







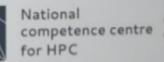






National competence centre for HPC







GPU Programming. When, Why and How?

2024

ENCCS Training





Non-portable kernel-based models. Fundamentals

CUDA and HIP



- CUDA and HIP are solely focused on GPUs.
- They provide all the necessary tools and advance features to write higly optimized applications for running on GPUs:
 - extensive libraries
 - low level APIs
 - compiler toolchains that optimize code execution on NVIDIA GPUs (in the case of CUDA) and both NVIDIA and AMD GPUs (in the case of HIP)
 - debugging and performance analysis tools

Hello World



```
#include <hip/hip_runtime.h>
#include <stdio.h>

int main(void){
   int count, device;
   hipGetDeviceCount(&count);
   hipGetDevice(&device);

   printf("Hello! I'm GPU %d out of %d GPUs in total.\n", device, count);
   return 0;
}
```

Vector Addition



```
__global__ void vector_add(float *A, float *B, float *C, int n){
   int tid = threadIdx.x + blockIdx.x * blockDim.x;
   if(tid<n){
       C[tid] = A[tid]+B[tid];
   }
}</pre>
```

```
#include <hip/hip_runtime.h>
#include <stdio.h>
#include <stlib.h>
#include <math.h>
int main(){
    const int N = 10000;
    float *Ah, *Bh, *Ch, *Cref;
    float *Ad, *Bd, *Cd;

    // Allocate the arrays on CPU
    Ah =(float*)malloc(n * sizeof(float));Bh =...;Ch =...; Cref =...;

    // Initialise data and calculate reference values on CPU
    for (i=0; i < n; i++) {
        Ah[i] = sin(i) * 2.3;
        Bh[i] = cos(i) * 1.1;
        Cref[i] = Ah[i] + Bh[i];}</pre>
```

```
// Allocate the arrays on GPU
hipMalloc((void**)&Ad, N * sizeof(float));
hipMalloc((void**)&Bd, N * sizeof(float));
hipMalloc((void**)&Cd, N * sizeof(float));
// Transfer the data from CPU to GPU
hipMemcpy(Ad, Ah, sizeof(float) * n, hipMemcpyHostToDevice);
hipMemcpy(Bd, Bh, sizeof(float) * n, hipMemcpyHostToDevice);
// define grid dimensions + launch the device kernel
im3 blocks, threads;
threads=dim3(256,1,1);
blocks=dim3((N+256-1)/256,1,1);
//Launch Kernel
//hipLaunchKernelGGL(vector_add, blocks, threads, 0, 0, Ad, Bd, Cd,
vector_add<<< blocks, threads,0,0>>(Ad, Bd, Cd, N);
// copy results back to CPU
hipMemcpy(Ch, Cd, sizeof(float) * N, hipMemcpyDeviceToHost);
// Free the GPU arrays
hipFree(Ad); hipFree(Bd); hipFree(Cd);
printf("Result: %f %f %f \n", Ch[3], Ch[n-2], Ch[n-1]);
```





```
int main(){
   const int N = 10000;
   float *Ah, *Bh, *Ch, *Cref;

   // Allocate the arrays using Unified Memory
   hipMallocManaged((void **)&Ah, N * sizeof(float));
   hipMallocManaged((void **)&Bh, N * sizeof(float));
   hipMallocManaged((void **)&Ch, N * sizeof(float));
   hipMallocManaged((void **)&Cref, N * sizeof(float));

   // Initialise data and calculate reference values on CPU
   for (i=0; i < n; i++) {
        Ah[i] = sin(i) * 2.3;
        Bh[i] = cos(i) * 1.1;
        Cref[i] = Ah[i] + Bh[i];}
   // All data at this point is on CPU</pre>
```

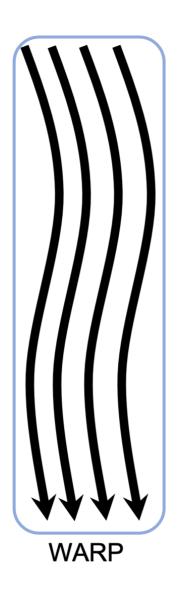
```
// define grid dimensions + launch the device kernel
dim3 blocks, threads;
threads=dim3(256,1,1);
blocks=dim3((N+256-1)/256,1,1);
//Launch Kernel
//hipLaunchKernelGGL(vector_add, blocks, threads, 0, 0, Ad, Bd, Cd,
vector_add<<< blocks, threads,0,0>>(Ad, Bd, Cd, N);
hipDeviceSynchronize(); // Wait for the kernel to complete
//Access the data on the CPU
printf("Result: %f %f %f \n", Ch[3], Ch[n-2], Ch[n-1]);
// Free the Unified Memory arrays
hipFree(Ah);
hipFree(Bh);
hipFree(Ch);
hipFree(Cref);
```

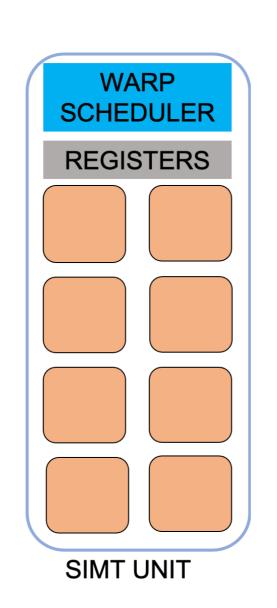


Non-portable kernel-based models. Memory optimizations

Coalesced Access







- The CUDA threads are physically locked into warps, currently of size 64 for AMD and 32 for Nvidia.
- All threads in the warp have to execute the same instruction.
- The memory accesses are done per warp.

Coalesced vs Non-coalesced Memory Access



```
__global__ void vector_add(float *A, float *B, float *C, int n, int stride, int shift){
  int tid = threadIdx.x + blockIdx.x * blockDim.x;
  if(tid<n){
      C[tid] = A[tid]+B[tid]; //Coalesced

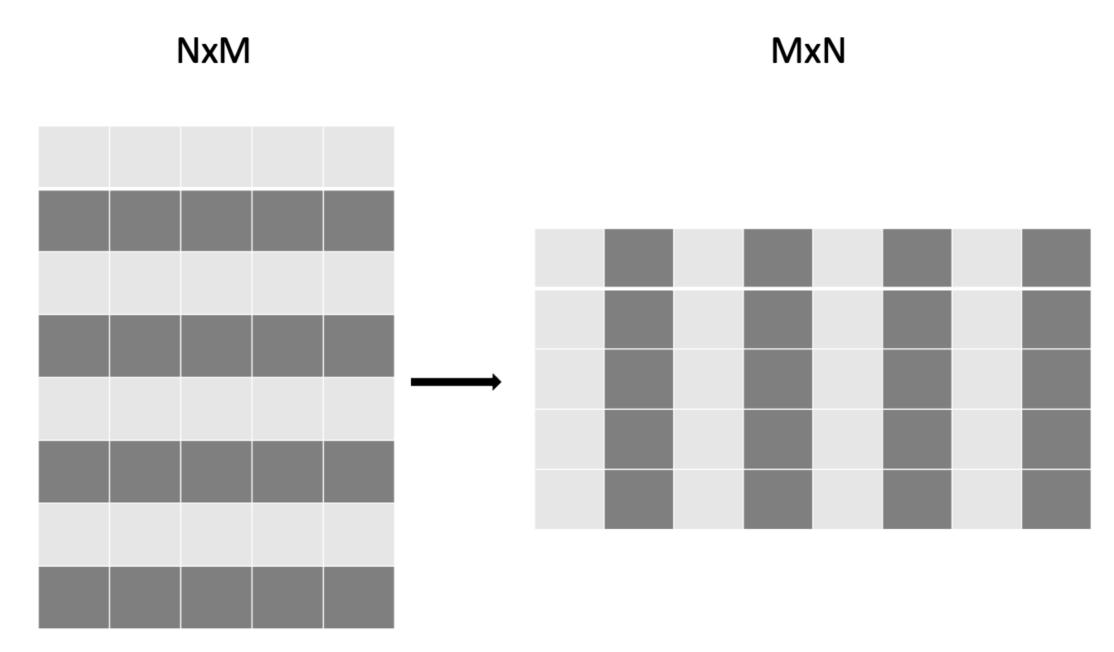
      C[tid*stride] = A[tid*stride]+B[tid*stride]; //Strided Non-Coalesced

      C[tid+shift] = A[tid+shift]+B[tid+shift]; //Shifted Non-Coalesced
  }
}</pre>
```

Local shared memory can be used to improve the memory accesses.

Optimizing matrix operations. B(i,j)=A(j,i)





Copy operation as base



```
__global__ void copy_kernel(float *in, float *out, int width, int height) {
  int x_index = blockIdx.x * tile_dim + threadIdx.x;
  int y_index = blockIdx.y * tile_dim + threadIdx.y;

  int index = y_index * width + x_index;

  out[index] = in[index];
}
```

The effective bandwidth is 717 GB/s, out of the theoretical peak 900 GB/s.

Matrix transpose naive



```
__global__ void transpose_kernel(float *in, float *out, int width, int height) {
   int x_index = blockIdx.x * tile_dim + threadIdx.x;
   int y_index = blockIdx.y * tile_dim + threadIdx.y;

   int in_index = y_index * width + x_index;
   int out_index = x_index * height + y_index;

   out[out_index] = in[in_index];
}
```

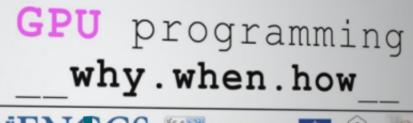
The effective bandwidth is 311 GB/s.

Matrix transpose with shared memory



```
_global__ void transpose_lds_kernel(float *in, float *out, int width,
                                   int height) {
__shared__ float tile[tile_dim][tile_dim];
int x_tile_index = blockIdx.x * tile_dim;
int y_tile_index = blockIdx.y * tile_dim;
int in_index =
    (y_tile_index + threadIdx.y) * width + (x_tile_index + threadIdx.x);
int out_index =
    (x_tile_index + threadIdx.y) * height + (y_tile_index + threadIdx.x);
tile[threadIdx.y][threadIdx.x] = in[in_index];
__syncthreads();
out[out_index] = tile[threadIdx.x][threadIdx.y];
```

The effective bandwidth is 674 GB/s.



Matrix transpose with shared memory without banks where because of the contract of the contrac conflicts

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```
_global__ void transpose_lds_kernel(float *in, float *out, int width,
                                   int height) {
__shared__ float tile[tile_dim][tile_dim+1];
int x_tile_index = blockIdx.x * tile_dim;
int y_tile_index = blockIdx.y * tile_dim;
int in index =
    (y_tile_index + threadIdx.y) * width + (x_tile_index + threadIdx.x);
int out index =
    (x_tile_index + threadIdx.y) * height + (y_tile_index + threadIdx.x);
tile[threadIdx.y][threadIdx.x] = in[in_index];
__syncthreads();
out[out_index] = tile[threadIdx.x][threadIdx.y];
```

On a NVIDIA V100 GPU this final code achieves an effective bandwidth 697 GB/s, out of the theoretical peak 900 GB/s.

Reductions



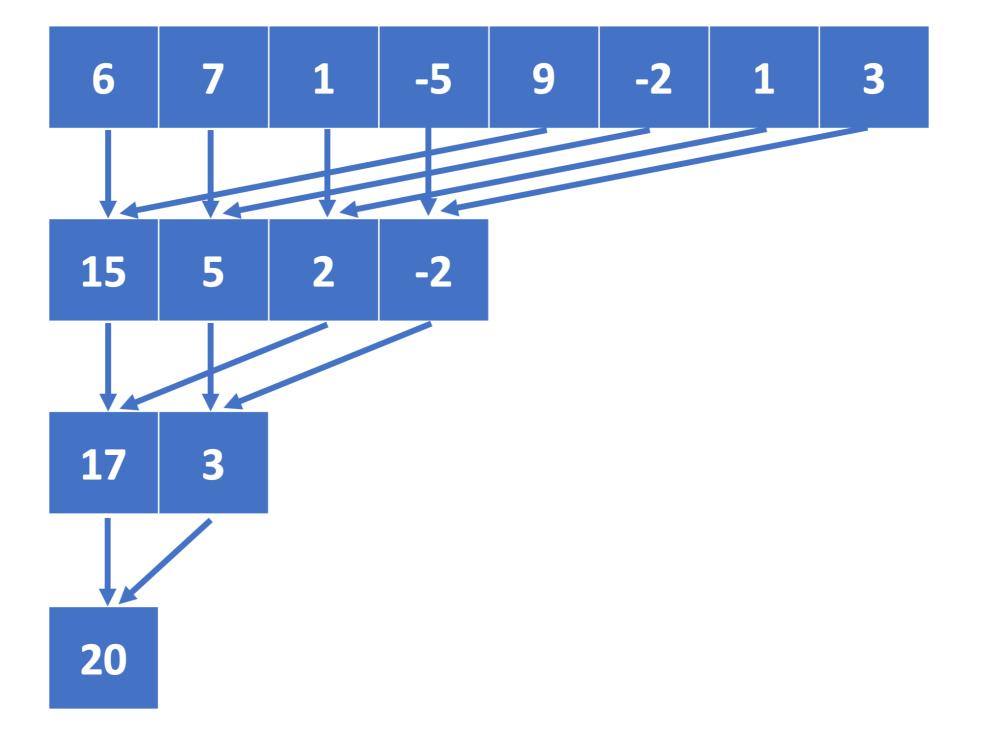
- Reductions refer to operations in which the elements of an array are aggregated in a single value through binary operations.
- Examples include summing, finding the maximum or minimum, or performing logical operations.
- Assumptions for being able to perform them in parallel:
 - the order of the iterations does not change the results. a+b=b+a
 - doing the operations on small subsets, and then on the results does not change the results. (a+b)+c=a+(b+c)

Reduction in parallel



- Divide the problem in subsets which can be processed in parallel.
- Each susbset is processed by a block of threads.
- Have an efficient reduction at block level:
 - keep as many as possible threads doing work.
 - o avoid global memory accesses using local data share.

Tree Reduction Inside a Block of Threads





CUDA/HIP Reduction



```
#define tpb 512 // size in this case has to be known at compile time
// this kernel has to be launched with at least N/2 threads
 _global__ void reduction_one(double x, double *sum, int N){
 int ibl=blockIdx.y+blockIdx.x*gridDim.y;
 int ind=threadIdx.x+blockDim.x*ibl;
  __shared__ double shtmp[2*tpb];
 shtmp[threadIdx.x]=0; // for sums we initiate with 0, for other operations should be different
 if(ind<N/2){
     shtmp[threadIdx.x]=x[ind];
 if(ind+N/2<N){
     shtmp[threadIdx.x+tpb]=x[ind+N/2];
  __syncthreads();
 for(int s=tpb; s>0; s>>=1){
   if(threadIdx.x<s){</pre>
       shtmp[threadIdx.x]+=shtmp[threadIdx.x+s];}
    __syncthreads();
 if(threadIdx.x==0){
   sum[ibl]=shtmp[0]; // each block saves its partial result to an array
    // atomicAdd(&sum[0], shene[0]); // alternatively could aggregate everything together at index 0. Only use when there not many partial sums lef
```

CUDA/HIP Streams



- Modern GPUs can overlap independent operations.
- CPU-GPU data transfers can be overllapped with lernel execution.
- **CUDA/HIP streams** are independent execution units, a sequence of operations that execute in issue-order on the GPU.
- The operations issue in different streams can be executed concurrently.
- Utilizing multiple streams, the GPU can avoid idle time, especially for problems with frequent CPU communication or multi-GPU setups.

GPU programming __why.when.how__

Overlapping Computations and Memory transferences M



Time

Serial Execution			
H2D Engine	0		
Kernel Engine		0	
H2D Engine			0

Concurrent Execution									
H2D Engine		1	2	3	4	5			
Kernel Engine			1	2	3	4	5		
H2D Engine				1	2	3	4	5	

Vector Addition with Streams



```
// Distribute kernel for 'n_streams' streams, and record each stream's timing
for (int i = 0; i < n_streams; ++i) {
    int offset = i * (N/stream_size);
    hipEventRecord(start_event[i], stream[i]); // stamp the moment when the kernel is submitted on stream i

    hipMemcpyAsync( &Ad[offset], &Ah[offset], N/n_streams*sizeof(float), hipMemcpyHostToDevice, stream[i]);
    hipMemcpyAsync( &Bd[offset], &Bh[offset], N/n_streams*sizeof(float), hipMemcpyHostToDevice, stream[i]);
    vector_add<<<gridsize / n_streams, blocksize, 0, stream[i]>>>(&Ad[offset], &Cd[offset], N/n_streams); //each call processes N/n_streams.
    hipMemcpyAsync( &Ch[offset], &Cd[offset], N/n_streams*sizeof(float), hipMemcpyDeviceToHost, stream[i]);
    hipEventRecord(stop_event[i], stream[i]); // stamp the moment when the kernel on stream i finished
}
...
```

Summary



- CUDA and HIP are the native programming models for NVIDIA and AMD GPUs.
- The programmer can take advantage of all GPU features.
- NVIDIA has a very extensive ecosystem. AMD is catching up.
- HIP is open source and can be used on both NVIDIA and AMD platforms.
- They are exclusive for NVIDIDA and AMD GPUs.
- Both CUDA and HIP require learning GPU programming concepts.
- Memory optimizitions are very important.