CS 207: Discrete Structures

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Lecture 07 – Uncountable sets, relations, equivalence relation

Recap: Countable and countably infinite sets

Definition

- ▶ For a given set C, if there is a bijection from C to N, then C is called countably infinite. A set is countable if it is finite or countably infinite.
- ▶ Examples: \mathbb{Z} , $\mathbb{N} \times \mathbb{N}$, \mathbb{Q} .
- ▶ Properties: Closure under (countable) unions and (finite) products.

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Are all sets countable?

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Proof by contradiction: Suppose there is such a bijection, say f. This would imply that each $i \in \mathbb{N}$ maps to some set $f(i) \subseteq \mathbb{N}$.

		1	2	3	
f(0)	\checkmark	×	×	×	
f(1)	✓	×	\checkmark	\checkmark	
f(2)	×	×	×	×	
f(0) $f(1)$ $f(2)$ $f(3)$	×	\checkmark	×	\checkmark	

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- ▶ But S and f(j) differ at position j.
- ▶ Thus, $S \neq f(j)$ which is a contradiction!





Figure : Cantor and Russell

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- ▶ $S = \{i \in \mathbb{N} \mid i \notin f(i)\}$ is like the one from Russell's paradox.
- ▶ If $\exists j \in \mathbb{N}$ such that f(j) = S, then we have a contradiction.
 - ▶ If $j \in S$, then $j \notin f(j) = S$.
 - If $j \notin S$, then $j \in f(j) = S$.

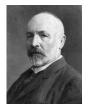




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In fact, using diagonalization Cantor showed that...

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- ▶ There is no bijection from \mathbb{R} to \mathbb{N} (H.W). Moreover, there is a bijection from \mathbb{R} to set of subsets of \mathbb{N} .
- ▶ What about the set of all sets??

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There is no bijection between $\mathbb N$ and the set of all subsets of $\mathbb N.$

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- ▶ Thus, the "size" of $\mathcal{P}(\mathbb{N})$ is strictly greater that \mathbb{N} !

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Cantor's Continuum hypothesis

There is no set whose "cardinality" is strictly between \mathbb{N} and $\mathcal{P}(\mathbb{N})$ (i.e., between naturals and reals).





Figure: 1st of Hilbert's 23 problems for the 20th century in 1900.

What did the world think about these proofs (in 1890s?)







(a) Kronecker (b) Poincare

(c) Theologians

- ► Kronecker: Only constructive proofs are proofs! "Scientific Charlatan", "Corruptor of youth"!
- ▶ Poincare: Set theory is a "disease" from which mathematics will be cured.
- ► Christian Theologians: God=Uniqueness of an absolute infinity. So, what is all this different infinities...?!

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- ▶ Hilbert: No one can expel us from the paradise that Cantor has created for us.

Summary and moving on...

- ▶ Finite and infinite sets.
- ▶ Using functions to compare sets: focus on bijections.
- ▶ Countable, countably infinite and uncountable sets.
- ► Cantor's diagonalization argument (A new powerful proof technique!).

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Next: Basic Mathematical Structures – Relations

Relations

Definition: Function

Let A, B be two sets. A function f from A to B is a subset R of $A \times B$ such that

- (i) $\forall a \in A, \exists b \in B \text{ such that } (a, b) \in R, \text{ and }$
- (ii) if $(a, b) \in R$ and $(a, c) \in R$, then b = c.
 - Now, suppose A is the set of all Btech students and B is the set of all courses. Clearly, we can assign to each student the set of courses he/she is taking. Is this a function?

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 - ▶ Now, suppose A is the set of all Btech students and B is the set of all courses. Clearly, we can assign to each student the set of courses he/she is taking. Is this a function?
 - ▶ By removing the two extra assumptions in the defn, we get:

Definition: Relation

- ▶ A relation R from A to B is a subset of $A \times B$. If $(a,b) \in R$, we also write this as a R b.
- ▶ Thus, a relation is a way to relate the elements of two (not necessarily different) sets.

Examples and representations of relations

We write R(A, B) for a relation from A to B and just R(A) if A = B. Also if A is clear from context, we just write R.

Examples of relations

- ▶ All functions are relations.
- $R_1(\mathbb{Z}) = \{(a,b) \mid a,b \in \mathbb{Z}, a-b \text{ is even } \}.$
- $R_2(\mathbb{Z}) = \{(a,b) \mid a,b \in \mathbb{Z}, a \le b\}.$
- ▶ Let S be a set, $R_3(\mathcal{P}(S)) = \{(A, B) \mid A, B \subseteq S, A \subseteq B\}.$
- ▶ Relational databases are practical examples.

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- ▶ Let S be a set, $R_3(\mathcal{P}(S)) = \{(A, B) \mid A, B \subseteq S, A \subseteq B\}.$
- ▶ Relational databases are practical examples.

Representations of a relation from A to B.

- \blacktriangleright As a set of ordered pairs of elements, i.e., subset of $A \times B$.
- ► As a directed graph.
- ► As a (database) table.

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- ▶ Are there other special relations? What are they useful for?

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- ► Functions were special kinds of relations that were useful to compare sets.
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 - ► Equivalence relations
 - ▶ Partial orders

Partitions of a set – grouping "like" elements

Examples

- ▶ Natural numbers are partitioned into even and odd.
- ► This class is partitioned into set of all those who are taking the same set of courses.

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A partition of a set S is a set P of its subsets such that

- if $S' \in P$, then $S' \neq \emptyset$.
- $\bigcup_{S' \in P} S' = S$: its union covers entire set S.
- ▶ If $S_1, S_2 \in P$, then $S_1 \cap S_2 = \emptyset$: sets are disjoint.

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Can you think of two trivial partitions that any set must have?

Interpreting partitions as relations

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What properties does this relation have?

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We will see that any relation that satisfies these three properties defines a partition!

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Definition

A relation which satisfies all these three properties is called an equivalence relation.

Examples

▶ Reflexive: $\forall a \in S, aRa$.

▶ Symmetric: $\forall a, b \in S$, aRb implies bRa.

▶ Transitive: $\forall a, b, c \in S$, aRb, bRc implies aRc.

▶ Equivalence: Reflexive, Symmetric and Transitive.

Relation	Refl.	Symm.	Trans.	Equiv.
aRb if students a and b take	✓	✓	✓	√
same set of courses				
$ \overline{\{(a,b) \mid a,b \in \mathbb{Z}, (a-b)\}} $	√	√	√	√
$\mod 2 = 0\}$				
$\overline{\{(a,b) \mid a,b \in \mathbb{Z}, a \le b\}}$				
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$\{(a,b) \mid a,b \in \mathbb{R}, a-b < 1\}$				
$\overline{\{((a,b),(c,d)) \mid (a,b),(c,d) \in A\}}$				
$\mathbb{Z} \times (\mathbb{Z} \setminus \{0\}), (ad = bc)\}$				