

Functional Data Structures

Exercise Sheet 11

Exercise 11.1 Insert for Leftist Heap

- Define a function to directly insert an element into a leftist heap. Do not construct an intermediate heap like insert via merge does!
- Show that your function is correct
- Define a timing function for your insert function, and show that it is linearly bounded by the rank of the tree.

```

fun lh_insert :: "'a::ord  $\Rightarrow$  'a heap  $\Rightarrow$  'a heap"
lemma set_lh_insert: "set_tree (lh_insert x t) = set_tree t  $\cup$  {x}"
lemma "heap t  $\Longrightarrow$  heap (lh_insert x t)"
lemma "ltree t  $\Longrightarrow$  ltree (lh_insert x t)"
fun t_lh_insert :: "'a::ord  $\Rightarrow$  'a heap  $\Rightarrow$  nat"
lemma "t_lh_insert x t  $\leq$  rank t + 1"
    
```

Exercise 11.2 Bootstrapping a Priority Queue

Given a generic priority queue implementation with $O(1)$ *empty*, *is_empty* operations, $O(f_1\ n)$ *insert*, and $O(f_2\ n)$ *get_min* and *del_min* operations.

Derive an implementation with $O(1)$ *get_min*, and the asymptotic complexities of the other operations unchanged!

Hint: Store the current minimal element! As you know nothing about f_1 and f_2 , you must not use *get_min*/*del_min* in your new *insert* operation, and vice versa!

For technical reasons, you have to define the new implementations type outside the locale!

```

datatype ('a,'s) bs_pq =
locale Bs_Priority_Queue =
  orig: Priority_Queue where
    empty = orig_empty and
    is_empty = orig_is_empty and
    insert = orig_insert and
    get_min = orig_get_min and
    del_min = orig_del_min and
    
```

```

    invar = orig_invar and
    mset = orig_mset
  for orig_empty orig_is_empty orig_insert orig_get_min orig_del_min orig_invar
  and orig_mset :: "'s ⇒ 'a::linorder multiset"
begin

```

In here, the original implementation is available with the prefix *orig*, e.g.

```

term orig_empty term orig_invar
thm orig.invar_empty

definition empty :: "('a,'s) bs_pq"

fun is_empty :: "('a,'s) bs_pq ⇒ bool"

fun insert :: "'a ⇒ ('a,'s) bs_pq ⇒ ('a,'s) bs_pq"

fun get_min :: "('a,'s) bs_pq ⇒ 'a"

fun del_min :: "('a,'s) bs_pq ⇒ ('a,'s) bs_pq"

fun invar :: "('a,'s) bs_pq ⇒ bool"

fun mset :: "('a,'s) bs_pq ⇒ 'a multiset"

lemmas [simp] = orig.is_empty orig.mset_get_min orig.mset_del_min
  orig.mset_insert orig.mset_empty
  orig.invar_empty orig.invar_insert orig.invar_del_min

```

Show that your new implementation satisfies the priority queue interface!

```

sublocale Priority_Queue
  where empty = empty
  and is_empty = is_empty
  and insert = insert
  and get_min = get_min
  and del_min = del_min
  and invar = invar
  and mset = mset
  apply unfold_locales
proof goal_cases
  case 1
  then show ?case
  next
  case (2 q)
  qed
end

```

Homework 11.1 Converting a 2-3 tree into a heap

Submission until Friday, 10. 7. 2020, 10:00am.

The following predicate describes the heap property for a binary tree.

```
fun heap:: "'a::linorder tree  $\Rightarrow$  bool" where  
  "heap Leaf = True"  
| "heap (Node l x r) = (( $\forall y \in \text{set\_tree } l. x \leq y$ )  $\wedge$  ( $\forall y \in \text{set\_tree } r. x \leq y$ )  $\wedge$  heap l  $\wedge$  heap r)"
```

Recall the function *sift_down* from the AFP entry *Priority_Queue_Braun*. Define an equivalent function for sifting the root of a 2-3 tree. Hint: you can do that by firstly converting the 2-3 tree into a binary tree, and then defining a function similar to the one in the AFP entry, but for binary trees. That function has to include extra cases that account for the fact that, unlike a Braun tree, a binary tree is not necessarily balanced.

Define a function *heapify* which, given a 2-3 tree, reorders the elements of the 2-3 tree into a heap. That function has to use the function *sift_down*. Show that the function indeed creates a heap and that it preserves the elements in the given binary tree.

```
fun heapify:: "'a::linorder tree23  $\Rightarrow$  'a::linorder tree" where  
lemma "heap (heapify t)"  
lemma "mset (inorder (heapify t)) = mset (Tree23.inorder t)"
```

Homework 11.2 Be Original!

Submission until Friday, 17. 7. 2020, 10:00am.

Develop a nice Isabelle formalisation yourself!

- This homework goes in parallel to other homeworks for the rest of the lecture period. From next sheet on, we will reduce regular homework load a bit, such that you have a time-frame of 3 weeks with reduced regular homework load.
- This homework will yield 15 points (for minimal solutions). Additionally, up to 15 bonus points may be awarded for particularly nice/original/etc solutions.
- You may develop a formalisation from all areas, not only functional data structures.
- Document your solution, such that it is clear what you have formalised and what your main theorems state!
- Set yourself a time frame and some intermediate/minimal goals. Your formalisation needs not be universal and complete after 3 weeks.
- You are welcome to discuss the realisability of your project with the tutor or ask him for possible ideas!
- Should you need inspiration to find a project: Sparse matrices, skew binary numbers, arbitrary precision arithmetic (on lists of bits), interval data structures (e.g. interval lists), spatial data structures (quad-trees, oct-trees), Fibonacci heaps, prefix tries/arrays and BWT, etc.