

1 Cumulative Abstract Semantics

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4 TODO

5 1 Introduction(both)

6 Desired contributions of paper:

- 7 - we introduce cumulative semantics as an idea
- 8 - we give a formalization for: concrete, collecting, and abstract interpreter in the context of cumulative semantics
- 9 - a comparison between effect and typeclass based implementation of cumulative semantics
- 10 - soundness proofs?

11 2 Overview(both)

12 Often when one is crafting an analysis, they begin with a fixed target language in mind, and analysis
13 they would like to apply to that language. We begin this example in a different way: with a fixed
14 analysis we would like to do on a growing language. This analysis will be a concrete interpretation
15 of a simple language. To start this language we will introduce the idea of numbers:

$$22 \quad \begin{array}{ll} \text{numbers} & n \in \mathbb{Z} \\ 23 \quad \text{expressions} & \text{exp} := \text{cst}(n) \mid e_1 + e_2 \end{array}$$

24 Evaluation of this language can be done trivially in any number of ways, and we will choose one
25 that is most convenient for our purposes. We will define a judgment form $\llbracket \text{expr} \rrbracket$

26 3 Technical(both, maybe more andrew)

27 talk through formalization of it in here, step through ok this is the step 0 unsubstantiated semantics,
28 this is the collecting, this is the concrete, this is the symbolic. Look at how nice they play together.

29 3.1 Standard abstract interpretation recipe (review)

30 3.2 New recipe with these semantics

31 3.3 concrete

32 3.4 collecting

33 3.5 abstract(backwards)?

34 4 Evaluation

35 Two case studies:

36 4.1 Effect based (cade)

37 In here show concrete and backwards symbolic in effekt/koka

38 4.2 Typeclass based (andrew)

39 In here show concrete and backwards symbolic in lean

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50 5 Conclusion(both)

51 compare the two methodologies, talk about how great both are and the future of both: for the effect
 52 stuff (assuming it's easier to build) we can make a larger repo for many languages and accomplish
 53 a lot. For the lean one talk about how you get to reason about it for free.

54 6 Related Works(both)

55 *Van Horn Abstracting abstract machines.* [2]: Sergey built on this idea, it seems to allow for the
 56 isolation of more complex ideas, allowing for their simplification and independent reasoning. They
 57 accomplish this by abstracting the store (and making it finite), turning an infinite execution tree
 58 into a finite state graph.

59 *Sergey Monadic Abstract Interpreters.* [7]: This work seems to mostly be based off of a previous
 60 work (abstracting abstract machines) but also using a monad to capture state. Their essential claim
 61 seems to be similar to ours in that they are making the semantics orthogonal (and reusable) across
 62 different interpreters with the monadic approach. By using a monadic state you delay any choices
 63 in your interpreter until you actually assemble it. In my monadic attempt we take the idea of using
 64 a monadic state to track it easily through everything, keeping its representation orthogonal to
 65 what we want. The monadic implementation of intro and elim is similar to this paper, but slightly
 66 different because we retain a more precise control over control flow. Also, our version is done in
 67 a more modern and dependently typed language leaving room to reason and prove things if we
 68 wanted to.

69 *Michelland Monadic Modular Verification.* [5]: They advanced the usage of monadic modularity
 70 with their ITree-based framework for abstract interpretation, composing state and control flow
 71 monads for sound meta-theory in Roq. Their approach supports modular analyses (e.g., binding-
 72 time analysis) but requires complex transformer stacks. Their work is also done in a dependently
 73 type language like ours.

74 *Keidel Sturdy.* [3, 4]: The goal of this work is automatic soundness proofs for components, so
 75 not centrally aligned with what we are going for. Their approach provides soundness guarantees
 76 for free but is fixed to non-relational domains and requires an implementation of a monolithic,
 77 generic interpreter for every language. It accomplishes modularity with arrow transformers which
 78 are generally more rigid than effects or monads with control flow manipulation.

79 *Reinder Modular Effect Interpreter.* [6]: This short paper introduced the idea to make an effect
 80 based interpreter in theory, and played with a small toy example of it.

81 *Bunkenburg Making a Curry Interpreter.* [1]: They started using effect handlers embedded in
 82 Haskell to build a modular interpreter, one of the first of its kind. Its focus was on concrete
 83 interpretation of the Curry language, but Curry already supports non-determinism which is the
 84 same as non determinism in abstract interpretation. This demonstrated the usefulness of effects for
 85 control flow manipulation, which we use and extend to more generic domains in our work.

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