



University of
Pittsburgh

Introduction to Operating Systems CS 1550



Spring 2023
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(Some slides are from **Silberschatz, Galvin and Gagne ©2013**)

Announcements

- Upcoming deadlines
 - Homework 3 is due **this Friday**
 - Lab 1 is due on Tuesday 2/7 at 11:59 pm
 - Project 1 is due on Friday 2/17 at 11:59 pm
 - Discussed in this week's recitations

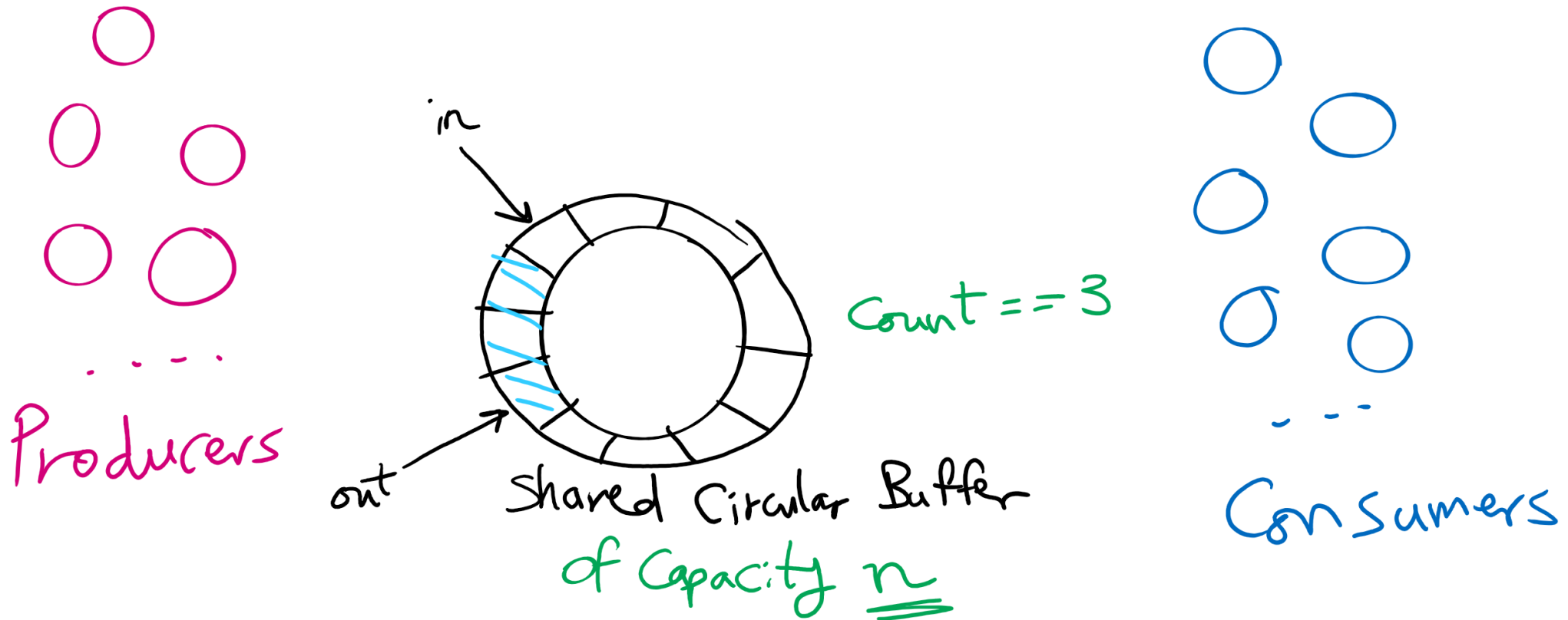
Previous lecture ...

- Bounded buffer problem
 - semaphore-based solution

Problem of the Day

It is easy to make mistakes when using semaphores

Producers Consumers Problem



Solving Producers Consumers using Semaphores

Semaphore $\text{empty}(\underline{n}), \text{full}(0)$
Mutex $\text{Sem}(1);$

Producer

$\text{down}(\text{empty})$

$\text{down}(\text{sem})$

$\text{buffer}[\text{in}] = \text{new item}$

$\text{in} += 1 \% n$

$\text{Count}++$

$\text{up}(\text{sem})$

$\text{up}(\text{full})$

Consumer

$\text{down}(\text{full})$

$\text{down}(\text{sem})$

$\text{item} = \text{buffer}[\text{out}]$

$\text{out} += 1 \% n$

$\text{Count}--$

$\text{up}(\text{sem})$

$\text{up}(\text{empty})$

Some thoughts

- Do we need the count variable?
- If we have one producer and one consumer
 - do we need the mutex?
- For multiple producers and consumers
 - Why do we need the mutex?
 - what benefit would we get if we have one mutex for producers and one for consumers?

Semaphore empty(n), full(0)
Mutex Sem(1);

Producer

```
down(empty)
down(sem)
buffer[in] = new item
in += 1 % n
Count++
up(sem)
up(full)
```

Consumer

```
down(full)
down(sem)
item = buffer[out]
out += 1 % n
Count--
up(sem)
up(empty)
```

Condition Variables

- Yet another construct for process/thread synchronization
 - (Add to Spinlock and Semaphore)
- Has 3 operations
 - wait()
 - signal()
 - broadcast()
- Not foreign to us at all
 - Every object variable in Java is a Condition Variable

Condition Variables Operations

- These 3 operations must be called while holding a mutex lock
- `wait()`
 - unlock mutex
 - block process
 - when awake, relock mutex
 - when successful, return
- `signal()`
 - wakeup one waiting process in the condition variable's queue if any
- `broadcast()`
 - wakeup all waiting processes in the condition variable's queue if any

Solving Bounded Buffer Using Condition Variables

Mutex sem;
ConditionVariable CV;

Producer

```
down(sem)
while(count == n)
    CV.wait()
buffer[in] = new item
in = (in + 1) % n
count ++
CV.broadcast()
up(sem)
```

Consumer

```
down(sem)
while(count == 0)
    CV.wait()
item = buffer[out]
out = (out + 1) % n
count --
CV.broadcast()
up(sem)
```

Readers & Writers Problem

- Many processes that may read and/or write
- Only one writer allowed at any time
- Many readers allowed, but not while a process is writing
- Real-world Applications
 - Database queries
 - We have this problem in Project 1

Semaphore-based Solution

Shared variables

```
int nreaders;  
Semaphore mutex(1), writing(1);
```

Reader process

```
...  
mutex.down();  
nreaders += 1;  
if (nreaders == 1) // wait if  
    writing.down(); // 1st reader  
mutex.up();  
// Read some stuff  
mutex.down();  
nreaders -= 1;  
if (nreaders == 0)    // signal if  
    writing.up();      // last reader  
mutex.up();  
...
```

Writer process

```
...  
writing.down();  
// Write some stuff  
writing.up();  
...
```

Solution Tracing

- enterRead

Reader process

...

```
mutex.down();
```

```
nreaders += 1;
```

```
if (nreaders == 1) // wait if
```

```
    writing.down(); // 1st reader
```

```
mutex.up();
```

```
// Read some stuff
```

```
mutex.down();
```

```
nreaders -= 1;
```

```
if (nreaders == 0) // signal if
```

```
    writing.up(); // last reader
```

```
mutex.up();
```

...

Solution Tracing

- read

Reader process

```
...
mutex.down();
nreaders += 1;
if (nreaders == 1) // wait if
    writing.down(); // 1st reader
mutex.up();
// Read some stuff
mutex.down();
nreaders -= 1;
if (nreaders == 0) // signal if
    writing.up(); // last reader
mutex.up();
...
```

Solution Tracing

- doneRead

Reader process

```
...
mutex.down();
nreaders += 1;
if (nreaders == 1) // wait if
    writing.down(); // 1st reader
mutex.up();
// Read some stuff
mutex.down();
nreaders -= 1;
if (nreaders == 0) // signal if
    writing.up(); // last reader
mutex.up();
...
```

Writer Events

- enterWrite

Writer process

...

writing.down();

// Write some stuff

writing.up();

...

Writer Events

- write

Writer process

...

```
writing.down();
```

```
// Write some stuff
```

```
writing.up();
```

...

Writer Events

- doneWrite

Writer process

...

```
writing.down();
```

```
// Write some stuff
```

```
writing.up();
```

...