# Final Project Checkpoint Formalized Algebra from Text

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#### 1 Abstract

In this project, we are proposing to solve algebraic word problems in a way similar to Kushman et al (Kushman et al., 2014). However, rather than relying on a rule based approach (Mukherjee and Garain, 2008), or a mapping function to a set of templates (as in Kushman et al), we are proposing a slightly different method. The method that we are proposing is much closer to the solution in (Hosseini et al., 2014). In this method, we will learn a mapping from the input sequence of words, to the set of verbs and their syntactic dependents, that include some type of quantifier, via a dependency parse of the text. We will then categorize the verb expression to be equivalent to some command, expression, or set of commands and expressions in the semantics of a programming language that we define for solving basic algebraic problems.

The categorization task will be based on a semantic representation in a vector space model, using fasttext (Bojanowski et al., 2016). Each command or expression, will then have its arguments filled by one of the syntactic dependents of the verb, which will then map to the arguments in the described in the semantics of our target language. The categorization parameters are learned over the same data that (Hosseini et al., 2014) use, and will be assessed in order to compare results directly. The target language will be the simple **IMP-1** language.

### 2 Architecture

The general workflow of our system is pictured in figure 1

# 3 Example

**Example 1.** Pooja has 3 apples. She eats one apple. How many apples does Pooja have now?

We attempt to translate the sample problem's text into the following commands:

$$\begin{array}{lll} \textit{Pooja has 3 apples} & -> & a_p := 3 \\ \textit{She eats one apple} & -> & a_p := a_p - 1 \\ \textit{How many apples does Pooja have now?} & -> & \texttt{print } a_p \end{array}$$

And since these are sequence of steps we end up with:

$$a_p := 3; a_p := a_p - 1; print a_p$$

More formally:

$$seq(set[a_p](3); seq(set[a_p](minus(a_p; 1)); print(a_p)))$$

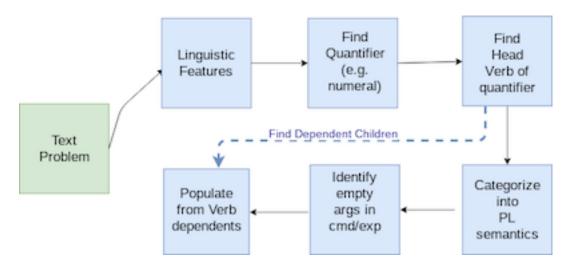


Figure 1: High level overview of the architecture

## 4 Language: IMP-1

 $e:\tau$ 

 $e_1$ : float  $e_2$ : float  $e_1$ : float  $e_2$ : float  $\overline{\mathsf{float}[\mathsf{f}]:\mathsf{float}}$  $\overline{\mathrm{addr}[a]:\mathrm{float}}$  $plus(e_1; e_2)$ : float  $minus(e_1;e_2):float$  $e_1$ : float  $e_2$ : float  $e_1$ : float  $e_2$ : float  $\overline{\mathtt{times}(e_1;e_2):\mathsf{float}}$  $\overline{\text{divide}(e_1;e_2):\text{float}}$ c:ok $c_1$  ok  $c_2$  ok e: float e: float  $\overline{\operatorname{\mathtt{set}}[a](e)\operatorname{\mathsf{ok}}}$  $\overline{\operatorname{seq}(c_1;c_2)\operatorname{ok}}$ print(e) ok skipok

### Syntax chart:

```
numbers
Тур
       	au ::= num
                                 num
                                                          booleans
                 bool
                                 bool
                                                          addresses (or "assignables")
Exp
       e ::= addr[a]
                                 a
                 num[n]
                                                          numeral
                                 n
                                                          boolean
                 bool[b]
                                 b
                 plus(e_1;e_2)
                                                          addition
                                 e_1 + e_2
                 times(e_1;e_2)
                                                          multiplication
                                 e_1 * e_2
                                                          equal
                 eq(e_1;e_2)
                                 e_1 == e_2
                                                          less-than-or-equal
                 le(e_1;e_2)
                                 e_1 <= e_2
                 not(e_1)
                                 !e_1
                                                          negation
                 and(e_1;e_2)
                                 e_1 \&\& e_2
                                                          conjunction
                 or(e_1;e_2)
                                                          disjunction
                                 e_1 || e_2
                 set[a](e)
\mathsf{Cmd} \quad c ::=
                                 a := e
                                                          assignment
                                                          skip
                 skip
                                 skip
                 seq(c_1;c_2)
                                                          sequencing
                                 c_1; c_2
                                                         conditional
                 if(e;c_1;c_2)
                                 if e then c_1 else c_2
                 while(e; c_1)
                                 while e \ \mathsf{do} \ c_1
                                                          looping
                 print(e)
                                 print e
                                                          printing
Addr a
```

### References

- Piotr Bojanowski, Edouard Grave, Armand Joulin, and Tomas Mikolov. Enriching word vectors with subword information. *arXiv preprint arXiv:1607.04606*, 2016.
- M. J. Hosseini, H. Hajishirzi, O. Etzioni, , and N. Kushman. Learning to solve arithmetic word problems with verb categorization. In *Empirical Methods in Natural Language Processing (EMNLP)*, pages 523–533, 2014.
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- Anirban Mukherjee and Utpal Garain. A review of methods for automatic understanding of natural language mathematical problems. *Artificial Intelligence Review*, 29, 2008.