Axiomatic Potentialism

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Overview

- Background
- 2 Warm Up: Height Potentialism
- Height and Width Potentialism
- Concluding Remarks

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Potentialism

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is the idea that a mathematical object (e.g. a set) is the sort of thing that may *merely possibly* exist.

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- Still deeper roots in mathematics in general.

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- E.g. Hamkins and Linnebo showed in MT potentialism with the structures initial segments of V that S5 at a world V_{κ} is equivalent to Σ_3 correctness of κ .
- Here we will be focused on axiomatic potentialism, and on relations between potentialist axiom systems and their first order counterparts.

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■ Extensionality, ∈-rigidity, foundation

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Instantiation requires assumption of existence.



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Theorem

 $\mathsf{L} \vdash \varphi \mathsf{ implies } \mathit{ZFC} \vdash \mathit{Tran}(X) \to t(\varphi)(X)$



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Linnebo Interpretation Theorem

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Proof: use mirroring.

Axioms for the theory L

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We have an exact proof-theoretic equivalence, $L \equiv ZFC$.

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Bi-interpretation

In fact, ZFC $\vdash Tran(X) \rightarrow t(\phi^{\diamond})(X) \leftrightarrow \phi$, and conversely.

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- M proves it is possible for the continuum to exist and have a cardinality at least as great as any ℵ number whose existence is provable in ZFC.

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But then the rigidity/extensionality imply w = z after all, so we have a contradiction.



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Amounts to restricting ourselves to parameters that exist at the world of evaluation.



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- Our translation will be doubly parameterized, once by a real and once by a transitive set.
- Our interpretation for \Diamond_{\uparrow} will involve holding r fixed and climbing transitive sets in L[r];
- while our interpretation for \Diamond_{\leftarrow} will involve allowing new reals to be added but not extending the height of the transitive set parameter.

$$t: \mathcal{L}_1 \times M \times \mathbb{R}^M \to \mathcal{L}_{\in}, (\varphi, T, r) \mapsto \psi(T, r)$$

Let
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- $T(\square_{\leftarrow}\varphi)(T,r) := \forall s[r \in L[s] \to \forall S[S \leq T \land S \in L[s] \to t(\varphi)(S,s)]]$

Theorem

 $\mathsf{M} \vdash \varphi \text{ implies } \mathsf{T} \vdash \mathbb{R}(r) \land \mathit{Tran}(X) \land X \in \mathit{L}[r] \to \mathit{t}(\varphi)(X,r)$

Theorem

 $\mathsf{M} \vdash \mathsf{\Pi}_1^1 \mathit{PSP}^{\diamond}$

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What happens with stronger regularity properties? How does ${\sf PD}/{\sf AD}$ affect things?

Thanks!