CS 3110

Functions

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Fall 2019

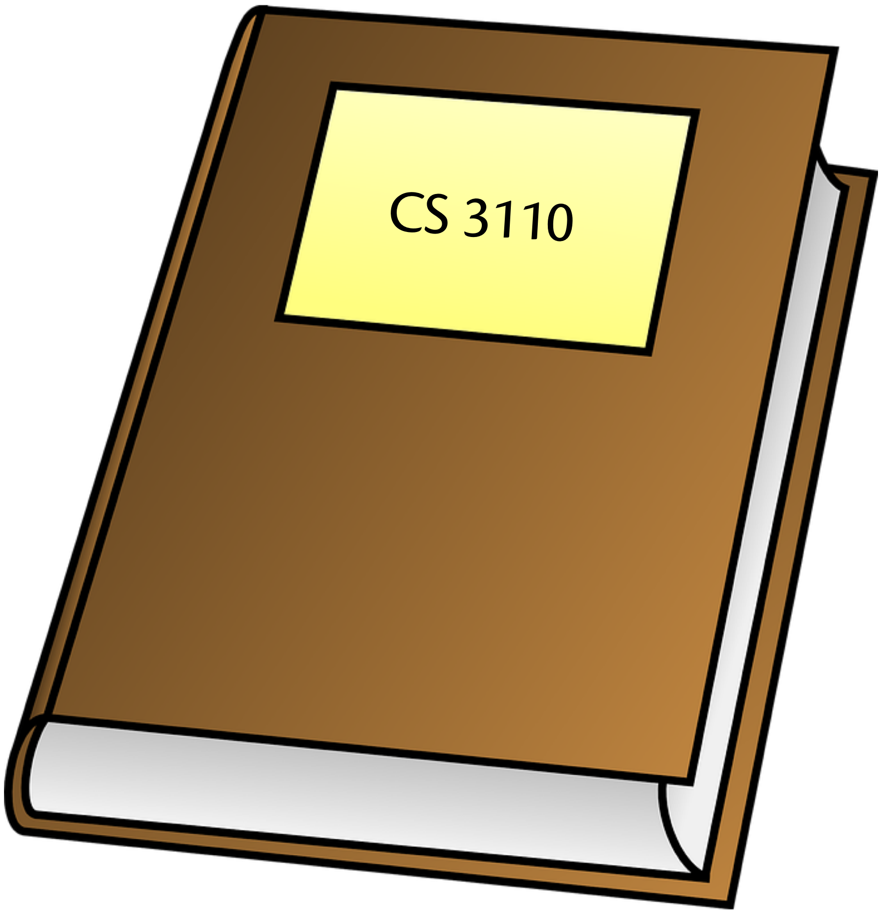
Today's music: Function by E-40 (Clean remix)

CLICKER QUESTION 1

Search

- Website
- Textbook
- Campuswire

Textbook vs. lecture



PRIMARY

summary

Review

Previously in 3110:

- **Syntax and semantics**
- **Expressions:** if, let
- **Definitions:** let

Today:

- **Functions**

ANONYMOUS FUNCTION EXPRESSIONS & FUNCTION APPLICATION EXPRESSIONS

Demo

Anonymous function expression

Syntax: `fun x1 ... xn -> e`

`fun` is a keyword

Evaluation:

- A function is a value: no further computation to do
- In particular, body **`e`** is not evaluated until function is applied

Lambda

- Anonymous functions a.k.a. *lambda expressions*
- Math notation: $\lambda x \; . \; e$
- The lambda means "what follows is an anonymous function"

Lambda

- Python
- Java 8
- A popular PL blog
- Lambda style

Functions are values

Can use them **anywhere** we use values:

- Functions can **take** functions as arguments
- Functions can **return** functions as results

This is an incredibly powerful language feature!

Function application

Syntax: `e0 e1 ... en`

No parentheses required!

(unless you need to force particular order of evaluation)

Function application

What is the evaluation rule for

`e0 e1 ... en` ?

Challenge: invent it right now!

Function application

Evaluation of `e0 e1 ... en`:

1. Evaluate subexpressions:

   `e0 ==> v0, ..., en ==> vn`

   Note that `v0` is guaranteed to be a function:

   `fun x1 ... xn -> e`

2. Substitute `vi` for `xi` in `e` yielding new expression `e'`. Evaluate it: `e' ==> v`

3. Result is `v`

Example

Let vs. function

These two expressions are **syntactically different** but semantically equivalent:

```
let x = 2 in x+1
(fun x -> x+1) 2
```

FUNCTION DEFINITIONS

Demo

Two syntaxes to define functions

These definitions are **syntactically different** but semantically equivalent:

```
let inc = fun x -> x+1
let inc x = x + 1
```

- First is fundamentally no different from `let` definitions we saw last lecture
- Second is syntactic sugar: not necessary, makes language "sweeter"

Recursive function definition

Must explicitly state that function is recursive:

```
let rec f ...
```

Demo

Reverse application

- Instead of `f e` can write `e |> f`
- Use: pipeline a value through several functions
  `5 |> inc |> square (* ==> 36*)`

assuming

```
let inc x = x + 1
let square x = x * x
```

FUNCTIONS AND TYPES

Function types

Type `t -> u` is the type of a function that takes input of type `t` and returns output of type `u`

Type `t1 -> t2 -> u` is the type of a function that takes input of type `t1` and another input of type `t2` and returns output of type `u`

etc.

Note dual purpose for `->` syntax:

- Function types
- Function values

Function application

Type checking:

If `e0 : t1 -> ... -> tn -> u`

And `e1 : t1,`

`...,`

`en : tn`

Then `e0 e1 ... en : u`

Anonymous function expression

Type checking:

If `x1:t1, ..., xn:tn`

And `e:u`

Then `(fun x1 ... xn -> e) :`

`t1 -> ... -> tn -> u`

PARTIAL APPLICATION

Demo

More syntactic sugar

Multi-argument functions do not exist

```
fun x y -> e
```

is syntactic sugar for

```
fun x -> (fun y -> e)
```

More syntactic sugar

Multi-argument functions do not exist

```
let add x y = x + y
```

is syntactic sugar for

```
let add = fun x ->
            fun y ->
              x + y
```

CLICKER QUESTION 2

More syntactic sugar

Multi-argument functions do not exist

```
fun x y z -> e
```

is syntactic sugar for

```
fun x -> (fun y -> (fun z -> e))
```

Again: Functions are values

Can use them **anywhere** we use values:

- Functions can **take** functions as arguments
- Functions can **return** functions as results

This is an incredibly powerful language feature!

Upcoming events

- [today] A0 released, R1 released
- [Monday] R1 due
- [Wednesday] A0 due

This is **`fun`***!*

THIS IS 3110