

The background image is a photograph of a large, modern server room. It features rows of server racks with glowing lights, a complex network of overhead cables, and a high ceiling with industrial lighting. The overall color palette is dominated by blues and greys, with some warm light from the server racks.

Version 2.15

UCSD CSE 30

Computer Organization and Systems Programming

Arm Assembly – Part 4

Keith Muller

Function Calls

Branch with Link (function call) instruction

`bl label`

`bl`

`imm24`

- Function call to the instruction with the address `label` (no local labels for functions)
 - `imm24` number of instructions from pc+8 (24-bits)
 - `label` any function label in the current file, any function label that is defined as `.global` in any file that it is linked to, any C function that is not static

Branch with Link Indirect (function call) instruction

`blx Rm`

`blx`

`Rm`

- Function call to the instruction whose address is stored in Rm (Rm is a function pointer)
- `bl` and `blx` both save the address of the instruction immediately following the `bl` or `blx` instruction in register `lr` (link register is also known as r14)
- The contents of the link register is the return address in the calling function

- (1) Branch to the instruction with the label f1
- (2) copies the address of the instruction AFTER the `bl` in `lr`

main:

•

`bl f1`

→

`f1:`

•

•

Function Call Return

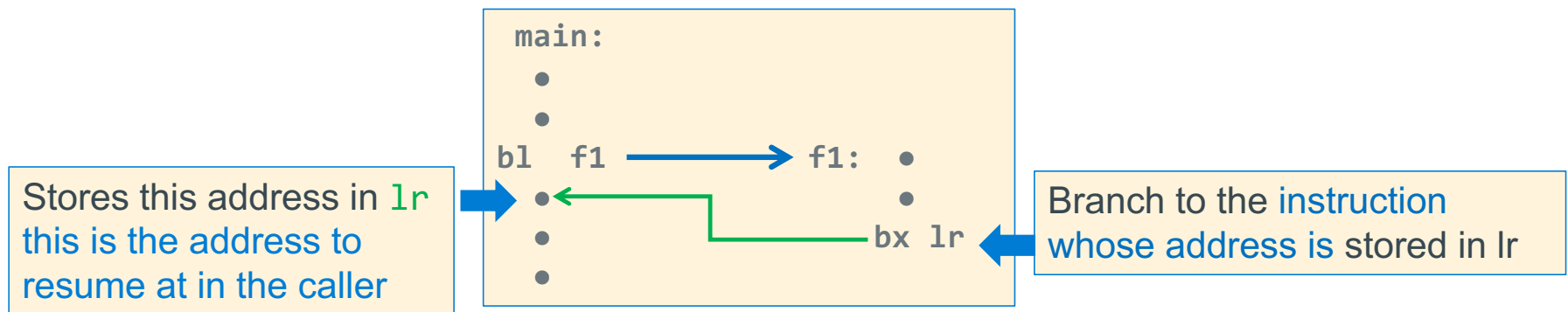
Branch & exchange (function return) instruction

`bx lr`

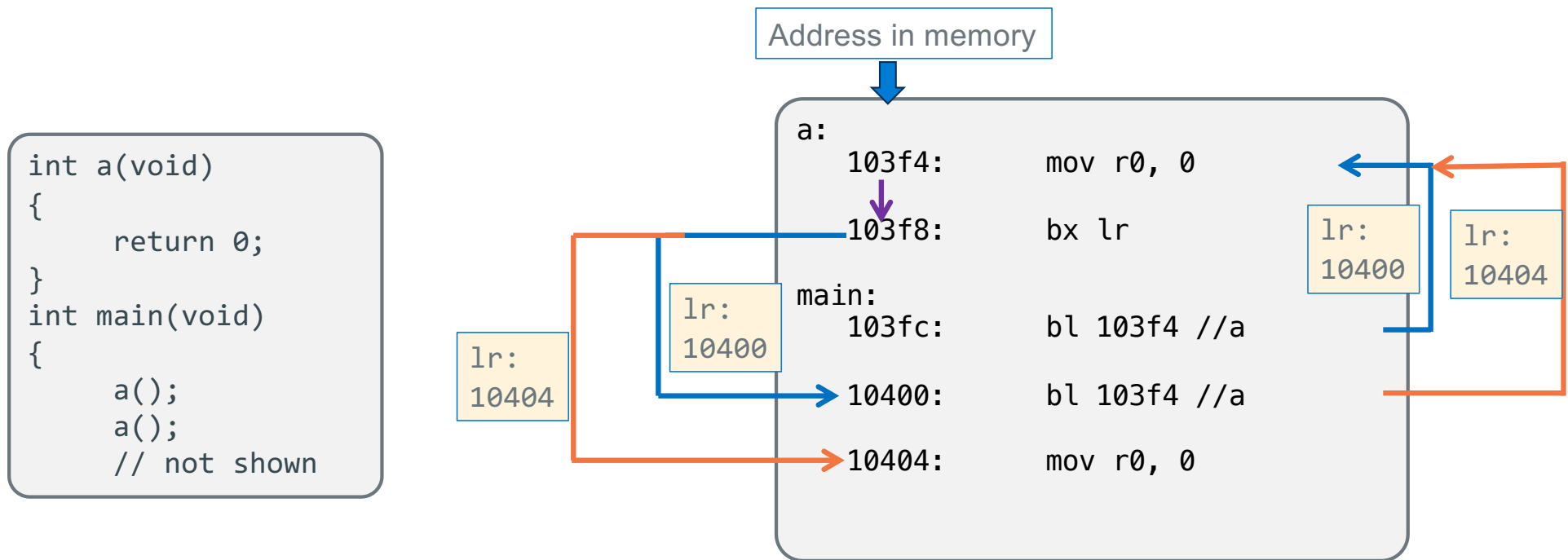


// we will always use lr

- Causes a branch to the instruction whose address is stored in register <lr>
 - It copies lr to the PC
- This is often used to implement a return from a function call (exactly like a C return) when the function is called using either `bl label`, or `blx Rm`



Understanding bl and bx - 1

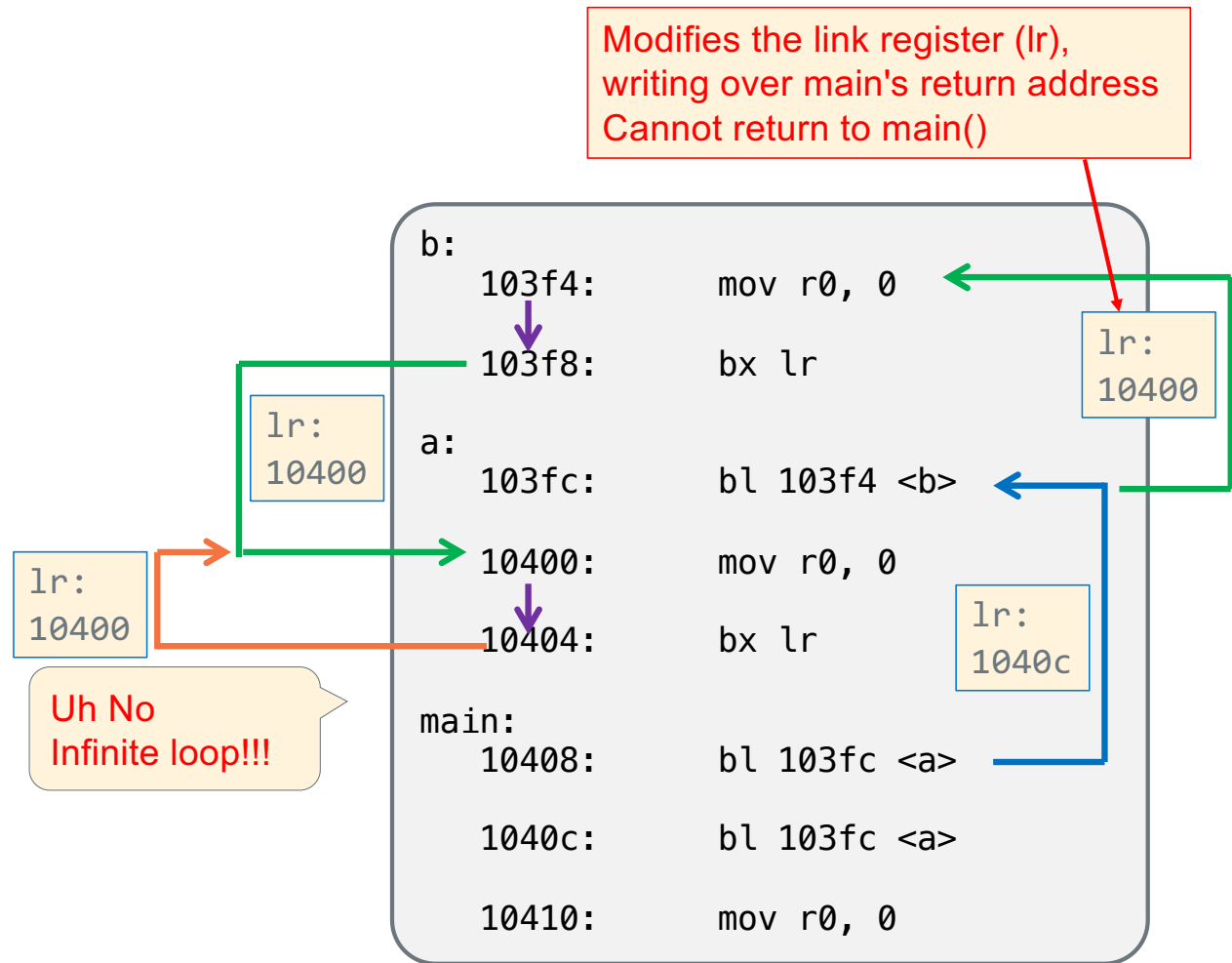


But there is a problem we must address here – next slide

Understanding bl and bx - 2

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
    // not shown
}
```

We need to preserve the lr!



Understanding bl and blx - 3

```
int a(void)
{
    return 0;
}

int (*func)() = a;

int main(void)
{
    (*func)();
    // not shown
}
```

But this has the same infinite loop problem when main() returns!

```
.data
func: .word a // func initialized with address of a()

.text
.global a
.type a, %function
.equ FP_OFF, 4

a:
    mov     r0, 0
    bx      lr
    .size a, (. - a)

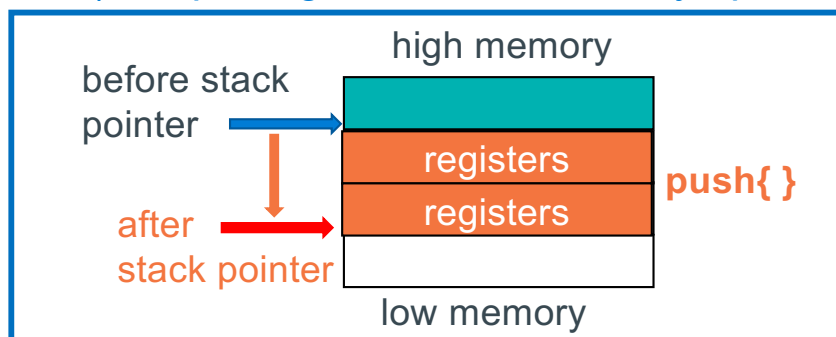
.global main
.type main, %function
.equ FP_OFF, 4

main:
    ldr     r4, =func // load address of func in r4
    ldr     r4, [r4]   // load contents of func in r4
    blx     r4         // we lose the lr for main!
    // not shown
    bx      lr         // infinite loop!
```

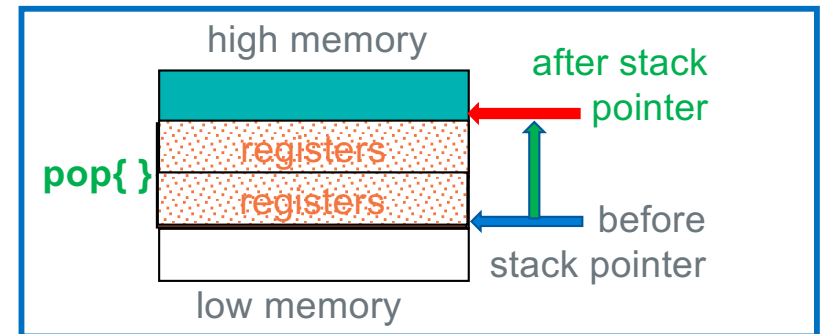
Preserving and Restoring Registers on the stack - 1

Operation	Pseudo Instruction	Operation
Push registers Function Entry	<code>push {reg list}</code>	$sp = sp - 4 \times \text{\#registers}$ Copy registers to <code>mem[sp]</code>
Pop registers Function Exit	<code>pop {reg list}</code>	Copy <code>mem[sp]</code> to registers, $sp = sp + 4 \times \text{\#registers}$

push (multiple register **str** to memory operation)



push (multiple register **ldr** from memory operation)

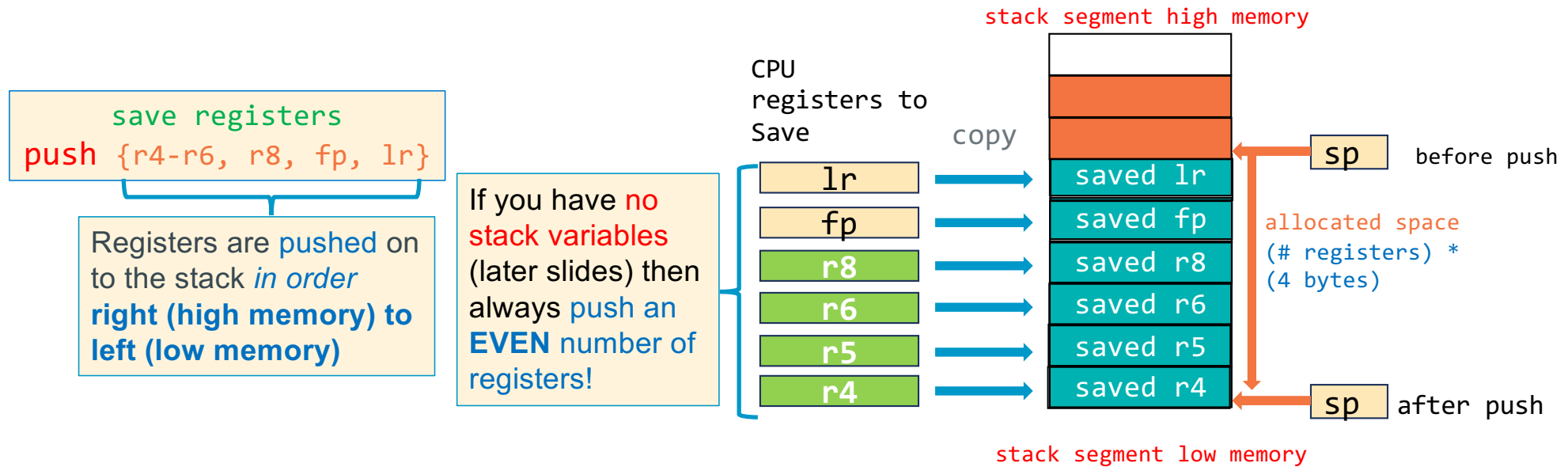


Preserving and Restoring Registers on the Stack - 2

Operation	Pseudo Instruction	Operation
Push registers Function Entry	<code>push {reg list}</code>	$sp = sp - 4 \times \text{\#registers}$ Copy registers to <code>mem[sp]</code>
Pop registers Function Exit	<code>pop {reg list}</code>	Copy <code>mem[sp]</code> to registers, $sp = sp + 4 \times \text{\#registers}$

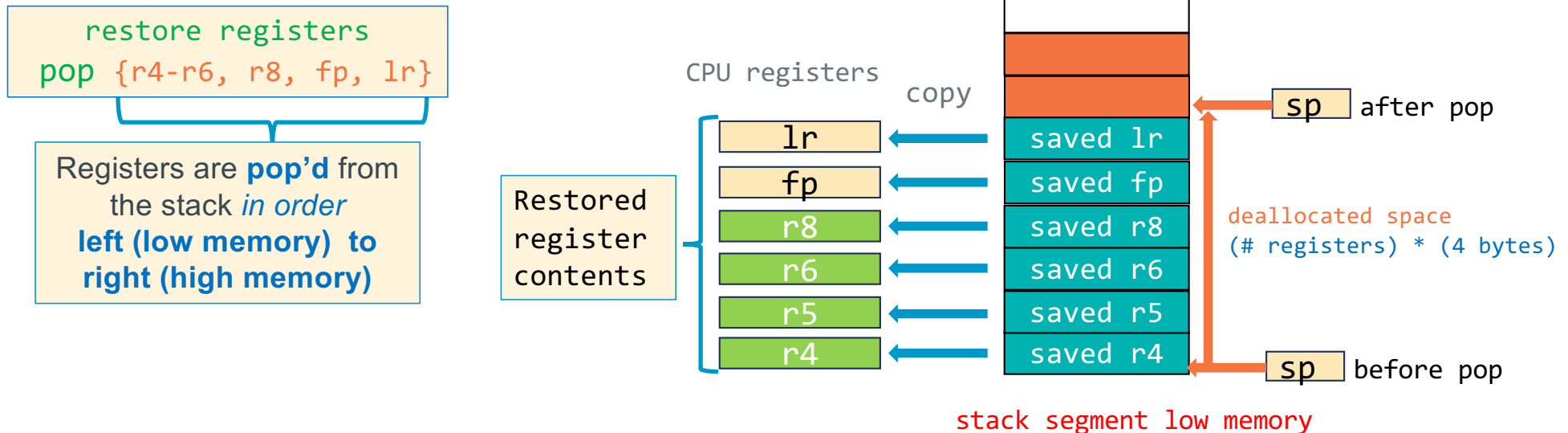
- `{reg list}` is a **list of registers in numerically increasing order, left to right**
`push {r4-r10, fp, lr}` // *fp is r11, lr is r14*
- Registers **cannot be**:
 1. duplicated in the list
 2. listed out of increasing numeric order (left to right)
- Register ranges can be specified `{r4, r5, r8-r10, fp, lr}`
- **Never!** push/pop `r12, r13, or r15`
 - the top two registers on the stack must always be `fp, lr` // ARM function spec – later slides

push: Multiple Register Save to the stack



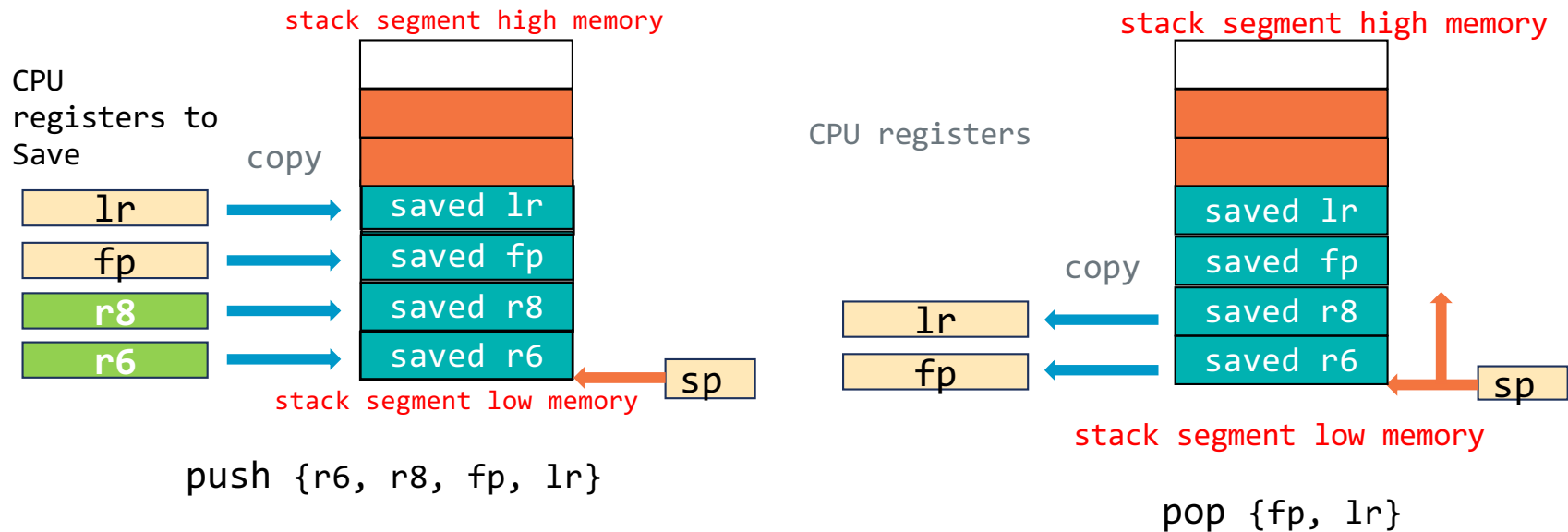
- **push** copies the contents of the **{reg list}** to stack segment memory
- **push** subtracts $(\text{\# of registers saved}) * (4 \text{ bytes})$ from the **sp** to **allocate** space on the stack
 - $sp = sp - (\text{\# registers_saved} * 4)$
- **this must always be true: $sp \% 8 == 0$**

pop: Multiple Register Restore from the stack



- **pop** copies the contents of stack segment memory to the **{reg list}**
- **pop adds:** (# of registers restored) * (4 bytes) to **sp** to **deallocate** space on the stack
 - $sp = sp + (\text{\# registers restored} * 4)$
- **Remember:** **{reg list}** must be the same in both the **push** and the corresponding **pop**

Consequences of inconsistent push and pop operands



- `lr` gets contents of `saved r8`, likely causing a segmentation fault when the `bx lr` is executed at function exit

Registers: Rules For Use

<i>Register</i>	<i>Function Call Use</i>	<i>Function Body Use</i>	<i>Save before use Restore before return</i>
r0	arg1 and return value	scratch registers	No
r1-r3	arg2 to arg4	scratch registers	No
r4-r10	preserved registers	contents preserved across function calls	Yes
r11 / fp	stack frame pointer	Use to locate variables on the stack	Yes
r12 / ip	may used by assembler with large text file	can be used as a scratch if really needed	No
r13 / sp	stack pointer	stack space allocation	Yes
r14 / lr	link register	contains return address for function calls	Yes
r15	Do not use	Do not use	No

Return Value and Passing Parameters to Functions

(Four parameters or less)

Register	Function Call Use	Function Body Use	Save before use Restore before return
r0	arg1 and return value	scratch registers	No
r1-r3	arg2 to arg4	scratch registers	No

- Where **r0**, **r1**, **r2**, **r3** are arm registers, the function declaration is (first four arguments):
`r0 = function(r0, r1, r2, r3) // 32-bit return`
- Each **parameter and return value** is limited to data that **can fit in 4 bytes or less**
- **Calling function:**
 - copy up to the first four parameters into these four registers before calling a function
 - **MUST assume** that the called function will **alter the contents of all four registers: r0-r3**
 - **In terms of C runtime support, these registers contain the copies given to the called function**
 - **C allows the copies to be changed in any way by the called function**
- **Called function:**
 - you receive **the first four parameters in these four registers (r0 – r3)**

Return Value and Passing Parameters to Functions

(Four parameters or less)

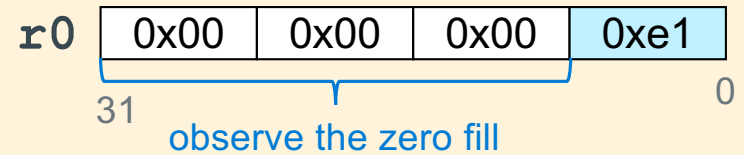
Register	Function Call Use	Function Body Use	Save before use Restore before return
r0	arg1 and return value	scratch registers	No
r1-r3	arg2 to arg4	scratch registers	No

- Where **r0**, **r1**, **r2**, **r3** are arm registers, the function declaration is (first four arguments):
`r0 = function(r0, r1, r2, r3) // 32-bit return`
- For parameters, whose size is larger than 4 bytes, pass a pointer to the parameter (we will cover this later)
- One arg value per register!** – NO arrays across multiple registers
 - chars, shorts and ints are directly stored
 - Structs (not always), and arrays (always) are passed via a pointer
 - Pointers** passed as output parameters contain an address *that points at* the stack, BSS, data, or heap

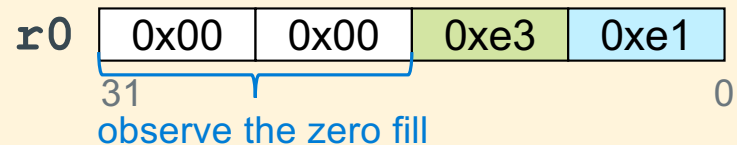
Register Arguments and Return Values

- When passing or returning values from a function you must do the following:
 - Make sure that the values in the registers r0-r3 are in their **properly aligned position in the register based on data type**
 - Upper bytes in byte and halfword values in registers r0-r3 when passing arguments and returning values **are**
 - zero filled for unsigned values
 - sign extended for signed values

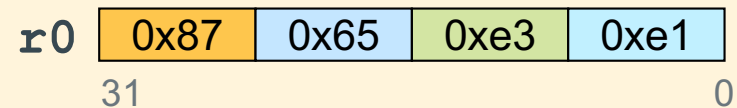
Single Byte (unsigned char)



Single Halfword (unsigned short)



Full Word (int or pointer)



What it means to be a Temporary/argument register

```
int a(void)
{
    // not shown
}
int main(void)
{
    int r0 = 0;
    int r1 = 1;
    int r2 = 2;
    int r3 = 3;
    r0 = a();
    // in C r1 and r3 would have the same values
    // after the call
```

```
// main()
// code not shown
mov r0, 0
mov r1, 1
mov r2, 2
mov r3, 3
bl a
// r0 = return value
// r1-r3 values are unknown as a() has right to change them as it wants
```

Preserved Registers

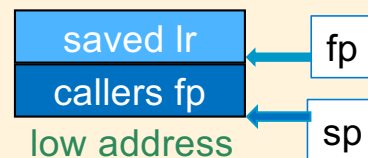
Register	Function Call Use	Function Body Use	Save before use Restore before return
r4-r10	preserved registers	contents preserved across function calls	Yes
r11/fp	stack frame pointer	Use to locate variables on the stack	Yes
r13/sp	stack pointer	stack space allocation	Yes
r14/lr	link register	contains return address for function calls	Yes

- **Any value** you have in a **preserved register before a function call will still be there after the function returns** (Contents are “preserved” across function calls)
- If the function **wants to use a preserved register** it must:
 1. **Save** the **value contained in the register** at function entry
 2. Use the register in the body of the function
 3. **Restore** the **original saved value** to the register at **function exit** (before returning to the caller)
- You use a preserved register when a function makes calls another function and you have:
 1. Local variables allocated to be in registers
 2. Parameters passed to you (in **r0-r3**) that **you need to continue to use after calling another function**

Minimum Stack Frame (Arm Arch32 Procedure Call Standards)

- Minimal frame: allocating at function entry: **push {fp, lr}**

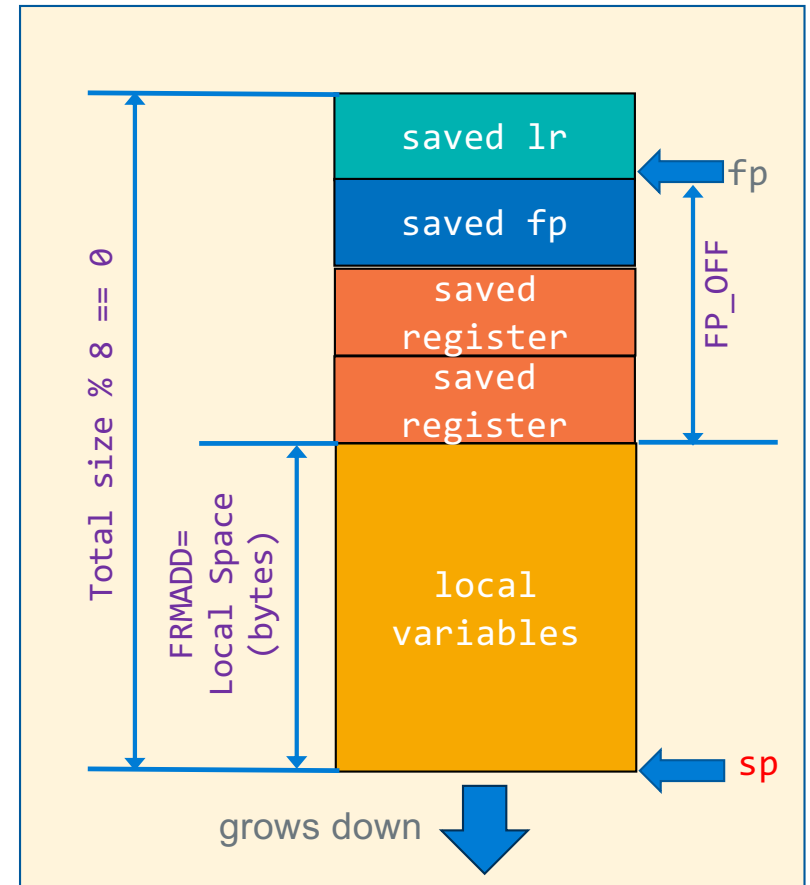
Minimum stack frame



- **sp** always points at top element in the stack (lowest byte address)
- **fp** always points at the bottom element in the stack
 - Bottom element is always the saved **lr** (contains the return address of caller)
 - A saved copy of **callers fp** is always the next element below the **lr**
 - **fp** will be used later when referencing stack variables
- Minimal frame: deallocating at function exit: **pop {fp, lr}**
- On function entry: **sp** must be 8-byte aligned (**sp % 8 == 0**)

First Look: A typical Stack Frame

- Saved lr and fp of the caller (so function calls work)
- Save values for any preserved registers this function will change
- Space (FRMADD) for local variables is allocated on the stack right below the lowest pushed register



Function Prologue and Epilogue

```
.global myfunc
.type    myfunc, %function
.equ     FP_OFF, 4           // fp distance to sp after push
.equ     FRMADD, 8          // number of bytes for local stack vars

myfunc:
{
    push    {fp, lr}        // push (save) fp and lr on stack
    add     fp, sp, FP_OFF  // set fp at bottom of stack
    add     sp, sp, -FRMADD // allocate FRMADD bytes for local vars
                           // by moving sp
    // your code here
}

{
    sub     sp, fp, FP_OFF  // deallocate local variables by moving sp
    pop     {fp, lr}       // pop (restore) fp and lr from stack
    bx      lr             // return to caller
}

.size myfunc, (. - myfunc)
```

Function Prologue creates stack frame

Function Epilogue removes stack frame

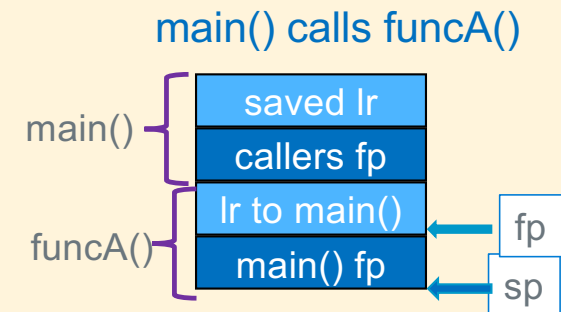
- **Only one prologue** right after the function label (name)
- **Only one epilogue** at the bottom of the function right above the `.size` directive

Minimum Stack Frame (Arm Arch32 Procedure Call Standards)

- **Function entry (Function Prologue):**

1. save lr and fp registers (push)
2. set fp to top entry in stack
3. allocate space for local vars – later slides

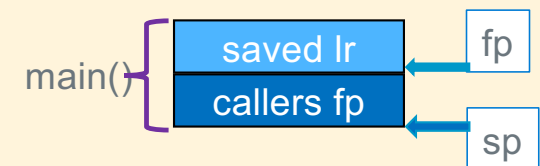
allocate stack space
 $SP = SP - \text{"space"}$
grows "down"



- **Function return (Function Epilogue):**

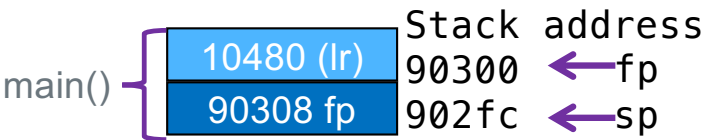
1. deallocate space for locals -later
2. restores lr and fp registers (pop)
3. Return To Caller

deallocate stack space
 $SP = SP + \text{"space"}$
shrinks "up"



Using Minimal Stack Frames

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
}
```



b:
103f4: push {fp, lr}
103f8: add fp, sp, 4
103fc: mov r0, 0
10400: sub sp, fp, 4
10404: pop {fp, lr}
10408: bx lr

a:
1040c: push {fp, lr}
10410: add fp, sp, 4
10414: bl 103f4
10418: mov r0, 0
1041c: sub sp, fp, 4
10420: pop {fp, lr}
10424: bx lr

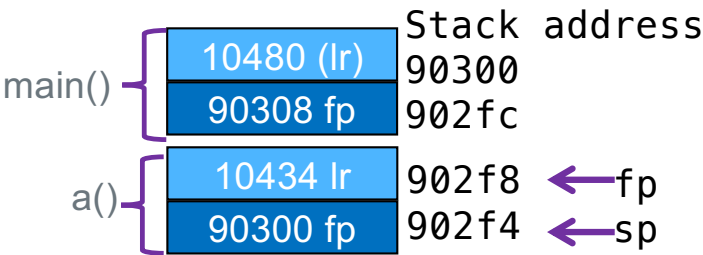
main:
10428: push {fp, lr}
1042c: add fp, sp, 4
10430: bl 1040c <a>
10434: bl 1040c <a>

// not shown

Memory address

Using Minimal Stack Frames

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
}
```



b:

```
103f4:    push {fp, lr}
103f8:    add fp, sp, 4
103fc:    mov r0, 0
10400:    sub sp, fp, 4
10404:    pop {fp, lr}
10408:    bx lr
```

a:

```
1040c:    push {fp, lr}
10410:    add fp, sp, 4
10414:    bl 103f4 <b>
10418:    mov r0, 0
1041c:    sub sp, fp, 4
10420:    pop {fp, lr}
10424:    bx lr
```

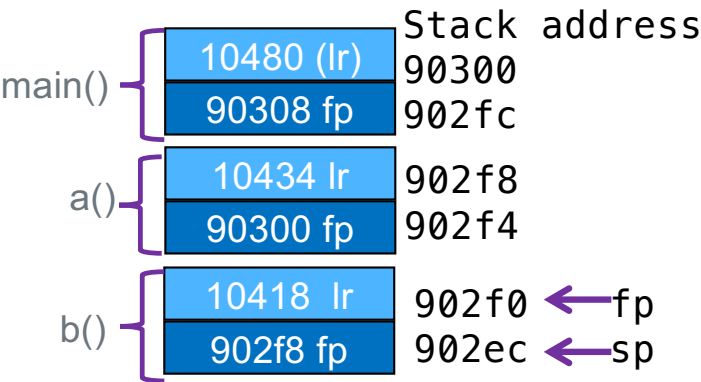
main:

```
10428:    push {fp, lr}
1042c:    add fp, sp, 4
10430:    bl 1040c <a>
10434:    bl 1040c <a>
// not shown
```

lr:
10434

Using Minimal Stack Frames

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
}
```



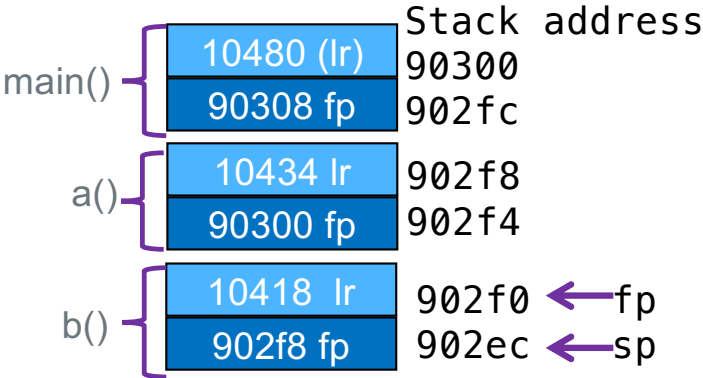
```
b:
103f4:  push {fp, lr}
103f8:  add fp, sp, 4
103fc:  mov r0, 0
10400:  sub sp, fp, 4
10404:  pop {fp, lr}
10408:  bx lr

a:
1040c:  push {fp, lr}
10410:  add fp, sp, 4
10414:  bl 103f4 <b>
10418:  mov r0, 0
1041c:  sub sp, fp, 4
10420:  pop {fp, lr}
10424:  bx lr

main:
10428:  push {fp, lr}
1042c:  add fp, sp, 4
10430:  bl 1040c <a>
10434:  bl 1040c <a>
// not shown
```

Using Minimal Stack Frames

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
}
```



lr:
10418

b:

103f4: push {fp, lr} ←

103f8: add fp, sp, 4

103fc: mov r0, 0

10400: sub sp, fp, 4

10404: **pop** {fp, lr} ←

10408: bx lr

lr: 10418

a:

1040c: push {fp, lr} ←

10410: add fp, sp, 4

10414: bl 103f4 ←

10418: mov r0, 0

1041c: sub sp, fp, 4

10420: **pop** {fp, lr} ←

10424: bx lr

lr: 10434

main:

10428: push {fp, lr}

1042c: add fp, sp, 4

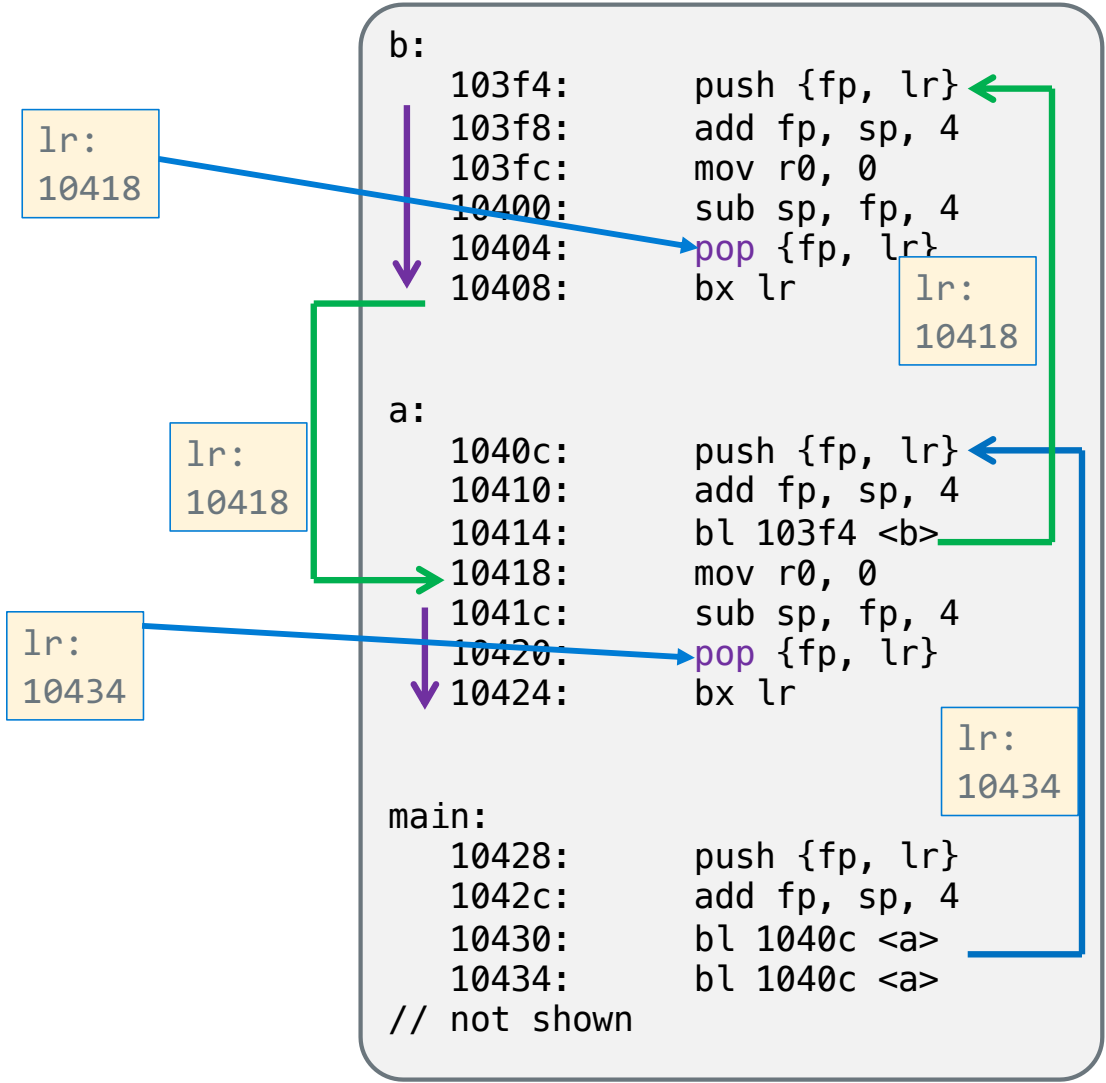
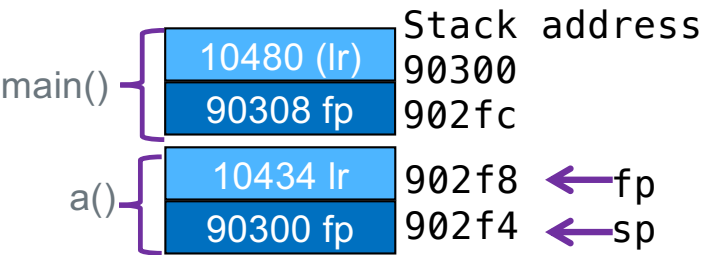
10430: bl 1040c <a> ←

10434: bl 1040c <a> ←

// not shown

Using Minimal Stack Frames

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
}
```



Using Minimal Stack Frames

```
int b(void)
{
    return 0;
}
int a(void)
{
    b();
    return 0;
}
int main(void)
{
    a();
    a();
}
```

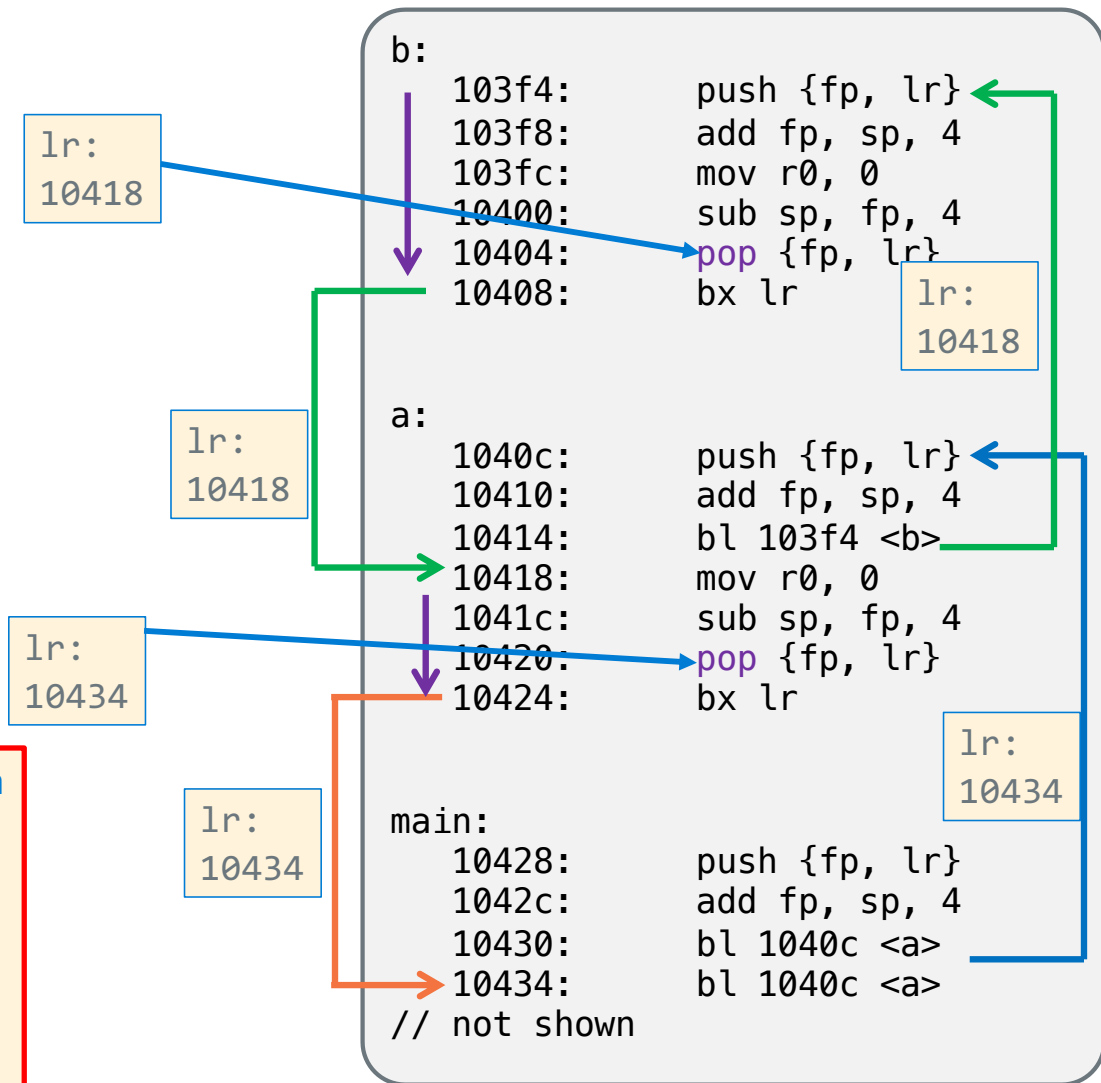
main() {

	Stack address	
10480 (lr)	90300	← fp
90308 fp	902fc	← sp

}

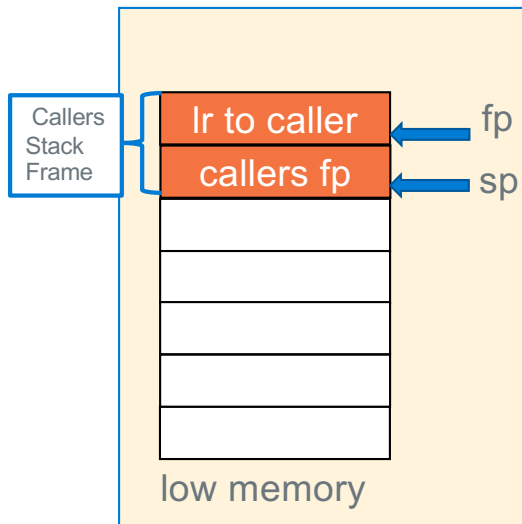
We are saving the lr on the stack on each function call and restoring it before returning.

Result: NO infinite loop and we return to the correct instruction in the caller no matter how many functions we call.
Even recursion will work!



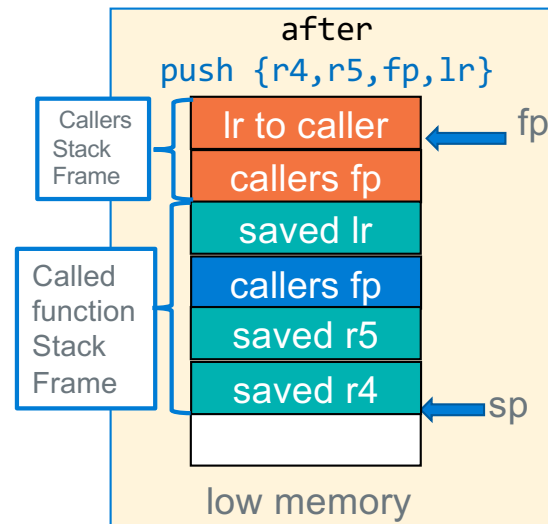
Function Prologue: Allocating the Stack Frame -1

at function entry



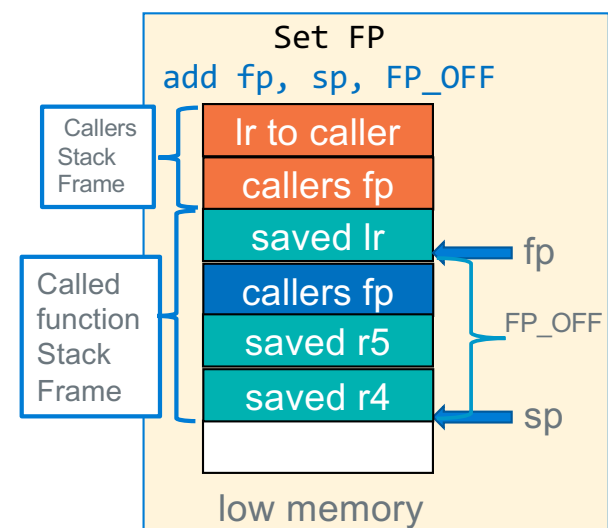
Function was just called this how the stack looks
The orange blocks are part of the caller's stack frame

Prologue Step 1 of 3



using a push, save lr, fp and those preserved registers it wants to use on the stack

Prologue Step 2 of 3



move the fp to point at the saved lr as required by the Aarch32 spec

myfunc:

Function Prologue

```
push    {fp, lr}
add     fp, sp, FP_OFF
add     sp, sp, -FRMADD
```

```
// push (save) fp and lr on stack
// set fp for this function
// allocate FRMADD bytes for local vars
// by moving sp
```


Function Prologue: Allocating the Stack Frame - 2 Prologue Step 3 of 3

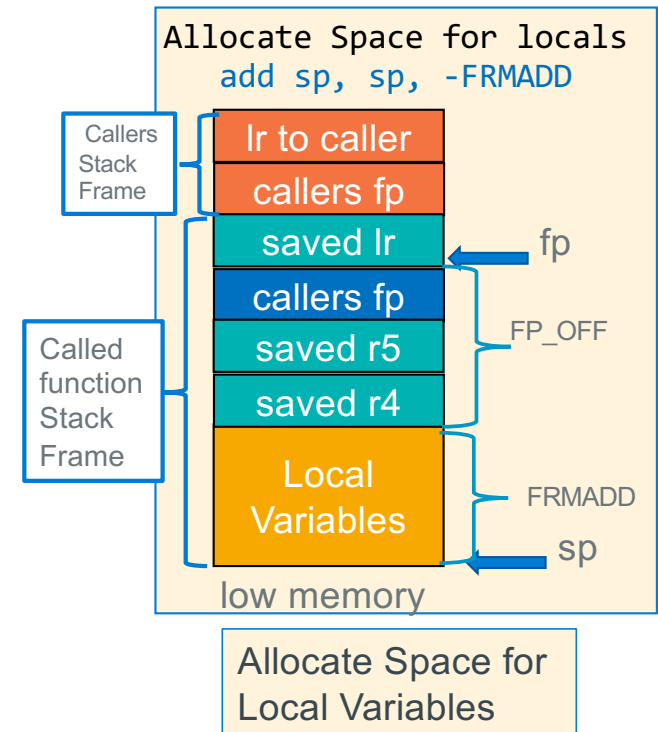
- Space for local variables is allocated on the stack right below the lowest pushed register
- Add memory to the stack frame for local variables** by **moving** the **sp** **towards low memory**
- The amount moved is the total size of all local variables in bytes **plus** memory alignment **padding**

$FRMADD = \text{total local var space (bytes)} + \text{padding}$

- Allocate the space after the register push by

```
add    sp, sp, -FRMADD
```

- fp** (frame pointer) is used as a **pointer (base register)** to access all stack variables – later slides

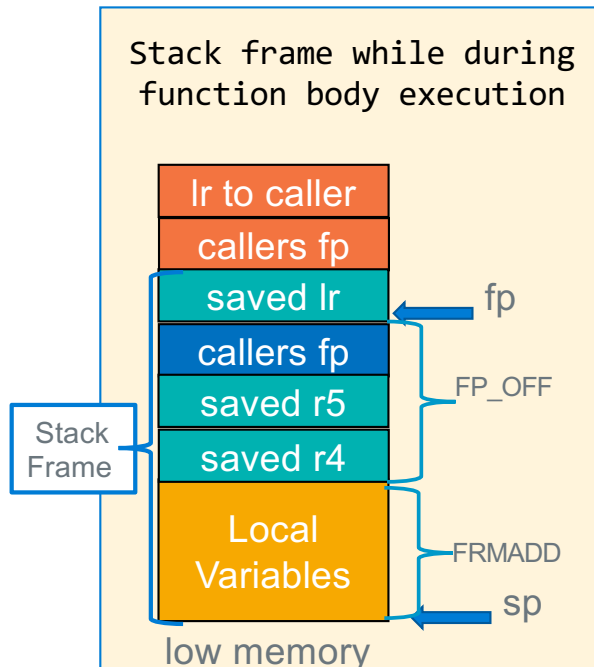


myfunc:

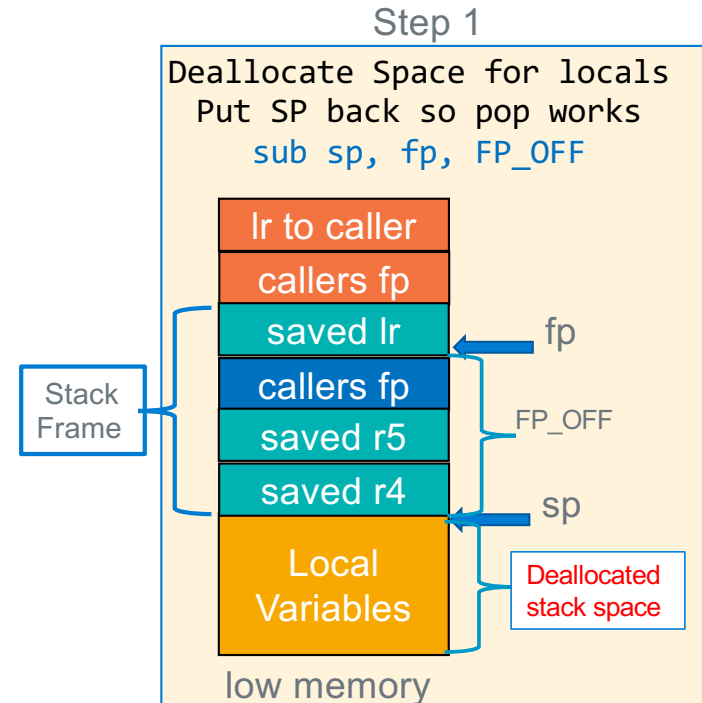
Function Prologue

```
push    {fp, lr}           // push (save) fp and lr on stack
add     fp, sp, FP_OFF      // set fp for this function
add     sp, sp, -FRMADD     // allocate FRMADD bytes for local vars
                                // by moving sp
```

Function Epilogue: Deallocating the Stack Frame - 1



Use fp as a pointer to find local variables on the stack



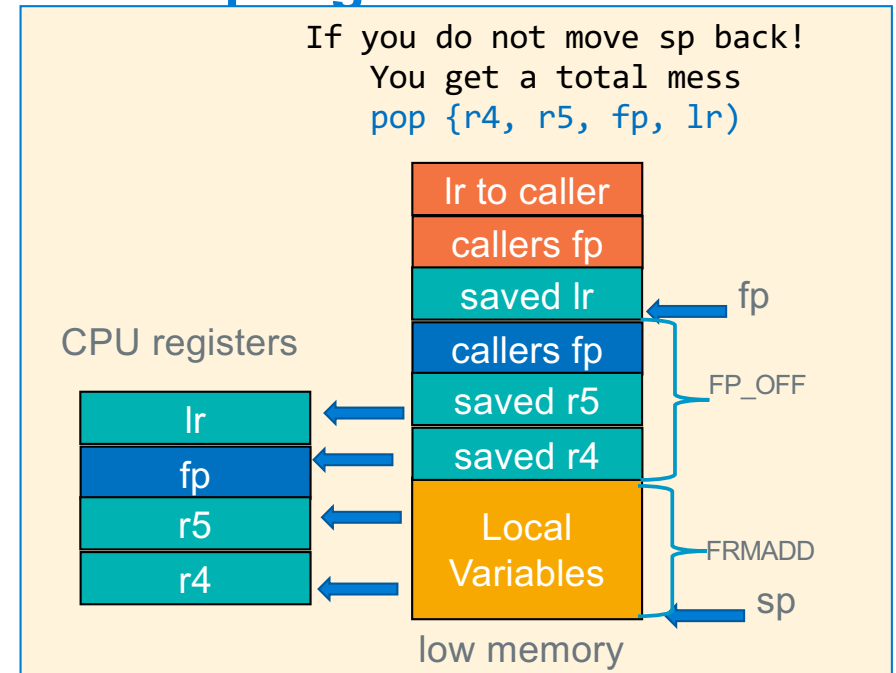
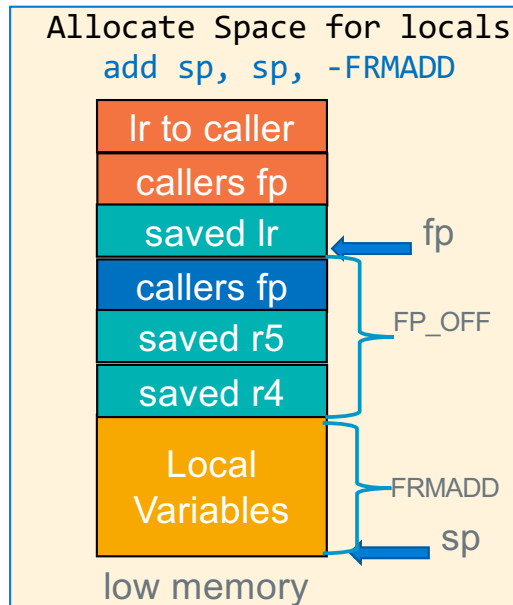
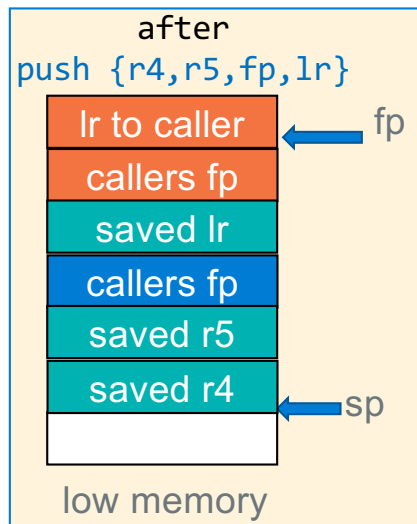
Move SP back to where it was after the push in the prologue.
So, pop works properly (this also deallocates the local variables)

function
Epilogue

```
sub    sp, fp, FP_OFF
pop    {fp, lr}
bx     lr
```

```
// deallocate local variables by moving sp
// pop (restore) fp and lr from stack
// return to caller
```

Why You must move SP before POP in the Epilogue



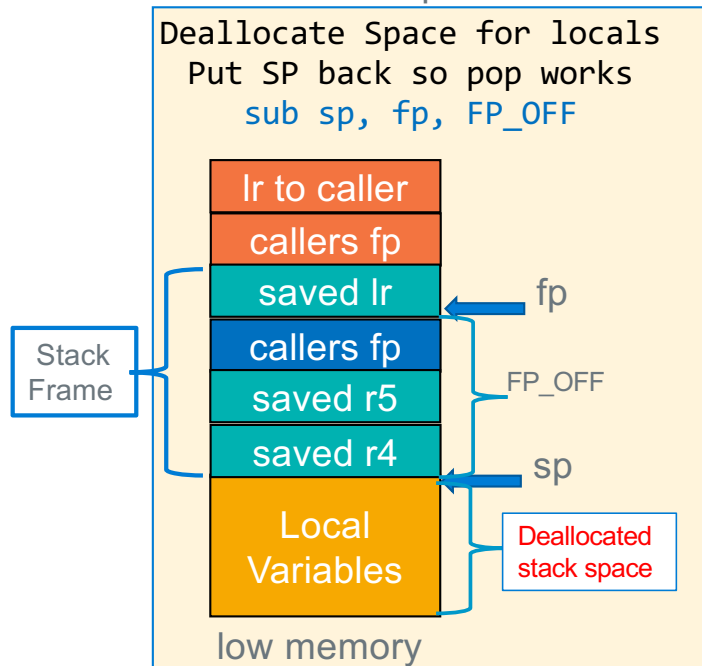
function
Epilogue

```

sub    sp, fp, FP_OFF    // deallocate local variables by moving sp
pop    {fp, lr}          // pop (restore) fp and lr from stack
bx     lr                // return to caller
    
```

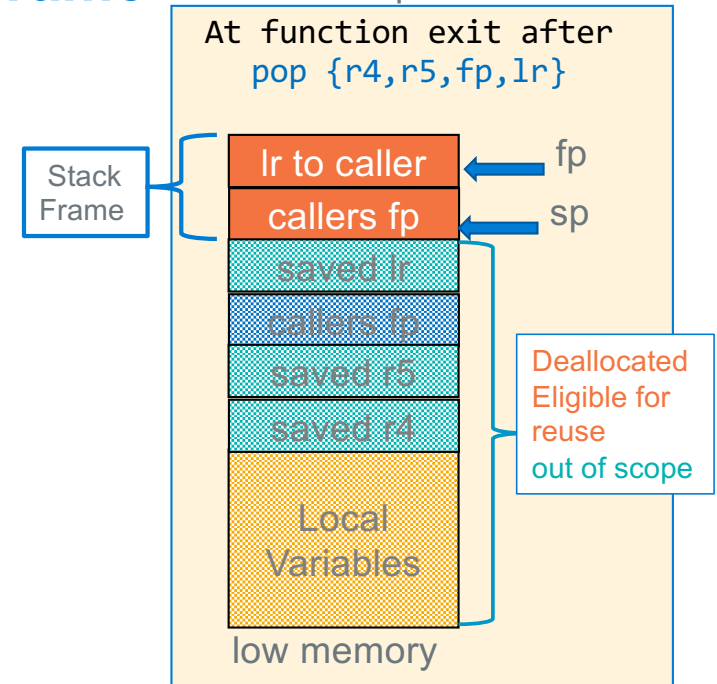
Function Epilogue: Deallocating the Stack Frame

Step 1



Move SP back to where it was after the push in the prologue.
So, pop works properly (this also deallocates the local variables)

Step 2



Use `pop` to restore the registers to the values they had at function entry

function
Epilogue

```
sub    sp, fp, FP_OFF    // deallocate local variables by moving sp
pop    {fp, lr}          // pop (restore) fp and lr from stack
bx     lr                // return to caller
```

How to Set FP

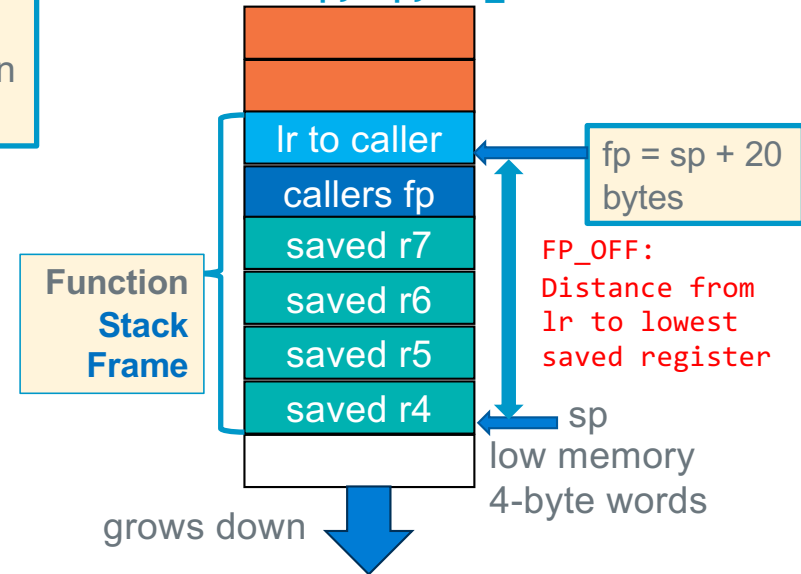
```
// other code etc
.equ    FP_OFF, 20

main:
  push   {r4-r7, fp, lr}
  add    fp, sp, FP_OFF
  .....
  sub    sp, fp, FP_OFF
  pop    {r4-r7, fp, lr}
  bx     lr
```

Function Prologue
always at top of function
saves regs and **sets fp**

Function Epilogue
always at bottom of function
restores regs including the sp

after push {r4-r7, fp, lr}
add fp, sp, FP_OFF



$$FP_OFF = (\#regs\ saved - 1) * 4$$



Means Caution, odd number of saved regs!

If odd number pushed, make sure frame is 8-byte aligned (later)

this must always be true: **sp % 8 == 0**

# regs saved	FP_OFF in Bytes Distance from lr to lowest saved register
2	4
3	8
4	12
5	16
6	20
7	24
8	28
9	32

Reference Table: Global Variable access

var	global variable address into r0 (lside)	global variable contents into r0 (rside)	contents of r0 into global variable
x	ldr r0, =x	ldr r0, =x ldr r0, [r0]	ldr r1, =x str r0, [r1]
*x	ldr r0, =x ldr r0, [r0]	ldr r0, =x ldr r0, [r0] ldr r0, [r0]	ldr r1, =x ldr r1, [r1] str r0, [r1]
**x	ldr r0, =x ldr r0, [r0] ldr r0, [r0]	ldr r0, =x ldr r0, [r0] ldr r0, [r0] ldr r0, [r0]	ldr r1, =x ldr r1, [r1] ldr r1, [r1] str r0, [r1]
stderr	ldr r0, =stderr	ldr r0, =stderr ldr r0, [r0]	<do not write unless you really know what you are doing>
.Lstr	ldr r0, =.Lstr	ldr r0, =.Lstr ldrb r0, [r0]	<read only>

```
.bss // from libc
stderr:.space 4 // FILE *
```

```
.data
x: .data y //x = &y
```

```
.section .rodata
.Lstr: .string "HI\n"
```

stdin, stdout and stderr are global variables

Assembler Directives: Label Scope Control (Normal Labels only)

```
.extern printf
.extern fgets
.extern strcpy
.global fbuf
```

.extern <label>

- **Imports** label (function name, symbol or a static variable name);
- An address associated with the label from another file can be used by code in this file

.global <label>

- **Exports** label (or symbol) to be visible outside the source file boundary (other assembly or c source)
 - label is either a function name or a global variable name
 - Only use with function names or static variables
- **Without** .global, labels are usually (depends on the assembler) **local to the file**

Passing global variables as a parameter: fprintf()

- `r0 = function(r0, r1, r2, r3)`
`fprintf(stderr, "arg2", arg3, arg4)`
- create a literal string for `arg2` which tells `fprintf()` how to interpret the remaining arguments
- `stdin`, `stdout`, `stderr` are all **global variable** and are **part of libc**
 - these **names are their lside (label names)**
 - **get their contents** and pass that to `fprintf()`, `fread()`, `fwrite()`

```
#include <stdio.h>
#include <stdlib.h>
int
main(void)
{
```

```
    int a = 2;
    int b = 3;
    int c;
```

We are going to
put these
variables in
temporary
registers

```
    c = a + b;
    fprintf(stderr, "c=%d\n", c);
```

`r0, r1, r2`

```
    return EXIT_SUCCESS;
}
```

three passed
args in this
use of `fprintf`

```
.extern fprintf           //declare fprintf
.section .rodata          // note the dots "."
.Lfst: .string "c=%d\n"
```

// part of the **text segment** below

```
mov     r2, 2             // int a = 2;
mov     r3, 3             // int b = 3;
add     r2, r2, r3        // arg 3: int c = a + b;

ldr     r0, =stderr       // get stderr address
ldr     r0, [r0]          // arg 1: get stderr contents
ldr     r1, =.Lfst        // arg 2: =literal address
bl      fprintf
```


Example: using preserved registers for local variables

```
#include <stdio.h>
#include <stdlib.h>
int
main(void)
{
```

```
    int c; // use r0
    int count = 0; // use r4
```

r0

```
    while ((c = getchar()) != EOF) {
```

```
        putchar(c);
        count++;
```

r0

r0

r1

```
    }
    printf("Echo count: %d\n", count);
    return EXIT_SUCCESS;
}
```

You must assume that both getchar() and putchar() alter r0-r3

Push two registers to keep stack 8-byte aligned (sp % 8 == 0)

```
.extern getchar
.extern putchar
.section .rodata
.Lst: .string "Echo count: %d\n"
```

```
.text
.type main, %function
.global main
.equ EOF, -1
.equ FP_OFF, 12
.equ EXIT_SUCCESS, 0
```

main:

```
push    {r4, r5, fp, lr}
add     fp, sp, FP_OFF
mov     r4, 0 //r4 = count
```

/* while loop code will go here */

```
mov     r0, EXIT_SUCCESS
sub     sp, fp, FP_OFF
pop     {r4, r5, fp, lr}
bx      lr
.size main, (. - main)
```

Putchar/getchar: The while loop

```
#include <stdio.h>
#include <stdlib.h>
int
main(void)
{
    int c;
    int count = 0;

    while ((c = getchar()) != EOF) {
        putchar(c);
        count++;
    }
    printf("Echo count: %d\n", count);
    return EXIT_SUCCESS;
}
```

initialize count

pre loop test with a call to getchar()
if it returns EOF in r0 we are done

echo the character read with getchar and
then read another and increment count

did getchar() return EOF if not loop

saw EOF, print count

```
mov    r4, 0    //count
bl     getchar
cmp    r0, EOF
beq    .Ldone
```

.Lloop:

```
bl     putchar
bl     getchar
add    r4, r4, 1
cmp    r0, EOF
bne    .Lloop
```

.Ldone:

```
mov    r1, r4    //arg2
ldr    r0, =.Lst //arg1
bl     printf
```

address of string literal variable

.Lst: .string "Echo count: %d\n"

File header and footers are not shown

Accessing Pointers (argv) in ARM assembly

```
.extern printf
.extern stderr
.section .rodata
.Lstr: .string "argv[%d] = %s\n"
.text
.global main // main(r0=argc, r1=argv)
.type main, %function
.equ FP_OFF, 20
main:
push {r4-r7, fp, lr}
add fp, sp, FP_OFF
mov r7, r1 // save argv!
ldr r4, =stderr // get the address of stderr
ldr r4, [r4] // get the contents of stderr
ldr r5, =.Lstr // get the address of .Lstr
mov r6, 0 // set indx = 0;

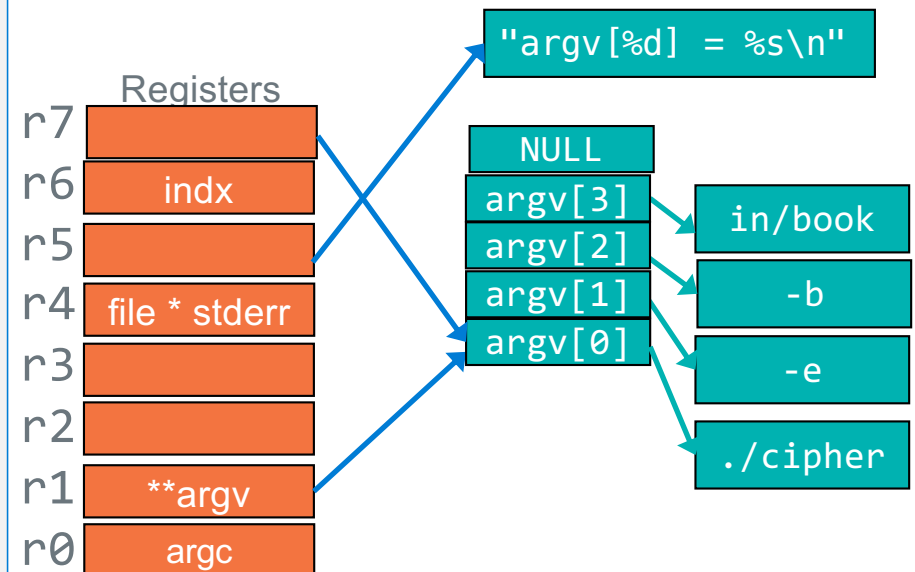
// see next slide

.Ldone:
mov r0, 0
sub sp, fp, FP_OFF
pop {r4-r7, fp, lr}
bx lr
```

need to save r1 as
we are calling a
function - fprintf

```
% ./cipher -e -b in/B00K
argv[0] = ./cipher
argv[1] = -e
argv[2] = -b
argv[3] = in/B00K
```

r0-r3 lost due to fprintf call



```
fprintf(stderr, "argv[%d] = %s\n", indx, *argv);
```

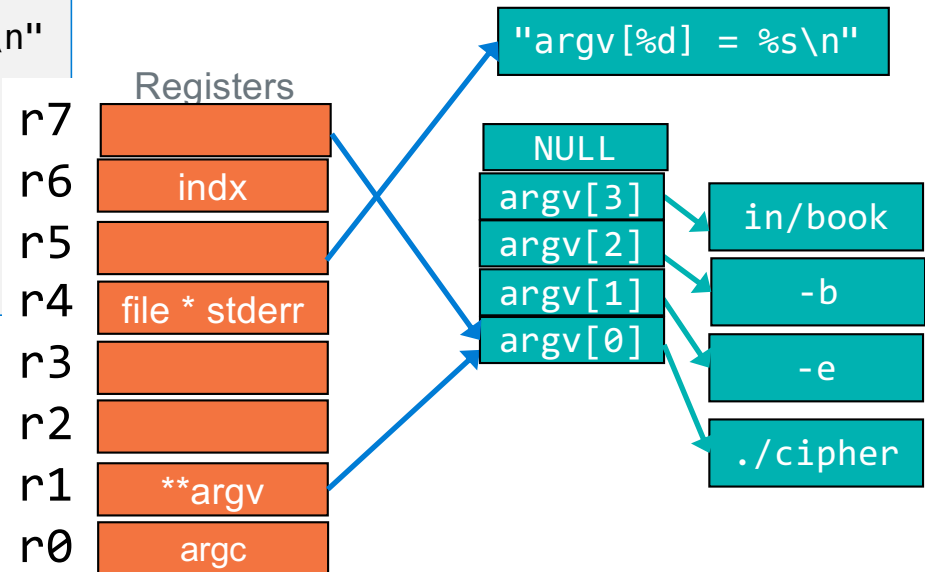
Accessing Pointers (argv) in ARM assembly

```
.Lloop:
    // fprintf(stderr, "argv[%d] = %s\n", indx, *argv)
    ldr    r3, [r7]           // arg 4: *argv
    cmp    r3, 0              // check *argv == NULL
    beq    .Ldone             // if so done
    mov     r2, r6             // arg 3: indx
    mov     r1, r5             // arg 2: "argv[%d] = %s\n"
    mov     r0, r4             // arg 1: stderr
    bl     fprintf
    add     r6, r6, 1          // indx++ for printing
    add     r7, r7, 4          // argv++ pointer
    b      .Lloop
.Ldone:
```

observe the
different
increment sizes

```
% ./cipher -e -b in/BOOK
argv[0] = ./cipher
argv[1] = -e
argv[2] = -b
argv[3] = in/BOOK
```

r0-r3 lost due to fprintf call

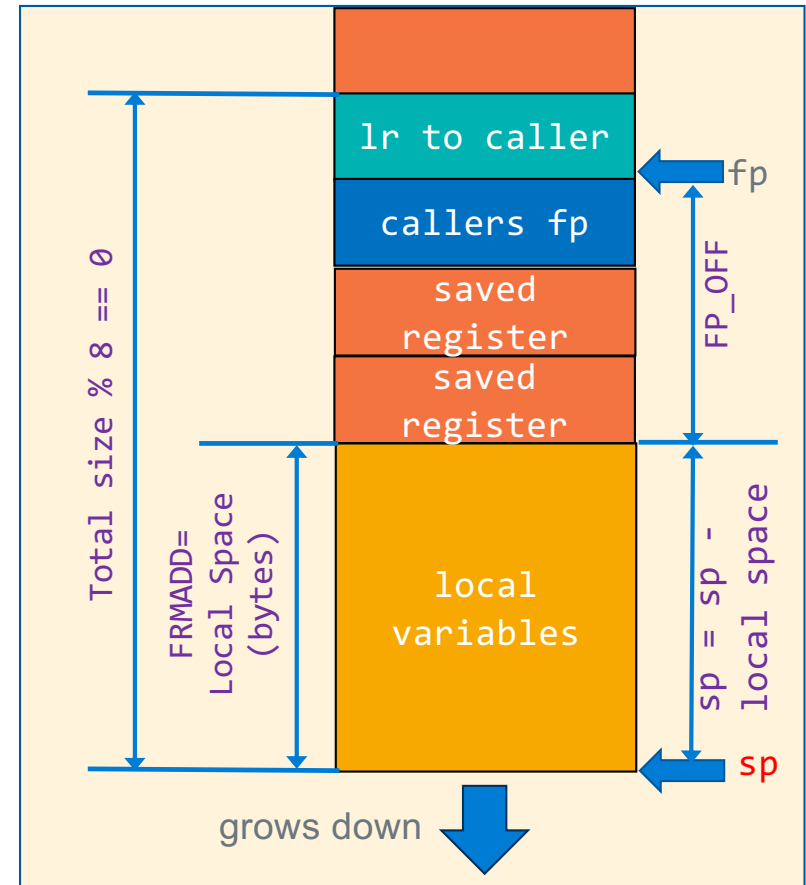


Allocating Space For Locals on the Stack

- Space for local variables is allocated on the stack right below the lowest pushed register
 - Move the **sp** towards low memory by the total size of all local variables in bytes **plus padding**

$FRMADD = \text{total local var space (bytes)} + \text{padding}$

- Allocate the space after the register push by
`add sp, sp, -FRMADD`
- Requirement:** on function entry, **sp** is always 8-byte aligned
 $sp \% 8 == 0$
- Padding (as required):**
 - Additional space between variables on the stack to meet memory alignment requirements
 - Additional space so the frame size is evenly divisible by 8
- fp** (frame pointer) is used as a **pointer (base register)** to access all stack variables – later slides



Review Variables: Size

- **Integer types**

- `char` (unspecified default)
- `int` (signed default)

- **Floating Point**

- `float`, `double`

- Optional Modifiers for each base type

- `short` [int]
- `long` [int, double]
- `signed` [char, int]
- `unsigned` [char, int]
- `const`: variable read only

- **char type**

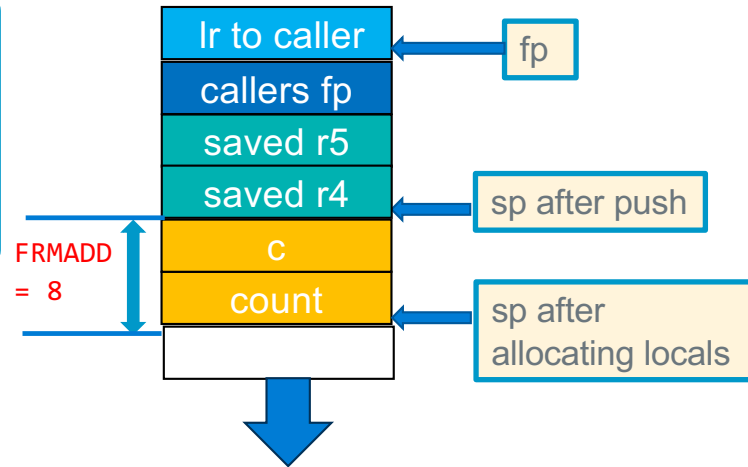
- One byte in a byte addressable memory
- **Be careful** `char` is unsigned on arm and signed on other HW like intel

C Data Type	AArch-32 contiguous Bytes	printf specification
unsigned char	1	%c
signed char	1	%c
short int	2	%hd
unsigned short int	2	%hu
int	4	%d / %i
unsigned int	4	%u
long int	4	%ld
long long int	8	%lld
float	4	%f
double	8	%lf
long double	8	%Lf
pointer *	4	%p

Local Variables on the stack

after push {r4-r5,fp,lr}
add fp, sp, FP_OFF

```
int main(void)
{
    int c;
    int count = 0;
    // rest of code
}
```



```
.text
.type    main, %function
.global  main
.equ     FP_OFF,    12
.equ     FRMADD,    8
main:
    push    {r4, r5, fp, lr}
    add     fp, sp, FP_OFF
    add     sp, sp, -FRMADD
    // but we are not done yet!
```

```
// when FRMADD values fail to assemble
ldr r3, =-FRMADD
add sp, sp, r3
```

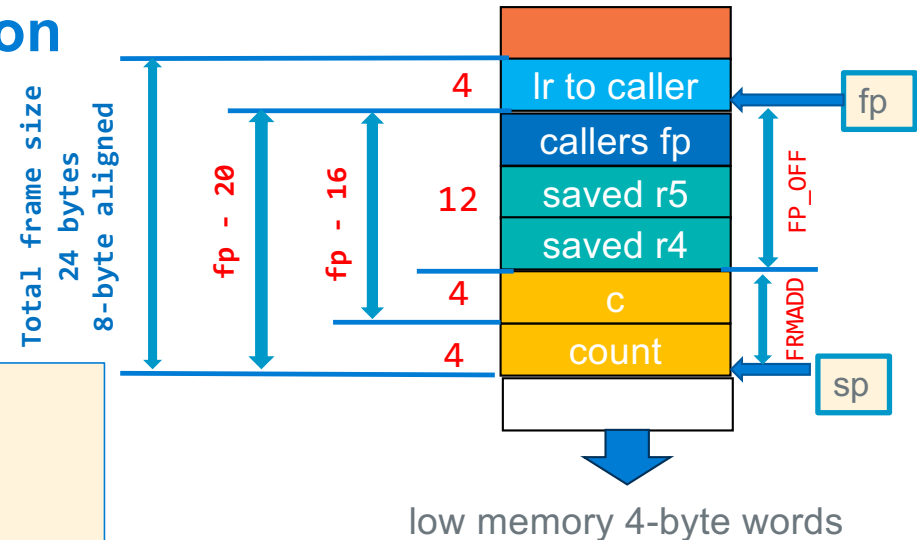
- In this example we are **allocating two variables on the stack**
- When writing assembly functions, in many situations **you may choose allocate these to registers instead**

- Add space on the stack for each local
 - we will allocate space in same order the locals are listed the C function shown from high to low stack address
 - gcc compiler allocates from low to high stack addresses
 - Order does not matter for our use

Accessing Stack Variables: Introduction

```
int main(void)
{
    int c;
    int count = 0;
    // rest of code
}
```

- To Access data stored in the stack
 - use the `ldr/str` instructions
- Use register `fp` with offset (**distance in bytes**) addressing (use either register offset or immediate offset)
- No matter what address the stack frame is at**, `fp` always points at saved `lr`, so you can find a local stack variable by using an offset address from the contents of `fp`

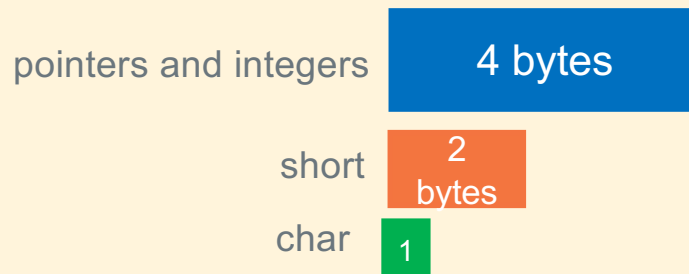


```
.text
.type    main, %function
.global  main
.equ     FP_OFF,    12
.equ     FRMADD,    8
main:
    push    {r4, r5, fp, lr}
    add     fp, sp, FP_OFF
    add     sp, sp, -FRMADD
    // but we are not done yet!
```

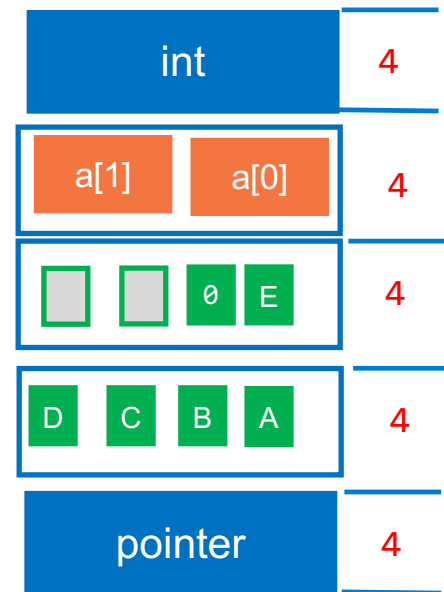
Variable	distance from fp	Read variable	Write Variable
int c	-16	ldr r0, [fp, -16]	str r0, [fp, -16]
int count	-20	ldr r0, [fp, -20]	str r0, [fp, -20]

Stack Frame Design – Local Variables

- When writing an ARM equivalent for a C program, for CSE30 we will not re-arrange the order of the variables to optimize space (covered in the compiler course)
- **Arrays** start at a 4-byte boundary (even arrays with only 1 element)
 - Exception: double arrays [] start at an 8-byte boundary
 - **struct** arrays are **aligned to the requirements of largest member**
- Single chars (and shorts) can be grouped together in same 4-byte word (following the alignment for the short)
- Padding may be required (see next slide)

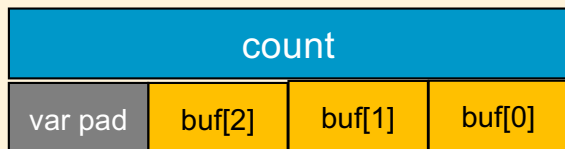


Rule: When the function is entered the stack is already 8-byte aligned

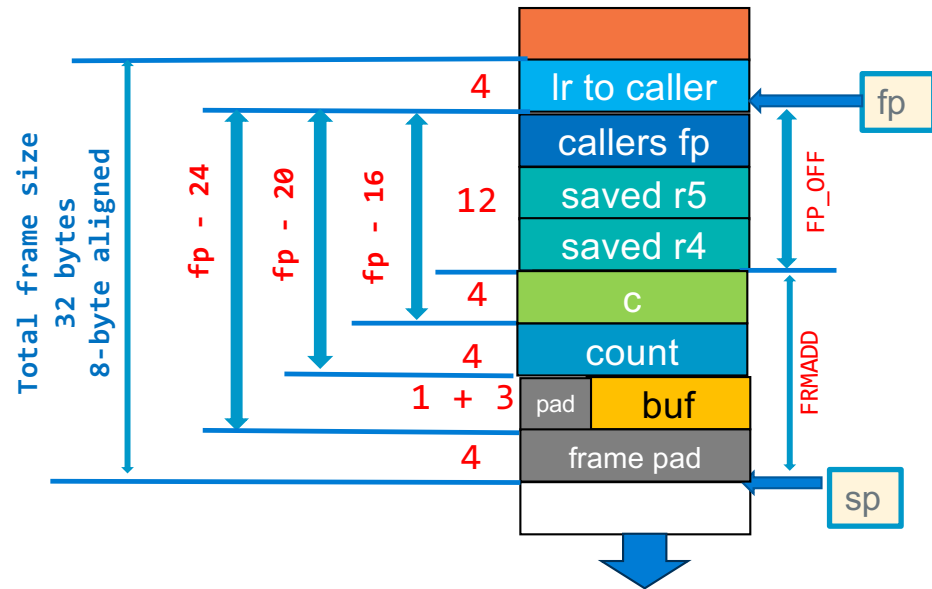
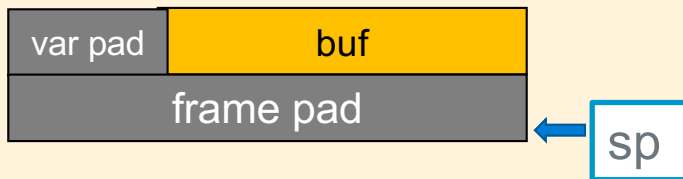


Stack Variables: Padding

- **Variable padding** – start arrays at 4-byte boundary and **leave unused space at end** (high side address) before the variable higher on the stack



- **Frame padding** – add space below the last local variable to keep 8-byte alignment



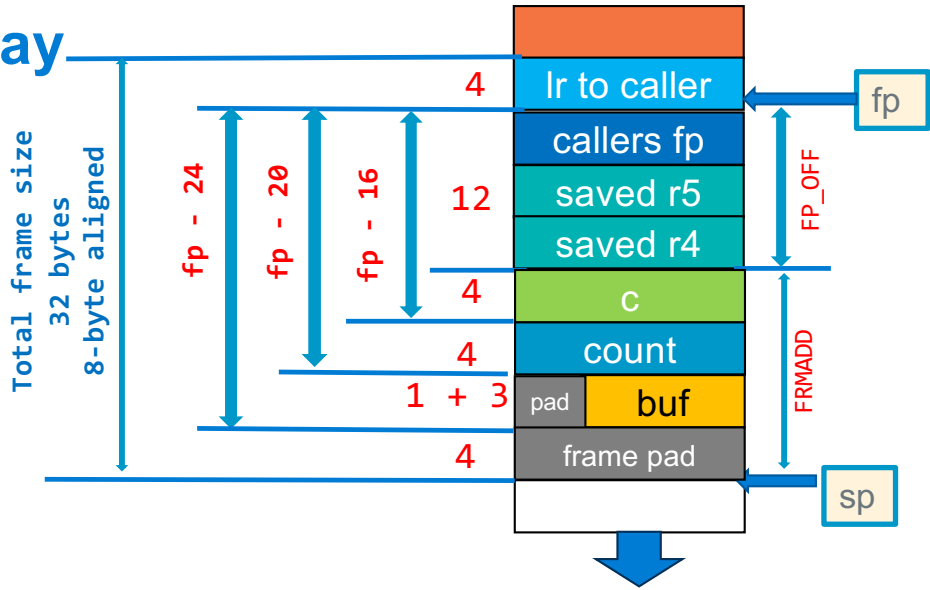
```
int main(void)
{
    int c;
    int count = 0;
    char buf[] = "hi";
    // rest of code
}
```

```
.text
.type    main, %function
.global main
.equ     FP_OFF,    12
.equ     FRMADD,    16
main:
    push    {r4, r5, fp, lr}
    add     fp, sp, FP_OFF
    add     sp, sp, -FRMADD
    // but we are not done yet!
```

Accessing Stack Variables, the hard way

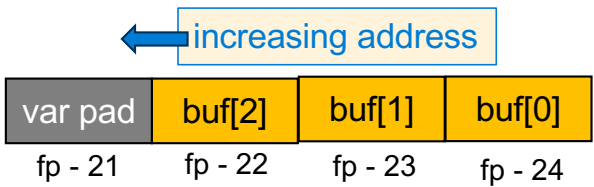
```
int main(void)
{
    int c;
    int count = 0;
    char buf[] = "hi";
    // rest of code
}
```

```
.text
.type    main, %function
.global  main
.equ     FP_OFF,    12
.equ     FRMADD,    16
main:
    push    {r4, r5, fp, lr}
    add     fp, sp, FP_OFF
    add     sp, sp, -FRMADD
    // but we are not done yet!
```



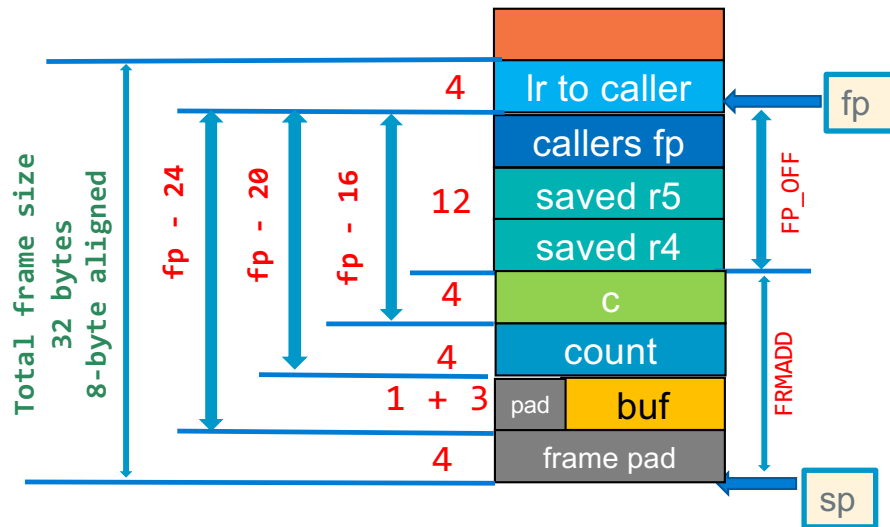
char buf[] by usage with ASCII chars we will use strb (or make it unsigned char)

Variable	distance from fp	Read variable	Write Variable
int c	16	ldr r0, [fp, -16]	str r0, [fp, -16]
int count	20	ldr r0, [fp, -20]	str r0, [fp, -20]
char buf[0]	24	ldrb r0, [fp, -24]	strb r0, [fp, -24]
char buf[1]	23	ldrb r0, [fp, -23]	strb r0, [fp, -23]
char buf[2]	22	ldrb r0, [fp, -22]	strb r0, [fp, -22]



- Calculating offsets is a lot of work to get it correct
- It is also hard to debug
- There is a better way!

Best Practice: Assembler Generated FP Distance Table



FP Distance Table one For each function

```
.type    main, %function
.global  main
.equ     FP_OFF, 12

.equ     C, 4 + FP_OFF
.equ     COUNT, 4 + C
.equ     BUF, 4 + COUNT
.equ     PAD, 4 + BUF
.equ     FRMADD, PAD - FP_OFF
// FRMADD = 28 - 12 = 16
```

Annotations in the diagram:

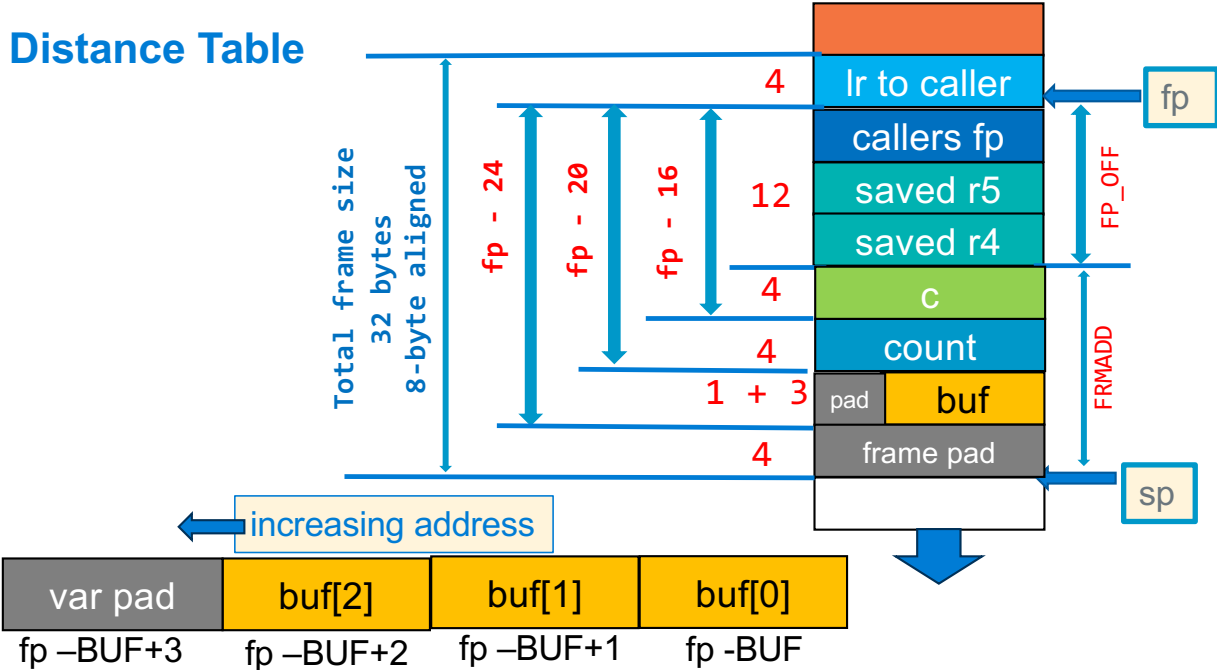
- pushed reg fp distance**: Points to the `FP_OFF` definition.
- variable size in bytes**: Points to the `4` in the `C` definition.
- Prior allocation distance**: Points to the `FP_OFF` definition.

- For each **stack variable** create a **.equ symbol** whose **value** is the **distance in bytes from the FP** after the **prologue**
- After the last variable add a name **PAD** for the size of the frame padding (if any). if no padding, PAD will be set to the same value as the variable above it
- The **value of the symbol** is an **expression that calculates the distance from the FP** based on the distance of the variable above it on the stack. The first variable will use **SP_OFF** as the starting distance
.equ VAR, size_of var + variable_padding + previous_var_symbol *// previous_var_symbol distance of the var above*
- Calculate the size of the local variable area that needs to be added to the sp in bytes
FRMADD = distance PAD minus distance of the SP to the FP (FP_OFF) after the prologue push

Best Practice: Assembler Generated FP Distance Table

FP Distance Table For each function

```
.type    main, %function
.global main
.equ     FP_OFF,      12
.equ     C,           4 + FP_OFF
.equ     COUNT,       4 + C
.equ     BUF,         4 + COUNT
.equ     PAD,         4 + BUF
.equ     FRMADD,      PAD - FP_OFF
// FRMADD = 28 - 12 = 16
```



Variable	distance from fp	Address on Stack	Read variable	Write Variable
int c	C	add r0, fp, -C	ldr r0, [fp, -C]	str r0, [fp, -C]
int count	COUNT	add r0, fp, -COUNT	ldr r0, [fp, -COUNT]	str r0, [fp, -COUNT]
char buf[0]	BUF	add r0, fp, -BUF	ldrb r0, [fp, -BUF]	strb r0, [fp, -BUF]
char buf[1]	BUF-1	add r0, fp, -BUF+1	ldrb r0, [fp, -BUF+1]	strb r0, [fp, -BUF+1]
char buf[2]	BUF-2	add r0, fp, -BUF+2	ldrb r0, [fp, -BUF+2]	strb r0, [fp, -BUF+2]

Initializing and Accessing Stack variables

```
.section .rodata
.Lmess: .string "%d %d %s\n"
.extern printf
```

```
main:
    push    {r4, r5, fp, lr}
    add     fp, sp, FP_OFF
    add     sp, sp, -FRMADD
    // nothing to do for C
    mov     r2, 0
    str     r2, [fp, -COUNT]
    strb    r2, [fp, -BUF+2]
    mov     r2, 'h'
    strb    r2, [fp, -BUF]
    mov     r2, 'i'
    strb    r2, [fp, -BUF+1]

    ldr     r0, =.Lmess // arg1
    ldr     r1, [fp, -C] // arg2
    ldr     r2, [fp, -COUNT] // arg3
    add     r3, fp, -BUF // arg4
    bl      printf
```

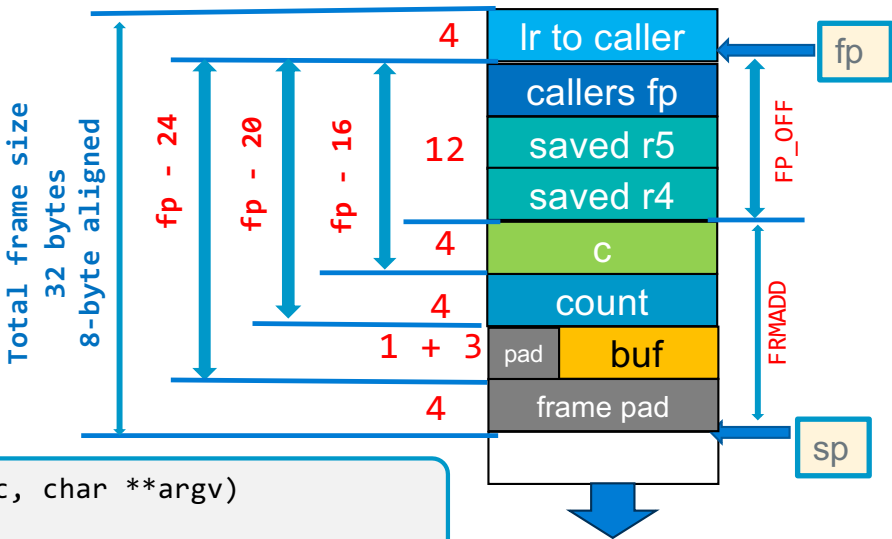
passes contents of stack var C and COUNT

passes address of a stack variable buf

```
int main(int argc, char **argv)
{
    int c;
    int count = 0;
    char buf[] = "hi";
    printf("%d %d %s\n", c, count, buf);
    // rest of code
```

pass stack address

./a.out
-136572160 0 hi



Variable	distance from fp	Address on Stack	Read variable	Write Variable
int c	C	add r0, fp, -C	ldr r0, [fp, -C]	str r0, [fp, -C]
int count	COUNT	add r0, fp, -COUNT	ldr r0, [fp, -COUNT]	str r0, [fp, -COUNT]
char buf[0]	BUF	add r0, fp, -BUF	ldrb r0, [fp, -BUF]	strb r0, [fp, -BUF]
char buf[1]	BUF-1	add r0, fp, -BUF+1	ldrb r0, [fp, -BUF+1]	strb r0, [fp, -BUF+1]
char buf[2]	BUF-2	add r0, fp, -BUF+2	ldrb r0, [fp, -BUF+2]	strb r0, [fp, -BUF+2]

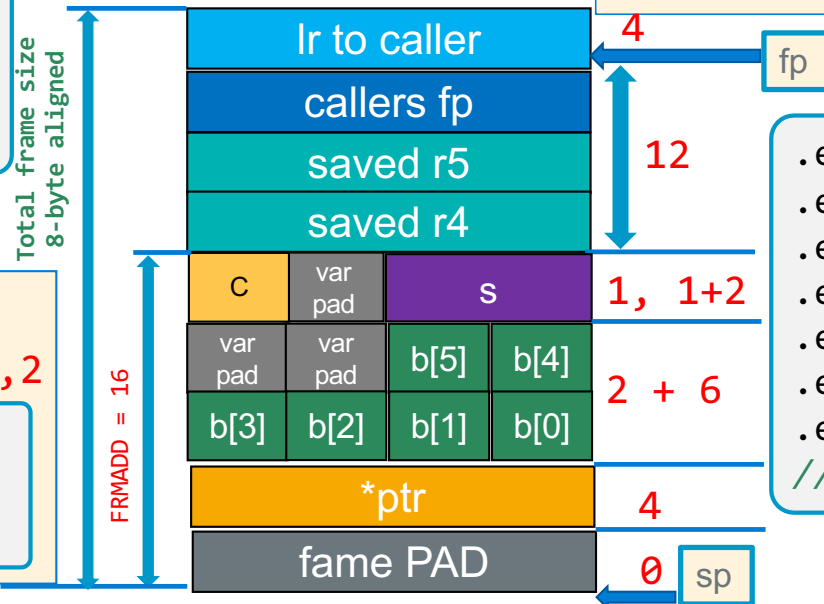
Stack Frame Design Practice

```
void func(void)
{
    signed char c;
    signed short s;
    unsigned char b[] = "Stack";
    unsigned char *ptr = &b;
    // rest of code
}
```

Alternative design



```
.equ  FP_OFF,    12
.equ   C,        2 + FP_OFF
.equ   S,        2 + C
...
```



1. Write the variables in C
2. Draw a picture of the stack frame
3. Write the code to generate the offsets
4. create the distance table to the variables

```
.equ  FP_OFF,    12
.equ   C,        1 + FP_OFF
.equ   S,        3 + C
.equ   B,        8 + S
.equ   PTR,      4 + B
.equ   PAD,      0 + PTR
.equ   FRMADD,   PAD - FP_OFF
// FRMADD = 28 - 12 = 16
```

Variable	distance from fp	Address on Stack	Read variable	Write Variable
signed char c	C	add r0, fp, -C	ldrsb r0, [fp, -C]	strsb r0, [fp, -C]
signed short s	S	add r0, fp, -S	ldrsh r0, [fp, -S]	strsh r0, [fp, -S]
unsigned char b[0]	B	add r0, fp, -B	ldrb r0, [fp, -B]	strb r0, [fp, -B]
unsigned char *ptr	PTR	add r0, fp, -PTR	ldr r0, [fp, -PTR]	str r0, [fp, -PTR]

Working with Pointers on the stack

```
int sum(int j, int k)
{
    return j + k;
}
void testp(int j, int k, int (*func)(int, int), int *i)
{
    *i = func(j,k);
    return;
}
int main()
{
    int i;                                // NOTICE: i must be on stack as you pass the address!
    int (*pf)(int, int) = sum;           // pf could be in a register

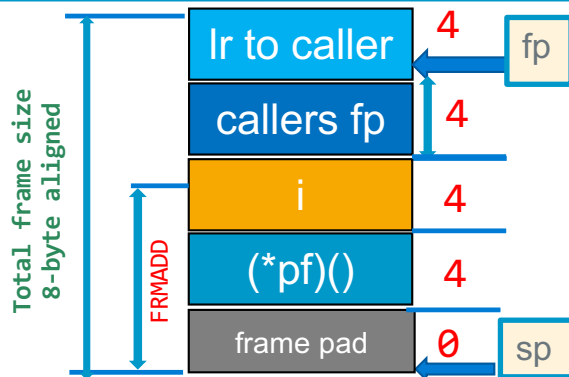
    testp(1, 2, pf, &i);
    printf("%d\n", i);
    return EXIT_SUCCESS;
}
```

Output Parameters (like i) you pass a pointer to them, **must be on the stack!**

Working with Pointers on the stack

```
int main()
{
    int i; // NOTICE: i must be on stack as you pass the address!
    int (*pf)(int, int) = sum; // pf could be in a register

    testp(1, 2, pf, &i);
    printf("%d\n", i);
    return EXIT_SUCCESS;
}
```



```
.section .rodata
.Lmess: .string "%d\n"
.extern printf
.text
.global main
.type    main, %function
.equ     FP_OFF, 4
.equ     I,      4 + FP_OFF
.equ     PF,     4 + I
.equ     PAD,    0 + PF
.equ     FRMADD, PAD - FP_OFF
// FRMADD = 12 - 4 = 8
```

Variable	distance from fp	Address on Stack	Read variable	Write Variable
int i	I	add r0, fp, -I	ldr r0, [fp, -I]	str r0, [fp, -I]
int (*pf)()	PF	add r0, fp, -PF	ldr r0, [fp, -PF]	str r0, [fp, -PF]

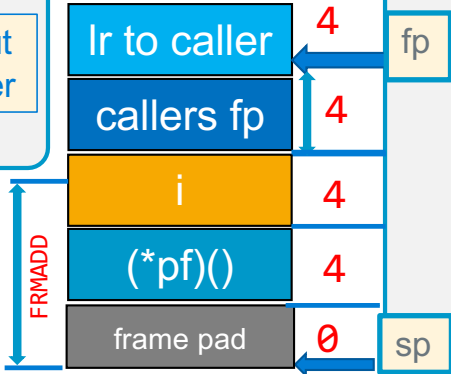
Working with Pointers on the stack

```
int main()
{
    int i;
    int (*pf)(int, int) = sum;

    testp(1, 2, pf, &i);
    printf("%d\n", i);
    return EXIT_SUCCESS;
}
```

I is Output Parameter

```
.section .rodata
.Lmess: .string "%d\n"
.extern printf
.text
.global main
.type main, %function
.equ FP_OFF, 4
.equ I, 4 + FP_OFF
.equ PF, 4 + I
.equ PAD, 0 + PF
.equ FRMADD, PAD-FP_OFF
// FRMADD = 12 - 4 = 8
```



```
main:
    push    {fp, lr}
    add     fp, sp, FP_OFF
    add     sp, sp, -FRMADD

    ldr     r2, =sum           // func address
    add     r1, fp, -PF        // PF address
    str     r2, [r1]           // store in pf

    mov     r0, 1              // arg 1: 1
    mov     r1, 2              // arg 2: 2
    ldr     r2, [fp, -PF]      // arg 3: (*pf)()
    add     r3, fp, -I         // arg 4: &I
    bl      testp

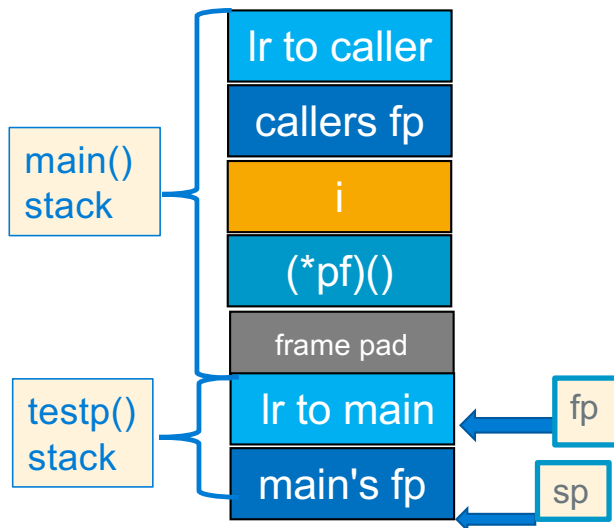
    ldr     r0, =.Lmess        // arg 1: "%d\n"
    ldr     r1, [fp, -I]       // arg 2: I
    bl      printf

    sub     sp, fp, FP_OFF
    pop     {fp, lr}
    bx      lr
```

Variable	distance from fp	Address on Stack	Read variable	Write Variable
int i	I	add r0, fp, -I	ldr r0, [fp, -I]	str r0, [fp, -I]
int (*pf)()	PF	add r0, fp, -PF	ldr r0, [fp, -PF]	str r0, [fp, -PF]

Working with Pointers on the stack

```
void
testp(int j, int k, int (*func)(int, int), int *i)
{
    *i = func(j, k);
    return;
}
```



`r0,r1,r2` already set

```
.global testp
.type testp, %function
.equ FP_OFF, 12
testp:
    push    {r4, r5, fp, lr}
    add     fp, sp, FP_OFF

    mov     r4, r3           // save i
    blx     r2               // r0=func(r0,r1)
    str     r0, [r4]         // *i = r0

    sub     sp, fp, FP_OFF
    pop     {r4, r5, fp, lr}
    bx     lr
.size testp, (. - testp)
```

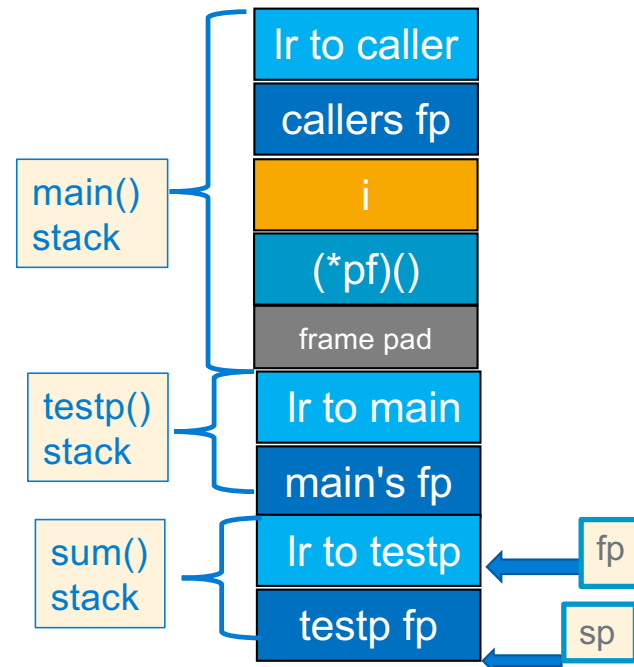
Working with Pointers on the stack

```
int
sum(int j, int k)
{
    return j + k;
}
```

```
.global sum
.type    sum, %function
.equ    FP_0FF, 4
sum:
    push    {fp, lr}
    add     fp, sp, FP_0FF

    add     r0, r0, r1

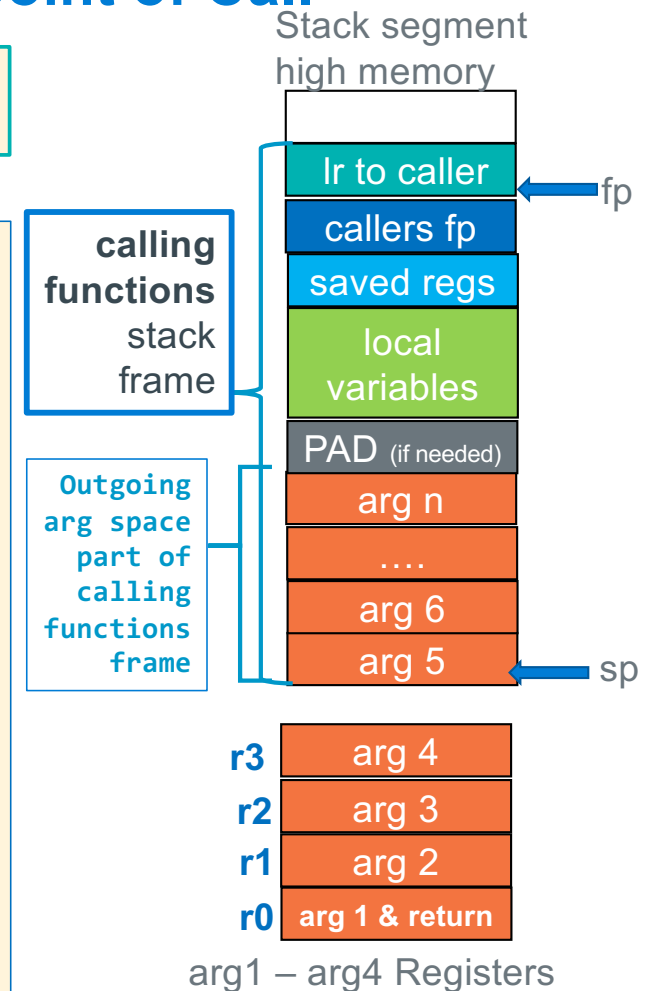
    sub     sp, fp, FP_0FF
    pop     {fp, lr}
    bx     lr
.size sum, (. - sum)
```



Passing More Than Four Arguments – At the point of Call

```
r0 = function(r0, r1, r2, r3, arg5, arg6, ... argn)
      arg1, arg2, arg3, arg4, ...
```

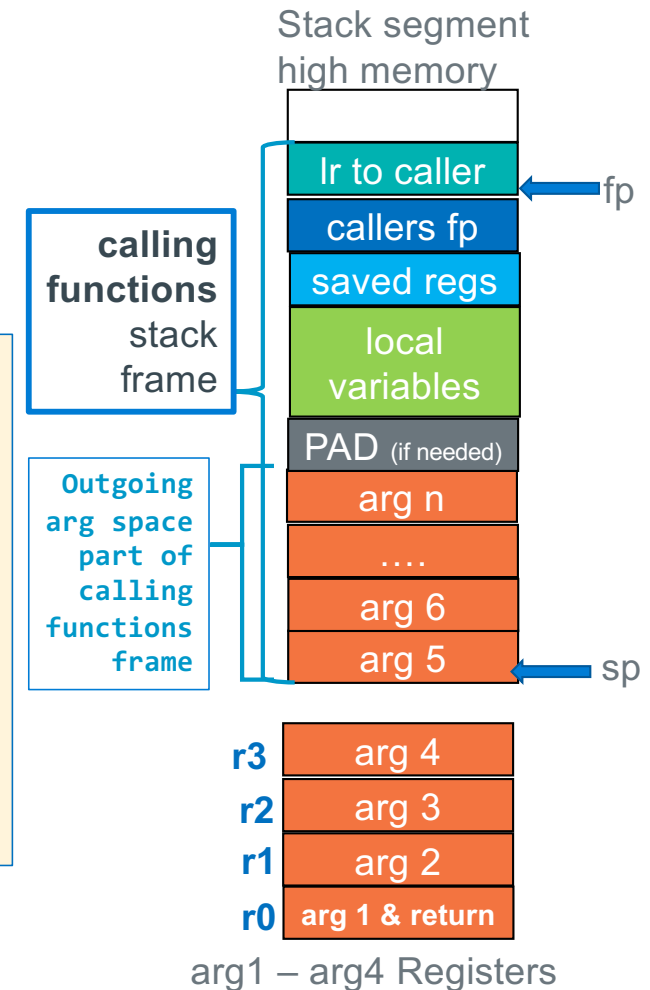
- **Approach: Increase stack frame size to include space for args# > 4**
 - Arg5 and above are in caller's stack frame at the **bottom of the stack**
- **Arg5** is always at the **bottom (at sp)**, arg6 and greater are above it
- **One arg value per slot!** – NO arrays across multiple slots
 - chars, shorts and ints are directly stored
 - Structs (not always), and arrays (always) are passed via a pointer
- **Output parameters** contain an **address that points at** the **stack**, **BSS**, **data**, or **heap**
- Prior to any function call (and obviously at the start of the called function):
 1. sp must point at arg5
 2. sp and therefore **arg5** must be at an 8-byte boundary,
 3. **Add padding** to force arg5 alignment if needed is **placed above** the last **argument the called function is expecting**



Passing More Than Four Arguments – At the point of Call

```
r0 = function(r0, r1, r2, r3, arg5, arg6, ... argn)
      arg1, arg2, arg3, arg4, ...
```

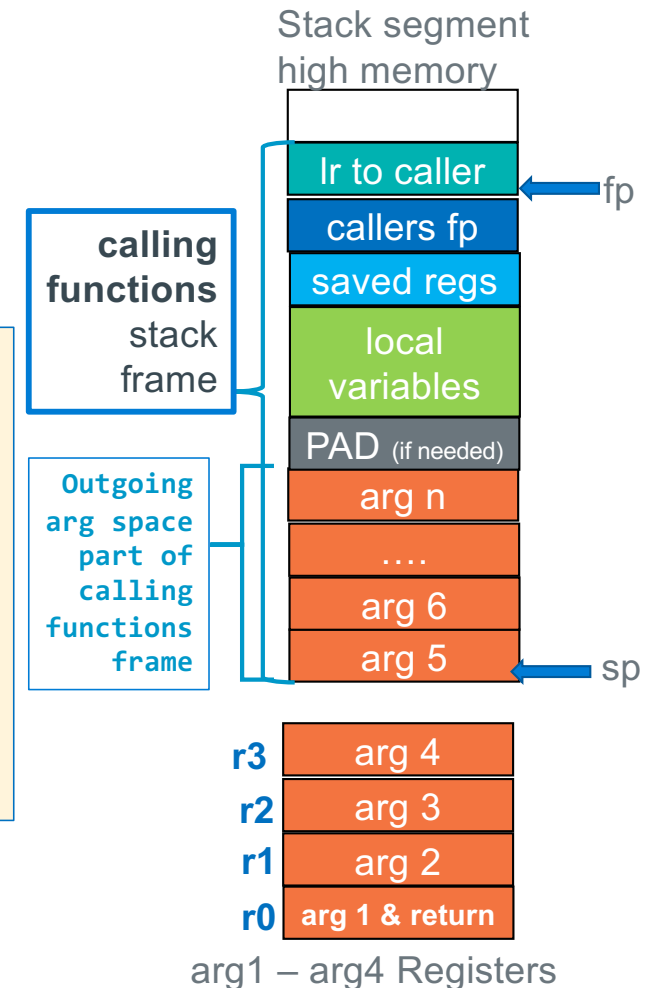
- **Called functions** have the **right to change stack args** just like they can change the register args!
 - Caller **must always assume all args including ones on the stack are changed by the caller**
- Calling function prior to making the call you must
 1. Evaluate **first four args**: place the resulting **values in r0-r3**
 2. Evaluate Arg 5 and greater and place the resulting values on the stack



Passing More Than Four Arguments – At the point of Call

```
r0 = function(r0, r1, r2, r3, arg5, arg6, ... argn)
      arg1, arg2, arg3, arg4, ...
```

- **Approach:** Extend the stack frame to include enough space for stack arguments for the called function that has the greatest number of args
 1. Examine every function call in the body of a function
 2. Find the function call with greatest arg count, this determines space needed for outgoing args
 3. Add the greatest arg count space as needed to the frame layout
 4. Adjust PAD as required to keep the sp 8-byte aligned



Determining Size of the Passed Parameter Area on The Stack

- Find the function called by main with the largest number of parameters
- That function determines the size of the Passed Parameter allocation on the stack

```
int main(void)
{
    /* code not shown */
    a(g, h);

    /* code not shown */
    sixsum(a1, a2, a3, a4, a5, a6);

    /* code not shown */

    b(q, w, e, r);
    /* code not shown */
}
```

← largest arg count is 6
allocate space for $6 - 4 = 2$ arg slots

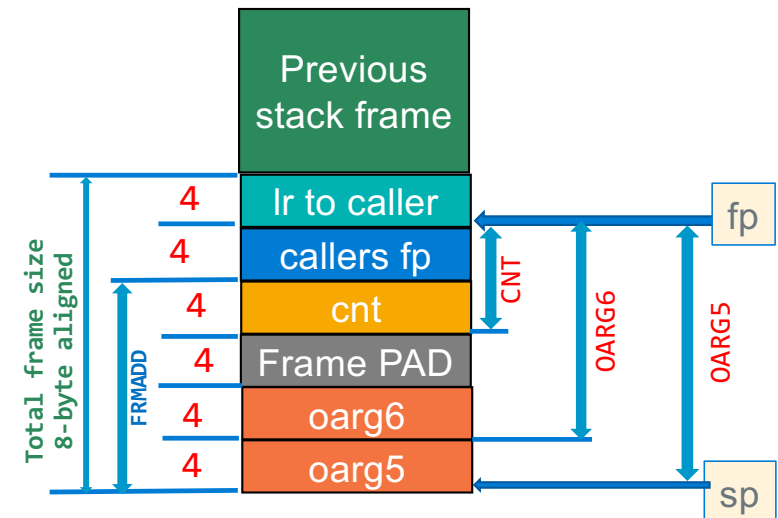
Calling Function Stack Frame: Pass ARG 5 and higher

Rules: At point of call

1. OARG5 must be pointed at by sp
2. SP must be 8-byte aligned at function call

```
int cnt;
r0 = func(r0, r1, r2, r3, OARG5, OARG6);
```

```
.equ  FP_OFF, 4
.equ  CNT,      4 + FP_OFF    // int cnt
.equ  PAD,      4 + CNT      // added as needed
.equ  OARG6,    4 + PAD      // 4 bytes
.equ  OARG5,    4 + OARG6    // 4 bytes
.equ  FRMADD,   OARG5 - FP_OFF
// FRMADD = 20 - 4 = 16
```



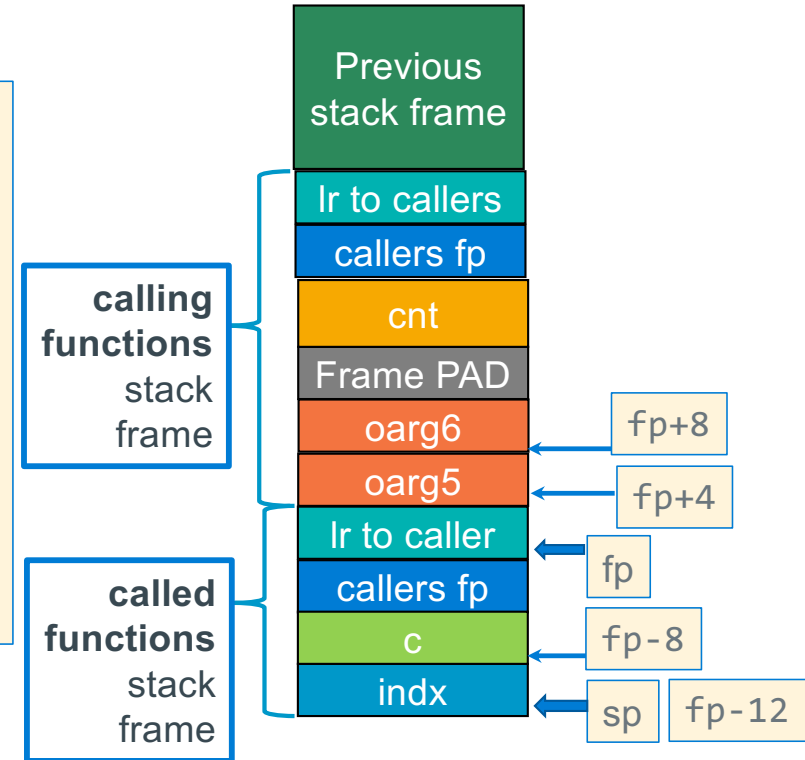
Variable	distance from fp	Address on Stack	Read variable	Write Variable
int cnt	CNT	add r0, fp, -CNT	ldr r0, [fp, -CNT]	str r0, [fp, -CNT]
int oarg6	OARG6	add r0, fp, -OARG6	ldr r0, [fp, -OARG6]	str r0, [fp, -OARG6]
int oarg5	OARG5	add r0, fp, -OARG5	ldr r0, [fp, -OARG5]	str r0, [fp, -OARG5]

Called Function: Retrieving Args From the Stack

```
r0 = func(r0, r1, r2, r3, ARG5, ARG6);
```

- At function start and before the push{} the **sp** is at an 8-byte boundary
 - Args > 4 in caller's stack frame** and **arg 5 always starts at fp+4**
 - Additional args are higher up the stack, with one "slot" every 4-bytes
- ```
.equ ARGN, (N-4)*4 // where n must be > 4
```
- This "algorithm" for finding args was designed to enable **variable arg count functions** like `printf("conversion list", arg0, ... argn);`
  - No limit to the number of args (except running out of stack space)

**Rule:**  
**Called functions** always access stack args using a **positive offset to the fp**

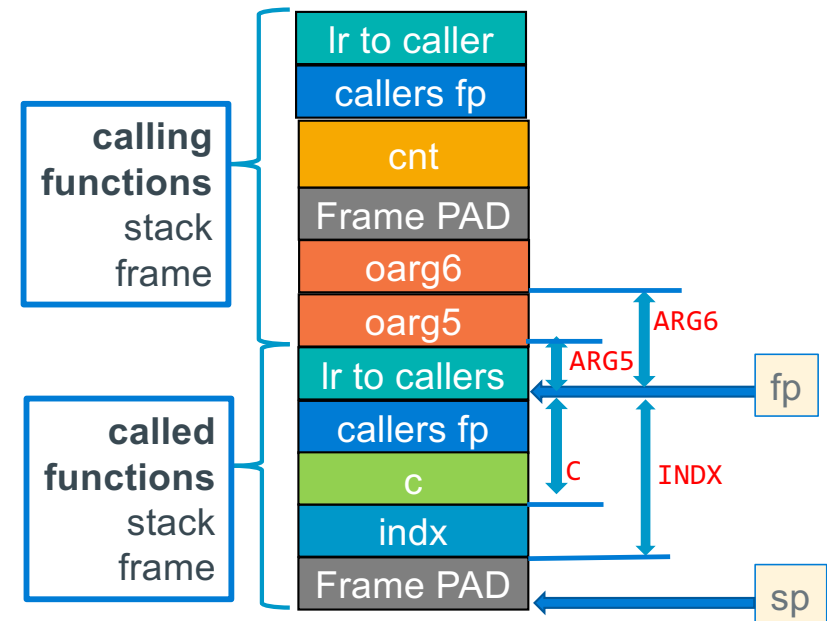


## Called Function: Retrieving Args From the Stack

```
.equ FP_OFF, 4
.equ C, 4 + FP_OFF
.equ INDX, 4 + C
.equ PAD, 0 + INDX
.equ FRMADD, PAD - FP_OFF
// below are distances into the caller's stack frame
.equ ARG6, 8
.equ ARG5, 4
```

```
r0 = func(r0, r1, r2, r3, r4, ARG5, ARG6);
```

**Rule:**  
**Called functions** always access stack args  
 using a **positive offset to the fp**



| Variable or Argument   | distance from fp | Address on Stack               | Read variable                    | Write Variable                   |
|------------------------|------------------|--------------------------------|----------------------------------|----------------------------------|
| <code>int arg6</code>  | ARG6             | <code>add r0, fp, ARG6</code>  | <code>ldr r0, [fp, ARG6]</code>  | <code>str r0, [fp, ARG6]</code>  |
| <code>int arg5</code>  | ARG5             | <code>add r0, fp, ARG5</code>  | <code>ldr r0, [fp, ARG5]</code>  | <code>str r0, [fp, ARG5]</code>  |
| <code>int c</code>     | C                | <code>add r0, fp, -C</code>    | <code>ldr r0, [fp, -C]</code>    | <code>str r0, [fp, -C]</code>    |
| <code>int count</code> | INDX             | <code>add r0, fp, -INDX</code> | <code>ldr r0, [fp, -INDX]</code> | <code>str r0, [fp, -INDX]</code> |

Observe the positive offsets

## Example: Passing Stack Args, Calling Function

```
int sum(int j, int k)
{
 return j + k;
}

void

arg1	arg2	arg3	arg4	arg5	arg6
------	------	------	------	------	------

testp(int j, int k, int l, int m, int (*func)(int, int), int *i)
{
 *i = func(j,k) + func(l, m); // notice two func() calls

 return;
}

int main()
{
 int i; // NOTICE: i must be on stack as you pass the address!
 int (*pf)(int, int) = sum; // pf could be in a register

 testp(1, 2, 3, 4, pf, &i);
 printf("%d\n", i);

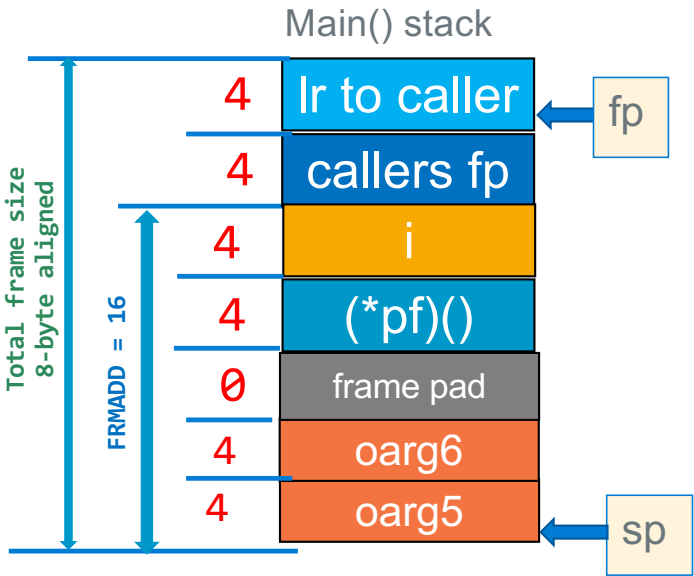
 return EXIT_SUCCESS;
}
```

# Example: Passing Stack Args, Calling Function

```
int main()
{
 int i; // NOTICE: i must be on stack as you pass the address!
 int (*pf)(int, int) = sum; // pf could be in a register

 testp(1, 2, 3, 4, pf, &i);
 printf("%d\n", i);
 return EXIT_SUCCESS;
}
```

```
.equ FP_OFF, 4
.equ I, 4 + FP_OFF
.equ PF, 4 + I
.equ PAD, 0 + PF
.equ OARG6, 4 + PAD
.equ OARG5, 4 + OARG6
.equ FRMADD, OARG5 - FP_OFF
// FRMADD = 20 - 4 = 16
```



| Variable or Argument | distance from fp | Address on Stack   | Read variable        | Write Variable       |
|----------------------|------------------|--------------------|----------------------|----------------------|
| int i                | I                | add r0, fp, -I     | ldr r0, [fp, -I]     | str r0, [fp, -I]     |
| int (*pf)()          | PF               | add r0, fp, -PF    | ldr r0, [fp, -PF]    | str r0, [fp, -PF]    |
| int oarg6            | OARG6            | add r0, fp, -OARG6 | ldr r0, [fp, -OARG6] | str r0, [fp, -OARG6] |
| int oarg5            | OARG5            | add r0, fp, -OARG5 | ldr r0, [fp, -OARG5] | str r0, [fp, -OARG5] |

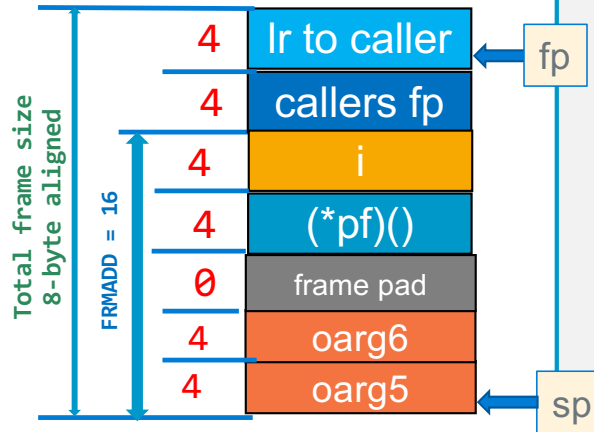
## Example: Passing Stack Args, Calling Function

```
int main()
{
 int i;
 int (*pf)(int, int) = sum;

 testp(1, 2, 3, 4, pf, &i);
 printf("%d\n", i);
 return EXIT_SUCCESS;
}
```

```
.equ FP_OFF, 4
.equ I, 4 + FP_OFF
.equ PF, 4 + I
.equ PAD, 0 + PF
.equ OARG6, 4 + PAD
.equ OARG5, 4 + OARG6
.equ FRMADD, OARG5 - FP_OFF

// FRMADD = 20 - 4 = 16
```



```
main:
 push {fp, lr}
 add fp, sp, FP_OFF
 add sp, sp, -FRMADD

 ldr r0, =sum // get func address
 add r1, fp, -PF // PF address on stack
 str r0, [r1] // store sum in var pf

 add r0, fp, -I // get address of I
 add r1, fp, -OARG6 // address of OARG6
 str r0, [r1] // arg 6: store address of I

 ldr r0, [fp, -PF] // get PF from stack
 add r1, fp, -OARG5 // address of OARG5
 str r0, [r1] // arg 5: store sum() address

 mov r0, 1 // arg 1: 1
 mov r1, 2 // arg 2: 2
 mov r2, 3 // arg 3: 3
 mov r3, 4 // arg 4: 4

 bl testp

 ldr r0, =.Lmess // arg 1: "%d\n"
 ldr r1, [fp, -I] // arg 2: i
 bl printf

 sub sp, fp, FP_OFF
 pop {fp, lr}
 bx lr
```

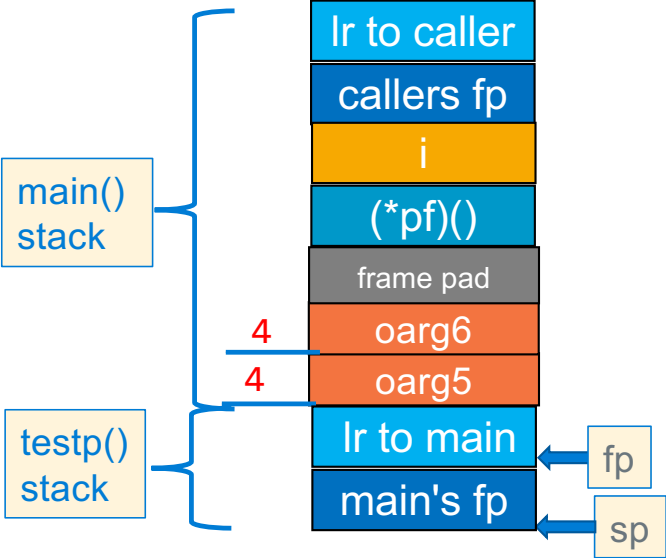
| Variable or Argument | distance from fp | Address on Stack   | Read variable        | Write Variable       |
|----------------------|------------------|--------------------|----------------------|----------------------|
| int i                | I                | add r0, fp, -I     | ldr r0, [fp, -I]     | str r0, [fp, -I]     |
| int (*pf)()          | PF               | add r0, fp, -PF    | ldr r0, [fp, -PF]    | str r0, [fp, -PF]    |
| int oarg6            | OARG6            | add r0, fp, -OARG6 | ldr r0, [fp, -OARG6] | str r0, [fp, -OARG6] |
| int oarg5            | OARG5            | add r0, fp, -OARG5 | ldr r0, [fp, -OARG5] | str r0, [fp, -OARG5] |

Example: Passing Stack Args, Called Function

arg1    arg2    arg3    arg4    arg5    arg6

```
void
testp(int j, int k, int l, int m, int (*func)(int, int), int *i)
{
 *i = func(j, k) + func(l, m);
 return;
}
```

↑  
short circuit: make this call first



```
.equ FP_OFF, 20
.equ ARG6, 8
.equ ARG5, 4

testp:
 push {r4-r7, fp, lr}
 add fp, sp, FP_OFF

 mov r4, r2 // save l
 mov r5, r3 // save m
 ldr r6, [fp, ARG5] // load func
 ldr r7, [fp, ARG6] // load i
 blx r6 // r0 = func(j, k)

 mov r1, r5 // arg 2 saved m
 mov r5, r0 // save func return value
 mov r0, r4 // arg 1 saved l
 blx r6 // r0 = func(l, m)
 add r0, r0, r5 // func(l,m) + func(j,k)
 str r0, [r7] // store sum to *i

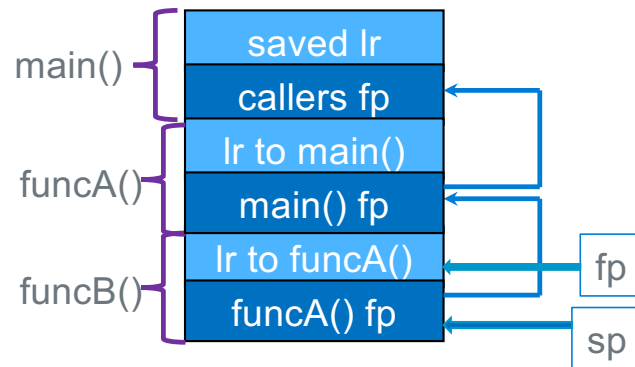
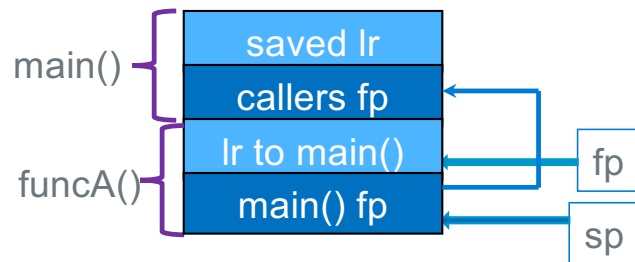
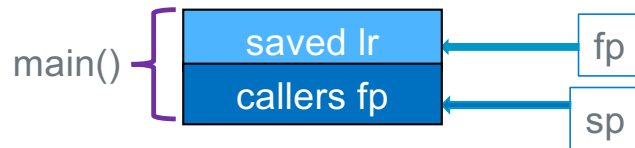
 sub sp, fp, FP_OFF
 pop {r4-r7, fp, lr}
 bx lr
```

| Argument    | distance | Address on Stack | Read variable      | Write Variable     |
|-------------|----------|------------------|--------------------|--------------------|
| int *i      | ARG6     | add r0, fp, ARG6 | ldr r0, [fp, ARG6] | str r0, [fp, ARG6] |
| int (*pf)() | ARG5     | add r0, fp, ARG5 | ldr r0, [fp, ARG5] | str r0, [fp, ARG5] |

## Extra Slides



## By following the saved fp, you can find each stack frame



How gdb finds stack frames

