



University of
Pittsburgh

Algorithms and Data Structures 2

CS 1501



Fall 2022

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(Slides are adapted from Dr. Ramirez's and Dr. Farnan's CS1501 slides.)

Announcements

- Upcoming Deadlines
 - Homework 3: this Friday @ 11:59 pm
 - Lab 2: next Monday @ 11:59 pm
 - Assignment 1: Monday Oct 10th @ 11:59 pm
- Lecture recordings are available on Canvas under Panopto Video
- Please use the “Request Regrade” feature on GradeScope if you have any issues with your grades
- TAs student support hours available on the syllabus page

Previous lecture

- Binary Search Tree
 - How to delete
- Runtime of BST operations
 - add, search, delete
- Red-Black BST (Balanced BST)
 - definition and basic operations

Muddiest Points

- **Q: Was just curious of the real-life applications of red-black tree**
- **A: Example applications:**
 - TreeMap and TreeSet in Java
 - CPU Scheduling inside the Linux Kernel
- **Q: what does the triangle mean at the bottom of tree drawings?**
- **A: It means a subtree**

Muddiest Points

- **Q: when deleting interior nodes how does the parent node know where its new child is**
- **A:**
 - The removal method (`removeFromRoot`) returns the new root
 - The return value of the removal method is assigned to the right or left child of the parent
- **Q: how does the new node know where it's child is?**
- **A:** The new root has the same two children that the old root has; we are just replacing the data of the root node.

Muddiest Points

- Q: Can you explain how all the different functions for the removal of a node connect together into remove(T)
- A:
 - remove(T) creates an empty wrapper object and calls removeEntry, passing to it the entry to delete and the empty wrapper
 - removeEntry searches for *entry* by either moving left or right (depending on the comparisons between entry and the data portion of the current node)
 - Once the node that contains entry is found, the data inside the node is put inside the wrapper object
 - Then, removeFromRoot is called on the found node
 - removeFromRoot removes the node (the three cases) and may need to call findLargest and removeLargest on the left subtree of the node (if it has two children)
 - removeFromRoot returns the new root of the tree to removeEntry
 - removeEntry then sets the right or left child of current node to the return value of removeFromRoot and returns the node after this modification
 - Eventually removeEntry returns back to remove the (possibly) new root of the tree after removing entry
 - Finally, remove sets the root of the tree to the return value of removeEntry and returns the entry inside the wrapper object

Muddiest Points

- **Q: Why do you need a wrapper class for removeEntry?**
- A: Because we want removeEntry to return the root after removing entry.
- But we want it also to return a reference to the deleted data object, which is achieved using the wrapper object.

Muddiest Points

- **Q: Would the theoretical worst case BST operation be when the tree is only left nodes/right nodes?**
- **A:** Yes. For example, when searching for the deepest node in such a degenerate tree, the depth of the node will be $n-1$ and the run time is $\Theta(n)$.

Muddiest Points

- Q: A red node can be anywhere in the right subtree just so long as it's leaning to the left side of that subtree, correct?
- A: Correct

Muddiest Points

- **Q: what is the difference between a red node and black node?**
- **A:** A red node has a red edge to its parent, whereas a black node has a black edge to its parent. The root is always a black node.
- **Q: Where is this going? What utility do we get out of labelling the tree this way?**
- **A:** By labeling the tree nodes this way with the other constraints, the height of the tree is at most $2 \log n$

Muddiest Points

- **Q: What purpose does key serve in Red-Black BST?**
- A: The data objects inside the nodes are compared based on the key field.

Muddiest Points

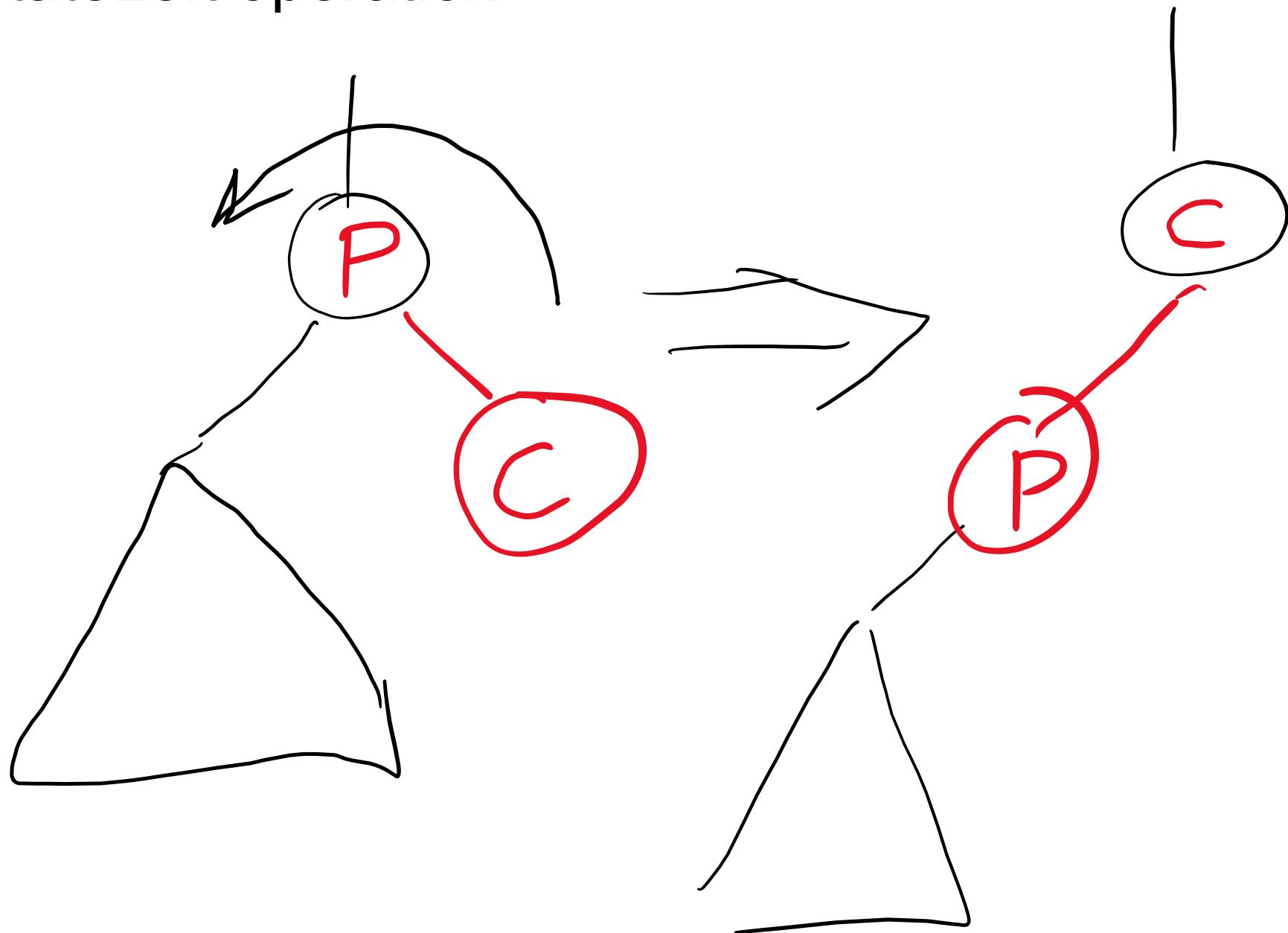
- **Q: How do you tell if a node is red/black based on the color of the link?**
- A: A node has the same color as the link to its parent.

Adding to a RB-BST

- Ok, so we add a red leaf node!
- What can go wrong then?
 - The new node is a right child
 - Fix it by left rotation

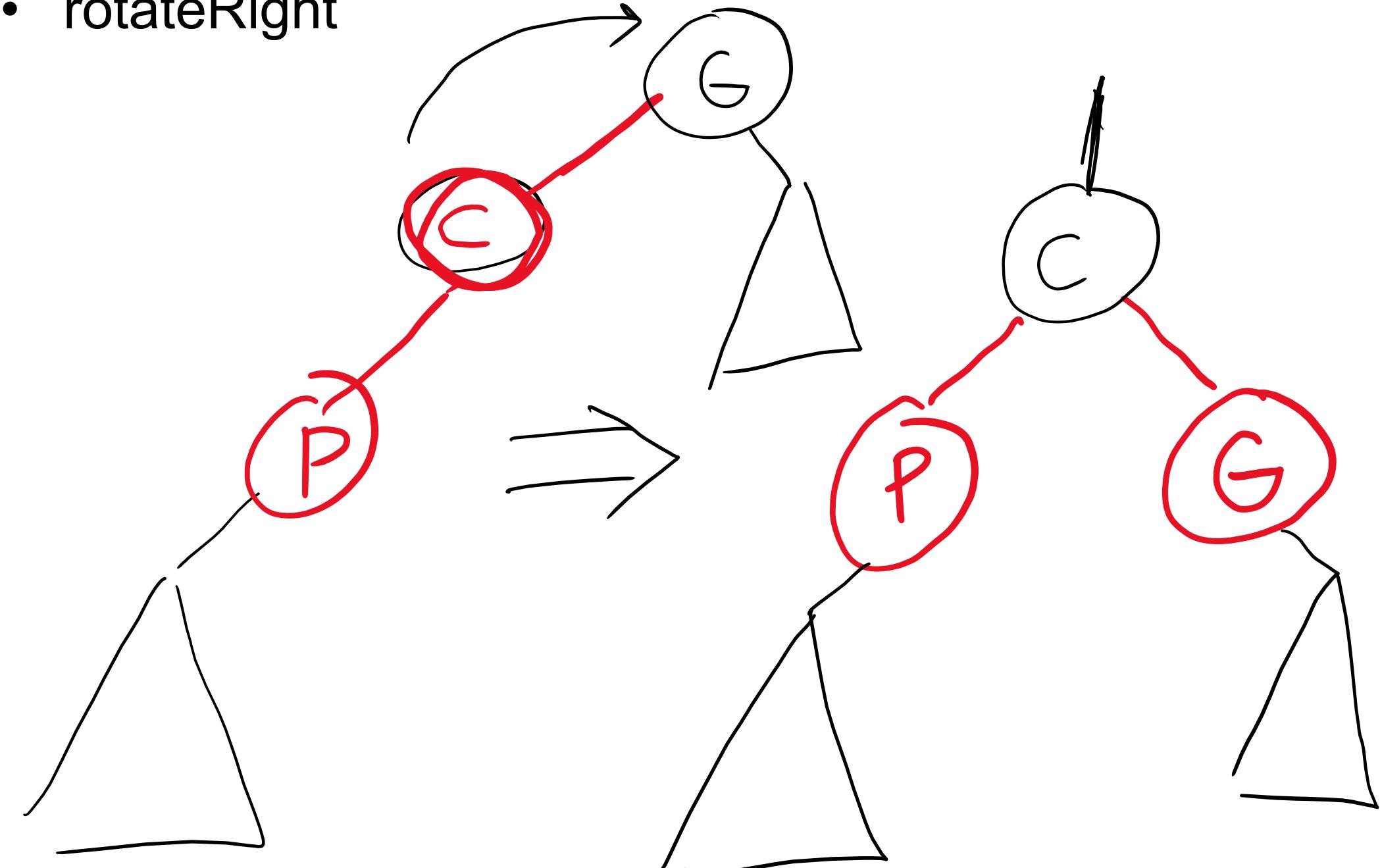
What if the new red node is a right child?

- `rotateLeft` operation



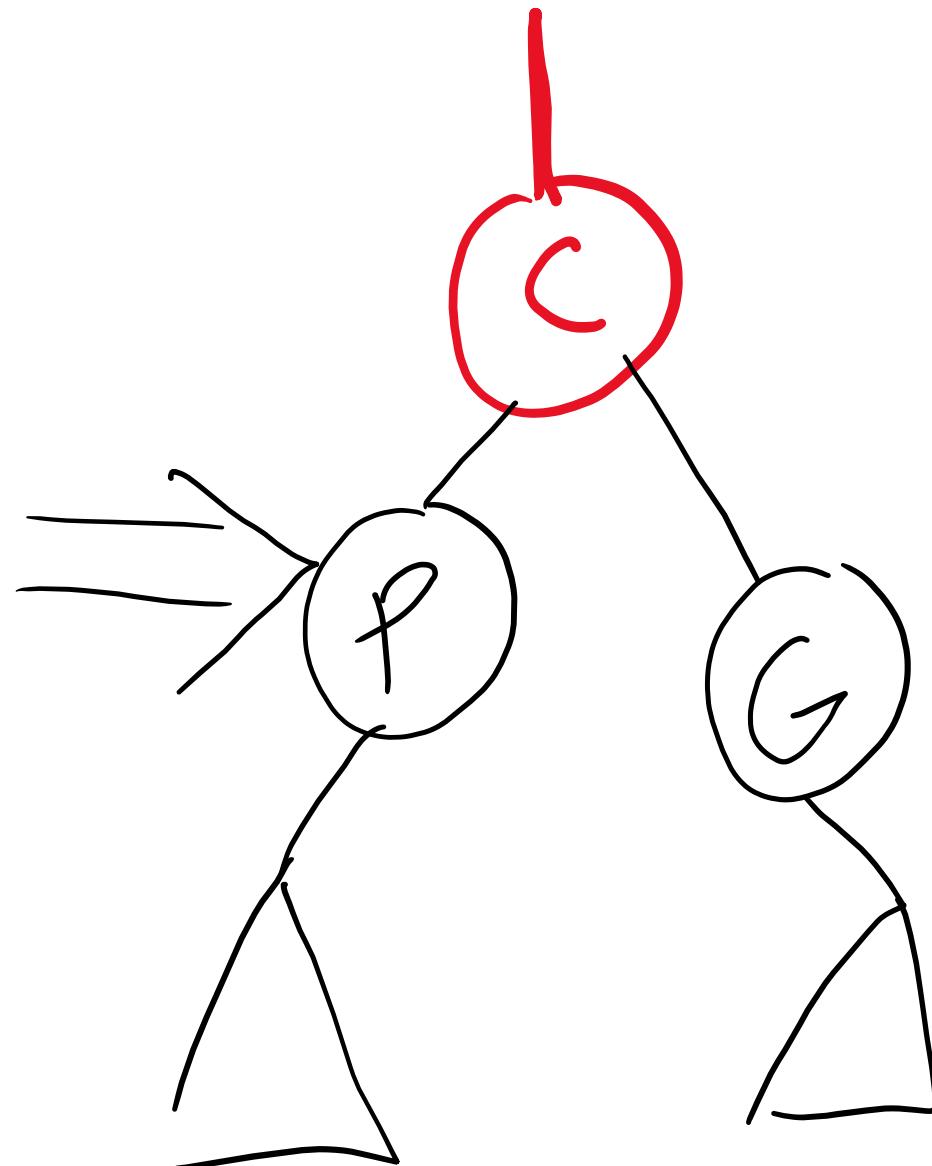
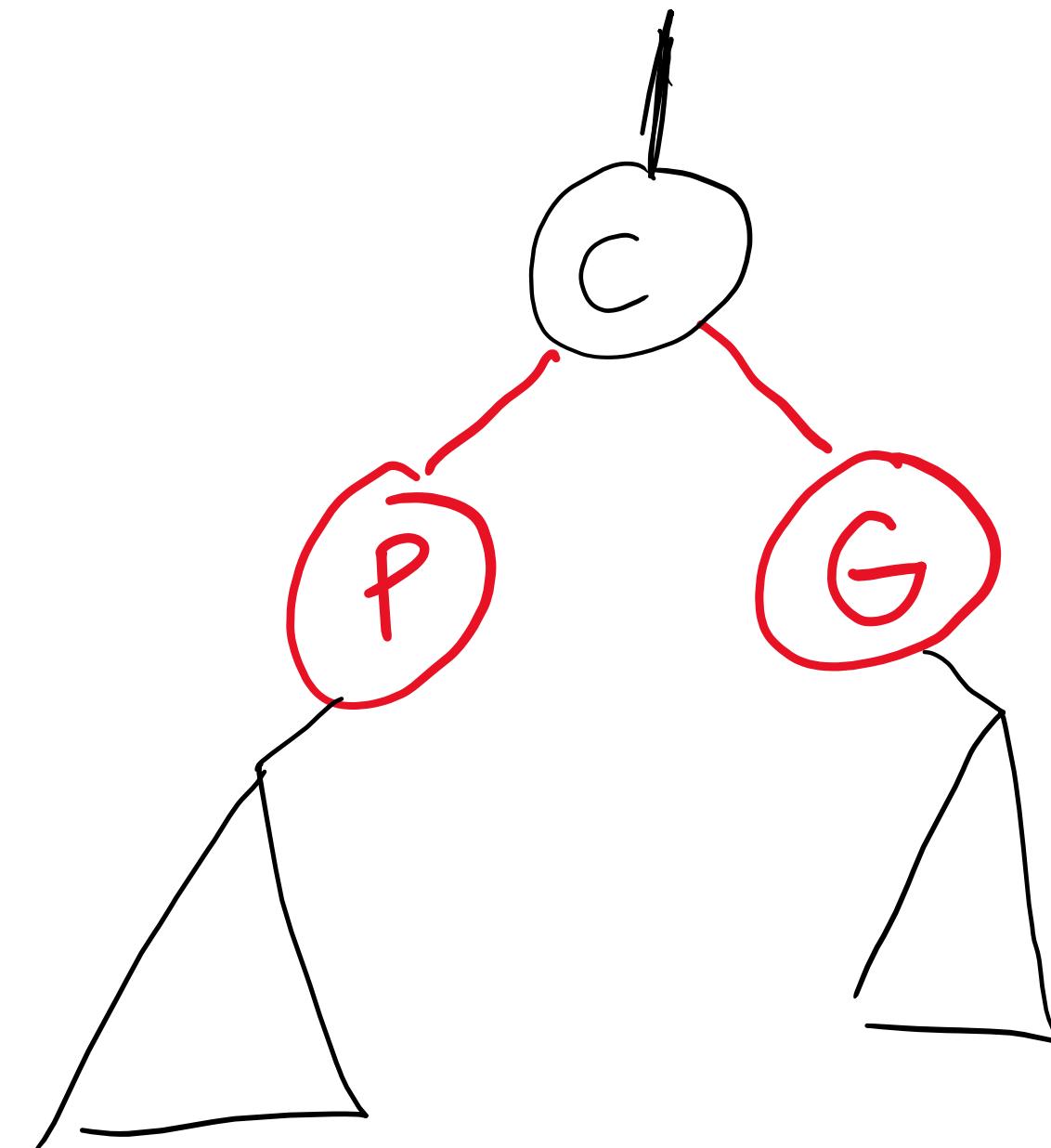
What if the new node becomes red?

- `rotateRight`



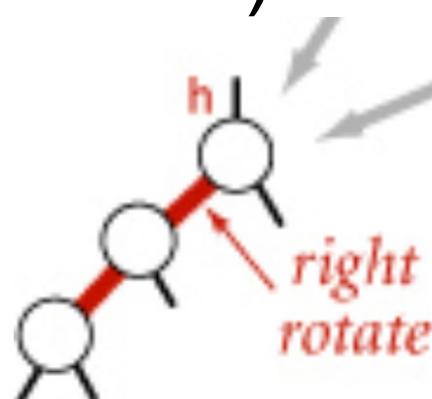
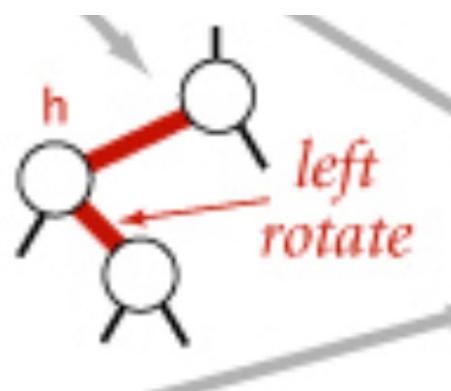
What if both children of a node are red?

- flipColors()



Muddiest Points

- Q: how to fix violations in B-R BST in terms of recursive call
- A: Violations are checked and corrected as we climb back up the tree.
- The code for violation corrections is after the recursive calls.
- Violations are detected and corrected in the following order (h is the current node)



Muddiest Points

- **Q: when do we add red/black nodes**
- **Q: How do we construct an RB BST if each node we add has to be red?**
- A: We always add red nodes then potentially change color of new node
- **Q: I found the transient right red edges confusing**
- **Q: I didn't quite understand how the form other than a red node being added or a rotation/color flip.**
- **Q: How does a Red black BST with red edges eventually turn into an all black BST**
- A: Let's see an example

Adding to a RB-BST: Example

- Let's add the numbers:
- 1, 2, 3, 4, 5, 6, 7

1

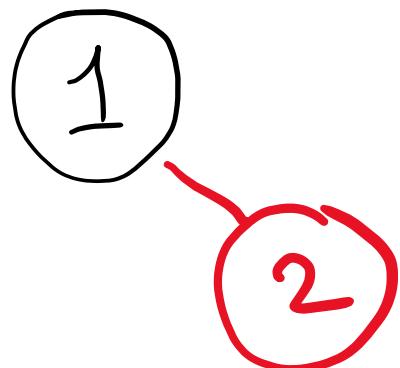
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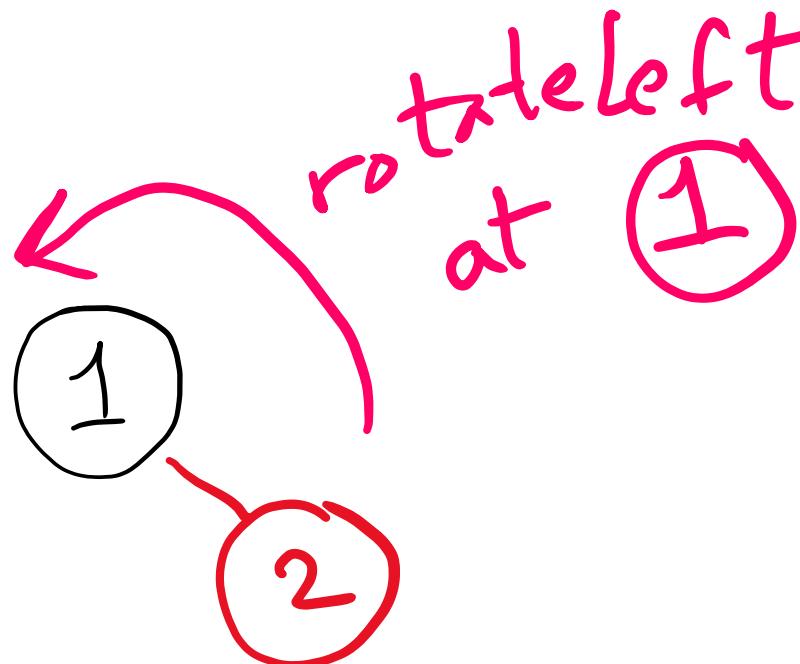
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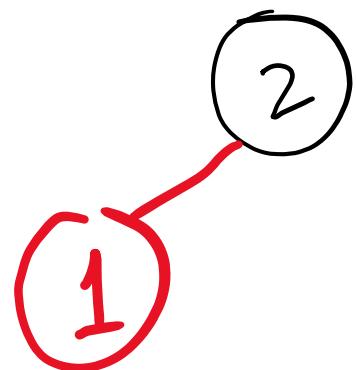
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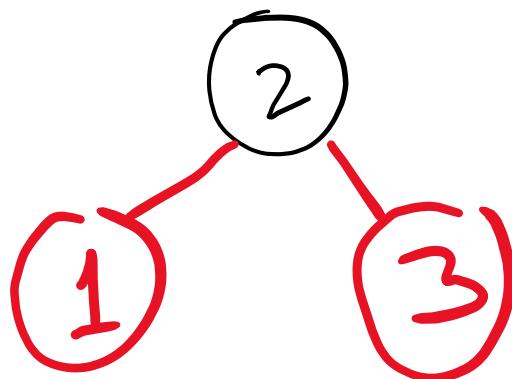
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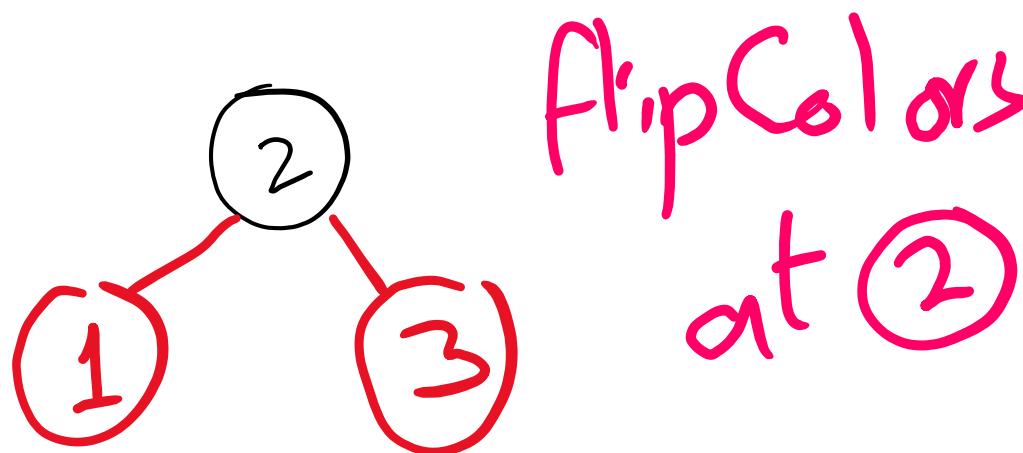
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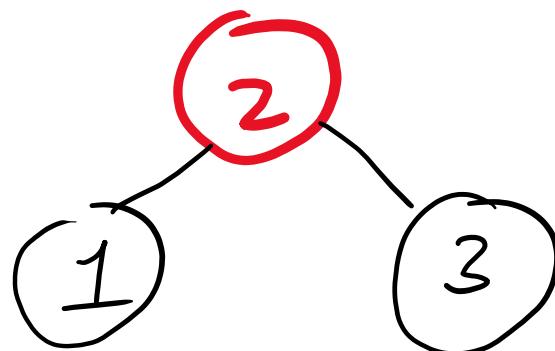
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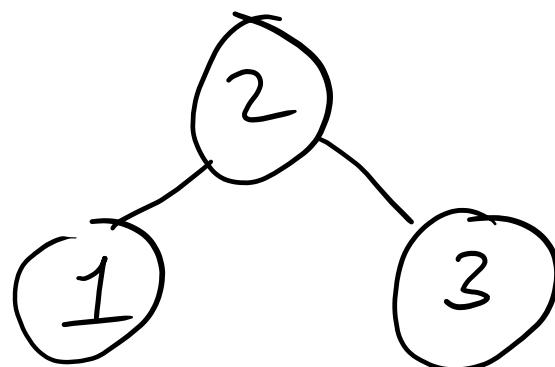
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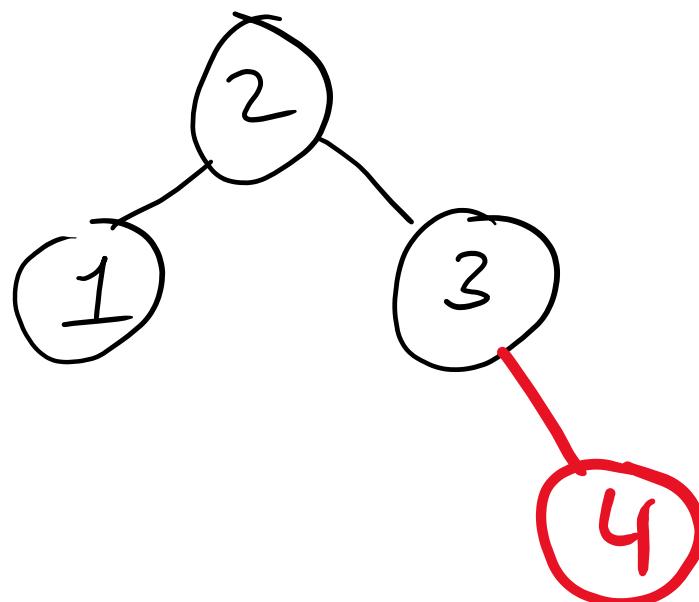
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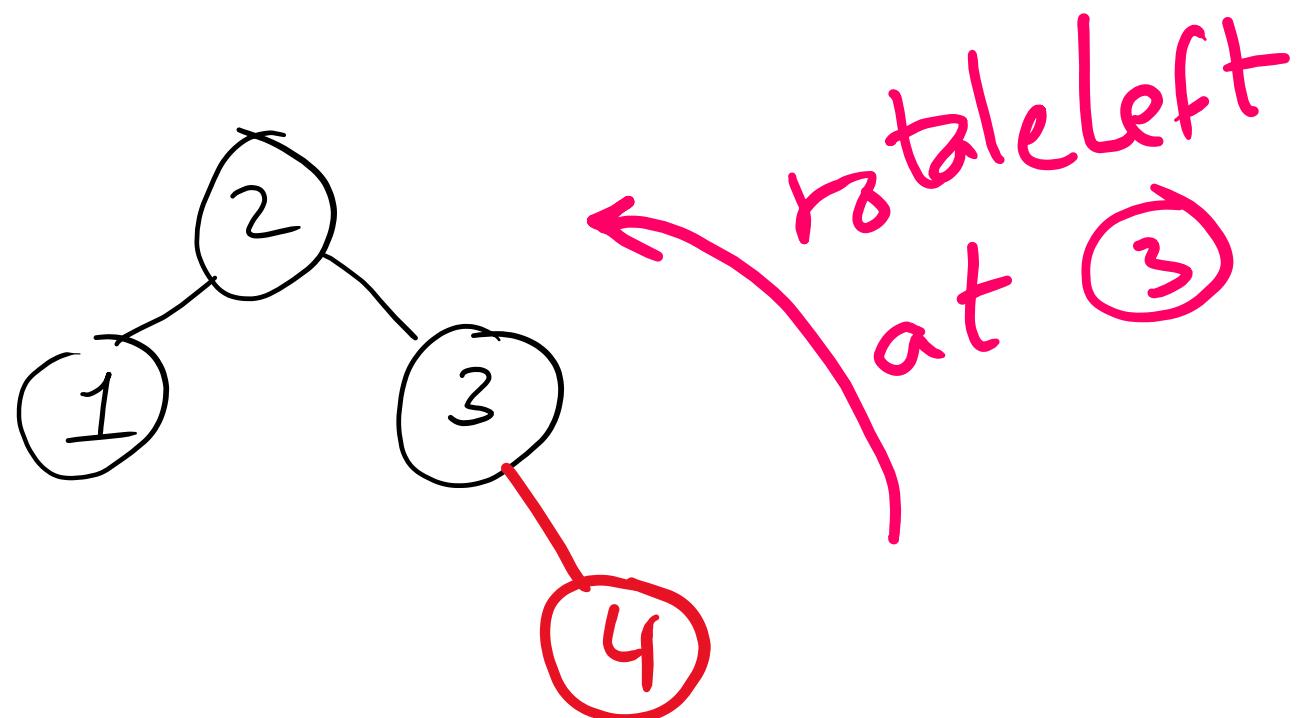
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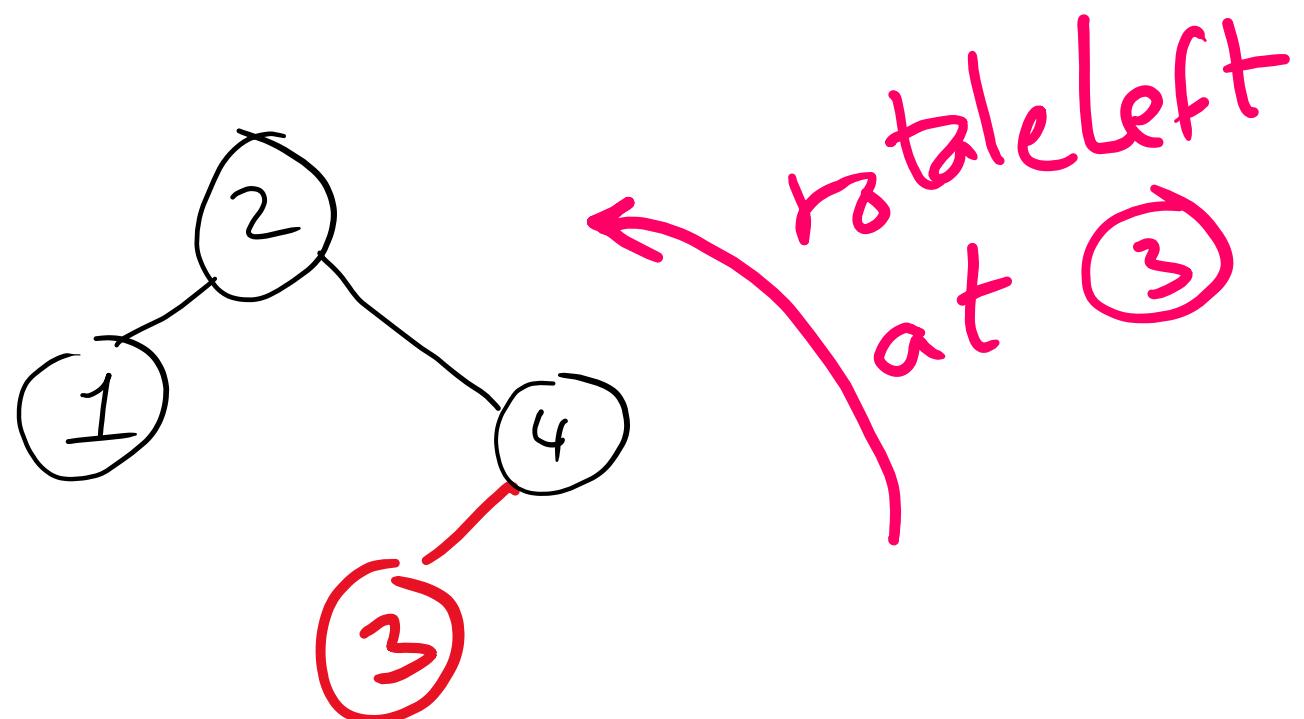
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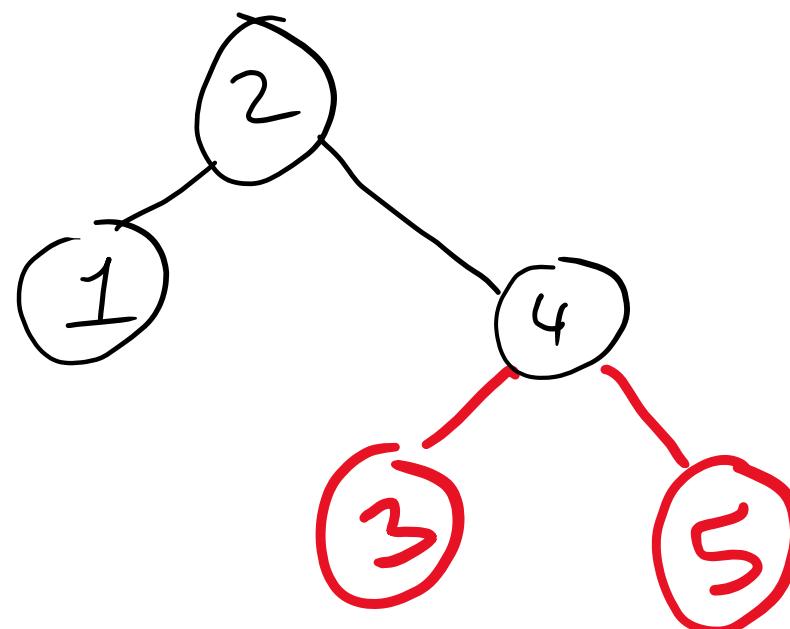
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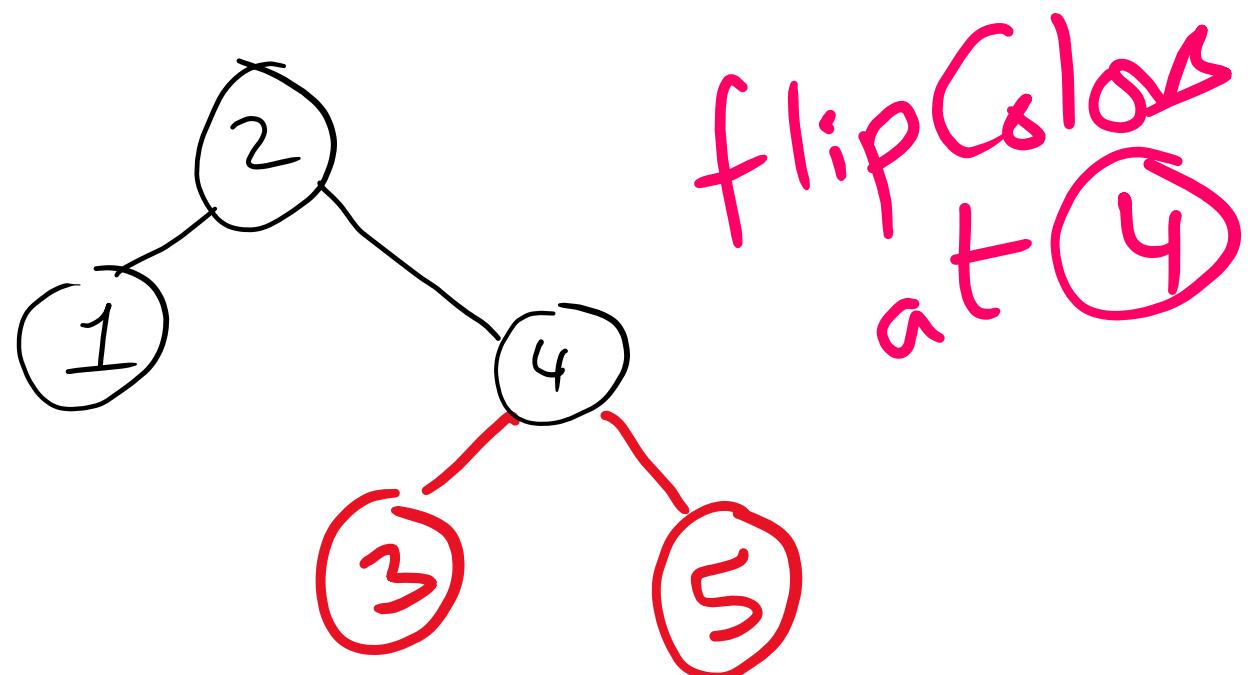
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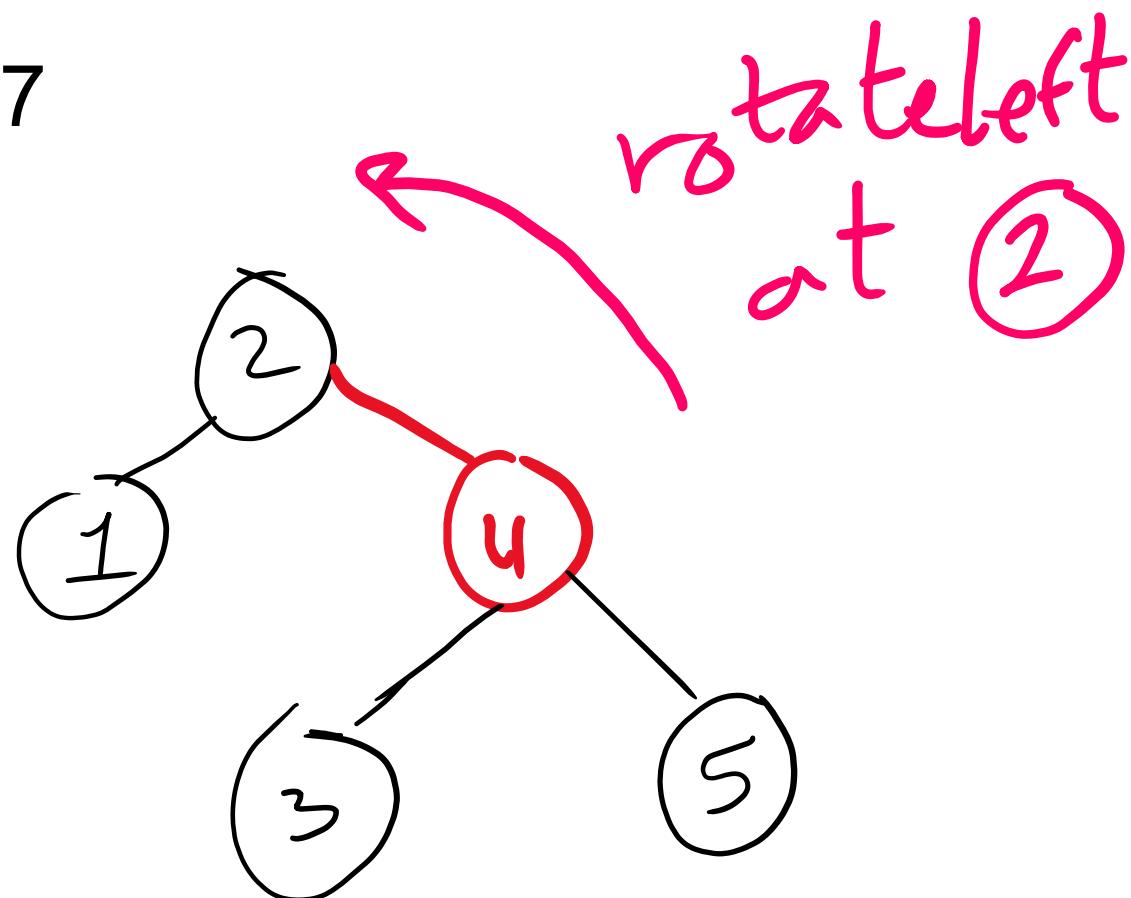
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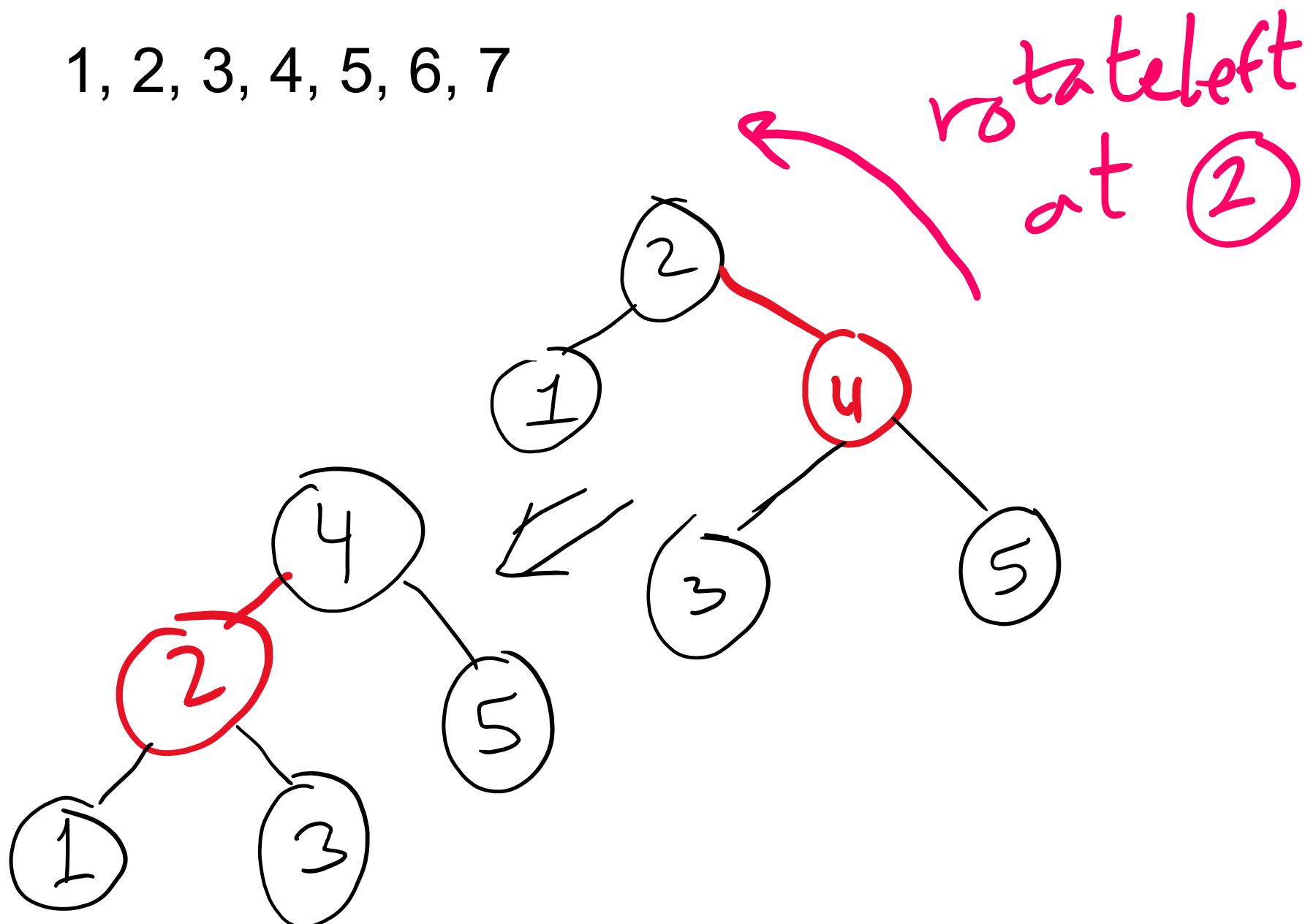
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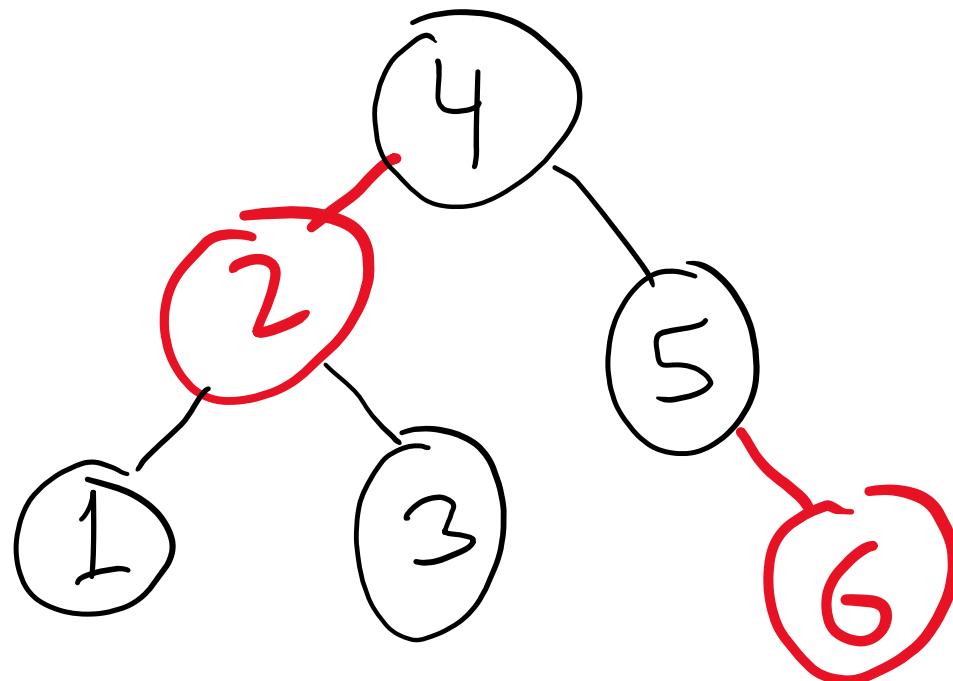
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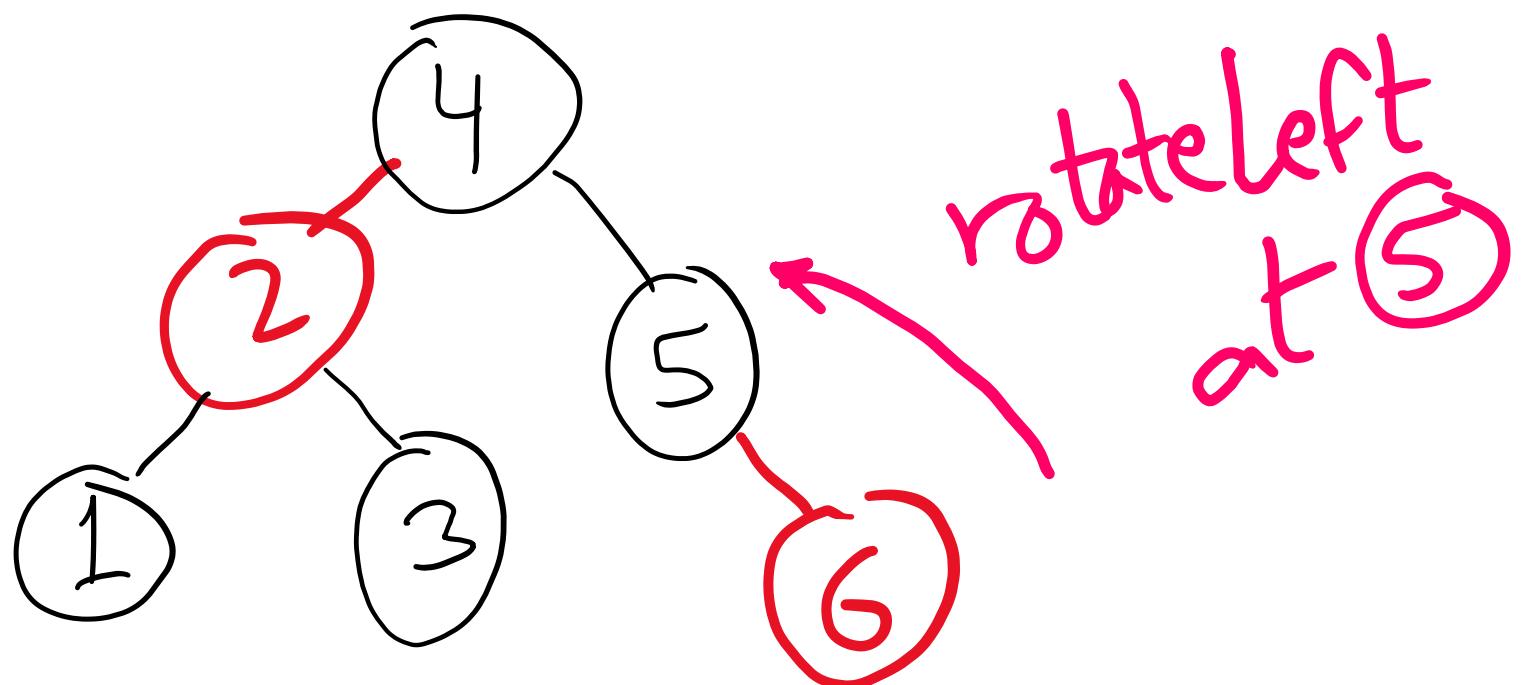
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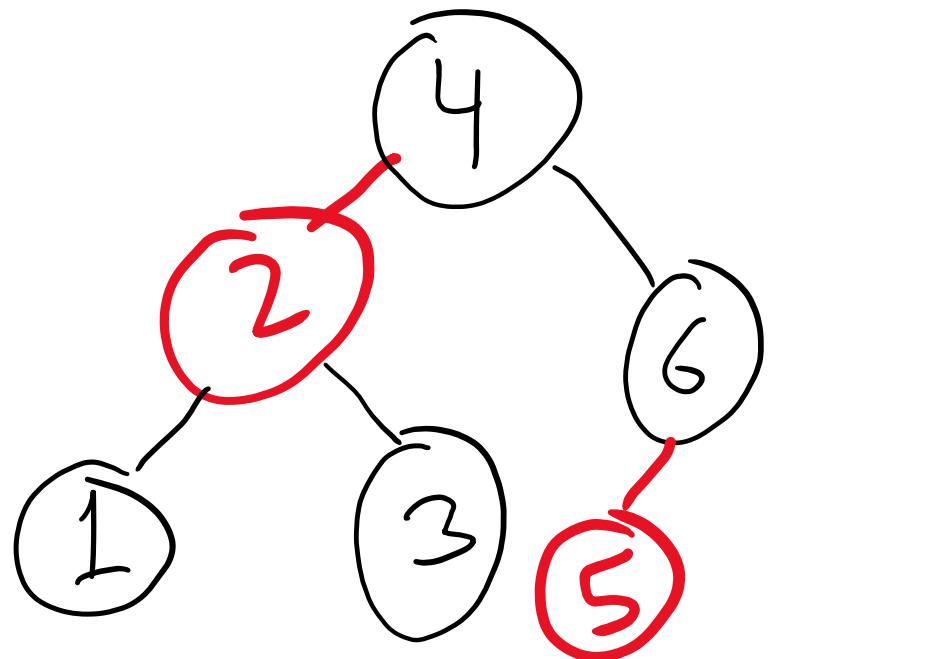
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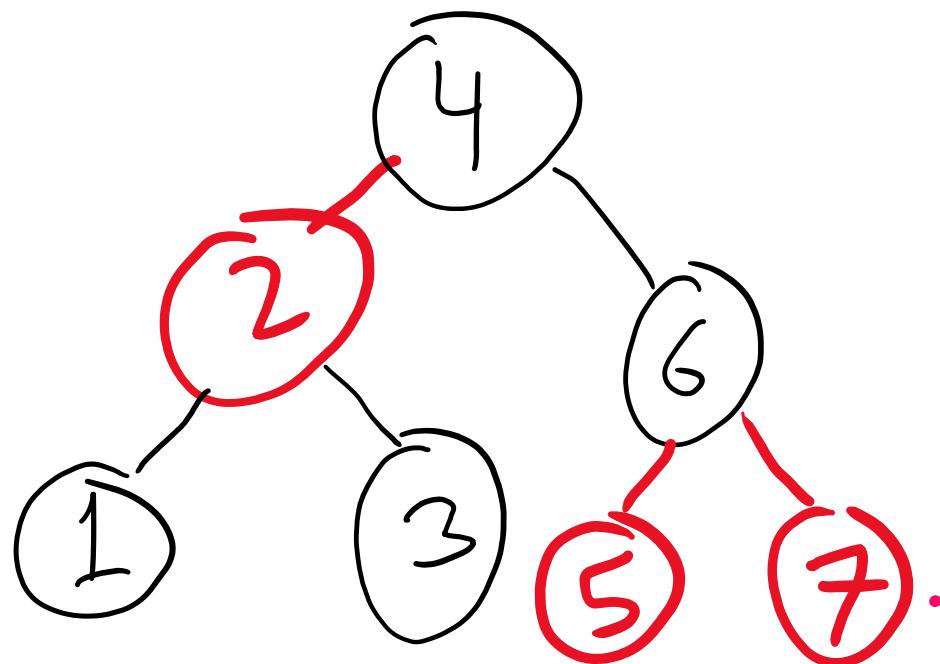
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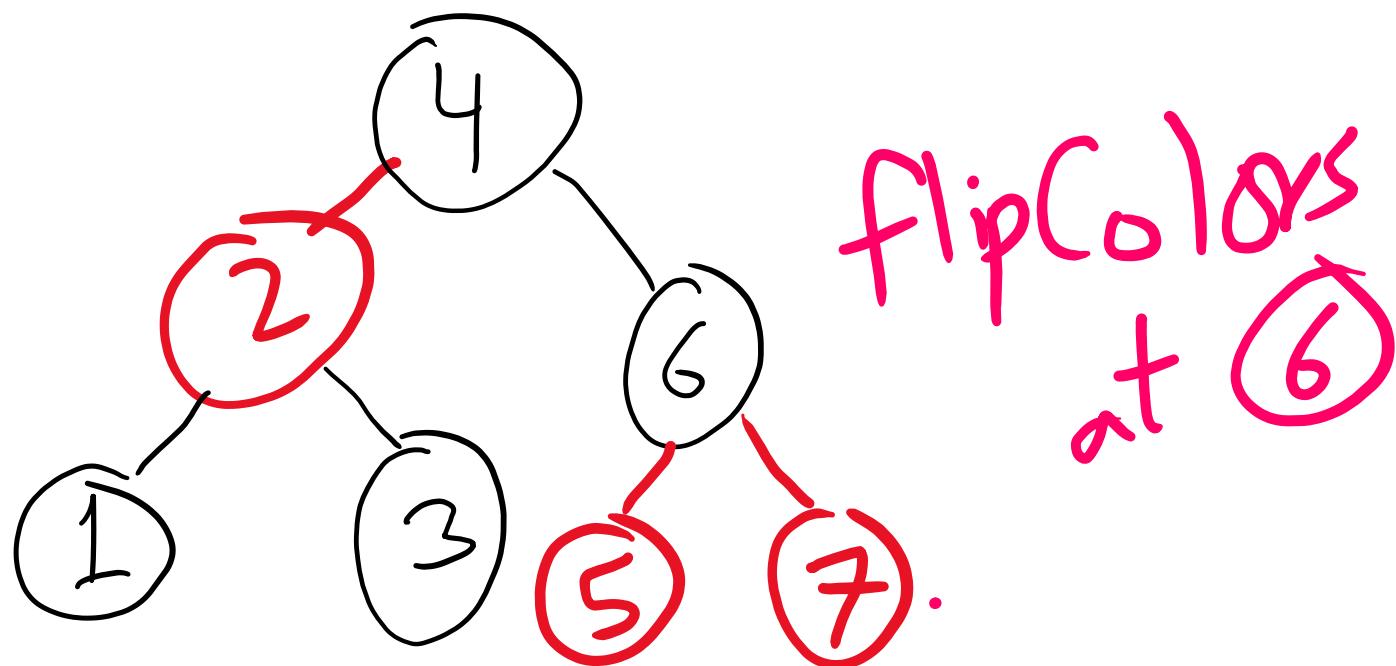
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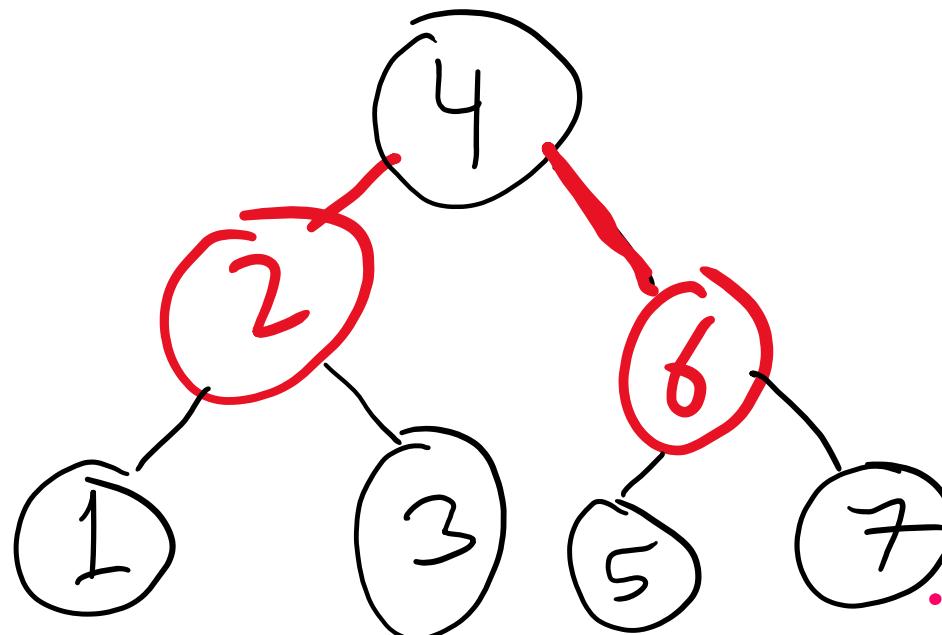
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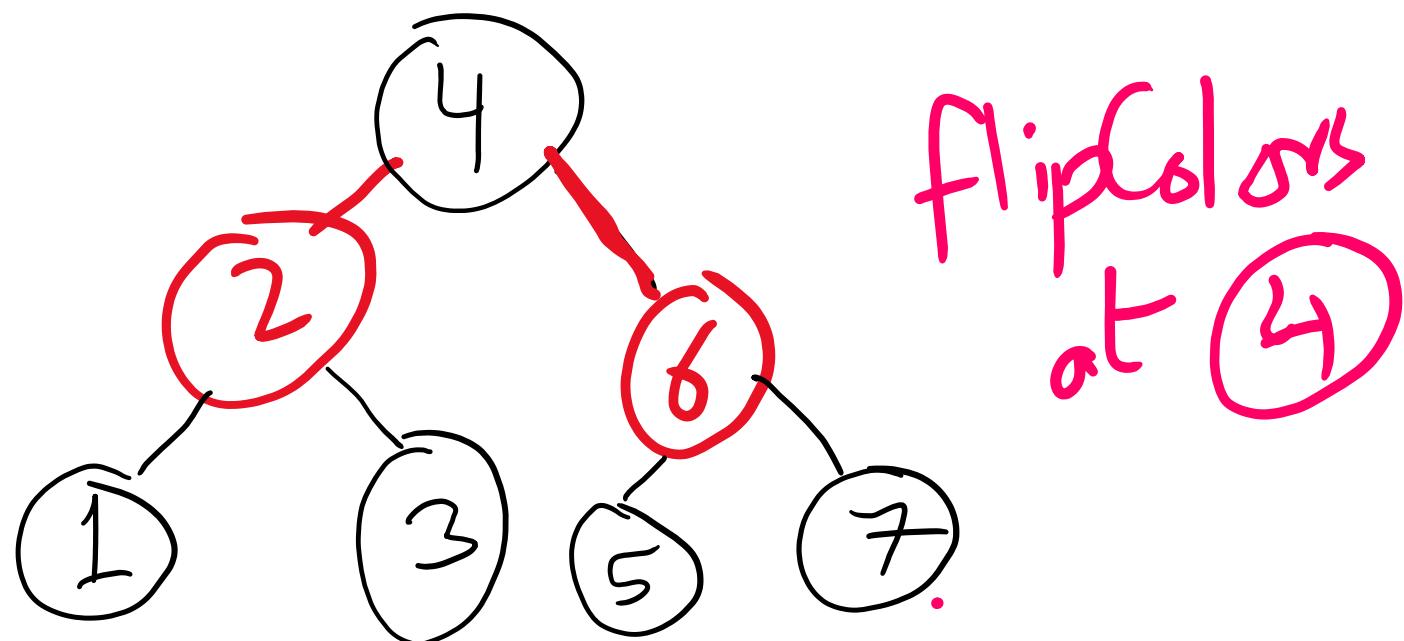
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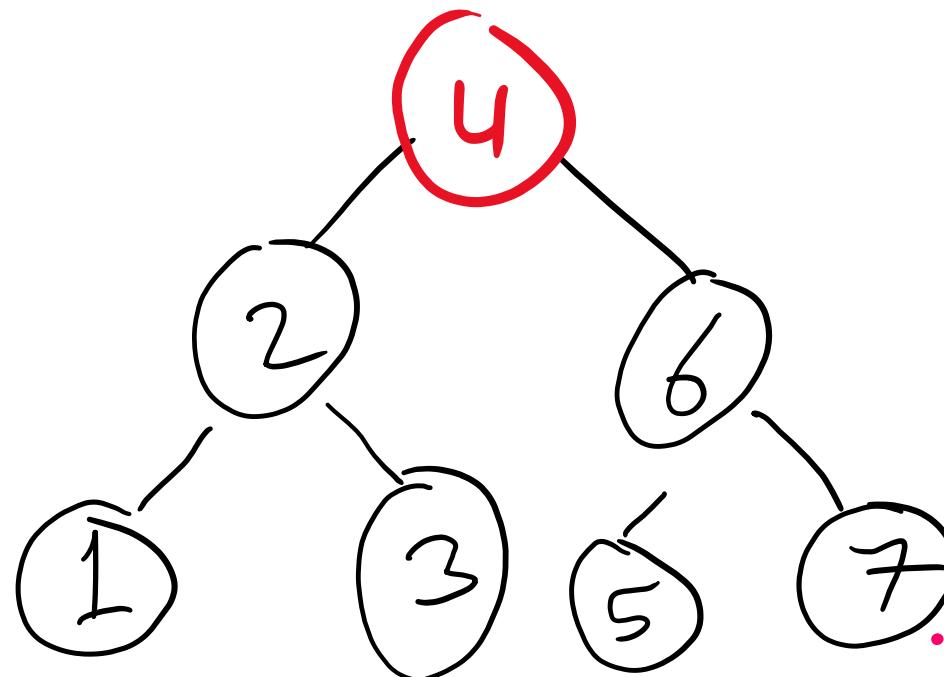
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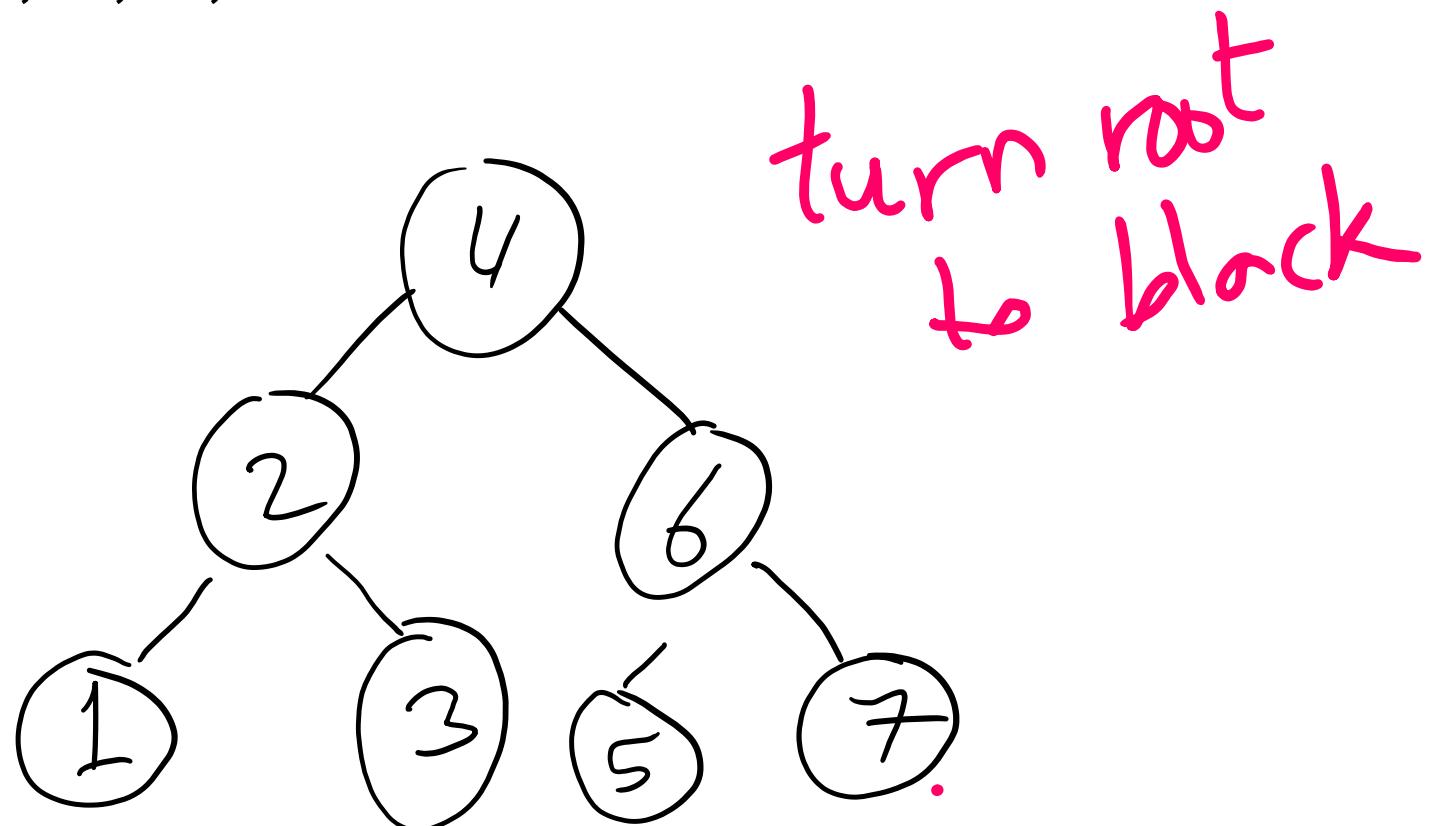
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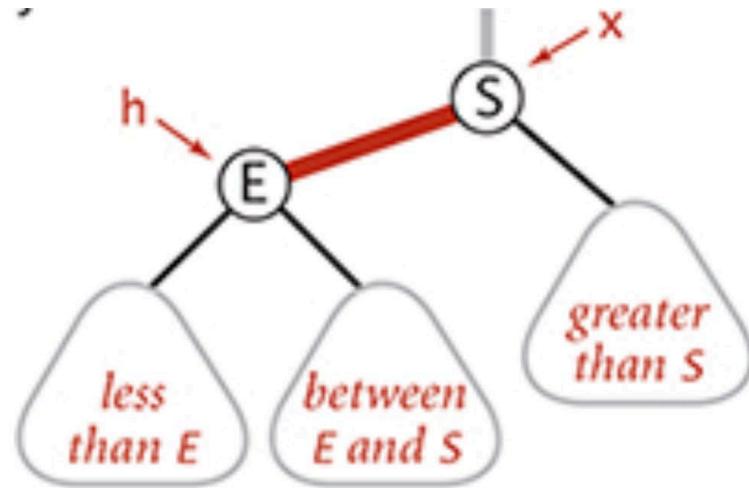
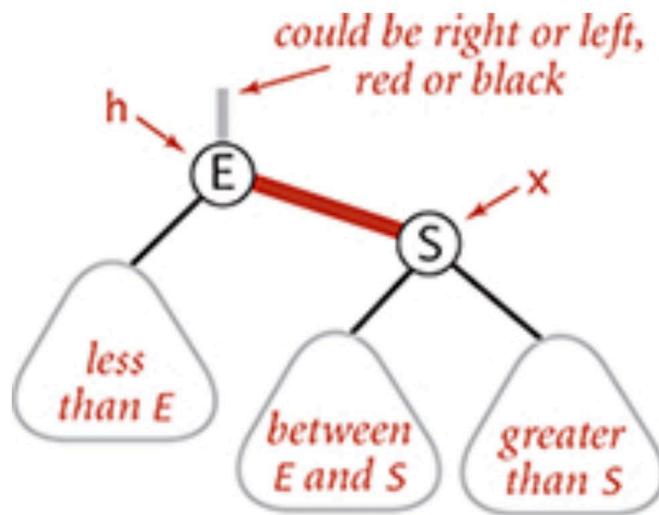
Adding to a RB-BST: Example

- Let's add the numbers:
- 1, 2, 3, 4, 5, 6, 7



Left Rotate

- Q: How to perform left rotations?
- A: **BinaryNode<T> rotateLeft(BinaryNode<T> h)** performs left rotation at node h . The code emerges from contrasting the before and after pictures

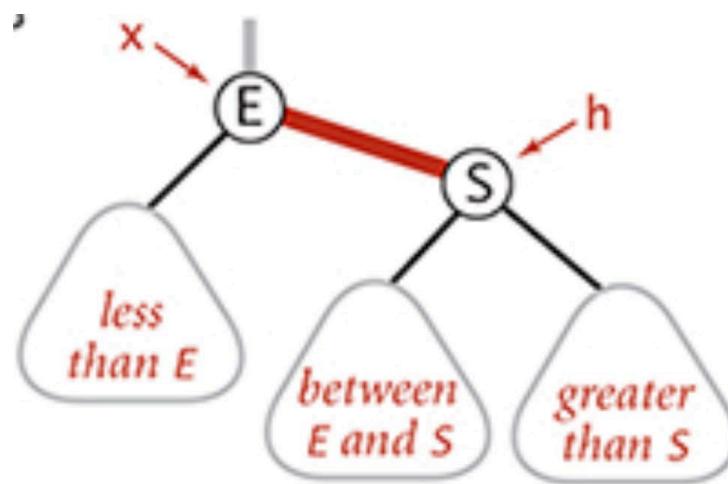
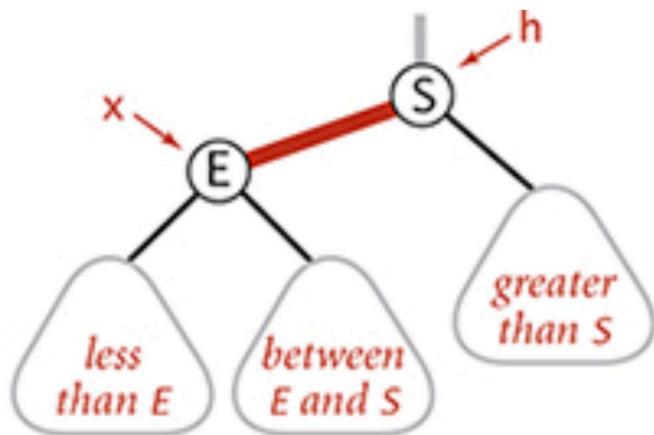


```
x = h.right;  
h.right = x.left;  
x.left = h;  
x.color = h.color;  
h.color = RED;  
return x;
```

- `rotateLeft` maintains the search tree property
 - why?
- It also maintains the perfect balance property
 - why?

Right Rotate

- Q: How to perform right rotations?
- A: **BinaryNode<T> rotatRight(BinaryNode<T> h)** performs right rotation at node h . The code emerges from contrasting the before and after pictures

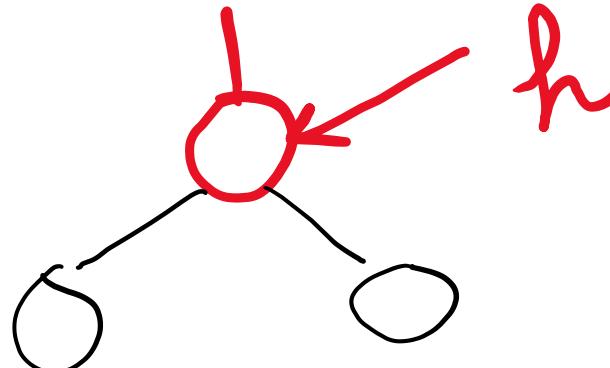
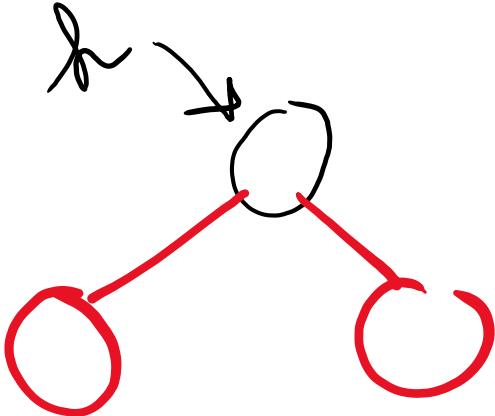


```
x = h.left;  
h.left = x.right;  
x.right = h;  
x.color = h.color;  
h.color = RED;  
return x;
```

- `rotateRight` maintains the search tree property
 - why?
- It also maintains the perfect balance property
 - why?

Color Flip

- Q: How to perform color flips?
- A: **BinaryNode<T> flipColors(BinaryNode<T> h)** performs color flip at node h . The code emerges from contrasting the before and after pictures



```
x = h.left;  
h.left = x.right;  
x.right = h;  
x.color = h.color;  
h.color = RED;  
return h;
```

- flipColors maintains the search tree property
 - why?
- It also maintains the perfect balance property
 - why?

Muddiest Points

- **Q: If correcting one violation for the red-black BST, leads to another violation, how is that corrected?**
- A: One violation may lead to another violation of a different type, which may lead to another violation of a different type.
- However, this process stops after at most three corrections passing the violation to the next level up.
- Violations keep propagating up the tree until we have a violation at root, which is root node being red instead of black
 - This can be easily corrected by coloring root black again

Muddiest Points

- Q: Confused on the benefits of adding red nodes and what that would look like in an application
- A: Adding a black node would violate the perfect balance constraint (all paths from root to null-edges have the same number of black edges)
 - Hard to fix!

Muddiest Points

- **Q: can you explain the implementations again or point me to a resource that explains them if that would take too much time?**
- A: The code is available in the code handouts. The lecture recordings are available on Canvas.
- **Q: It is hard to visualize a red black bst without pictures and diagrams for me**
- A: I agree! When you study, use diagrams of example trees.
- **Q: PLEASE GO THROUGH THE HOMEWORKS**
- A: Please use Piazza to ask about any homework questions that you are confused about

Muddiest Points

Anonymous

a day ago



Comments 0

1



This Lecture

- Red-Black BST (self-balancing BST)
 - definition and basic operations
 - delete
 - runtime of operations
- Turning recursive tree traversals to iterative

Deleting a node

- Make sure that we are not deleting a black node
 - as we go down the tree, make sure that the next node down is red
 - using a different set of operations
 - as we go back up the tree, correct any violations
 - same as we did while adding
 - if deleting a node with 2 children
 - replace with minimum of right subtree
 - delete minimum of right subtree
 - similar trick to delete in regular BST

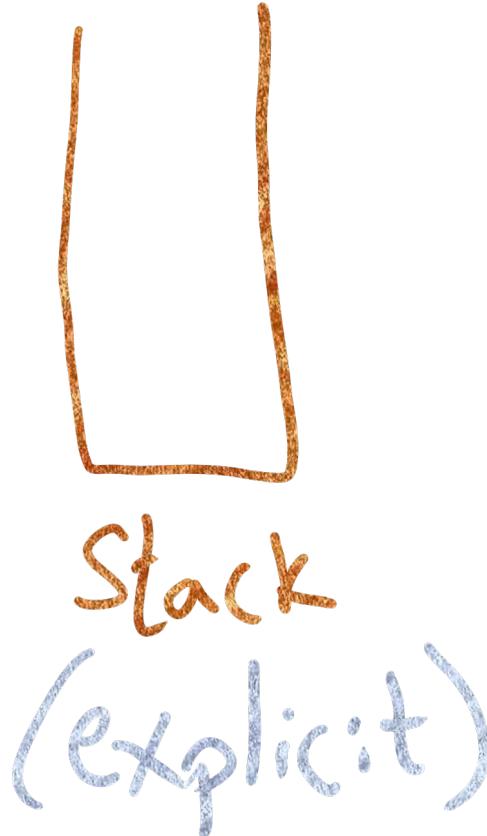
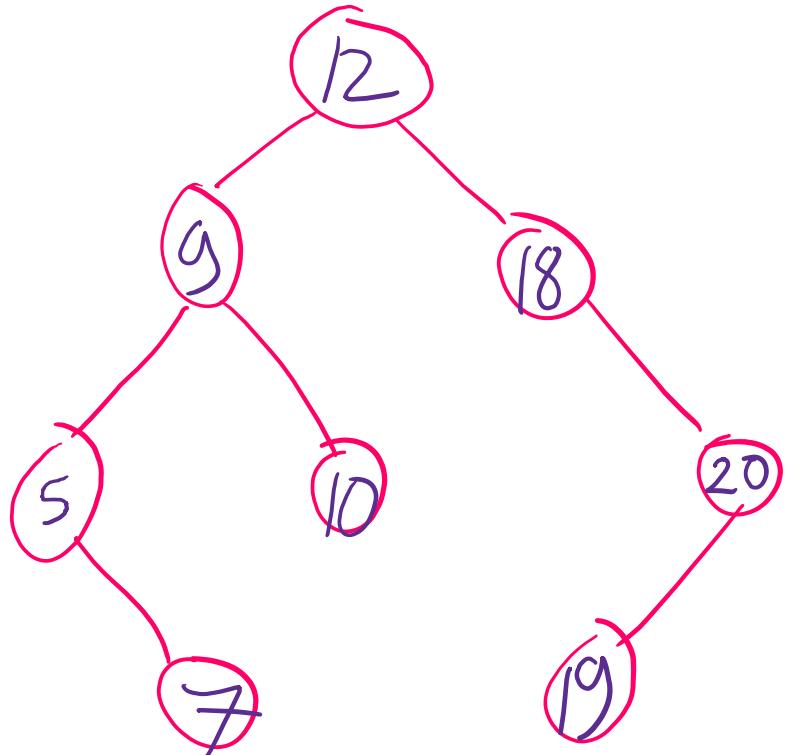
Other BST operations

- Find successor and predecessor of an item
 - Lab 3
- Find all items within a specific range
 - Please check the keys methods in RedBlackBST.java inside the TreeADT folder in the code handouts
- Same code as regular BST!
- **worst-case runtime = Theta(log n)**

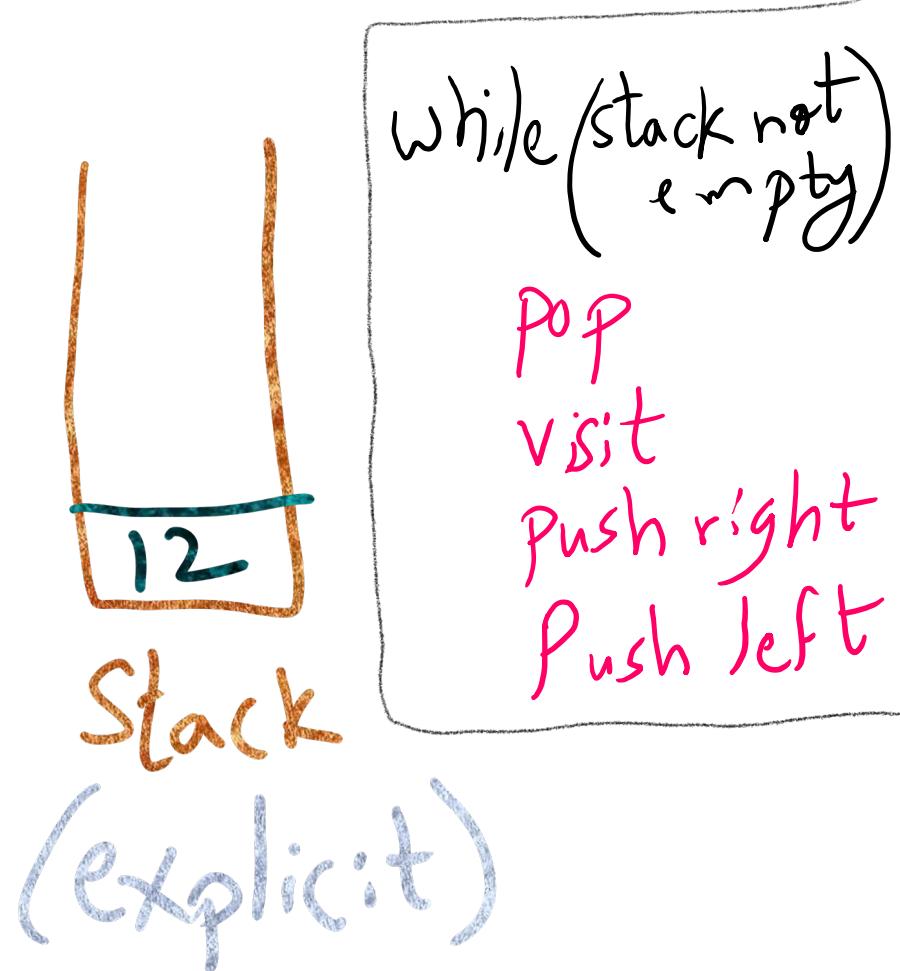
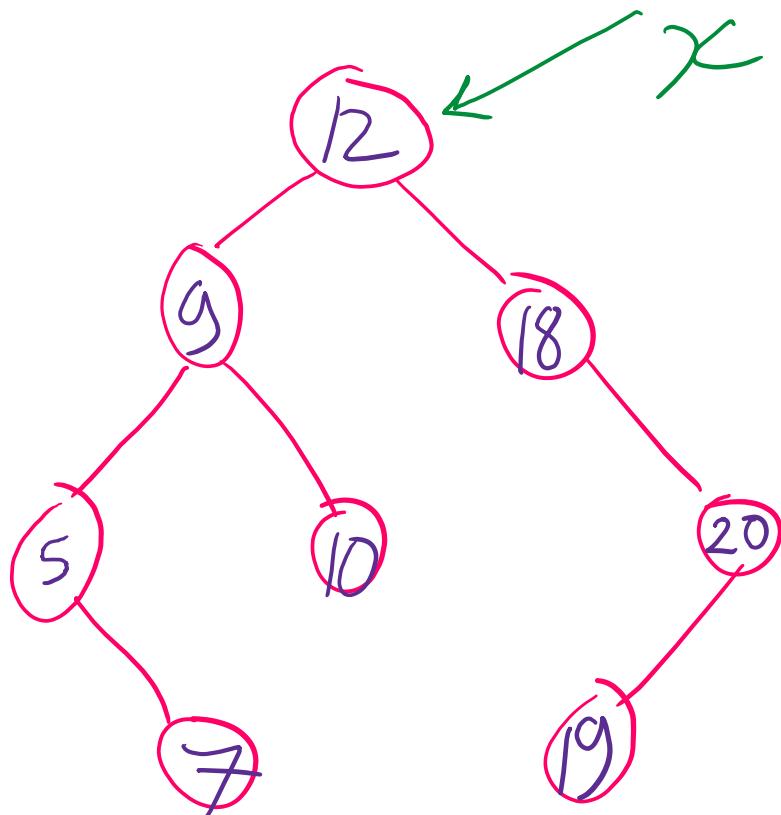
Why use iteration instead of recursion?

- Iteration is faster than recursion
 - No function call overhead (stack allocation)
- Especially useful for frequently used operations
 - Search is a good example of these operations

Iterative Preorder traversal

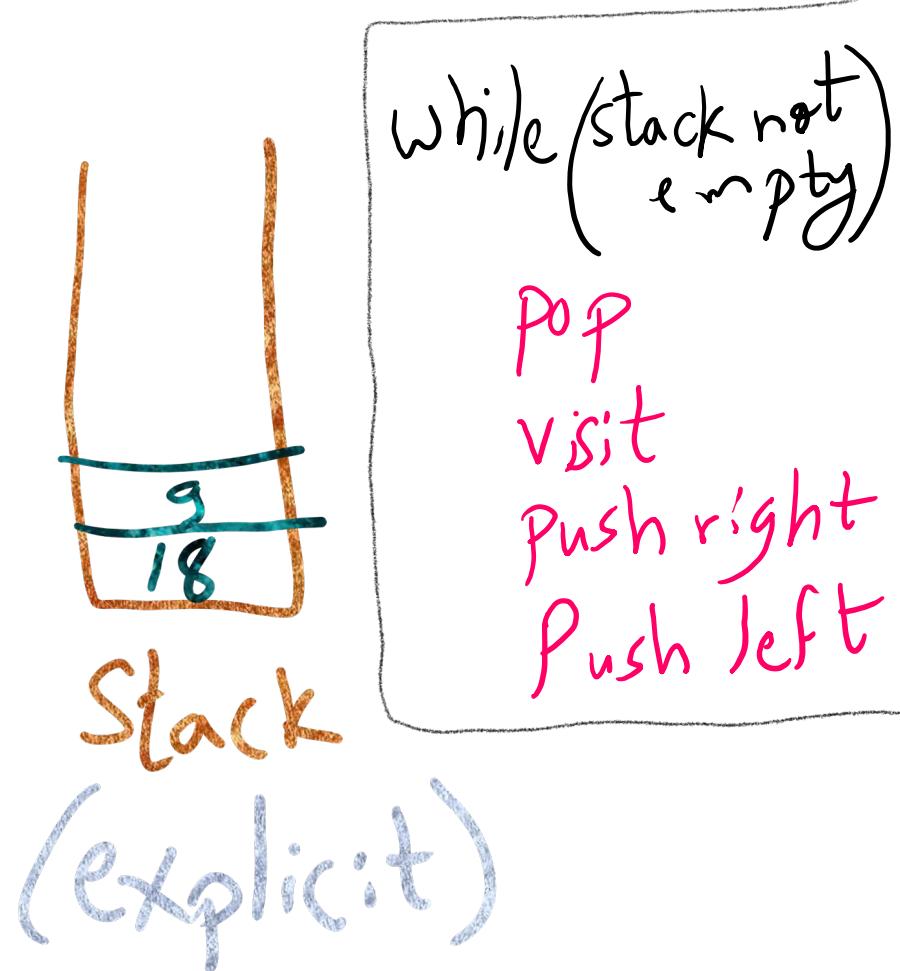
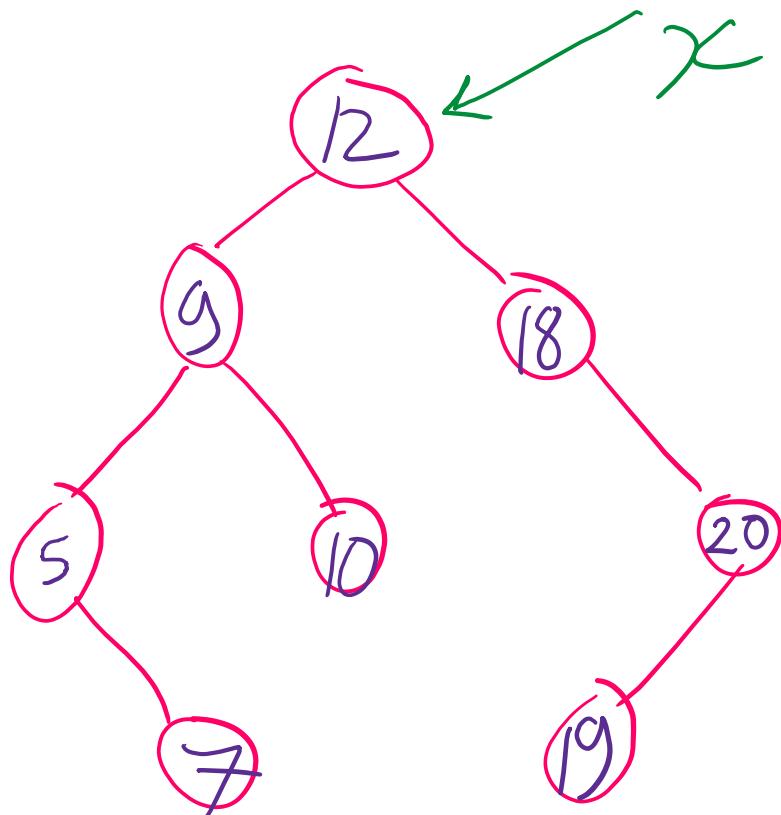


Iterative Preorder traversal



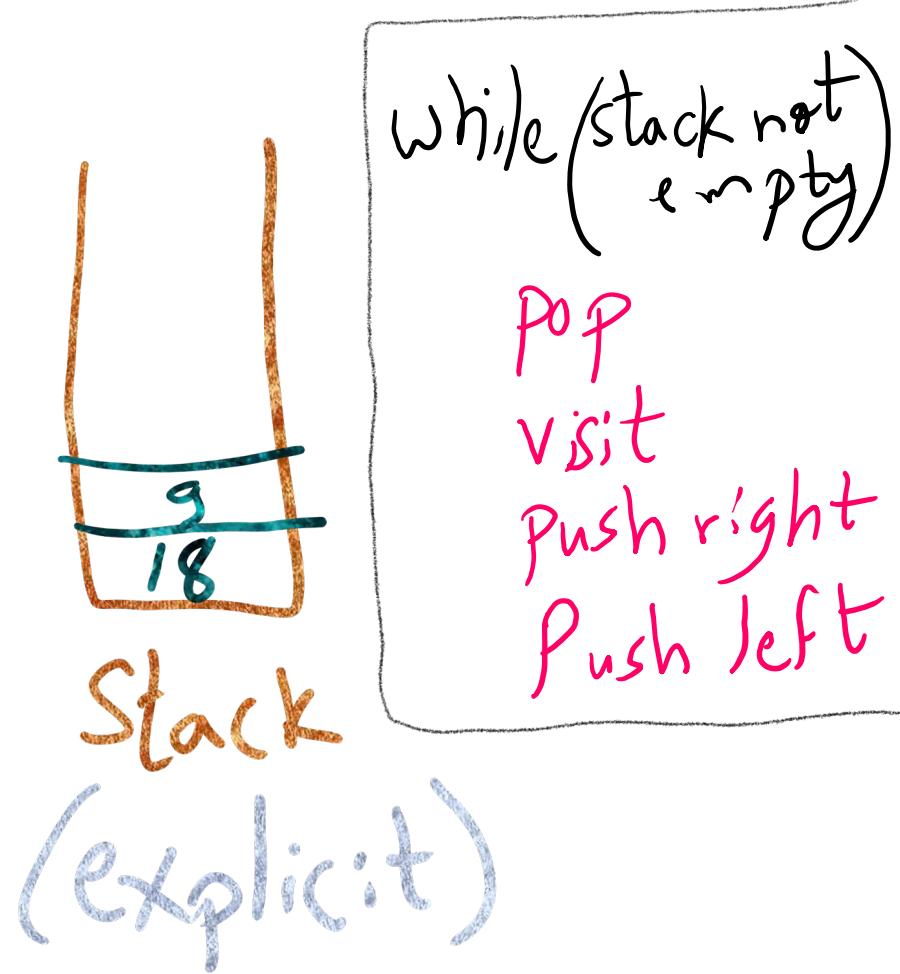
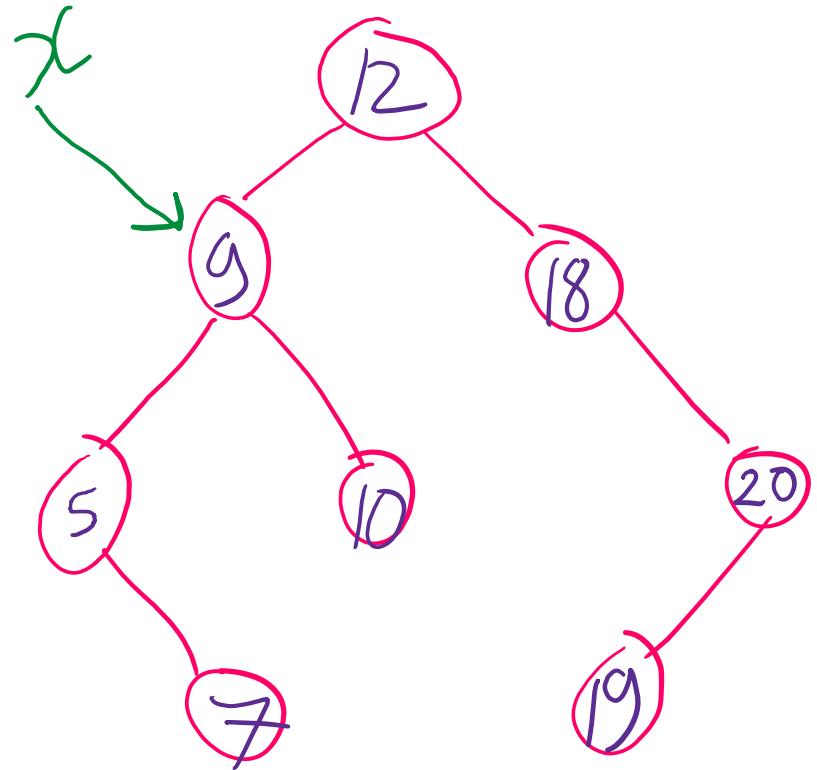
Visit order : 12

Iterative Preorder traversal



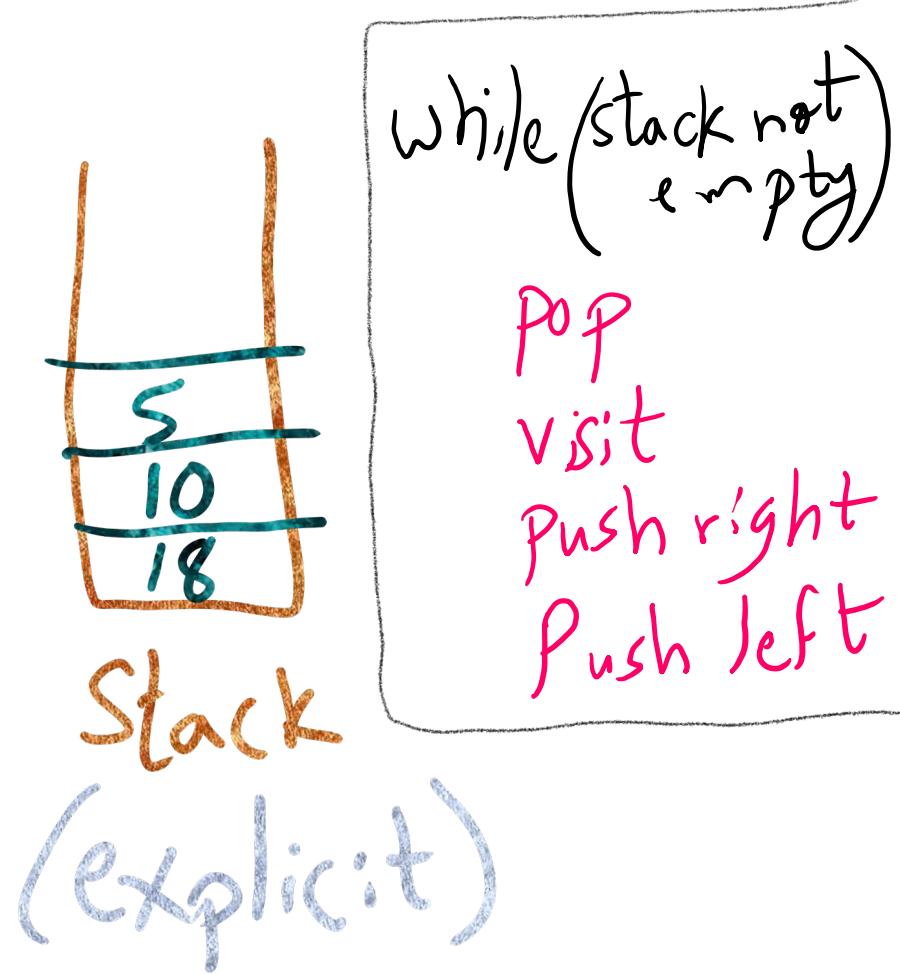
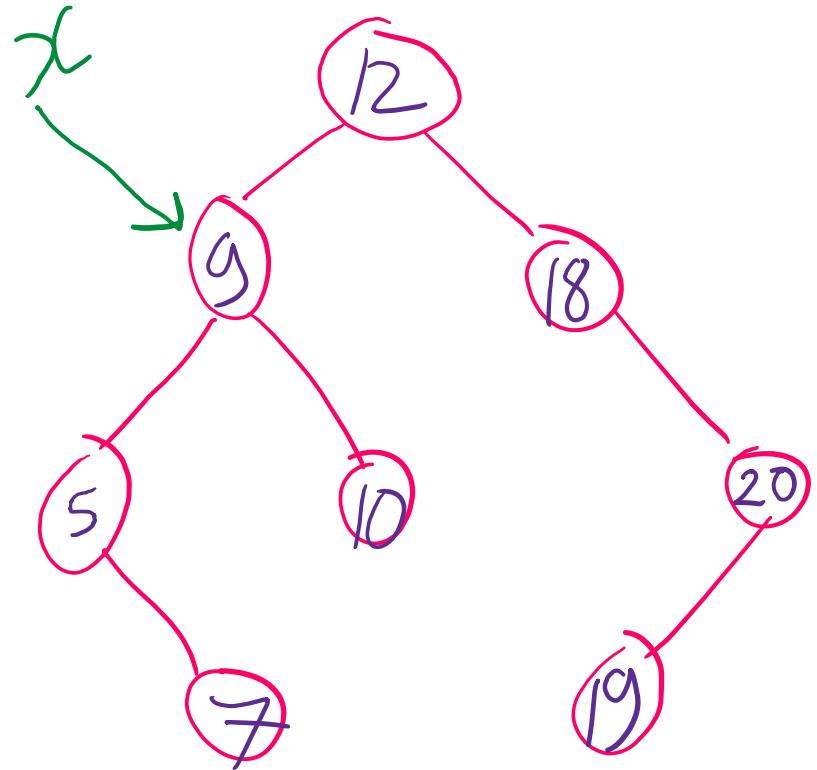
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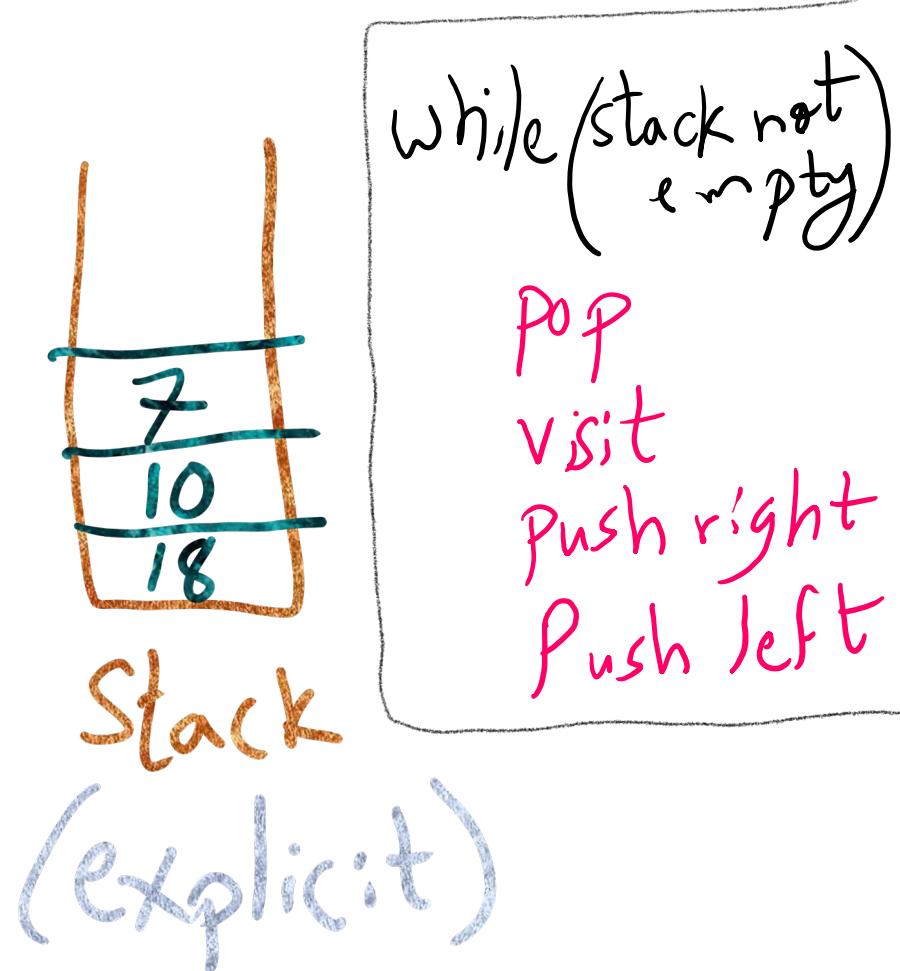
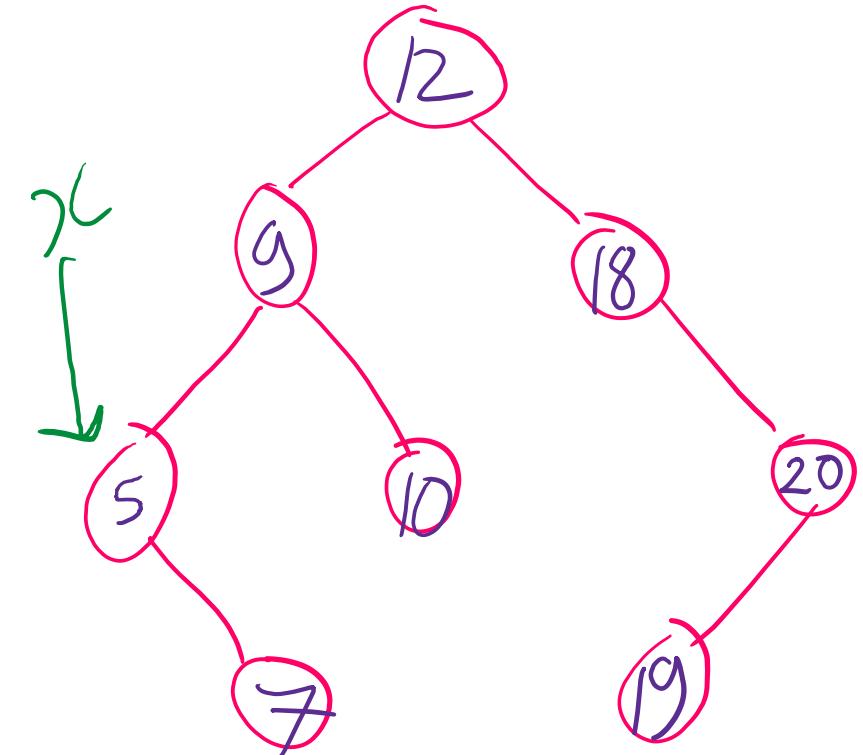
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Iterative Preorder traversal



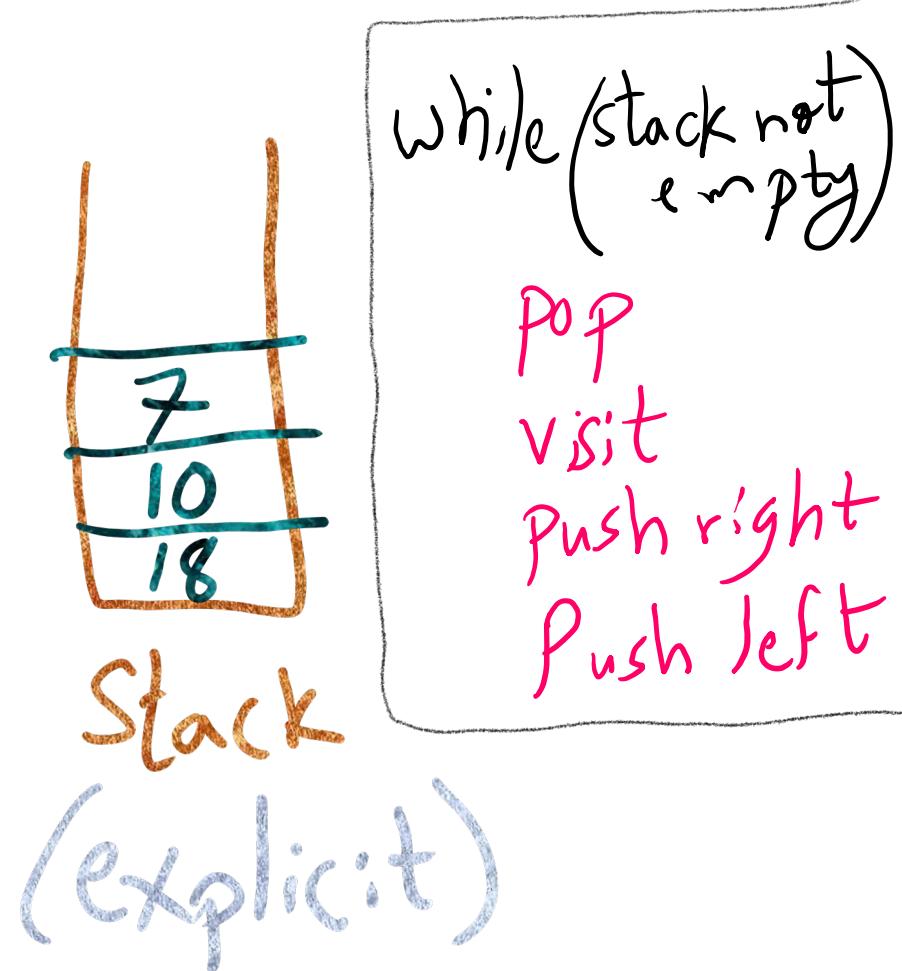
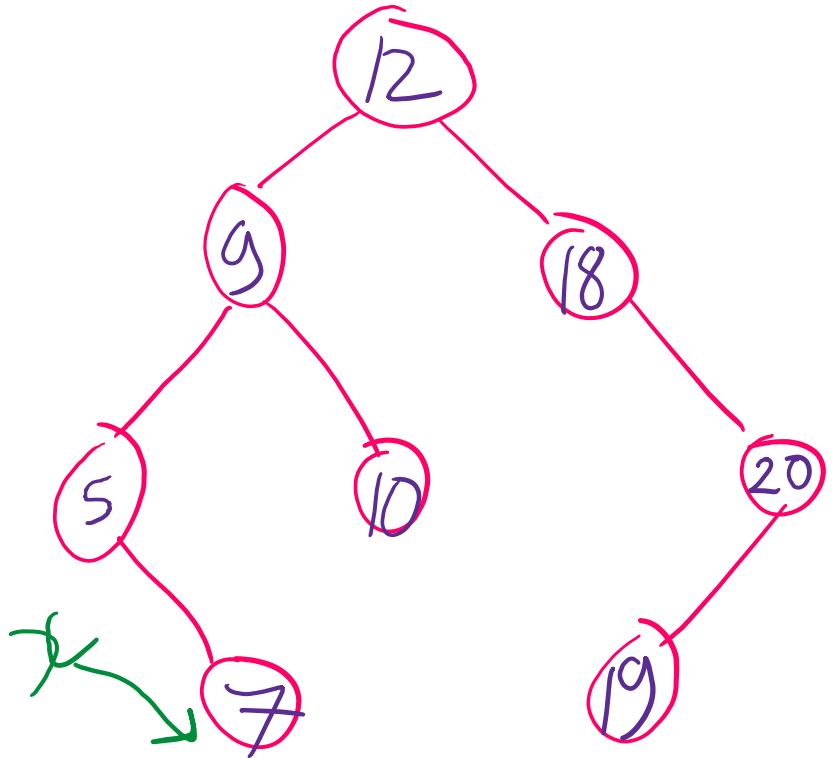
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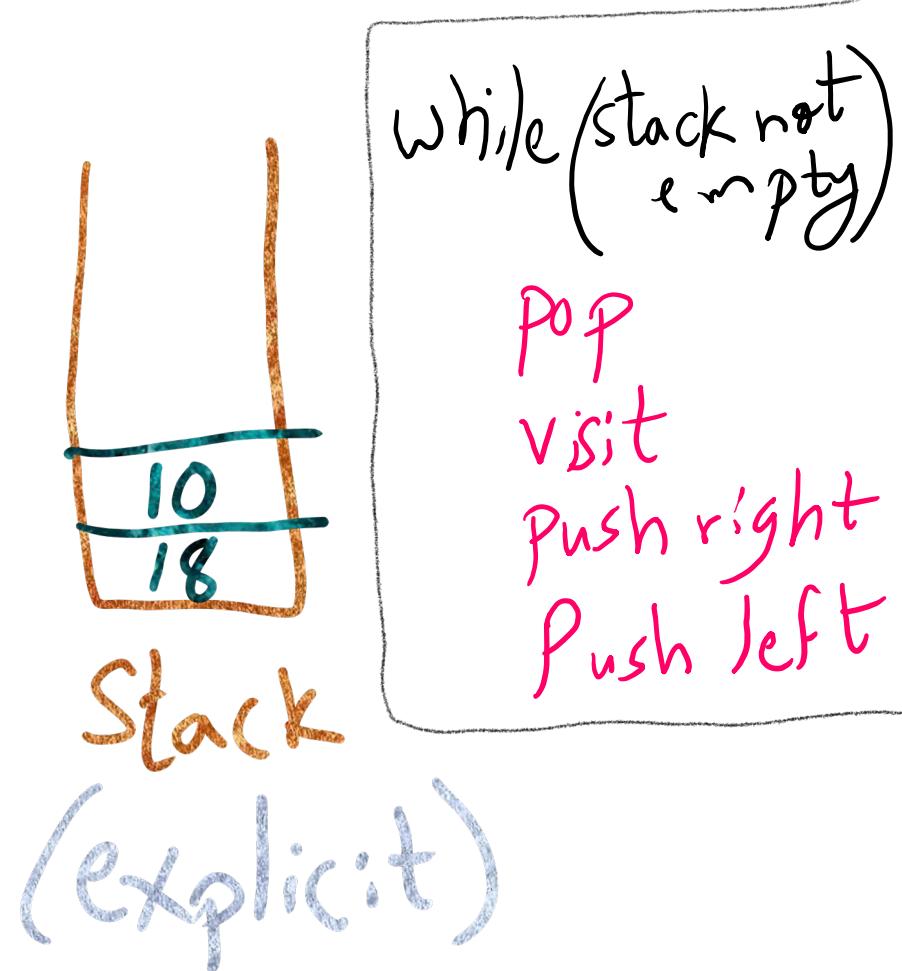
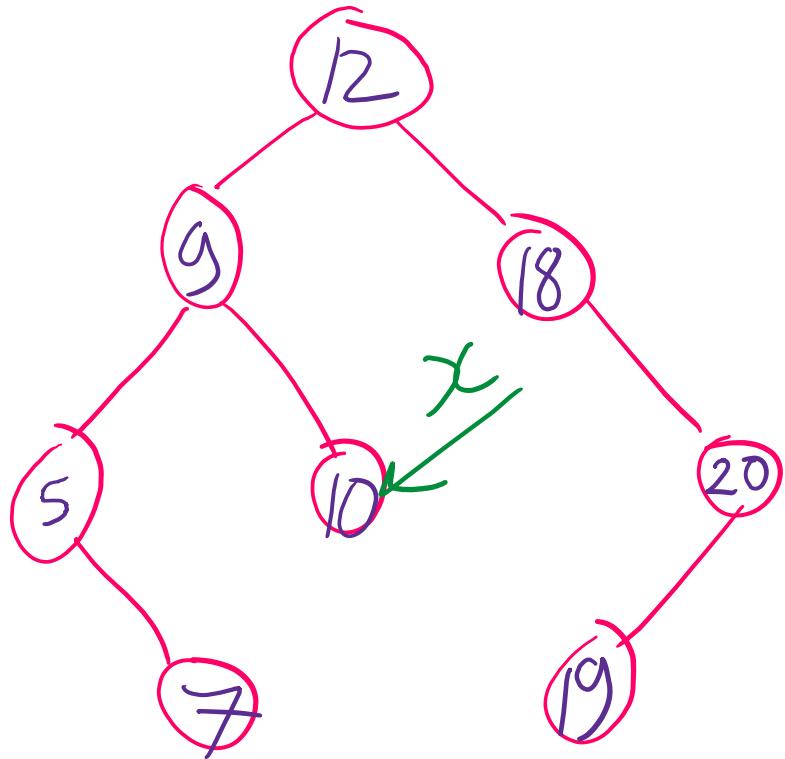
Visit order : 12, 9, 5

Iterative Preorder traversal



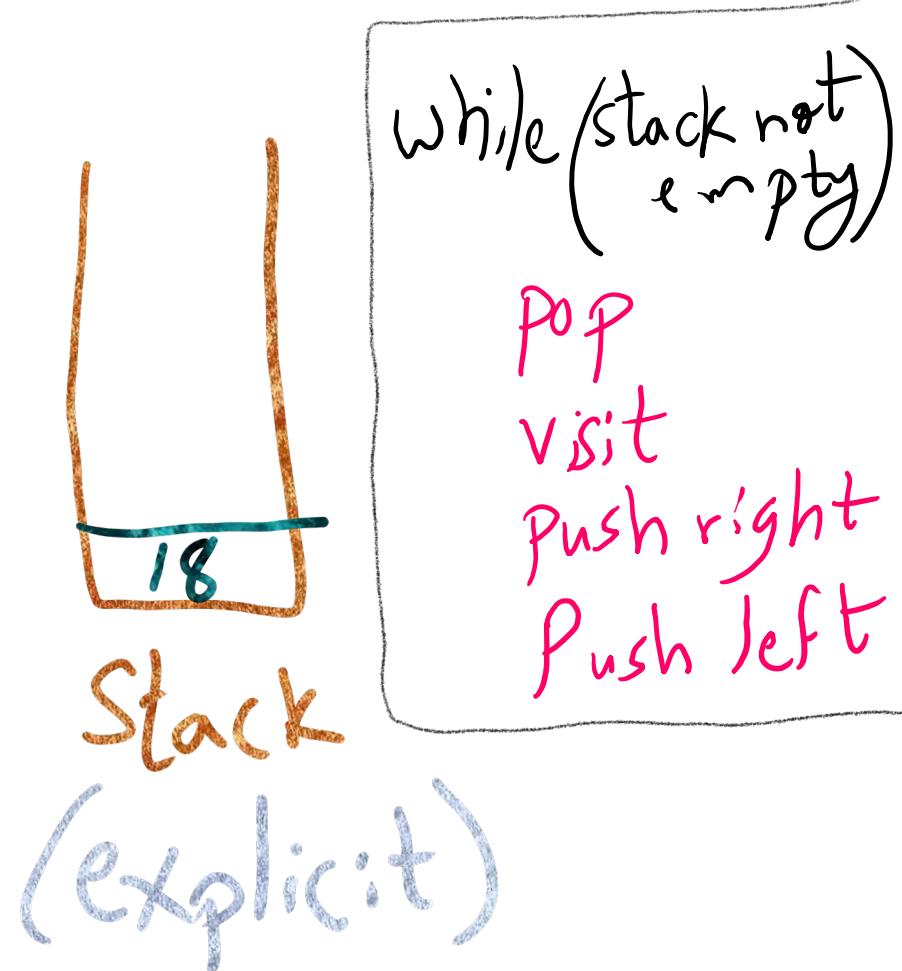
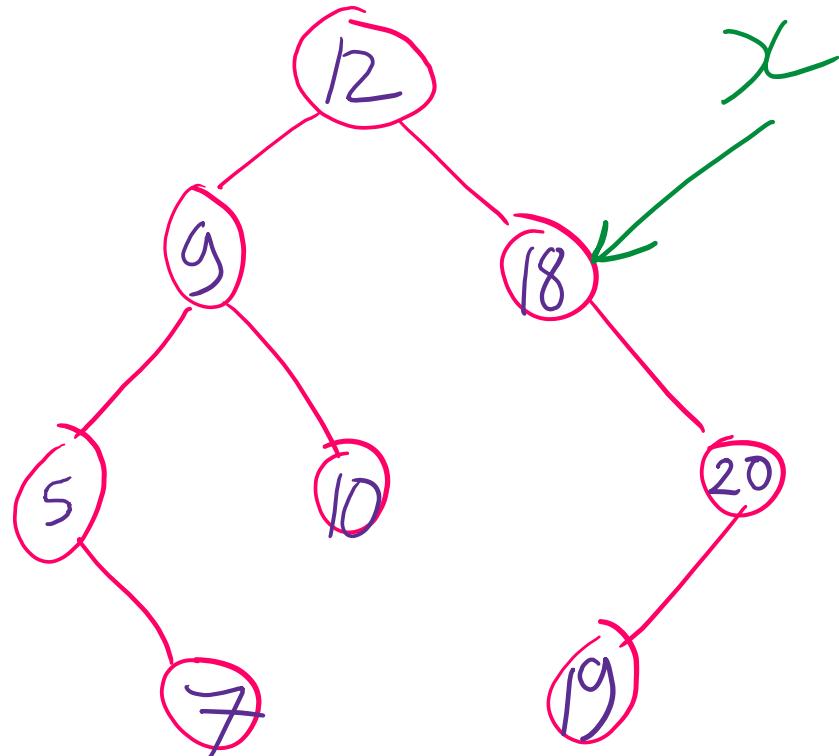
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Iterative Preorder traversal



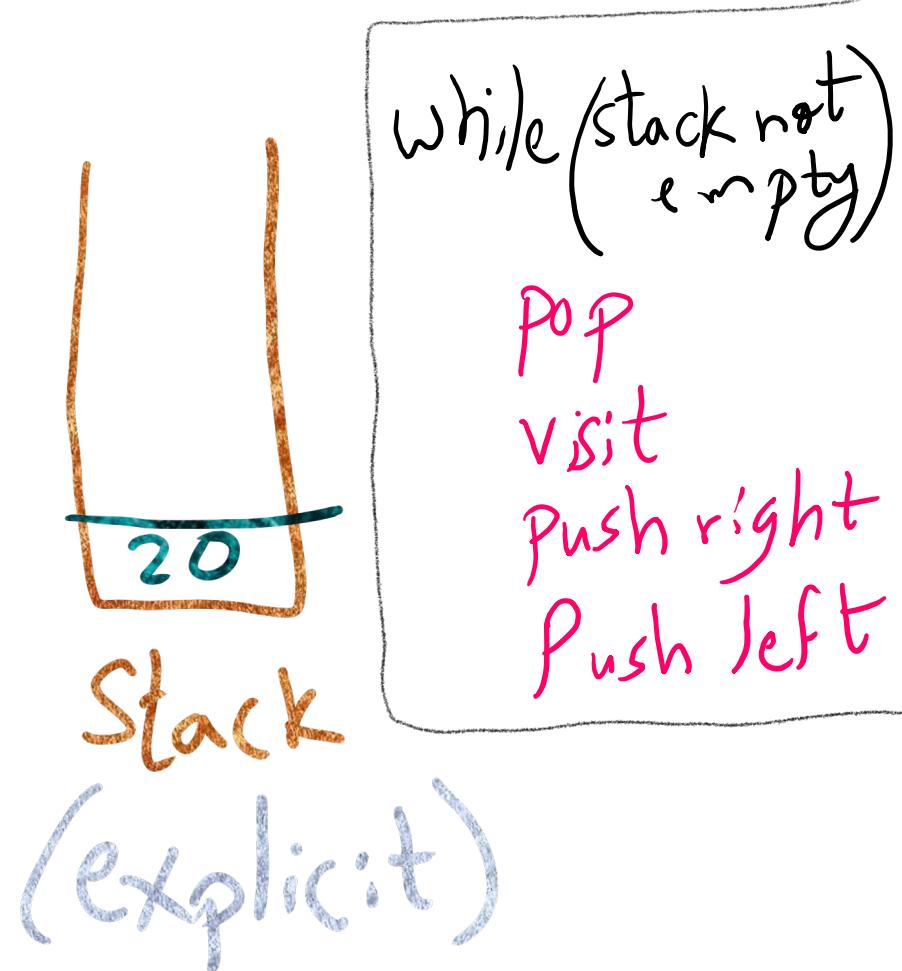
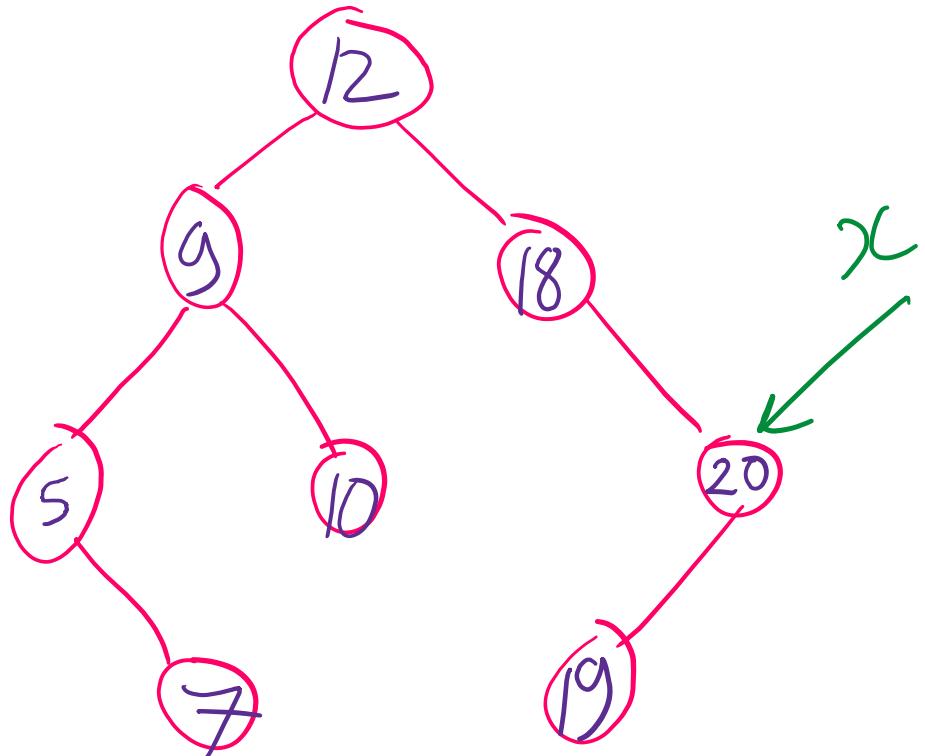
Visit order : 12, 9, 5, 7, 10

Iterative Preorder traversal



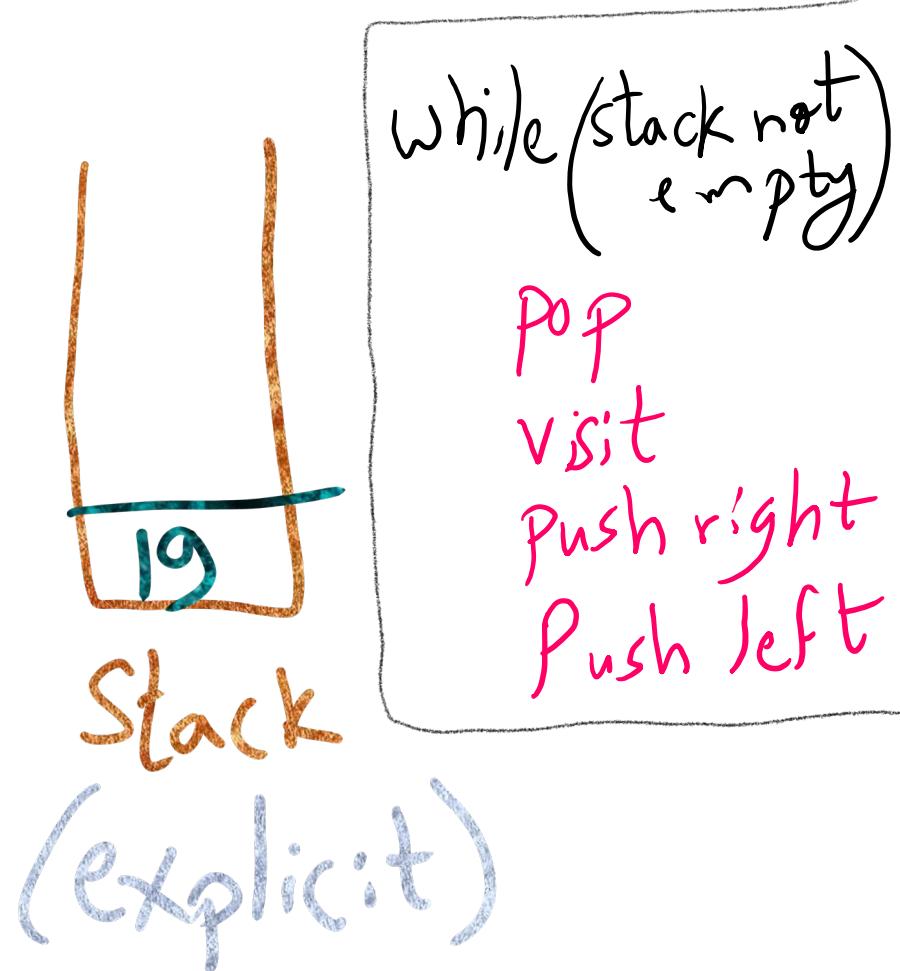
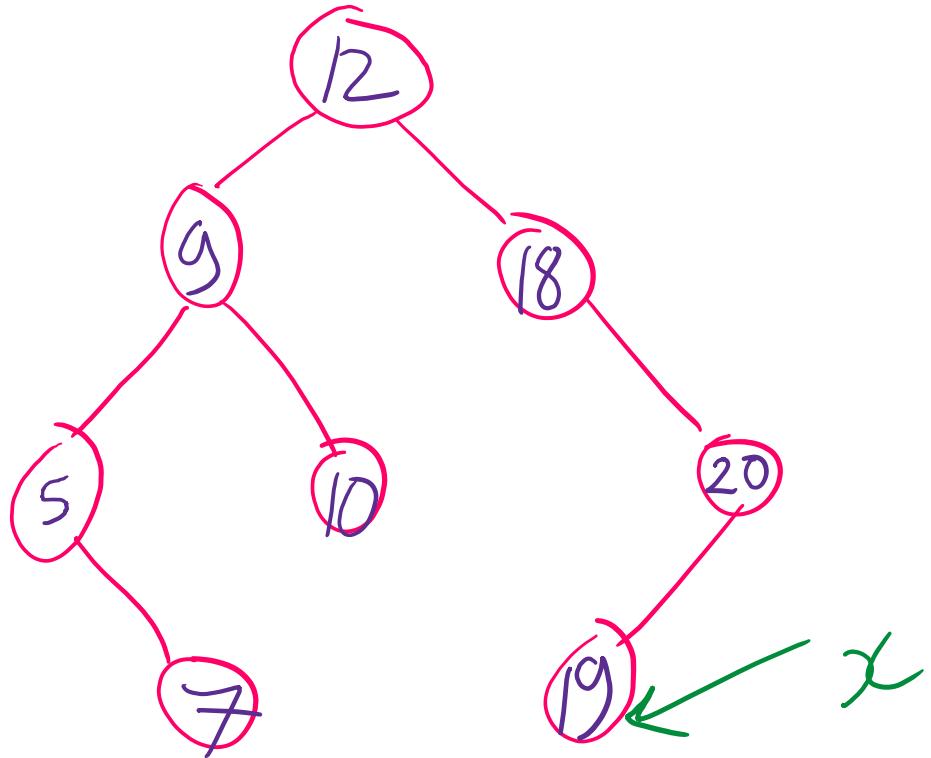
Visit order : 12, 9, 5, 7, 10, 18

Iterative Preorder traversal



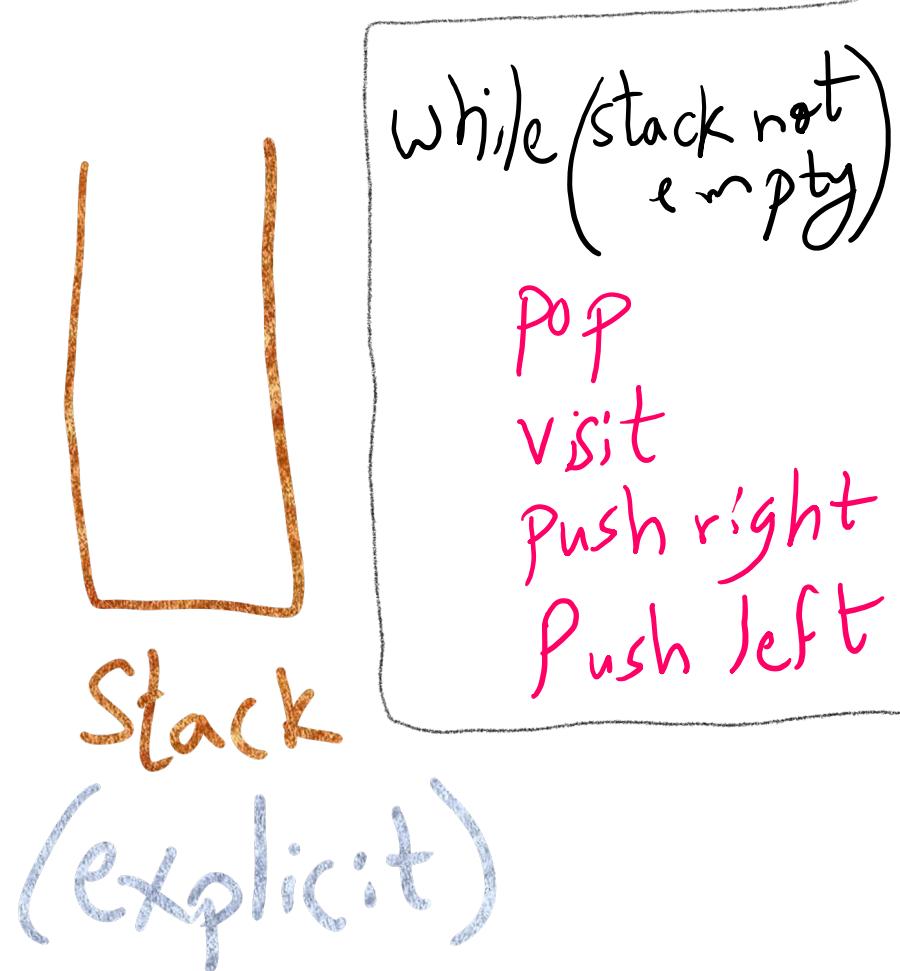
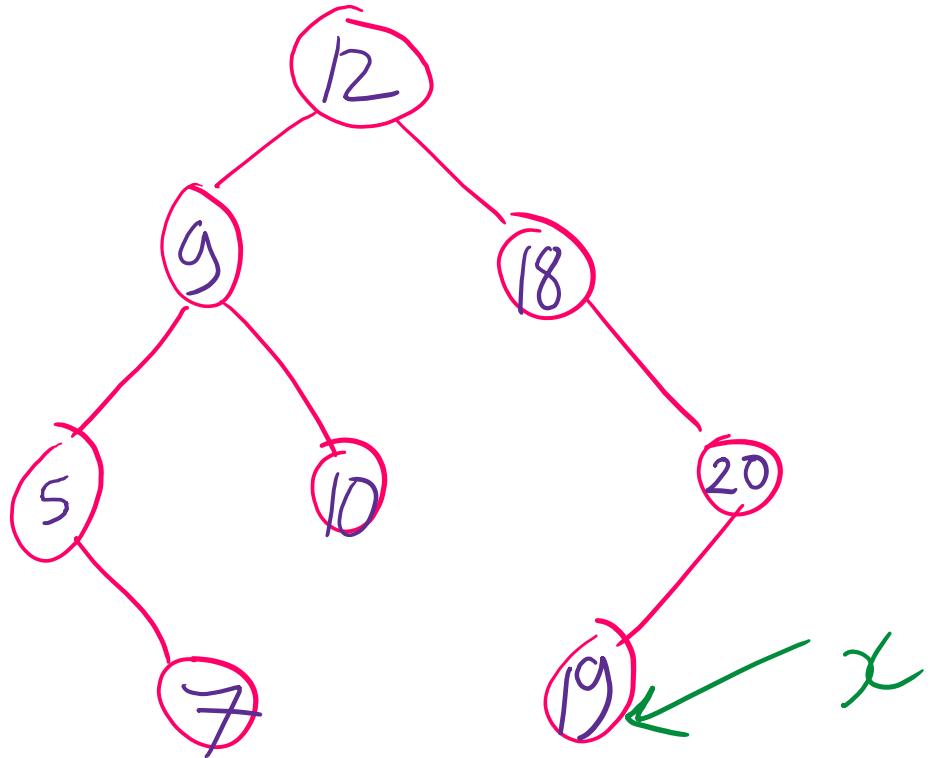
Visit order : 12, 9, 5, 7, 10, 18, 20

Iterative Preorder traversal



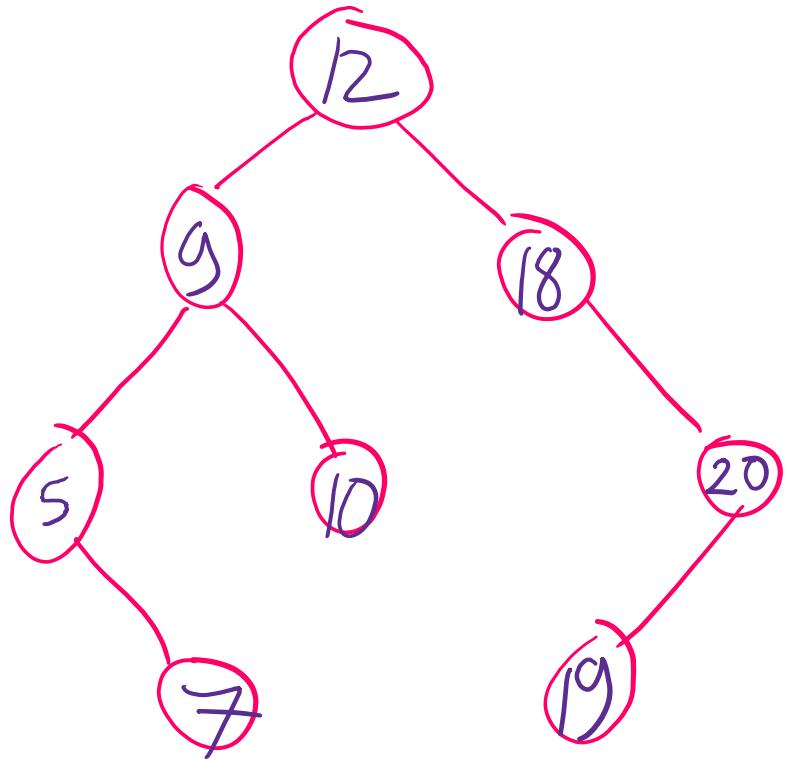
Visit order : 12, 9, 5, 7, 10, 18, 20, 19

Iterative Preorder traversal



Visit order : 12, 9, 5, 7, 10, 18, 20, 19

Iterative Inorder traversal



Visit order :

repeat until stack empty & $x == \text{null}$

$x = \text{leftmost node}$

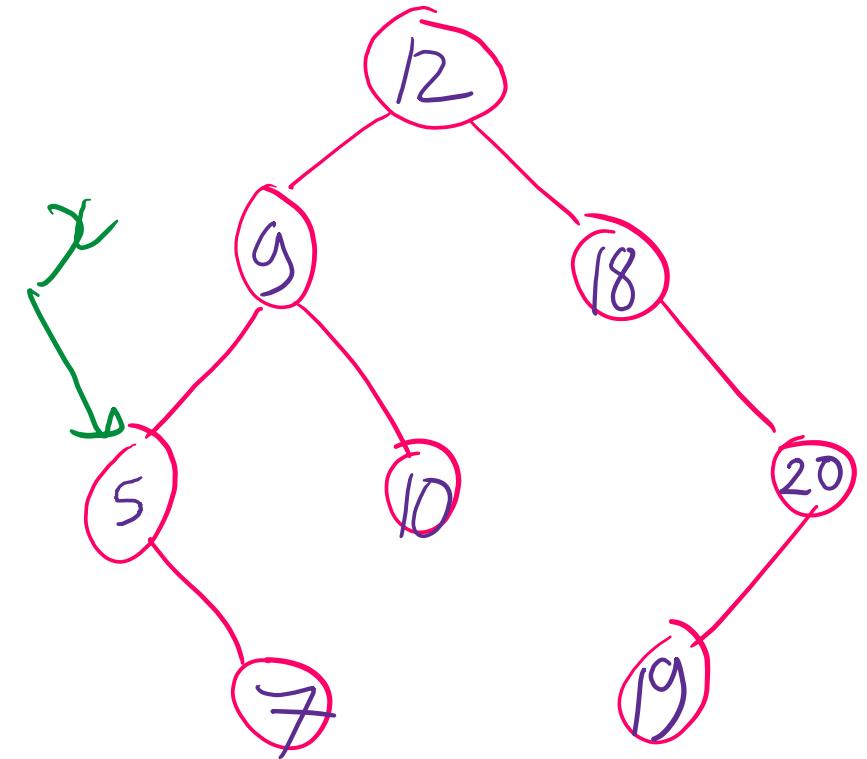
push to stack while moving down

pop & visit

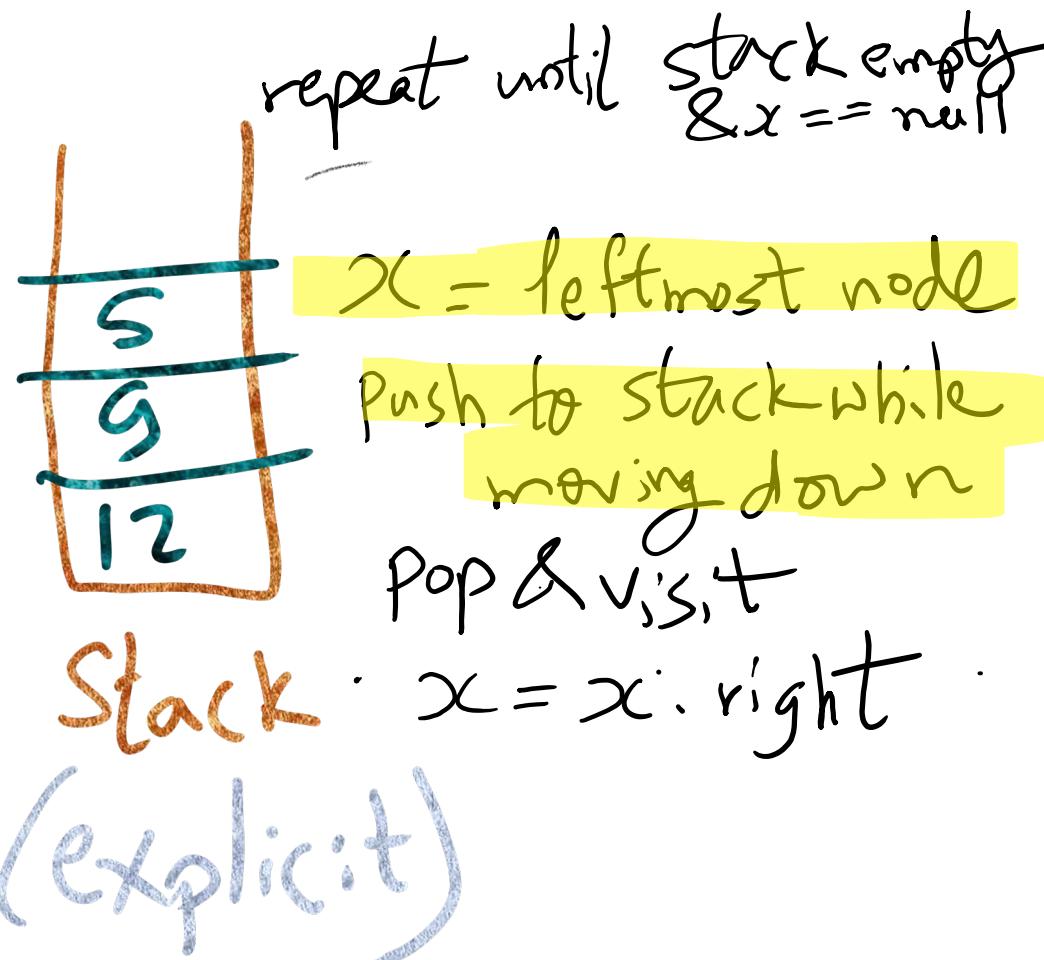
$x = x.\text{right}$

Stack (explicit)

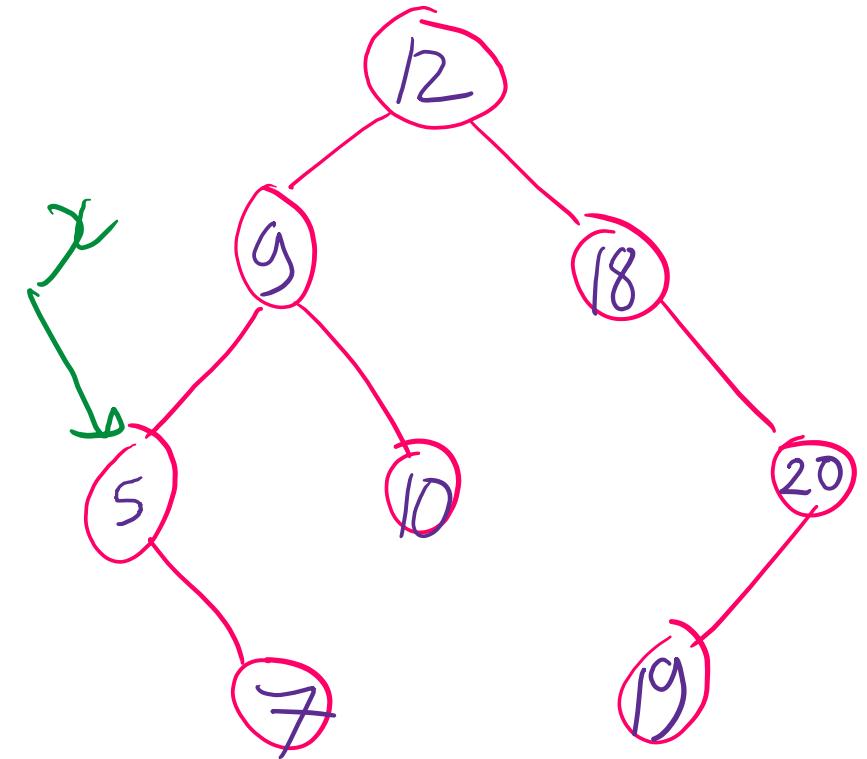
Iterative Inorder traversal



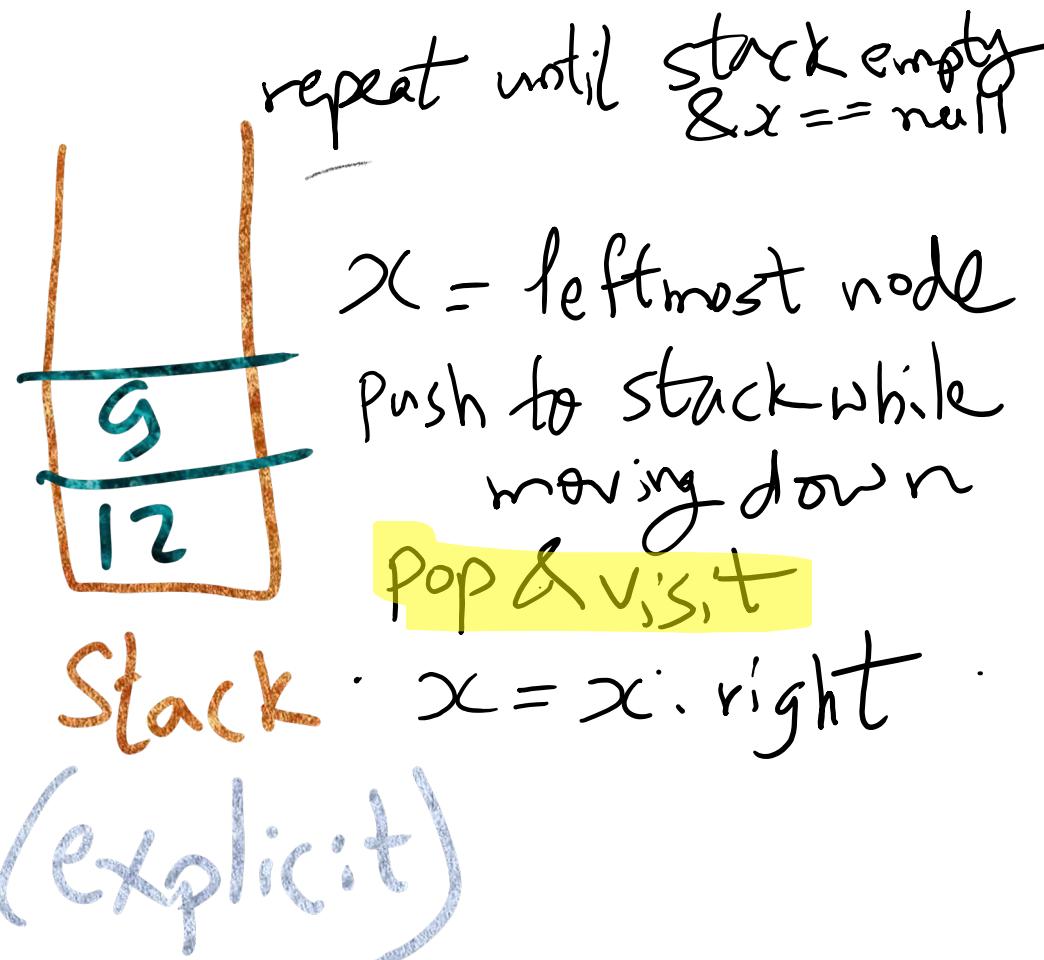
Visit order :



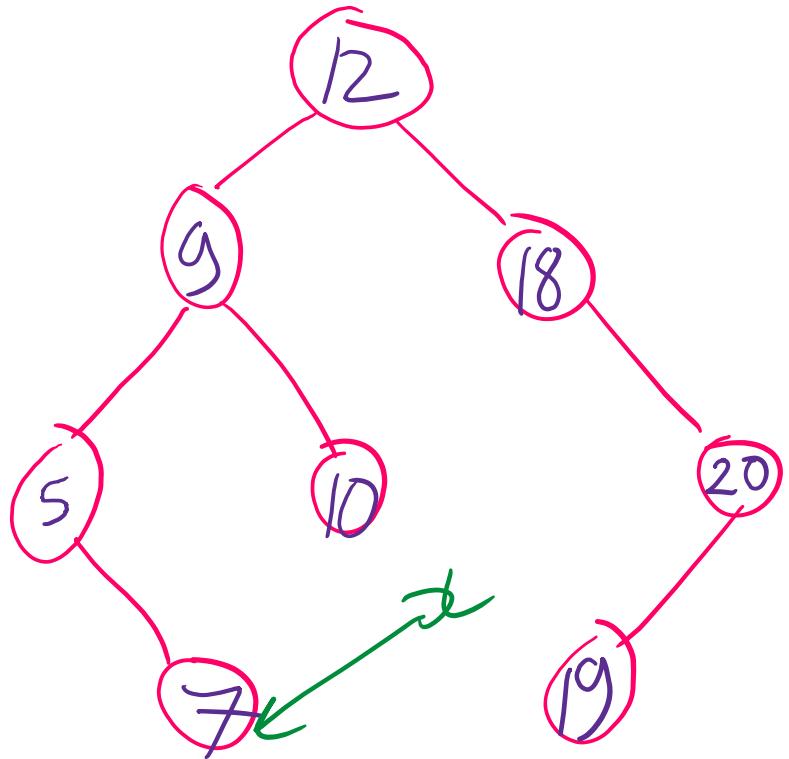
Iterative Inorder traversal



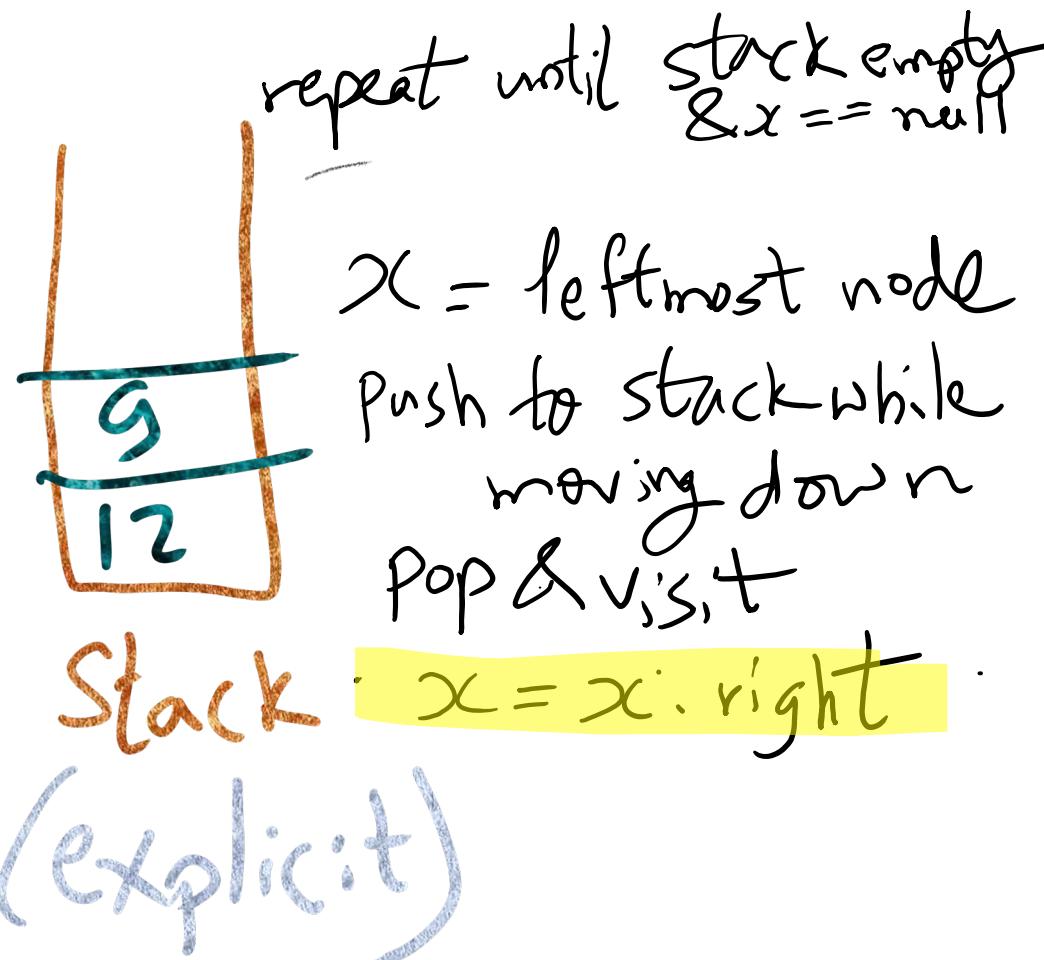
Visit order : 5



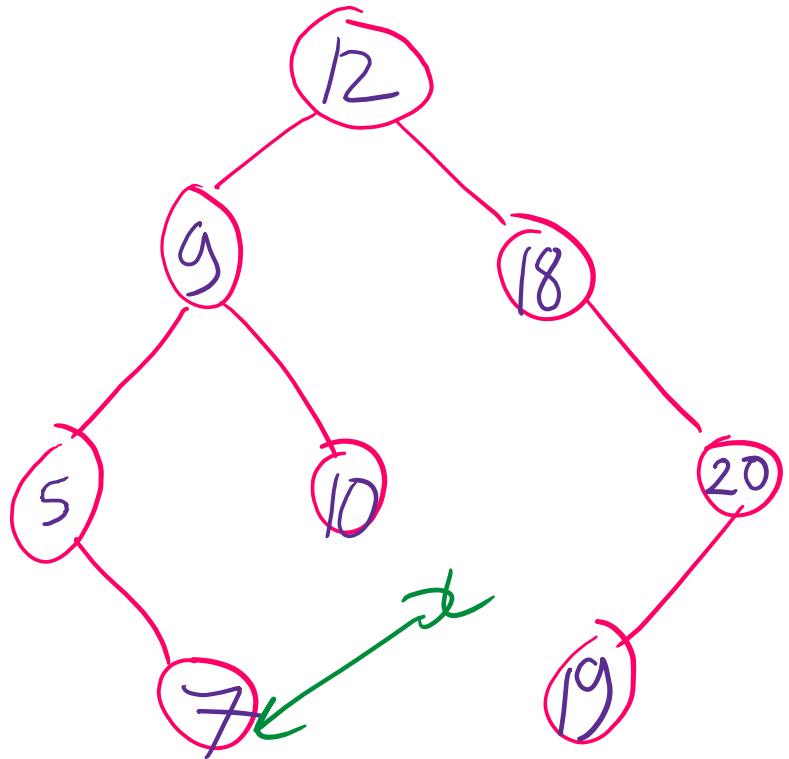
Iterative Inorder traversal



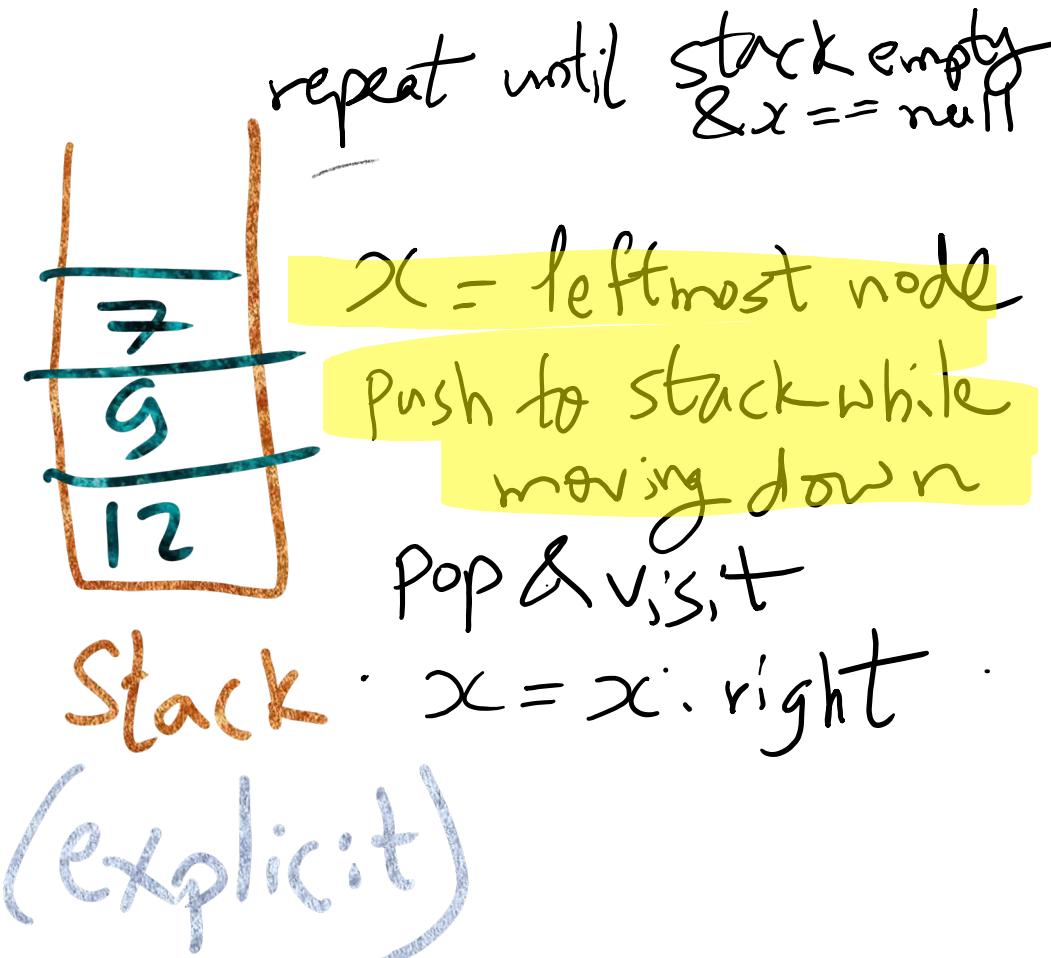
Visit order : 5



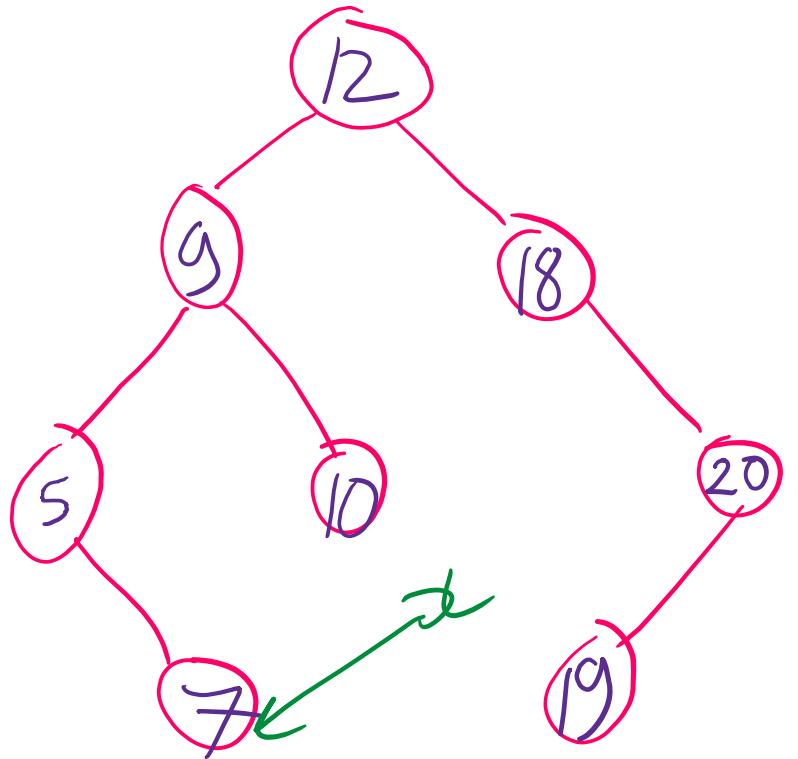
Iterative Inorder traversal



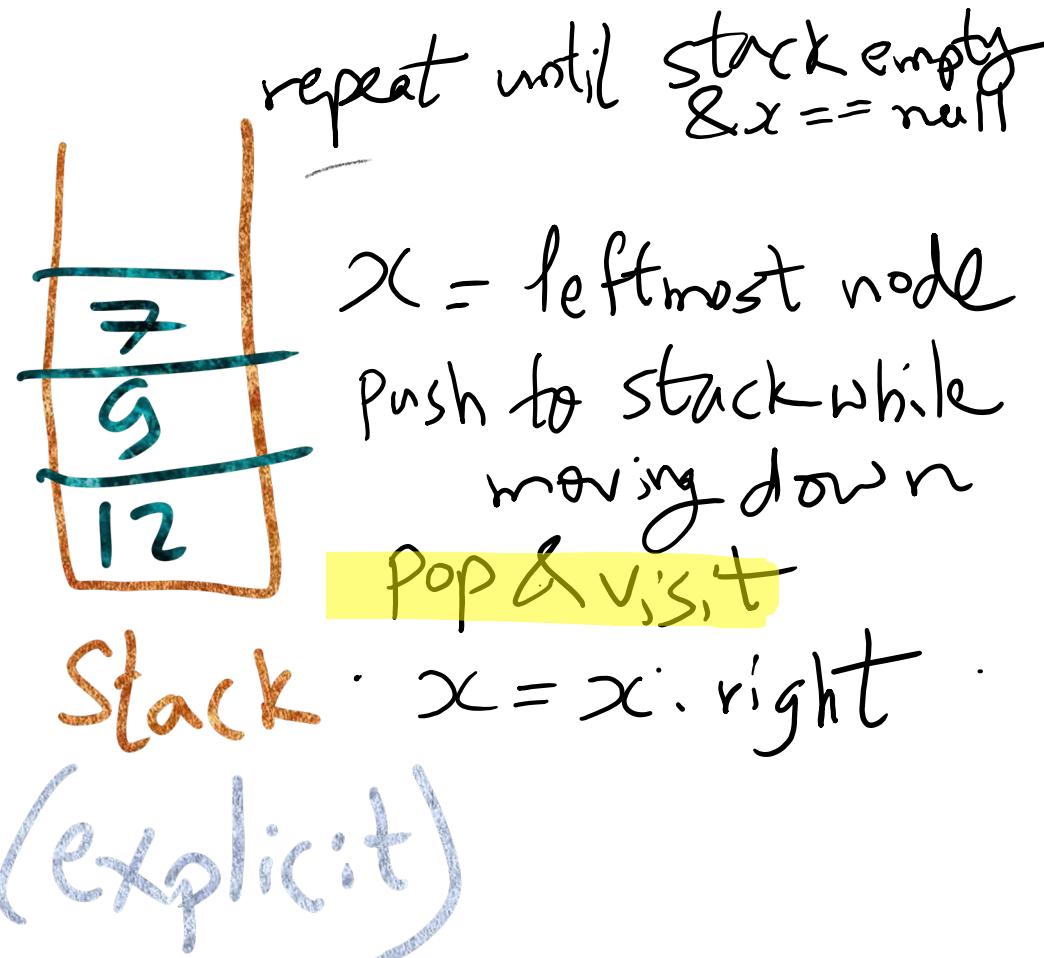
Visit order : 5



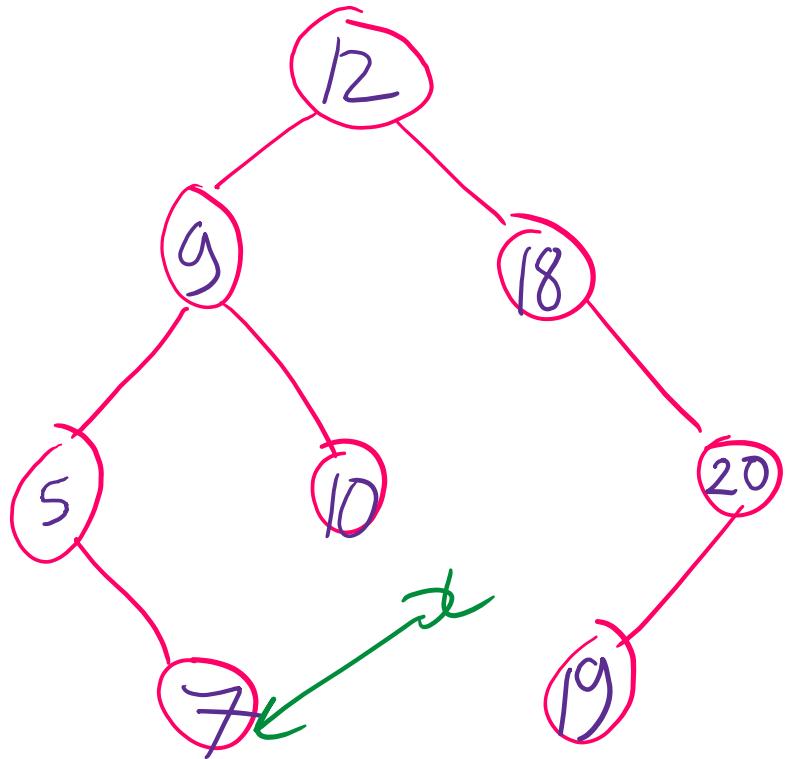
Iterative Inorder traversal



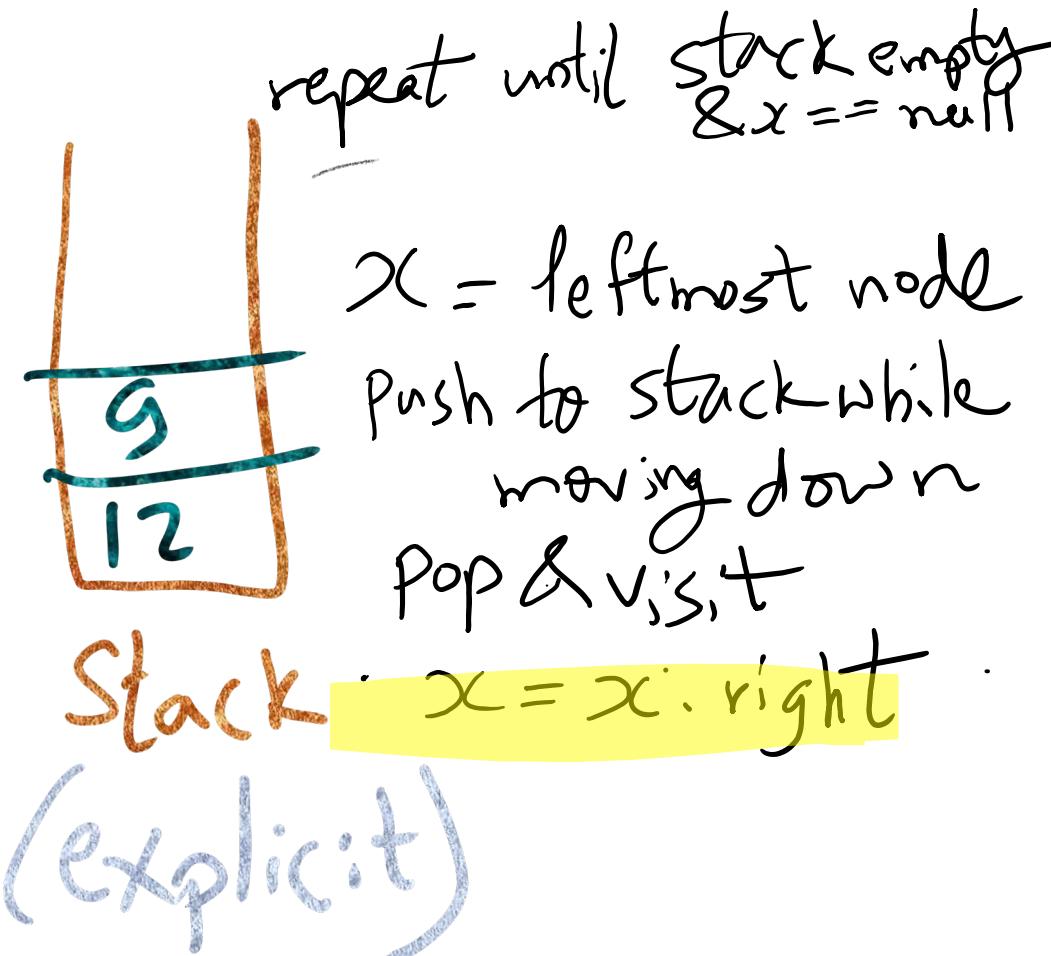
Visit order : 5, 7



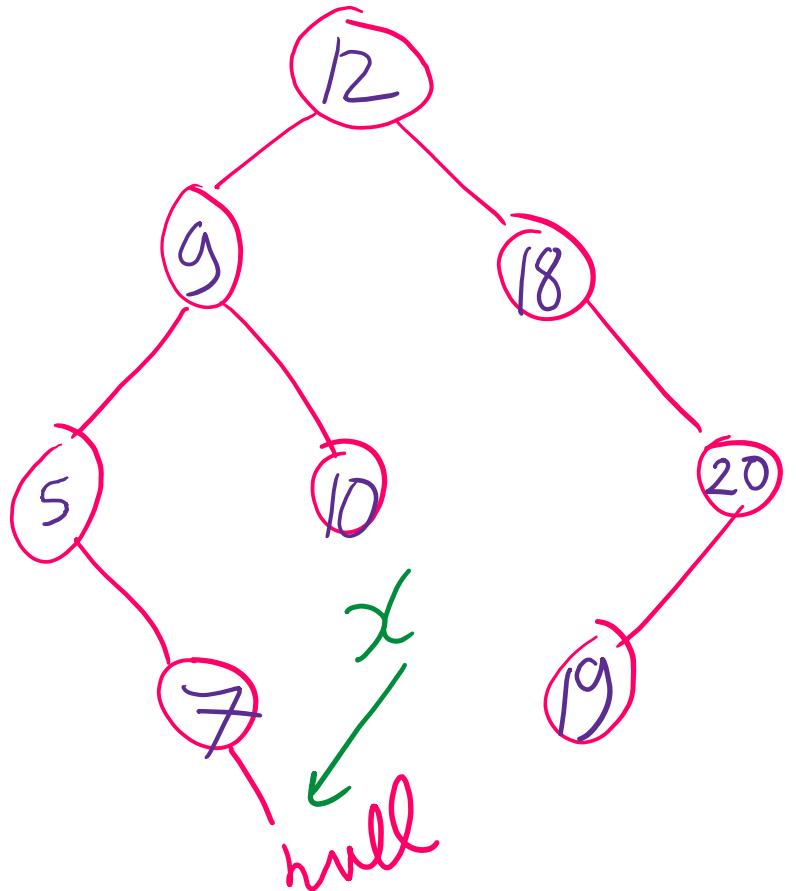
Iterative Inorder traversal



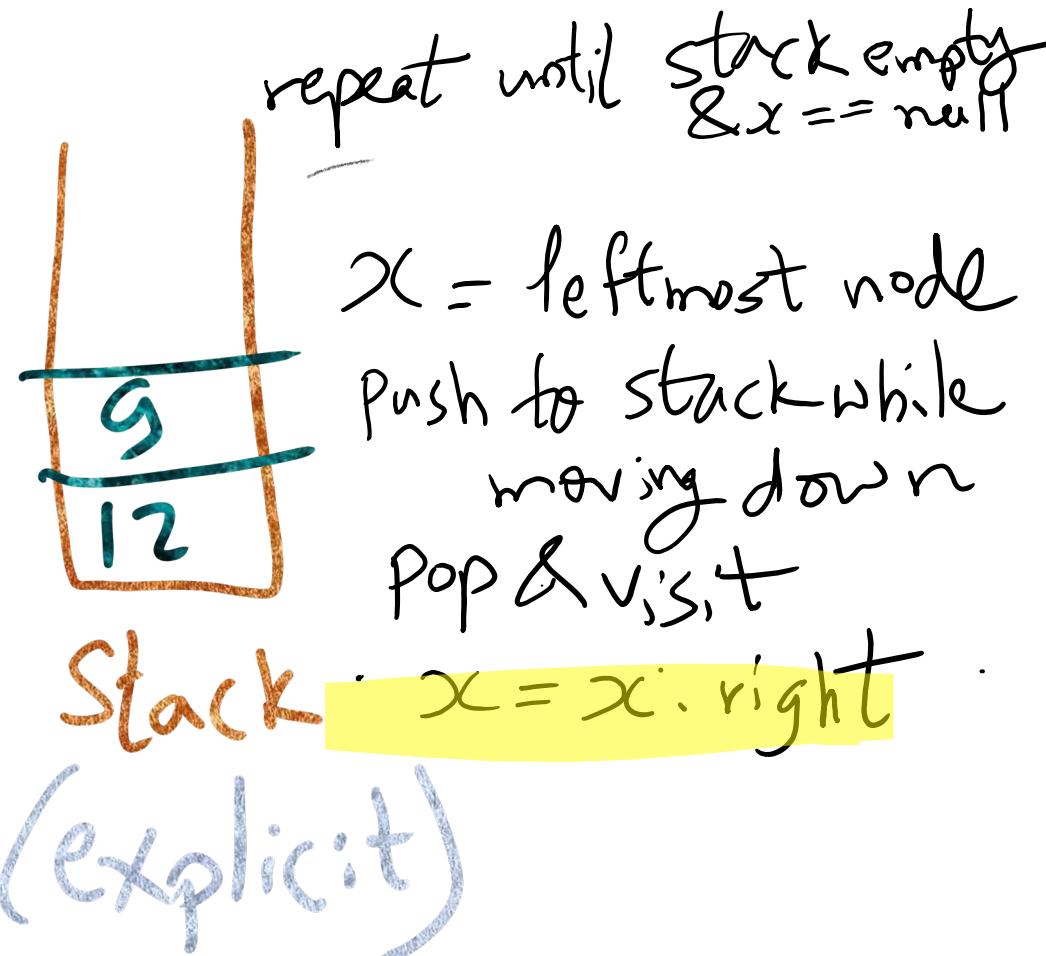
Visit order : 5, 7



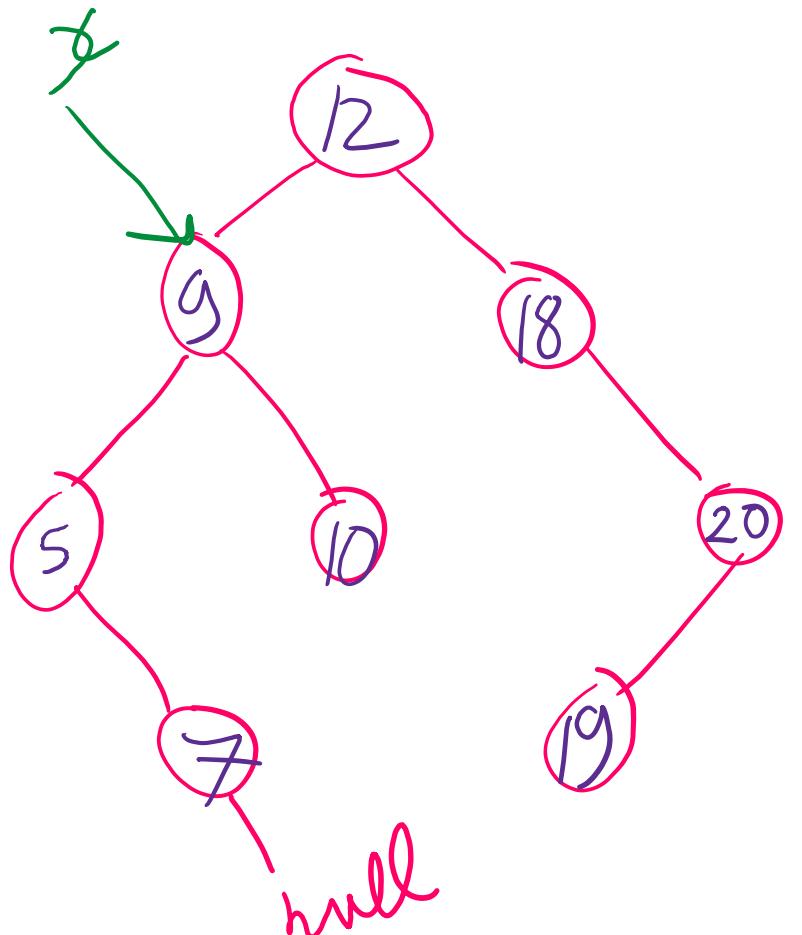
Iterative Inorder traversal



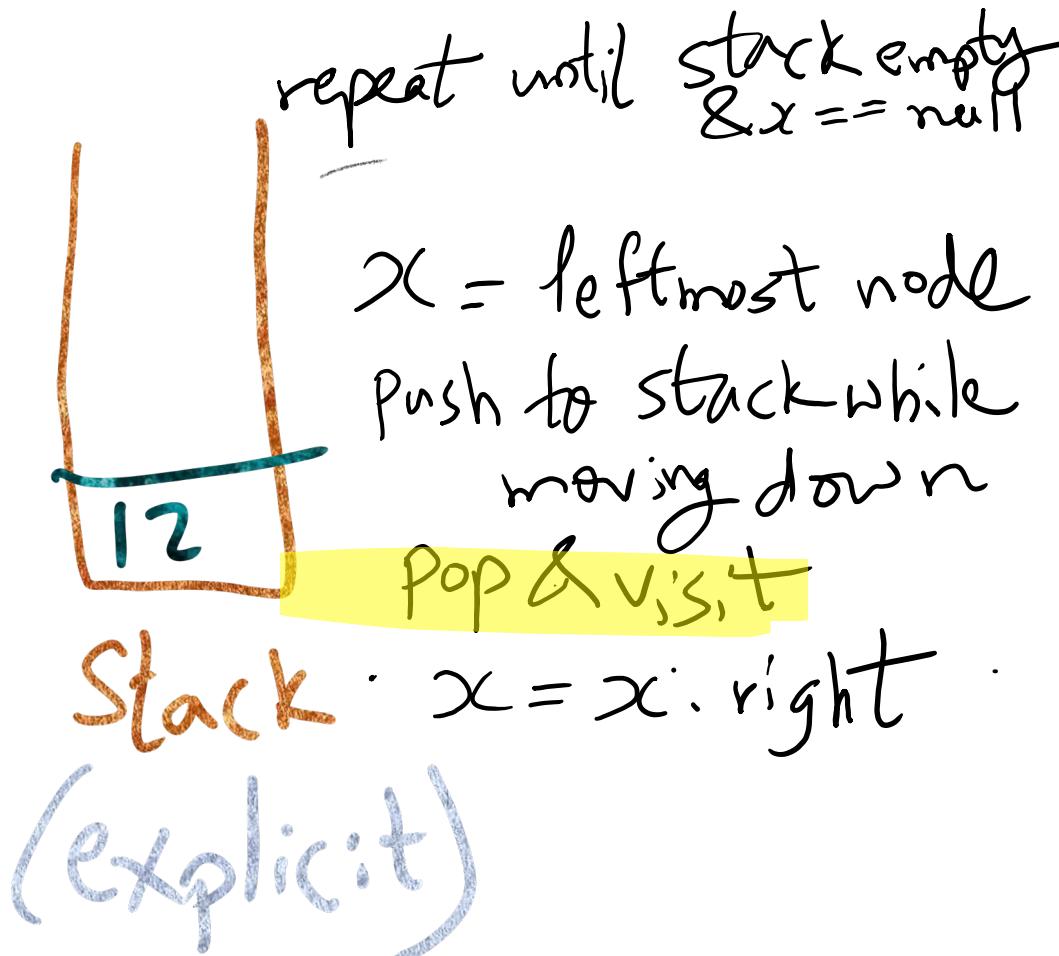
Visit order : 5, 7



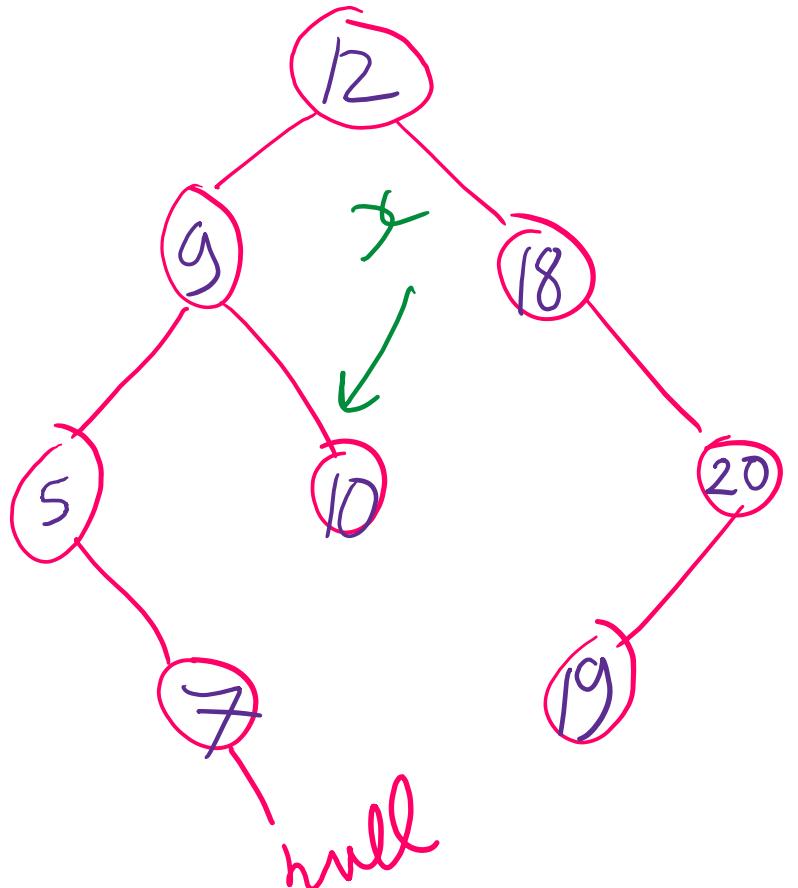
Iterative Inorder traversal



Visit order : 5, 7, 9



Iterative Inorder traversal



Visit order : 5, 7, 9

repeat until stack empty & $x == \text{null}$

$x = \text{leftmost node}$

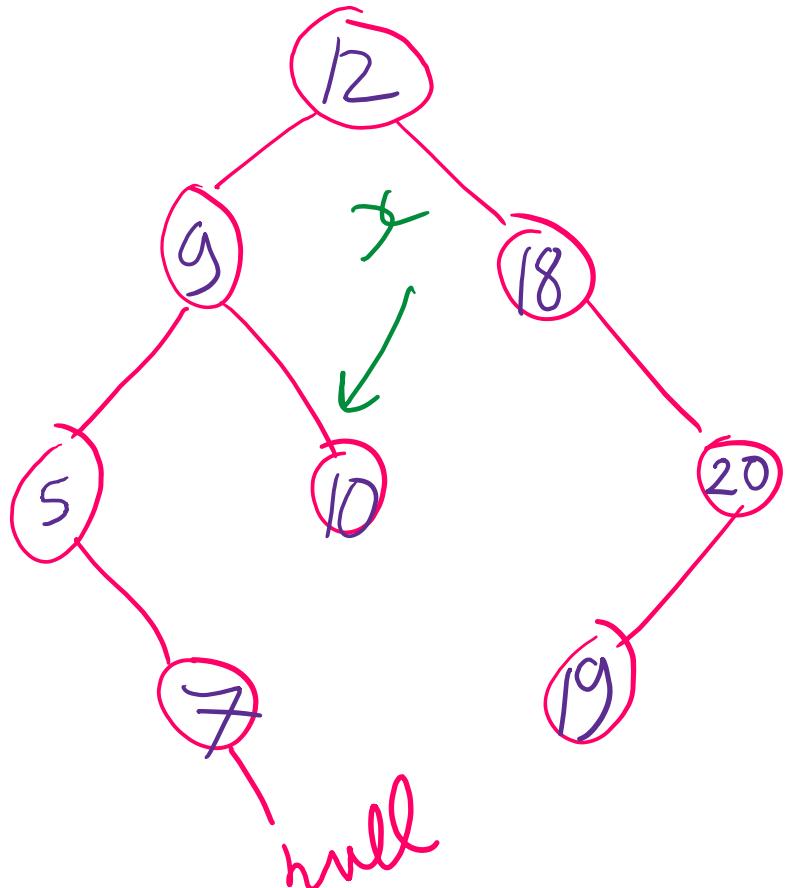
push to stack while moving down

pop & visit

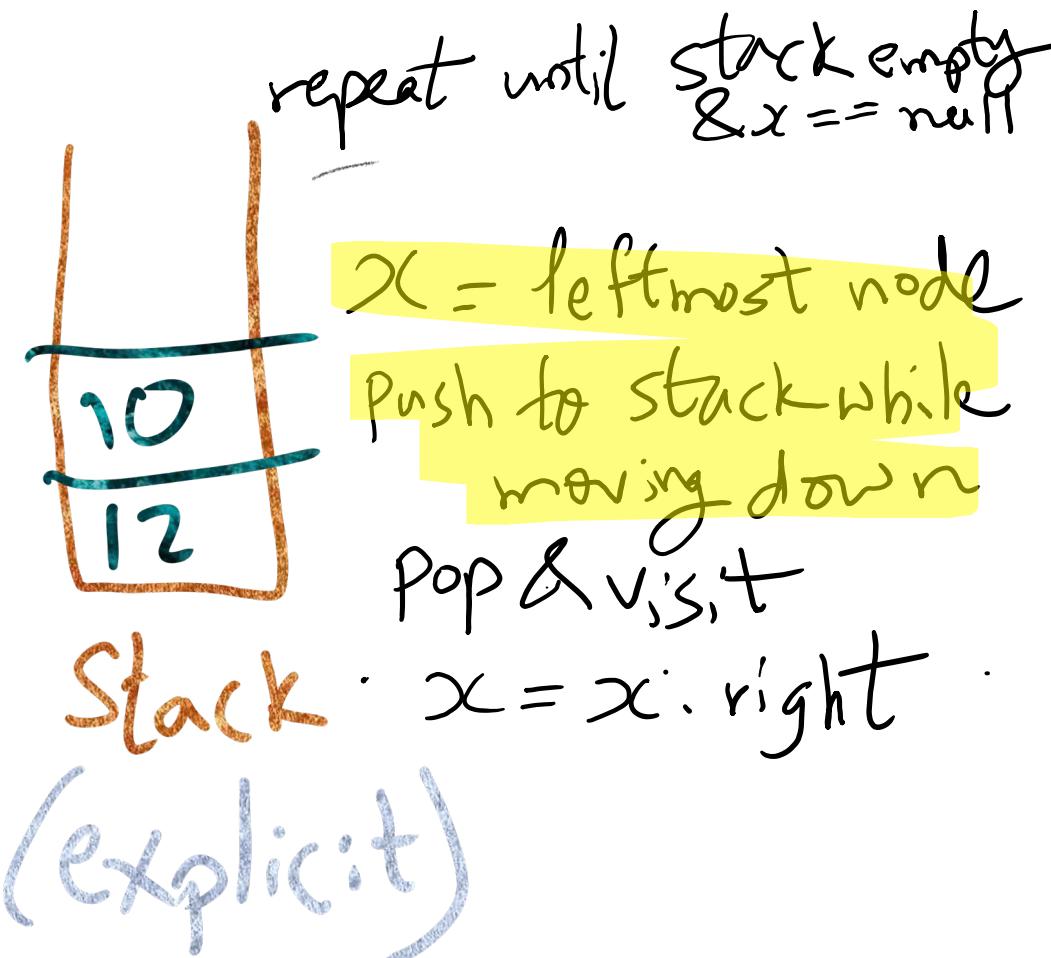
Stack : $x = x.\text{right}$

(explicit)

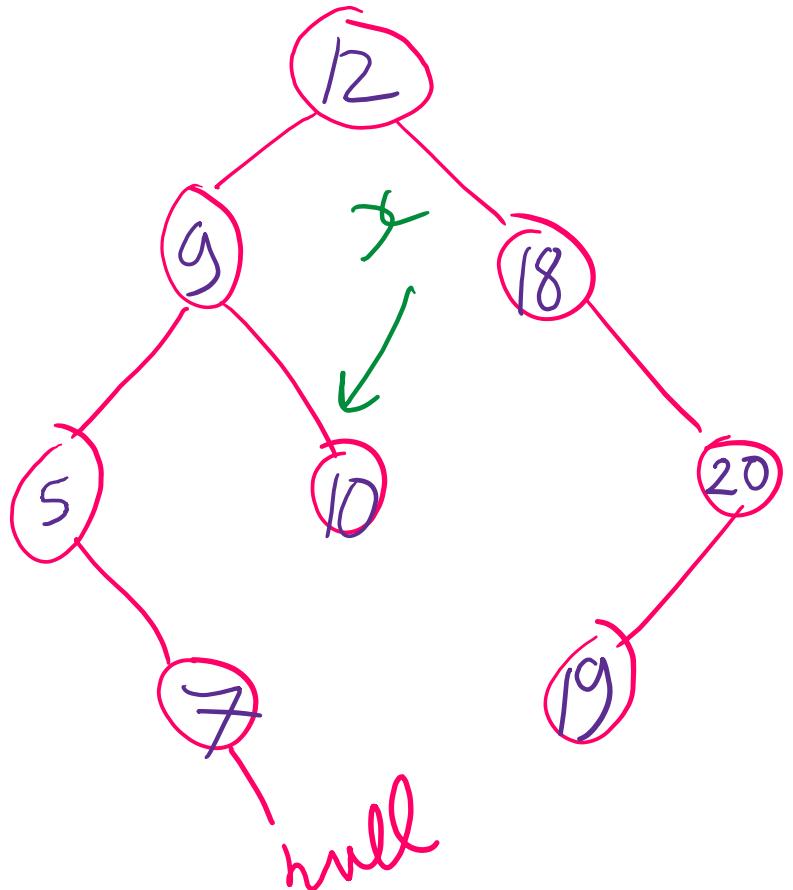
Iterative Inorder traversal



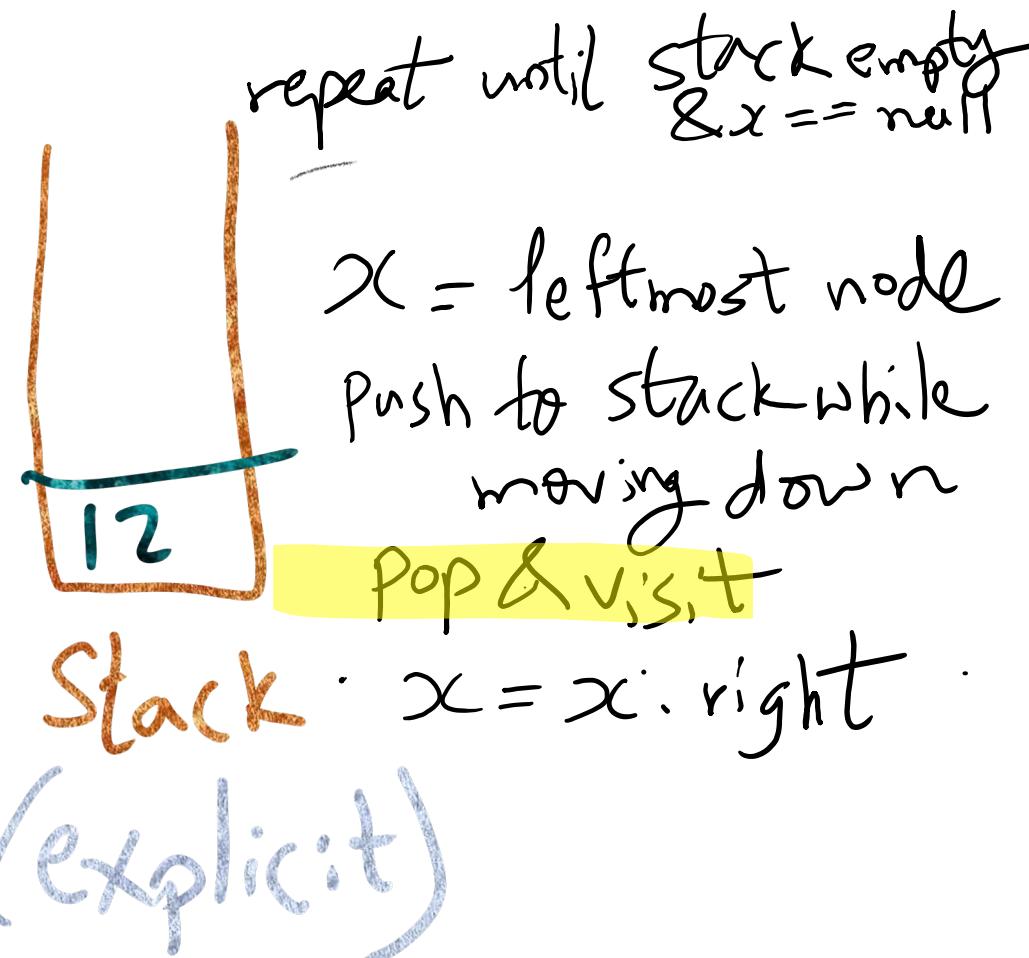
Visit order : 5, 7, 9



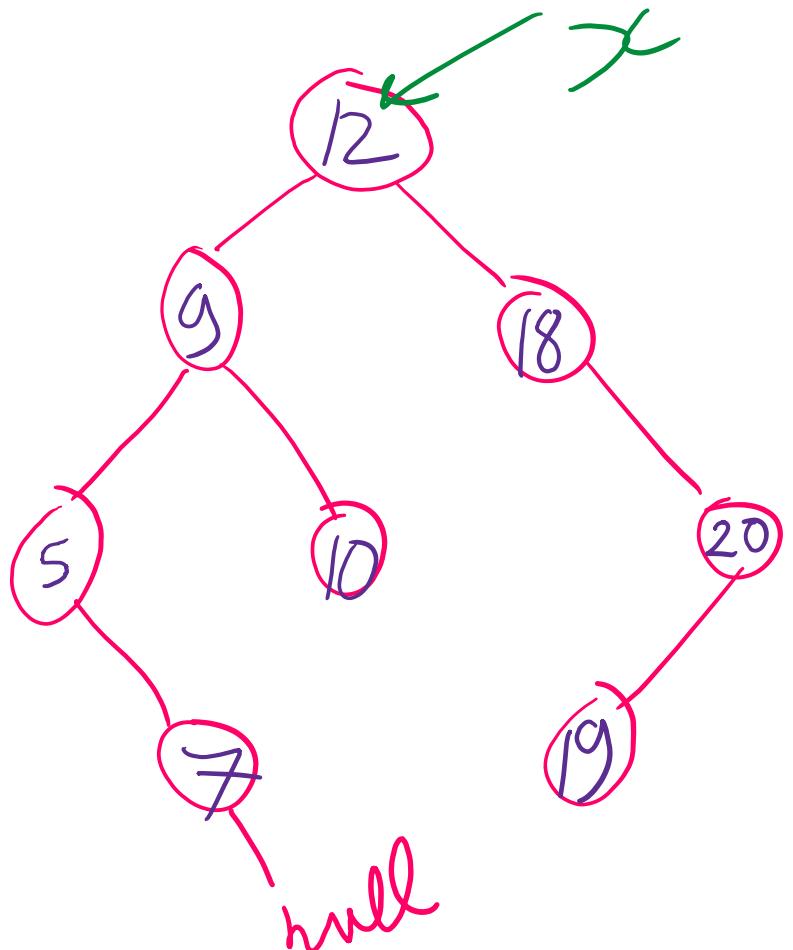
Iterative Inorder traversal



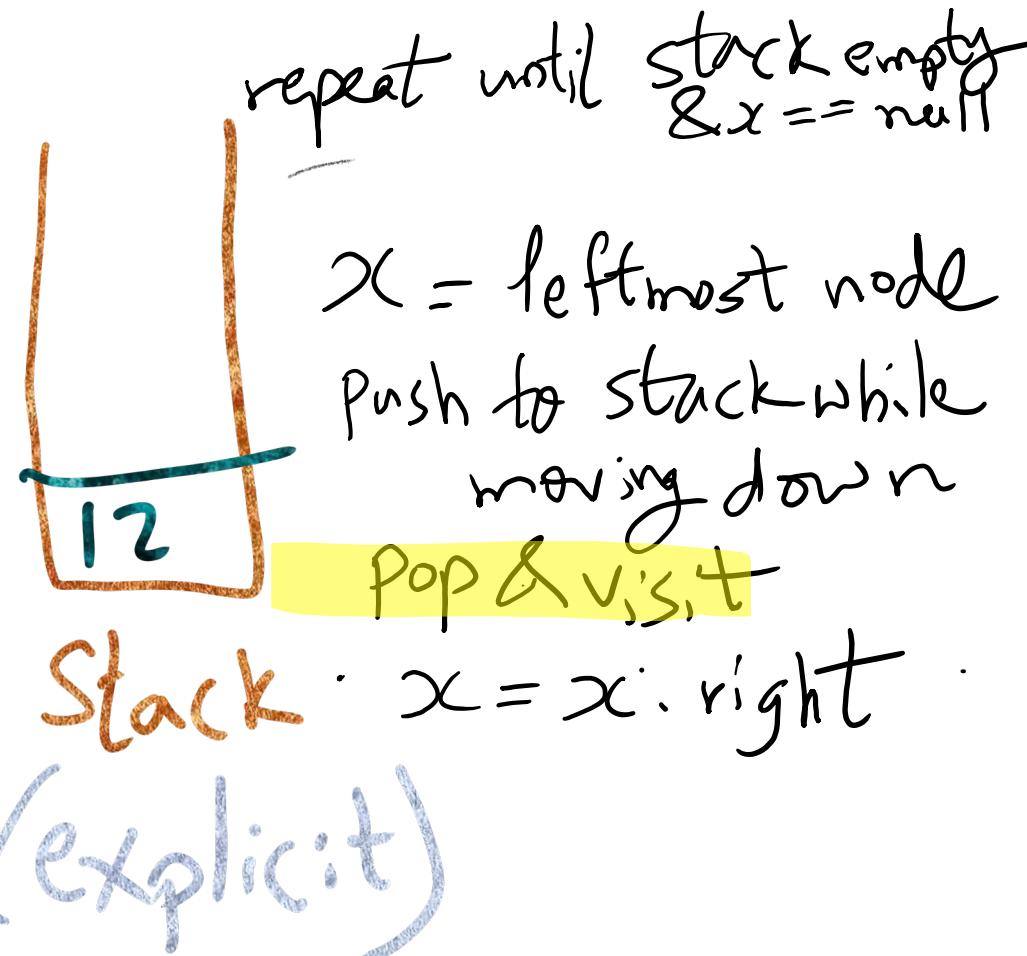
Visit order : 5, 7, 9, 10



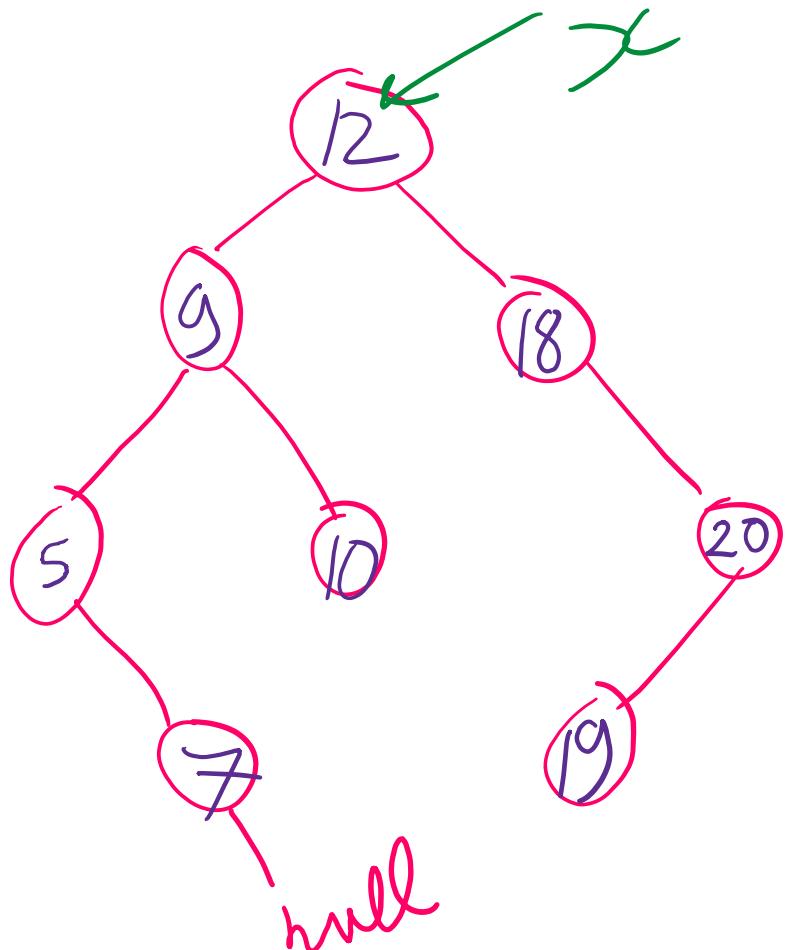
Iterative Inorder traversal



Visit order : 5, 7, 9, 10



Iterative Inorder traversal



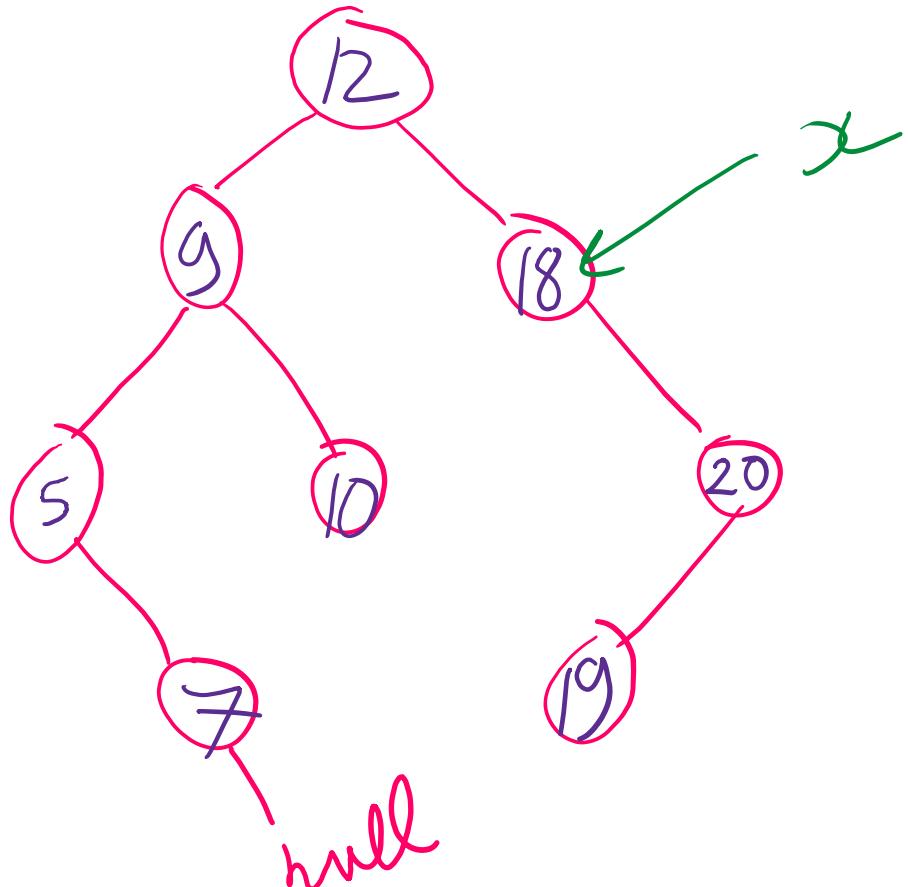
Visit order : 5, 7, 9, 10, 12

repeat until stack empty
& $x == \text{null}$

$x = \text{leftmost node}$
push to stack while moving down
pop & visit

Stack : $x = x.\text{right}$
(explicit)

Iterative Inorder traversal



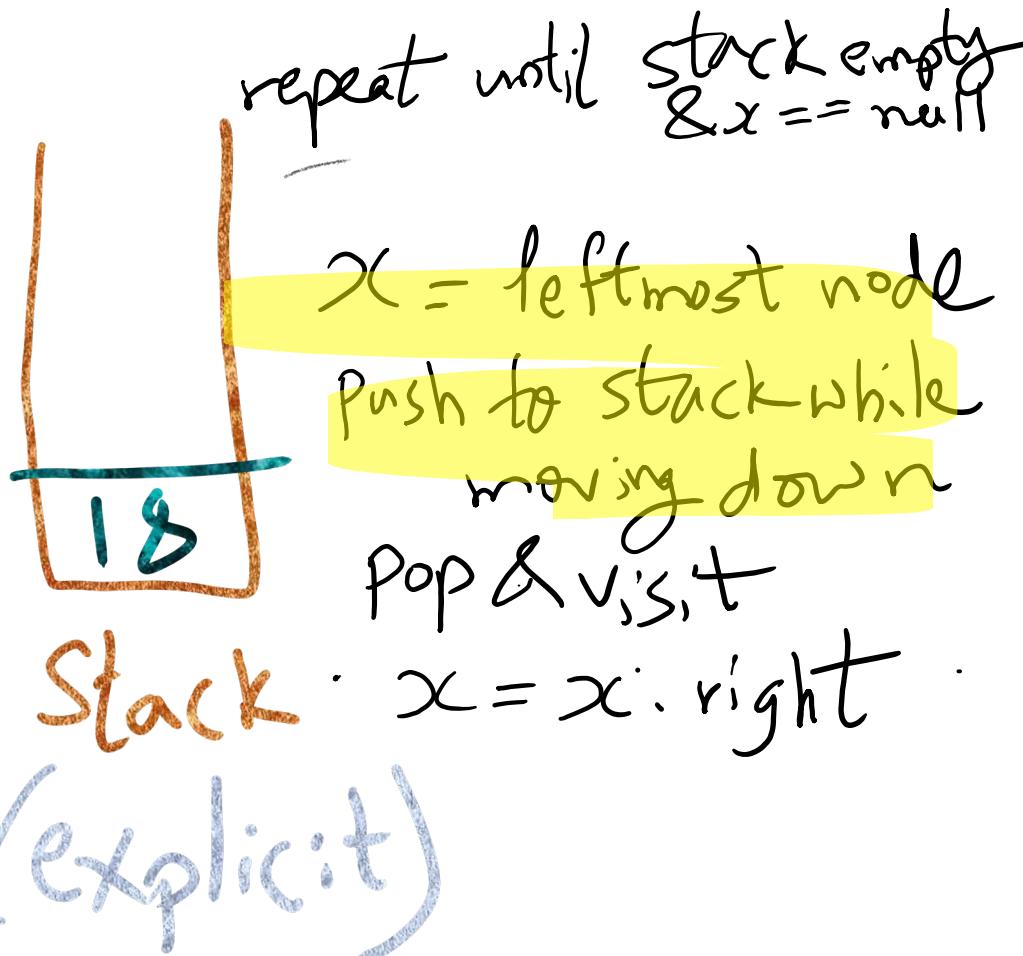
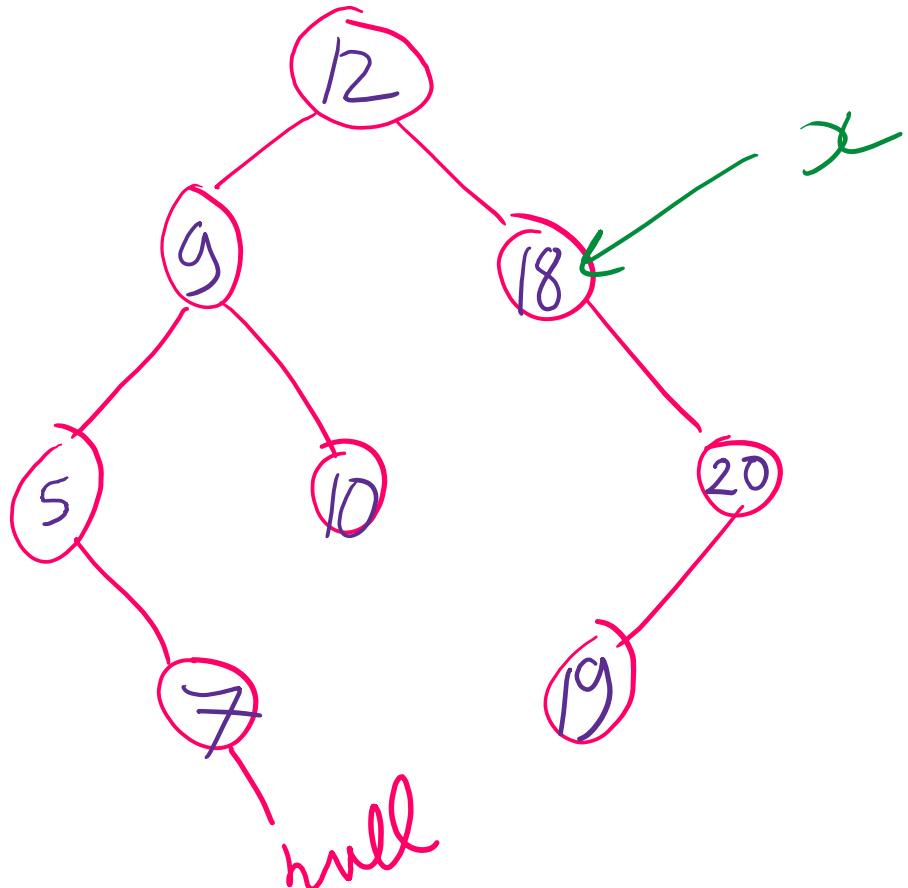
repeat until stack empty
& $x == \text{null}$

$x = \text{leftmost node}$
push to stack while moving down
pop & visit

Stack: $x = x.\text{right}$
(explicit)

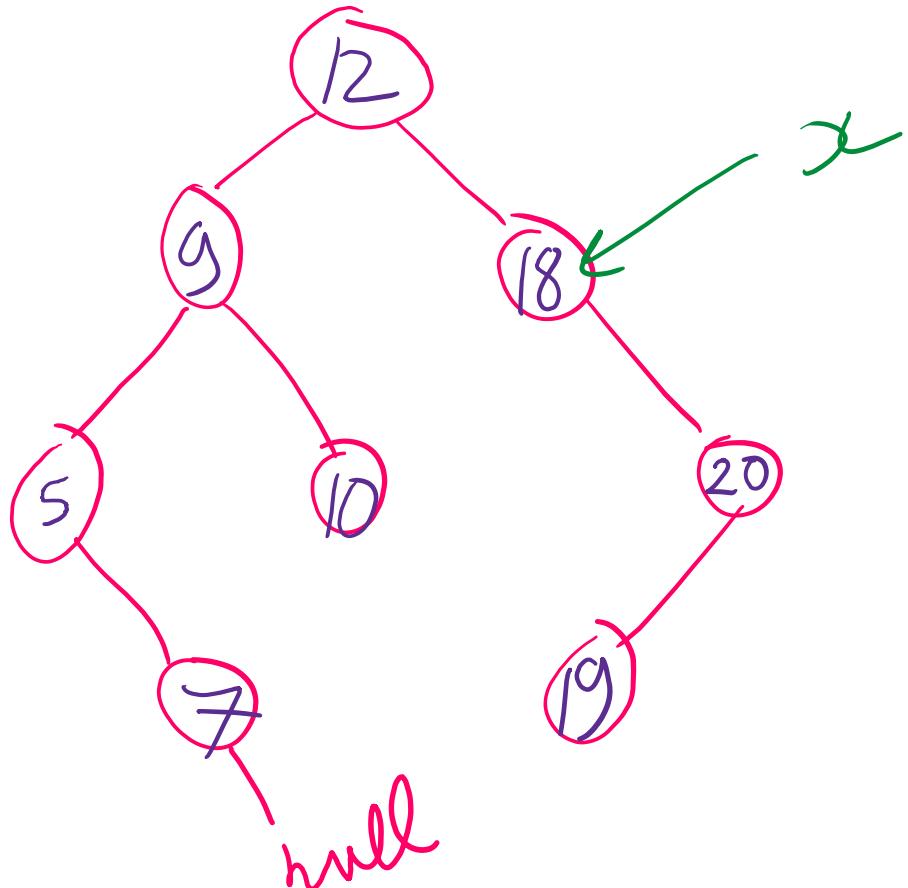
Visit order: 5, 7, 9, 10, 12

Iterative Inorder traversal



Visit order : 5, 7, 9, 10, 12

Iterative Inorder traversal



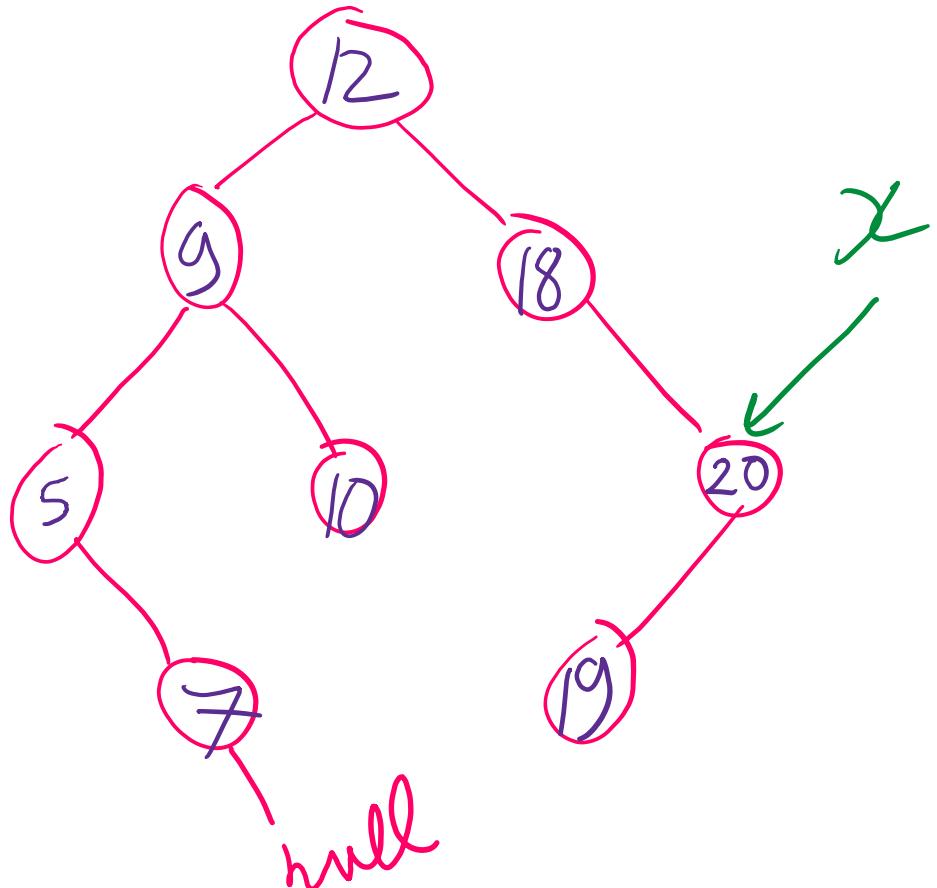
Visit order : 5, 7, 9, 10, 12, 18

repeat until stack empty
 $\& x == \text{null}$

$x = \text{leftmost node}$
push to stack while moving down
pop & visit

Stack : $x = x.\text{right}$
(explicit)

Iterative Inorder traversal



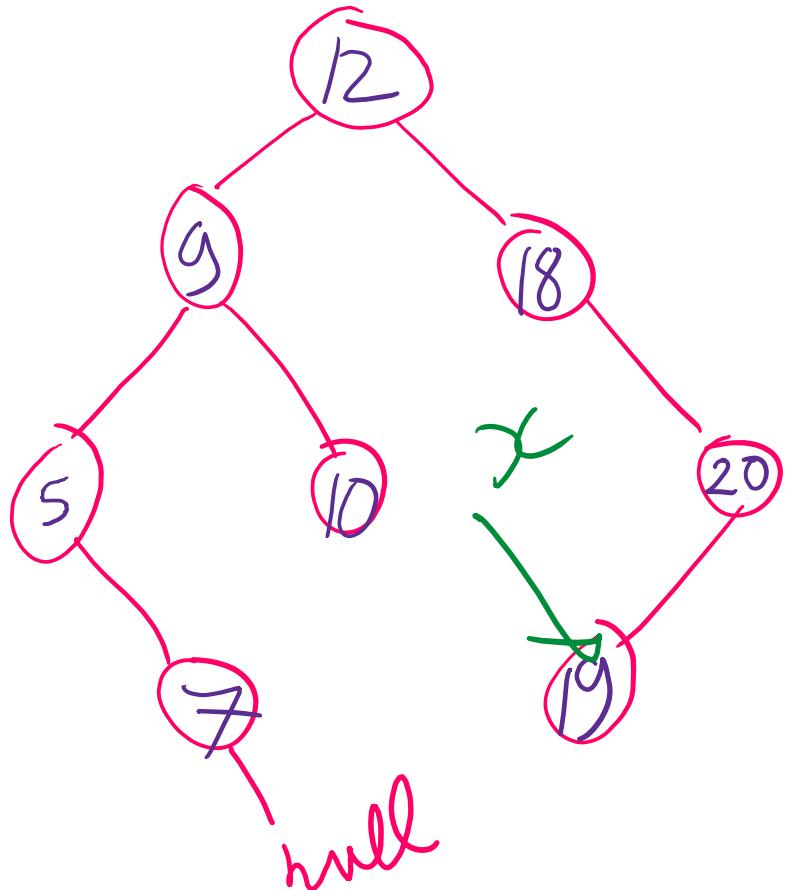
repeat until stack empty
& $x == \text{null}$

$x = \text{leftmost node}$
push to stack while moving down
pop & visit

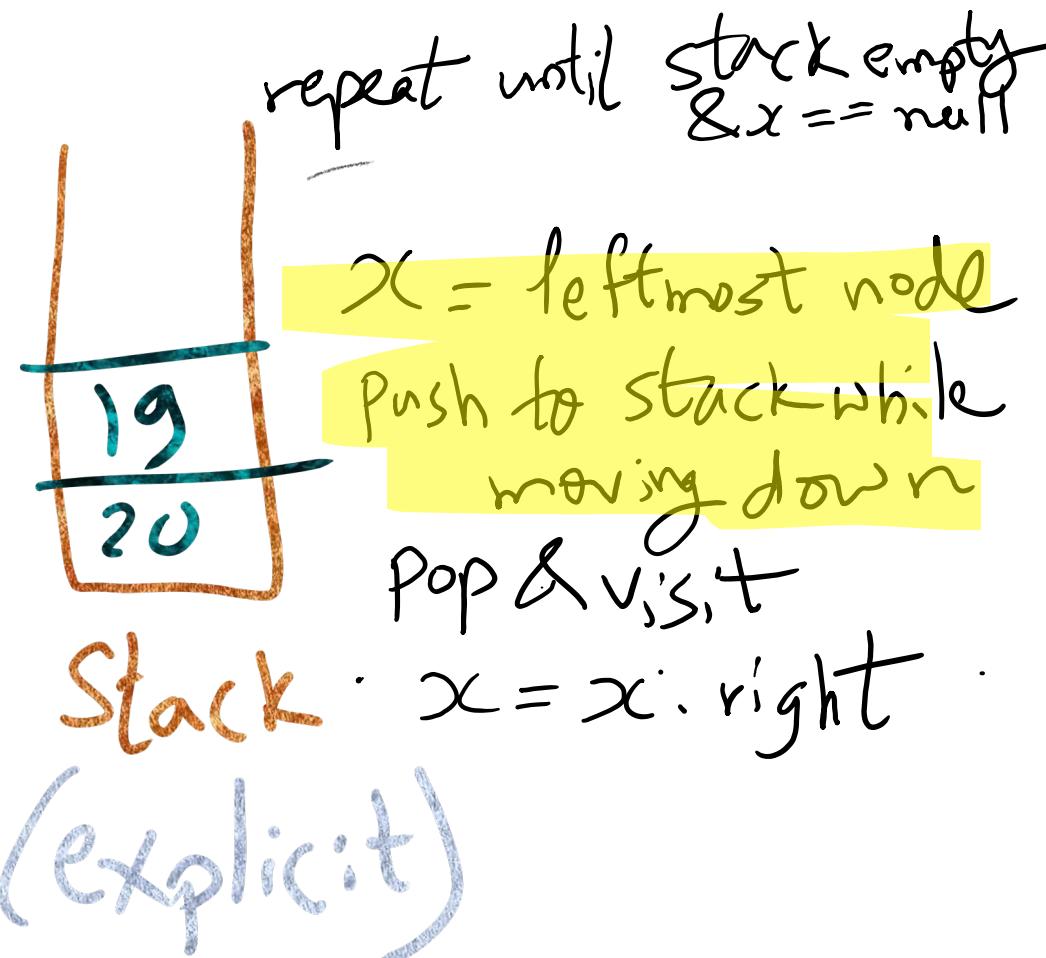
Stack : $x = x.\text{right}$
(explicit)

Visit order : 5, 7, 9, 10, 12, 18

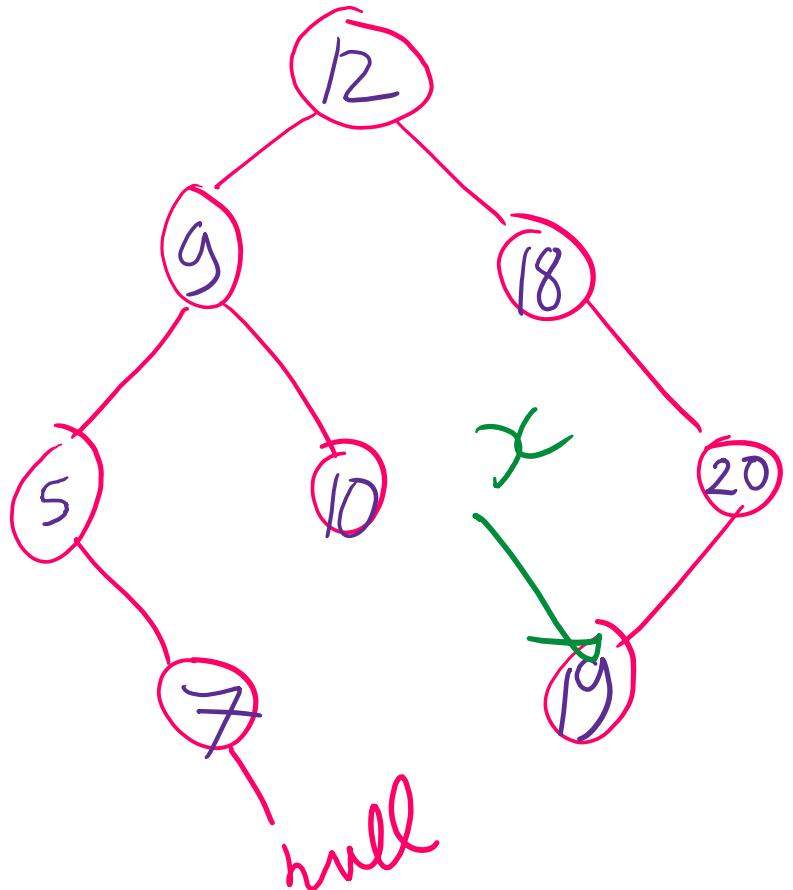
Iterative Inorder traversal



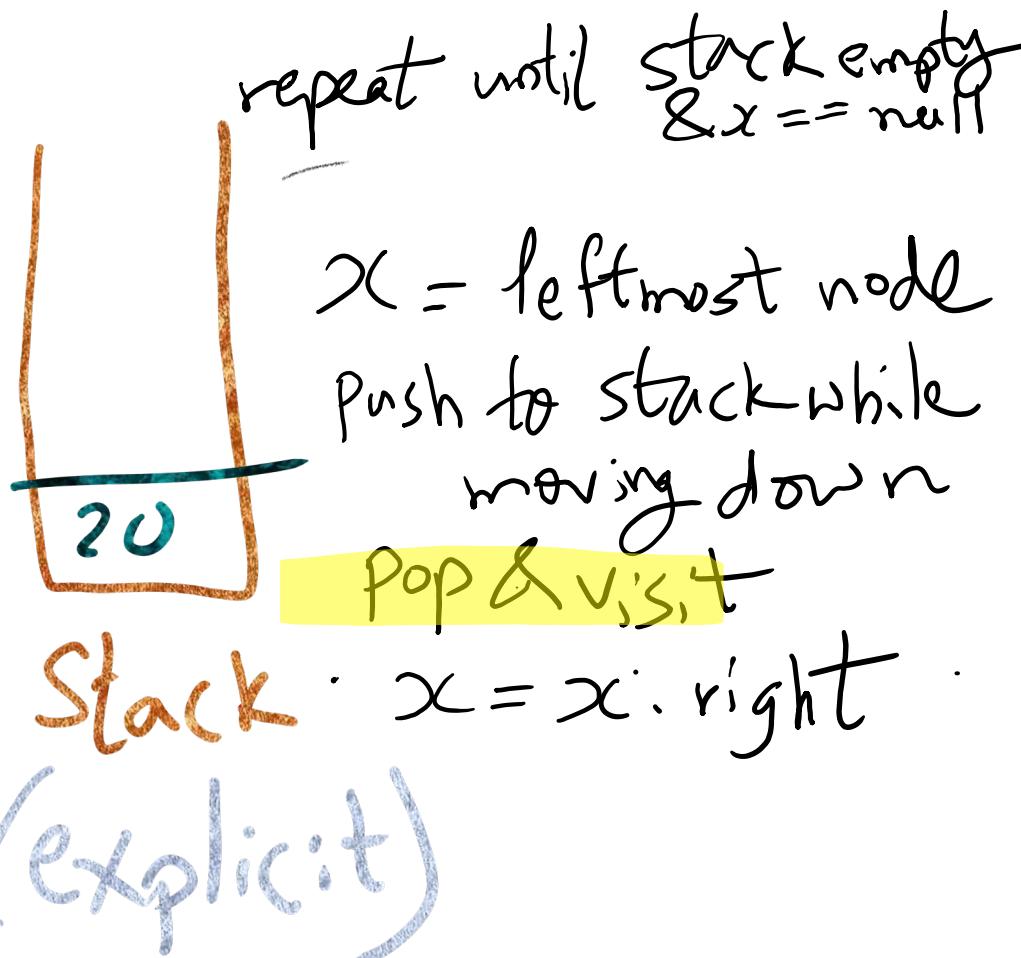
Visit order : 5, 7, 9, 10, 12, 18



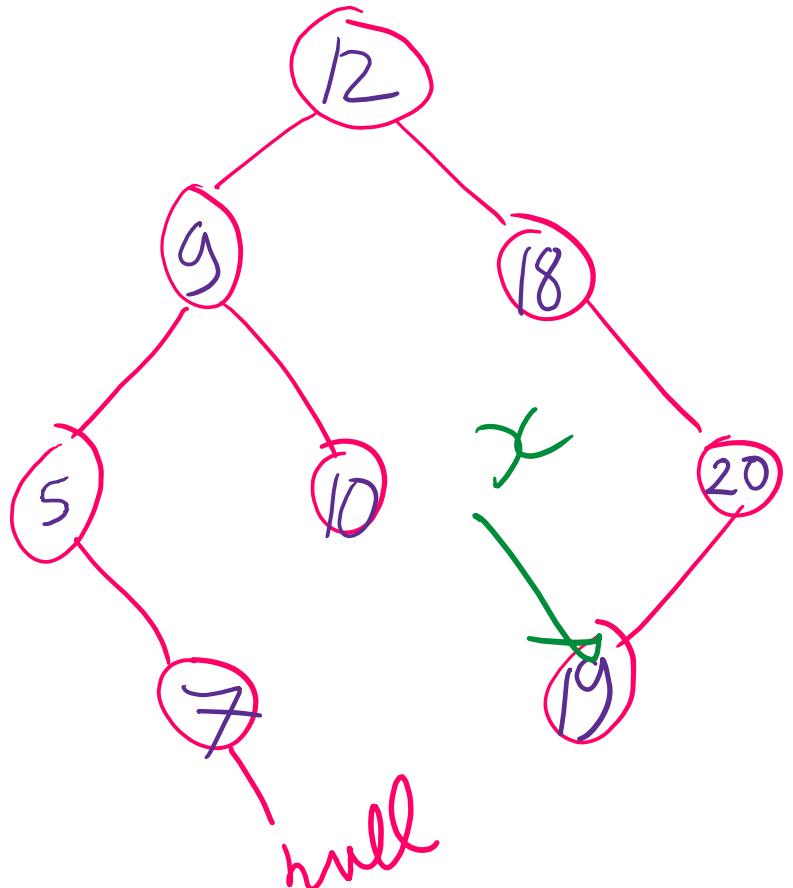
Iterative Inorder traversal



Visit order : 5, 7, 9, 10, 12, 18, 19



Iterative Inorder traversal



Visit order : 5, 7, 9, 10, 12, 18, 19

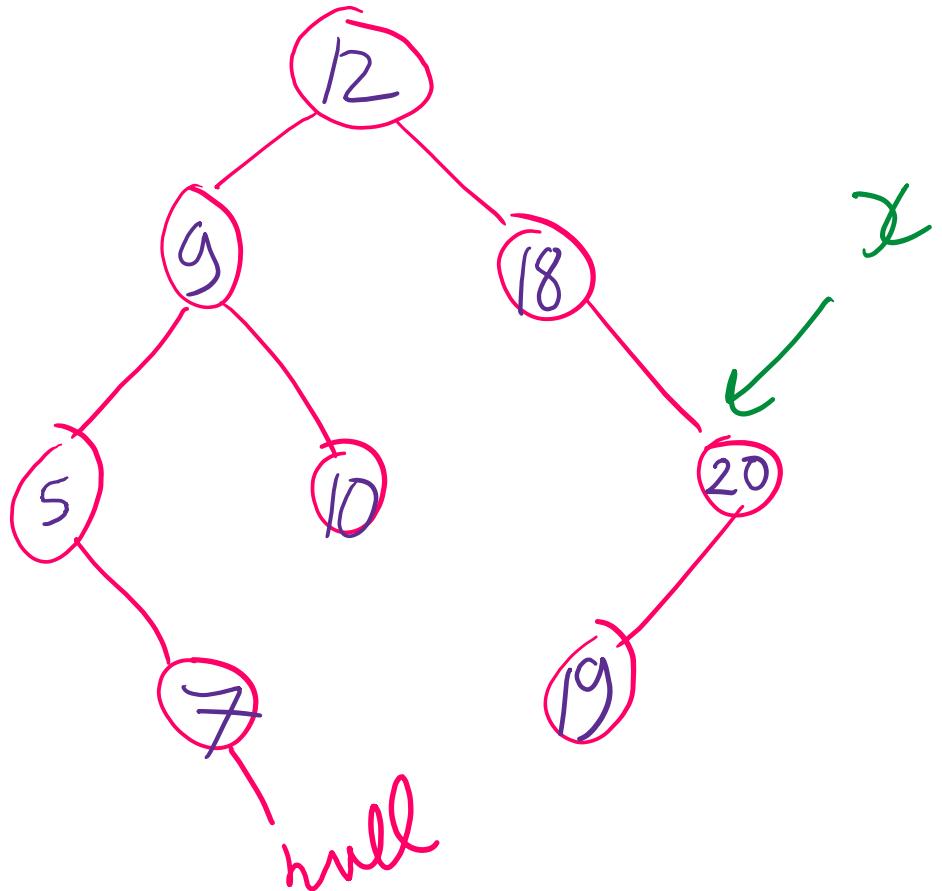
repeat until stack empty
 $\& x == \text{null}$



$x = \text{leftmost node}$
push to stack while moving down
pop & visit
 $x = x.\text{right}$

Stack (explicit)

Iterative Inorder traversal



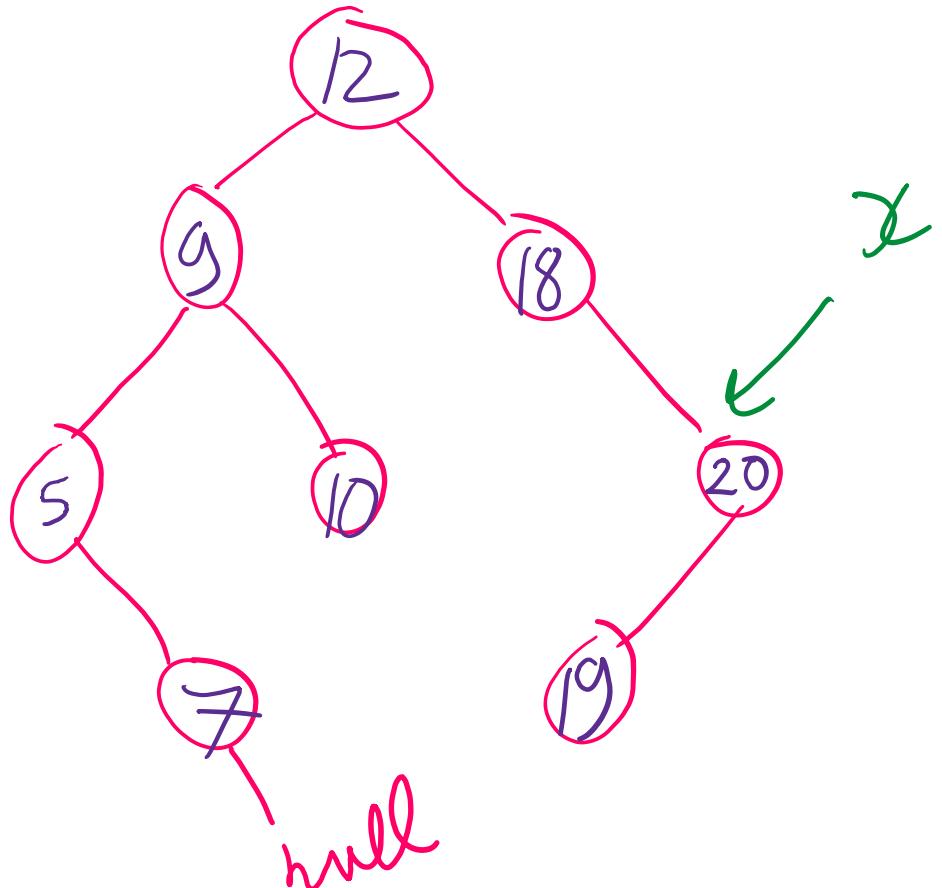
Visit order : 5, 7, 9, 10, 12, 18, 19, 20

repeat until stack empty
 $\& x == \text{null}$

$x = \text{leftmost node}$
push to stack while moving down
pop & visit

Stack : $x = x.\text{right}$
(explicit)

Iterative Inorder traversal



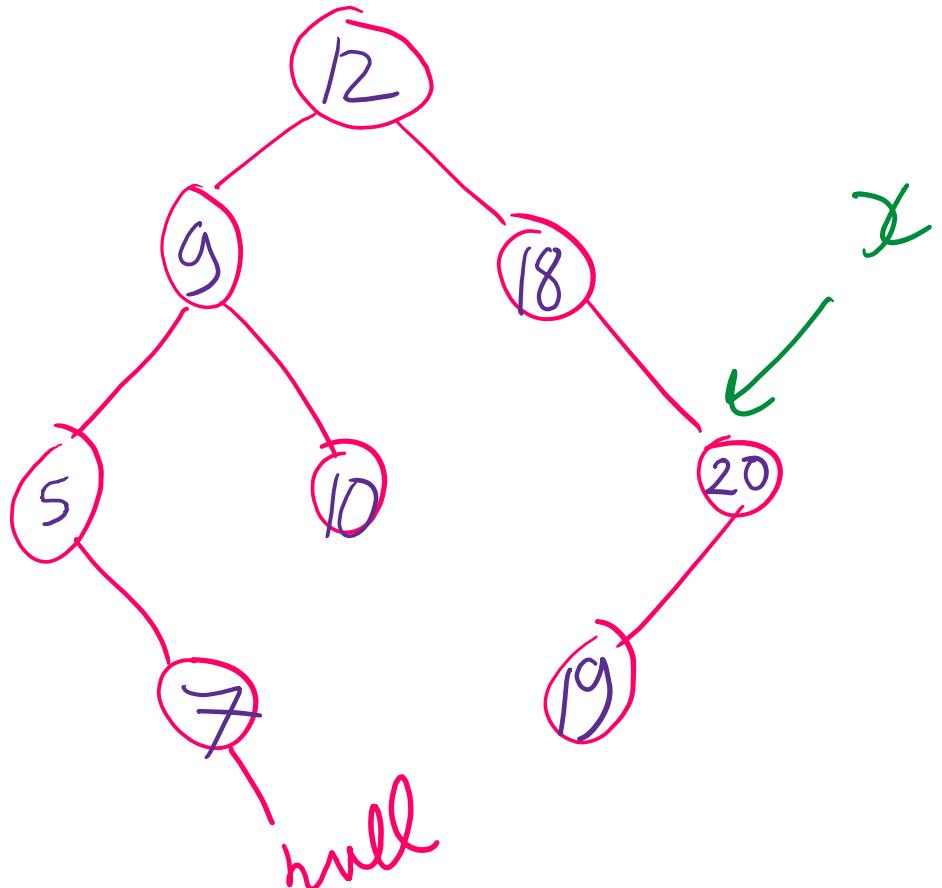
repeat until stack empty
& $x == \text{null}$

$x = \text{leftmost node}$
push to stack while moving down
pop & visit

Stack : $x = x.\text{right}$
(explicit)

Visit order : 5, 7, 9, 10, 12, 18, 19, 20

Iterative Inorder traversal



Visit order : 5, 7, 9, 10, 12, 18, 19, 20

repeat until stack empty & $x == \text{null}$

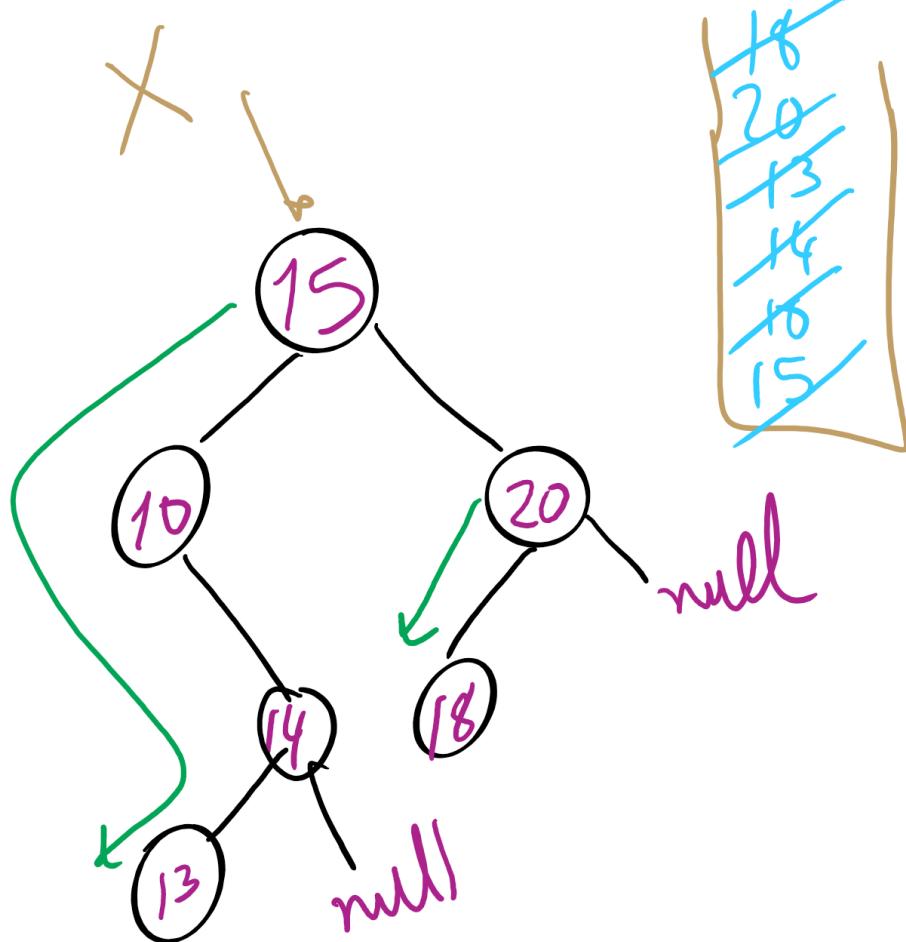
$x = \text{leftmost node}$
push to stack while moving down
pop & visit
 $x = x.\text{right}$

Stack (explicit)

Traversals of a Binary Tree

- Preorder traversal
 - Visit root before we visit root's subtrees
- Inorder traversal
 - Visit root of a binary tree between visiting nodes in root's subtrees.
- Postorder traversal
 - Visit root of a binary tree after visiting nodes in root's subtrees
 - left then right then root

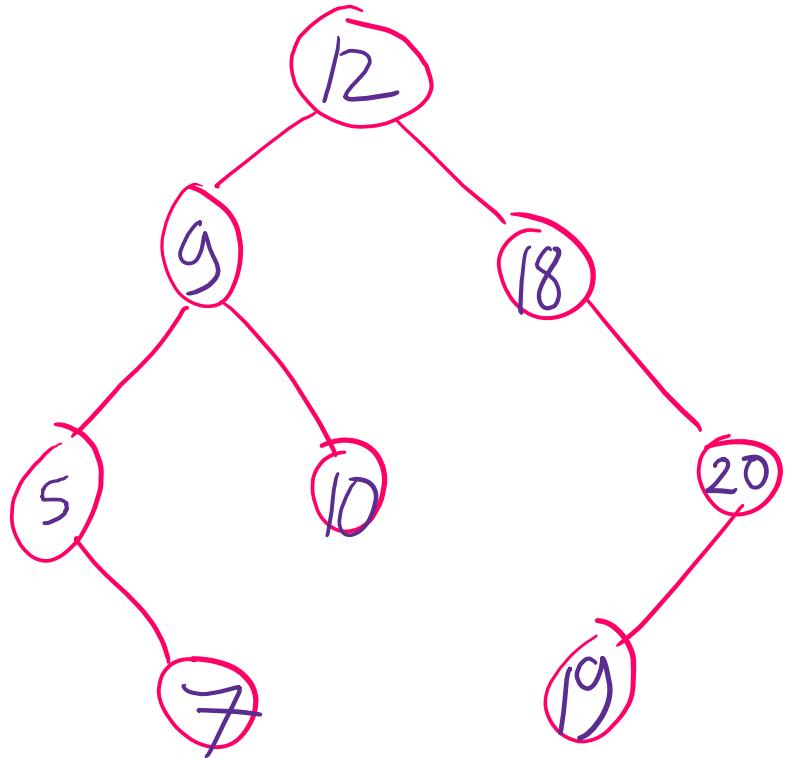
Iterative Post-order Traversal



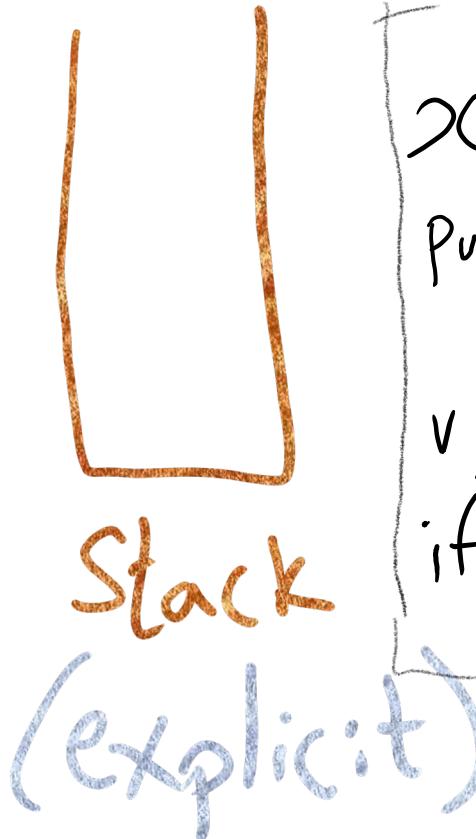
13, 14, 10, 18, 20 15

stack = empty stack
x = root
while(x != null || stack is not empty){
 - find leftmost leaf
 - push nodes as we go down the tree
 - x = pop() & visit x
 - if x is a left child
 x = right sibling
 else
 x = null
y

Iterative Postorder traversal



Visit order :

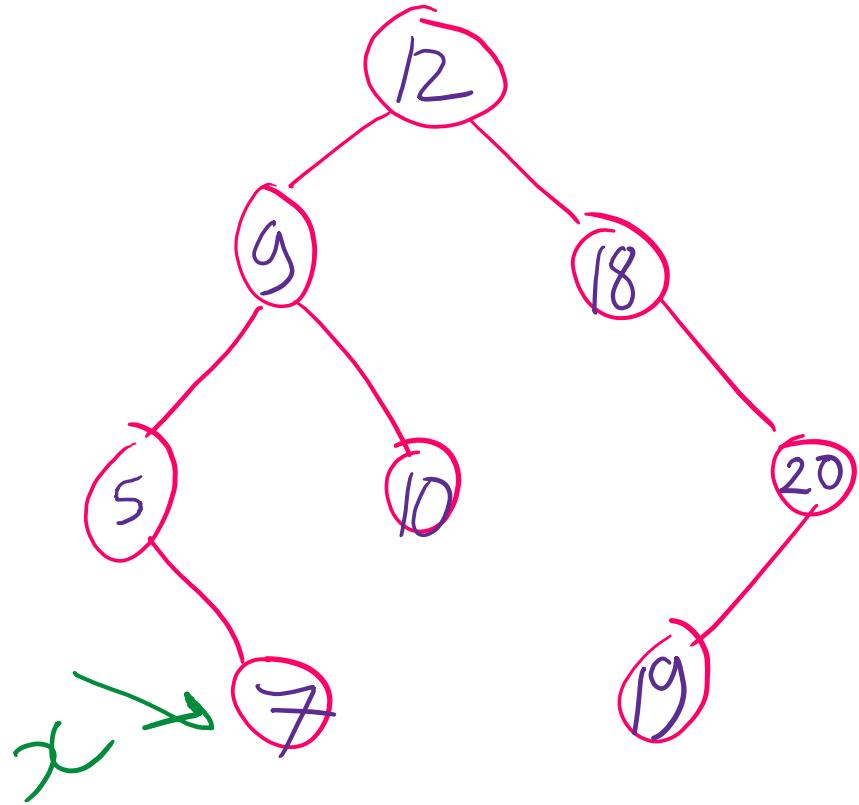


repeat until stack empty
& $x = \text{null}$

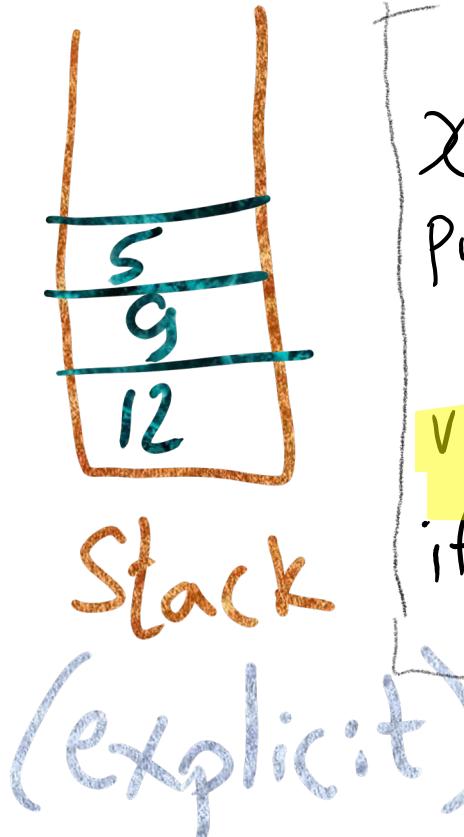
$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

else $x = \text{parent, right}$

Iterative Postorder traversal



Visit order : 7



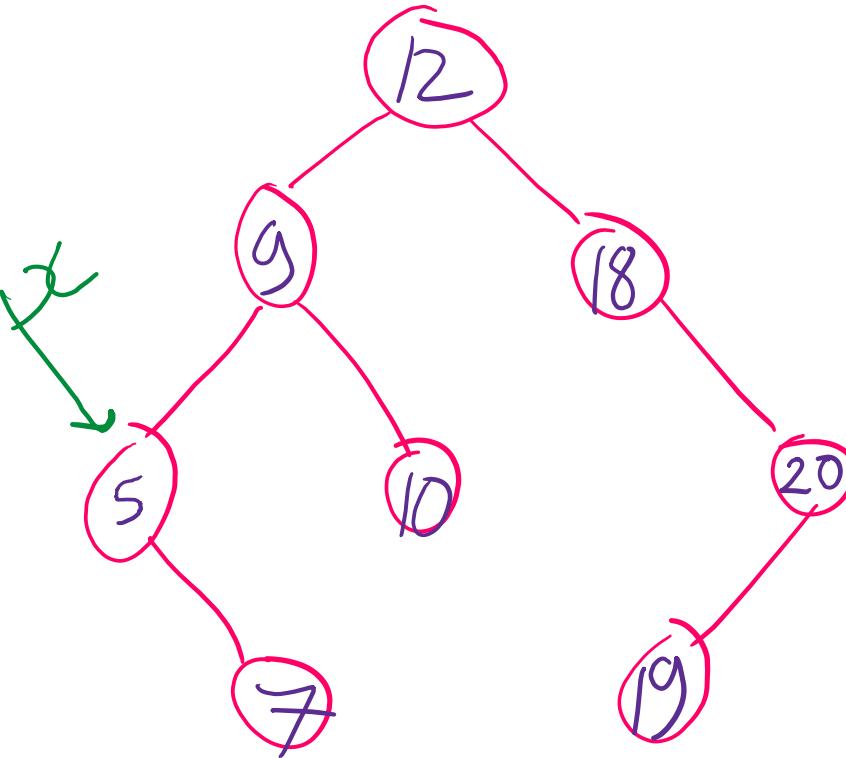
repeat until stack empty & $x = \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down

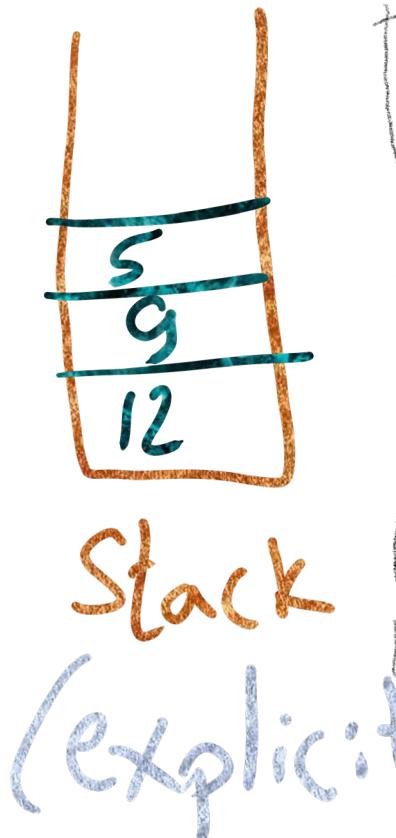
visit x
if x is right child
pop & visit
skip leftmost leaf

else $x = \text{parent. right}$

Iterative Postorder traversal



Visit order : 7, 5

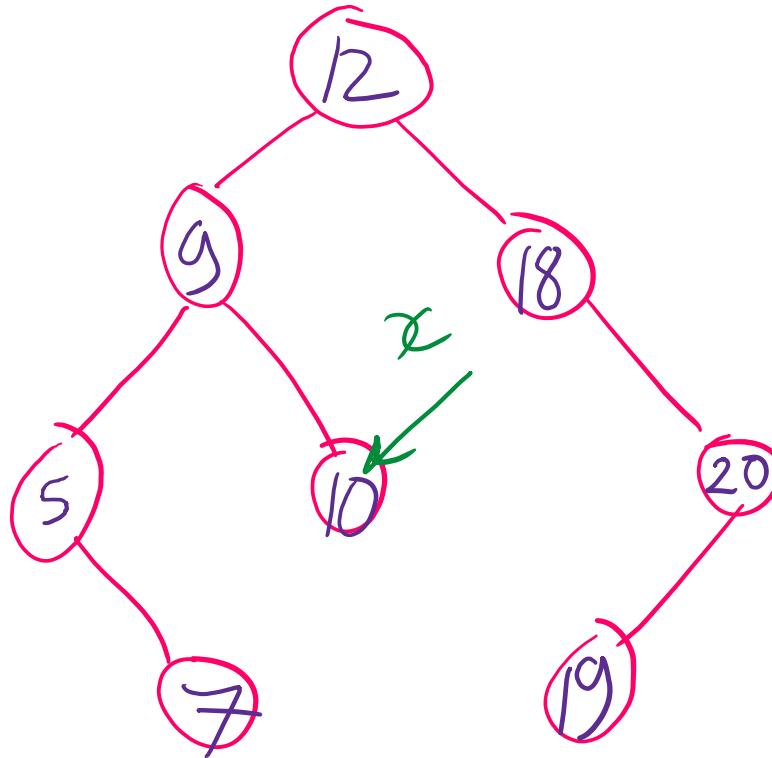


repeat until stack empty & $x = \text{null}$

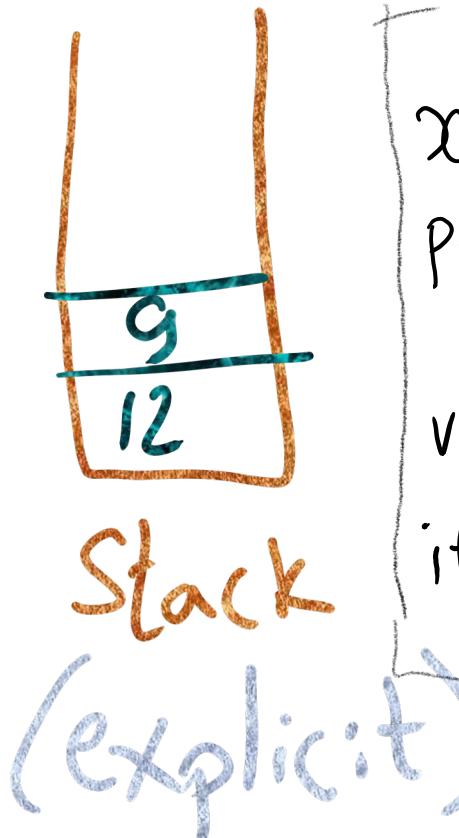
$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

else $x = \text{parent. right}$

Iterative Postorder traversal



Visit order : 7, 5

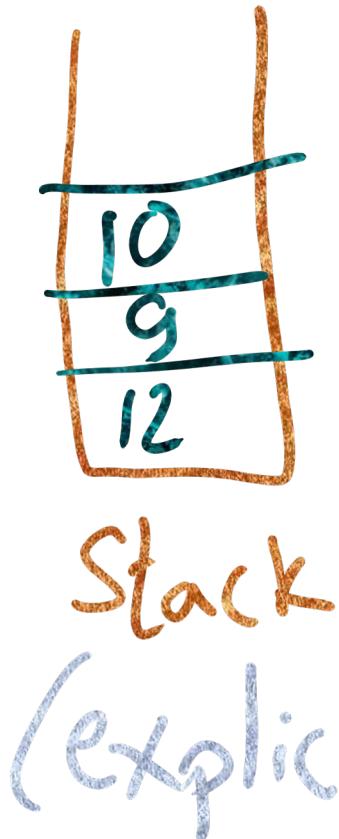
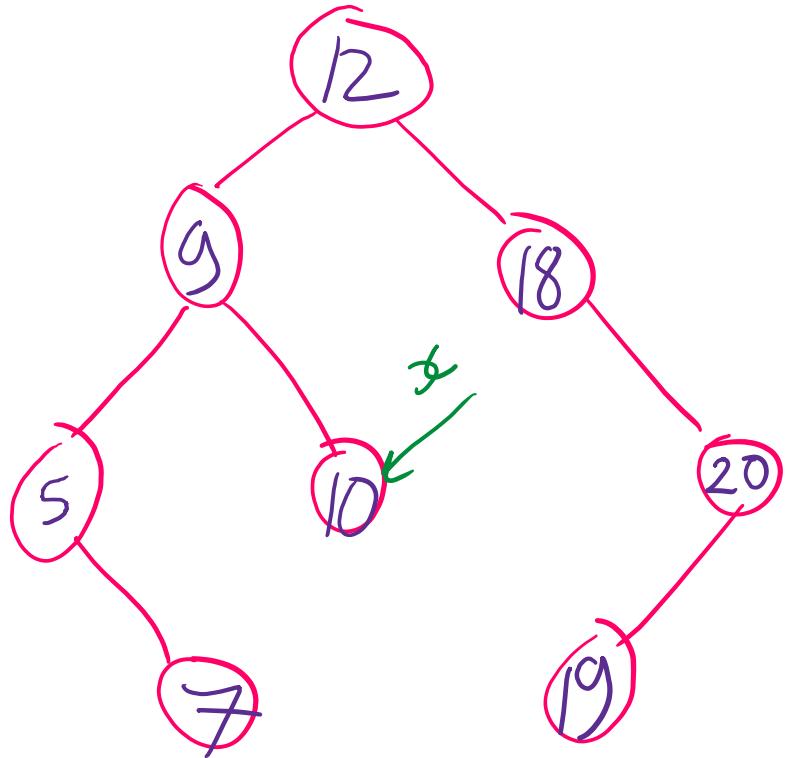


repeat until stack empty & $x = \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

else $x = \text{parent. right}$

Iterative Postorder traversal



repeat until stack empty & $x = \text{null}$

$x = \text{leftmost leaf}$

Push to stack while moving down

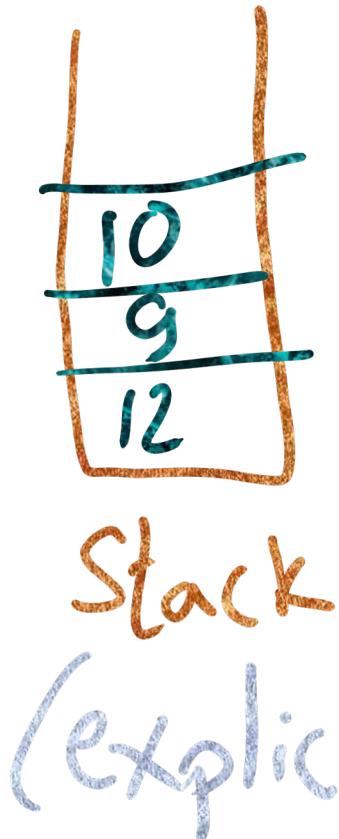
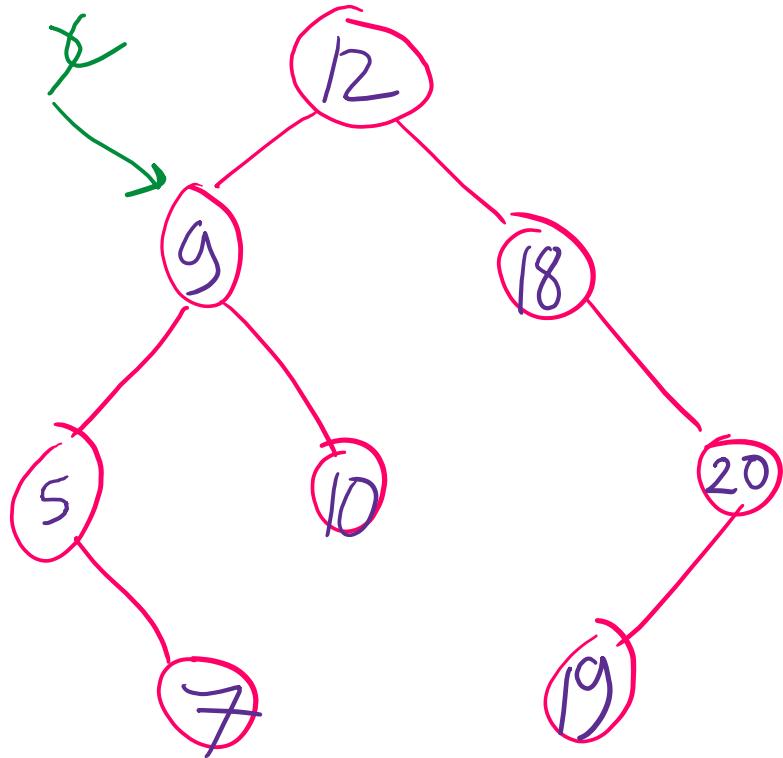
visit x

if x is right child
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10

else $x = \text{parent. right}$

Iterative Postorder traversal



repeat until stack empty & $x = \text{null}$

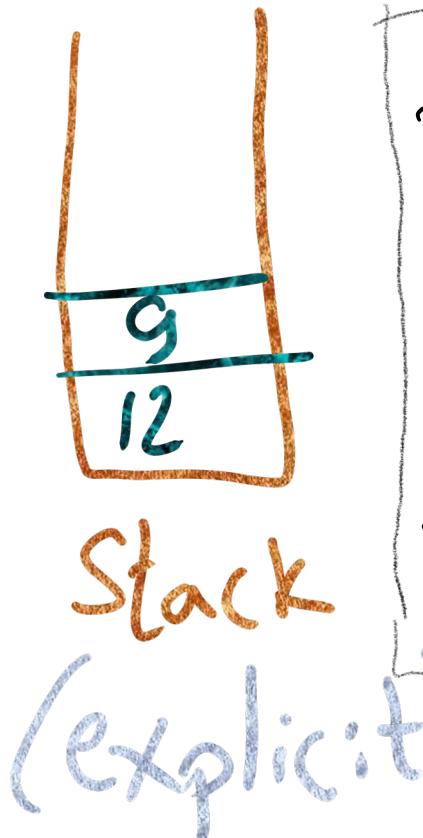
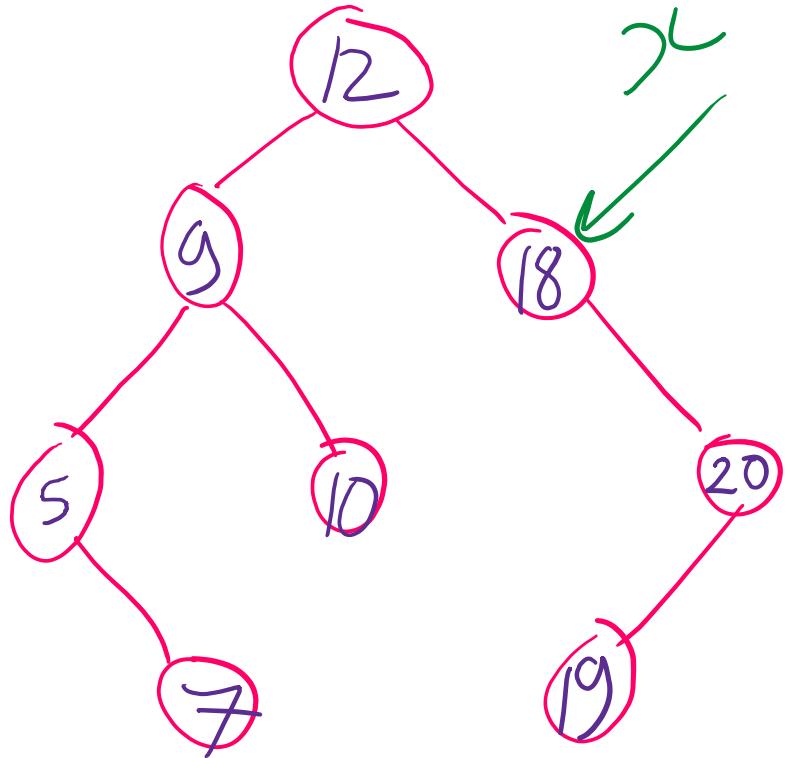
$x = \text{leftmost leaf}$
push to stack while moving down
visit x

if x is rightchild
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10, 9

else $x = \text{parent. right}$

Iterative Postorder traversal



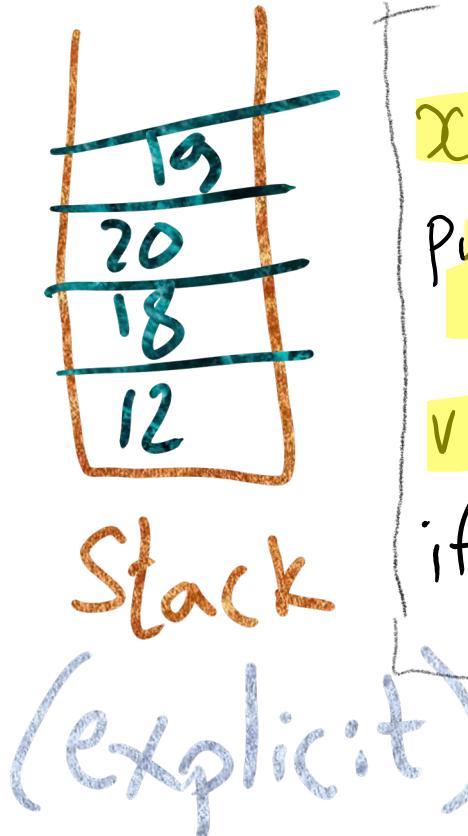
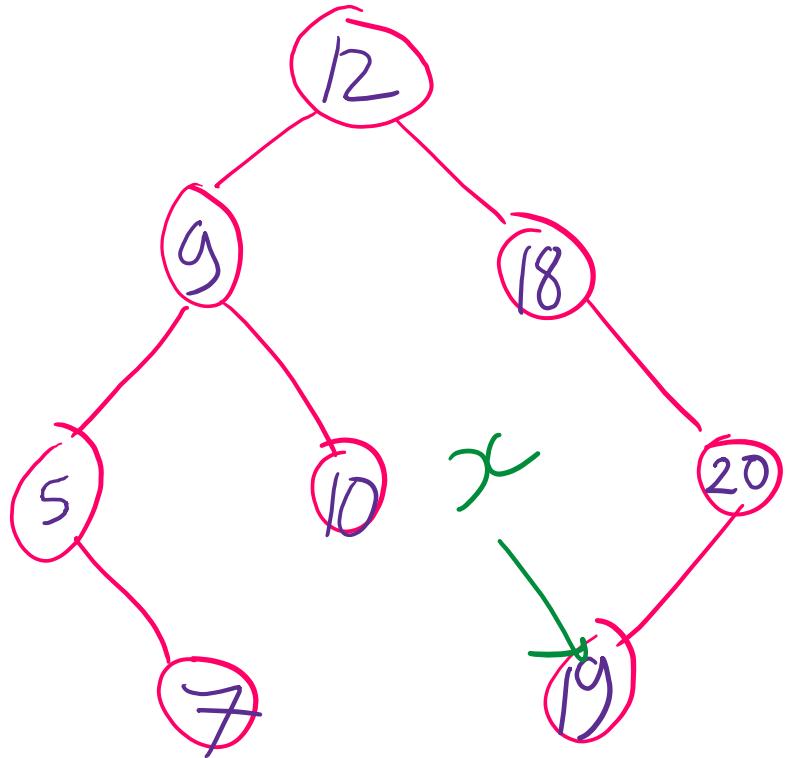
repeat until stack empty
& $x == \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10, 9

else $x = \text{parent. right}$

Iterative Postorder traversal



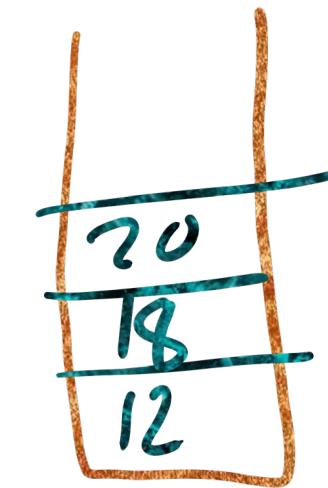
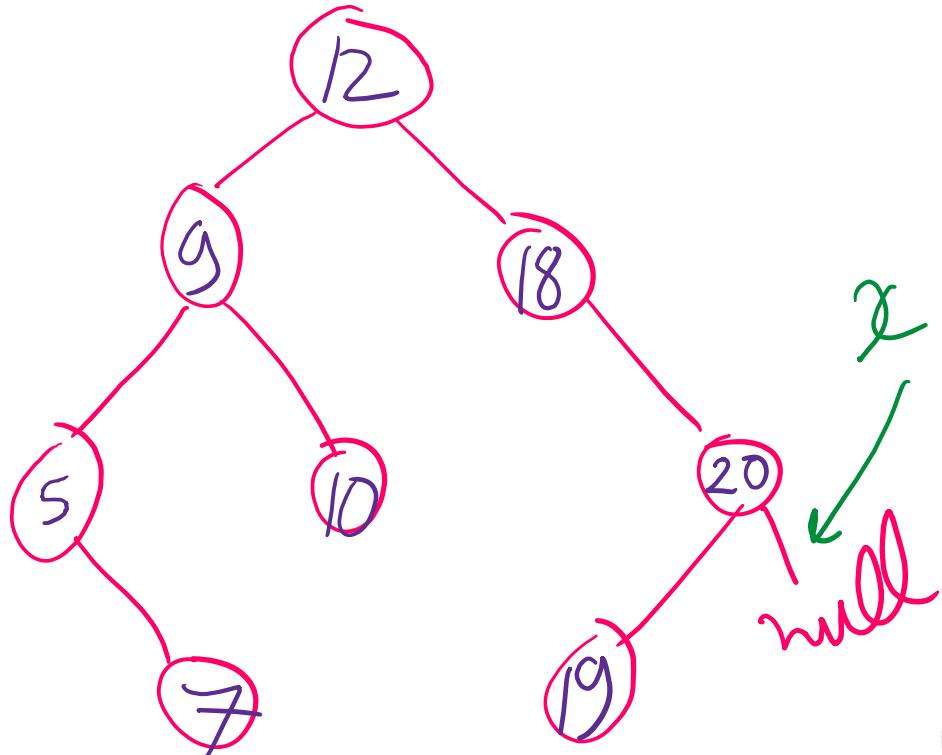
repeat until stack empty
 $\& x == \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

else $x = \text{parent. right}$

Visit order: 7, 5, 10, 9, 19

Iterative Postorder traversal



Stack
(explicit)

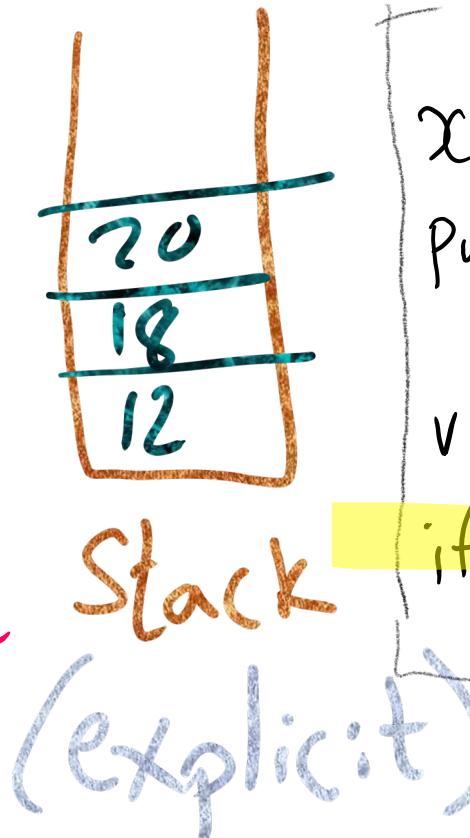
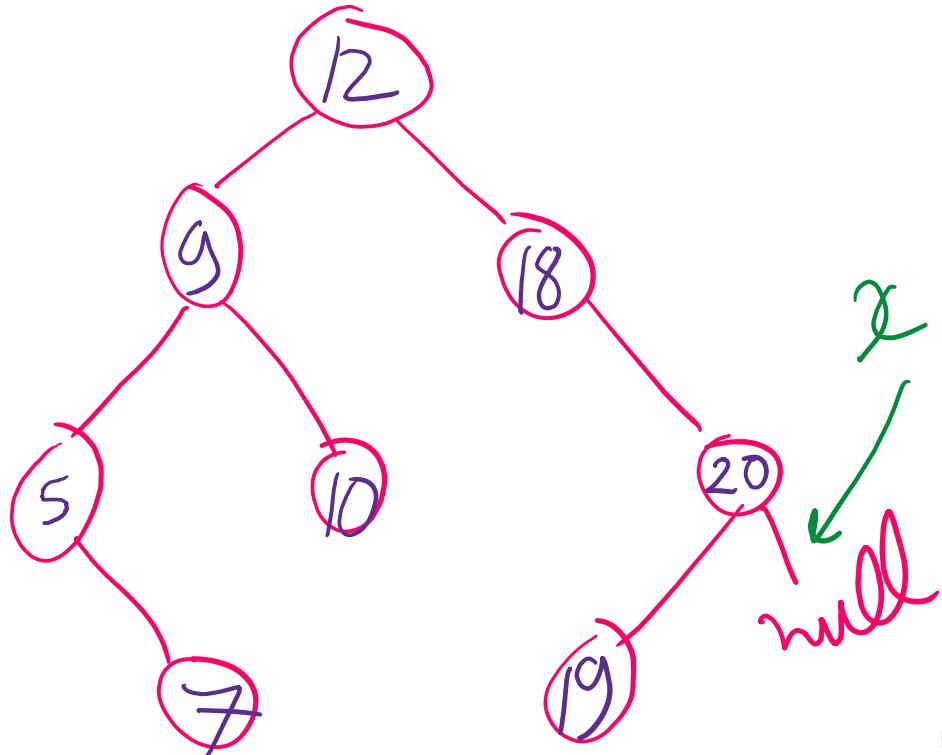
repeat until stack empty
& $x = \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

Visit order : 7, 5, 10, 9, 19

else $x = \text{parent. right}$

Iterative Postorder traversal



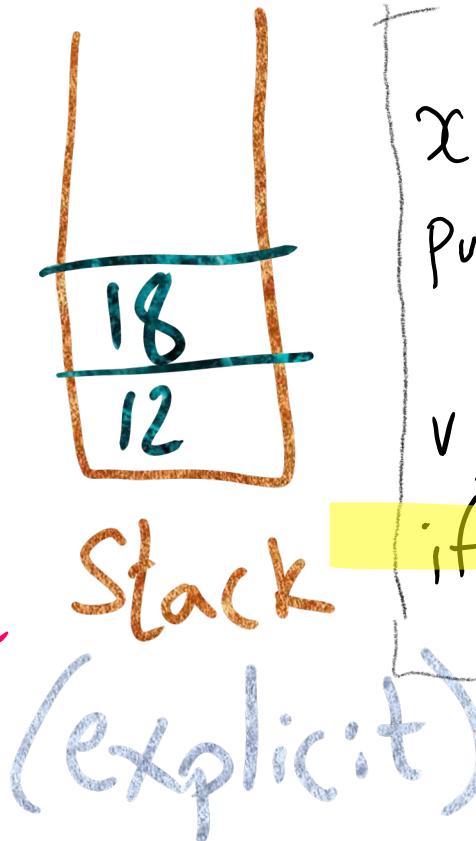
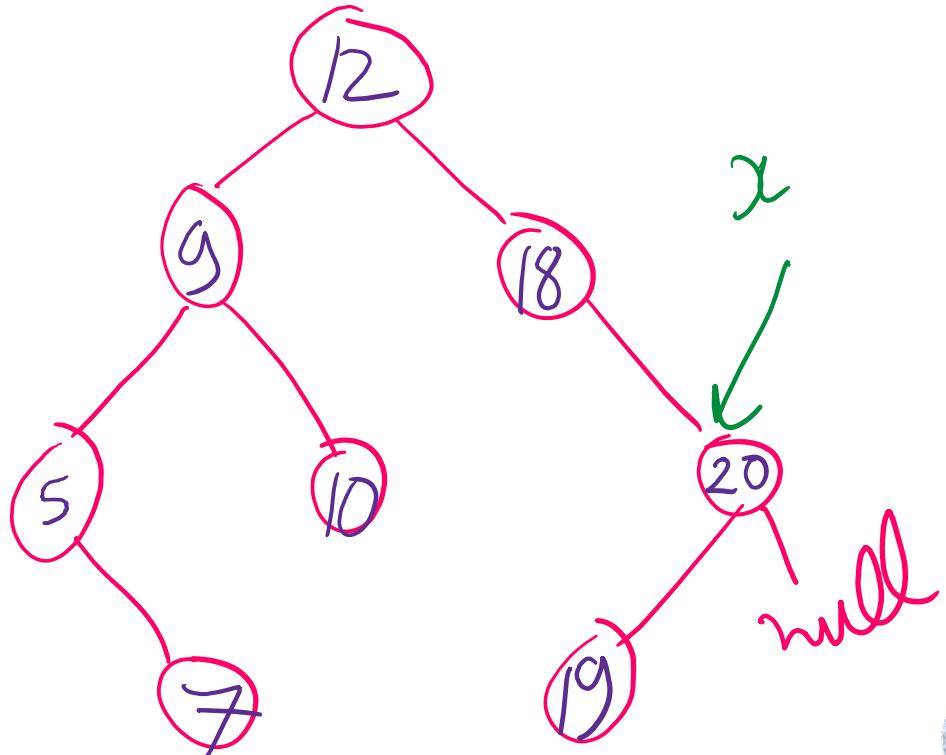
repeat until stack empty & $x = \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down
visit x
if x is right child
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10, 9, 19

else $x = \text{parent. right}$

Iterative Postorder traversal



repeat until stack empty & $x = \text{null}$

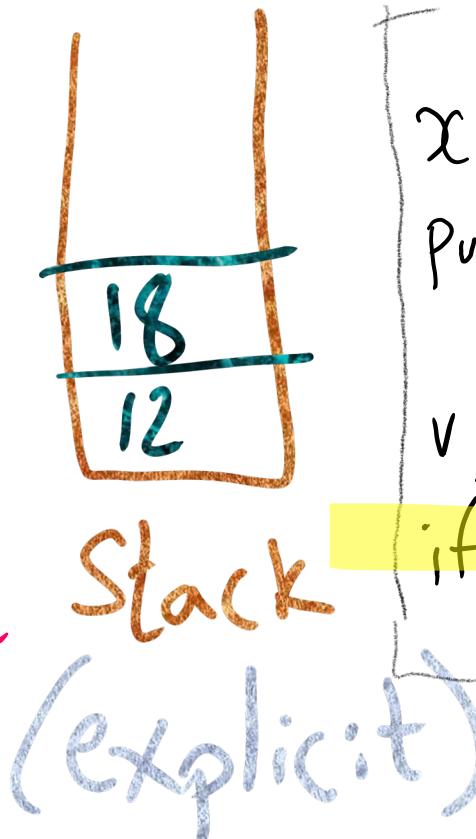
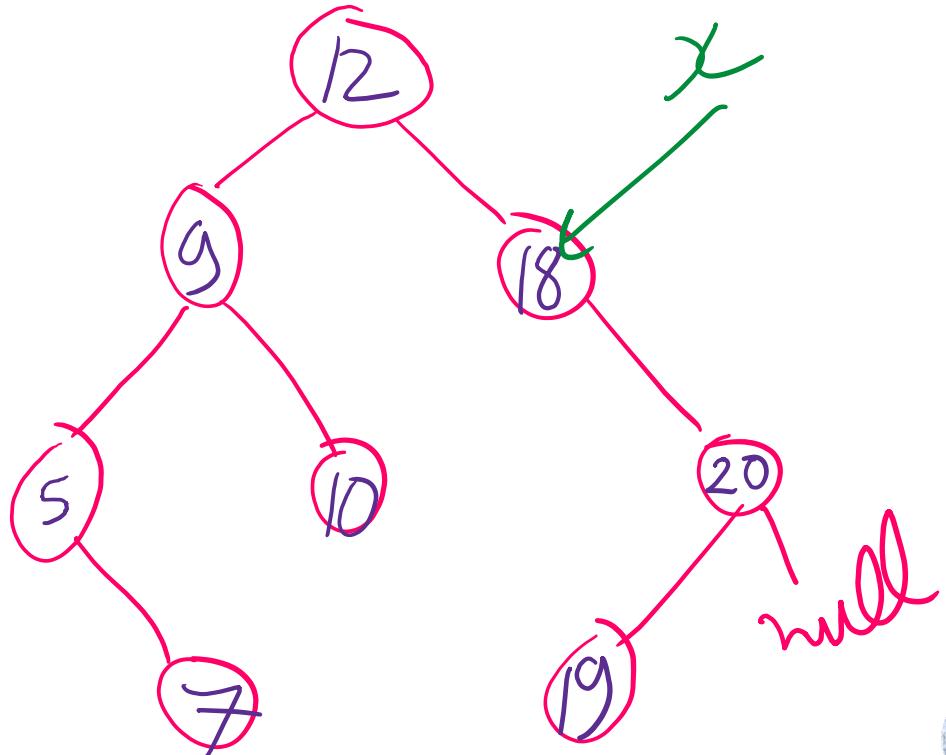
$x = \text{leftmost leaf}$
push to stack while moving down
visit x

if x is right child
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10, 9, 19, 20

else $x = \text{parent. right}$

Iterative Postorder traversal



repeat until stack empty & $x = \text{null}$

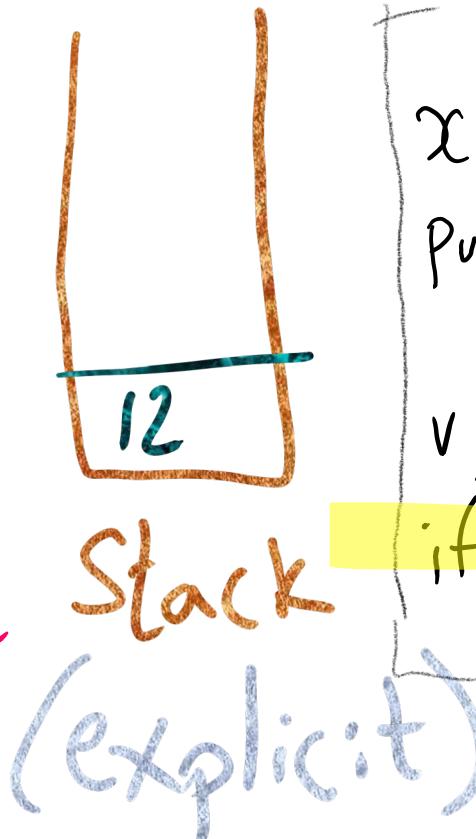
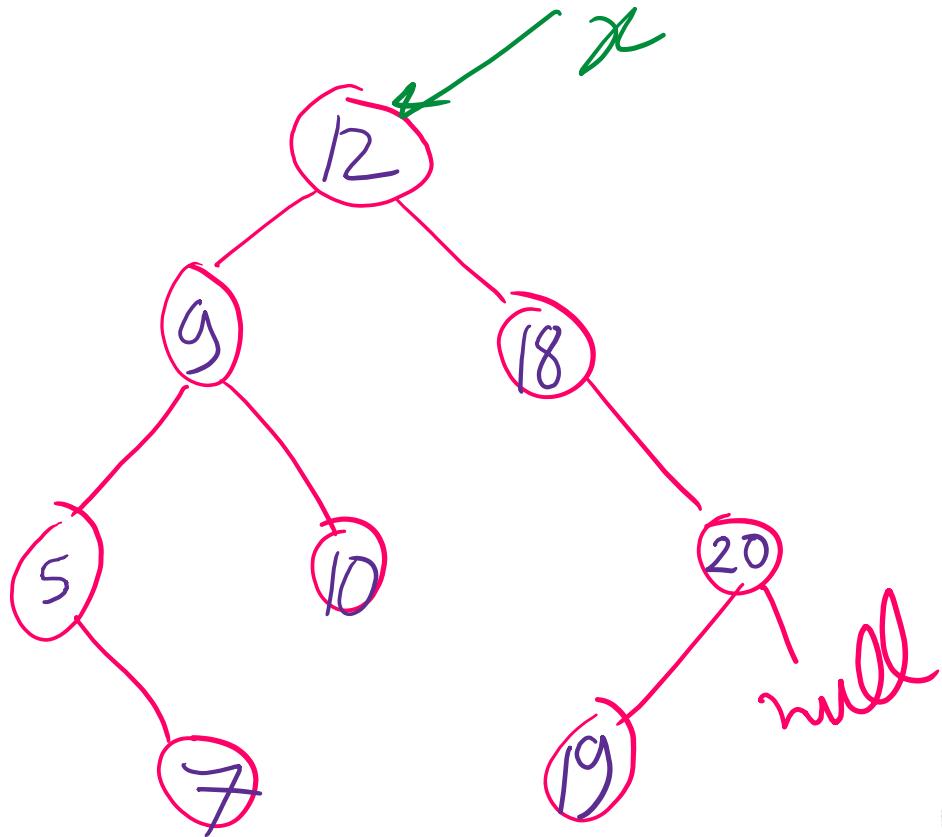
$x = \text{leftmost leaf}$
push to stack while moving down
visit x

if x is right child
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10, 9, 19, 20, 18

else $x = \text{parent. right}$

Iterative Postorder traversal



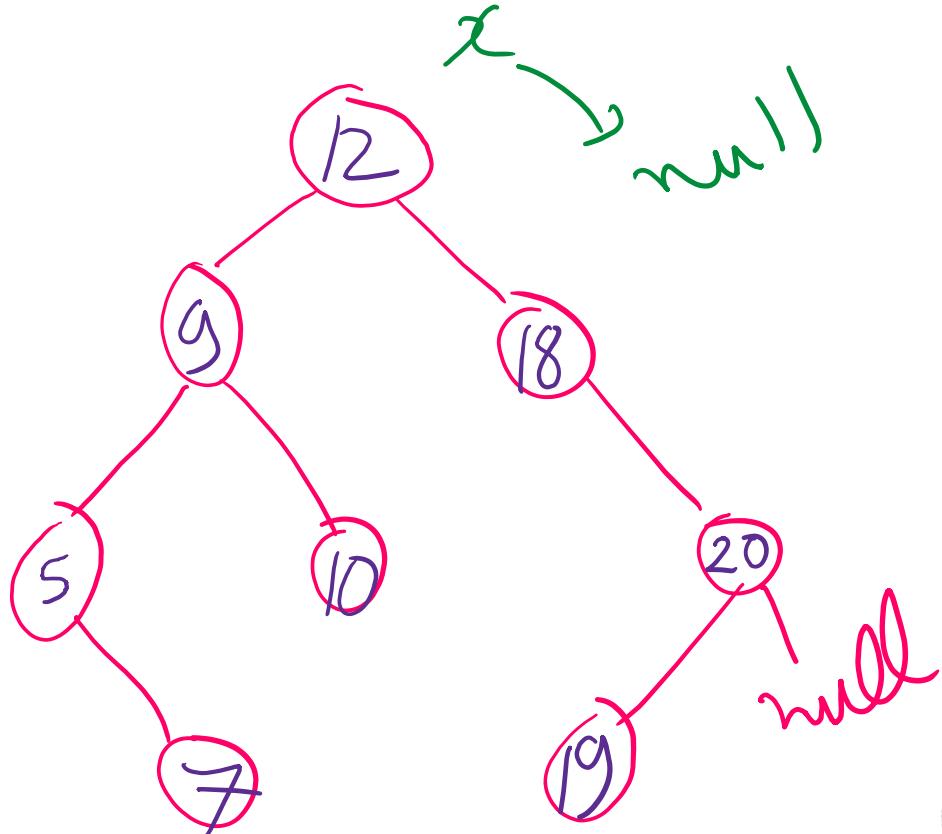
repeat until stack empty & $x = \text{null}$

$x = \text{leftmost leaf}$
push to stack while moving down
visit x

if x is right child
pop & visit
skip leftmost leaf

Visit order: 7, 5, 10, 9, 19, 20, 18, 12
else $x = \text{parent. right}$

Iterative Postorder traversal



Stack
(explicit)

repeat until stack empty
& $x = \text{null}$

$x = \text{leftmost leaf}$
push to stack while
moving down
visit x

if x is right child
pop & visit
skip leftmost
leaf

Visit order : 7, 5, 10, 9, 19, 20, 18, 12
else $x = \text{parent. right}$

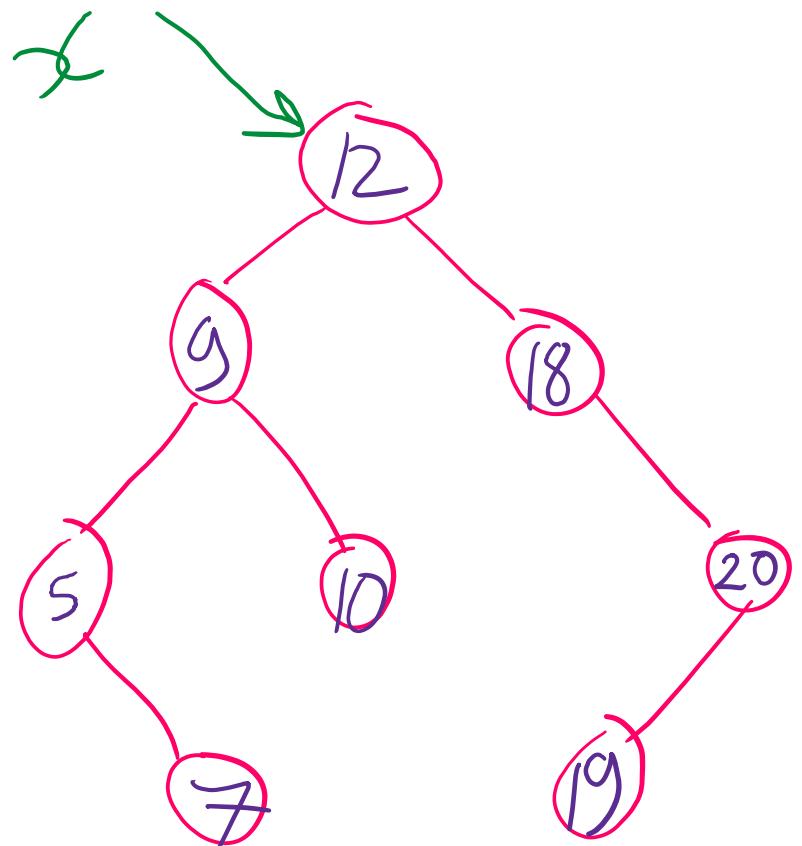
Traversals of a Binary Tree

- Preorder traversal
 - Visit root before we visit root's subtrees
- Inorder traversal
 - Visit root of a binary tree between visiting nodes in root's subtrees.
- Postorder traversal
 - Visit root of a binary tree after visiting nodes in root's subtrees
- Level-order traversal
 - Begin at root and visit nodes one level at a time
 - We will see the implementation when we learn Breadth-First Search of Graphs

Traversals of Binary Search Trees

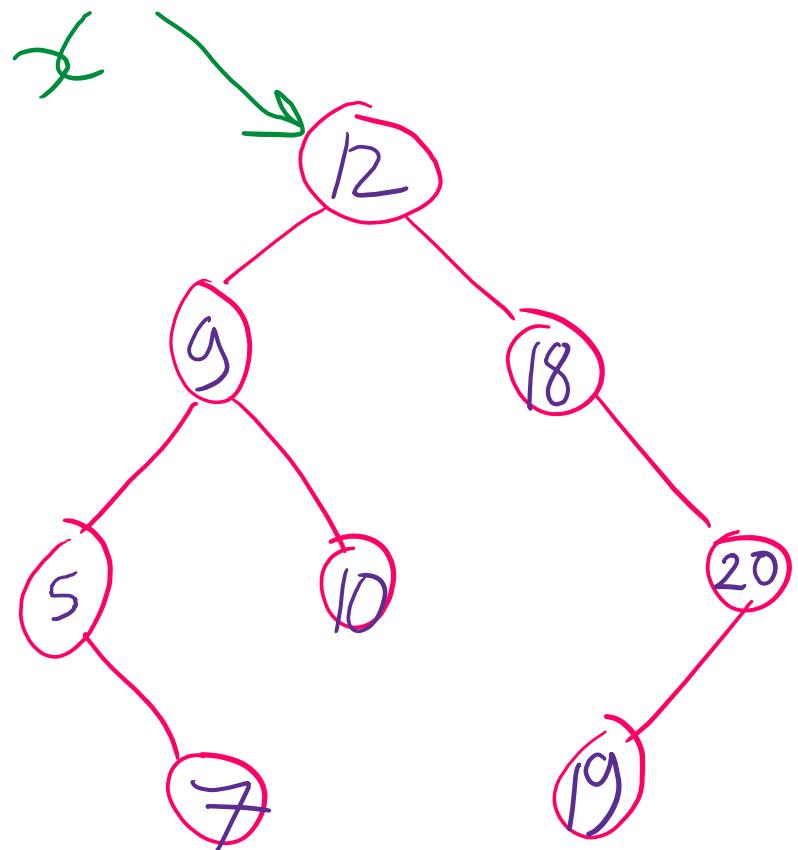
- Traversal for enumerating all items
- Traversal for finding an item

BST search using iteration



$x = \text{root}$
while($x \neq \text{null}$)
if equal break;
if $<$ $x = x.\text{left}$
if $>$ $x = x.\text{right}$
}

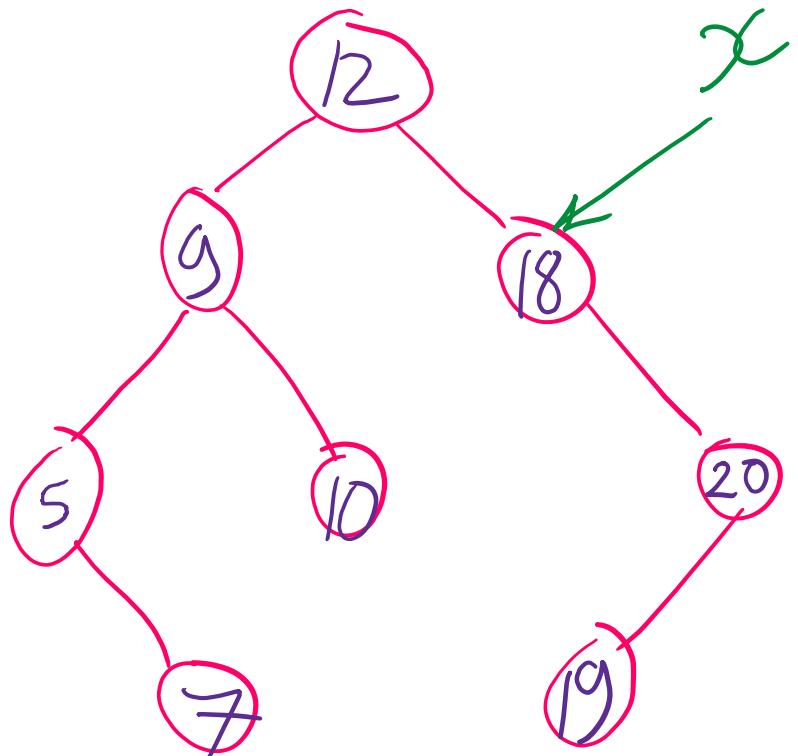
BST search using iteration



```
x = root  
while(x != null){  
    if equal break;  
    if < x=x.left  
    if > x=x.right  
}
```

Search for 19

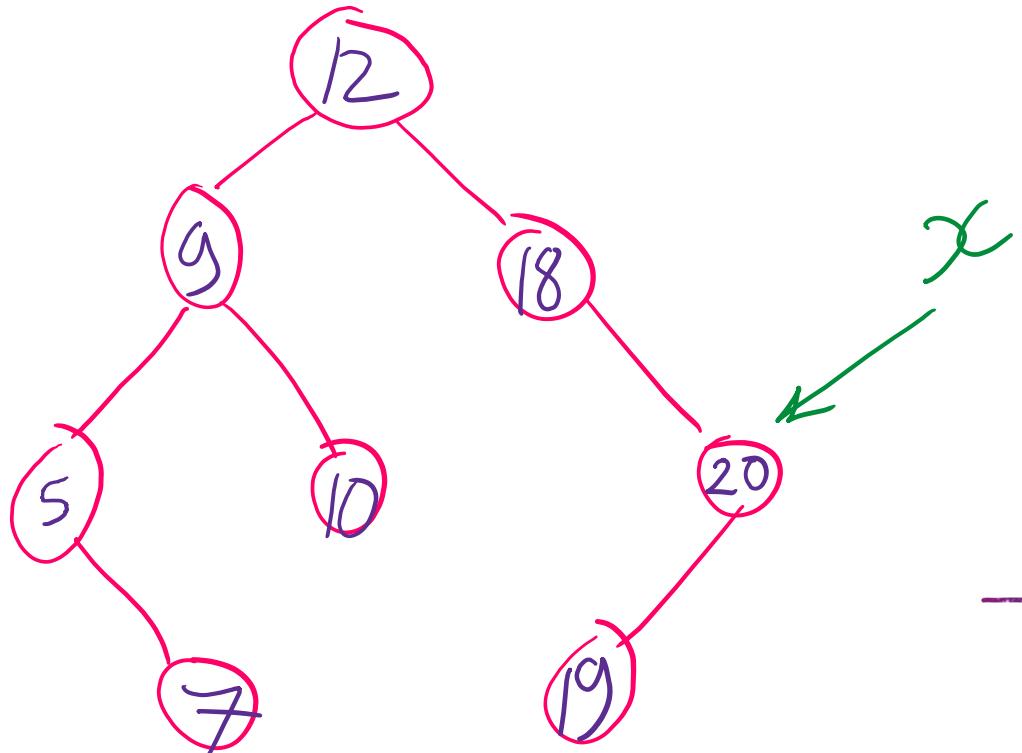
BST search using iteration



```
x = root  
while(x != null){  
    if equal break;  
    if < x=x.left  
    if > x=x.right  
}
```

Search for 19

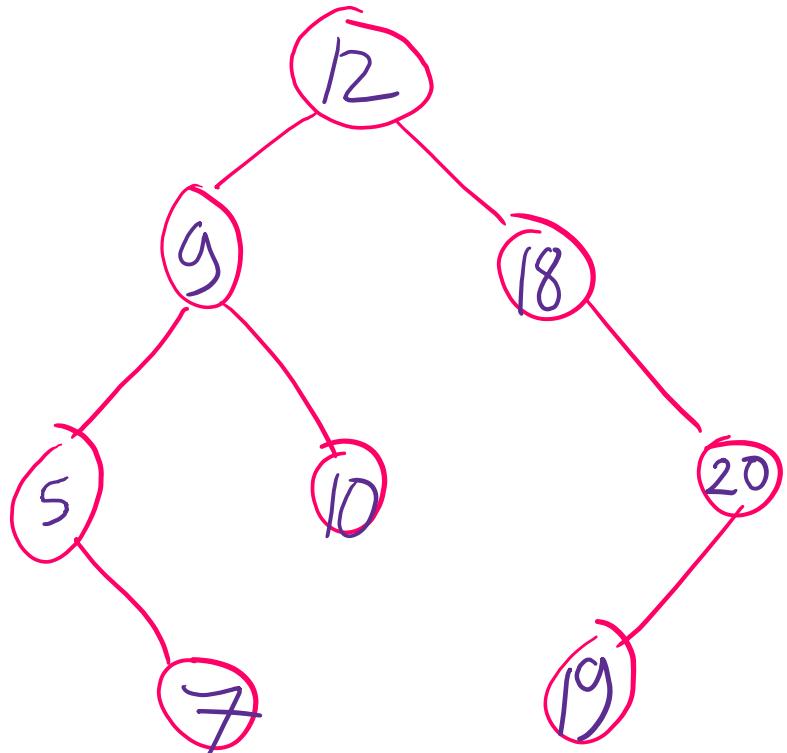
BST search using iteration



```
x = root  
while(x != null){  
    if equal break;  
    if < x = x.left  
    if > x = x.right  
}
```

Search for 19

BST search using iteration

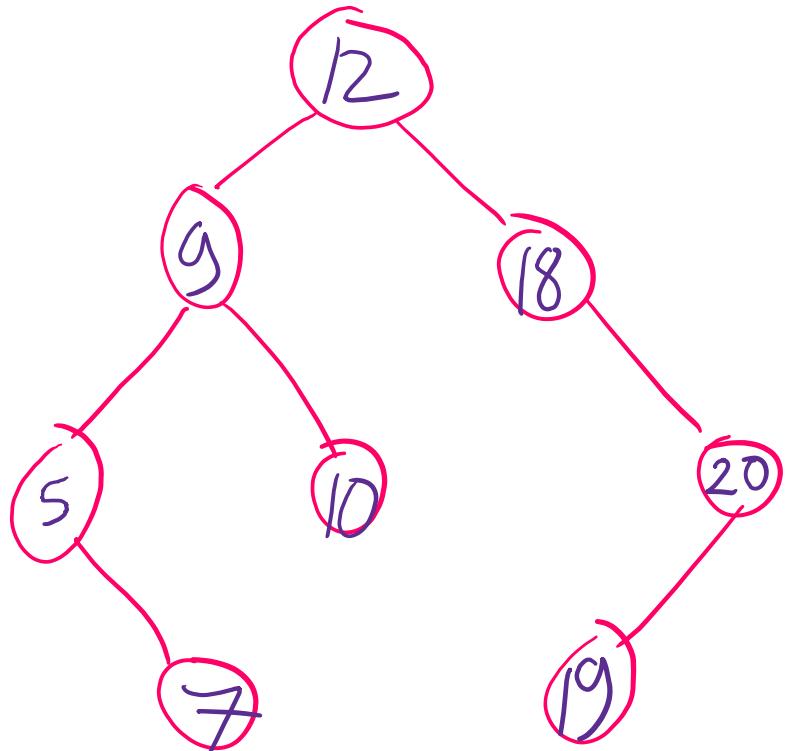


x

Search for 19

```
 $x = \text{root}$ 
while( $x \neq \text{null}$ ){
    if equal break;
    if  $<$   $x = x.\text{left}$ 
    if  $>$   $x = x.\text{right}$ 
}
```

BST search using iteration



x Search for 19

```
 $x = \text{root}$ 
while( $x \neq \text{null}$ ){
    if equal break;
    if  $<$   $x = x.\text{left}$ 
    if  $>$   $x = x.\text{right}$ 
}
```

Symbol Table Implementations

	Unsorted Array	Sorted Array	Unsorted L.L.	Sorted L.L.	BST	RB BST
add	$\Theta(1)$	$\Theta(n)$	$\Theta(1)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\log n)$
Search	$\Theta(n)$	$\Theta(\log n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\log n)$

Note: The original table includes "Binary Search" in blue text under the Sorted Array column, which is crossed out in purple. The "Search" row for the Sorted Array column is also crossed out in purple.