BELLMAN-FORD



Shortest path Algorithm



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Overview:

- 1 <u>non-greedy algorithm</u>
- 2 Time complexity : O(V.E)

Assumptions

Cycles are non-negative

All paths are simple

At most V-1 iterations

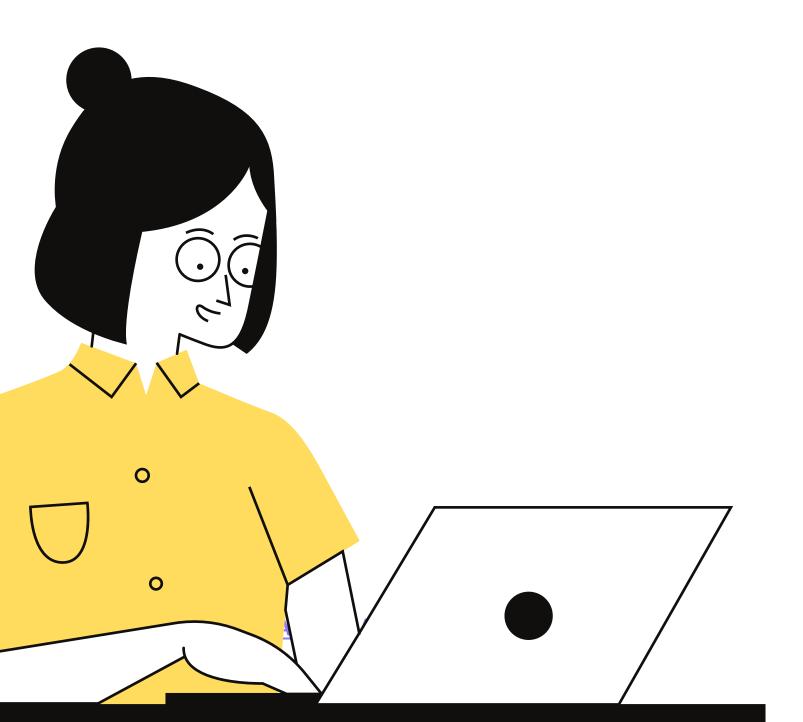
When to use

Works for negative edge weights

Dosen't work for negative cycles



Pseudo code



```
function bellmanFord(G, S)
for each vertex V in G
 distance[V] <- infinite
  previous[V] <- NULL</pre>
distance[S] <- 0
for each vertex V in G
 for each edge (U,V) in G
  tempDistance <- distance[U] + edge_weight(U, V)
  if tempDistance < distance[V]</pre>
   distance[V] <- tempDistance
   previous[V] <- U
for each edge (U,V) in G
 If distance[U] + edge_weight(U, V) < distance[V]
  Error: Negative Cycle Exists
return distance[], previous[]
```

Example

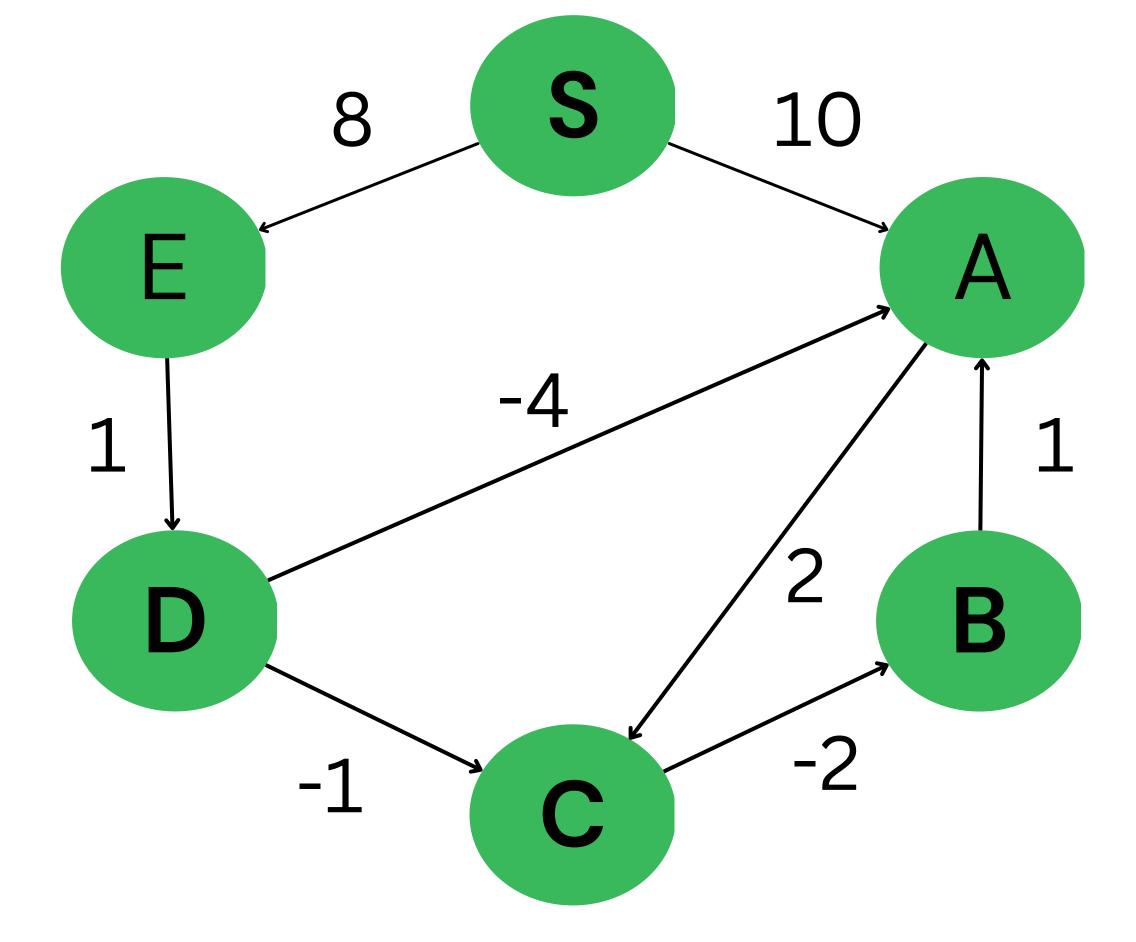
A cab driver in a city want to know which path is most profitable so that he can increase his rides in that particular path. He even choose a source S and wants the most profitable paths

Plan

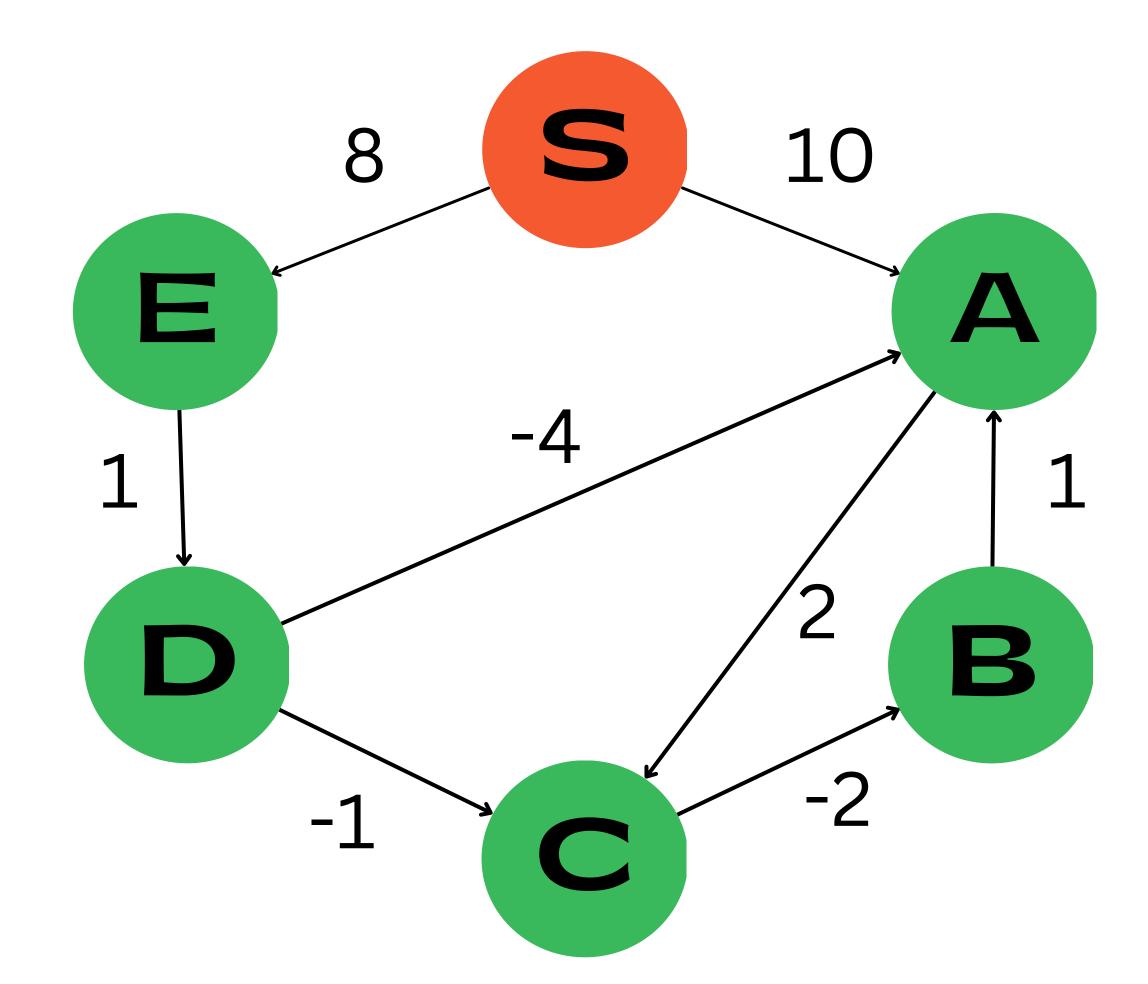
so to solve this problem we build a graph of possible cost and when ever there is profit we indicate it as negative weight because cost reduce in that particular path

Example

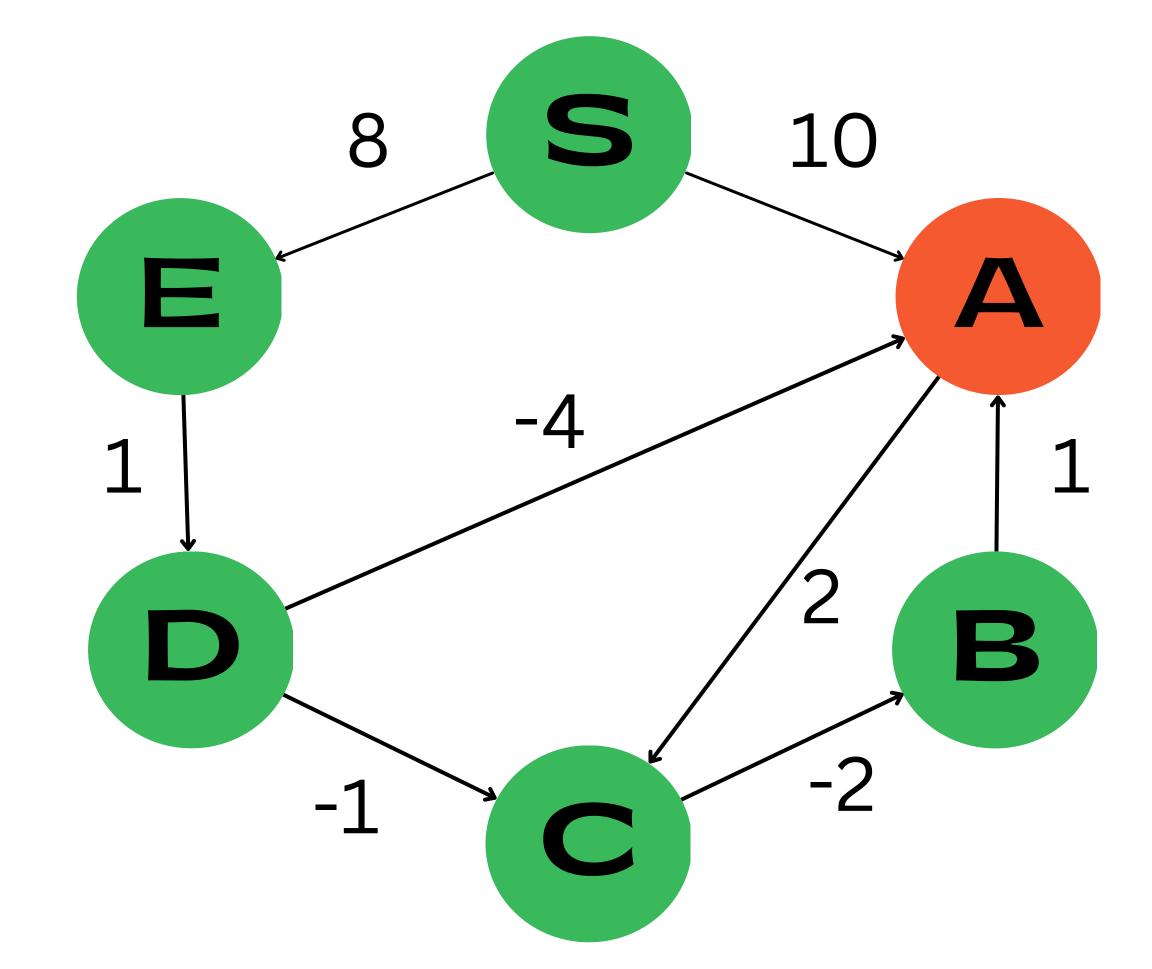
B ∞ D



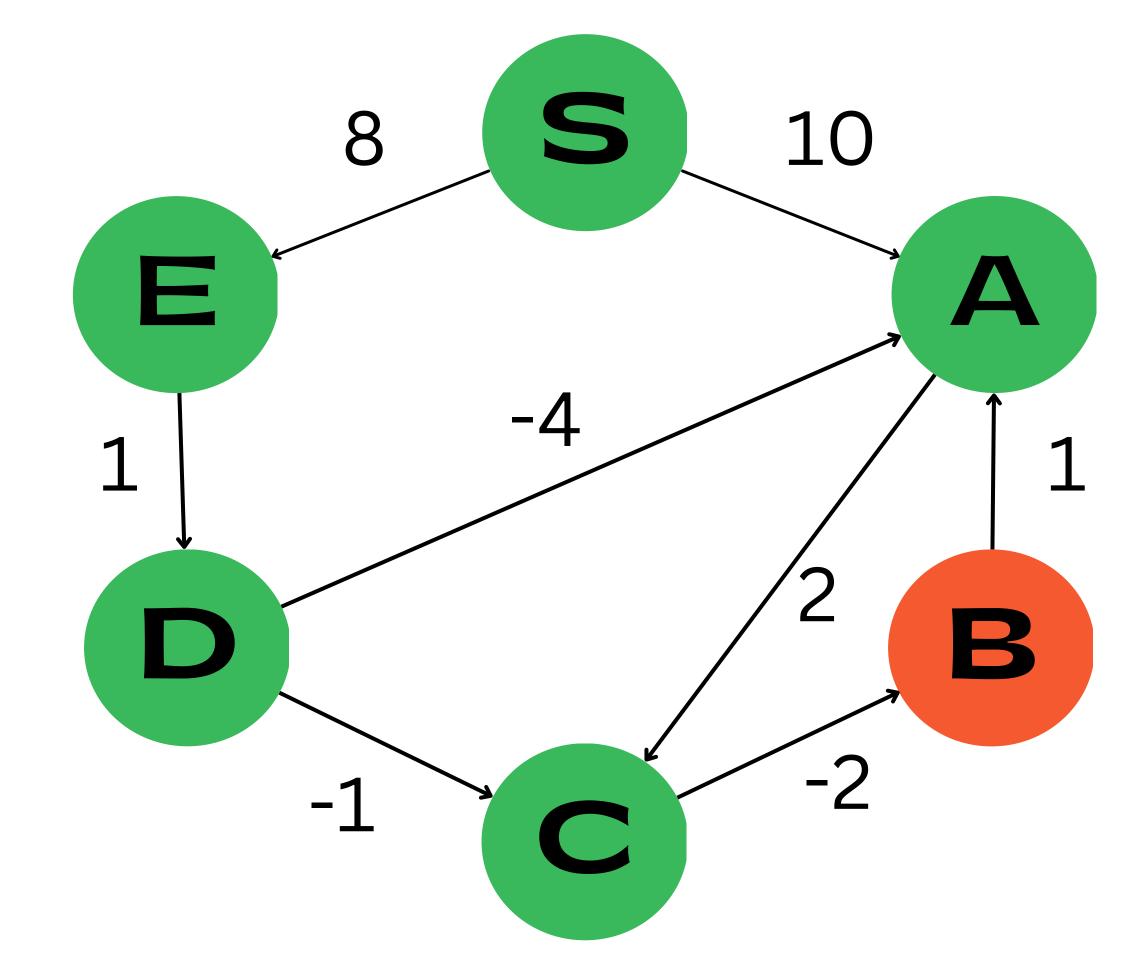
S	0
A	10
В	00
C	00
D	00
E	8



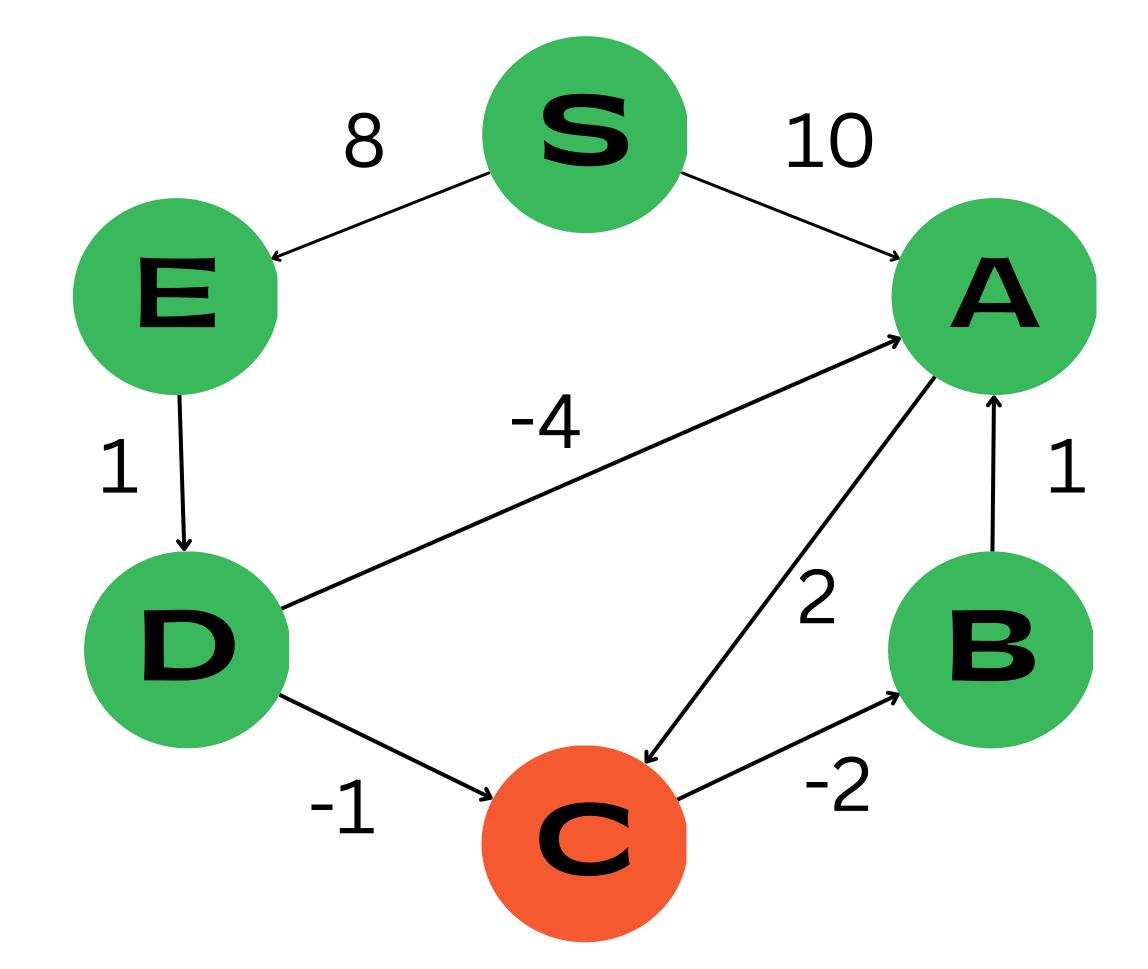
S	0
A	10
В	00
C	12
D	00
E	8



S	0
A	10
В	00
C	12
D	00
E	8

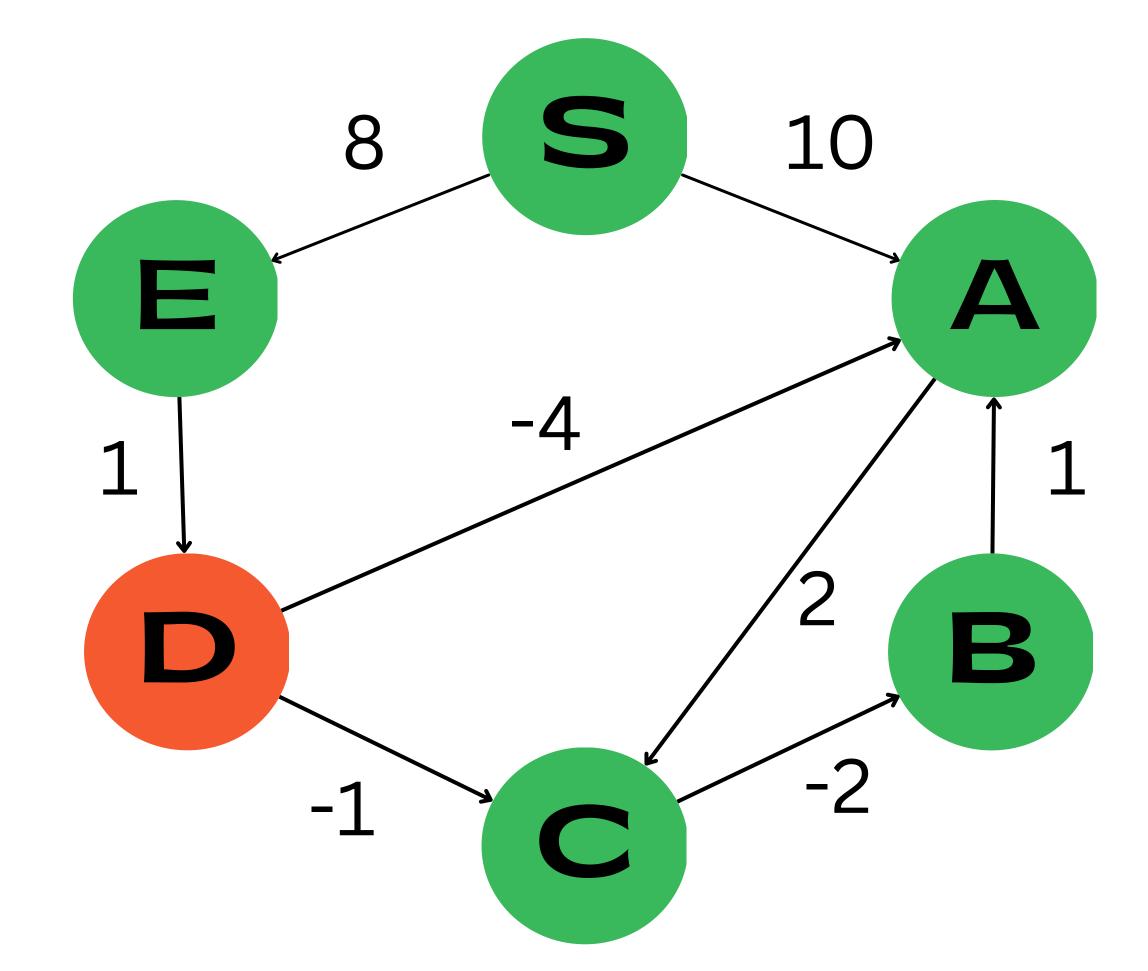


S	0
A	10
В	10
C	12
D	00
E	8

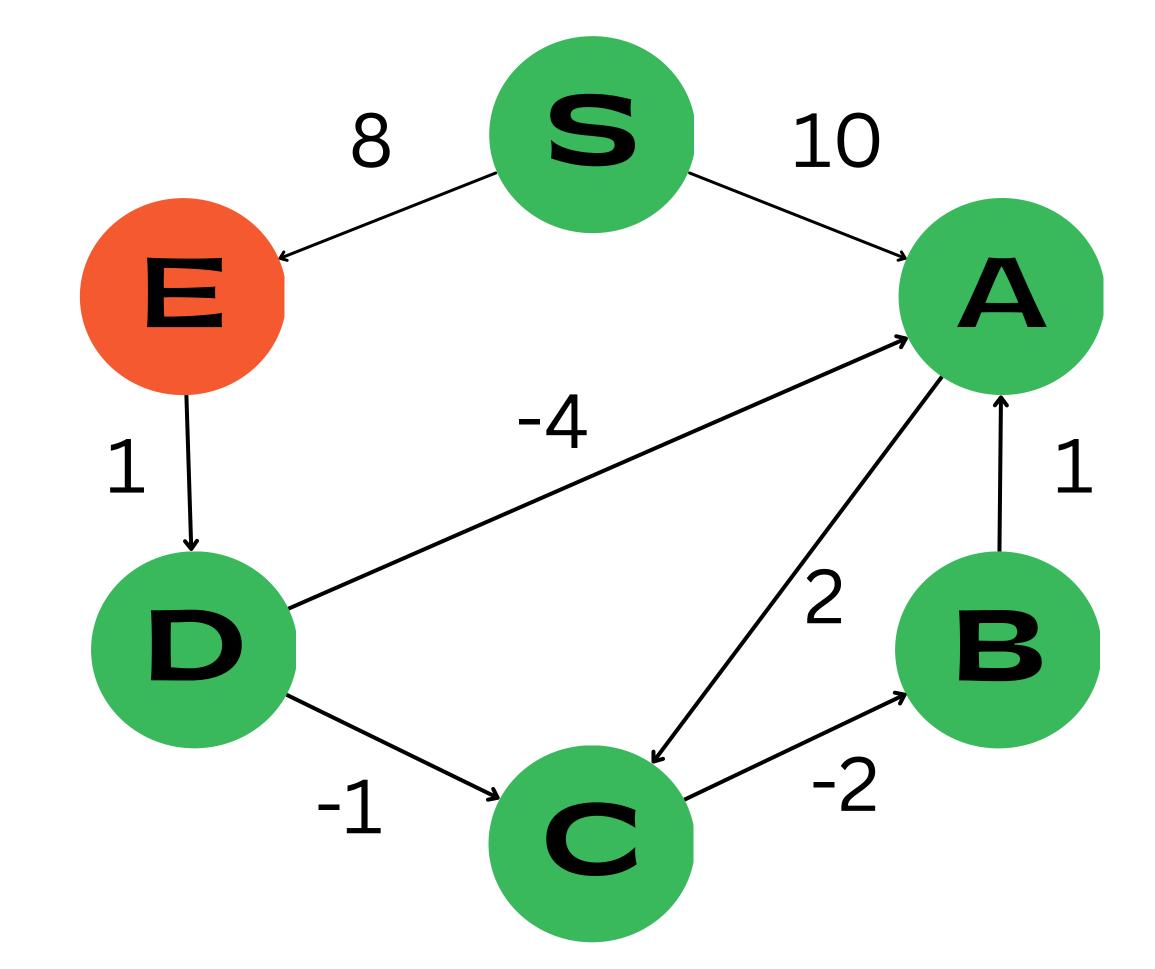


AGENDA PAGE

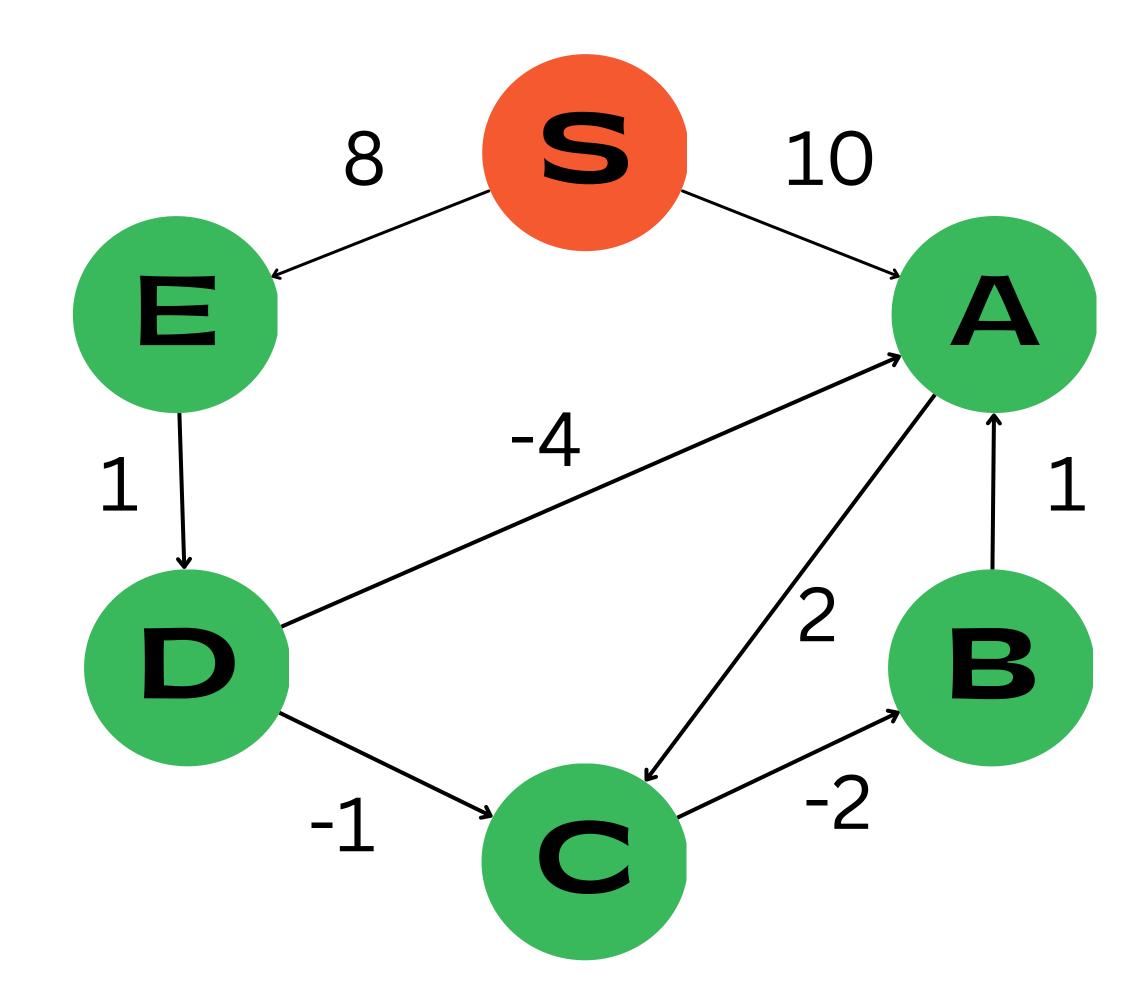
S	0
A	10
В	10
C	12
D	00
E	8



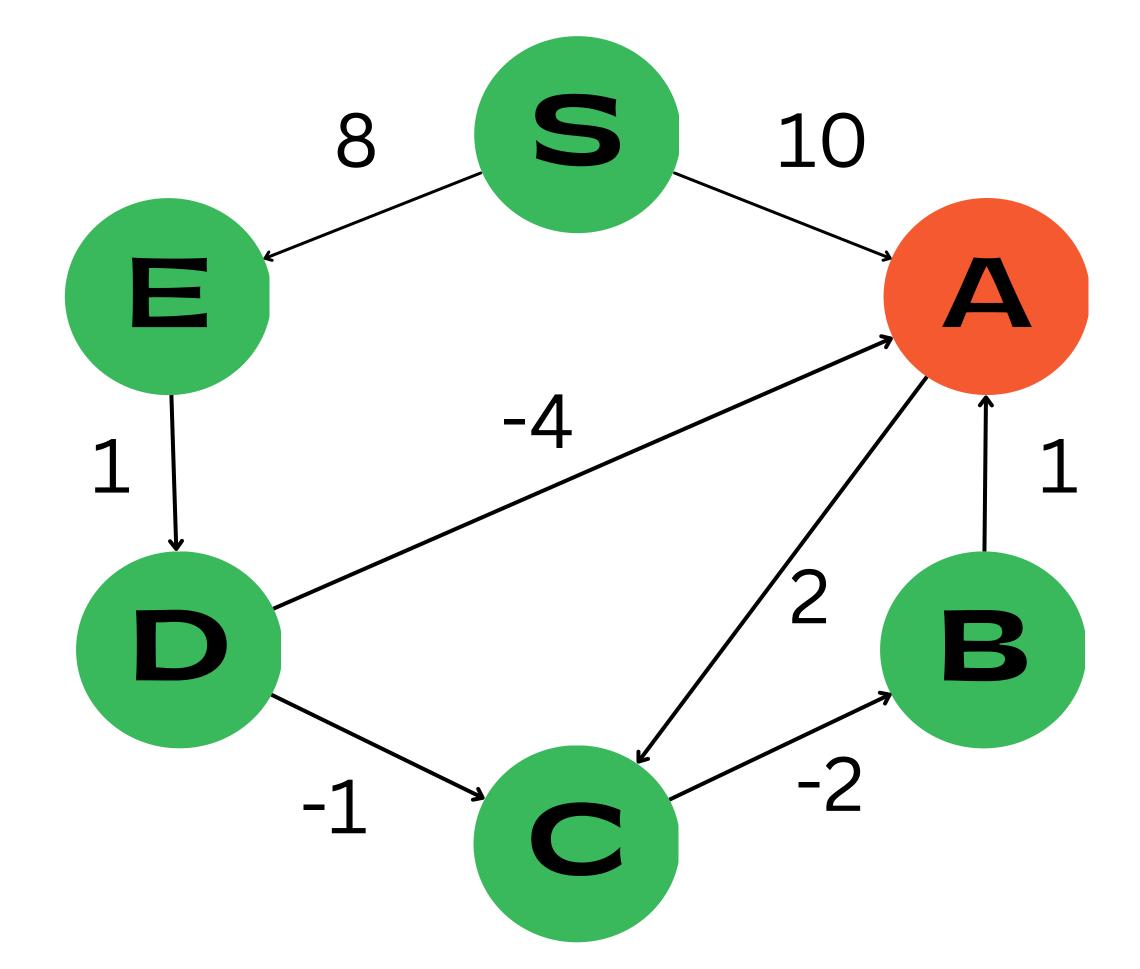
S	0
A	10
В	10
C	12
D	9
E	8



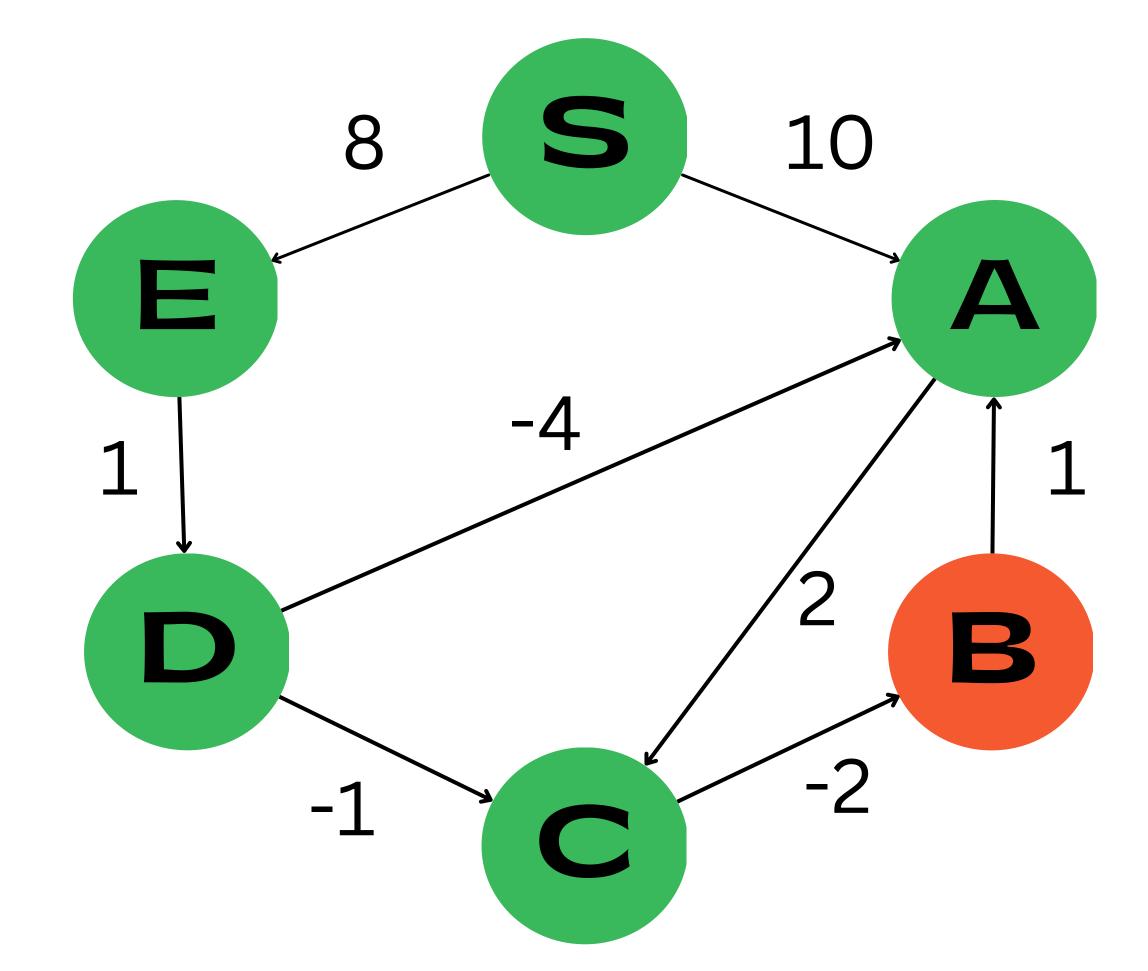
S	0
A	10
В	10
C	12
D	9
E	8



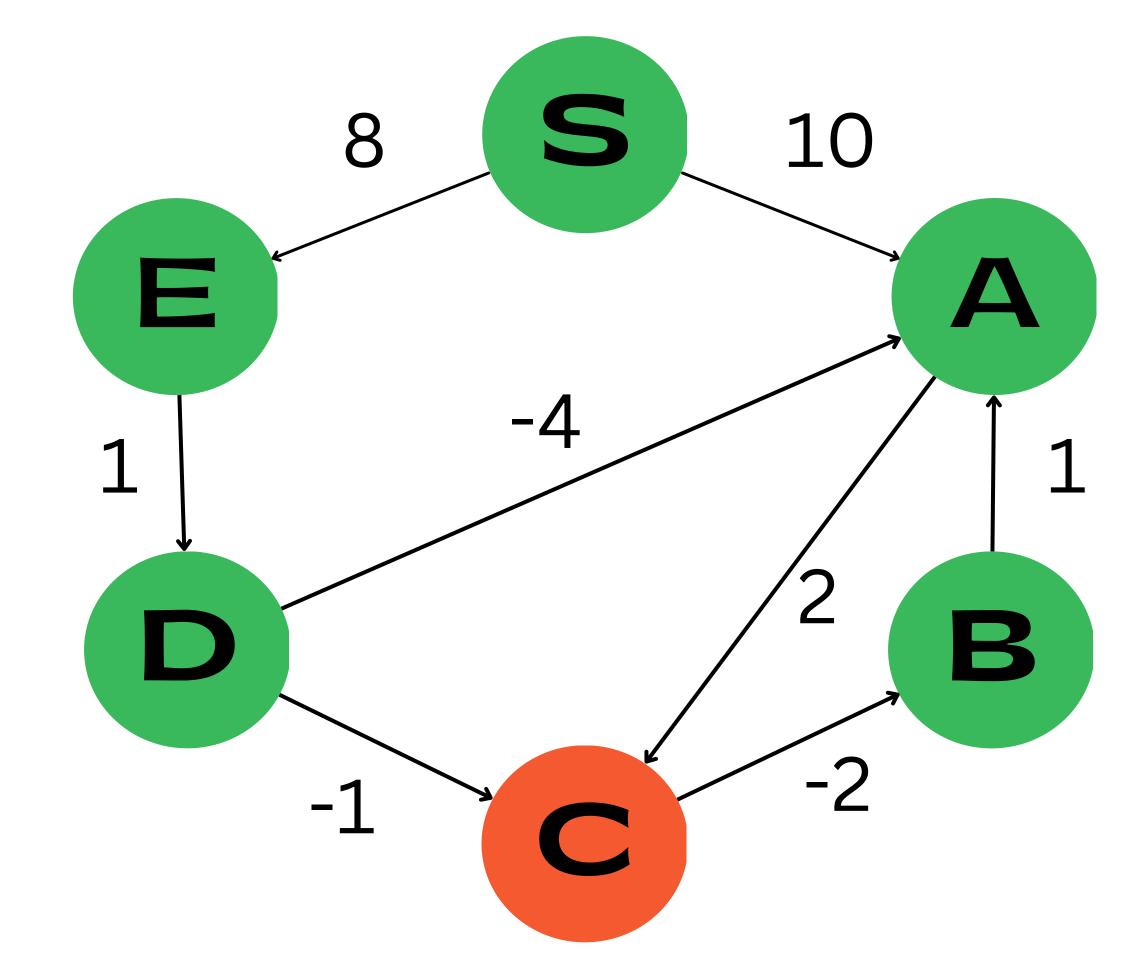
8	0
A	10
В	10
C	12
D	9
E	8



S	0
A	10
В	10
C	12
D	9
E	8

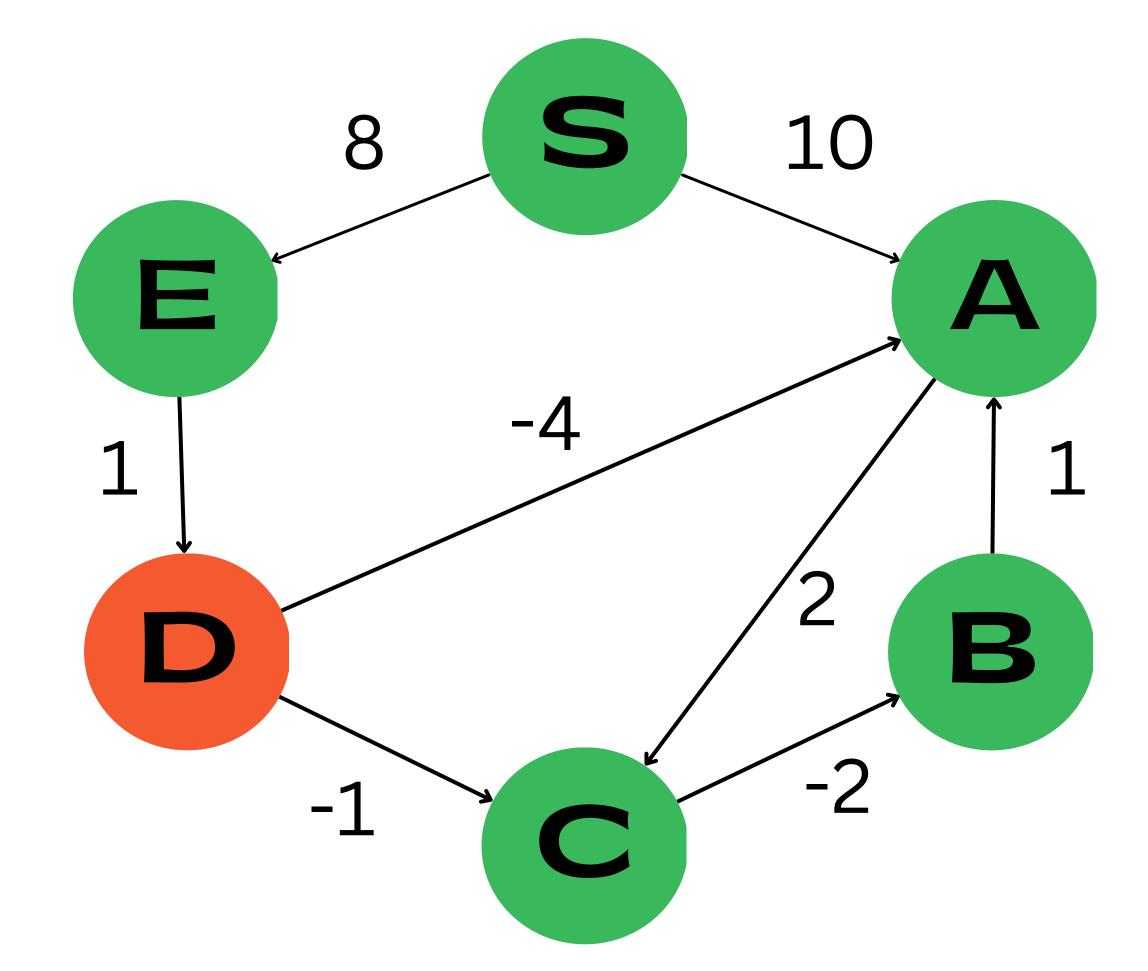


S	0
A	10
В	10
C	12
D	9
E	8

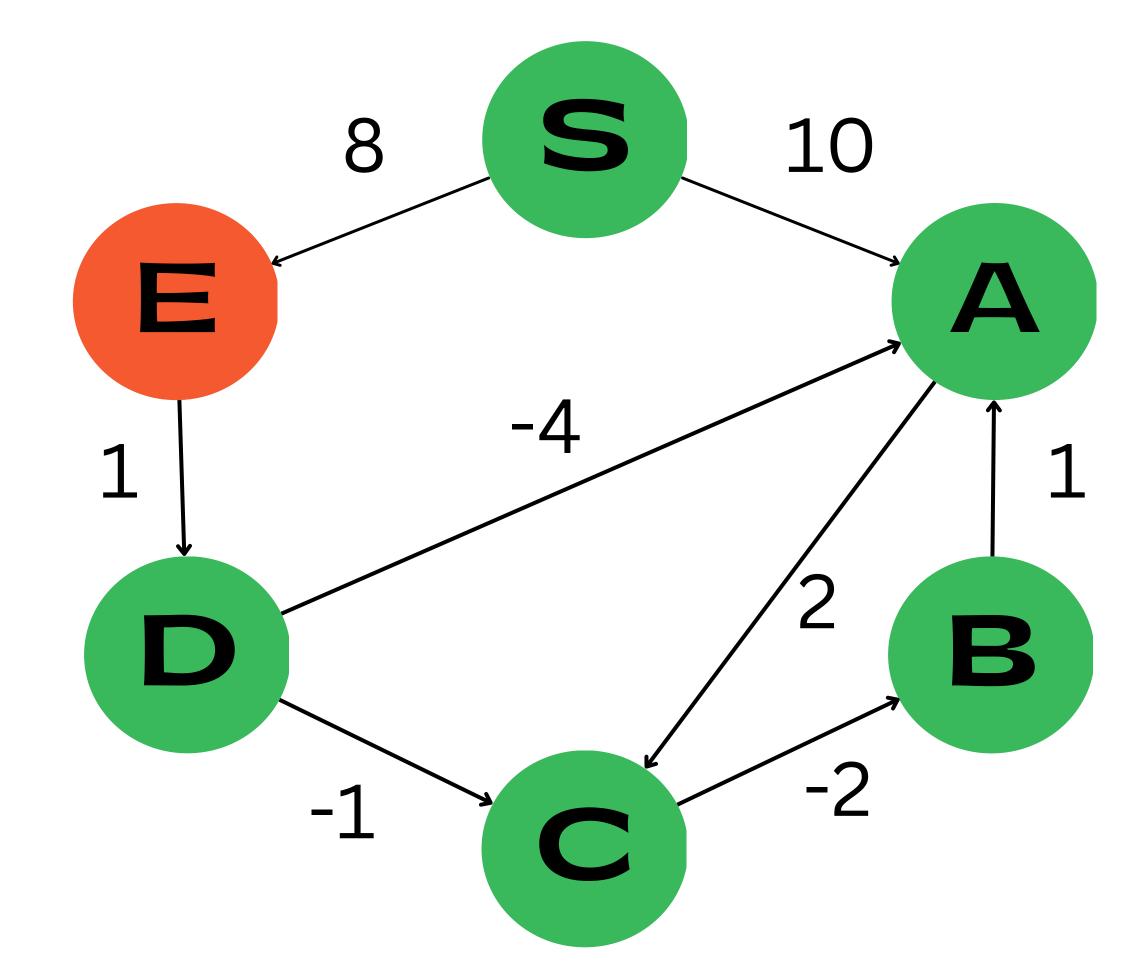


AGENDA PAGE

S	0
A	5
В	10
C	8
D	9
E	8

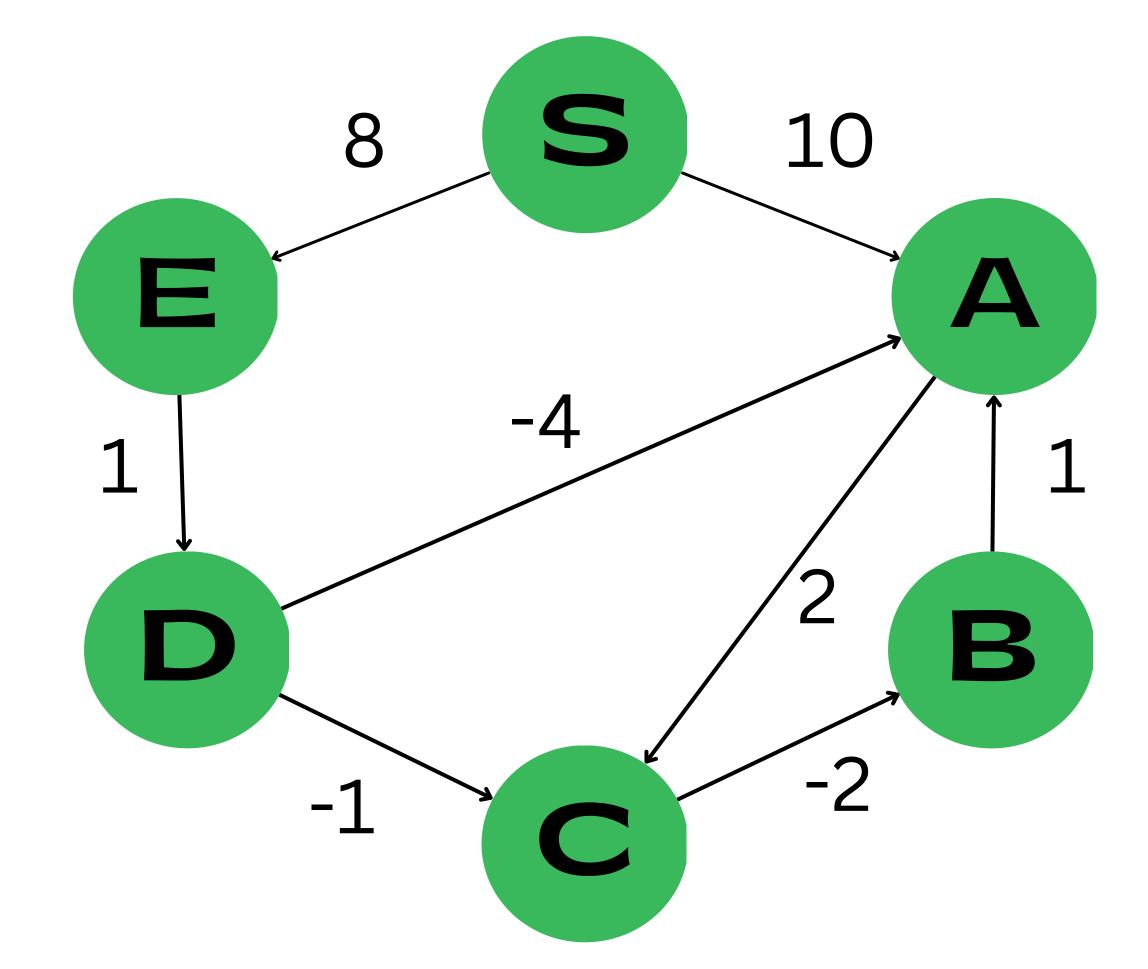


S	0
A	5
В	10
C	8
D	9
E	8



AGENDA PAGE

S	0
A	5
В	5
C	7
D	9
E	8



S	0
A	5
В	5
C	7
D	9
E	8

