

Axiomatic Potentialism

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Overview

- 1 Background
- 2 Warm Up: Height Potentialism
- 3 Height and Width Potentialism
- 4 Concluding Remarks

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- Still deeper roots in mathematics in general.

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- E.g. Hamkins and Linnebo showed in MT potentialism with the structures initial segments of V that $S5$ at a world V_κ is equivalent to Σ_3 correctness of κ .
- Here we will be focused on axiomatic potentialism, and on relations between potentialist axiom systems and their first order counterparts.

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Instantiation requires assumption of existence.

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Theorem

$L \vdash \varphi$ implies $ZFC \vdash t(\varphi)(\emptyset)$

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Proof: use mirroring.

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Bi-interpretation

In fact, $ZFC \vdash t(\phi^\diamond)(\emptyset) \leftrightarrow \phi$, and conversely.

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- M interprets ZFC^- under the translation $\varphi \mapsto \varphi^\diamond$.
- M proves $\neg Pow^\diamond$.
- M proves $V = HC^\diamond$ and hence SOA^\diamond .
- M proves it is possible for the continuum to exist and have a cardinality at least as great as any \aleph number whose existence is provable in ZFC.

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But then the rigidity/extensionality imply $w = z$ after all, so we have a contradiction.

The argument just sketched uses comprehension with arbitrary parameters:

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And in the crucial application, it applies when we have no *a priori* guarantee z even exists (indeed this is what we are trying to establish.)

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Amounts to restricting ourselves to parameters that exist at the world of evaluation.

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- Our interpretation for \diamond_{\uparrow} will involve holding r fixed and climbing transitive sets in $L[r]$;
- while our interpretation for \diamond_{\leftarrow} will involve allowing new reals to be added but not extending the height of the transitive set parameter.

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Theorem

$M \vdash \varphi$ implies $T \vdash t(\varphi)(\emptyset, 0)$

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What happens with stronger regularity properties? How does PD/AD affect things?

Thanks!