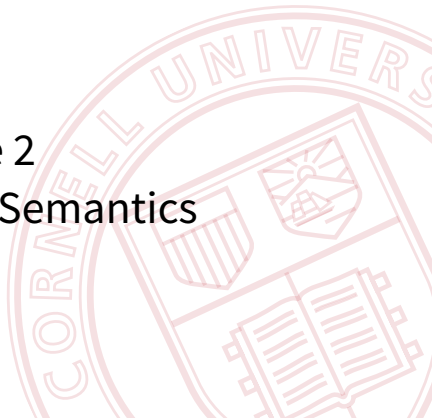# CS 4110

Programming Languages & Logics
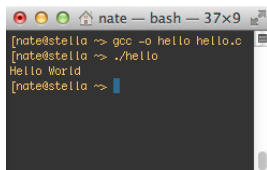
Lecture 2
Introduction to Semantics

Semantics

Question: What is the meaning of a program?

Semantics

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Answer: We could execute the program using an interpreter or a compiler, or we could consult a manual...

A8.7 Void

The (nonexistent) value of a void object may not be used in any way, and neither explicit nor implicit conversion to any non-void type may be applied. Because a void expression denotes a nonexistent value, such an expression may be used only where the value is not required, for example as an expression statement (§A9.2) or as the left operand of a comma operator (§A7.18).

An expression may be converted to type void by a cast. For example, a void cast documents the discarding of the value of a function call used as an expression statement.

void did not appear in the first edition of this book, but has become common since.

...but none of these is a satisfactory solution.

Formal Semantics

Three Approaches

- Operational $\qquad\qquad\qquad\qquad\qquad\qquad \langle \sigma, e \rangle \longrightarrow \langle \sigma', e' \rangle$
  - ▶ Model program by execution on abstract machine
  - ▶ Useful for implementing compilers and interpreters

- Denotational: $\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad [\![e]\!]$
  - ▶ Model program as mathematical objects
  - ▶ Useful for theoretical foundations

- Axiomatic $\qquad\qquad\qquad\qquad\qquad\qquad\qquad \vdash \{\phi\}\, e\, \{\psi\}$
  - ▶ Model program by the logical formulas it obeys
  - ▶ Useful for proving program correctness

Arithmetic Expressions

Syntax

A language of integer arithmetic expressions with assignment.

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BNF Grammar:

$$
\begin{aligned}
e ::=\ & x \\
| \ & n \\
| \ & e_1 + e_2 \\
| \ & e_1 * e_2 \\
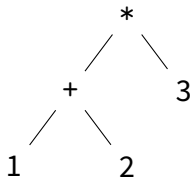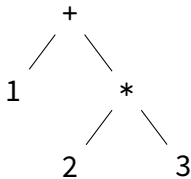| \ & x := e_1 ;\ e_2
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$$

Ambiguity

What expression does the string "$1 + 2 * 3$" describe?

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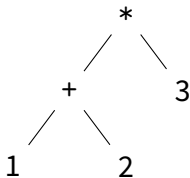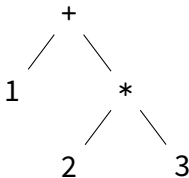
In this course, we will distinguish abstract syntax from concrete syntax, and focus primarily on abstract syntax (using conventions or parentheses at the concrete level to disambiguate as needed).

Representing Expressions

BNF Grammar:

$$e ::= x$$
$$| \; n$$
$$| \; e_1 + e_2$$
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$$| \; x := e_1 \; ; \; e_2$$

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BNF Grammar:

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$$\mid n$$
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$$\mid x := e_1 \; ; \; e_2$$

OCaml:

```
type exp = Var of string
         | Int of int
         | Add of exp * exp
         | Mul of exp * exp
         | Assgn of string * exp * exp
```

Example: Mul(Int 2, Add(Var "foo", Int 1))

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Java:

```
abstract class Expr { }
class Var extends Expr { String name; ... }
class Int extends Expr { int val; ... }
class Add extends Expr { Expr exp1, exp2; ... }
class Mul extends Expr { Expr exp1, exp2; ... }
class Assgn extends Expr { String var, Expr exp1, exp2; ... }
```

Example: new Mul(new Int(2), new Add(new Var("foo"), new Int(1)))

Quiz

- $7 + (4 * 2)$ evaluates to ...?

Quiz

- $7 + (4 * 2)$ evaluates to 15

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- $x + 1$ evaluates to error?

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- $i := 6 + 1 ; 2 * 3 * i$ evaluates to 42
- $x + 1$ evaluates to error?

The rest of this lecture will make these intuitions precise...

Mathematical Preliminaries

Binary Relations

The *product* of two sets A and B, written $A \times B$, contains all ordered pairs (a, b) with $a \in A$ and $b \in B$.

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Some Important Relations

- empty – $\emptyset$
- total – $A \times B$
- identity on A – $\{(a, a) \mid a \in A\}$.
- composition $R; S$ – $\{(a, c) \mid \exists b. (a, b) \in R \wedge (b, c) \in S\}$

Functions

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The *image* of *f* is the set of elements $b \in B$ that are mapped to by at least one $a \in A$. More formally: $\text{image}(f) \triangleq \{f(a) \mid a \in A\}$

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Given two functions $f : A \to B$ and $g : B \to C$, the composition of f and g is defined by: $(g \circ f)(x) = g(f(x))$ Note order!

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A function $f : A \to B$ is said to be *surjective* (or *onto*) if and only if the image of f is B.

Operational Semantics

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- and the expression e being evaluated

More formally,

$$\begin{aligned} \textbf{Store} &\triangleq \textbf{Var} \rightharpoonup \textbf{Int} \\ \textbf{Config} &\triangleq \textbf{Store} \times \textbf{Exp} \end{aligned}$$

Note that a store is a *partial* function from variables to integers.

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Intuitively, if facts above the line hold, then facts below the line hold. More formally, "$\rightarrow$" is the smallest relation "closed" under the inference rules.

Variables

$$\frac{n = \sigma(x)}{\langle \sigma, x \rangle \to \langle \sigma, n \rangle} \text{ VAR}$$

Addition

$$\frac{\langle \sigma, e_1 \rangle \rightarrow \langle \sigma', e_1' \rangle}{\langle \sigma, e_1 + e_2 \rangle \rightarrow \langle \sigma', e_1' + e_2 \rangle} \; \text{LADD}$$

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Multiplication

$$\frac{\langle \sigma, e_1 \rangle \rightarrow \langle \sigma', e_1' \rangle}{\langle \sigma, e_1 * e_2 \rangle \rightarrow \langle \sigma', e_1' * e_2 \rangle} \ \text{LMUL}$$

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$$\frac{p = m \times n}{\langle \sigma, m * n \rangle \rightarrow \langle \sigma, p \rangle} \text{ MuL}$$

Assignment

$$\frac{\langle \sigma, e_1 \rangle \rightarrow \langle \sigma', e_1' \rangle}{\langle \sigma, x := e_1 \, ; \, e_2 \rangle \rightarrow \langle \sigma', x := e_1' \, ; \, e_2 \rangle} \text{ ASSGN1}$$

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$$\frac{\sigma' = \sigma[x \mapsto n]}{\langle \sigma, x := n \, ; \, e_2 \rangle \to \langle \sigma', e_2 \rangle} \; \text{Assgn}$$

Notation: $\sigma[x \mapsto n]$ maps x to n and otherwise behaves like σ

Operational Semantics

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