Functional Data Structures

Exercise Sheet 12

Exercise 12.1 Sparse Binary Numbers

lemma size_snat:

Implement operations carry, inc, and add on sparse binary numbers, analogously to the operations link, ins, and meld on binomial heaps.

Show that the operations have logarithmic worst-case complexity.

```
type\_synonym \ rank = nat
type\_synonym \ snat = "rank \ list"
abbreviation invar :: "snat \Rightarrow bool" where "invar s \equiv sorted\_wrt (<) s"
definition \alpha :: "snat \Rightarrow nat" where "\alpha s = sum_list (map ((^) 2) s)"
\mathbf{lemmas}\ [\mathit{simp}] = \ \mathit{sorted\_wrt\_append}
fun carry :: "rank <math>\Rightarrow snat \Rightarrow snat"
lemma carry\_invar[simp]:
 assumes "invar rs"
 shows "invar (carry r rs)"
lemma carry\_\alpha:
 assumes "invar rs"
 assumes "\forall r' \in set \ rs. \ r \leq r'"
 shows "\alpha (carry r rs) = 2\hat{r} + \alpha rs"
definition inc :: "snat \Rightarrow snat"
lemma inc\_invar[simp]: "invar\ rs \implies invar\ (inc\ rs)"
lemma inc_{-}\alpha[simp]: "invar\ rs \Longrightarrow \alpha\ (inc\ rs) = Suc\ (\alpha\ rs)"
\mathbf{fun} \ add :: \ ``snat \Rightarrow snat"
lemma add\_invar[simp]:
 assumes "invar rs<sub>1</sub>"
 assumes "invar rs<sub>2</sub>"
 shows "invar (add rs<sub>1</sub> rs<sub>2</sub>)"
lemma add\_\alpha[simp]:
 assumes "invar rs<sub>1</sub>"
 assumes "invar rs_2"
 shows "\alpha (add rs_1 \ rs_2) = \alpha \ rs_1 + \alpha \ rs_2"
\mathbf{thm}\ sorted\_wrt\_less\_sum\_mono\_lowerbound
```

```
assumes "invar rs" shows "2^length rs \le \alpha \ rs + 1" fun t\_carry :: "rank \Rightarrow snat \Rightarrow nat" definition t\_inc :: "snat \Rightarrow nat" lemma t\_inc\_bound: assumes "invar rs" shows "t\_inc rs \le log \ 2 \ (\alpha \ rs + 1) + 1" fun t\_add :: "snat \Rightarrow snat \Rightarrow nat" lemma t\_add\_bound: fixes rs_1 \ rs_2 defines "n_1 \equiv \alpha \ rs_1" defines "n_2 \equiv \alpha \ rs_2" assumes INVARS: "invar rs_1" "invar rs_2" shows "t\_add rs_1 \ rs_2 \le 4*log \ 2 \ (n_1 + n_2 + 1) + 2"
```

Homework 12.1 Modified Binomial Heaps

Submission until Friday, 17. 7. 2020, 10:00am.

In its simplest form, a binomial heap can be implemented using binomial trees that store the the rank of every tree in its root. Such an Isabelle/HOL implementation can be found at: "src/HOL/Data_Structures/Binomial_Heaps.thy"

One optimisation is to eliminate the redundancy of storing ranks in the roots of every tree, and instead store the ranks only at the top level by pairing every tree with its rank in the heap. The following types describe a binomial heap with this optimisation.

```
datatype 'a tree = Node (root: 'a) (children: "'a tree list")

type_synonym 'a heap = "(nat*'a tree) list"
```

For such a heap to be a binomial heap it has to conform to the invariant *invar* defined as follows:

```
fun invar\_btree :: "nat \Rightarrow 'a::linorder tree \Rightarrow bool" where "invar\_btree \ r \ (Node \ x \ ts) \longleftrightarrow length \ ts = r \land (\forall (r',t) \in set \ (zip \ (rev \ [0..< r]) \ ts). \ invar\_btree \ r' \ t)" definition invar\_bheap :: "'a::linorder \ heap \Rightarrow bool" where "invar\_bheap \ ts \longleftrightarrow (\forall (r,t) \in set \ ts. \ invar\_btree \ r \ t) \land (sorted\_wrt \ (<) \ (map \ fst \ ts))" fun invar\_otree :: "'a::linorder \ tree \Rightarrow bool" where "invar\_otree :: "'a::linorder \ tree \Rightarrow bool" where "invar\_otree \ (Node \ x \ ts) \longleftrightarrow (\forall \ t \in set \ ts. \ invar\_otree \ t \land \ x \leq root \ t)"
```

```
definition invar\_oheap :: "'a::linorder tree list <math>\Rightarrow bool" where "invar\_oheap ts \longleftrightarrow (\forall t \in set ts. invar\_otree t)"
```

```
definition invar :: "'a::linorder \ heap \Rightarrow bool" where "invar \ ts \longleftrightarrow invar\_bheap \ ts \land invar\_oheap \ (map \ snd \ ts)"
```

In this homework you are required to define an insertion and a merging functions for this heap and show that they preserve the elements of their inputs as well as produce heaps that conform to the invariant *invar*.

```
definition insert :: "'a::linorder \Rightarrow 'a heap \Rightarrow 'a heap" where
lemma invar_insert[simp]: "invar t \Longrightarrow invar (insert x t)"
lemma mset_heap_insert[simp]: "mset_heap (insert x t) = {#x#} + mset_heap t"
fun merge :: "'a::linorder heap \Rightarrow 'a heap \Rightarrow 'a heap" where
lemma invar_merge[simp]: "[ invar ts_1; invar ts_2 ] \Longrightarrow invar (merge ts_1 ts_2)"
lemma mset_heap_merge[simp]: "mset_heap (merge ts_1 ts_2) = mset_heap ts_1 + mset_heap ts_2"
```

Hint: you can start with the theory file "src/HOL/Data_Structures/Binomial_Heaps.thy" that has an implementation of binomial heaps without this optimisation, and then edit it.

Homework 12.2 Be Original!

Submission until Friday, 17. 7. 2020, 10:00am. Develop a nice Isabelle formalisation yourself!

- This homework goes in parallel to other homeworks for the rest of the lecture period. From next sheet on, we will reduce regular homework load a bit, such that you have a time-frame of 3 weeks with reduced regular homework load.
- This homework will yield 15 points (for minimal solutions). Additionally, up to 15 bonus points may be awarded for particularly nice/original/etc solutions.
- You may develop a formalisation from all areas, not only functional data structures.
- Document your solution, such that it is clear what you have formalised and what your main theorems state!
- Set yourself a time frame and some intermediate/minimal goals. Your formalisation needs not be universal and complete after 3 weeks.
- You are welcome to discuss the realisability of your project with the tutor or ask him for possible ideas!

• Should you need inspiration to find a project: Sparse matrices, skew binary numbers, arbitrary precision arithmetic (on lists of bits), interval data structures (e.g. interval lists), spatial data structures (quad-trees, oct-trees), Fibonacci heaps, prefix tries/arrays and BWT, etc.