

Problem Set 7

Problem 1: RSA Signatures

- (a) Consider the plain RSA signature scheme with modulus $N = 85 = 5 \cdot 17$.
1. Compute $\phi(N)$.
 2. Say the public exponent is $e = 5$. Find the private exponent d . (Hint: look at the sequence $\phi(N)+1, 2 \cdot \phi(N)+1, \dots$ until you find a multiple of e .)
 3. Compute the signature on the message 3.
Hint: Note that $[3^4 \bmod 85] = [81 \bmod 85] = [(-4) \bmod 85]$.
- (b) Consider a *padded* RSA signature scheme, where to sign a message $m \in \{0, 1\}^{80}$ the signer chooses random r with $r \| m < N$ (where $\|$ denotes concatenation), computes $\sigma = [(r \| m)^d \bmod N]$, and outputs signature σ .
1. How would verification be done?
 2. Is this scheme secure? If yes, give a 1-2 sentence explanation; if not, show an attack.

Problem 2: Not a PRF

Consider the keyed function $H : \{0, 1\}^n \times \{0, 1\}^n \rightarrow \{0, 1\}^{2n}$ defined as: $H_k(x) = G(k) \oplus G(x)$, where $G : \{0, 1\}^n \rightarrow \{0, 1\}^{2n}$ is a pseudorandom generator.

- (a) Describe and formally analyze an explicit attack showing that H is not a PRF.
- (b) Is there a successful attack making a single query that distinguishes H_k (for random k) from a random function $f : \{0, 1\}^n \rightarrow \{0, 1\}^{2n}$? Why or why not?

Problem 3: A randomized variable-length MAC from a PRF

Let F be a pseudorandom function. Show that the following MAC is insecure for variable-length messages. Gen outputs a uniform $k \in \{0, 1\}^n$. Let $\langle i \rangle$ denote an $n/2$ -bit encoding of the integer i .

To authenticate a message $m = m_1 \| \dots \| m_\ell$, where $m_i \in \{0, 1\}^{n/2}$, choose a uniform $r \leftarrow \{0, 1\}^n$, compute $t := F_k(r) \oplus F_k(\langle 1 \rangle \| m_1) \oplus \dots \oplus F_k(\langle \ell \rangle \| m_\ell)$ and let the tag be (r, t) .

Problem 4: Cryptographic Mechanisms

For each of the following, identify the most appropriate cryptographic mechanism(s) (from among private-key encryption, pseudorandom generators, pseudorandom functions, message authentication codes, hash functions, public-key encryption, or digital signatures) for addressing the problem. Points will be deducted if you list extraneous mechanisms. **Explain your answer in 1-2 sentences.**

- (a) A company wants to distribute authenticated software updates to its customers.
- (b) A user wants to ensure secrecy of the files stored on his hard drive.
- (c) A customer wants to send his credit card number (confidentially) to a merchant over the web to complete a purchase.
- (d) A general wants to send a message to a lieutenant, and wants to ensure both confidentiality and integrity.
- (e) A client wants to store a short record of a large file he uploads to a server, so that the client can verify that the file has not been altered when it downloads the file later.
- (f) A user needs 1,000,000 random bits in order to run a simulation, but obtaining truly random bits is expensive.

Problem 5: Mode of Encryption

Let $F : \{0, 1\}^n \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ be a block cipher, and consider the following mode of encryption: to encrypt an ℓ -block message m_1, \dots, m_ℓ using key k , choose uniform $c_0 \in \{0, 1\}^n$ and then for $i = 1, \dots, \ell$ set $c_i := F_k(m_i) \oplus c_{i-1}$. Output the ciphertext c_0, \dots, c_ℓ .

- (a) How would decryption of a ciphertext c_0, \dots, c_ℓ be done?
- (b) Describe and analyze an explicit attack showing that this scheme is not EAV-secure.
- (c) Is this scheme CPA-secure? Provide a brief justification of your answer.