(Un)linkable Pseudonyms for Governmental Databases

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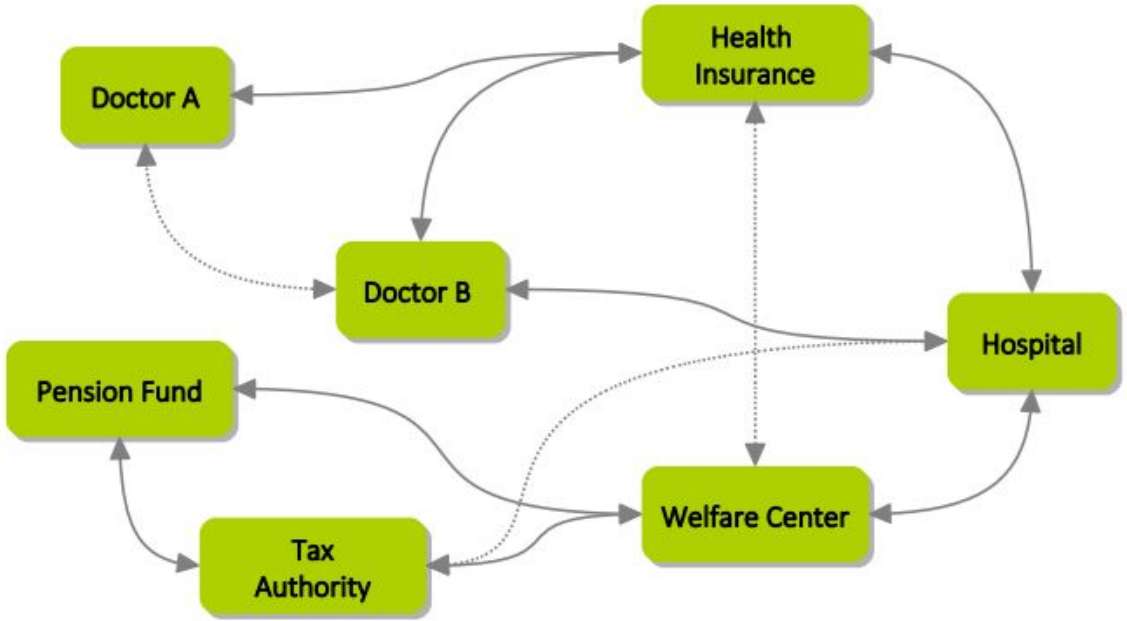
Introduction

Decentralised system:

- Large data is distributed over several databases and organisations
- Eventually data needs to be exchanged or co-related



Global Identifier

Unique Identifier for each user

Advantages:

- Allows all entities to easily share and link data records

Disadvantages:

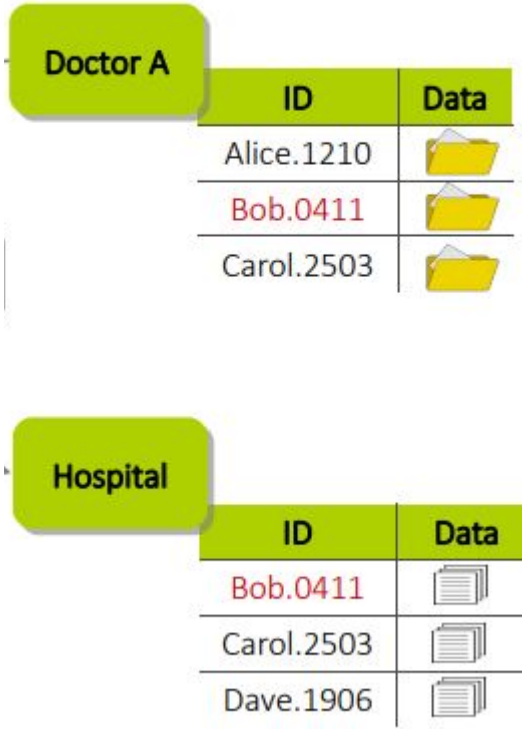
- Significant Privacy threat (In Data breach)
- Data can hardly be controlled & authorized

Solution:

- A certain control to limit dataflow(Central Authority)
- Data exchange need to authorized by central authority

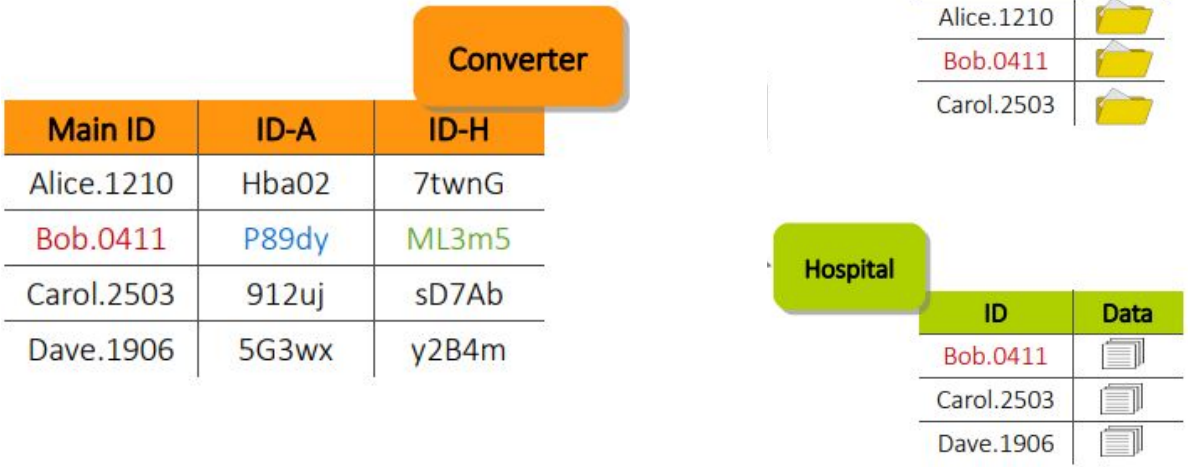
Drawback:

- Central Authority knows requests
- Can reveal sensitive information



Pseudonym

- User data is associated with (unlinkable) server-local identifiers aka “pseudonyms”
- Only converter can link & convert pseudonyms ➜ central hub for data exchange



Drawback:

➔ Data sets maintained by the entities do not contain other unique identifying information which allows linkage without using the pseudonyms

➔ Converter still needs to be trusted(Learns from requests & knows all co-relations)

Existing Solutions

Use block cipher for encryption of unique identifier

$$P_A = Enc(K_A, uid_i) \quad K_A, K_B, K_C \text{ .... are server keys known only to converter}$$

Converter

Main ID	ID-A	ID-H
Alice.1210	Hba02	7twnG
Bob.0411	P89dy	ML3m5
Carol.2503	912uj	sD7Ab
Dave.1906	5G3wx	y2B4m

- ❑ All keys are different
- ❑ Pseudonyms are unlinkable

Drawback:

- ➔ Converter still needs to be trusted
- ➔ Indirect identification by profession, age etc

Converter

Main ID	ID-A	ID-H
Alice.1210	Hba02	Hba02
Bob.0411	P89dy	P89dy
Carol.2503	912uj	912uj
Dave.1906	5G3wx	5G3wx

- ❑ If all keys are same (K_A=K_B=K_C)
- ❑ Pseudonyms are linkable
- ❑ Protocol fails

Problem Statement

Main aim : Unlinkable pseudonyms but without trusted converter
Solution: Converter and server both contribute to the derivation of pseudonym

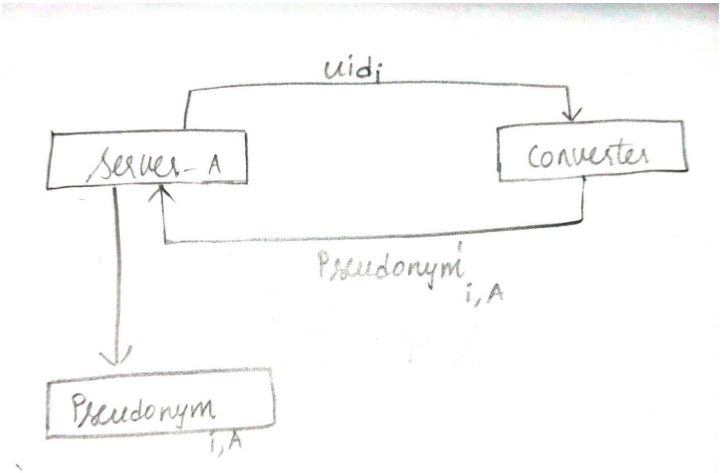
Advantages:

- Control about data exchange
- If records are lost, pieces cannot be linked together without the converter
- Converter cannot tell if requests are for the same pseudonym or not & Knows there's a data request from S_A to S_B

Protocol

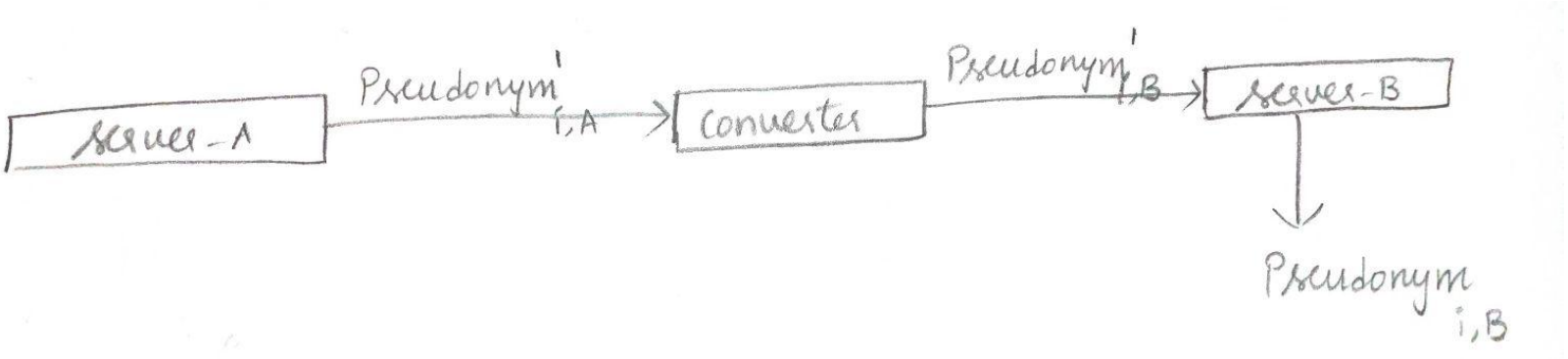
- Pseudonym Generation
- Conversion Request
- Conversion Response

Pseudonym Generation



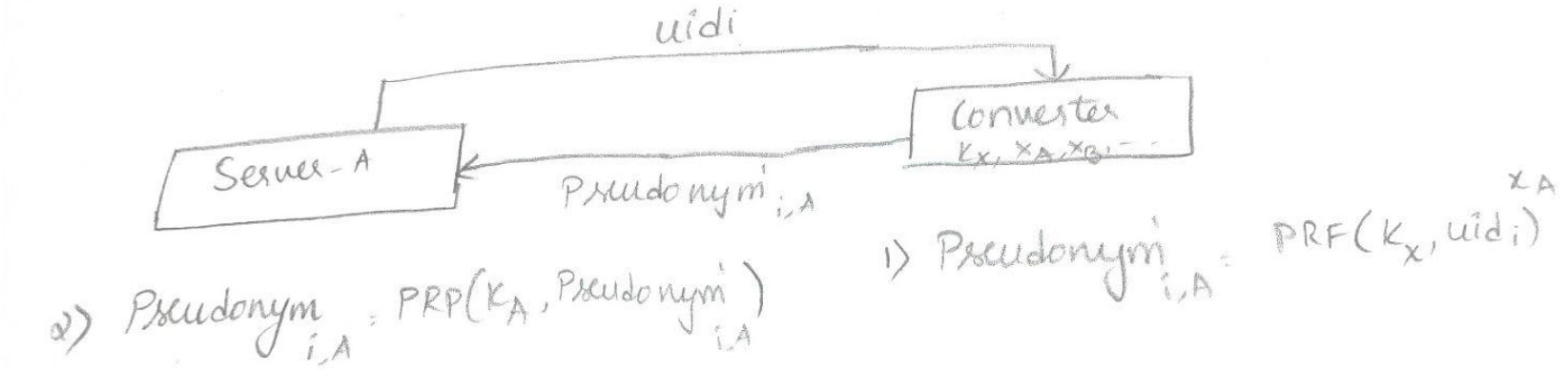
- Converter will use different keys to different servers corresponding to uid_i so that pseudonyms are unlinkable

Conversion Request & Response

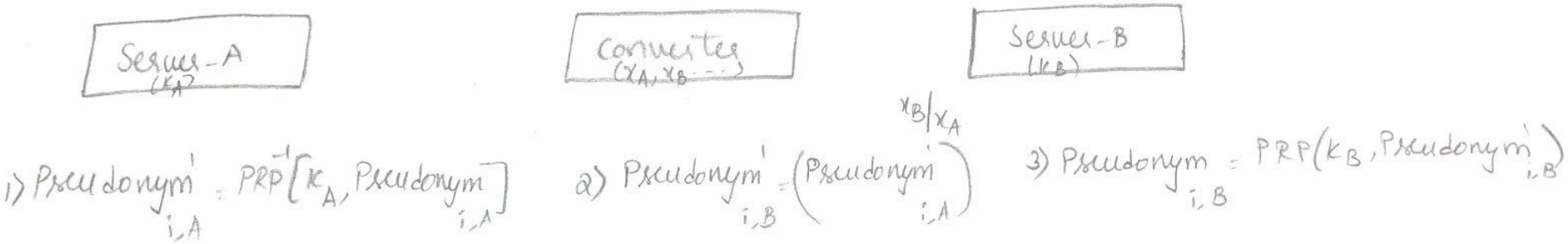


- S_A wants some information from S_B
- We convert $Pseudonym_{i,A}$ to $Pseudonym_{i,B}$ with the help of S_A,Converter & S_B

Pseudonym Generation



Pseudonym Conversion



Drawbacks

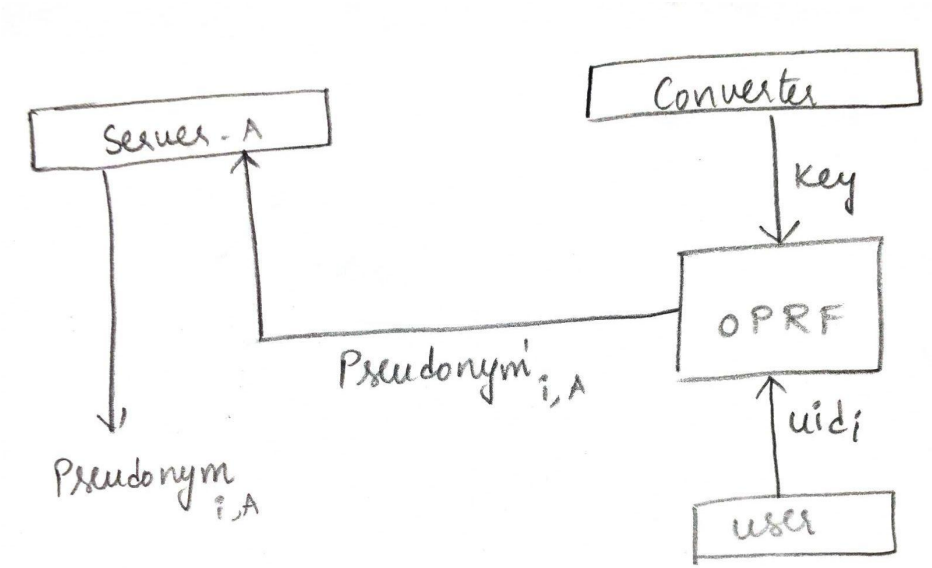
- Server knows uid_i at the generation process.
- A corrupted converter and a corrupted server can link pseudonyms

Solution:-

- To involve user in the process of conversion as well as generation of pseudonym
- To not fully involve converter in the process of conversion from $Pseudonym_A$ to $Pseudonym_B$

Proposed Solution-1

Pseudonym Generation:-

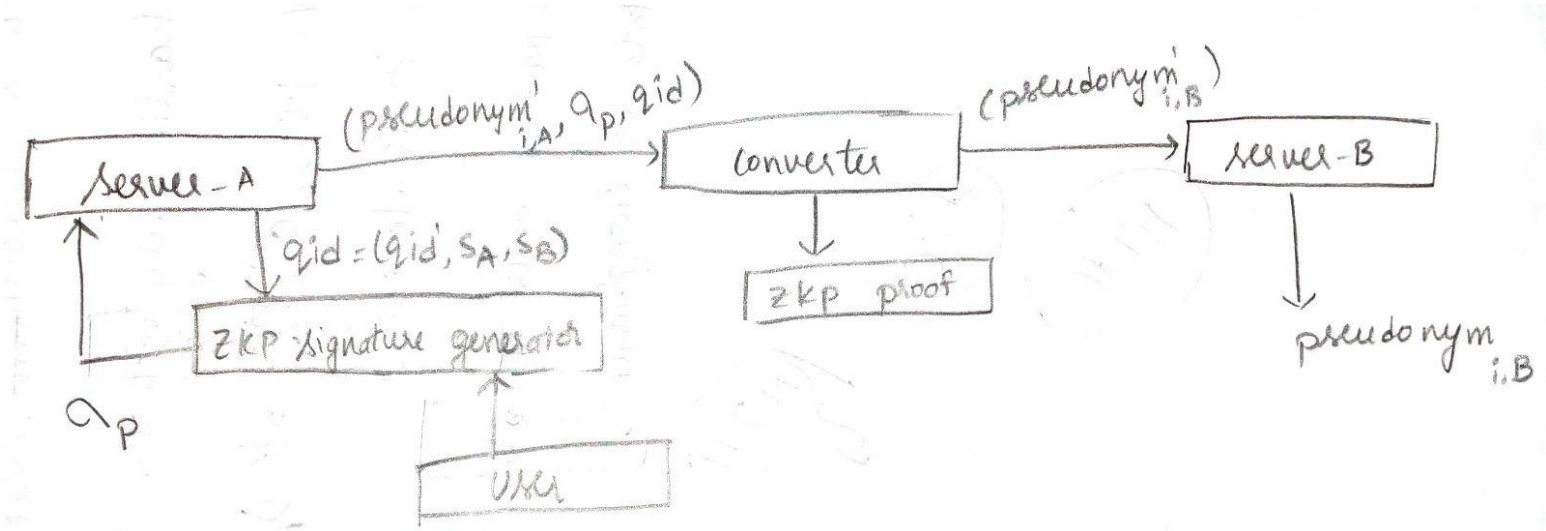


Oblivious Pseudorandom function(OPRF) generates output of PRF(m) without knowing message m to converter

Oblivious Pseudorandom functions

- PRF $f_k(m)=g^{1/(k+m)}$
- Encryption scheme additively homomorphic on message domain Z_n
- Converter blindly computes $z_i=PRF_G(k,uid_i)$
- User initiates pseudonym generation unlike previous where server triggers pseudonym generation

Conversion



- ZKP Signature generator will output user's signature on query identifier(qid)
- Converter will give zkp that it indeed verified the qid and converted the pseudonym honestly

Drawback:-

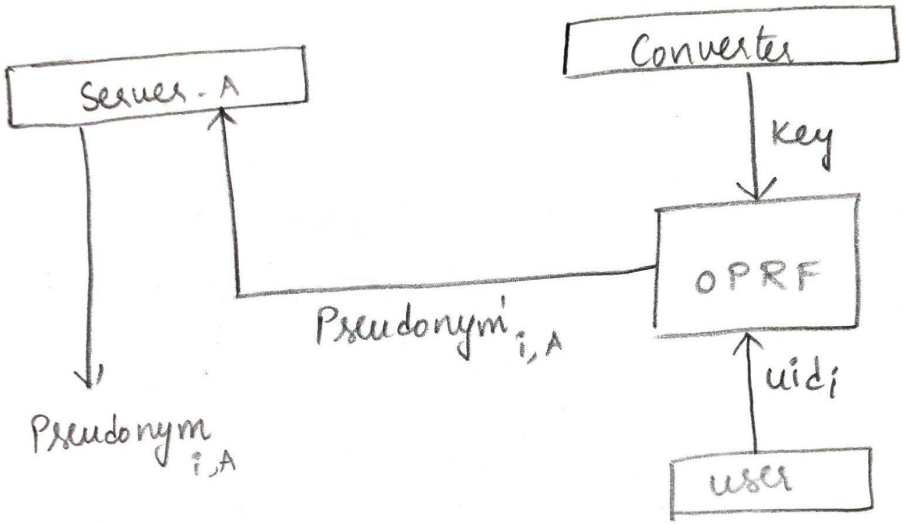
User Needs to involve in every conversion of pseudonym

Anonymous Credentials

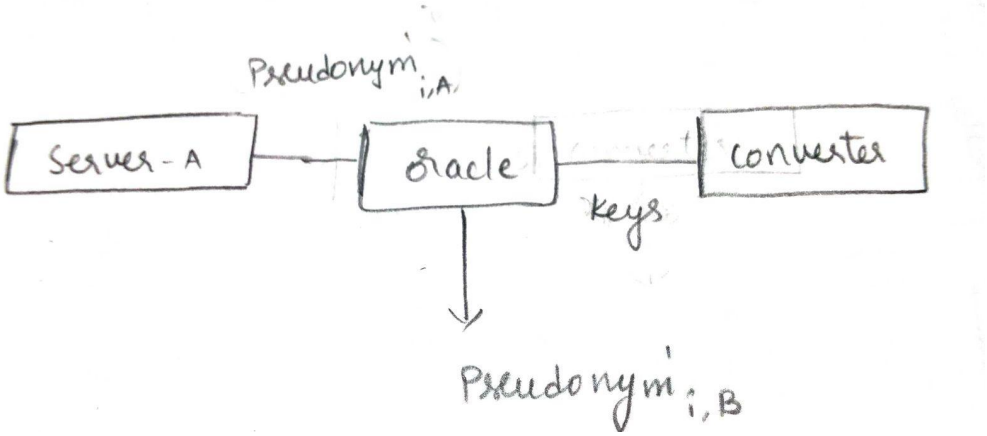
- Digital Credentials which one can obtain from issuers and can Verify without revealing any identifiable information.
- So even if the Verifiers collude, they cannot pinpoint the identity of the Credential presenter
- It allows a user to get a signature σ on a message m by sending a commitment of x to the signer
- User is basically using a zero-knowledge proof to convince the verifier of possessing a signature generated by the issuer

Proposed Solution-2

Pseudonym Generation:-



Conversion



- The oracle function blindly converts Pseudonym'$_{i,A}$ to Pseudonym'$_{i,B}$ without converter knowing Pseudonym'$_{i,A}$ & Pseudonym'$_{i,B}$

Oracle Function

Server - A

1) $\text{Pseudonym}'_{i,A} = \left[\text{PRF}(k_x, \text{uid}_i)\right]^{x_A}$

2) $\gamma \in G_q$

3) Sends $(\text{qid}', \text{Pseudonym}'^{\gamma}_{i,A})$ to converter

6) $\text{Enc}(\text{epk}_B, (\text{Pseudonym}'^{\gamma}_{i,B})^{1/\gamma})$

7) output $\text{Enc}(\text{epk}_B, \text{Pseudonym}'_{i,B})$

Converter

4) Calculate $(\text{Pseudonym}'^{\gamma * x_B / x_A}_{i,A})$

ZKP that converter indeed calculated correct.

5) Sends $\text{Enc}_{G_q}(\text{epk}_B, (\text{Pseudonym}'^{\gamma}_{i,B}))$ to S_A

References

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THANK YOU!!