

Mutable Data Types

A New Despair

Mutability Strikes Back

Return of Imperative Programming

Prof. Clarkson Fall 2019

Today's music: *The Imperial March* from the soundtrack to *Star Wars, Episode V: The Empire Strikes Back*

CLICKER QUESTIONS 1 AND 2

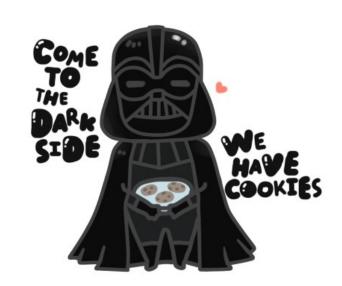
Review

Previously in 3110:

Efficiency: Big Oh

Lectures 13-15:

- Efficiency of data structures
- Running example: maps



Today: THE DARK SIDE ARRIVES

Mutable data types: refs, mutable fields, (arrays)

REFS

References

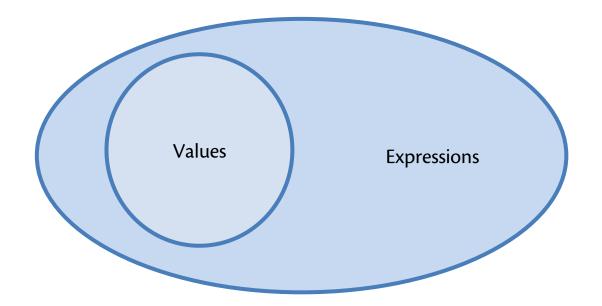
- Aka "ref" or "ref cell"
- Pointer to a typed location in memory
- Binding of variable to pointer: immutable
- Contents of memory location: mutable

References

- Syntax: ref e
- Evaluation:
 - Evaluate e to a value v
 - Allocate a new location loc in memory to hold v
 - Store v in loc
 - Return **loc**
- Type checking:
 - New type constructor: t ref where t is a type
 - Note: ref is used as keyword in type and as keyword in value
 - -ref e : t ref if e : t

Locations

- Locations are values
- Locations are not expressions
- So this picture [lec02] was a lie:



References

- Syntax: e1 := e2
- Evaluation:
 - Evaluate e2 to a value v2
 - Evaluate e1 to a location loc
 - Store v2 in loc
 - Return ()
- Type checking:
 - -lfe2:t
 - and e1 : t ref
 - then **e1:=e2** : **unit**

Unit

- unit is a type
 - Its only value is (), also pronounced "unit"
 - There are no interesting operations on unit
- Analogy: Booleans
 - bool is a type
 - It has just two values, true and false
 - So unit is like a Boolean with one fewer value
- Analogy: void
 - When a procedure in Java or C has no interesting value to return, its return type is void
 - Similar with print and assert in OCaml:
 - print_string : string -> unit
 - assert b : unit (assuming b : bool)

References

- Syntax: !e
 - note: not negation
- Evaluation:
 - Evaluate e to loc
 - Return contents of **loc**
- Type checking:
 - -lfe : t ref
 - -then !e : t

CLICKER QUESTION 3

Equality

- Suppose we have two refs...
 - let r1 = ref 3110
 - let r2 = ref 3110
- Double equals is physical equality
 - -r1 == r1
 - -r1 != r2
- Single equals is *structural equality*
 - -r1 = r1
 - -r1 = r2
 - ref 3110 <> ref 2110
- You usually want single equals

Semicolon

- Syntax: e1; e2
- Evaluation:
 - Evaluate e1 to a value v1
 - Then throw away that value (note: e1 could have side effects)
 - evaluate **e2** to a value **v2**
 - return v2
- Type checking:
 - If e1 : unit
 - and e2: t
 - then **e1**; **e2** : **t**

MUTABLE FIELDS

Implementing refs

Ref cells are essentially implemented as records with a mutable field:

```
type 'a ref = { mutable contents : 'a }
let ref x = { contents = x }
let (!) r = r.contents
let (:=) r newval = r.contents <- newval</pre>
```

- That type is declared in **Stdlib**
- The functions are compiled down to something equivalent



BEWARE

Immutable lists

We have never needed to worry about aliasing with lists!

let
$$x = [2;4]$$

let $y = [5;3;0]$
let $z = x @ y$

$$x \rightarrow 2 \rightarrow 4$$

$$y \rightarrow 5 \rightarrow 3 \rightarrow 0$$

$$y \rightarrow 5 \rightarrow 3 \rightarrow 0$$

$$y \rightarrow 5 \rightarrow 3 \rightarrow 0$$

$$z \rightarrow 2 \rightarrow 4 \rightarrow 5 \rightarrow 3 \rightarrow 0$$

(no code you write could ever tell, but OCaml implementation uses the first one)

OCaml:

blissfully unaware of aliasing

Java:

obsession with aliasing

Faulty code

```
class ProtectedResource {
   private Resource theResource = ...;
   private String[] allowedUsers = ...;
   public String[] getAllowedUsers() {
      return allowedUsers;
   public String currentUser() { ... }
   public void useTheResource() {
      for(int i=0; i < allowedUsers.length; i++) {</pre>
         if (currentUser().equals(allowedUsers[i])) {
             ... // access allowed: use it
             return;
      throw new IllegalAccessException();
```

Have to make copies

The exploit:

```
p.getAllowedUsers()[0] = p.currentUser();
p.useTheResource();
```

The fix:

```
public String[] getAllowedUsers() {
    ... return a copy of allowedUsers ...
}
```

Similar errors as recent as Java 1.7beta

Pros and cons of immutability

Pros:

- Programmer doesn't have to think about aliasing; can concentrate on other aspects of code
- Language implementation is free to use aliasing, which is cheap
- Often easier to reason about whether code is correct

Cons:

- I/O is fundamentally about mutation
- Some data abstractions (dictionaries, arrays, ...) are more efficient if imperative

Try not to abuse your new-found power!

Upcoming events

- [last night] R5 due
- [Wed] A4 due

This is (reluctantly) imperative.

THIS IS 3110