### Axiomatic Potentialism

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### Overview

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3 Height and Width Potentialism

#### **Potentialism**

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is the idea that a mathematical object (e.g. a set) is the sort of thing that may *merely possibly* exist.

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- Still deeper roots in mathematics in general.

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- Here we will be focussed on axiomatic potentialism, and on relations between potentialist axiom systems and their first order counterparts.

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- Axiomatize this conception and relate it to standard set theory.

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Standard modal model theory validates the rule

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Hence the need for free logic.

$$t: \mathcal{L}_0 \times V \to \mathcal{L}_{\in}, (\varphi, T) \mapsto \psi(T)$$

In fact ZFC interprets L.

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#### Theorem

 $\mathsf{L} \vdash \varphi \text{ implies } \mathit{ZFC} \vdash \mathit{t}(\varphi)(\emptyset)$ 

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### Linnebo Interpretation Theorem

 $L \vdash ZFC^{\diamond}$ .

Proof: use mirroring.

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## Mini-conclusion

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## Equivalence

We have an exact proof-theoretic equivalence,  $L \equiv ZFC$ .

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**M** interprets ZFC under the translation  $\varphi \mapsto \varphi^{\diamond \uparrow}$ .

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- M interprets ZFC<sup>-</sup> under the translation  $\varphi \mapsto \varphi^{\diamond}$ .
- M proves  $\neg Pow^{\diamond}$ .

- M interprets ZFC under the translation  $\varphi \mapsto \varphi^{\diamond_{\uparrow}}$ .
- M interprets ZFC<sup>-</sup> under the translation  $\varphi \mapsto \varphi^{\diamond}$ .
- M proves ¬Pow<sup>5</sup>.
- M proves  $V = HC^{\diamond}$  and hence  $SOA^{\diamond}$ .
- M proves it is possible for the continuum to exist and have a cardinality at least as great as any N number whose existence is provable in ZFC.

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But then the rigidity/extensionality imply w = z after all, so we have a contradiction.

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$$\Box \forall z \Box \forall Z \exists X \forall y [Xy \leftrightarrow \varphi(y, z, Z)]$$

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Amounts to restricting ourselves to parameters that exist at the world of evaluation.



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- Our translation will be doubly parameterized, once by a real and once by a transitive set.
- Our interpretation for  $\Diamond_{\uparrow}$  will involve holding r fixed and climbing transitive sets in L[r];
- while our interpretation for  $\Diamond_{\leftarrow}$  will involve allowing new reals to be added but not extending the height of the transitive set parameter.

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- $t(\Box_{\leftarrow}\varphi)(T,r) := \forall s[r \in L[s] \to \forall S[S \leq T \land S \in L[s] \to t(\varphi)(S,s)]]$

Let  $M \models SOA + \Pi_1^1 PSP$ .

$$t: \mathcal{L}_1 \times M \times \mathbb{R}^M \to \mathcal{L}_{\in}, (\varphi, T, r) \mapsto \psi(T, r)$$

- $\blacksquare$  assign plural variables odd numbered variables t(X).
- Membership claims on sets, propositional connectives = id
- $t(Xx)(T,r) := x \in t(X)$
- $t(\forall x\varphi)(T,r) := \forall x \in T[t(\varphi)(T,r)]$
- $t(\forall X\varphi)(T,r) := \forall x \subseteq T[x \in L[r] \to t(\varphi)(T,r)]$
- $t(\Box_{\uparrow}\varphi)(T,r) := \forall S \supseteq T[Tran(S) \land S \in L[r] \rightarrow t(\varphi)(S,r)]$
- $t(\Box_{\leftarrow}\varphi)(T,r) := \forall s[r \in L[s] \to \forall S[S \leq T \land S \in L[s] \to t(\varphi)(S,s)]]$

#### Theorem

 $M \vdash \varphi \text{ implies } T \vdash t(\varphi)(\emptyset, 0)$ 

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Given any possible real r, one can show  $\lozenge \exists x[x = \mathbb{R}^{L[r]}]$  using the  $\lozenge_{\uparrow}$  translation of ZFC and some absoluteness lemmas.

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#### References



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# Thanks!