

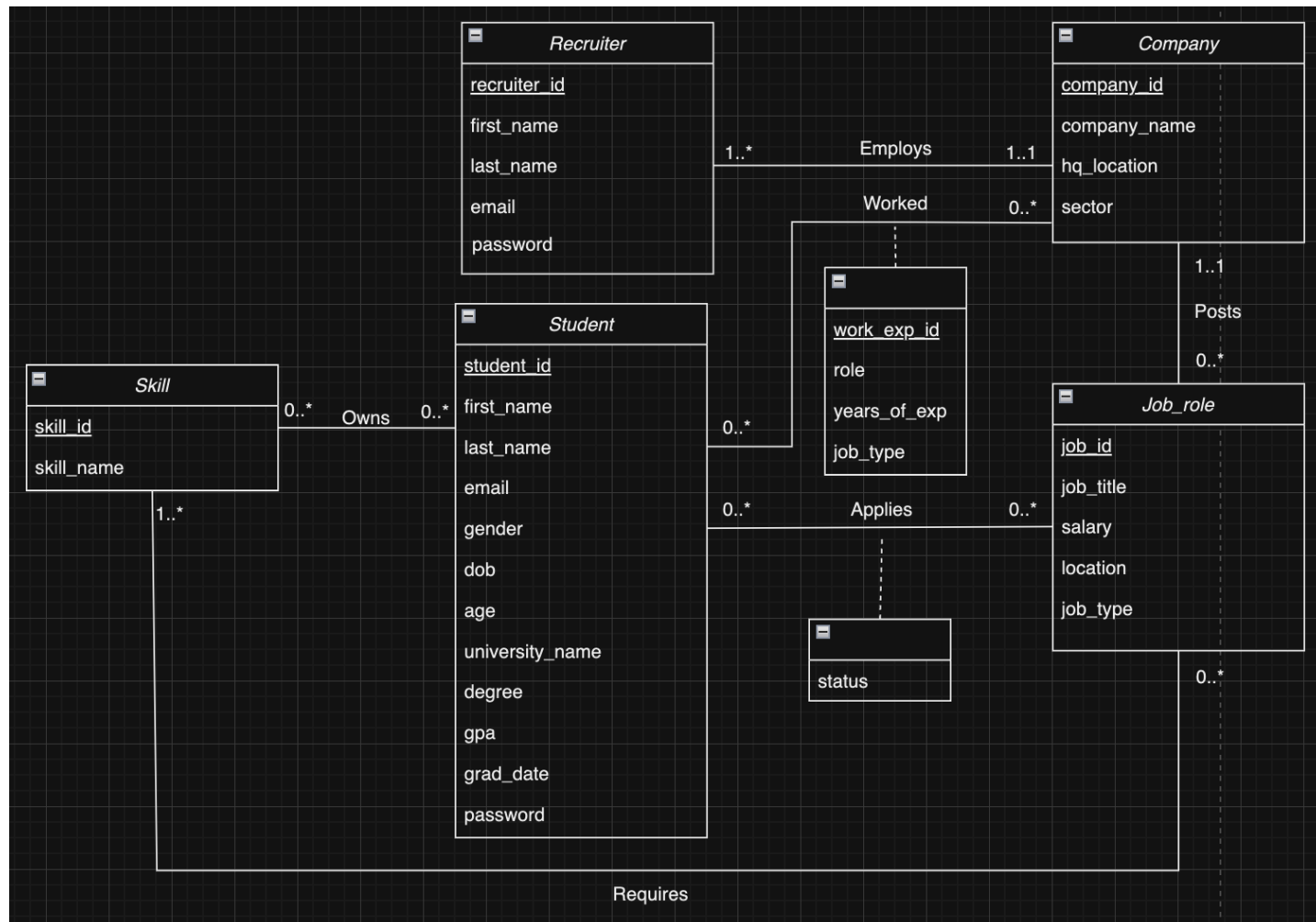
HireIt

Stage 2: Database Design

Team No: 081

Team Name: ACID

UML Diagram:



Assumptions:

- We assume that the Company table will have the list of all the companies that are recruiting through our platform. This also includes companies that a student may have worked for previously. For example, if a student has worked for Amazon previously and is currently applying for Microsoft, we assume that both these companies exist in the Company table.
- We assume that each job role posting will require a minimum of one skill.

- We assume that the number of years of work experience(years_of_exp) is a whole number. Ie, We consider years of experience in terms of years and not months.
- The dob attribute in the student table determines the value of age

Description of each relationship and its cardinality

- **Employs:**
This is a one-many relationship table between the company and recruiter table.
1 company can have 1 or more employees
1 recruiter will be employed by one company
- **Worked:**
This is a many-many relationship table between the company and student table.
This relationship table also contains attributes for job role (role), years of experience (years_of_exp) and work experience ID (work_exp_id).
A student could have worked at 0 or many of the companies in the past
A company could have been the workplace for 0 or many students
- **Posts:**
This is a one-many relationship table between the company and job_role table.
1 company can have 0 or more job role postings
A job role posting will be will posted by 1 company
- **Applies:**
This is a many-many relationship table between the student and job_role table.
This relationship table also has an attribute to track the status of the job application (status).
A student can apply to 0 or many job roles
A job role can be applied by 0 or many students
- **Requires:**
This is a one-many relationship table between the job_role and skill table.
1 job role will require 1 or many skills
A skills maybe required by 0 or many job roles
- **Owns:**
This is a many-many relationship table between the student and skills table.
A student can own 1 or many skills
A skill maybe owned by 1 or many students

Relational Schema:

- **Student** (student_id:INT NOT NULL [PK], first_name:VARCHAR(255) , last_name:VARCHAR(255), email:VARCHAR(255), gender:VARCHAR(255),dob:DATETIME, age: INT, university_name:VARCHAR(255), degree:VARCHAR(255), gpa:DECIMAL(1,2), grad_date: DATETIME, password: VARCHAR(20))
- **Recruiter** (recruiter_id: INT NOT NULL [PK] , first_name: VARCHAR(255), last_name: VARCHAR(255), email: VARCHAR(255), password: VARCHAR(20))
- **Company** (company_id: INT NOT NULL [PK] L, company_name: VARCHAR(255), hq_location: VARCHAR(255), sector: VARCHAR(255))
- **Job_role** (job_id: INT [PK], company_id:INT [FK to Company.company_id], job_title: VARCHAR(255), salary: INT, location: VARCHAR(255), job_type: VARCHAR(255))
- **Skill** (skill_id: INT NOT NULL [PK], skill_name: VARCHAR(255))
- **Applies** (status: VARCHAR(255), student_id:INT [FK to Student.student_id], recruiter_id:INT [FK to Recruiter.recruiter_id], student_id,recruiter_id:INT,INT [PK])
- **Worked** (company_id:INT [FK to Company.company_id], student_id :INT [FK to Student.student_id], work_exp_id:INT [PK], role: VARCHAR(255), years_of_exp: INT, job_type: VARCHAR(255))
- **Requires** (student_id:INT [FK to Student.student_id, job_id:INT [FK to Job_role.job_id], job_id:INT, student_id,job_id: INT,INT [PK])
- **Owns** (student_id:INT [FK to Student.student_id], skill_id:INT [FK to Skill.skill_id], student_id,skill_id: INT,INT [PK])

Functional Dependencies:

Recruiter

- recruiter_id -> first_name, last_name, email, password

- email -> recruiter_id, first_name, last_name, password

Company

- company_id -> company_name, hq_location, sector

Job_role

- job_id -> job_title, salary, location, job_type

Student

- student_id -> first_name, last_name, email, gender, dob, age, university_name, degree, gpa, grad_date, password
- email -> first_name, last_name, student_id, gender, dob, age, university_name, degree, gpa, grad_date, password
- student_id, dob -> age
- email_id, dob -> age

Skill

- skill_id -> skill_name

Applies

- student_id, job_id -> status

Worked

- work_exp_id -> student_id, company_id, role, job_type, years_of_exp

Normalization:

1. Recruiter Relation:

- recruiter_id -> A
- first_name -> B
- last_name -> C
- email -> D
- password -> E

A -> BCDE

D -> ABCE

a) Making RHS of every FD as singleton

A -> B

A → C

A → D

A → E

D → A

D → B

D → C

D → E

b) There are no redundant attributes in the LHS

c) Removing unnecessary FD's

A → D

D → B

D → C

D → E

D → A

Finally,

A → D

D → ABCE

Final FD's:

- recruiter_id → email
- email → recruiter_id, first_name, last_name, password

Candidate Key - {recruited_id, email}

In the above relation, the LHS of the minimal basis FD is a super key. Therefore the relation is in both BCNF and 3NF

2. Company Relation:

- company_id → A
- company_name → B
- hq_location → C
- sector → D

A → BCD

a. Making RHS of every FD as singleton

A → B

A → C

$A \rightarrow D$

b. There is no redundant attributes in the LHS

c. Removing unnecessary FD's

$A \rightarrow B$

$A \rightarrow C$

$A \rightarrow D$

Finally,

$A \rightarrow BCD$

Final FD:

- $\text{company_id} \rightarrow \text{company_name}, \text{hq_location}, \text{sector}$
- Candidate Key - $\{\text{company_id}\}$

In the above relation, the LHS of the minimal basis FD is a super key. Therefore the relation is in both BCNF and 3NF

3. Job Relation:

- $\text{job_id} \rightarrow A$
- $\text{job_title} \rightarrow B$
- $\text{salary} \rightarrow C$
- $\text{location} \rightarrow D$
- $\text{job_type} \rightarrow E$

$A \rightarrow BCDE$

a. Making RHS of every FD as singleton

$A \rightarrow B$

$A \rightarrow C$

$A \rightarrow D$

$A \rightarrow E$

b. There is no redundant attributes in the LHS

c. Removing unnecessary FD's

$A \rightarrow B$

$A \rightarrow C$

$A \rightarrow D$

Finally,

$A \rightarrow BCDE$

Final FD:

- $\text{job_id} \rightarrow \text{job_title, salary, location, job_type}$
- Candidate Key - $\{\text{job_id}\}$

In the above relation, the LHS of the minimal basis FD is a super key. Therefore the relation is in both BCNF and 3NF

4. Skill Relation:

- $\text{skill_id} \rightarrow A$
- $\text{skill_name} \rightarrow B$

$A \rightarrow B$

a. Making RHS of every FD as singleton

$A \rightarrow B$

b. There is no redundant attributes in the LHS

c. Removing unnecessary FD's

$A \rightarrow B$

Final FD:

- $\text{skill_id} \rightarrow \text{skill_name}$

Candidate Key - $\{\text{skill_id}\}$

In the above relation, the LHS of the minimal basis FD is a super key. Therefore the relation is in both BCNF and 3NF

5. Student Relation

- $\text{student_id} \rightarrow A$
- $\text{first_name} \rightarrow B$
- $\text{last_name} \rightarrow C$
- $\text{email} \rightarrow D$
- $\text{gender} \rightarrow E$
- $\text{dob} \rightarrow F$
- $\text{age} \rightarrow G$
- $\text{university_name} \rightarrow H$
- $\text{degree} \rightarrow I$
- $\text{gpa} \rightarrow I$

- grad_date->J
- password->K

A -> BCDEFGHIJK

D -> ABCEFGHIJK

AF->G

DF->G

a. Making RHS of every FD as singleton

A ->B

A ->C

A ->D

A ->E

A ->F

A ->G

A ->H

A ->I

A ->J

A ->K

D ->B

D ->C

D ->A

D ->E

D ->F

D ->G

D ->H

D ->I

D ->J

D ->K

AF -> G

DF- > G

b. Redundant attributes in the LHS

AF -> G is redundant because A->G

Therefore it can be AF->G can be reduced to A->G

DF -> G is redundant because D -> G

Therefore DF-> G can be reduced to D -> G

c. Removing unnecessary FD's

A → D
D → A
D → B
D → C
D → E
D → F
D → G
D → H
D → I
D → J
D → K

Finally,
A → D
D → ABCEFGHIJK

Final FD:

- student_id → email
- email → first_name, last_name, student_id, gender, dob, age, university_name, degree, gpa, grad_date, password

Candidate Keys - {student_id, email}

In the above relation, the LHS of all the minimal basis FD is a super key.
Therefore the relation is in both BCNF and 3NF

6. Applies Relation

- student_id → A
- job_id → B
- Status → C

AB → C

a. Making RHS of every FD as singleton

AB → C

b. There is no redundant attributes in the LHS

c. Removing unnecessary FD's

AB → C

Final FD:

- Student_id, job_id → Status

Candidate Key - {student_id}

In the above relation, the LHS of the minimal basis FD is a super key. Therefore the relation is in both BCNF and 3NF

7. Worked Relation

- work_exp_id -> A
- Student_id -> B
- Company_id -> C
- Role -> D
- job_type -> E
- Years_of_exp -> F

A -> BCDEF

a) Making RHS of every FD as singleton

A ->B

A ->C

A ->D

A ->E

A ->F

b) There is no redundant attributes in the LHS

c) Removing unnecessary FD's

A ->B

A ->C

A ->D

A ->E

A ->F

Finally,

A -> BCDEF

Final FD's:

- work_exp_id -> student_id, company_id, role, job_type, years_of_exp

Candidate Key - {work_exp_id}

In the above relation, the LHS of the minimal basis FD is a super key. Therefore the relation is in both BCNF and 3NF

Normalized FD's:

- recruiter_id -> email
- email -> recruiter_id, first_name, last_name, password
- company_id -> company_name, hq_location, sector
- job_id -> job_title, salary, location, job_type
- skill_id -> skill_name
- student_id -> email
- email -> first_name, last_name, student_id, gender, dob, age, university_name, degree, gpa, grad_date, password
- Student_id, job_id -> Status
- work_exp_id -> student_id, company_id, role, job_type, years_of_exp

3NF vs BCNF:

We have meticulously designed our database to consist of relations that are both 3NF and BCNF. We reduced redundancy in the FDs by identifying their corresponding minimal basis.