# project3\_tree

## October 13, 2020

```
[]: # Please do not change this cell because some hidden tests might depend on it.
     import os
     # Otter grader does not handle ! commands well, so we define and use our
     # own function to execute shell commands.
     def shell(commands, warn=True):
         """Executes the string `commands` as a sequence of shell commands.
            Prints the result to stdout and returns the exit status.
            Provides a printed warning on non-zero exit status unless `warn`
           flag is unset.
         file = os.popen(commands)
         print (file.read().rstrip('\n'))
         exit_status = file.close()
         if warn and exit_status != None:
             print(f"Completed with errors. Exit status: {exit_status}\n")
         return exit_status
     shell("""
     ls requirements.txt >/dev/null 2>&1
     if [ ! $? = 0 ]; then
     rm -rf .tmp
     git clone https://github.com/cs187-2020/project3.git .tmp
     mv .tmp/requirements.txt ./
     rm -rf .tmp
     fi
     pip install -q -r requirements.txt
```

```
[]: # Initialize Otter
import otter
grader = otter.Notebook()
```

# 1 Project 3: Parsing Using the Cocke-Kasami-Younger (CKY) Algorithm

Constituency Parsing is an intermediary task in Natural Language Processing, where the goal is to extract the syntactic parse tree given a sentence.

In this project, you will implement the CKY algorithm for context-free grammars (CFG). You will implement parsing algorithms for both a determinstic grammar (CFG) and a probabilistic grammar (PCFG).

#### 1.1 Goals

- 1. Finish a CFG for the ATIS dataset.
- 2. Implement an algorithm for determining whether a parse exists or not given a sentence.
- 3. Implement the CKY algorithm for parsing using CFGs.
- 4. Construct a probabilistic context-free grammar (PCFG) based on a CFG.
- 5. Implement the CKY algorithm for parsing using PCFGs.

# 1.2 Setup

[]: !pip install nltk

```
[]: import nltk
     nltk.download('large_grammars')
     from nltk.grammar import ProbabilisticProduction, PCFG, Nonterminal
     from nltk.tree import Tree
     from nltk import treetransforms
     from collections import defaultdict, Counter
[]: # Grammar Coverage
     !wget -nv -N -P data https://raw.githubusercontent.com/nlp-course/data/master/
     →ATIS/train.nl
     # PCFG Construction
     !wget -nv -N -P data https://raw.githubusercontent.com/nlp-course/data/master/
     →ATIS/train.trees
     !wget -nv -N -P data https://raw.githubusercontent.com/nlp-course/data/master/
     →ATIS/dev.trees
     !wget -nv -N -P data https://raw.githubusercontent.com/nlp-course/data/master/
     →ATIS/test.trees
     # Tree Utils
```

## 1.2.1 Helper Functions

The standard CKY algorithm requires the grammar to be in Chomsky normal form (CNF), i.e., only rules of the form  $A \rightarrow B$  C,  $A \rightarrow a$  are allowed where A, B, C are nonterminals and a is a terminal symbol. However, in some downstream applications (such as our next project) we want to use grammar rules of more general forms, such as  $A \rightarrow B$  C D. To satisfy both of these constraints, we will convert the grammar to CNF, parse using CKY, and then convert back to the form of the original grammar. We have provided helper functions to perform those conversions.

To convert a grammar to CNF:

```
cnf_grammar, cnf_grammar_wunaries = get_cnf_grammar(grammar)
```

To convert a tree output from CKY back to the original form of the grammar in place:

```
un_cnf(tree, cnf_grammar_wunaries)
```

Note that above we need to pass in cnf\_grammar\_wunaries, an intermediate version of the grammar before removing unary rules. It's the second returned value of get cnf grammar.

```
[]: from scripts.tree_utils import get_cnf_grammar, un_cnf
```

#### 1.3 A custom ATIS grammar

To parse, we need a grammar. In this project, we will use a hand-crafted grammar for the ATIS dataset. It captures a large fraction of the language in the ATIS dataset while including a minimal number of rules. This tiny grammar will be used again in the next project for a question answering application.

#### 1.3.1 Goal 1: Finish the CFG for the ATIS dataset

Fill in the missing rule for PREJ. This is a production rule that captures all the "junk" before the actual semantically meaningful question. For example, in "i would like a flight between boston and dallas", "i would like" carries no meaningful information hence is considered "junk".

Hint: you will make use of the JUNK preduction

```
[]: #TODO
miniATIS = nltk.CFG.fromstring("""
```

```
S -> PREJ DET ADJS FLIGHT PPS | DET ADJS FLIGHT PPS | PREJ ADJS FLIGHT PPS | _ _
→ADJS FLIGHT PPS | PREJ DET FLIGHT PPS | DET FLIGHT PPS | PREJ FLIGHT PPS |
\hookrightarrowFLIGHT PPS | PREJ DET ADJS FLIGHT | DET ADJS FLIGHT | PREJ ADJS FLIGHT |_{1.1}
→ADJS FLIGHT | PREJ DET FLIGHT | DET FLIGHT | PREJ FLIGHT | FLIGHT
PPS -> PP | PP PPS
ADJS -> ADJ | ADJ ADJS
FLIGHT -> 'flights' | 'flight' | 'to' 'fly'
PREJ ->_
JUNK -> 'me' | 'show' | 'now' | 'only' | 'can' | 'you' | 'the' | 'itinerary' |
→"what's" | 'are' | 'my' | 'choices' | 'for' | 'i' | 'would' | 'like' | "i'd"□
\hookrightarrow | 'to' | 'see' | 'have' | 'make' | 'book' | 'find' | 'information' | 'on' |
→'know' | 'some' | 'hello' | 'yes' | 'please' | 'repeat' | 'do' | 'have' | __

→'there' | 'need' | 'hi' | 'get' | 'may' | 'listing' | 'listings' | 'travel' |

□
→ | 'able' | 'put' | 'requesting' | "i'm" | 'looking' | 'display' | UNK
ADJ -> AIRLINE | SIMPLEWEEKDAY | FLIGHTTYPE | SIMPLETIME | "monday's" |

→ "tuesday's" | "wednesday's" | "thursday's" | "friday's" | "saturday's" |

□
-- "sunday's" | 'available' | 'possible' | 'first' 'class' | 'economy' | ...
_{\hookrightarrow} 'thrift' 'economy' | 'cheapest' | 'lowest' 'cost' | 'least' 'expensive' | _{\sqcup}
→'most' 'expensive' | 'weekday' | 'daily' | 'last' | 'first' | 'dinner' |
DET -> 'all' 'the' | 'all' | A | 'an' | THE | 'any' | 'all' 'of' 'the'
BETWEEN -> 'between'
AND -> 'and'
OR -> 'or'
EITHER -> 'either'
OF -> 'of'
THE -> 'the'
A -> 'a'
PP -> PPLACE PLACE OR PLACE | PPLACE EITHER PLACE OR PLACE | PPLACE PLACE |
→EITHER PLACE OR PLACE | PLACE | PLACE | BETWEEN PLACE AND PLACE |
→BETWEEN TIME AND TIME | PTIME TIME | TIME | PDAY WEEKDAY | WEEKDAY TIME |
→WEEKDAY | PDAY WEEKDAY TIME | PDAY DATE | DATE | PAIRLINE AIRLINE | AIRCRAFT |
\hookrightarrow | FLIGHTTYPE | FARETYPE | PRICE | FOOD | AVAIL | POSTJ
```

```
PPLACE -> 'to' | 'that' 'arrive' 'at' | 'that' 'arrives' 'in' | 'coming' 'back'
\rightarrow 'going' 'to' | 'into' | 'for' | 'with' 'the' 'destination' 'city' 'of' |
_{\hookrightarrow}'in' | 'and' 'then' | 'arriving' 'to' | 'from' | 'leaving' | 'return' 'from'
→| 'leaving' 'from' | 'departing' 'from' | 'departing' | 'go' 'from' | 'going'

→ 'leaves' 'from' | 'which' 'leave' | 'that' 'leave' | 'originating' 'in' |
□
→ | 'and' 'a' 'stopover' 'in' | 'stop' 'in' | 'stopping' 'in' | 'make' 'a'

→ 'stops' 'in' | 'that' 'stops' 'over' 'in' | 'by' 'way' 'of' | 'connecting'
□
PLACE -> 'anywhere' | 'atlanta' | 'austin' | 'baltimore' | 'boston' | 'boston' |
→'logan' | 'burbank' | 'bwi' | 'charlotte' | 'chicago' | 'cincinnati' | 

→'cleveland' | 'cleveland' 'ohio' | 'columbus' | 'dallas' 'fort' 'worth' |
□

→'general' 'mitchell' 'international' | 'general' 'mitchell' | 'houston' |
□
→'indianapolis' | 'jfk' | 'kansas' 'city' | 'laguardia' 'airport' | 'las'

_{\hookrightarrow}'vegas' | 'long' 'beach' | 'los' 'angeles' | 'love' 'field' | 'memphis' |_{\sqcup}
_{\hookrightarrow}'miami' | 'milwaukee' | 'minneapolis' | 'montreal' | 'montreal' 'quebec' | _{\sqcup}
→'montreal' 'canada' | 'nashville' | 'new' 'york' 'city' | 'new' 'york' | □
→'newark' | 'newark' 'new' 'jersey' | 'oakland' | 'oakland' 'california' | □

→'ontario' | 'orlando' | 'orlando' 'florida' | 'philadelphia' | 'philly' |

□
→'phoenix' | 'pittsburgh' | 'salt' 'lake' 'city' | 'san' 'diego' | 'san' 

→'diego' 'california' | 'san' 'francisco' | 'san' 'jose' | 'seattle' | 'st.'

□
→'louis' | 'st.' 'paul' | 'st.' 'petersburg' | 'tacoma' | 'tacoma' |
→'washington' | 'tampa' | 'toronto' | 'various' 'cities' | 'washington' |
\hookrightarrow 'washington' 'dc' | 'dc' | 'westchester' 'county' | UNK
PTIME -> 'that' 'arrive' 'before' | 'that' 'arrives' 'before' | 'arriving'

→ 'before' | 'arrival' 'by' | 'arrives' | 'before' | 'departing' 'before' |

□
→ 'around' | 'arriving' 'about' | 'that' 'arrive' 'soon' 'after' | 'leaving' |
→'at' | 'leaving' | 'which' 'leave' 'after' | 'leaving' 'after' | 'after' | ___
→'departing' 'after' | 'that' 'depart' 'after' | 'departing' 'at' | 'arriving' |
```

```
TIME -> SIMPLETIME | 'the' 'afternoon' | 'the' 'late' 'afternoon' | 'the'
_{\hookrightarrow}'evening' | 'the' 'morning' | 'the' 'early' 'am' | 'the' 'day' | 'mornings' | _{\sqcup}
→'afternoons' | 'evenings' | 'as' 'early' 'as' 'possible' | 'earliest'
→'possible' 'time' | 'as' 'soon' 'thereafter' 'as' 'possible' | UNK
SIMPLETIME -> 'one' | 'two' | 'three' | 'four' | 'five' | 'six' | 'seven' | |
_{\hookrightarrow}'eight' | 'nine' | 'ten' | 'eleven' | 'twelve' | '1' | '2' | '3' |'4' | '5' | _{\sqcup}
→'6' | '7' | '8' | '9' | '10' | '11' | '12' | '1pm' | '2pm' | '3pm' | '4pm' | ⊔
→'2am' | '3am' | '4am' | '5am' | '6am' | '7am' | '8am' | '9am' | '10am' |
→'11am' | '12am' | '230' | '1505' | '630am' | '720am' | '723pm' | '819' | '845'ц

      → | '1026' | '1145am' | '1' "o'clock" | '1' "o'clock" 'am' | "1" "o'clock" "pm"ц

→ | "2" "o'clock" | "2" "o'clock" "am" | "2" "o'clock" "pm" | "3" "o'clock" |
→"3" "o'clock" "am" | "3" "o'clock" "pm" | "4" "o'clock" | "4" "o'clock" "pm" |
→"4" "o'clock" "am" | "5" "o'clock" | "5" "o'clock" "am" | "5" "o'clock" "pm",
→ | "6" "o'clock" | "6" "o'clock" "pm" | "6" "o'clock" "am" | "7" "o'clock" | _
→ | "8" "o'clock" "pm" | "9" "o'clock" | "9" "o'clock" "pm" | "9" "o'clock" |
→ "am" | "10" "o'clock" | "10" "o'clock" "am" | "10" "o'clock" "pm" | "11",
--"o'clock" | "11" "o'clock" "am" | "11" "o'clock" "pm" | "12" "o'clock" | "12" |
→"o'clock" "pm" | "12" "o'clock" "am" | 'noon' | '12' 'noon' | "12" "o'clock",
→"noon" | 'midnight' | '12' 'midnight' | 'lunch' 'time' | 'dinnertime' | |
→ | 'earliest' 'possible' | 'earliest' | 'latest' 'possible' | 'latest'
PDAY -> 'on' | 'returning' 'on' | 'of' | 'for' | 'next' | 'the' 'next' | 'in',,
_{\hookrightarrow} 'the' 'next' | 'of' 'next' | 'leaving' | 'arriving' | 'arriving' 'on' | _{\sqcup}
⇔'which' 'leave' | 'that' 'arrive' 'on' | 'leaving' 'on' | 'which' 'arrive'⊔
→'on' | 'a' 'week' 'from'
SIMPLEWEEKDAY -> 'saturday' | 'sunday' | 'monday' | 'tuesday' | 'wednesday' |
WEEKDAY -> SIMPLEWEEKDAY | A SIMPLEWEEKDAY | 'saturdays' | 'sundays' | 'mondays'
→ | 'tuesdays' | 'wednesdays' | 'thursdays' | 'fridays' | 'this' 'saturday' |
_{\hookrightarrow} 'this' 'sunday' | 'this' 'monday' | 'this' 'tuesday' | 'this' 'wednesday' | _{\sqcup}
DATE -> MONTH DAY YEAR | MONTH DAY | THE DAY | MONTH DAY OR DAY | THE DAY OR DAY
→ | EITHER MONTH DAY OR DAY | EITHER THE DAY OR THE DAY | THE DAY OF MONTH | DAY
\hookrightarrow OF MONTH | EITHER THE DAY OR THE DAY OF MONTH | THE DAY OF MONTH OR THE DAY OF_\sqcup
→MONTH | EITHER THE DAY OF MONTH OR THE DAY OF MONTH
```

```
MONTH -> 'january' | 'february' | 'march' | 'april' | 'may' | 'june' | 'july' |
→'august' | 'september' | 'october' | 'november' | 'december'
DAY -> 'first' | 'second' | 'third' | 'fourth' | 'fifth' | 'sixth' | 'seventh' |
_{\hookrightarrow}'eighth' | 'ninth' | 'tenth' | 'eleventh' | 'twelfth' | 'thirteenth' | _{\sqcup}
_{\hookrightarrow} 'fourteenth' | 'fifteenth' | 'sixteenth' | 'seventeenth' | 'eighteenth' | _{\sqcup}
_{\hookrightarrow} 'twenty-second' | 'twenty' 'second' | 'twenty-third' | 'twenty' 'third' | _{\sqcup}
_{\hookrightarrow}'twenty-sixth' | 'twenty' 'sixth' | 'twenty-seventh' | 'twenty' 'seventh' | _{\sqcup}
_{\hookrightarrow}'thirtieth' | 'thirty-first' | 'thirty' 'first' | 'one' | 'two' | 'three' | _{\sqcup}
→'twelve' | 'thirteen' | 'fourteen' | 'fifteen' | 'sixteen' | 'seventeen' |
_{\hookrightarrow} 'eighteen' | 'nineteen' | 'twenty' | 'twenty' 'one' | 'twenty' 'two' | _{\square}
→ 'one'
YEAR -> '1991' | '1992'
PAIRLINE -> 'on' | 'using' | 'of' | 'with'
AIRLINE -> 'continental' | 'continental' 'airline' | 'continental' 'airlines' |
→'american' 'airlines' | 'american' 'airlines' | 'american' | 'united' |
_{\hookrightarrow}'airlines' | 'united' | 'united' 'airline' | 'northwest' 'and' 'united' _{\sqcup}
→'northwest' 'airlines' | 'northwest' 'airline' | 'northwest' | 'us' 'air' |<sub>□</sub>
_{\hookrightarrow} 'us' 'airlines' | 'delta' | 'twa' | 'air' 'canada' | 'eastern' 'airlines' | _{\sqcup}
→'midwest' 'express' | 'trans' 'world' 'airline'
AIRCRAFT -> 'using' 'a' 'j31' 'aircraft'
FLIGHTTYPE -> 'round' 'trip' | 'round-trip' | 'nonstop' | 'one' 'way' | 'return'
→ | 'direct' | 'connecting' | 'direct' 'and' 'connecting' | 'nonstop' 'or'
FARETYPE -> 'with' 'economy' 'fares' | 'qualify' 'for' 'fare' 'code' 'qx' | |

→'with' 'first' 'class' 'service' | 'that' 'offer' 'first' 'class' | 'on'
□
_{\hookrightarrow}'fares' | 'with' 'a' 'class' 'of' 'service' 'code' 'f' | 'with' 'q' 'fares' | _{\sqcup}
PRICE -> 'less' 'than' '1100' 'dollars' | 'with' 'the' 'highest' 'fare' | 'the'
→ | 'least' 'expensive' | 'most' 'expensive' | 'cheapest'
FOOD -> 'dinner' | 'a' 'meal' | 'lunch' | 'breakfast'
```

```
AVAIL -> 'available' | 'i' 'can' 'get'
POSTJ -> 'please' | 'there' | 'are' | 'currently' | 'do' | 'you' | 'have' | ...
→ | 'be' | 'go' | 'departures' | 'is' | 'such' | 'a' | 'that' | 'serves' | ___
→'both' | 'and' | 'along' | 'with' | 'can' | 'get' | "i'd" | 'traveling' | □
11111
```

```
Solution
[]: #Solution
    miniATIS = nltk.CFG.fromstring("""
    S -> PREJ DET ADJS FLIGHT PPS | DET ADJS FLIGHT PPS | PREJ ADJS FLIGHT PPS | L
     ADJS FLIGHT PPS | PREJ DET FLIGHT PPS | DET FLIGHT PPS | PREJ FLIGHT PPS |
     →FLIGHT PPS | PREJ DET ADJS FLIGHT | DET ADJS FLIGHT | PREJ ADJS FLIGHT |
     ADJS FLIGHT | PREJ DET FLIGHT | DET FLIGHT | PREJ FLIGHT | FLIGHT
    PPS -> PP | PP PPS
    ADJS -> ADJ | ADJ ADJS
    FLIGHT -> 'flights' | 'flight' | 'to' 'fly'
    PREJ -> JUNK PREJ | JUNK
    JUNK -> 'me' | 'show' | 'now' | 'only' | 'can' | 'you' | 'the' | 'itinerary' |
     →"what's" | 'are' | 'my' | 'choices' | 'for' | 'i' | 'would' | 'like' | "i'd"□
     → | 'to' | 'see' | 'have' | 'make' | 'book' | 'find' | 'information' | 'on' | |
     \rightarrow'know' | 'some' | 'hello' | 'yes' | 'please' | 'repeat' | 'do' | 'have' |

¬'arrangements' | 'okay' | 'want' | 'tell' | 'about' | 'how' | 'would' | 'be'

     → | 'able' | 'put' | 'requesting' | "i'm" | 'looking' | 'display' | UNK
    ADJ -> AIRLINE | SIMPLEWEEKDAY | FLIGHTTYPE | SIMPLETIME | "monday's" | |
     _{\hookrightarrow}"tuesday's" | "wednesday's" | "thursday's" | "friday's" | "saturday's" | _{\sqcup}
     -- "sunday's" | 'available' | 'possible' | 'first' 'class' | 'economy' | ...
     _{\hookrightarrow}'thrift' 'economy' | 'cheapest' | 'lowest' 'cost' | 'least' 'expensive' | _{\sqcup}

¬'most' 'expensive' | 'weekday' | 'daily' | 'last' | 'first' | 'dinner' |

     DET -> 'all' 'the' | 'all' | A | 'an' | THE | 'any' | 'all' 'of' 'the'
    BETWEEN -> 'between'
    AND -> 'and'
    OR -> 'or'
    EITHER -> 'either'
    OF -> 'of'
    THE -> 'the'
    A -> 'a'
```

```
PP -> PPLACE PLACE OR PLACE | PPLACE EITHER PLACE OR PLACE | PPLACE PLACE |
 →EITHER PLACE OR PLACE | PLACE OR PLACE | BETWEEN PLACE AND PLACE |
 →BETWEEN TIME AND TIME | PTIME TIME | TIME | PDAY WEEKDAY | WEEKDAY TIME |
 →WEEKDAY | PDAY WEEKDAY TIME | PDAY DATE | DATE | PAIRLINE AIRLINE | AIRCRAFT |
 \hookrightarrow | FLIGHTTYPE | FARETYPE | PRICE | FOOD | AVAIL | POSTJ
PPLACE -> 'to' | 'that' 'arrive' 'at' | 'that' 'arrives' 'in' | 'coming' 'back'
 _{\hookrightarrow} 'to' | 'that' 'go' 'to' | 'and' 'then' 'to' | 'arriving' 'in' | 'and' _{\square}
 →'going' 'to' | 'into' | 'for' | 'with' 'the' 'destination' 'city' 'of' |
 →'arriving' | 'goes' 'to' | 'flying' 'into' | 'goes' 'on' 'to' | 'reaching' |<sub>□</sub>
 →'in' | 'and' 'then' | 'arriving' 'to' | 'from' | 'leaving' | 'return' 'from'
 _{\hookrightarrow}| 'leaving' 'from' | 'departing' | 'go' 'from' | 'going' _{\sqcup}
 →'from' | 'back' 'from' | 'that' 'goes' 'from' | 'that' 'departs' | 'which'
 _{\hookrightarrow} 'leaves' 'from' | 'which' 'leave' | 'that' 'leave' | 'originating' 'in' | _{\sqcup}
 →'a' 'stopover' 'in' | 'with' 'a' 'layover' 'in' | 'with' 'a' 'stopover' 'at'□
 → | 'and' 'a' 'stopover' 'in' | 'stop' 'in' | 'stopping' 'in' | 'make' 'a'

→ 'through' | 'which' 'go' 'through' | 'makes' 'a' 'stopover' 'in' | 'that'
□

→ 'stops' 'in' | 'that' 'stops' 'over' 'in' | 'by' 'way' 'of' | 'connecting'||

| The stops of the stop of the stops of the stops of the stops of the stops of the stop 
 PLACE -> 'anywhere' | 'atlanta' | 'austin' | 'baltimore' | 'boston' | 'boston' |
 - 'logan' | 'burbank' | 'bwi' | 'charlotte' | 'chicago' | 'cincinnati' | ...

→'cleveland' | 'cleveland' 'ohio' | 'columbus' | 'dallas' 'fort' 'worth' |
□

→ 'dallas' | 'denver' | 'denver' 'colorado' | 'detroit' | 'fort' 'worth' |
□

→'general' 'mitchell' 'international' | 'general' 'mitchell' | 'houston' |
□
 →'indianapolis' | 'jfk' | 'kansas' 'city' | 'laguardia' 'airport' | 'las'
 →'vegas' | 'long' 'beach' | 'los' 'angeles' | 'love' 'field' | 'memphis' |
 \hookrightarrow 'miami' | 'milwaukee' | 'minneapolis' | 'montreal' | 'montreal' 'quebec' |
 → 'montreal' 'canada' | 'nashville' | 'new' 'york' 'city' | 'new' 'york' |
 →'newark' | 'newark' 'new' 'jersey' | 'oakland' | 'oakland' 'california' | □
 -- 'ontario' | 'orlando' | 'orlando' 'florida' | 'philadelphia' | 'philly' |

→ 'phoenix' | 'pittsburgh' | 'salt' 'lake' 'city' | 'san' 'diego' | 'san' |

□
 - diego' 'california' | 'san' 'francisco' | 'san' 'jose' | 'seattle' | 'st.'⊔
 →'louis' | 'st.' 'paul' | 'st.' 'petersburg' | 'tacoma' | 'tacoma' |

→ 'washington' | 'tampa' | 'toronto' | 'various' 'cities' | 'washington' |
□
```

```
PTIME -> 'that' 'arrive' 'before' | 'that' 'arrives' 'before' | 'arriving'
_{\hookrightarrow}'before' | 'arrival' 'by' | 'arrives' | 'before' | 'departing' 'before' | _{\sqcup}
→'around' | 'arriving' 'about' | 'that' 'arrive' 'soon' 'after' | 'leaving'
_{\hookrightarrow}'at' | 'leaving' | 'which' 'leave' 'after' | 'leaving' 'after' | 'after' | _{\sqcup}
→'departing' 'after' | 'that' 'depart' 'after' | 'departing' 'at' | 'arriving'
TIME -> SIMPLETIME | 'the' 'afternoon' | 'the' 'late' 'afternoon' | 'the'
_{\hookrightarrow}'evening' | 'the' 'morning' | 'the' 'early' 'am' | 'the' 'day' | 'mornings' | _{\sqcup}
_{\hookrightarrow} 'afternoons' | 'evenings' | 'as' 'early' 'as' 'possible' | 'earliest' _{\sqcup}
→'possible' 'time' | 'as' 'soon' 'thereafter' 'as' 'possible' | UNK
SIMPLETIME -> 'one' | 'two' | 'three' | 'four' | 'five' | 'six' | 'seven' |
_{\hookrightarrow}'eight' | 'nine' | 'ten' | 'eleven' | 'twelve' | '1' | '2' | '3' |'4' | '5' | _{\sqcup}
_{\hookrightarrow} '6' | '7' | '8' | '9' | '10' | '11' | '12' | '1pm' | '2pm' | '3pm' | '4pm' | _{\sqcup}
-- '5pm' | '6pm' | '7pm' | '8pm' | '9pm' | '10pm' | '11pm' | '12pm' | '1am' | □
→'2am' | '3am' | '4am' | '5am' | '6am' | '7am' | '8am' | '9am' | '10am' |
→'11am' | '12am' | '230' | '1505' | '630am' | '720am' | '723pm' | '819' | '845'ц
→ | '1026' | '1145am' | '1' "o'clock" | '1' "o'clock" 'am' | "1" "o'clock" "pm" |
\hookrightarrow | "2" "o'clock" | "2" "o'clock" "am" | "2" "o'clock" "pm" | "3" "o'clock" |
→"3" "o'clock" "am" | "3" "o'clock" "pm" | "4" "o'clock" | "4" "o'clock" "pm" |
→"4" "o'clock" "am" | "5" "o'clock" | "5" "o'clock" "am" | "5" "o'clock" "pm" |
→ | "6" "o'clock" | "6" "o'clock" "pm" | "6" "o'clock" "am" | "7" "o'clock" | _
→"7" "o'clock" "pm" | "7" "o'clock" "am" | "8" "o'clock" | "8" "o'clock" "am" |
→ | "8" "o'clock" "pm" | "9" "o'clock" | "9" "o'clock" "pm" | "9" "o'clock"
→ "am" | "10" "o'clock" | "10" "o'clock" "am" | "10" "o'clock" "pm" | "11"
-- "o'clock" "pm" | "12" "o'clock" "am" | 'noon' | '12' 'noon' | "12" "o'clock" ⊔

→ "noon" | 'midnight' | '12' 'midnight' | 'lunch' 'time' | 'dinnertime' | □
\hookrightarrow | 'earliest' 'possible' | 'earliest' | 'latest' 'possible' | 'latest'
PDAY -> 'on' | 'returning' 'on' | 'of' | 'for' | 'next' | 'the' 'next' | 'in'

→'which' 'leave' | 'that' 'arrive' 'on' | 'leaving' 'on' | 'which' 'arrive'
□
→'on' | 'a' 'week' 'from'
SIMPLEWEEKDAY -> 'saturday' | 'sunday' | 'monday' | 'tuesday' | 'wednesday' |
```

```
WEEKDAY -> SIMPLEWEEKDAY | A SIMPLEWEEKDAY | 'saturdays' | 'sundays' | 'mondays'
→ | 'tuesdays' | 'wednesdays' | 'thursdays' | 'fridays' | 'this' 'saturday' |
→'coming' 'sunday' | 'this' 'coming' 'monday' | 'this' 'coming' 'tuesday' | □
→'this' 'coming' 'wednesday' | 'this' 'coming' 'thursday' | 'this' 'coming'
- 'friday' | 'day' | 'week' | 'today' | 'tomorrow' | 'the' 'day' 'after'
→'tomorrow' | 'a' 'weekday' | 'weekdays'
DATE -> MONTH DAY YEAR | MONTH DAY | THE DAY | MONTH DAY OR DAY | THE DAY OR DAY
→ | EITHER MONTH DAY OR DAY | EITHER THE DAY OR THE DAY | THE DAY OF MONTH | DAY
→OF MONTH | EITHER THE DAY OR THE DAY OF MONTH | THE DAY OF MONTH OR THE DAY OF L
→MONTH | EITHER THE DAY OF MONTH OR THE DAY OF MONTH
MONTH -> 'january' | 'february' | 'march' | 'april' | 'may' | 'june' | 'july' |
→'august' | 'september' | 'october' | 'november' | 'december'
DAY -> 'first' | 'second' | 'third' | 'fourth' | 'fifth' | 'sixth' | 'seventh' |
_{\hookrightarrow}'eighth' | 'ninth' | 'tenth' | 'eleventh' | 'twelfth' | 'thirteenth' | _{\sqcup}
→'fourteenth' | 'fifteenth' | 'sixteenth' | 'seventeenth' | 'eighteenth' |
_{\hookrightarrow}'nineteenth' | 'twentieth' | 'twenty-first' | 'twenty' 'first' | _{\sqcup}
_{\hookrightarrow}'twenty-second' | 'twenty' 'second' | 'twenty-third' | 'twenty' 'third' | _{\sqcup}
_{\hookrightarrow}'twenty-fourth' | 'twenty' 'fourth' | 'twenty-fifth' | 'twenty' 'fifth' | _{\sqcup}
_{\hookrightarrow}'twenty-sixth' | 'twenty' 'sixth' | 'twenty-seventh' | 'twenty' 'seventh' | _{\sqcup}
→'thirtieth' | 'thirty-first' | 'thirty' 'first' | 'one' | 'two' | 'three' |<sub>□</sub>
_{\hookrightarrow}'twelve' | 'thirteen' | 'fourteen' | 'fifteen' | 'sixteen' | 'seventeen' | _{\sqcup}
- 'twenty' 'seven' | 'twenty' 'eight' | 'twenty' 'nine' | 'thirty' | 'thirty'

    one¹

YEAR -> '1991' | '1992'
PAIRLINE -> 'on' | 'using' | 'of' | 'with'
AIRLINE -> 'continental' | 'continental' 'airline' | 'continental' 'airlines' | |
_{\rm \hookrightarrow}'american' 'airlines' | 'american' 'airlines' | 'american' | 'united' _{\rm \sqcup}
→'airlines' | 'united' | 'united' | 'airline' | 'northwest' 'and' 'united' |
_{\hookrightarrow} 'northwest' 'airlines' | 'northwest' 'airline' | 'northwest' | 'us' 'air' | _{\sqcup}

→ 'us' 'airlines' | 'delta' | 'twa' | 'air' 'canada' | 'eastern' 'airlines' |
□
→'midwest' 'express' | 'trans' 'world' 'airline'
AIRCRAFT -> 'using' 'a' 'j31' 'aircraft'
FLIGHTTYPE -> 'round' 'trip' | 'round-trip' | 'nonstop' | 'one' 'way' | 'return'
→ | 'direct' | 'connecting' | 'direct' 'and' 'connecting' | 'nonstop' 'or' |
```

```
FARETYPE -> 'with' 'economy' 'fares' | 'qualify' 'for' 'fare' 'code' 'qx' |
→'with' 'first' 'class' 'service' | 'that' 'offer' 'first' 'class' | 'on'
_{\hookrightarrow}'fares' | 'with' 'a' 'class' 'of' 'service' 'code' 'f' | 'with' 'q' 'fares' |_{\sqcup}
PRICE -> 'less' 'than' '1100' 'dollars' | 'with' 'the' 'highest' 'fare' | 'the' _{\sqcup}
→ | 'least' 'expensive' | 'most' 'expensive' | 'cheapest'
FOOD -> 'dinner' | 'a' 'meal' | 'lunch' | 'breakfast'
AVAIL -> 'available' | 'i' 'can' 'get'
POSTJ -> 'please' | 'there' | 'are' | 'currently' | 'do' | 'you' | 'have' |
_{\hookrightarrow}'fares' | 'information' | 'i' | 'want' | 'would' | 'like' | 'the' | 'flight'
→ | 'be' | 'go' | 'departures' | 'is' | 'such' | 'a' | 'that' | 'serves' | □
_{\hookrightarrow} 'both' | 'and' | 'along' | 'with' | 'can' | 'get' | "i'd" | 'traveling' | _{\sqcup}
→'for' | 'me' | UNK | '.' | '?'
11111)
```

To iterate over the constructed grammar, we can use grammar.productions method.

```
[]: for production in miniATIS.productions():
    lhs = production.lhs() # left hand side
    rhs = production.rhs() # right hand side
    print (f'Rule: {lhs} -> {rhs}')
    print (f'{len(rhs)}')
    break
```

To convert the constructed grammar to CNF, we use get cnf grammar.

```
[ ]: normal_miniATIS, miniATIS_wunaries = get_cnf_grammar(miniATIS)
print (normal_miniATIS.is_chomsky_normal_form())
```

Now, we can use normal\_miniATIS for CKY recognition and parsing.

```
[]: for production in normal_miniATIS.productions():
    lhs = production.lhs() # left hand side
    rhs = production.rhs() # right hand side
    print (f'Rule: {lhs} -> {rhs}')
    print (f'{len(rhs)}')
    break
```

# 1.4 Deterministic parsing

# 1.4.1 Goal 2: CKY Recognition

Implement a recognizer using the CKY algorithm to determine if a sentence is parsable. See J&M Chapter 13 for notes on CKY and the pseudo-code of the algorithm.

```
def cky_recognize(grammar, s):
    """
    CKY algorithm for recognition.
    Arguments:
        grammar: a CFG in CNF
        s: the input string to parse
    Returns whether sentence is parsable
    """
    assert(grammar.is_chomsky_normal_form())

# TODO: Implement a CKY Recognizer
    return False
```

```
[]: #Solution
     def cky_recognize(grammar, s):
       CKY algorithm for recognition.
       Arguments:
             grammar: a CFG in CNF
             s: the input string to parse
       Returns whether sentence is parsable
       assert(grammar.is_chomsky_normal_form())
       nonterminals = set([])
       # Get the nonterminals
       for production in grammar.productions():
         nonterminals.add(production.lhs())
         assert(type(production.lhs()) == Nonterminal)
      num_nonterminals = len(nonterminals)
      nonterminals = list(nonterminals)
       # Setup lookup table
       dynamic_table = {}
       # Lookup table initialization
```

```
for i in range(len(s)):
  dynamic_table[i] = {}
  for j in range(len(s), i, -1):
    dynamic_table[i][j] = {}
    for k in nonterminals:
      dynamic_table[i][j][k] = False
 # Iterate over each token
for j in range(1, len(s)+1):
  # Unary productions of the form A -> token[j-1]
  \rightarrows[j-1], grammar.productions()):
    dynamic_table[j-1][j][production.lhs()] = True
  # Spans from size j-2 to 0
  for i in range(j-2, -1, -1):
    # Iterate over possible midpoints
    for k in range(i+1, j):
      for production in list(filter(lambda p : len(p.rhs()) == 2,grammar.
→productions())):
        if dynamic_table[i][k][production.rhs()[0]] and__

→dynamic_table[k][j][production.rhs()[1]]:
          dynamic_table[i][j][production.lhs()] = True
 # Is the sentence in the grammar?
return dynamic_table[0][len(s)][grammar.start()]
```

You can use the following test sentences to test your code:

```
[]: # A grammartical sentence

test_sentence = "show me the flights on united that arrive before 4 between_

→atlanta and new york".split()

result = cky_recognize(normal_miniATIS, test_sentence)

print (result)

# An ungrammatical sentence

test_sentence = "show me the flights on united that between atlanta and new york

→arrive before 4 ".split()

result = cky_recognize(normal_miniATIS, test_sentence)

print (result)
```

## 1.4.2 Goal 3: CKY Parsing

Implement the CKY algorithm for parsing under CFGs. You only need to add a few lines of code to your CKY recognizer to achieve this, using back-pointers. The expected return value is an NLTK

tree, which can be constructed using Tree.fromstring.

A standard tree string will be of the following form:

```
(S (A B) (C (D E) (F G)))
```

which corresponds to the following tree (drawn using tree.pretty\_print()):

```
S

---| C

| C

A D F

| | |
```

```
def cky_parse(grammar, s):
    """
    CKY algorithm for parsing.
    Arguments:
        grammar: a CFG in CNF
        s: the input string to parse
    Returns an nltk Tree if parsable or None if not parsable
    """
    assert(grammar.is_chomsky_normal_form())

# TODO: Implement a CKY Parser
    return Tree.fromstring('(S (A B) (C (D E) (F G)))')
```

```
[]: #Solution
def cky_parse(grammar, s):
    """
    CKY algorithm for parsing.
    Arguments:
        grammar: a CFG in CNF
        s: the input string to parse
    Returns an nltk Tree if parsable or None if not parsable
    """
    assert(grammar.is_chomsky_normal_form())
    nonterminals = set([])
    # Get the nonterminals
    for production in grammar.productions():
        nonterminals.add(production.lhs())
        assert(type(production.lhs()) == Nonterminal)
    num_nonterminals = len(nonterminals)
```

```
nonterminals = list(nonterminals)
 # Setup lookup table
 dynamic_table = {}
 # bottom up parse tree
 tree = {}
 # Initialization of lookup table
 for i in range(len(s)):
   dynamic table[i] = {}
   tree[i] = {}
   for j in range(len(s), i, -1):
     dynamic_table[i][j] = {}
     tree[i][j] = {}
     for k in nonterminals:
       dynamic_table[i][j][k] = False
       tree[i][j][k] = ''
 # For each token
 for j in range(1, len(s)+1):
   # Unit productions of the form A -> token[j - 1]
   for production in list(filter(lambda p : len(p.rhs()) == 1 and p.rhs()[0] ==_u
\rightarrows[j-1], grammar.productions()):
     dynamic_table[j-1][j][production.lhs()] = True
     tree[j-1][j][production.lhs()] = '(' + production.lhs()._str_() + ' ' +_{\square}
→s[j-1] + ')'
   # For spans from size j-2 to 0
   for i in range(j-2, -1, -1):
     # Iterate over possible midpoints
     for k in range(i+1, j):
       # Binary productions
       for production in list(filter(lambda p : len(p.rhs()) == 2,grammar.
→productions())):
         if dynamic_table[i][k][production.rhs()[0]] and__
→dynamic_table[k][j][production.rhs()[1]]:
           dynamic_table[i][j][production.lhs()] = True
           tree[i][j][production.lhs()] = '(' + production.lhs()._str__() + 'u
\rightarrow' + tree[i][k][production.rhs()[0]] + ' ' + tree[k][j][production.rhs()[1]] +
\hookrightarrow 1) 1
 # If the sentence is in the grammar return the constructed tree
 if dynamic_table[0][len(s)][grammar.start()]:
   print (tree[0][len(s)][grammar.start()])
   return Tree.fromstring(tree[0][len(s)][grammar.start()])
```

```
else:
return None
```

You can use the following test sentence to test your code:

You can also compare against a built-in nltk parser:

```
[]: parser = nltk.parse.BottomUpChartParser(miniATIS)
    parses = [p for p in parser.parse(test_sentence)]
    print('Reference parse:')
    for ptree in parses:
        print(ptree)

print ('Predicted parse:')
    print (tree)

print("Parses match (at least one):", tree in parses)
```

Grammar Coverage To check your implementation, let's compute the fraction of the original ATIS dataset that parses under the ATIS grammar. This grammar should cover 2116 out of the 4379 sentences.

```
[]: # This script might take up to 30min.
    # We recommend to make sure the above checks work first.
    all_sents_big_ATIS = []
    with open('data/train.nl') as f:
        for line in f:
            all_sents_big_ATIS.append(line.split())

parsed = 0
    total = 0
    for sent in all_sents_big_ATIS:
        total += 1
        sent = [tok.lower() for tok in sent]

    tree = cky_parse(normal_miniATIS, sent)
    if tree:
        parsed += 1

print(parsed, total)
```

## 1.5 Probabilistic parsing

In practice, we want to work with grammars that cover nearly all the language we expect to come across for a given application. This leads to an explosion of rules and a large number of possible parses for any one sentence. To remove ambiguity between the different parses, it's desirable to move to PCFGs. In this part, you will construct a PCFG from data, parse using a probabilistic version of CKY, and evaluate the quality of the resulting parses against gold trees.

#### 1.5.1 Goal 4: PCFG construction

Compared to CFGs, PCFGs need to assign probabilities to grammar rules. Fill in the missing pieces of the below method to build a PCFG in CNF by counting using the *train* split (train.trees that we downloaded in Setup). Note that we want the PCFG to be in CNF format because the probabalistic version of CKY also requires the grammar to be in CNF. However, the gold trees are not in CNF form, so in this case you will need to convert the gold *trees* to CNF before building the PCFG. To accomplish this, we will use the treetransforms package from nltk, which includes functions for converting to and from CNF.

Finish the below code to construct a PCFG grammar. You will use it later for parsing.

```
[ ]: #TODO
     def construct_pcfg_rules(gold_trees):
       Construct PCFG grammar rules from gold parse trees.
       Arguments:
             gold_trees: a list of gold parse trees
       Returns a set of ProbabilisticProduction objects constructed by
             ProbabilisticProduction(lhs, rhs, prob=prob)
       productions with probs = set([])
      pcfg_data = defaultdict(Counter)
       for tree in gold_trees:
         # Convert to CNF
         treetransforms.collapse_unary(tree, collapsePOS=True)
         treetransforms.chomsky_normal_form(tree)
         # TODO: implement here
         # Construct a PCFG by counting rule frequencies from these trees
         # Hint: You can get the list of productions from the tree via tree.
      →productions()
       # TODO: implement here
       # Compute probabilities using the counts
       return productions_with_probs
```

```
[]: #Solution
     def construct_pcfg_rules(gold_trees):
       Construct PCFG grammar rules from gold parse trees.
       Arguments:
             gold_trees: a list of gold parse trees
       Returns a set of ProbabilisticProduction objects constructed by
             ProbabilisticProduction(lhs, rhs, prob=prob)
       productions_with_probs = set([])
      pcfg_data = defaultdict(Counter)
       for tree in gold trees:
         # Convert to CNF
         treetransforms.collapse_unary(tree, collapsePOS=True)
         treetransforms.chomsky_normal_form(tree)
         # Count rule frequencies
         for production in tree.productions():
           pcfg_data[production.lhs()][production.rhs()] += 1
       # Compute probabilities
       for lhs, rhs_counter in pcfg_data.items():
         total count = sum(rhs counter.values())
         for rhs, count in rhs_counter.items():
           prob = count/total count
           productions_with_probs.add(ProbabilisticProduction(lhs, rhs, prob=prob))
       return productions with probs
```

```
[]: # Load gold trees
gold_trees = []
with open('data/train.trees') as f:
    for line in f:
        gold_trees.append(nltk.Tree.fromstring(line.strip()))

productions_with_probs = construct_pcfg_rules(gold_trees)
pgrammar = PCFG(Nonterminal('TOP'), productions_with_probs)
print (pgrammar.is_chomsky_normal_form())
```

# 1.5.2 Goal 5: Probabalistic CKY parsing

Finally, we are ready to implement CKY parsing under PCFGs. Adapt the CKY parser from Goal 3 to return the most likely parse and its probability given a PCFG.

```
[]: #TODO

def cky_parse_probabalistic(grammar, s):
    """

CKY algorithm for parsing using PCFG.
```

```
Arguments:
    grammar: a PCFG in CNF
    s: the input string to parse
Returns twov values:
    first: the most likely nltk Tree if parsable or None if not parsable second: the highest probability
"""
assert(grammar.is_chomsky_normal_form())

# TODO: Implement a probabilisite CKY parser
return Tree.fromstring('(S (A B) (C (D E) (F G)))'), 0
```

```
[]: #Solution
     def cky_parse_probabalistic(grammar, s):
       CKY algorithm for parsing using PCFG.
       Arguments:
             grammar: a PCFG in CNF
             s: the input string to parse
       Returns twov values:
             first: the most likely nltk Tree if parsable or None if not parsable
             second: the highest probability
       .....
       assert(grammar.is_chomsky_normal_form())
       nonterminals = set([])
       # Get the nonterminals
      for production in grammar.productions():
         nonterminals.add(production.lhs())
         assert(type(production.lhs()) == Nonterminal)
       num_nonterminals = len(nonterminals)
       nonterminals = list(nonterminals)
       # Setup lookup table
       dynamic_table = {}
       # bottom up parse tree
       tree = {}
       # Dynamic table initialization
      for i in range(len(s)):
         dynamic_table[i] = {}
         tree[i] = {}
         for j in range(len(s), i, -1):
```

```
dynamic_table[i][j] = {}
     tree[i][j] = {}
     for k in nonterminals:
       dynamic_table[i][j][k] = None
       tree[i][j][k] = ''
 # Iterate over each token
for j in range(1, len(s)+1):
   # Unit productions of the form A -> tokens[j-1]
   for production in list(filter(lambda p : len(p.rhs()) == 1 and p.rhs()[0] ==_{\sqcup}
\rightarrows[j-1], grammar.productions()):
     dynamic_table[j-1][j][production.lhs()] = production.logprob()
     tree[j-1][j][production.lhs()] = '(' + production.lhs()._str_() + ' ' +_

→s[j-1] + ')'
   # Span length from j-2 to 0
   for i in range(j-2, -1, -1):
     # Iterate over all midpoints
     for k in range(i+1, j):
       # Iterate over each binary production
       for production in list(filter(lambda p : len(p.rhs()) == 2,grammar.
→productions())):
         A = dynamic_table[i][j][production.lhs()]
         B = dynamic_table[i][k][production.rhs()[0]]
         C = dynamic_table[k][j][production.rhs()[1]]
         if B is not None and C is not None and (A is None or B + C +
→production.logprob() > A):
           dynamic_table[i][j][production.lhs()] = B + C + production.logprob()
           tree[i][j][production.lhs()] = '(' + production.lhs()._str_() + 'u
\rightarrow' + tree[i][k][production.rhs()[0]] + ' ' + tree[k][j][production.rhs()[1]] + \Box
\hookrightarrow 1) 1
 # If the sentence is in the grammar return the constructed tree
if dynamic_table[0][len(s)][grammar.start()] is not None:
   return Tree.fromstring(tree[0][len(s)][grammar.start()]),
→2**dynamic_table[0][len(s)][grammar.start()]
   return None, None
```

Evaluation There are a number of ways to evaluate parsing algorithms. In this assignment we will use the bracket scoring program evalb which we have downloaded during setup. You are expected to achieve precision of 0.67, recall of 0.41 and F1 of 0.51.

Read in the test data:

```
with open('data/test.trees') as f:
   for line in f:
     t = Tree.fromstring(line.strip())
     all_sents.append(t.leaves())
```

Parse the test sentences using your implementation, and write output to a file:

```
[]: trees_out = []
for sent in all_sents:
    tree, _ = cky_parse_probabalistic(pgrammar, sent)
    if tree is not None:
        trees_out.append(tree.pformat(margin=9999999999))
    else:
        trees_out.append('()')

with open('outp.trees', 'w') as f:
    for line in trees_out:
        f.write(line + '\n')
```

Compare the predicted trees to the ground truth trees.

```
[]: !python scripts/evalb.py outp.trees data/test.trees
```