

Midterm 1

Due: October. 28, 2021

Reminder: Your name should not appear anywhere on your handin; each individual page of the homework should include your Banner ID only. For your digital submission, each page should include work for only one problem (i.e., make a new page/new pages for each problem).

Collaboration or any kind or use of any third-party resource is strictly **not allowed** on this assignment. The TA staff will only answer questions regarding clarifications on the statement of the problems in the assignment. Questions should be asked via private posts on Ed. Please monitor Ed, as we will post clarifications of frequently asked questions there.

Hand in your solutions by 14:29pm to Gradescope. See the course syllabus for the late policy. Usage of any materials outside of course notes, the course textbook, lecture slides and Ed posts is strictly forbidden.

You may apply up to two late days to this assignment.

Problem 1 - 10 points

Let $X[1, \dots, n]$ and $Y[1, \dots, n]$ be two arrays, each containing n numbers already in sorted - non decreasing - order. For simplicity assume n is a power of 2, Give an $O(\log n)$ -time algorithm to find the median of all $2n$ elements in arrays X and Y .

- (a) Provide a succinct, but clear description of your algorithm. You may provide a pseudocode.
- (b) Provide a proof of the correctness of the algorithm.
- (c) Provide and analysis of the running time (asymptotic analysis is correct) and memory utilization of the algorithm.

Hint: Note that the given arrays are already sorted and of the **same size!** You may want to use log search to exploit this fact :)

Problem 2 - 10 points

Provide an efficient greedy algorithm that given a set $\{x_1, x_2, \dots, x_n\}$ of points on the real line, determines the smallest set of closed intervals of length 3 that contains all of the given points.

- (a) Provide a succinct, but clear description of your algorithm. You may provide a pseudocode.
- (b) Provide a proof that the algorithm returns a correct (i.e., does not violate the rules) and optimal guest list (i.e., a list with maximum possible conviviality score among the possible correct guest lists).
- (c) Provide and analysis of the running time (asymptotic analysis is correct) and memory utilization of the algorithm.

Problem 3 - 10 points

You have been hired by a top 500 corporation...Congratulations (I guess)!!! Your first job is to plan a company party. The company has a hierarchical structure, that is the supervisor relation forms a tree rooted at the president. The HR department has assigned each company employee with a conviviality ranking which is a positive real number. In order to avoid awkwardness, the president of the company does not want both an employee and their immediate supervisor to attend. You are given the tree that describes the company's structure and the conviviality scores of all employees. Describe a Dynamic Programming algorithm that compiles a guest list which maximizes the sum of conviviality scores of the guests.

- (a) Provide a succinct, but clear description of your algorithm. You may provide a pseudocode.
- (b) Provide a proof that the algorithm returns a correct (i.e., does not violate the rules) and optimal guest list (i.e., a list with maximum possible conviviality score among the possible correct guest lists).
- (c) Provide an analysis of the running time (asymptotic analysis is correct) and memory utilization of the algorithm.

Problem 4 -10 points

- (a) Give an efficient divide and conquer algorithm to convert a given n -bit binary integer to a decimal representation and prove its correctness. For simplicity, assume n is a power of 2.
- (b) Analyze the running time of your algorithm. Assume that it is possible to multiply two decimal integer numbers with n digits in $O(n^{\log_2 3})$ time.

Hint: Recall that a given n -digit integer x is expressed in binary as $x = (x_{n-1}, x_{n-2}, \dots, x_1, x_0)_2$ where x_0 is its least significant digit and x_{n-1} is its most significant digit. Further, let $x_l = (x_{n/2-1}, x_{n/2-2}, \dots, x_1, x_0)_2$ be the $n/2$ -digit integer in binary which corresponds to the $n/2$ least significant digits of x , and let $x_m = (x_{n-1}, x_{n-2}, \dots, x_{n/2+1}, x_{n/2})_2$ be the $n/2$ -digit integer in binary which corresponds to the $n/2$ most significant digits of x . We have $x = x_l + 2^{n/2} * x_m$. This should suggest us a way to set up a divide and conquer strategy...:) Careful about the number of subproblems!