

Synchronous Dataflow Programming

CS684: Embedded Systems

Topic 5

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Summary

Multi-mode controller as FSA + Dataflow Equations.

- Automaton can be in one **active state** at each cycle.
- Each state is a mode with associated set of equations. Equations of the active state are applied.
- Each state is a name space and clock domain. **$pre(x)$** refers to mode local copy of x . Construct **$last(x)$** allows values to be shared between modes.
- Each state has outgoing **until** (or weak) transitions. After executing the equations of the active state, the guards of its until transitions are evaluated to decide the next state.
- Thus, weak state transition takes effect from the next cycle.
- In a **continue** type transition, the state change occurs without resetting the equations of the target state.
- In a **then** type transition, the state change occurs with resetting of the equations of the target state.

Structure of a Node with Concurrent Automata

```
node myautomaton () returns (y:int)
let
  y = x + z;

  automaton --autA
    state S1 do x = 10 -> pre(x)+1;
    ...
  end

  automaton --autB
    state T1 do z = 20;
    ...
  end
tel
```

Equations and automata all execute in parallel in lock-step (synchronous) manner.

Concurrent Automata: Example

```
node myautomaton () returns (ping,pong : bool)
let
  automaton -- A_ping
    state S1a
      do ping = true
    until true then S2a
    state S2a
      do ping = false
    until pong then S1a
  end;

  automaton -- A_pong
    state S1b
      do pong = false
    until ping then S2b
    state S2b
      do pong = true
    until true then S1b
  end
tel
```

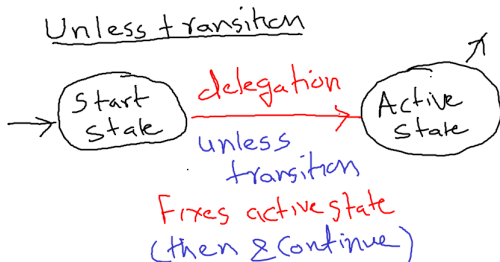
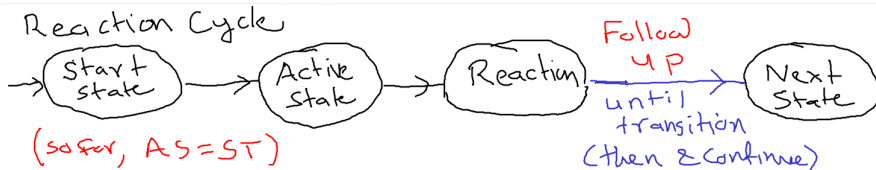
ST	(S1a, S1b)	(S2a, S2b)	(S1a, S1b)	(S2a, S2b)	(S1a, S1b)	...
ping	1	0	1	0	1	...
pong	0	1	0	1	0	...
NS	(S2a, S2b)	(S1a, S1b)	(S2a, S2b)	(S1a, S1b)	(S2a, S2b)	...

Automaton with until and reset transition (Revision)

```
node myautomaton(i : int; c: bool) returns (o: int; stup:bool)
  let
    automaton
      state Up
        do o = 60 -> i+1; stup = true;
      until c then Down
      state Down
        do o = 150 -> -2 * i; stup = false;
      until c then Up
    end
  tel
```

ST	<i>U</i>	<i>U</i>	<i>U</i>	<i>U</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>U</i>	<i>U</i>	<i>U</i>	<i>D</i>	<i>D</i>	...
i	4	4	3	3	3	3	3	3	3	3	3	3	3	...
c	0	0	0	1	0	0	0	1	0	0	1	0	0	...
o	60	5	4	4	150	-6	-6	-6	60	4	4	150	-6	...
stup	1	1	1	1	0	0	0	0	1	1	1	0	0	...
NS	<i>U</i>	<i>U</i>	<i>U</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>U</i>	<i>U</i>	<i>U</i>	<i>D</i>	<i>D</i>	<i>D</i>	...

Reaction Cycle of an Automaton



Unless Transitions: Motivation

Unless transitions are also called strong transitions.

- They allow **delegation** of reacting in current cycle to another state.
- A state can have one or more unless transitions.
- Guard of an unless transition is evaluated before looking at the equations (i.e. before reacting).

Hence, guard cannot use current values of equations.

- If guard is true, **control moves to the target state in the same cycle**.
- Thus, target state becomes the active state. Its equations are applied.
- If the guards of all unless transitions are false, the current state remains the active state.
- After the reaction, the until transitions of the active state decides the next state.
- At most one delegation is allowed per cycle.

Reaction Cycle

- **Start State**
- **Delegation** Apply **unless** transition to determine the **active state** of the current cycle.
- **Reaction** Equations of the active state are applied to compute output from input.
- **Followup** AFTER the reaction, apply the **until** transition of the **active state**:
 - Guard of each until transition is evaluated.
 - Guard can refer to outputs the equations.
 - If guard is true the transition is taken and next state is changed to target state.
- **Next State**: this is the start state of the **next cycle**.
- For both unless and until transitions, the equations are reset for **then** type transition. They are not reset for a **continue** type transition.

Unless Transitions: Example

```
node myautomaton(i : int; c: bool) returns (o: int; stup:bool)
```

```
  let
```

```
    automaton
```

```
      state Up
```

```
        do o = 60 -> i+1; stup = true;
```

```
→ unless c then Down
```

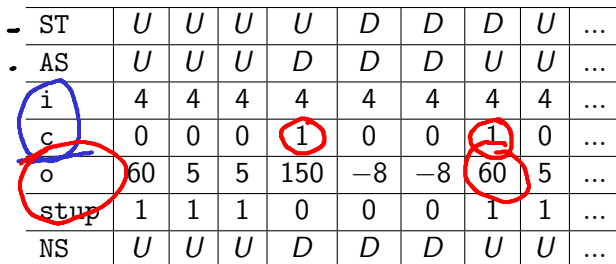
```
      state Down
```

```
        do o = 150 -> -2 * i; stup = false;
```

```
        unless c then Up
```



```
      end
```


































```
tel
```



- ST	U	U	U	U	D	D	D	U	...
. AS	U	U	U	D	D	D	U	U	...
i	4	4	4	4	4	4	4	4	...
c	0	0	0	1	0	0	1	0	...
o	60	5	5	150	-8	-8	60	5	...
stup	1	1	1	0	0	0	1	1	...
NS	U	U	U	D	D	D	U	U	...

Mixing unless and until Transitions: Example

```
node myautomaton(i : int; c: bool) returns (o: int; stup:bool)
  let
    automaton
      state Up
        do o = 60 -> i+1; stup = true; 
      unless c then Down
      state Down
        do o = 150 -> -2 * i; stup = true; 
      until c then Up
    end
  end
tel
```

ST												...
AS												...
i	4	4	4	4	4	4	4	4	4	4	4	...
c	0	0	0	1	0	0	0	1	1	0	0	...
o	60	5	5	<u>150</u>	60	5	5	150	150	60	5	...
stup	1	1	1	0	1	1	1	0	0	1	1	...
NS												...

Example: Farm Road Traffic Light Controller

A farm road (or side road) crosses a main road. Traffic light controller must turn on or off the lights `maingreen`, `mainred`, `sidegreen`, `sidered`. An input `carwait` is true if a car is waiting on the farm road. Input `second` is the timer input which becomes true for one clock cycle every one second. Thus the count of “second” gives how much time has elapsed.

Controller: Farm Road Traffic Light

```
node traffic(carwait,second:bool)
  returns (maingreen, mainred,sidegreen,sidered:bool)

let
  automaton
    state Maingreen
      do

        until ??
        state Sidegreen
          do

            until ??
          end
        end
      tel
```

Sidegreen

Maingreen

```

node traffic(carwait,second:bool)
    returns (maingreen, mainred,sidegreen,sidered:bool)
var timegreen:int;
let
    automaton
        state Maingreen
            do

                maingreen = true; mainred = false;
                sidegreen = false; sidered = true;
            until ?? then Sidegreen
        state Sidegreen
            do

                maingreen = false; mainred = true;
                sidegreen = true; sidered = false;
            until ?? then Maingreen
    end
tel

```

```

node traffic(carwait,second:bool)
    returns (maingreen, mainred,sidegreen,sidered:bool)
var timegreen:int;
let
    automaton
        state Maingreen
            do timegreen = 180 -> if (((pre(timegreen)) > 0) and second) then
                pre(timegreen)-1 else pre(timegreen);
                maingreen = true; mainred = false;
                sidegreen = false; sidered = true;
            until (timegreen <= 0) then Sidegreen
        state Sidegreen
            do
                timegreen = 60 ->
                    if (((pre(timegreen)) > 0) and second) then pre(timegreen)-
                        else pre(timegreen);
                    maingreen = false; mainred = true;
                    sidegreen = true; sidered = false;
            until (timegreen <=0) then Maingreen
        end
    tel

```

Controller: Farm Road Traffic Light

```
node traffic(carwait,second:bool)
  returns (maingreen, mainred,sidegreen,sidered:bool)
var timegreen:int;
let
  automaton
    state Maingreen
      do timegreen = 180 -> if (((pre(timegreen)) > 0) and second) then
        pre(timegreen)-1 else pre(timegreen);
        maingreen = true; mainred = false;
        sidegreen = false; sidered = true;
      until ((timegreen <= 0) and carwait) then Sidegreen
    state Sidegreen
      do
        timegreen = 60 ->
          if (((pre(timegreen)) > 0) and second) then pre(timegreen)-
            else pre(timegreen);
          maingreen = false; mainred = true;
          sidegreen = true; sidered = false;
        until ((timegreen <=0) and not carwait) then Maingreen
      end
  end
```

Synchronous Dataflow Programming using language **Heptagon**:

- Finite State Automata with Synchronous Data Flow Equations.
- Automata provide modes of operation.
Dataflow equations define equations for transforming input into output.
- Highly complex modes of operation can be programmed:
Hierarchy, concurrency, communication with shared flows.
- A heptagon program (node) compiles into a **reactive kernel**. Target C and Java.

A modelling language.

- Clean Semantics
- Deterministic execution: easy to certify.
- Simulation, testing and verification tools.
- Qualified code generator,
- Efficient Code
- Industrial Strength: Used widely in Nuclear reactor control, Aerospace and railway industries.