



CS 3110

Abstraction and Specification

Prof. Clarkson
Fall 2019

Today's music: Never Change by JAY-Z

CLICKER QUESTION 1



OCaml



Review

Previously in 3110:

- OCaml module system

Today:

- Abstract and specification

Abstraction

(verb)

Forgetting information, so that different things can be treated as the same

(noun)

Artifacts that result from that process

Specification

(noun)

Intended behavior of abstraction

(verb)

The act of creating such an artifact

Audience of specification

- Clients

- What they must guarantee (preconditions)
- What they can assume (postconditions)

- Implementers

- What they can assume (preconditions)
- What they must guarantee (postconditions)

Benefits of specifications

- **Locality:** understand abstraction without needing to read implementation
- **Modifiability:** change implementation without breaking client code
- **Accountability:** clarify who is to blame

Specifications are contracts



Satisfaction

An implementation **satisfies** a specification if it provides the described behavior

Many implementations can satisfy the same specification

- **Client** has to assume it could be any of them
- **Implementer** gets to pick one

What if spec is ambiguous?

Ambiguity is a fact of life.

Do the most reasonable thing you can.

Probably not 

Who wrote it?

- **You:** improve it
- **Client:** seek clarification

but if you make 500 requests they probably won't hire you again

SPECIFYING FUNCTIONS

Template

```
(** [f x] is ...  
    Example: ...  
    Requires: ...  
    Raises: ... *)  
  
val f : t -> u
```

Based on *Abstraction and Specification in Program Development*

(Now *Program Development in Java: Abstraction, Specification, and Object-Oriented Design*)

By Barbara Liskov and John Guttag

Barbara Liskov



b. 1939

Turing Award Winner 2008

*For contributions to practical and theoretical foundations of programming language and system design, especially related to **data abstraction**, fault tolerance, and distributed computing.*

Data abstraction

- A **data abstraction** is a specification of operations on a set of values
 - e.g., stacks have push, pop, peek, etc.; we don't know what the values concretely are
- A **data structure** is an implementation of a data abstraction with a particular representation
 - e.g., `ListStack` implemented `StackSig` with `'a list, (::), etc.`

SPECIFYING DATA ABSTRACTIONS

Demo

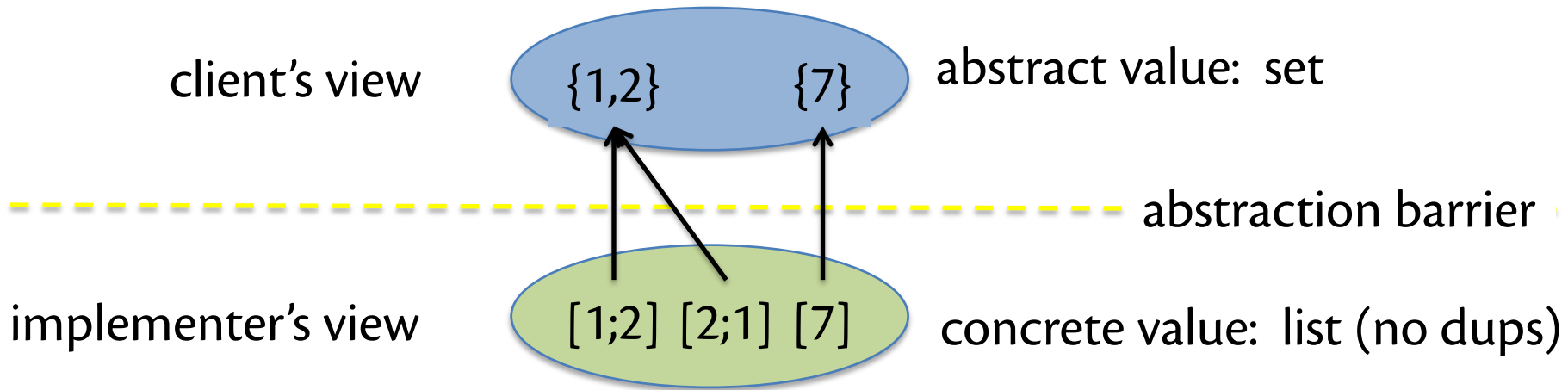
CLICKER QUESTION 2

Representation types

- Q: How to **interpret** the representation type as the data abstraction?
- A: Abstraction function
- Q: How to determine which values of representation type are **meaningful**?
- A: Representation invariant

ABSTRACTION FUNCTIONS

Abstraction function



the black arrows are the abstraction function

Abstraction function

maps

valid concrete values

to

abstract values

Documenting the AF

- Above rep type in implementation you write:
`(* AF: comment *)`
- Write it **first** before implementing operations

Implementing the AF

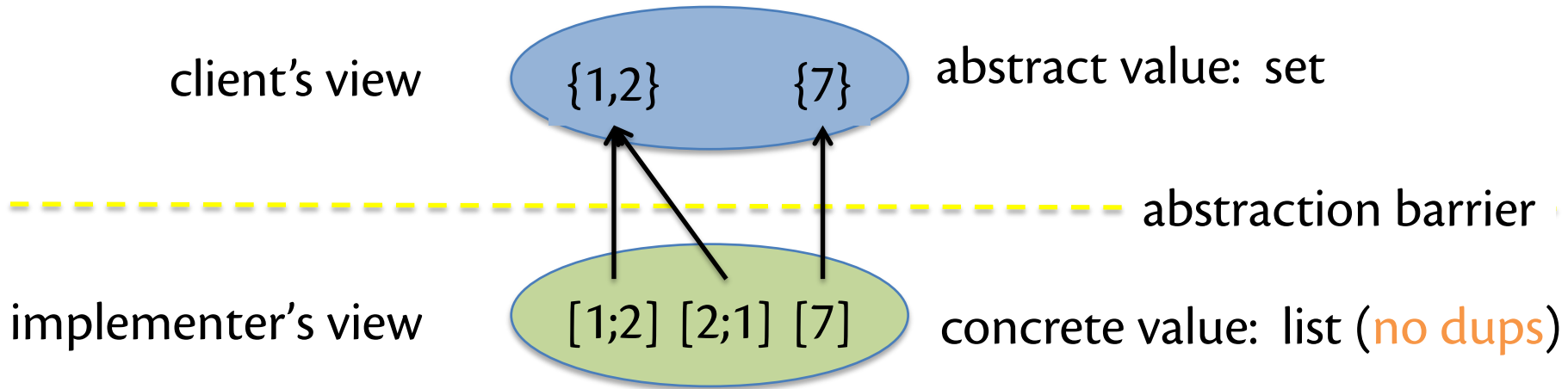
- You don't
 - Abstract values are a mathematical idea not a programming reality
 - Would need to have an OCaml type for abstract values
 - If you had that type, you wouldn't need whatever data abstraction you're working on
- But conversion to strings comes close

Representation types

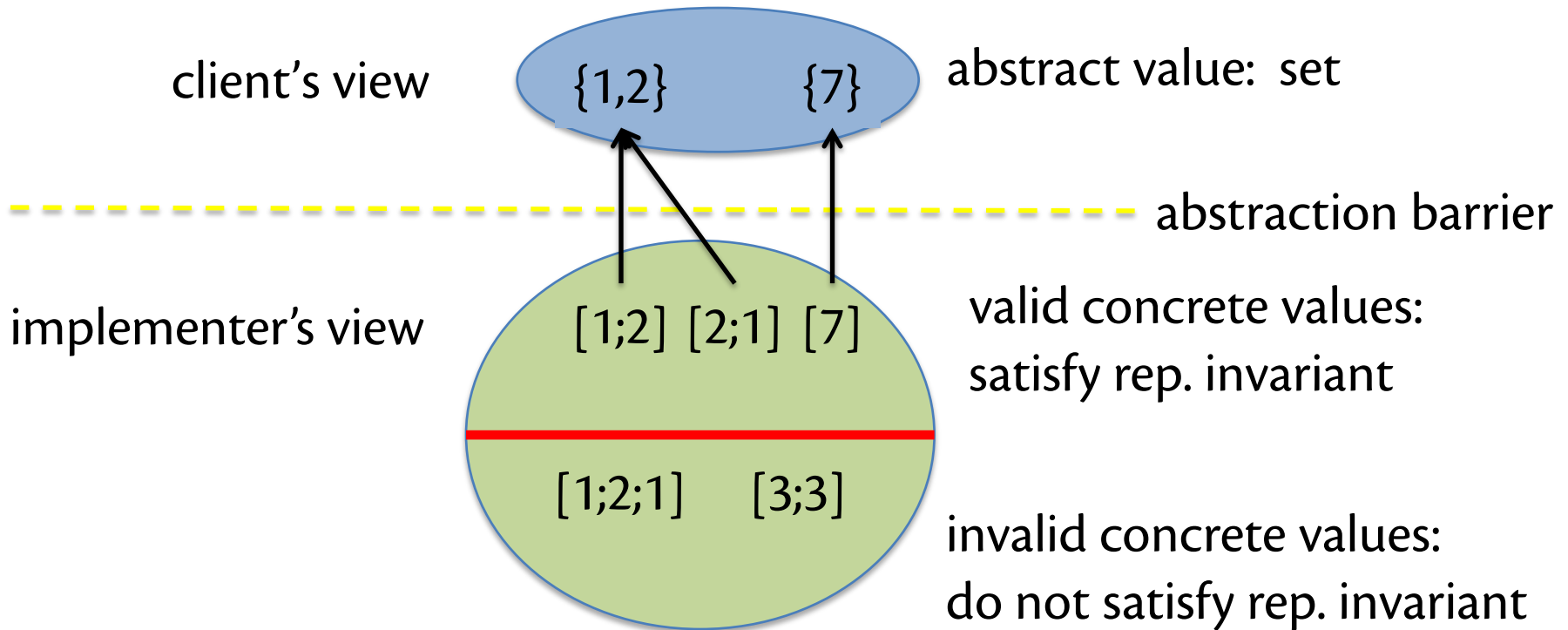
- Q: How to **interpret** the representation type as the data abstraction?
- A: Abstraction function
- Q: How to determine which values of representation type are **meaningful**?
- A: Representation invariant

REPRESENTATION INVARIANTS

Abstraction function



Representation invariant



the thick red line is the rep. invariant

Rep. invariant

distinguishes

valid concrete values

from

invalid concrete values

Documenting the RI

- Above rep type in implementation you write:
`(* RI: comment *)`
- Write it **first** before implementing operations

Rep. invariant
implicitly part of
every precondition and
every postcondition
in abstraction

Invariant may temporarily be violated

concrete
input



concrete
operation



concrete
output



RI holds



RI maybe violated



RI holds



Demo

Implementing the RI

Idiom:

- write `rep_ok` function
- call function on every input and output

```
let rep_ok (x : t) : t =  
  if (* check RI *) then t  
  else failwith "RI"
```

This saved a 3110 final project game tournament one year!

Upcoming events

- [last night] A2 due
- [later today] A3 out

This is invariant.

THIS IS 3110