CS3281 / CS5281

# Advanced Virtual Memory

CS3281 / CS5281

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**Some lecture slides borrowed and adapted from CMU's "Computer Systems: A Programmer's Perspective" and MIT's 6.S081 Course*

ISIS
Tel (615) 343-7472 | Fax (615) 343-7440
1025 16th Avenue South Nashville, TN 37212
www.isis.vanderbilt.edu

VANDERBILT UNIVERSITY

# Today

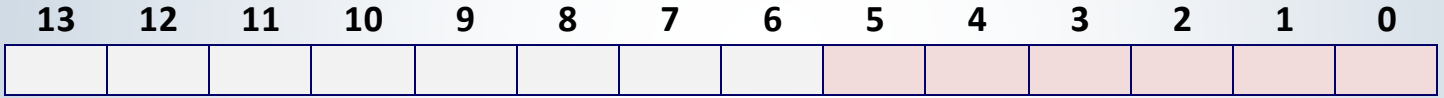
- Simple memory system example
- Memory mapping

# Review of Symbols

- Basic Parameters
  - **N = $2^n$** : Number of addresses in virtual address space
  - **M = $2^m$** : Number of addresses in physical address space
  - **P = $2^p$** : Page size (bytes)
- Components of the virtual address (VA)
  - **VPO**: Virtual page offset
  - **VPN**: Virtual page number
- Components of the physical address (PA)
  - **PPO**: Physical page offset (same as VPO)
  - **PPN:** Physical page number

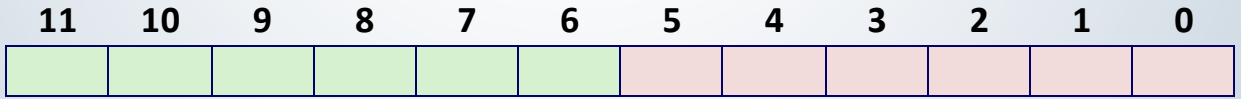
# Simple Memory System Example

- Addressing
  - 14-bit virtual addresses
  - 12-bit physical address
  - Page size = 64 bytes




Virtual Page Number

Virtual Page Offset




Physical Page Number

Physical Page Offset

# Simply Memory System Page Table

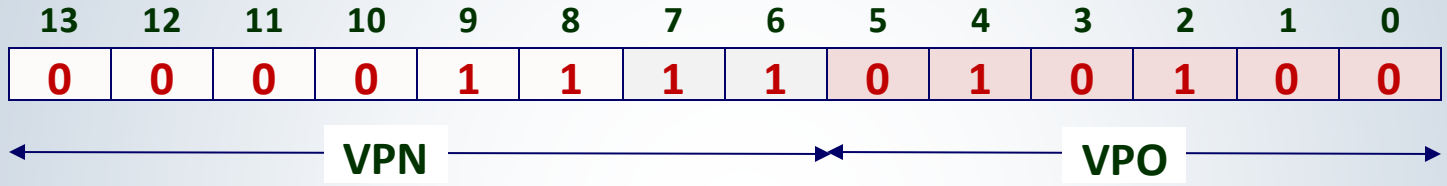
- Only show first 16 entries (out of )

| VPN | PPN | Valid |
|---|---|---|
| 00 | 28 | 1 |
| 01 | – | 0 |
| 02 | 33 | 1 |
| 03 | 02 | 1 |
| 04 | – | 0 |
| 05 | 16 | 1 |
| 06 | – | 0 |
| 07 | – | 0 |

| VPN | PPN | Valid |
|---|---|---|
| 08 | 13 | 1 |
| 09 | 17 | 1 |
| 0A | 09 | 1 |
| 0B | – | 0 |
| 0C | – | 0 |
| 0D | 2D | 1 |
| 0E | 11 | 1 |
| 0F | 0D | 1 |

# Address Translation Example #1

Virtual Address: 0x03D4



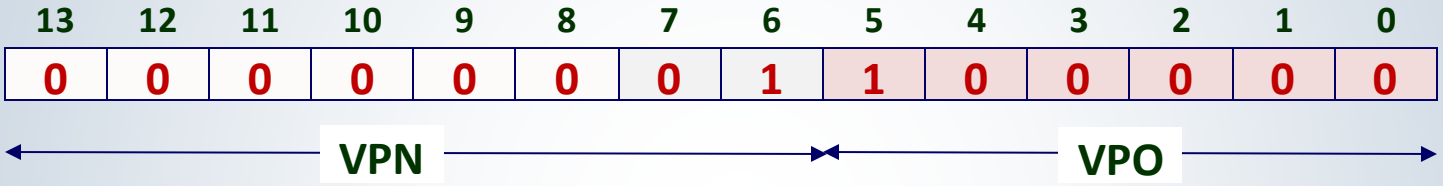

VPN:

Page Fault:

PPN:

# Address Translation Example #2

**Virtual Address: 0x0060**




VPN: **0x01**

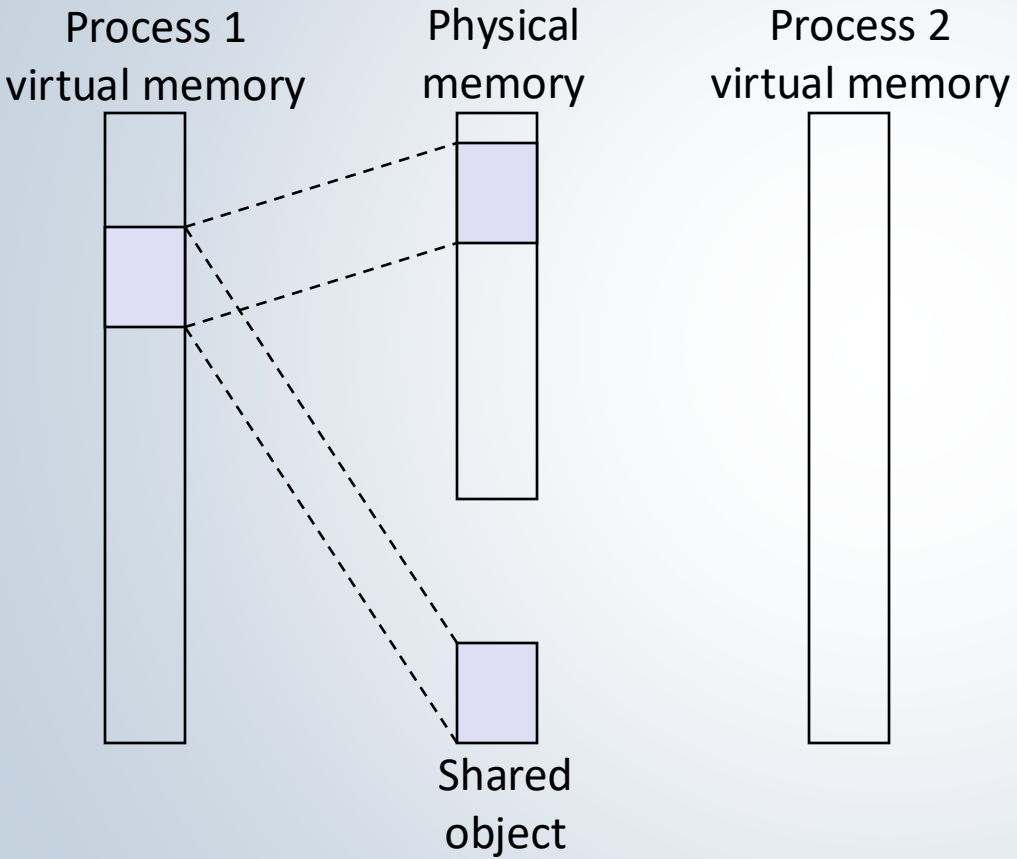
Page Fault: Y

PPN:

# Today

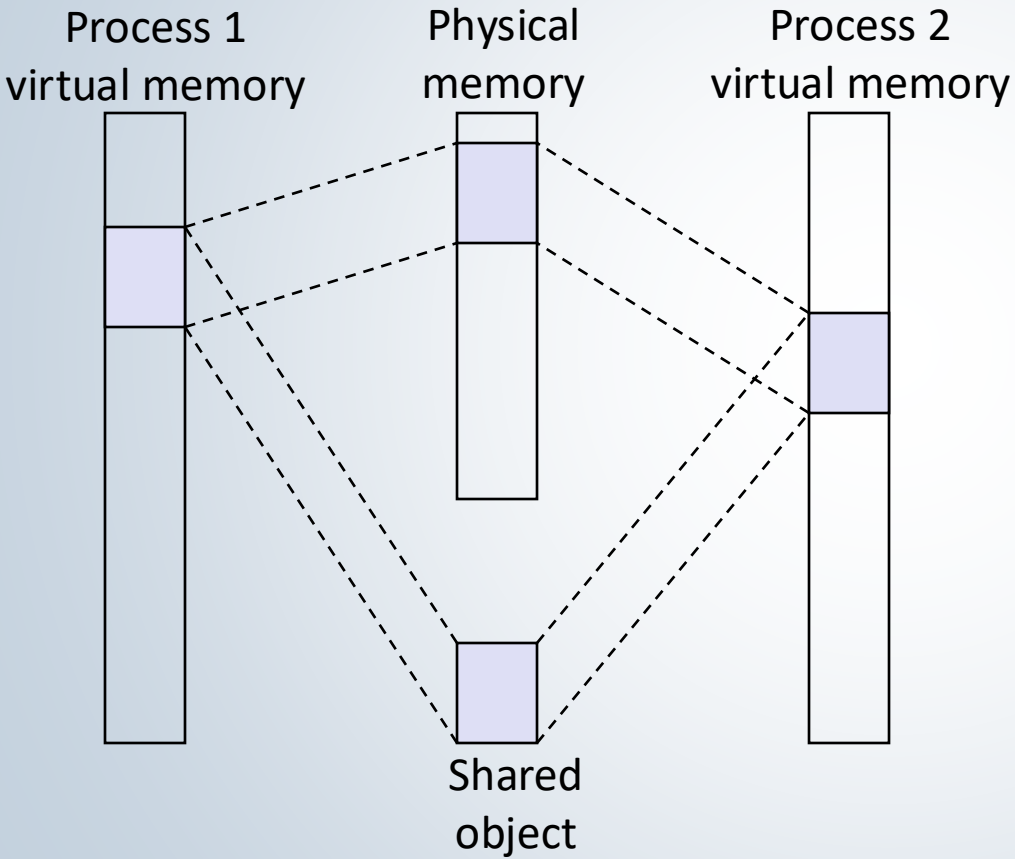
- Simple memory system example
- Shared Memory and Copy-on-Write
- Memory mapping

# Sharing Revisited: Shared Objects



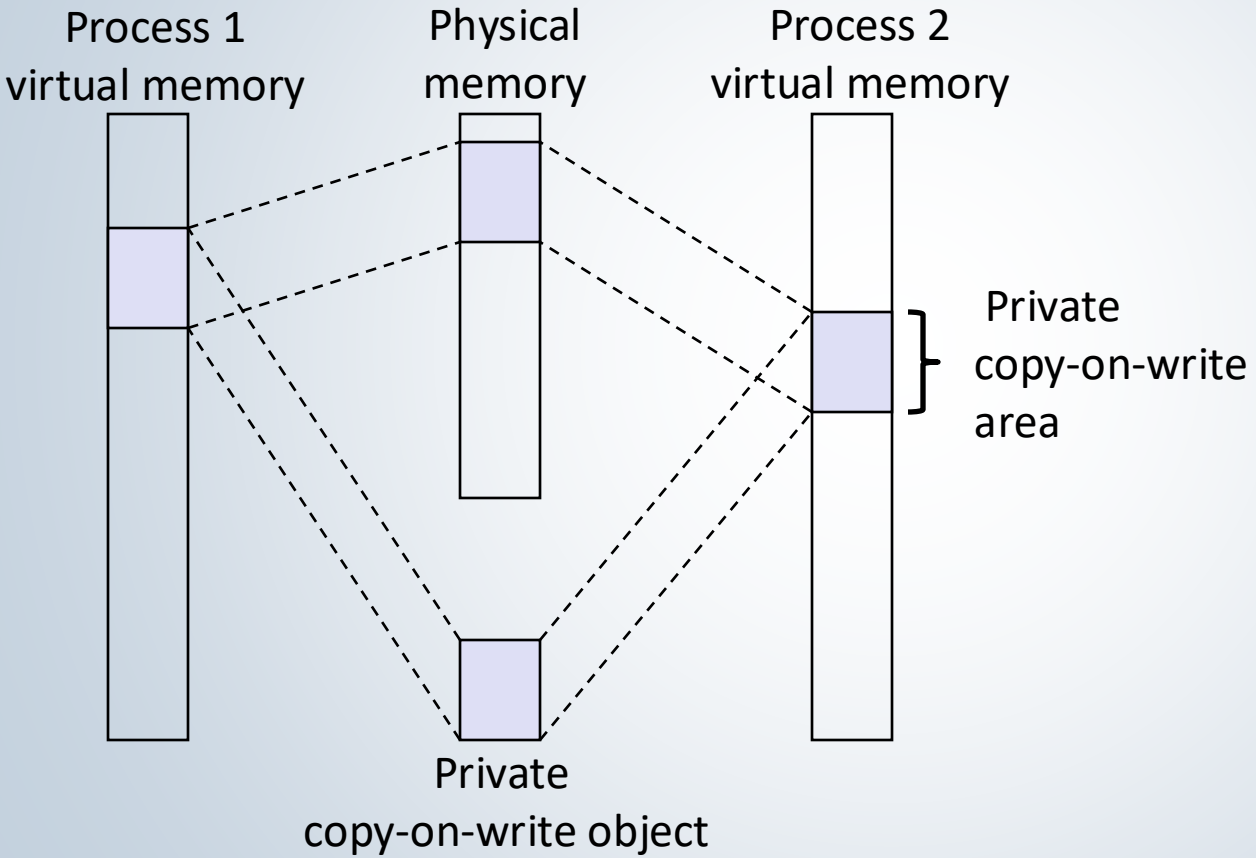

- **Process 2 maps the shared object.**

# Sharing Revisited: Shared Objects



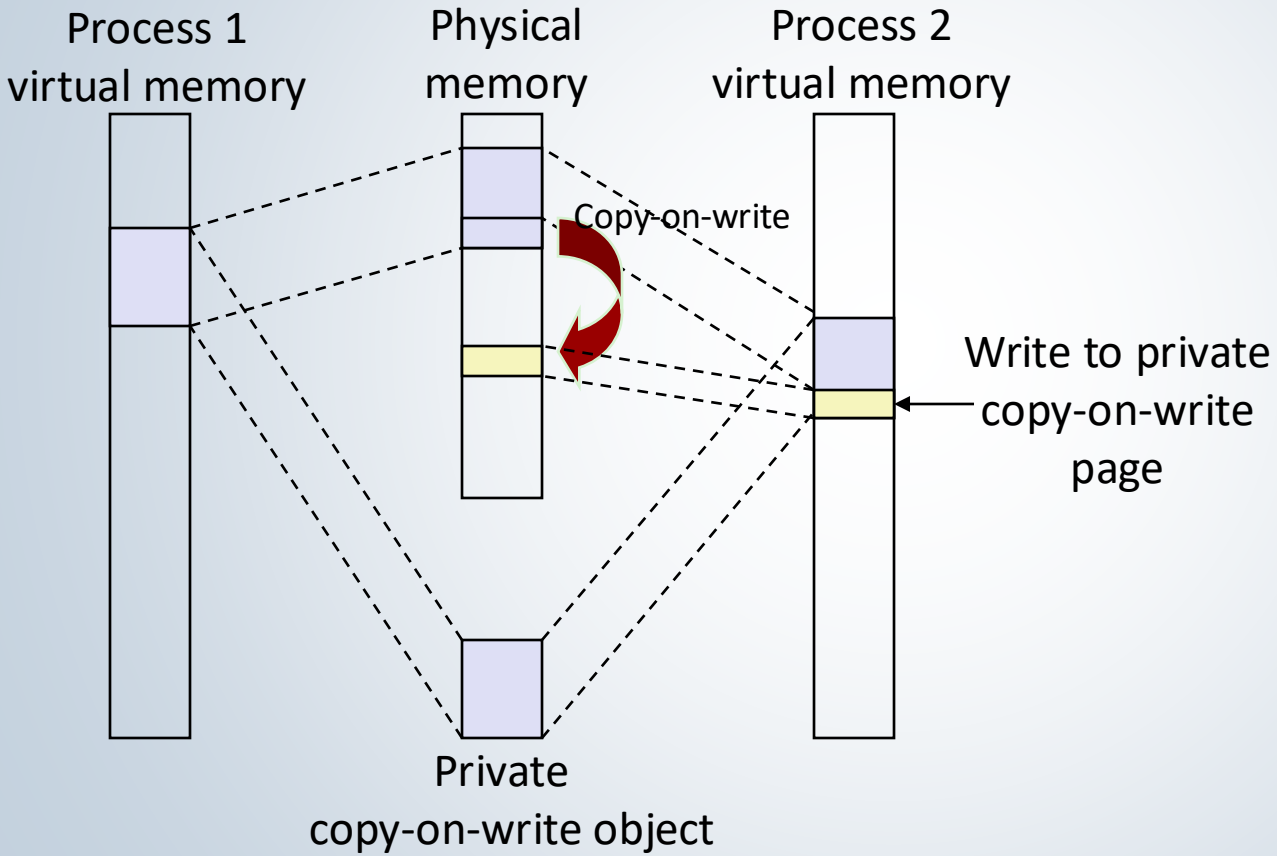

- **Process 2 maps the shared object.**
- Notice how the virtual addresses can be different.

# Sharing Revisited: Copy-On-Write (COW) Objects




- Two processes mapping a *private copy-on-write (COW)* object
- Initially, this object is like a shared object
- The access right for the pages of this object is read-only in PTEs

# Sharing Revisited: Copy-On-Write (COW) Objects




- Instruction writing to private page triggers protection fault
- Handler creates new R/W page
- Instruction restarts upon handler return
- Copying deferred as long as possible!

# The fork() Function Revisited

- Can use COW memory mapping in `fork()` to provides private address space for each process without duplicating physical memory unnecessarily

- To create virtual address for new process
  - Create exact copies of current page tables
  - Flag each page in PETs of both processes as read-only

- On return, each process has identical view of memory but only one copy of physical memory exists

- Subsequent writes, e.g., with exec(), trigger COW mechanism and force pages to be duplicated when needed

# Memory Mapping

- VM areas initialized by associating them with disk objects.
  - Process is known as ***memory mapping***.

- Area can be *backed by* (i.e., get its initial values from) :
  - ***Regular file*** on disk (e.g., an executable object file)
    - Initial page bytes come from a section of a file
  - ***Anonymous file*** (e.g., nothing)
    - First fault will allocate a physical page full of 0's (***demand-zero page***)
    - Once the page is written to (***dirtied***), it is like any other page

- Dirty pages are copied back and forth between memory and a special *swap file*.

# User-Level Memory Mapping

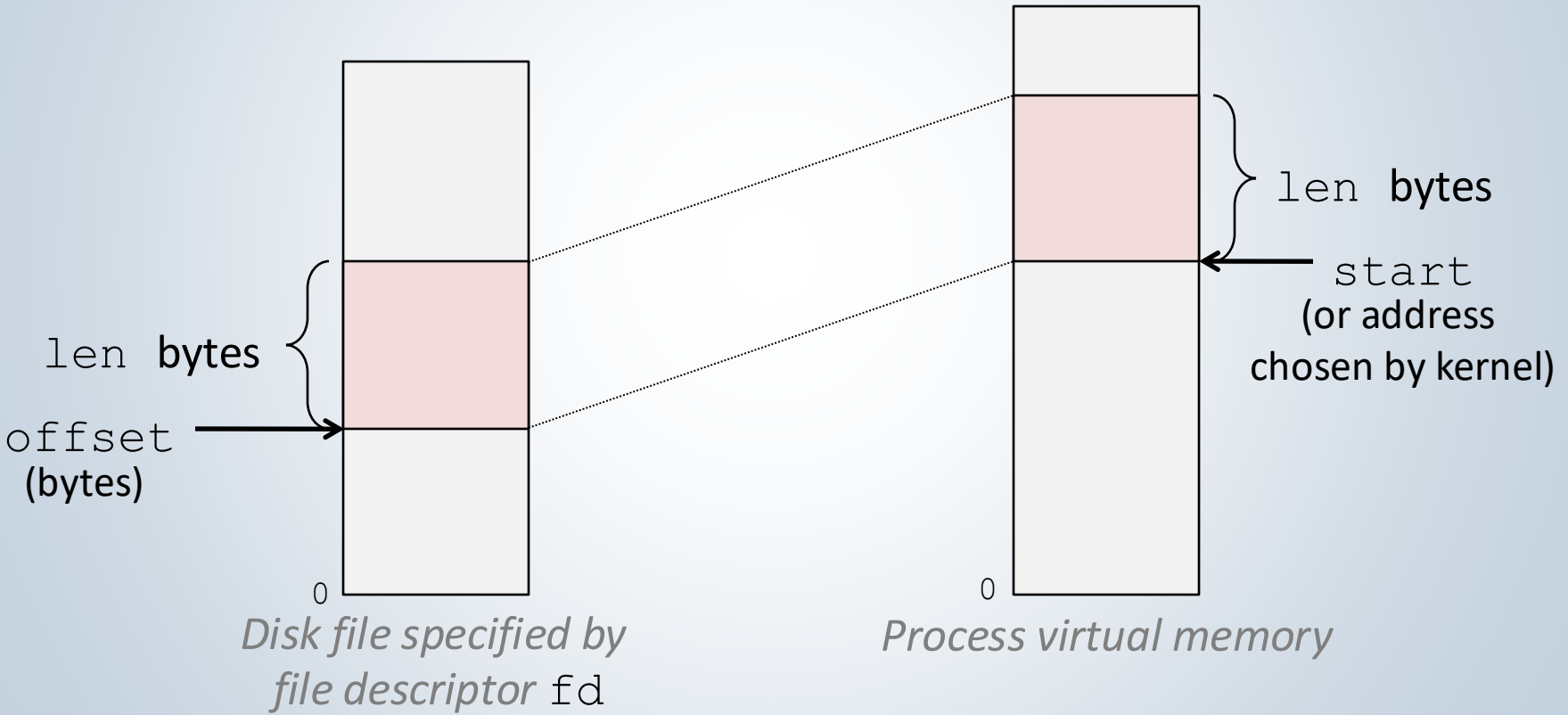
```
void *mmap(void *start, int len,
           int prot, int flags, int fd, int offset)
```

- Map **`len`** bytes starting at offset **`offset`** of the file specified by file description **`fd`**, preferably at address **`start`**
  - **`start`**: may be 0 for "pick an address"
  - **`prot`**: PROT_READ, PROT_WRITE, ...
  - **`flags`**: MAP_ANON, MAP_PRIVATE, MAP_SHARED, ...

- Return a pointer to start of mapped area (may not be **`start`**)
- **malloc()** calls mmap() to allocate new pages

# User-Level Memory Mapping

```
void *mmap(void *start, int len, int prot, int flags, int fd, int offset)
```



# Using mmap() to Copy Files (Linux)

- Copying a file to `stdout` without transferring data to user space

```
#include "csapp.h"

void mmapcopy(int fd, int size)
{

   /* Ptr to memory mapped area */
   char *bufp;

   bufp = mmap(NULL, size,
            PROT_READ,
            MAP_PRIVATE,
            fd, 0);
   Write(1, bufp, size);
   return;
}
```

```
/* mmapcopy driver */
int main(int argc, char **argv)
{
   struct stat stat;
   int fd;

   /* Check for required cmd line arg */
   if (argc != 2) {
      printf("usage: %s <filename>\n",
           argv[0]);
      exit(0);
   }

   /* Copy input file to stdout */
   fd = Open(argv[1], O_RDONLY, 0);
   Fstat(fd, &stat);
   mmapcopy(fd, stat.st_size);
   exit(0);
}
```