### Axiomatic Potentialism

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### Overview

- Background
- 2 Warm Up: Height Potentialism
- Height and Width Potentialism
- Concluding Remarks

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- Still deeper roots in mathematics in general.

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- Here we will be focussed on axiomatic potentialism, and on relations between potentialist axiom systems and their first order counterparts.

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Logical Axioms

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Standard modal model theory validates the rule

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Instantiation requires assumption of existence.



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In fact ZFC interprets L.

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#### **Theorem**

 $\mathsf{L} \vdash \varphi \text{ implies } \mathit{ZFC} \vdash \mathit{t}(\varphi)(\emptyset)$ 

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Proof: use mirroring.

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### Bi-interpretation

In fact, ZFC  $\vdash t(\phi^{\diamond})(\emptyset) \leftrightarrow \phi$ , and conversely.

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- M proves it is possible for the continuum to exist and have a cardinality at least as great as any ℵ number whose existence is provable in ZFC.

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But then the rigidity/extensionality imply w=z after all, so we have a contradiction.

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Amounts to restricting ourselves to parameters that exist at the world of evaluation.



#### Intuitive idea

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(From now on, I will ignore the difference between SOA and  $ZFC^- + V = HC$ . Replacement is formulated as collection.)

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- Our translation will be doubly parameterized, once by a real and once by a transitive set.
- Our interpretation for  $\Diamond_{\uparrow}$  will involve holding r fixed and climbing transitive sets in L[r];
- while our interpretation for  $\Diamond_{\leftarrow}$  will involve allowing new reals to be added but not extending the height of the transitive set parameter.

$$t: \mathcal{L}_1 \times M \times \mathbb{R}^M \to \mathcal{L}_{\in}, (\varphi, T, r) \mapsto \psi(T, r)$$

Let  $M \models SOA + \Pi_1^1 PSP$ .

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### **Theorem**

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The 'parsimony and order' perspective on the universe?

# Thanks!