

## CSL 356, Tutorial Sheet 5

1. Show how to implement Lagrange's interpolation formula in  $O(n^2)$  operations.
2. Describe an efficient algorithm to evaluate a degree  $n$  univariate polynomial  $P(x)$  at  $n$  arbitrary points  $x_1, x_2 \dots x_n$  (not necessarily roots of unity). You may assume that polynomial division takes the same asymptotic time as polynomial multiplication.  
Hint : Use the following observation which is similar to the remainder theorem. Let  $D_{i,j} = \prod (x - x_i)(x - x_{i+1}) \dots (x - x_j)$  and let  $P(x) = Q_{1,1}(x)D_{1,n} + R_{1,1}(x)$ . Then  $P(x_i) = R_{1,1}(x_i)$  where the degree of  $R_{1,1}$  is less than  $D_{1,n}$ . To compute  $R_{1,1}(x_i)$ , we can apply a similar decomposition, once with  $D_{1,n/2}$  and  $D_{n/2+1,n}$  recursively. This defines a tree where at each node we do a polynomial division (of degree  $n/2^i$  at distance  $i$  from the root). At the leaf nodes, we have the answers.
3. Let  $y$  be a string of length  $m$ . A substring  $x$  of  $y$  is called a period of  $y$  if  $y = (x^k)x'$ , where  $(x^k)$  is the string  $x$  repeated  $k$  times and  $x'$  is a prefix of  $x$ . The period is the shortest period of  $y$ . Design an efficient algorithm to determine the period of a string of length  $n$ .  
Hint: Prove that a string  $X$  is a period of a string  $Y$  iff  $Y$  is a prefix of  $XY$
4. If  $p$  and  $q$  are periods (not the shortest) and  $|p| + |q| < m$  then there is a period of length  $|p| - |q|$  (assuming  $p$  is larger).  
(This is the equivalent of Euclid's algorithm for strings).
5. Consider an intuitive extension to KMP algorithm for handling wild-card (single character) in pattern. To compute the failure function, we can assign  $f(i) = f(i-1) + 1$  where  $X_i = *$  since  $*$  can match any character. Similarly if  $f(j-1) = i-1$ , then  $f(j) = i$  as  $X_j$  can also match  $X_i = *$ .  
Now argue if this is consistent.  
Hint: Consider the pattern  $a \cdot b \cdot a \cdot * \cdot a \cdot a \cdot b \cdot a$ .
6. **Graph theory**
  - (i) Show that in any graph there are at least two vertices of the same degree.
  - (ii) Given a degree sequence  $d_1, d_2 \dots d_n$  such that  $\sum_i d_i = 2n - 2$ , construct a tree whose vertices have the above degrees.
  - (iii) Show that in a complete graph of six vertices where edges are colored red or blue, there is either a red or a blue triangle.
7. Given a directed acyclic graph, that has maximal path length  $k$ , design an efficient algorithm that partitions the vertices into  $k+1$  sets such that there is no path between any pair of vertices in a set.
8. Given an undirected graph, describe an algorithm to determine if it contains an even-length cycle. Can you do the same for odd-length cycle?
9. Given an undirected connected graph  $G$ , define the Biconnected Component Graph  $H$  as follows. For each BCC of  $G$  and articulation point of  $G$ , there is a vertex in  $H$ . There is an edge between vertices  $x$  and  $y$  in  $H$  if  $x$  is articulation point in the BCC  $y$ .
  - (a) Prove that  $H$  is a tree.
  - (b) Using  $H$  (or otherwise), design an efficient algorithm that adds the minimal number of edge to  $G$  to make it biconnected.
10. A directed graph is Eulerian if the in degree equals out degree for every vertex. Show that an Eulerian graph admits a tour where every edge is visited exactly once. Design an efficient (linear time) algorithm to find such a tour.

11. An  $n$ -vertex undirected graph is a scorpion if it has a vertex of degree 1 (the string) connected to a vertex of degree 2 (the tail) connected to a vertex of degree  $n - 2$  (the body) which is connected to the remaining  $n - 3$  vertices (the feet). Some of the feet may be connected among themselves. Give an  $O(n)$  algorithm to check if a given  $n \times n$  adjacency matrix represents a scorpion.
12. \* Instead of a DFS tree, starting from an arbitrary spanning tree, redesign the bi-connectivity algorithm. Your algorithm should run in linear time.
13. Given an undirected graph, orient the edges so that the resulting graph is strongly connected. When is it possible ? Design a linear time algorithm for this problem.
14. Find a maximum subgraph of  $G = (V, E)$  that has degrees of each vertex is at least  $k$ .
15. Describe an efficient algorithm to find the *girth* of a given undirected graph. The *girth* is defined as the length of the smallest cycle.