### Axiomatic Potentialism

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### Overview

- Background
- 2 Warm Up: Height Potentialism
- Height and Width Potentialism
- Concluding Remarks

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#### **Potentialism**

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is the idea that a mathematical object (e.g. a set) is the sort of thing that may *merely possibly* exist.

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- Still deeper roots in mathematics in general.

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- Here we will be focused on axiomatic potentialism, and on relations between potentialist axiom systems and their first order counterparts.

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- Axiomatize this conception and work out its non-modal counterpart.

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■ Extensionality, ∈-rigidity, foundation

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Instantiation requires assumption of existence.



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#### Theorem

 $\mathsf{L} \vdash \varphi \mathsf{ implies } \mathit{ZFC} \vdash \forall X[\mathit{Tran}(X) \rightarrow t(\varphi)(X)]$ 



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### Linnebo Interpretation Theorem

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Proof: use mirroring.

# Axioms for the theory L

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This is something like what Button calls near-synonymy (near-near-synonmy, I guess.)

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- M proves the universal possibility of forcing.
- e.g., it proves it is possible for the continuum to exist and have a cardinality at least as great as any N number whose existence is provable in ZFC.

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But then the rigidity/extensionality imply w = z after all, so we have a contradiction.



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Amounts to restricting ourselves to parameters that exist at the world of evaluation.



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- Our interpretation for  $\Diamond_{\uparrow}$  will involve holding r fixed and climbing transitive sets in L[r];
- while our interpretation for  $\Diamond_{\leftarrow}$  will involve allowing new reals to be added but not extending the height of the transitive set parameter.

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- Membership = id, commutes with propositional connectives
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#### **Theorem**

 $M \vdash \varphi \text{ implies } T \vdash \forall r, T \in L[r][t(\varphi)(X, r)]$ 

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What happens with stronger regularity properties? How does  ${\sf PD}/{\sf AD}$  affect things?

# Thanks!