

INFORMATION THEORY

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TEACHER: Christian Schaffner, TAs: Yfke Dulek, Esteban Landerreche, Kyah Smaal

Team homework set 4

Unless otherwise stated, you should provide exact answers rather than rounded numbers (e.g., $\log 3$ instead of 1.585) for non-programming exercises.

Problem 1: Calculation of the typical set (7pt)

To clarify the notion of a typical set $A_\varepsilon^{(n)}$ and the smallest set of high probability $B_\delta^{(n)}$, we will calculate these sets for a simple example. Consider a sequence of i.i.d. binary random variables X_1, X_2, \dots, X_n , where the probability that $P_X(1) = 0.6$ and $P_X(0) = 0.4$.

- (a) **(3pt)** With $n = 25$ and $\varepsilon = 0.1$, which sequences fall in the typical set $A_\varepsilon^{(n)}$? What is the probability of the typical set (three decimals)? How many elements are there in the typical set? (This involves computation of a table of probabilities for sequences with k 1's, $0 \leq k \leq 25$, and finding those sequences that are in the typical set.)

Hint: Here is the table: <http://goo.gl/sQCPMO>

- (b) **(2pt)** How many elements are there in the smallest set that has probability 0.9? In other words, what is $|B_\delta^{(n)}|$ for $n = 25$ and $\delta = 0.1$?
- (c) **(2pt)** How many elements are there in the intersection $|A_\varepsilon^{(n)} \cap B_\delta^{(n)}|$ of the sets computed in parts (b) and (c)? What is the probability of this intersection (three decimals)?

Problem 2: Three random variables (3pt)

Let A, B, C be random variables such that

$$\begin{aligned} I(A; B) &= 0 \\ I(A; C|B) &= I(A; B|C) \\ H(A|BC) &= 0 \end{aligned}$$

Prove one of the following possible relations between $H(A)$ and $H(C)$: $=, \leq, \geq, <, \text{ or } >$. If you show one of the weaker inequalities (\leq or \geq), also exhibit an example that shows why the inequality is not strict.

Problem 3: Piece of cake (3pt)

A big cake is repeatedly sliced into two pieces. At every step, the *smaller* part is discarded (or eaten). On the other part the process is continued. At every step, the remaining piece is cut randomly into two pieces with the following proportions:

$$\begin{aligned} \left(\frac{2}{3}, \frac{1}{3}\right) &\text{ with probability } \frac{3}{4}, \\ \left(\frac{3}{5}, \frac{2}{5}\right) &\text{ with probability } \frac{1}{4}. \end{aligned}$$

For example, three consecutive cuts might result in a piece of cake of size $\frac{3}{5} \cdot \frac{2}{3} \cdot \frac{2}{3}$. Let T_n be the fraction of the cake left after n cuts. Clearly, this fraction T_n decreases exponentially with n , i.e. $T_n \approx c^n$ for some real constant $0 < c < 1$. Determine c for large values of n .

Hint: Let C_i be a random variable describing the fraction of the cake that is cut (and kept) at the i th cut, i.e. $C_i = \frac{2}{3}$ with probability $\frac{3}{4}$, or $\frac{3}{5}$ otherwise. Then use the weak law of large numbers.

Problem 4: Programming (7pt)

In one of the quizzes for this week, you implemented the encryption and decryption procedure for the Vigenère cipher. (Note: we treat "A" as 1 instead of 0. That is, K+A=L)

- (a) **(1pt)** Consider this cipher for a fixed message length m and a fixed key length k . Suppose the possible messages are uniformly distributed over all monospace messages (excluding spaces; there are 26 possible letters) of length m . What is $I(M; C)$? When is the Vigenère cipher perfectly secure (i.e., $I(M; C) = 0$)?

In reality, the possible messages are not uniformly distributed. We can use this fact to break a ciphertext that was encoded with the Vigenère cipher.

Recall the definition of collision probability from the first homework:

$$\text{Coll}(P) := \sum_{x \in \mathcal{X}} P(x)^2.$$

Let $\text{Coll}(P_M)$, $\text{Coll}(P_K)$, and $\text{Coll}(P_C)$ denote the collision probabilities of (sampling a single letter from) the plaintext, key, and ciphertext, respectively. Suppose you know that the message is in English. In the first homework, you estimated that $\text{Coll}(P_M) \approx 0.0655$.

(b) (2pt) Show that

$$\text{Coll}(P_C) \approx \frac{1}{k} \text{Coll}(P_M) + \frac{k-1}{k} \text{Coll}(P_K)$$

by assuming that:

1. the message is very long compared to the key (m is very large), and
2. the letters in the message are i.i.d. according to P_M .
3. the letters in the key are i.i.d. according to a uniform P_K .

(c) (2pt) Use (b) to find the most likely key length k in the following ciphertext:

QMIXFFKFMVPZSVYGOMAPOTTUUBGIIPPOTPZEVNVJHIUBGSCDCIPXBKDBR
SIXFMVIUVJBCWYNPMVNSDFNDZDWQCVFLKZCGMXYJNJBAQPXISUEFVGGIT
FMQNOEIATBWG

(d) (2pt) The ciphertext in the previous subexercise was sent to Cleopatra by Caesar. That means it will very likely contain one or more of the following words: CLEOPATRA, CAESAR, QUEEN, GRACE, ALEXANDRIA. Use this side information, in combination with the key length you computed in (c), to figure out the original message. What was the encryption key?