CS 4110

Programming Languages & Logics

Lecture 5
The IMP Language

8 September 2014

Announcements

Office Hours

- Nitesh: Today at 10:30am-11:30am
- Nate: Today at 4-5pm
- Nitesh: Tuesday at 4:15pm 5:15pm
- Fran: Wednesday at 11-12pm

Homework #1

Due: Wednesday

We'll now consider a more realistic programming language...

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arithmetic expressions $a \in \mathbf{Aexp}$ $a := x \mid n \mid a_1 + a_2 \mid a_1 \times a_2$

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```
arithmetic expressions a \in \mathbf{Aexp} a :== x \mid n \mid a_1 + a_2 \mid a_1 \times a_2
boolean expressions b \in \mathbf{Bexp} b :== \mathbf{true} \mid \mathbf{false} \mid a_1 < a_2
```

We'll now consider a more realistic programming language...

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arithmetic expressions a \in \mathbf{Aexp} a :== x \mid n \mid a_1 + a_2 \mid a_1 \times a_2
boolean expressions b \in \mathbf{Bexp} b :== \mathbf{true} \mid \mathbf{false} \mid a_1 < a_2
commands c \in \mathbf{Com} c :== \mathbf{skip}
\mid x := a
\mid c_1; c_2
\mid \mathbf{if} \ b \ \mathbf{then} \ c_1 \ \mathbf{else} \ c_2
\mid \mathbf{while} \ b \ \mathbf{do} \ c
```

Three relations, one for each syntactic category:

$$\begin{array}{l} \rightarrow_{\mathsf{Aexp}} \subseteq (\mathsf{Store} \times \mathsf{Aexp}) \times (\mathsf{Store} \times \mathsf{Aexp}) \\ \\ \rightarrow_{\mathsf{Bexp}} \subseteq (\mathsf{Store} \times \mathsf{Bexp}) \times (\mathsf{Store} \times \mathsf{Bexp}) \\ \\ \rightarrow_{\mathsf{Com}} \subseteq (\mathsf{Store} \times \mathsf{Com}) \times (\mathsf{Store} \times \mathsf{Com}) \end{array}$$

$$\frac{n = \sigma(x)}{\langle \sigma, x \rangle \to \langle \sigma, n \rangle}$$

$$\frac{\langle \sigma, e_1 \rangle \to \langle \sigma, e'_1 \rangle}{\langle \sigma, e_1 + e_2 \rangle \to \langle \sigma, e'_1 + e_2 \rangle}$$

$$\frac{\langle \sigma, e_2 \rangle \to \langle \sigma, e'_2 \rangle}{\langle \sigma, n + e_2 \rangle \to \langle \sigma, n + e'_2 \rangle}$$

$$\frac{p = n + m}{\langle \sigma, n + m \rangle \to \langle \sigma, p \rangle}$$

$$\frac{\langle \sigma, e_1 \rangle \to \langle \sigma, e'_1 \rangle}{\langle \sigma, e_1 \times e_2 \rangle \to \langle \sigma, e'_1 \times e_2 \rangle}$$

$$\frac{\langle \sigma, e_2 \rangle \to \langle \sigma, e'_2 \rangle}{\langle \sigma, n \times e_2 \rangle \to \langle \sigma, n \times e'_2 \rangle}$$

$$\frac{p = n \times m}{\langle \sigma, n \times m \rangle \to \langle \sigma, p \rangle}$$

E

$$\frac{\langle \sigma, a_1 \rangle \to \langle \sigma, a_1' \rangle}{\langle \sigma, a_1 < a_2 \rangle \to \langle \sigma, a_1' < a_2 \rangle}$$

$$\frac{\langle \sigma, a_2 \rangle \to \langle \sigma, a_2' \rangle}{\langle \sigma, n < a_2 \rangle \to \langle \sigma, n < a_2' \rangle}$$

$$\frac{n < m}{\langle \sigma, n < m \rangle \to \langle \sigma, \mathbf{true} \rangle}$$

$$\frac{n \ge m}{\langle \sigma, n < m \rangle \to \langle \sigma, \mathbf{false} \rangle}$$

$$\frac{\langle \sigma, e \rangle \to \langle \sigma, e' \rangle}{\langle \sigma, x := e \rangle \to \langle \sigma, x := e' \rangle}$$

$$\langle \sigma, x := n \rangle \to \langle \sigma[x := n], \mathbf{skip} \rangle$$

$$\frac{\langle \sigma, c_1 \rangle \to \langle \sigma', c_1' \rangle}{\langle \sigma, c_1; c_2 \rangle \to \langle \sigma', c_1'; c_2 \rangle}$$

$$\frac{\langle \sigma, \mathsf{skip}; c_2 \rangle \to \langle \sigma, c_2 \rangle}{\langle \sigma, \mathsf{skip}; c_2 \rangle \to \langle \sigma, c_2 \rangle}$$

$$\frac{\langle \sigma, b \rangle \to \langle \sigma, b' \rangle}{\langle \sigma, \text{if } b \text{ then } c_1 \text{ else } c_2 \rangle \to \langle \sigma, \text{if } b' \text{ then } c_1 \text{ else } c_2 \rangle}$$

$$\langle \sigma, \text{if true then } c_1 \text{ else } c_2 \rangle \rightarrow \langle \sigma, c_1 \rangle$$

$$\langle \sigma, \text{ if false then } c_1 \text{ else } c_2 \rangle \rightarrow \langle \sigma, c_2 \rangle$$

 $\langle \sigma, \mathsf{while} \ b \ \mathsf{do} \ c \rangle \to \langle \sigma, \mathsf{if} \ b \ \mathsf{then} \ (c; \mathsf{while} \ b \ \mathsf{do} \ c) \ \mathsf{else} \ \mathsf{skip} \rangle$

 $\langle \sigma, \text{foo} := 3; \text{ while foo } < 4 \text{ do foo } := \text{foo} + 5 \rangle$

```
\langle \sigma, \text{foo} := 3; \text{ while foo} < 4 \text{ do foo} := \text{foo} + 5 \rangle
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\rightarrow \langle \sigma'', \text{ if false then (foo := foo + 5; } W \text{) else skip} \rangle
\rightarrow \langle \sigma'', \mathsf{skip} \rangle
```

Again, three relations, one for each syntactic category:

$$\label{eq:Aexp} \begin{array}{l} \Downarrow_{\mathsf{Aexp}} \subseteq \mathsf{Store} \times \mathsf{Aexp} \times \mathsf{Store} \\ \\ \Downarrow_{\mathsf{Bexp}} \subseteq \mathsf{Store} \times \mathsf{Bexp} \times \mathsf{Store} \\ \\ \Downarrow_{\mathsf{Com}} \subseteq \mathsf{Store} \times \mathsf{Com} \times \mathsf{Store} \end{array}$$

$$\frac{\sigma(x) = n}{\langle \sigma, n \rangle \Downarrow n}$$

$$\frac{\sigma(x) = n}{\langle \sigma, x \rangle \Downarrow n}$$

$$\langle \sigma, e_1 \rangle \Downarrow n_1 \qquad \langle \sigma, e_2 \rangle \Downarrow n_2 \qquad n = n_1 + n_2$$

$$\langle \sigma, e_1 + e_2 \rangle \Downarrow n$$

$$\langle \sigma, e_1 \rangle \Downarrow n_1 \qquad \langle \sigma, e_2 \rangle \Downarrow n_2 \qquad n = n_1 \times n_2$$

$$\langle \sigma, e_1 \rangle \Downarrow n_1 \qquad \langle \sigma, e_2 \rangle \Downarrow n_2 \qquad n = n_1 \times n_2$$

Skip
$$\overline{\langle \sigma, \mathbf{skip} \rangle \Downarrow \sigma}$$

Assgn
$$\frac{\langle \sigma, e \rangle \Downarrow n}{\langle \sigma, x := e \rangle \Downarrow \sigma[x \mapsto n]}$$

$$\mathsf{Seq} \frac{\langle \sigma, c_1 \rangle \Downarrow \sigma' \qquad \langle \sigma', c_2 \rangle \Downarrow \sigma''}{\langle \sigma, c_1; c_2 \rangle \Downarrow \sigma''}$$

If-T
$$\frac{\langle \sigma, b \rangle \Downarrow \mathsf{true}}{\langle \sigma, \mathsf{if} \ b \ \mathsf{then} \ c_1 \ \mathsf{else} \ c_2 \rangle \Downarrow \sigma'}$$
If-F $\frac{\langle \sigma, b \rangle \Downarrow \mathsf{false}}{\langle \sigma, \mathsf{if} \ b \ \mathsf{then} \ c_1 \ \mathsf{else} \ c_2 \rangle \Downarrow \sigma'}$

Command Equivalence

Intuitively, two commands are equivalent if they produce the same result under any store...

Definition (Equivalence of commands)

Two commands c and c' are equivalent (written $c \sim c'$) if, for any stores σ and σ' , we have

$$\langle \sigma, c \rangle \Downarrow \sigma' \iff \langle \sigma, c' \rangle \Downarrow \sigma'.$$

Command Equivalence

For example, we can prove that every **while** command is equivalent to its unfolding:

Theorem

For all $b \in \mathbf{Bexp}$ and $c \in \mathbf{Com}$ we have

while b do $c \sim$ if b then (c; while b do c) else skip.

Proof.

We show each implication separately...