



The Beauty and Joy of Computing



Lecture #19 Limits of Computing

Algorithms Determine "Character"

A new startup called "Upstart" is making a name by trying to use algorithms to determine "character" traits about who is mostly likely to pay a loan back. They try to give loans to people who are less likely to qualify using metrics like college graduation or SAT scores.



<http://mobile.nytimes.com/blogs/bits/2015/07/26/using-algorithms-to-determine-character/>



Admin Notes

- Schedule (see website)
- Next Week – Lots of guests!
- HKN Surveys
 - Please come, bonus points!



Introduction to Complexity Theory

Computer Science ... A UCB view

www.eecs.berkeley.edu/Research/Areas/

- CS research areas:
 - Artificial Intelligence
 - Biosystems & Computational Biology
 - Database Management Systems
 - Graphics
 - Human-Computer Interaction
 - Networking
 - Programming Systems
 - Scientific Computing
 - Security
 - Systems
 - Theory
 - Complexity theory
 - ...



UC Berkeley "The Beauty and Joy of Computing": **Limits of Computing** (4)



Ball



Let's revisit algorithm complexity

- Problems that...
 - are tractable with efficient solutions in reasonable time
 - are intractable
 - are solvable approximately, not optimally
 - have no known efficient solution
 - are not solvable





Tractable with efficient sols in reas time

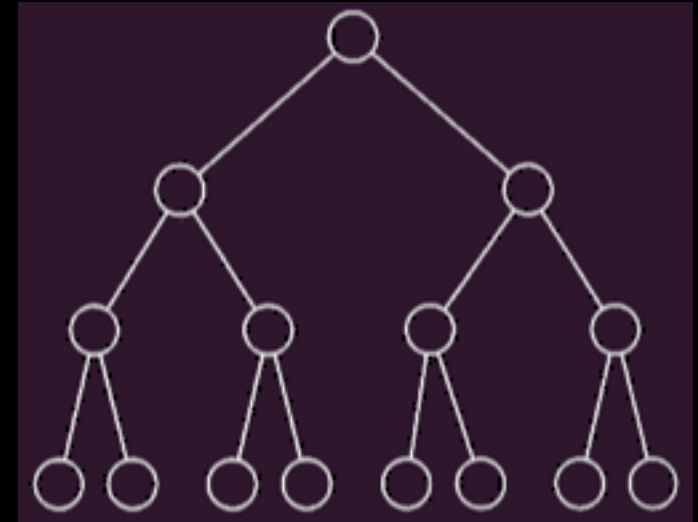
- Recall our algorithm complexity lecture we've got several common orders of growth
 - Constant
 - Logarithmic
 - Linear
 - Quadratic
 - Cubic
 - Exponential
- Order of growth is polynomial in the size of the problem
- E.g.,
 - Searching for an item in a collection
 - Sorting a collection
 - Finding if two numbers in a collection are same
- These problems are called being "in P" (for polynomial)





Intractable problems

- Problems that can be solved, but not solved fast enough
- This includes exponential problems
 - E.g., $f(n) = 2^n$
 - as in the image to the right
- This also includes poly-time algorithm with a huge exponent
 - E.g., $f(n) = n^{10}$
- Only solve for small n



Imagine a program that calculated something important at each of the bottom circles. This tree has height n , but there are 2^n bottom circles!





(Cal) Peer Instruction

What's the most you can put in your knapsack?

- a) \$10
- b) \$15
- c) \$33
- d) \$36
- e) \$40



Knapsack Problem

You have a backpack with a weight limit (here **15kg**), which boxes (with weights and values) should be taken to maximize value?

(any # of each box is available)

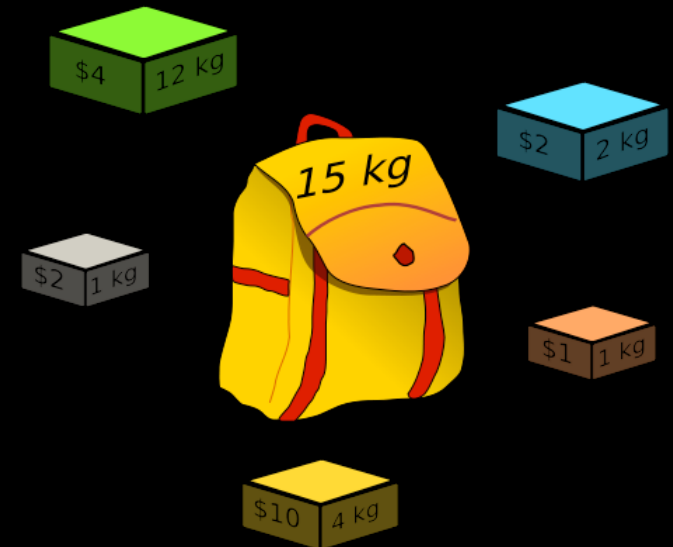


Heuristics, NP,
NP-Hard,
NP-Complete



Solvable approximately, not optimally in reasonable time

- A problem might have an optimal solution that cannot be solved in reasonable time
 - E.g., optimization problems such as “find the best/smallest”
- BUT if you don’t need to know the perfect solution, there might exist “approximation” algorithms which could give pretty good answers in reasonable time
- Heuristic: a technique that may allow us to find an **approximate solution** (e.g., valuable stuff first!)
- **Some problems cannot be solved using any algorithm.** (e.g., finding a robot path to a blocked-off area)



Knapsack Problem

You have a backpack with a weight limit (here 15kg), which boxes (with weights and values) should be taken to maximize value?





Have no known efficient solution

- Solving one of them would solve an entire class of them!
 - We can transform one to another, i.e., reduce
 - A problem P is “hard” for a class C if every element of C can be “reduced” to P
- If you’re “in NP” and “NP-hard”, then you’re “NP-complete”



-2 -3 15
14 7 -10

Subset Sum Problem

Are there a handful of these numbers (at least 1) that add together to get 0?

- If you guess an answer, can I verify it in polynomial time?
 - Called being “in NP”
 - Non-deterministic (the “guess” part)
Polynomial



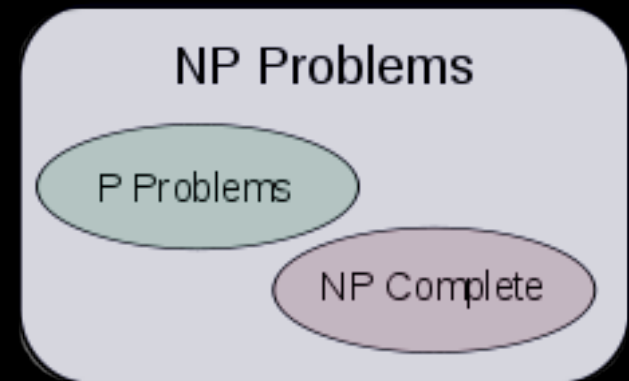


The fundamental question. **Is $P = NP$?**

- This is THE major unsolved problem in Computer Science!
 - One of 7 “millennium prizes” w/a \$1M reward
- All it would take is solving ONE problem in the NP-complete set in polynomial time!!
 - Huge ramifications for cryptography, others



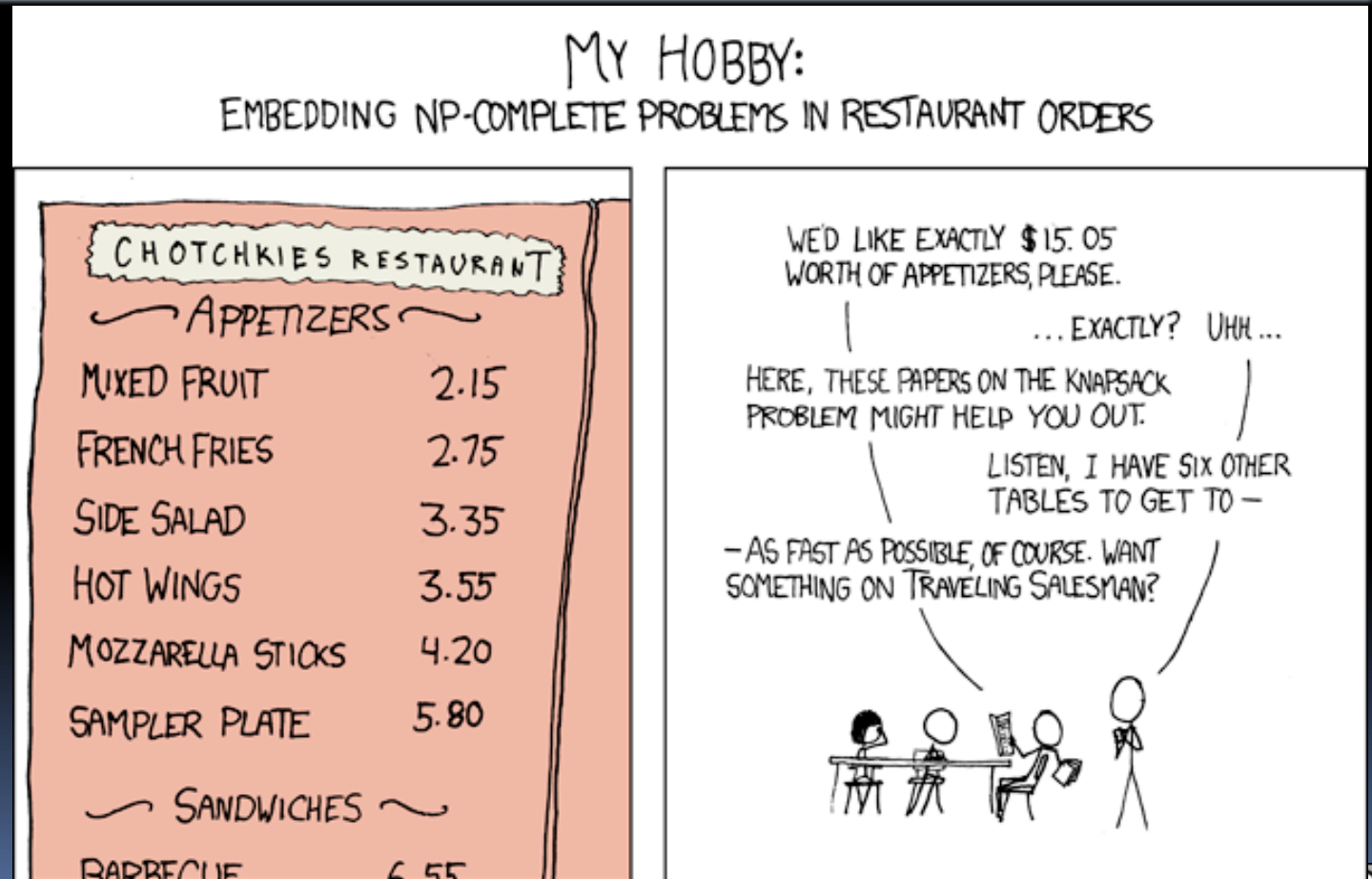
If $P \neq NP$, then



- Other NP-Complete
 - Traveling salesman who needs most efficient route to visit all cities and return home

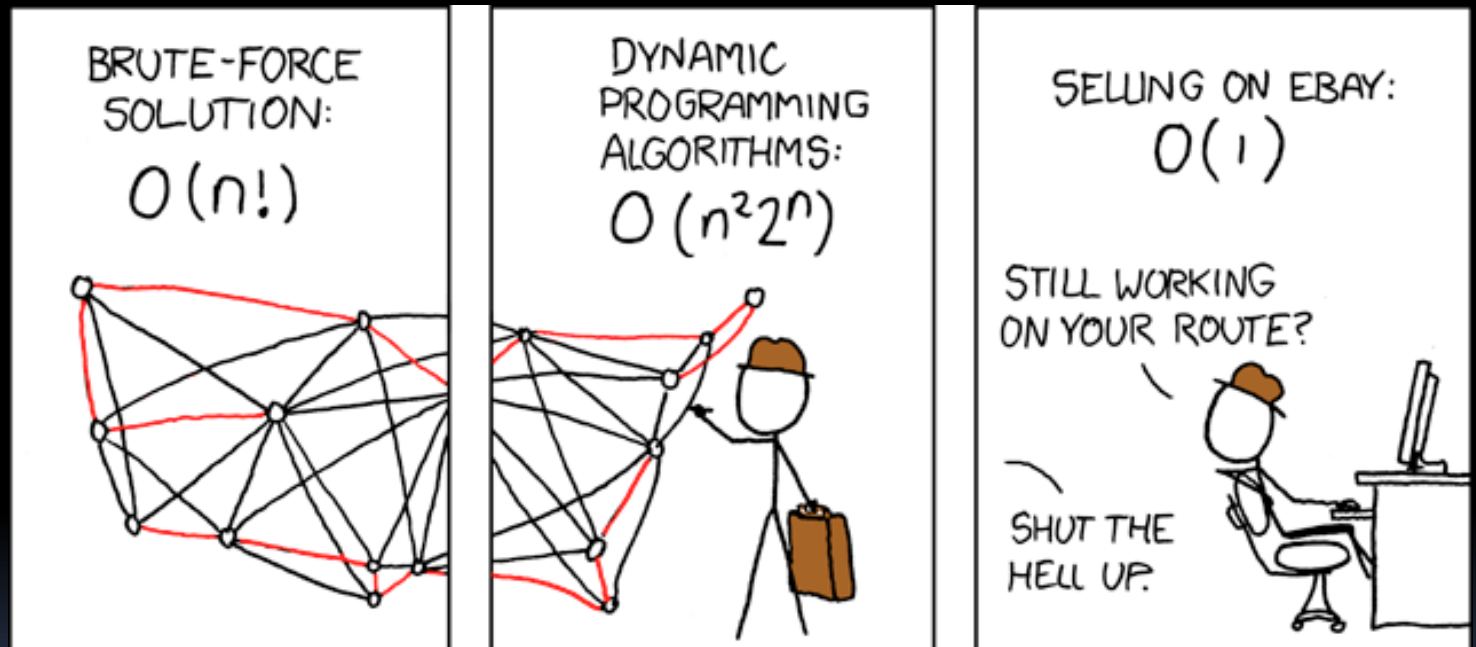


XKCD #287, NP-Complete





XKCD #399, Travelling Salesman

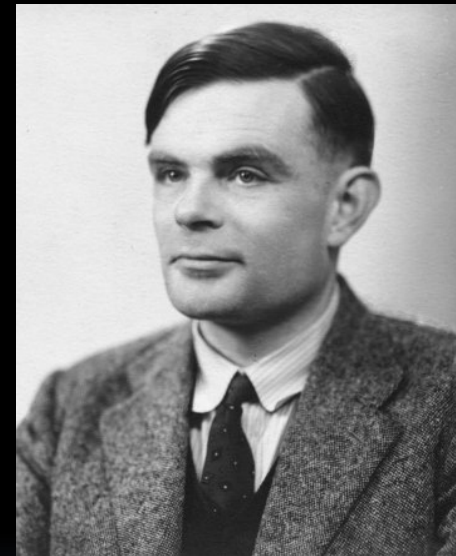


Decidability



Problems NOT solvable

- Decision problems answer YES or NO for an infinite # of inputs
 - E.g., is N prime?
 - E.g., is sentence S grammatically correct?
- An algorithm is a solution if it correctly answers YES/NO in a finite amount of time
- A problem is decidable if it has a solution



Alan Turing

He asked:

"Are all problems decidable?"
(people used to believe this was true)

Turing proved it wasn't for CS!



Review: Proof by Contradiction

- Infinitely Many Primes?
- Assume the contrary, then prove that it's impossible
 - Only a finite set of primes, numbered p_1, p_2, \dots, p_n
 - Consider $q = (p_1 \cdot p_2 \cdot \dots \cdot p_n) + 1$
 - Dividing q by p_i has remainder 1
 - q either prime or composite
 - If prime, q is not in the set
 - If composite, since no p_i divides q , there must be another p that does that is not in the set.
 - So there's infinitely many primes



Euclid

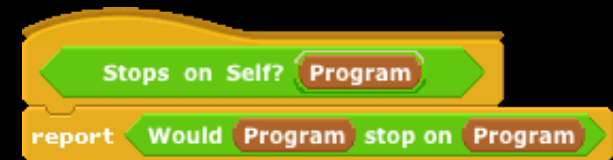
www.hisschemoller.com/wp-content/uploads/2011/01/euclides.jpg





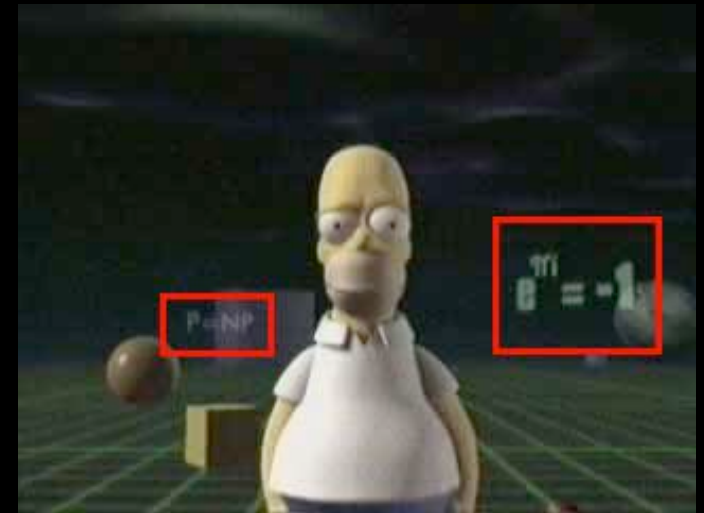
Turing's proof : The Halting Problem

- Given a program and some input, will that program eventually stop? (or will it loop)
- Assume we could write it, then let's prove a contradiction
 - 1. write Stops on Self?
 - 2. Write Weird
 - 3. Call Weird on itself



Conclusion

- Complexity theory
important part of CS
- If given a hard problem,
rather than try to solve it
yourself, see if others
have tried similar
problems
- If you don't need an
exact solution, many
approximation
algorithms help
- Some not solvable!



**P=NP question even made its way
into popular culture, here shown in
the Simpsons 3D episode!**

