Exercises 4

Exercise 1 (First-Order Formulas)**.** Among the following formulas of first order, classical logic, four are theorems and two are not. Prove two of them using usual pen-and-paper (but still reasonably formal) mathematical proofs, and two using formal resolution proofs. For the two that are not theorems, provide describe models (base set and interpretation) where they do not hold.

- $(\exists y.\forall x.P(x,y)) \rightarrow \forall x.\exists y.P(x,y)$
- $\exists x.\forall y.(Q(x) \rightarrow Q(y))$
- $(\neg\forall x.\exists y.P(x,f(y))) \vee \exists y.\forall x.P(f(y),x)$
- $(\forall x,y,z.P(x,y) \wedge P(y,z) \rightarrow P(x,z)) \wedge (\forall x,y.P(x,y) \vee P(y,x)) \rightarrow \forall x.P(x,x)$
- $\exists x.P(x) \wedge P(f(x)) \rightarrow \forall y.P(y) \vee P(f(y))$
- $\forall x.\forall y.P(x,y) \rightarrow \forall x.\forall y.P(y,x)$

Exercise 2 (Weakest Prediction)**.** Recall the definition of weakest precondition:

$$\mathrm{wp}(r,Q) = \{s \mid \forall s'.\ (s,s') \in r \rightarrow s' \in Q\}$$

1. Prove or disprove the following properties:
   - $\mathrm{wp}(r_1 \cup r_2, Q) = \mathrm{wp}(r_1,Q) \cup \mathrm{wp}(r_2,Q)$
   - $\mathrm{wp}(r_1 \cup r_2, Q) = \mathrm{wp}(r_1,Q) \cap \mathrm{wp}(r_2,Q)$
   - $\mathrm{wp}(r, Q_1 \cap Q_2) = \mathrm{wp}(r,Q_1) \cap \mathrm{wp}(r,Q_2)$
   - $\mathrm{wp}(r, Q_1 \cup Q_2) = \mathrm{wp}(r,Q_1) \cup \mathrm{wp}(r,Q_2)$

   For those that are wrong, do any of them hold if r is restricted to being functional, i.e. if r satisfies

   $$\forall x,y_1,y_2.\ ((x,y_1) \in r \wedge (x,y_2) \in r) \rightarrow (y_1 = y_2)$$

2. Let $r \subseteq S \times S$ and $Q \subseteq S$. Give an expression defining weakest precondition $\mathsf{wp}(r,Q)$ using operations of inverse of a relation, $(\_^{-1})$, set difference $(\setminus)$, and image of a relation under a set, $(\_[\_])$. Prove

that your expression is correct by expanding the definitions of wp as well as of relational and set operations.

Exercise 3 (Programing With Integers)**.** Consider the following program:

```
1   case class Container(var x: INT, var y: INT):
2     def fun: Unit = {
3       require(x > 0 && y > 0)
4       if x > y then
5         x = x+y
6         y = x−y
7         x = x−y
8       else
9         y = y−x
10    }.ensuring(2∗old(this).x + old(this).y > 2∗x + y &&
11                 x >= 0 && y >= 0)
12  end Container
```

1. Compute R(fun) formally, by expressing all intermediate formulas corresponding to subprograms.
2. Write the formula expressing the correctness of the ensuring clause. Is the formula valid when INT denotes mathematical integers with their usual operations (**type** INT = BigInt in Scala)?
3. Is the the verification condition formula valid for machine integers,

$$Z_{2^{32}} = \{-2^{31}, \ldots, -1, 0, 1 \ldots, 2^{31} - 1\}$$

and where operations are intepreted as the usual machine arithmetic (**type** INT = Int in Scala)?

Exercise 4 (Hoare Logic Proof)**.** Give a complete Hoare logic proof for the following program:

```
{n >= 0 && d > 0}
  q = 0
  r = n
  while ( r >= d ) {
    q = q + 1
    r = r - d
  }
{n == q * d + r && 0 <= r < d}
```

The proof should include step-by-step annotation for each line of the program, as in the example proof in the lecture.

Exercise 5 (Iterating a Relation)**.** Let $M = (S, I, r, A)$ be a transition system and $\bar{r} = \{(s, s') \mid (s, a, s') \in r\}$, as usual. Let $\Delta = \{(x, x) \mid x \in S\}$.

Let $\bar{r}^k$ denote the usual composition of relation $\bar{r}$ with itself k times.

Define the sequence of relations r_n, for all non-negative integers n, as follows:

- $r_0 = \Delta \cup \bar{r}$
- $r_{n+1} = r_n \circ r_n$

1. Prove that $r_n \subseteq r_{n+1}$ for every n.
2. Prove that for every n and every k where $0 \leq k \leq 2^n$ we have $\bar{r}^k \subseteq r_n$.
3. Suppose that S is finite. Find a bound B as a function of $|S|$ such that
$$\mathsf{Reach}(M) \subseteq r_B[I]$$
Aim to find as small bound as possible.

Exercise 6 (Approximating Relations)**.** Consider a guarded command language whose meanings are binary relations on the set states U.

Let $E(a_1, \ldots, a_n)$ denote an expression built from some atomic relations $a_1, \ldots, a_n$, as well as diagonal relations

$$\Delta_P = \{(x, x) \mid x \in P\}$$

for various sets $P \subseteq U$. The expression E is built from these relations using union (to model non-deterministic choice), relation composition (to represent sequential composition) and transitive closure (to represent loops).

Let us call a relation $s \subseteq U \times U$ an *effect* if it is reflexive (R) and transitive (T).

1. Prove that if s is an effect and $a_i \subseteq s$ for all $1 \leq i \leq n$, then
$$E(a_1, \ldots, a_n) \subseteq s$$

2. Let $U = \mathbb{Z}^2$ denote pairs of integers, denoted by integer variables x, y. Let s be a specification relation given by the formula:
$$s = \{((x, y), (x', y')) \mid y \geq 0 \to (x' \leq x \wedge y' \geq 0)\}$$
Show that s is an effect.