

```

expr := <number> | <name> | true | false
      | (if <expr> <expr> <expr>)
      | (let (<name> <expr>) <expr>)
      | (+ <expr> <expr>)
      | (< <expr> <expr>)
      | (set <name> <expr>)
      | (fun (<name> : <t>) : <t> <expr>)
      | (<expr> <expr>)

```

t := Num | Bool | (<t> → <t>)

prog := <expr>

Defs?! Replaced by fun!

```

type expr =
  | ENum of int | EBool of bool | EId of string
  | EIf of expr * expr * expr
  | ELet of string * expr * expr
  | EPlus of expr * expr
  | ELess of expr * expr
  | ESet of string * expr
  | EApp of string * expr * expr
  | EFun of string * typ * expr

```

and typ = TNum | TBool | TArrow of typ \* typ

type prog = expr

What's weird/  
interesting/  
wrong?

```

expr := <number> | <name> | true | false
      | (if <expr> <expr> <expr>)
      | (let (<name> <expr>) <expr>)
      | (+ <expr> <expr>)
      | (< <expr> <expr>)
      | (set <name> <expr>)
      | (<name> <expr> <expr>)

```

def := (def <name> (<name> : <t>) : <t>  
 <expr>)

t := Num | Bool

prog := def ... <expr>

```

type expr =
  | ENum of int | EBool of bool | EId of string
  | EIf of expr * expr * expr
  | ELet of string * expr * expr
  | EPlus of expr * expr
  | ELess of expr * expr
  | ESet of string * expr
  | EApp of string * expr

```

type def =  
 | DFun of string \* string \* typ \* typ \* expr

type typ = TNum | TBool

type prog = def list \* expr

(let (twx (fun (x : Num) : Num (+ x x)))

(twx 10) )

(let (and (fun (b1 : Bool) : (Bool → Bool)

(fun (b2 : Bool) : Bool (if b1 b2 false))))

((and true) false))

and : Bool → (Bool → Bool) type of and

closes over b1  
remembers b1

$$\frac{\Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash (e_1 e_2) : \tau_2}$$

$$\frac{(x, \tau) :: \Gamma \vdash e : \tau_R}{\Gamma \vdash (\text{fun } (x: \tau) : \tau_R \ e) : \tau \rightarrow \tau_R}$$

$((\text{fun } (x: \text{Num}) : \text{Num } (+ x 1)) 10)$

Did we lose anything meaningful?

A: Yes

B: No

[ Recursion?  
Mutual recursion?  
Program structure

$(\text{let } \boxed{\text{sum}} (\text{fun } (n: \text{Num}) : \text{Num}$   
 $(\text{if } (< n 1) 0$   
 $(+ (\boxed{\text{sum}} (- n 1) n))))$

unbound id!

X

$(\text{sum } 10))$

$(\text{letrec } (\text{sum } \dots) (\text{sum } 10))$

↓ means

$(\text{let } (\text{sum } (\text{null } (\text{Num} \rightarrow \text{Num})))$   
 $(\text{set sum } (\text{fun } \dots))$

$\vdots$   
 $(\text{sum } 10))$

$(\text{letrec } (\text{sum } \dots \text{sum } \dots))$

↑  
new error using  
sum before  
defined

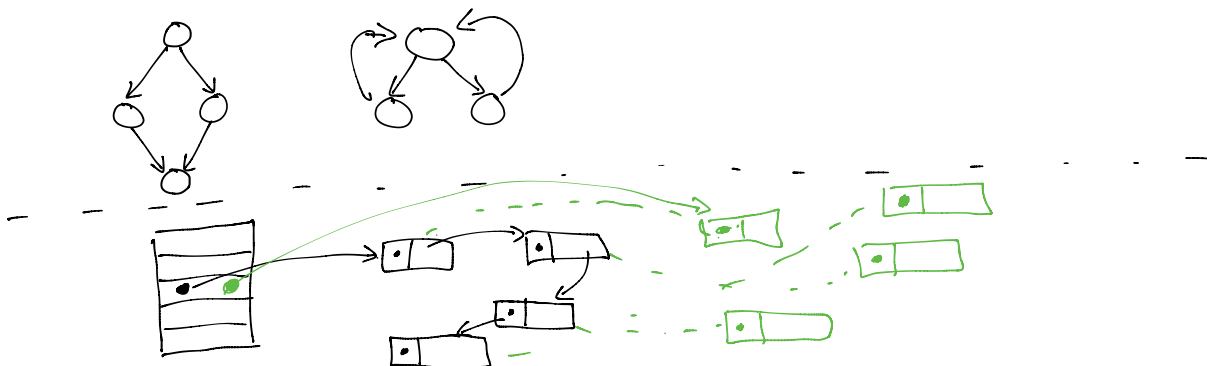
EFun( $wg, -, -, e$ )  $\rightarrow$   
 let body-is = compile e  $[(arg, 3)]$

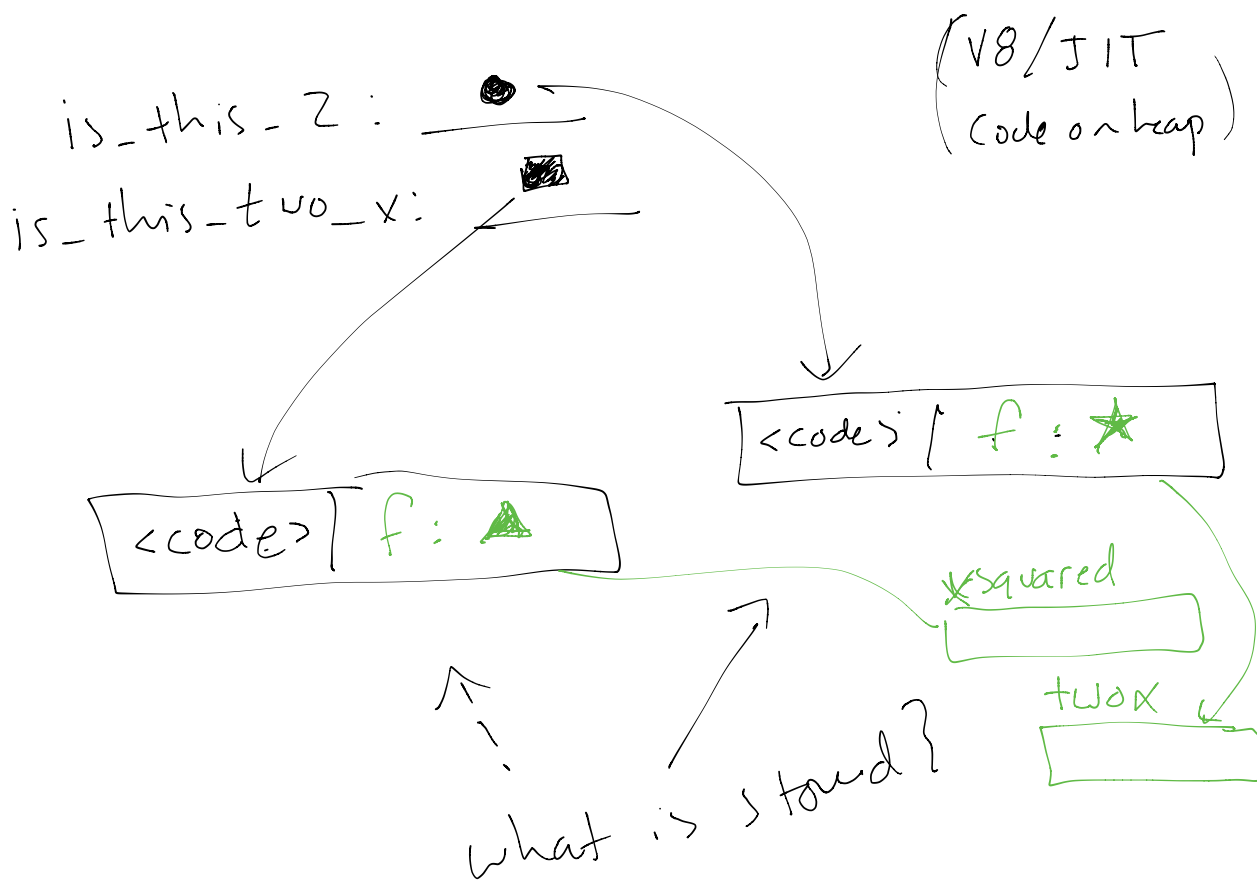
```
" jmp skip
label:
  body-is
  ret

skip:
  mov [r15], label
  mov [r15+8], unband/free id1
  mov [r15+16], unband/free id2
```

EApp( $ef, ea$ )  $\rightarrow$   
 let f-is = compile ef ... in  
 " mov rbx, [rax]  
 ... move args ...  
 jmp rbx "

Checking for cycles is not the same as checking for visiting same node twice.





<code>:  
isprime:  
~~~~~  
~~~~~  
ret

```
def ddx(f):
    def fprime(x):
        return (f(x + 0.001) - f(x)) / 0.001
    return fprime
```

f is free  
or unbound  
relative to  
fprime