CS 4110

Programming Languages & Logics

Lecture 5
The IMP Language

We'll now consider a more realistic programming language...

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arithmetic expressions $a \in \mathbf{Aexp}$ $a := x \mid n \mid a_1 + a_2 \mid a_1 \times a_2$

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arithmetic expressions a \in \mathbf{Aexp} a := x \mid n \mid a_1 + a_2 \mid a_1 \times a_2
Boolean expressions b \in \mathbf{Bexp} b := \mathbf{true} \mid \mathbf{false} \mid a_1 < a_2
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arithmetic expressions a \in \mathbf{Aexp} a := x \mid n \mid a_1 + a_2 \mid a_1 \times a_2 Boolean expressions b \in \mathbf{Bexp} b := \mathbf{true} \mid \mathbf{false} \mid a_1 < a_2 commands c \in \mathbf{Com} c := \mathbf{skip} \mid x := a \mid c_1; c_2 \mid if b then c_1 else c_2 \mid while b do c
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2

Three relations, one for each syntactic category:

$$egin{aligned} & \to_{\mathsf{Aexp}} \subseteq (\mathsf{Store} imes \mathsf{Aexp}) imes (\mathsf{Store} imes \mathsf{Aexp}) \\ & \to_{\mathsf{Bexp}} \subseteq (\mathsf{Store} imes \mathsf{Bexp}) imes (\mathsf{Store} imes \mathsf{Bexp}) \\ & \to_{\mathsf{Com}} \subseteq (\mathsf{Store} imes \mathsf{Com}) imes (\mathsf{Store} imes \mathsf{Com}) \end{aligned}$$

We'll typically just use \rightarrow where the specific relation we mean is clear from context.

For example:

$$\langle \sigma, \text{ if true then } x := 1 \text{ else } x := 2 \rangle$$

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 $\to_{\textbf{Com}} \langle \sigma, x := 1 \rangle$

For example:

$$\begin{split} &\langle \sigma, \text{if true then } x := 1 \text{ else } x := 2 \rangle \\ \to_{\textbf{Com}} &\langle \sigma, x := 1 \rangle \\ \to_{\textbf{Com}} &\langle \sigma[x \mapsto 1], \textbf{skip} \rangle \end{split}$$

Arithmetic expressions:

$$\frac{n = \sigma(x)}{\langle \sigma, x \rangle \to \langle \sigma, n \rangle}$$

Arithmetic expressions:

$$\begin{split} \frac{\langle \sigma, a_1 \rangle \to \langle \sigma, a_1' \rangle}{\langle \sigma, a_1 + a_2 \rangle \to \langle \sigma, a_1' + a_2 \rangle} & \frac{\langle \sigma, a_2 \rangle \to \langle \sigma, a_2' \rangle}{\langle \sigma, n + a_2 \rangle \to \langle \sigma, n + a_2' \rangle} \\ & \frac{p = n + m}{\langle \sigma, n + m \rangle \to \langle \sigma, p \rangle} \end{split}$$

Arithmetic expressions:

$$\begin{split} \frac{\langle \sigma, a_1 \rangle \to \langle \sigma, a_1' \rangle}{\langle \sigma, a_1 \times a_2 \rangle \to \langle \sigma, a_1' \times a_2 \rangle} & \frac{\langle \sigma, a_2 \rangle \to \langle \sigma, a_2' \rangle}{\langle \sigma, n \times a_2 \rangle \to \langle \sigma, n \times a_2' \rangle} \\ \frac{p = n \times m}{\langle \sigma, n \times m \rangle \to \langle \sigma, p \rangle} \end{split}$$

Boolean expressions:

Commands:

$$\frac{\langle \sigma, a \rangle \to \langle \sigma, a' \rangle}{\langle \sigma, x := a \rangle \to \langle \sigma, x := a' \rangle} \qquad \overline{\langle \sigma, x := n \rangle \to \langle \sigma[x := n], \mathsf{skip} \rangle}$$

Commands:

$$\frac{\langle \sigma, c_1 \rangle \to \langle \sigma', c_1' \rangle}{\langle \sigma, c_1; c_2 \rangle \to \langle \sigma', c_1'; c_2 \rangle} \qquad \qquad \overline{\langle \sigma, \mathbf{skip}; c_2 \rangle \to \langle \sigma, c_2 \rangle}$$

Commands:

$$\frac{\langle \sigma,b\rangle \to \langle \sigma,b'\rangle}{\langle \sigma,\text{if }b\text{ then }c_1\text{ else }c_2\rangle \to \langle \sigma,\text{if }b'\text{ then }c_1\text{ else }c_2\rangle}$$

$$\overline{\langle \sigma,\text{if true then }c_1\text{ else }c_2\rangle \to \langle \sigma,c_1\rangle}$$

$$\overline{\langle \sigma,\text{if false then }c_1\text{ else }c_2\rangle \to \langle \sigma,c_2\rangle}$$

7

Commands:

 $\langle \sigma, \mathsf{while}\, b \ \mathsf{do}\ c \rangle o \langle \sigma, \mathsf{if}\ b \ \mathsf{then}\ (c; \mathsf{while}\ b \ \mathsf{do}\ c) \ \mathsf{else}\ \mathsf{skip} \rangle$

Again three relations, one for each syntactic category:

$$\label{eq:Aexp} \begin{split} & \psi_{\mathsf{Aexp}} \subseteq (\mathsf{Store} \times \mathsf{Aexp}) \times \mathsf{Int} \\ & \psi_{\mathsf{Bexp}} \subseteq (\mathsf{Store} \times \mathsf{Bexp}) \times \mathsf{Bool} \\ & \psi_{\mathsf{Com}} \subseteq (\mathsf{Store} \times \mathsf{Com}) \times \mathsf{Store} \end{split}$$

And again, we'll typically just use \Downarrow where the specific relation we mean is clear from context.

$$\frac{\sigma(x) = n}{\langle \sigma, n \rangle \Downarrow n}$$

$$\frac{\langle \sigma, e_1 \rangle \Downarrow n_1 \qquad \langle \sigma, e_2 \rangle \Downarrow n_2 \qquad n = n_1 + n_2}{\langle \sigma, e_1 \rangle \Downarrow n_1}$$

$$\frac{\langle \sigma, e_1 \rangle \Downarrow n_1 \qquad \langle \sigma, e_2 \rangle \Downarrow n_2 \qquad n = n_1 \times n_2}{\langle \sigma, e_1 \rangle \Downarrow n_1}$$

$$\frac{\langle \sigma, e_1 \rangle \Downarrow n_1 \qquad \langle \sigma, e_2 \rangle \Downarrow n_2 \qquad n = n_1 \times n_2}{\langle \sigma, e_1 \times e_2 \rangle \Downarrow n}$$

S

C

SKIP

 $\overline{\langle \sigma, \mathsf{skip} \rangle \Downarrow \sigma}$

$$\frac{\langle \sigma, a \rangle \Downarrow n}{\langle \sigma, x := a \rangle \Downarrow \sigma[x \mapsto n]}$$

SEQ
$$\frac{\langle \sigma, c_1 \rangle \Downarrow \sigma' \qquad \langle \sigma', c_2 \rangle \Downarrow \sigma''}{\langle \sigma, c_1; c_2 \rangle \Downarrow \sigma''}$$

$$\frac{\text{IF-T}}{\langle \sigma, b \rangle \Downarrow \text{true}} \quad \langle \sigma, c_1 \rangle \Downarrow \sigma'}{\langle \sigma, \text{if } b \text{ then } c_1 \text{ else } c_2 \rangle \Downarrow \sigma'}$$

$$\begin{tabular}{ll} & \begin{tabular}{ll} & \begin{tabular}{ll}$$

C

$$\frac{\langle \sigma, b \rangle \Downarrow \mathbf{false}}{\langle \sigma, \mathbf{while} \ b \ \mathbf{do} \ c \rangle \Downarrow \sigma}$$