

Admin

YEAH Repeat, today 2pm (😘 Ishita)

Interest in `printf` code jam this weekend? When is best?

Today: Modules, Linking

Modules and libraries

What makes for good design?

Build process

How do source files become an executable?

What does the linker do?

Understanding and diagnosing build errors

How does a program begin executing?



Road map

Processor and memory architecture

Peripherals: GPIO, timers, UART

Assembly language and machine code

From C to assembly language

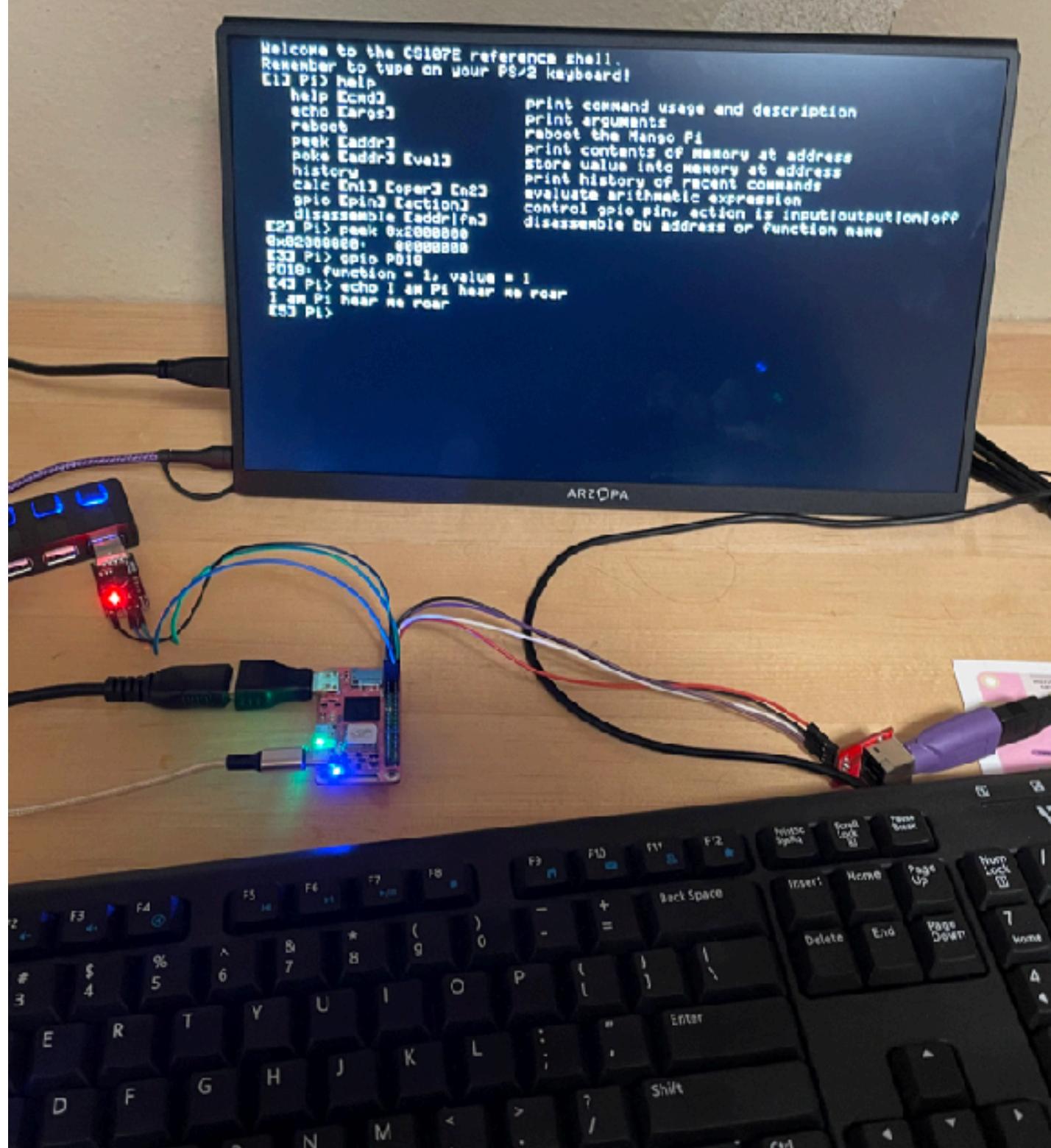
Function calls and stack frames

Serial communication and strings

→ **Modules and libraries: Building and linking**

Memory: Memory map, stack & heap

gpio ✓
timer ✓
uart ✓
strings
printf
malloc
keyboard
fb
gl
console
shell



Good software design

Decompose a system into smaller parts (modules)

- **Interface:** what the module does
- **Implementation:** how it does it

A good interface is:

- Easy-to-understand abstraction that simplifies client use
- Can be implemented easily and in different ways, portable

Module is tested independently with unit tests

Designing good interface boundaries is the "art" of software engineering

Module example: printf

Interface is single feature— make formatted string

```
int printf(const char* format, ...);
```

Implementation is complex, must:

- Parse format string
- Convert arguments into strings
- Concatenate conversions into one string
- Output string over uart

snprintf()

num_to_string()

strlcat()

uart_putstr()

Decompose problem into smaller, simpler tasks,
test independently, build up in layers

"The Build"



NASA Command Center during SpaceX mission



CS107e Command Center: **Makefile**

```
APPLICATION = clock
MY_MODULES = gpio.o timer.o

ARCH      = -march=rv64im -mabi=lp64
ASFLAGS   = $(ARCH)
CFLAGS    = $(ARCH) -Og -g -Wall -ffreestanding
LDFLAGS   = -nostdlib -T memmap.ld
LDLIBS    =
```

```
all : $(APPLICATION).bin
```

```
%.bin: %.elf
    riscv64-unknown-elf-objcopy $< -O binary $@
```

```
%.elf: %.o $(MY_MODULES) cstart.o start.o           LINKER
    riscv64-unknown-elf-ld $(LDFLAGS) $^ $(LDLIBS) -o $@
```

```
%.o: %.c                                         COMPILER
    riscv64-unknown-elf-gcc $(CFLAGS) -c $< -o $@
```

```
%.o: %.s                                         ASSEMBLER
    riscv64-unknown-elf-as $(ASFLAGS) $< -o $@
```

Build process

PREPROCESSOR: `cpp` (`code.c` -> `code.i`)

Input: C source file

Output: replace `#include` and `#define` with text expansions

COMPILER: `cc1` (`code.i` -> `code.s`)

Input: preprocessed C source file

Output: assembly file

ASSEMBLER: `as` (`code.s` -> `code.o`)

Input: assembly file

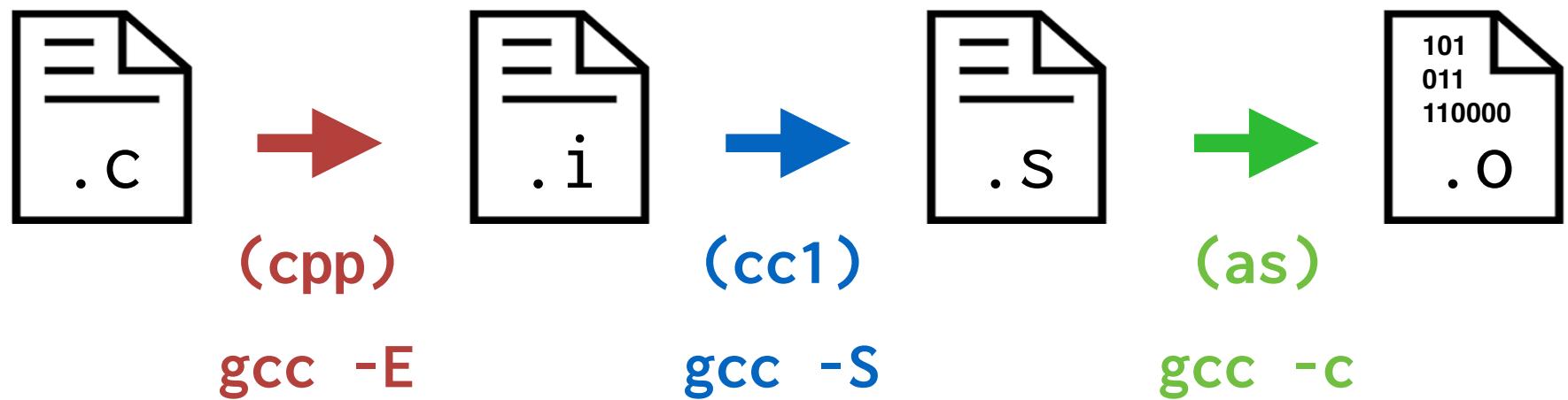
Output: machine code (binary-encoded)

LINKER: `ld` (`code.o` -> `code.elf`)

Input: multiple .o files, libraries

Output: executable

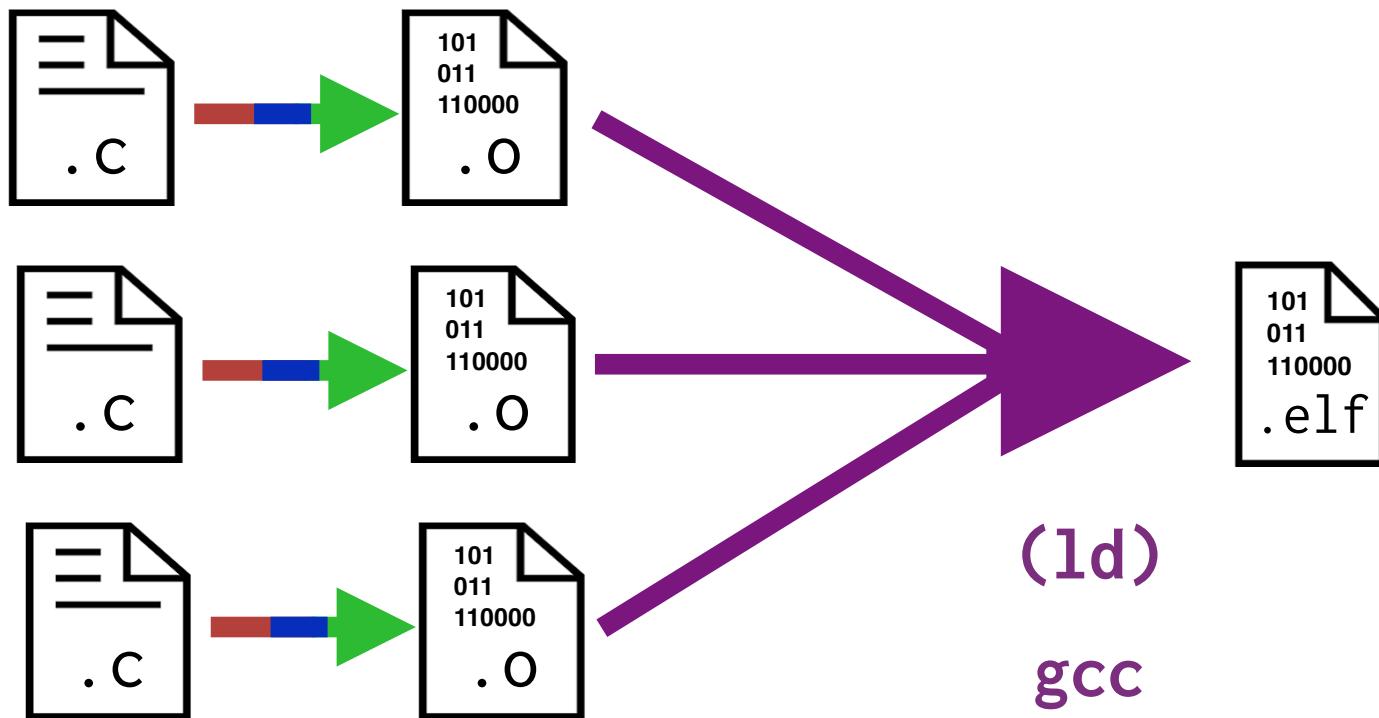
Compiling a single module



want to see intermediate files?

gcc --save-temp

Link modules into executable



Combines `clock.o` `timer.o` `gpio.o` `start.o` `cstart.o`

What does the linker do?

Combines one or more object files into executable

Resolve inter-module references

Consolidate code, data sections across all modules

Arrange sections in output file in proper order

Symbols

C has single global namespace

Each name can have only one definition

Need conventions to avoid name clashes

e.g. `gpio_init` versus `timer_init`

Qualifier dictates symbol visibility and link behavior

static private to module, not visible outside, no linking needed

extern public, visible to other modules, cross-module references need link

Linker only concerned with `extern` symbols

Symbol resolution

Rule: Every symbol referenced must be defined once and exactly once

Linker errors during symbol resolution:

Multiply-defined: two definitions for same symbol

Undefined: needed symbol never defined

Linker resolution process



Set D (symbols that have been defined)

Set U (symbols that have been referenced, but not yet defined)

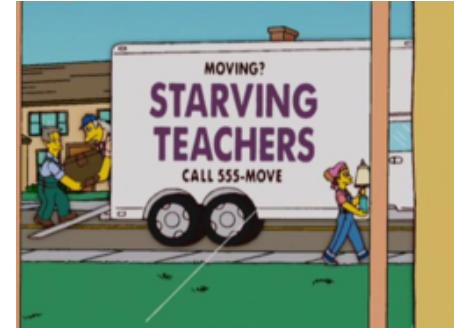
Linker processes .o files in order from left to right

Process each .o updates **Set U** and **Set D**

Continue until all .o processed

If end with **Set U** empty and **Set D** has no duplicates, link successful!

Symbol relocation



In .o file:

each symbol assigned a temporary address (offset within module)
reference to symbols within same module is PC-relative

In executable file:

modules consolidated into one file, laid out end to end
first module starts at address **0x40000000**
base for each subsequent module is offset by size of previous modules
final address of each symbol =
 $\text{module_base} + \text{symbol_offset_within_module}$

start.s

```
.globl _start
_start:
    lui      sp,0x50000
    jal      _cstart
hang: j hang
```

start.o disassembly

```
0: lui      sp,0x50000
4: jal      0
8: j       8 <hang>
```

Target address for `j hang` is 8, encoded as PC-relative offset (in this case, `pc+0`)

Target address for `jal _cstart` is 0 (placeholder!).
File `start.s` doesn't know where `_cstart` is! (Why not?)

```
// Disassembly of clock.elf
```

```
0000000040000000 <_start>:  
 4000000: lui sp,0x5000  
 4000004: jal 400000e4 <_cstart>
```

```
0000000040000008 <hang>:  
 400000c: j 4000008 <hang>
```

All addresses have been relocated/finalized
Address of `_cstart` is now known — **400000e4**
The linker worked it out and replaced target
placeholder with function's actual address

000000040000000	T	_start	
000000040000008	t	hang	start.o
00000004000000c	t	blink	
000000040000050	T	main	clock.o
000000040000150	T	_cstart	cstart.o
0000000400001a8	T	gpio_init	
0000000400001ac	T	gpio_set_input	
0000000400001b0	T	gpio_set_output	
0000000400001b4	T	gpio_set_function	
0000000400001b8	T	gpio_get_function	
0000000400001c0	T	gpio_write	
0000000400001c4	T	gpio_read	gpio.o
0000000400001cc	T	timer_get_ticks	
0000000400001d4	T	timer_init	
0000000400001d8	T	timer_delay_us	
000000040000210	T	timer_delay	
000000040000234	T	timer_delay_ms	timer.o
000000040000288	R	__bss_end	
000000040000288	R	__bss_start	

Handling of global data

Variable declared outside functions have global scope and program lifetime
(can be extern/static, read-only, initialized/not)

```
// uninitialized
int gNum;
static int sgNum;
```

```
// initialized
int iNum = 1;
static int siNum = 2;
```

```
// const
const int cNum = 3;
static const int scNum = 4;
```

```
$ riscv64-unknown-elf-nm -U -n data.o
0000000000000004 R cNum
0000000000000004 B gNum
0000000000000004 D iNum
000000000000002c T main
0000000000000000 r scNum
0000000000000000 b sgNum
0000000000000000 t show_var
0000000000000000 d siNum
```

LEGEND

T/t = text (code)

D/d = read-write data

R/r = read-only data (const)

B/b = bss (*Block Started by Symbol*)

C = common (instead of B)

lowercase letter means **static**, uppercase **extern**

Libraries

An archive .a is just a collection of modules (.o files)

Linker looks in library to find definitions for symbols in **Set U** (undefined). When symbol is found, entire module is linked in.

If linking this module results in new undefined symbols, they are added to **Set U**. Process repeats until no additional undefined symbols, move on to next library

Building with libmango

```
ARCH      = -march=rv64im -mabi=lp64
CFLAGS    = $(ARCH) -Og -I$$CS107E/include -Wall -ffreestanding
LDFLAGS   = -nostdlib -L$$CS107E/lib -T memmap.ld
LDLIBS    = -lmango

%.bin: %.elf
    riscv64-unknown-elf-objcopy $< -O binary $@

%.elf: %.o
    riscv64-unknown-elf-gcc $(LDFLAGS) $^ $(LDLIBS) -o $@

%.o: %.c
    riscv64-unknown-elf-gcc $(CFLAGS) -c $< -o $@
```

Diagnosing build errors

Which tool detects/reports... ?

mismatched type

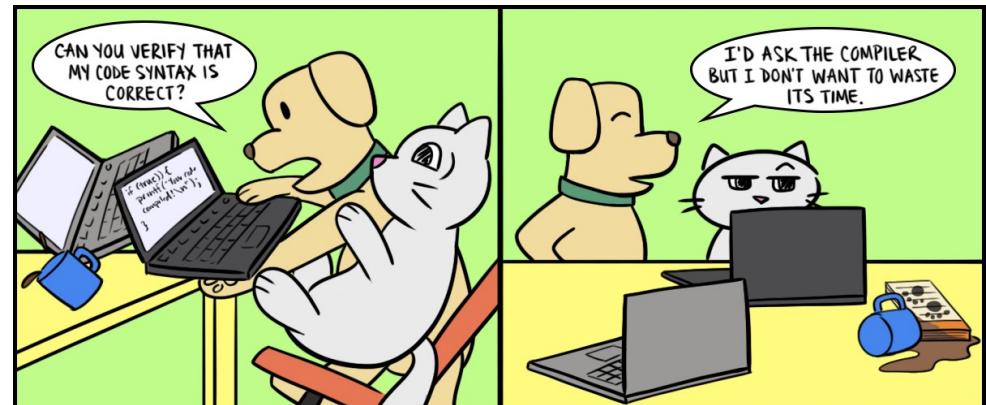
function call with wrong arguments

forgot #include

forgot to link with library

use variable uninitialized

two functions with same name



Modules and build process

Decompose into modules is skill to cultivate

- Separate system into coherent modules, minimize dependencies
- Clean abstractions, tidy interface
- Develop and test independently

Automate the build process

- Makefile is your command center
 - **Fast:** re-compile only what has changed
 - **Reliable:** track dependencies between modules
 - If file F changes, recompile everything that depends on F

Program start sequence

How is a program loaded into memory?

What must happen to start executing a program?

What is known about state of registers or contents of memory?

Does execution actually start at `main()`

•• **Read our files!** ••

[\\$CS107E/src/start.s](#)

[\\$CS107E/src/cstart.c](#)

[\\$CS107E/lib/memmap.ld](#)

memmap.ld (linker script)

```
SECTIONS
{
    .text 0x40000000 : { *(.text.start) *(.text*) }
    .rodata           : { *(.rodata*) }
    .data             : { *(.sdata*) }
    __bss_start       = .;
    .bss              : { *(.bss*) }
    __bss_end         = .;
}
```

What is special about **.text.start** and why must it go first?

What is the significance of **0x40000000**?

SECTIONS

```
{  
    .text 0x40000000 :{ *(.text.start)  
                         *(.text*)}  
    { *(.rodata*) }  
    { *(.data*) }  
    = .;  
    { *(.bss*) }  
    = . ;  
}
```

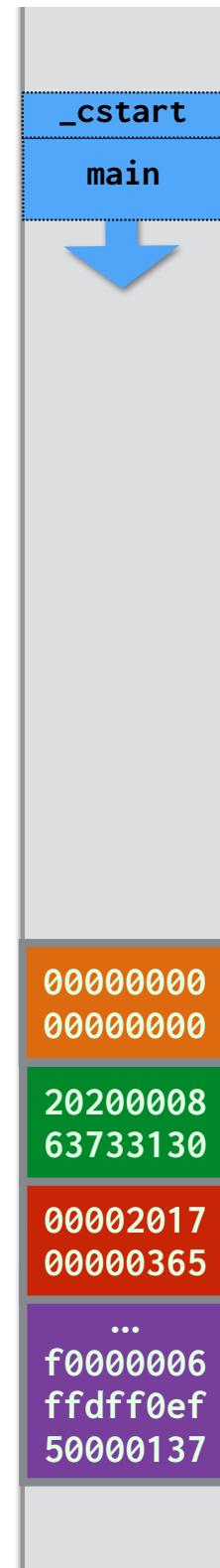
(zeroed data) .bss

(initialized data) .data

(read-only data) .rodata

.text

```
$ xfel write 0x40000000 clock.bin
```



0x50000000

```
_start:  
    lui    sp,0x50000  
    jal    _cstart
```

```
void _cstart(void) {  
    char *bss = &__bss_start__;  
    while (bss < &__bss_end__)  
        *bss++ = 0;  
}  
main();  
}
```

__bss_end

__bss_start

0x40000000

} clock.bin