

CS 207: Discrete Structures

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Lecture 08 – Basic mathematical structures
Equivalence relations and partitions, posets

Recap: Relations

Definition: Relation

- ▶ A **relation** R from A to B is a subset of $A \times B$. If $(a, b) \in R$, we also write this as $a R b$.

We write $R(A, B)$ for a relation from A to B and just $R(A)$ if $A = B$. Also if A is clear from context, we just write R .

- ▶ All functions are relations.
- ▶ $R_1(\mathbb{Z}) = \{(a, b) \mid a, b \in \mathbb{Z}, a - b \text{ is even} \}$.
- ▶ Relational databases are practical examples.

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Representations of a relation from A to B .

- ▶ As a set of **ordered pairs of elements**, i.e., subset of $A \times B$, as a **directed graph**, as a **(database) table**.

Special types of relations

- ▶ **Reflexive:** $\forall a \in S, aRa$.
- ▶ **Symmetric:** $\forall a, b \in S, aRb$ implies bRa .
- ▶ **Transitive:** $\forall a, b, c \in S, aRb, bRc$ implies aRc .
- ▶ **Equivalence:** Reflexive, Symmetric and Transitive.

| Relation | Refl. | Symm. | Trans. | Equiv. |
|--|-------|-------|--------|--------|
| aRb if students a and b take same set of courses | ✓ | ✓ | ✓ | ✓ |
| $\{(a, b) \mid a, b \in \mathbb{Z}, (a - b) \bmod 2 = 0\}$ | ✓ | ✓ | ✓ | ✓ |
| $\{(a, b) \mid a, b \in \mathbb{Z}, a \leq b\}$ | | | | |
| $\{(a, b) \mid a, b \in \mathbb{Z}, a < b\}$ | | | | |
| $\{(a, b) \mid a, b \in \mathbb{Z}, a \mid b\}$ | | | | |
| $\{(a, b) \mid a, b \in \mathbb{R}, a - b < 1\}$ | | | | |
| $\{((a, b), (c, d)) \mid (a, b), (c, d) \in \mathbb{Z} \times (\mathbb{Z} \setminus \{0\}), (ad = bc)\}$ | | | | |

Partitions of a set – grouping “like” elements

Definition

A partition of a set S is a set P of its subsets such that

- ▶ if $S' \in P$, then $S' \neq \emptyset$.
- ▶ $\bigcup_{S' \in P} S' = S$: its union covers entire set S .
- ▶ If $S_1, S_2 \in P$, then $S_1 \cap S_2 = \emptyset$: sets are disjoint.

Example: natural numbers partitioned into even and odd...

Theorem

Every partition of set S gives rise to a **canonical** equivalence relation R on S , namely,

- ▶ aRb if a and b belong to the same set in the partition of S .

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Is the converse true? Can we generate a partition from every equivalence relation?

Equivalence classes

Definition

- ▶ Let R be an equivalence relation on set S , and let $a \in S$.
- ▶ Then the **equivalence class** of a , denoted $[a]$, is the set of all elements related to it, i.e., $[a] = \{b \in S \mid (a, b) \in R\}$.

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Lemma

Let R be an equivalence relation on S . Let $a, b \in S$. Then, the following statements are equivalent:

1. aRb
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Proof Sketch: (1) to (2) symm and trans, (2) to (3) refl, (3) to (1) symm and trans. (H.W.: Do the proof formally.)

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Proof sketch of (1): Union, non-emptiness follows from reflexivity. The rest (pairwise disjointness) follows from the previous lemma.

(H.W.): Write the formal proofs of (1) and (2).

More “applications” of equivalence relations

Defining new objects using equivalence relations

Consider

$$R = \{((a, b), (c, d)) \mid (a, b), (c, d) \in \mathbb{Z} \times (\mathbb{Z} \setminus \{0\}), (ad = bc)\}.$$

- ▶ Then the equivalence classes of R define the rational numbers.
- ▶ e.g., $\left[\frac{1}{2}\right] = \left[\frac{2}{4}\right]$ are two names for the same rational number.
- ▶ Indeed, when we write $\frac{p}{q}$ we implicitly mean $\left[\frac{p}{q}\right]$.

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- ▶ With this definition, why are addition and multiplication “well-defined”?

Can we define **integers** and **real numbers** starting from naturals by using equivalence classes?

Geometrical objects using equivalence relations

Cut-and-paste

Consider the relation $R([0, 1]) = \{aRb \mid a, b \in [0, 1], \text{ either } a = b \text{ or } a = 1, b = 0, \text{ or } a = 0, b = 1\}$.

- Is R an equivalence relation? What does it define?

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- ▶ Is R an equivalence relation? What does it define?
- ▶ This is $[0, 1]$ in which the end-points have been related to each other.
- ▶ So the equivalence classes form a “loop”, since end-points are joined. If we imagine $[0, 1]$ as a 1-length string, we have glued its ends!

Geometrical objects using equivalence relations

Forming 2D objects

Consider a rectangular piece of the real plane, $[0, 1] \times [0, 1]$.

- ▶ Define $R_1([0, 1] \times [0, 1])$ by $(a, b)R_1(c, d)$ if
 - ▶ $(a, b) = (c, d)$ or
 - ▶ $b = d, a = 0, c = 1$ or
 - ▶ $b = d, c = 0, a = 1$.

Is R_1 an equivalence relation? What do its equivalence classes define?

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Is R_1 an equivalence relation? What do its equivalence classes define?

- ▶ Define $R_2([0, 1] \times [0, 1])$ by $(a, b)R_2(c, d)$ if
 - ▶ $(a, b) = (c, d)$ or
 - ▶ $a, b, c, d \in \{0, 1\}$.

Is R_2 an equivalence relation? What does it define?

Can you build even more interesting “shapes”? Torus? Mobius strip?!

Partial Orders

Consider $\{(a, b) \mid a, b \in \mathbb{Z}, a \leq b\}$.

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A relation R on S is **anti-symmetric** if for all $a, b \in S$ aRb and bRa implies $a = b$.

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Examples:

- ▶ $R_1(\mathbb{Z}) = \{(a, b) \mid a, b \in \mathbb{Z}, a \leq b\}$.
- ▶ $R_2(\mathcal{P}(S)) = \{(A, B) \mid A, B \in \mathcal{P}(S), A \subseteq B\}$.

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Definition

A **partial order** is a relation which is **reflexive**, **transitive** and **anti-symmetric**.

Partial orders and equivalences relations

- ▶ **Reflexive:** $\forall a \in S, aRa$.
- ▶ **Symmetric:** $\forall a, b \in S, aRb$ implies bRa .
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| | Reflexive | Transitive | Symmetric | Anti-symmetric |
|----------------------|-----------|------------|-----------|----------------|
| Equivalence relation | ✓ | ✓ | ✓ | |
| Partial order | ✓ | ✓ | | ✓ |

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| $\{(A, B) \mid A, B \in \mathcal{P}(S), A \subseteq B\}$ | ✓ | ✓ | ✓ | ✓ |
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 - ▶ i.e., $\forall a, b \in S$, either $a \preceq b$ or $b \preceq a$.

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- ▶ A total order is a partial order \preceq on S in which every pair of elements is comparable
- ▶ Qn: Can a relation be symmetric and anti-symmetric?
- ▶ (H.W): Can a relation be neither symmetric nor anti-symmetric?

Partially ordered sets (Posets)

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Examples

- ▶ (\mathbb{Z}, \leq) : integers with the usual less than or equal to relation.
- ▶ $(\mathcal{P}(S), \subseteq)$: powerset of any set with the subset relation.
- ▶ $(\mathbb{Z}^+, |)$: positive integers with divisibility relation.