Analysis Report

$maxwell_sgemm_128x64_raggedMn_tn$

Duration	15.392 μs
Grid Size	[7,2,1]
Block Size	[128,1,1]
Registers/Thread	120
Shared Memory/Block	12.5 KiB
Shared Memory Requested	96 KiB
Shared Memory Executed	96 KiB
Shared Memory Bank Size	4 B

[0] GeForce GTX 1080

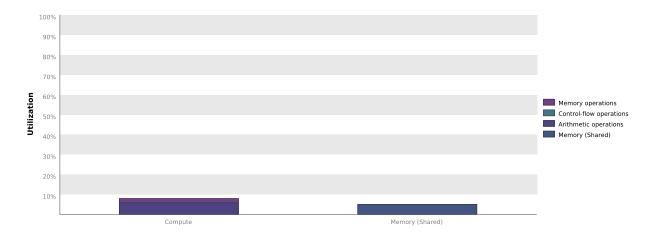
GPU UUID GPU-edd19385-a5f1-ce46-e1d9-61408827cccc								
	GPU-edd19385-a5f1-ce46-e1d9-61408827cccc							
Compute Capability	6.1							
Max. Threads per Block	1024							
Max. Threads per Multiprocessor	2048							
Max. Shared Memory per Block	48 KiB							
Max. Shared Memory per Multiprocessor	96 KiB							
Max. Registers per Block	65536							
Max. Registers per Multiprocessor	65536							
Max. Grid Dimensions	[2147483647, 65535, 65535]							
Max. Block Dimensions	[1024, 1024, 64]							
Max. Warps per Multiprocessor	64							
Max. Blocks per Multiprocessor	32							
Half Precision FLOP/s	72.9 GigaFLOP/s							
Single Precision FLOP/s	9.331 TeraFLOP/s							
Double Precision FLOP/s	291.6 GigaFLOP/s							
Number of Multiprocessors	20							
Multiprocessor Clock Rate	1.823 GHz							
Concurrent Kernel	true							
Max IPC	6							
Threads per Warp	32							
Global Memory Bandwidth	320.32 GB/s							
Global Memory Size	7.923 GiB							
Constant Memory Size	64 KiB							
L2 Cache Size	2 MiB							
Memcpy Engines	2							
PCIe Generation	3							
PCIe Link Rate	8 Gbit/s							
PCIe Link Width	16							

1. Compute, Bandwidth, or Latency Bound

The first step in analyzing an individual kernel is to determine if the performance of the kernel is bounded by computation, memory bandwidth, or instruction/memory latency. The results below indicate that the performance of kernel "maxwell_sgemm_128x64_ragged..." is most likely limited by instruction and memory latency. You should first examine the information in the "Instruction And Memory Latency" section to determine how it is limiting performance.

1.1. Kernel Performance Is Bound By Instruction And Memory Latency

This kernel exhibits low compute throughput and memory bandwidth utilization relative to the peak performance of "GeForce GTX 1080". These utilization levels indicate that the performance of the kernel is most likely limited by the latency of arithmetic or memory operations. Achieved compute throughput and/or memory bandwidth below 60% of peak typically indicates latency issues.



2. Instruction and Memory Latency

Instruction and memory latency limit the performance of a kernel when the GPU does not have enough work to keep busy. The results below indicate that the GPU does not have enough work because the kernel does not execute enough blocks.

2.1. Grid Size Too Small To Hide Compute And Memory Latency

The kernel does not execute enough blocks to hide memory and operation latency. Typically the kernel grid size must be large enough to fill the GPU with multiple "waves" of blocks. Based on theoretical occupancy, device "GeForce GTX 1080" can simultaneously execute 4 blocks on each of the 20 SMs, so the kernel may need to execute a multiple of 80 blocks to hide the compute and memory latency. If the kernel is executing concurrently with other kernels then fewer blocks will be required because the kernel is sharing the SMs with those kernels.

Optimization: Increase the number of blocks executed by the kernel.

2.2. GPU Utilization Is Limited By Register Usage

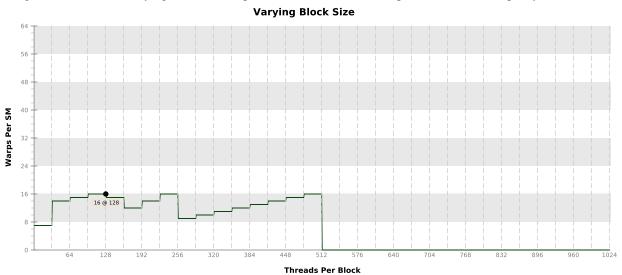
The kernel uses 120 registers for each thread (15360 registers for each block). This register usage is likely preventing the kernel from fully utilizing the GPU. Device "GeForce GTX 1080" provides up to 65536 registers for each block. Because the kernel uses 15360 registers for each block each SM is limited to simultaneously executing 4 blocks (16 warps). Chart "Varying Register Count" below shows how changing register usage will change the number of blocks that can execute on each SM.

Optimization: Use the -maxrregcount flag or the __launch_bounds__ qualifier to decrease the number of registers used by each thread. This will increase the number of blocks that can execute on each SM. On devices with Compute Capability 5.2 turning global cache off can increase the occupancy limited by register usage.

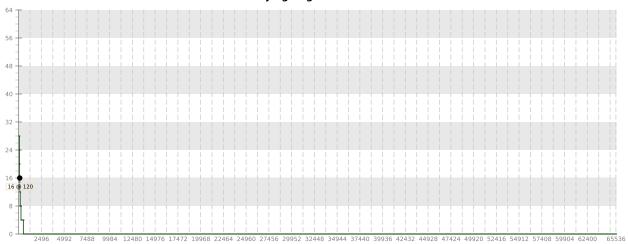
Variable	Achieved	Theoretical	Device Limit	Grid Siz	e: [7	7,2,1]	(14 k	olocks)	Block	Size:	[128	1,1](1
Occupancy Per SM												
Active Blocks		4	32	0	4	8	12	16	20	24	28	32
Active Warps	4	16	64	0	9	18	2	7 3	6	45	54	664
Active Threads		512	2048	0		512		1024		1536	5	2048
Occupancy	6.2%	25%	100%	0%		25%		50%)	759	%	100%
Warps												
Threads/Block		128	1024	0		256		512		768		1024
Warps/Block		4	32	0	4	8	12	16	20	24	28	32
Block Limit		16	32	0	4	8	12	16	20	24	28	32
Registers												
Registers/Thread		120	65536	0		16384		32768	3	4915	2	65536
Registers/Block		15360	65536	0		16k		32k		48k	<	64k
Block Limit		4	32	0	4	8	12	16	20	24	28	32
Shared Memory												
Shared Memory/Block		12800	98304	0			32k		.6	i4k		96k
Block Limit		7	32	0	4	8	12	16	20	24	28	32

2.3. Occupancy Charts

The following charts show how varying different components of the kernel will impact theoretical occupancy.

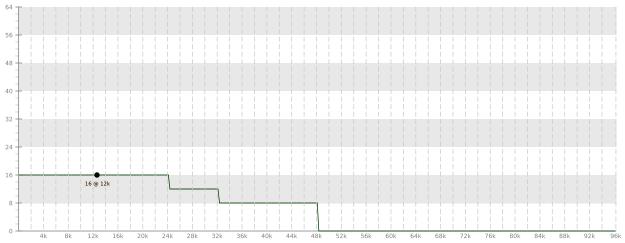


Varying Register Count



Registers Per Thread

Varying Shared Memory Usage



Shared Memory Per Block (bytes)

3. Compute Resources

GPU compute resources limit the performance of a kernel when those resources are insufficient or poorly utilized.

3.1. Function Unit Utilization

Different types of instructions are executed on different function units within each SM. Performance can be limited if a function unit is over-used by the instructions executed by the kernel. The following results show that the kernel's performance is not limited by overuse of any function unit.

Load/Store - Load and store instructions for shared and constant memory.

Texture - Load and store instructions for local, global, and texture memory.

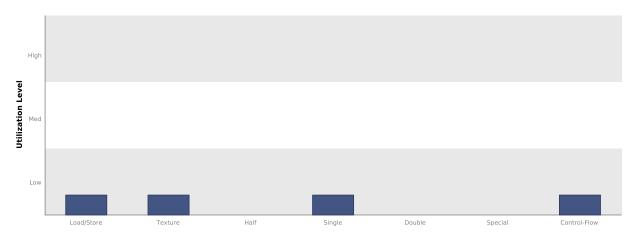
Half - Half-precision floating-point arithmetic instructions.

Single - Single-precision integer and floating-point arithmetic instructions.

Double - Double-precision floating-point arithmetic instructions.

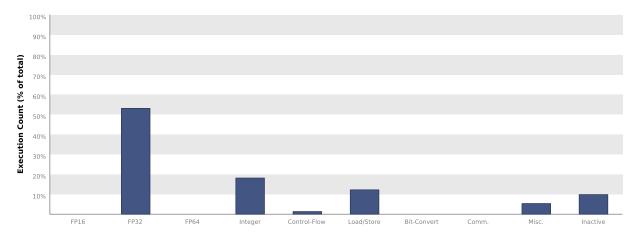
Special - Special arithmetic instructions such as sin, cos, popc, etc.

Control-Flow - Direct and indirect branches, jumps, and calls.



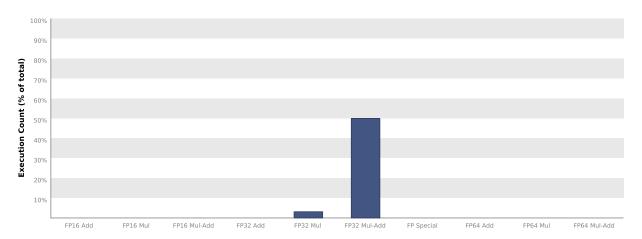
3.2. Instruction Execution Counts

The following chart shows the mix of instructions executed by the kernel. The instructions are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing instructions in that class. The "Inactive" result shows the thread executions that did not execute any instruction because the thread was predicated or inactive due to divergence.



3.3. Floating-Point Operation Counts

The following chart shows the mix of floating-point operations executed by the kernel. The operations are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing operations in that class. The results do not sum to 100% because non-floating-point operations executed by the kernel are not shown in this chart.



4. Memory Bandwidth

Memory bandwidth limits the performance of a kernel when one or more memories in the GPU cannot provide data at the rate requested by the kernel.

4.1. Memory Bandwidth And Utilization

The following table shows the memory bandwidth used by this kernel for the various types of memory on the device. The table also shows the utilization of each memory type relative to the maximum throughput supported by the memory.

