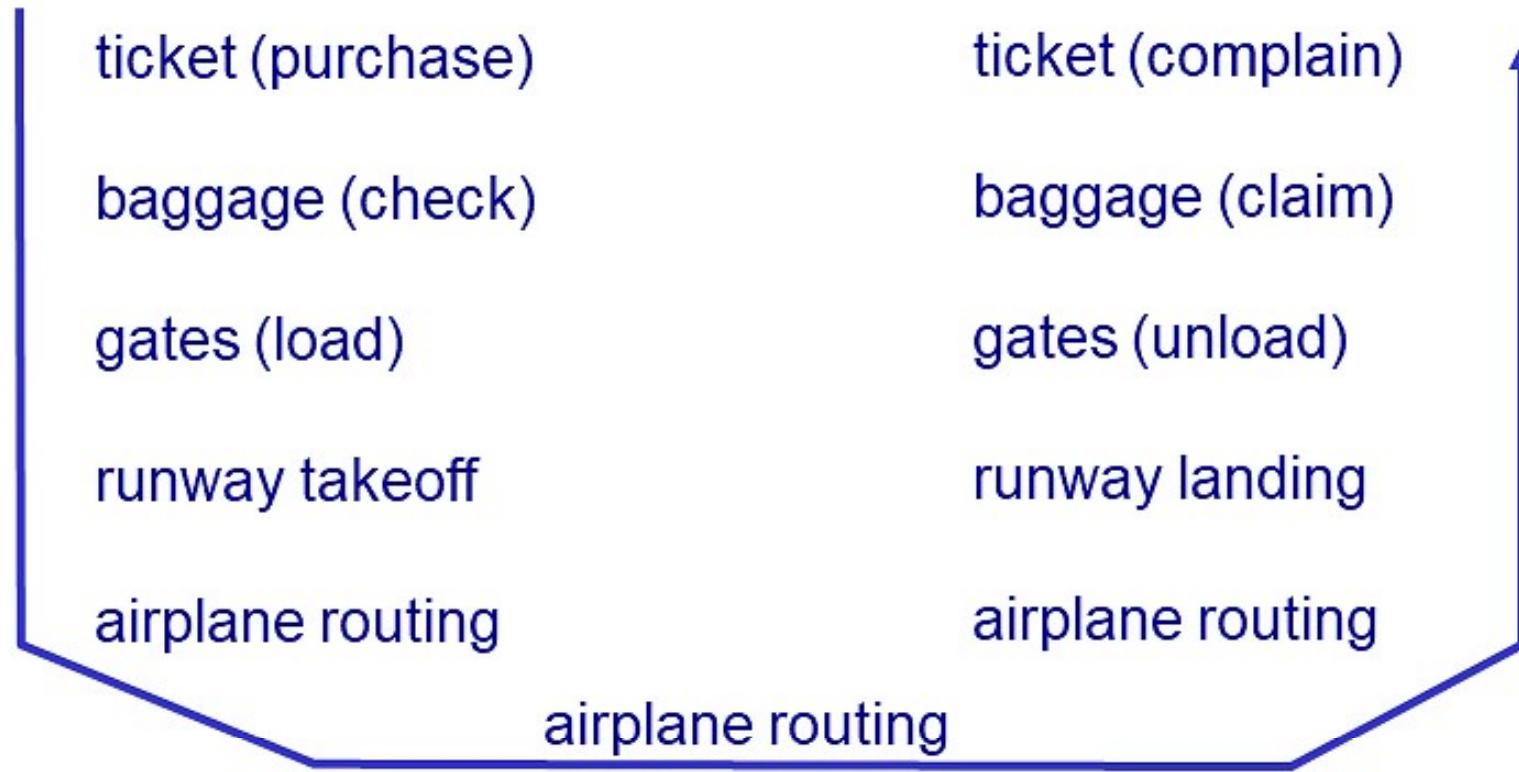
Networks are complex, with many “pieces”:

- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question:
is there any hope of
organizing structure of
network?

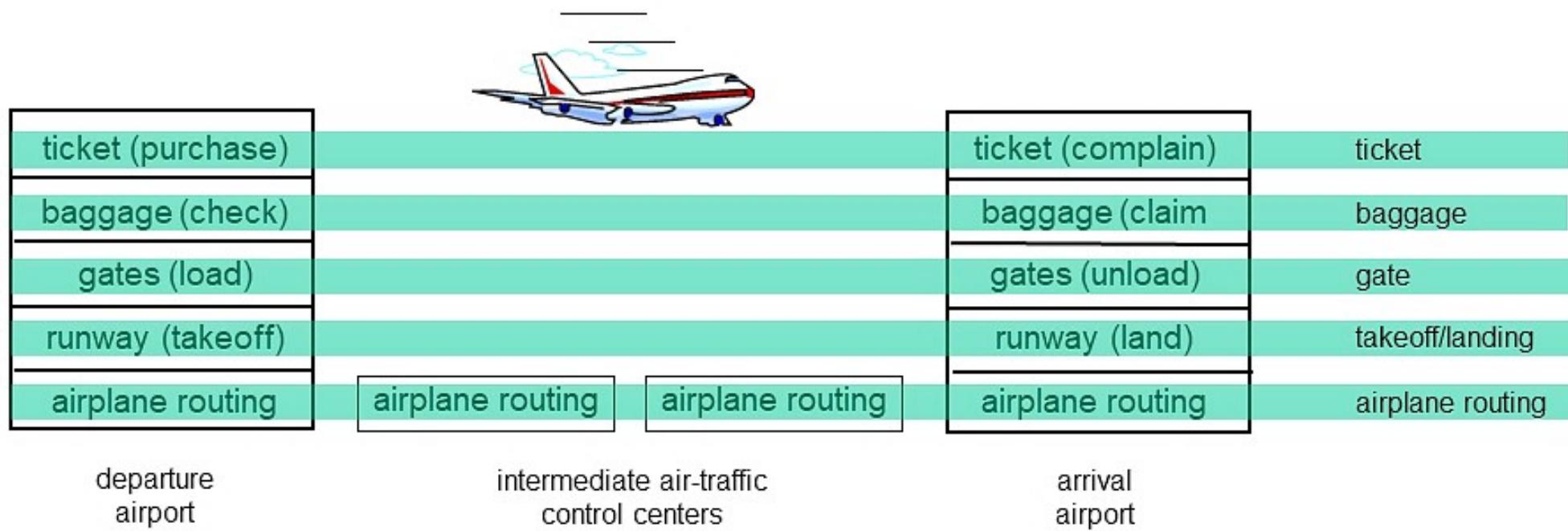
.... or at least our
discussion of networks?

Organization of air travel



- ❖ a series of steps

Layering of airline functionality



layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

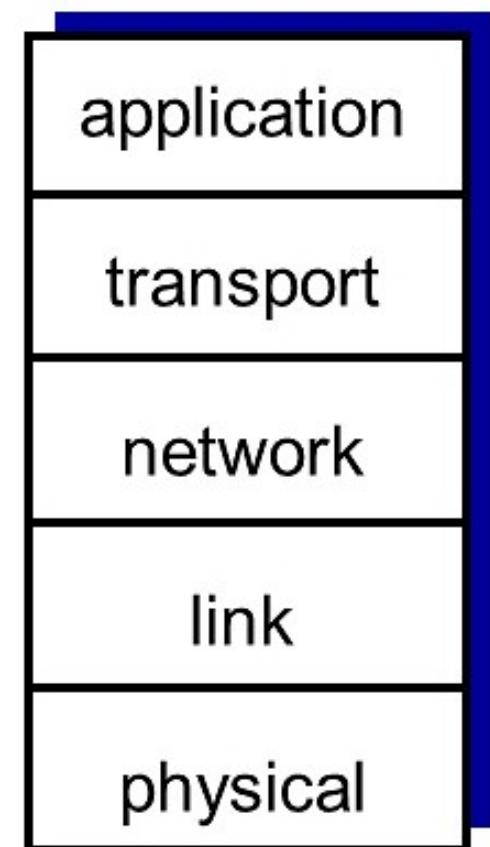
Why layering?

dealing with complex systems:

- ❖ explicit structure allows identification, relationship of complex system's pieces
 - layered *reference model* for discussion
- ❖ modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- ❖ layering considered harmful?

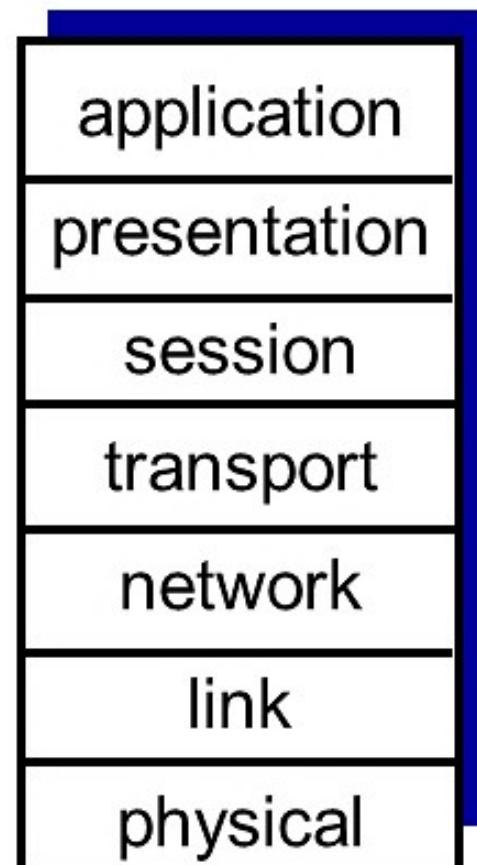
Internet protocol stack

- ❖ *application*: supporting network applications
 - FTP, SMTP, HTTP
- ❖ *transport*: process-process data transfer
 - TCP, UDP
- ❖ *network*: routing of datagrams from source to destination
 - IP, routing protocols
- ❖ *link*: data transfer between neighboring network elements
 - Ethernet, 802.111 (WiFi), PPP
- ❖ *physical*: bits “on the wire”



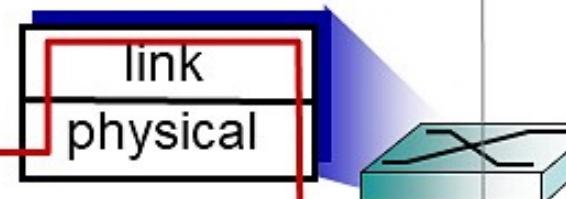
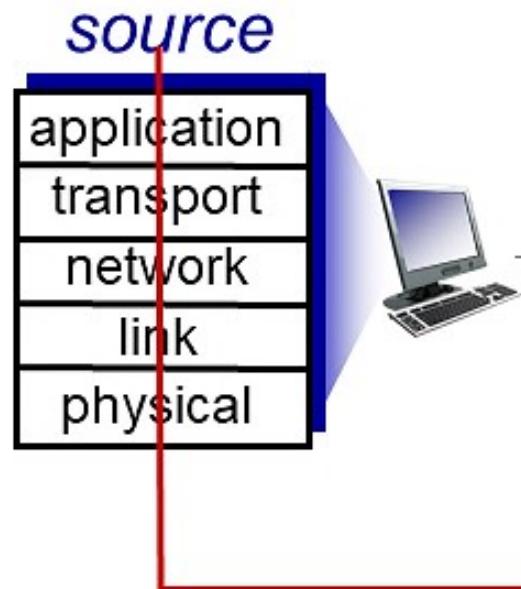
ISO/OSI reference model

- ❖ *presentation*: allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- ❖ *session*: synchronization, checkpointing, recovery of data exchange
- ❖ Internet stack “missing” these layers!
 - these services, *if needed*, must be implemented in application
 - needed?



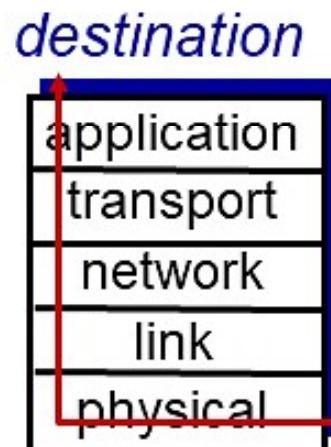
Encapsulation

message	M
segment	H _t M
datagram	H _n H _t M
frame	H _l H _n H _t M

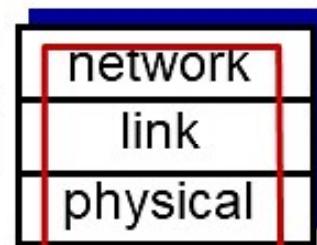


switch

M
H _t M
H _n H _t M
H _l H _n H _t M



H _n H _t M
H _l H _n H _t M



H _n H _t M



router

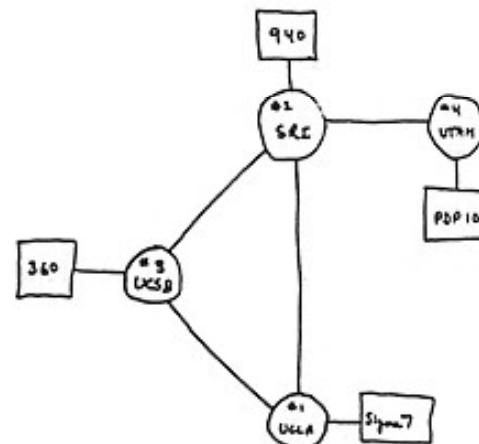
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Internet history

1961-1972: Early packet-switching principles

- ❖ 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- ❖ 1964: Baran - packet-switching in military nets
- ❖ 1967: ARPAnet conceived by Advanced Research Projects Agency
- ❖ 1969: first ARPAnet node operational
- ❖ 1972:
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



THE ARPA NETWORK

Internet history

1972-1980: Internetworking, new and proprietary nets

- ❖ 1970: ALOHAnet satellite network in Hawaii
- ❖ 1974: Cerf and Kahn - architecture for interconnecting networks
- ❖ 1976: Ethernet at Xerox PARC
- ❖ late 70' s: proprietary architectures: DECnet, SNA, XNA
- ❖ late 70' s: switching fixed length packets (ATM precursor)
- ❖ 1979: ARPAnet has 200 nodes

Cerf and Kahn's
internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet
architecture

Internet history

1980-1990: new protocols, a proliferation of networks

- ❖ 1983: deployment of TCP/IP
- ❖ 1982: smtp e-mail protocol defined
- ❖ 1983: DNS defined for name-to-IP-address translation
- ❖ 1985: ftp protocol defined
- ❖ 1988: TCP congestion control
- ❖ new national networks: Csnets, BITnet, NSFnet, Minitel
- ❖ 100,000 hosts connected to confederation of networks

Internet history

1990, 2000's: commercialization, the Web, new apps

- ❖ early 1990's: ARPAnet decommissioned
- ❖ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ❖ early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990's:
commercialization of the Web
- late 1990's – 2000's:
 - ❖ more killer apps: instant messaging, P2P file sharing
 - ❖ network security to forefront
 - ❖ est. 50 million host, 100 million+ users
 - ❖ backbone links running at Gbps

Internet history

2005-present

- ❖ ~750 million hosts
 - Smartphones and tablets
- ❖ Aggressive deployment of broadband access
- ❖ Increasing ubiquity of high-speed wireless access
- ❖ Emergence of online social networks:
 - Facebook: soon one billion users
- ❖ Service providers (Google, Microsoft) create their own networks
 - Bypass Internet, providing “instantaneous” access to search, email, etc.
- ❖ E-commerce, universities, enterprises running their services in “cloud” (eg, Amazon EC2)