

HW | Cryptography

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- 1) Yes, this modified scheme should still have one time secrecy. It only takes away one possible option out of all the rest. It's also the trivial option. If λ isn't small the scheme is secure. $\frac{1}{2} \approx \frac{1}{2^{n-1}}$.

2) $\text{Query}(2, 5)$

$K \leftarrow \text{KeyGen}$

$\text{Enc}(K, 2)$

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$\text{Enc}(K, 5)$

The left library will have an output of only even numbers in \mathbb{Z}_{10} , and the right library can only output 0 or 5. Therefore it would be easy to distinguish between libraries since they have different probabilities.

3) $\text{Dec}(K, C)$

return $(C - K) \% n$

Encrypt(K, m_L, m_R) \Rightarrow " "
return $(K + m_L) \% n$ return $\text{ctxt}(K, m_L)$

$\text{ctxt}(K, m)$ \Leftrightarrow $\text{ctxt}(K, m)$ so we can send
 $c \leftarrow (K + m) \% n$ $c \in \mathbb{Z}_n$ in any m into
return c return c ctxt since it won't use it

Encrypt(K, m_L, m_R) Return to the former ctxt
return $\text{ctxt}(K, m_R)$ and in line the code.

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return $(K + m_R) \% n$

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