

TCP Versions: Reno

Kameswari Chebrolu

All the figures used as part of the slides are either self created or from the public domain with either 'creative commons' or 'public domain dedication' licensing. The public sites from which some of the figures have been picked include:

<http://commons.wikimedia.org> (Wikipedia, Wikimedia and workbooks); <http://www.sxc.hu> and <http://www.pixabay.com>

Break

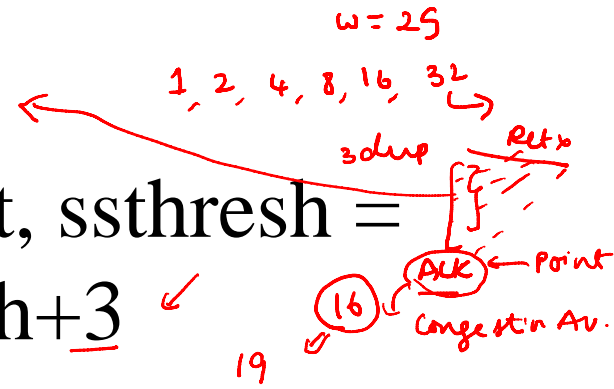


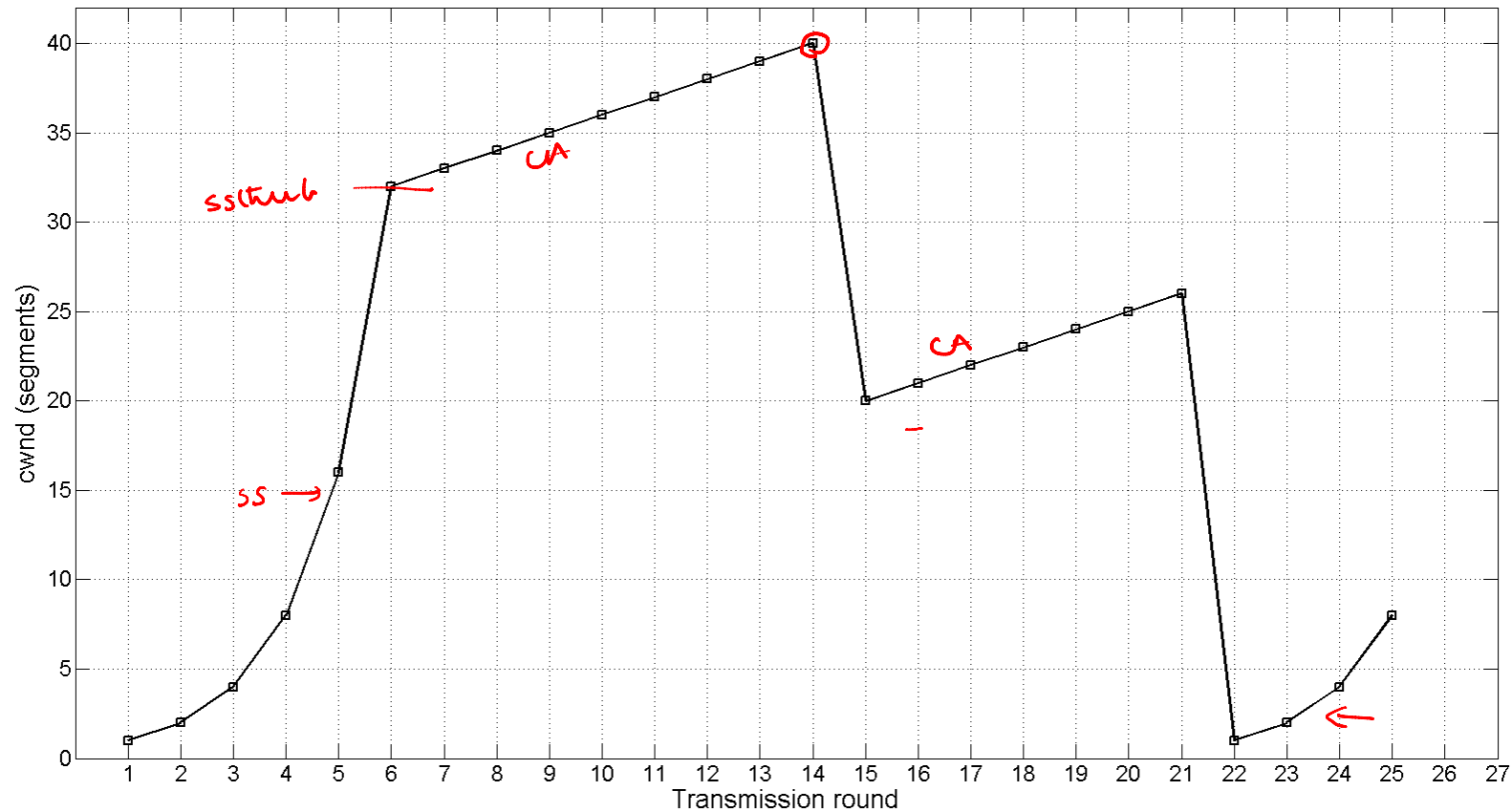
TCP Reno

- Incorporates two new mechanisms: Fast Retransmit and Fast Recovery
- Fast Retransmit: Retransmit packet at sender after 3 duplicate acks → Packet is lost
reorder ← 1 or 10?
- Cut the window by half (loss event)
- Avoids having to time-out which keep the link idle for longer duration

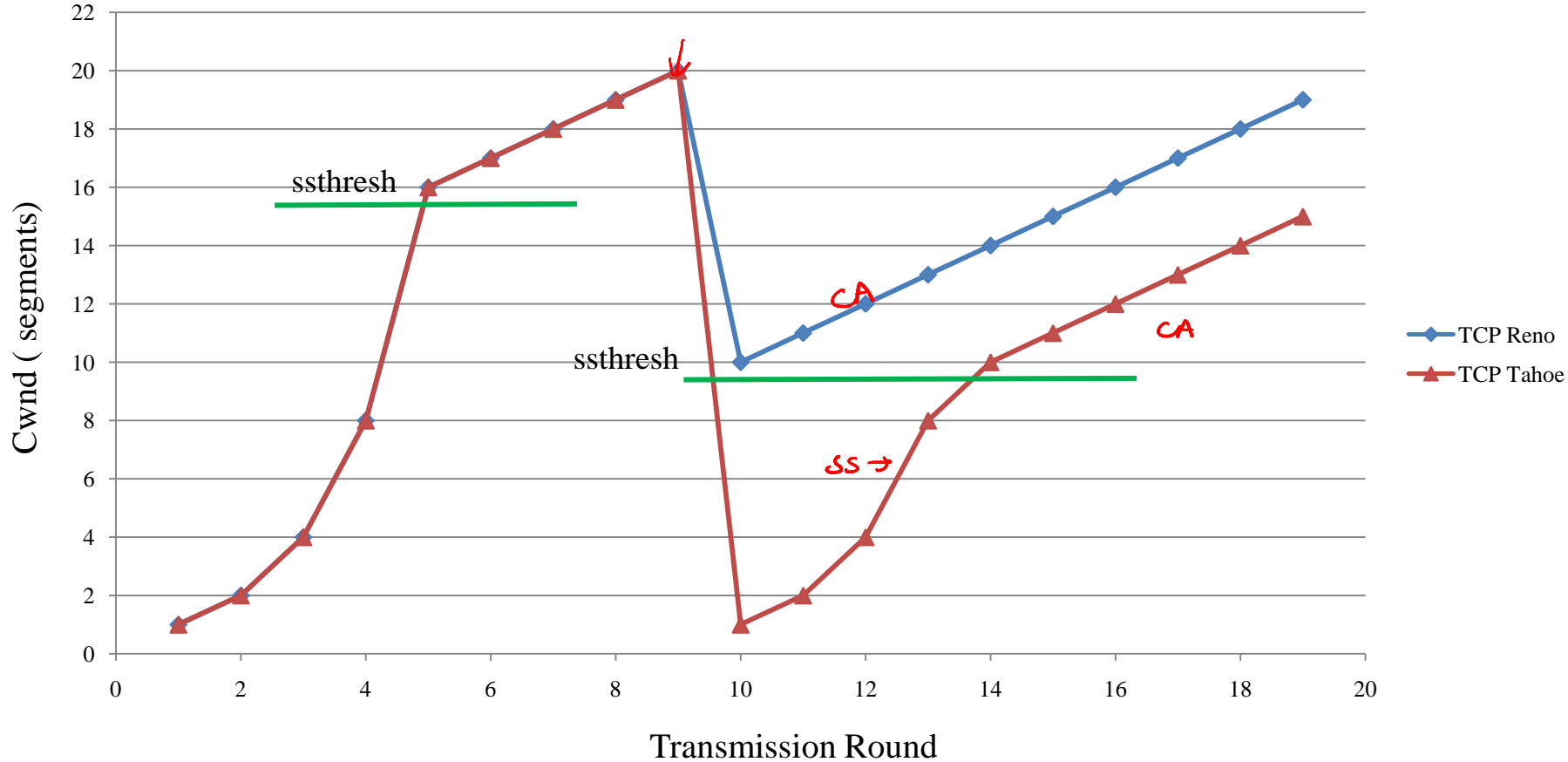
Fast Recovery

- On 3rd dupack, retransmit packet, $ssthresh = \max(\underline{cwnd/2}, 2)$; $cwnd = ssthresh + \underline{3}$
- Another dupack, $cwnd = \underline{cwnd} + 1$; transmit packet if allowed by $cwnd$
- On ack acknowledging new data, $cwnd = \underline{ssthresh}$, invoke congestion avoidance (linear increase in $cwnd$ now on)



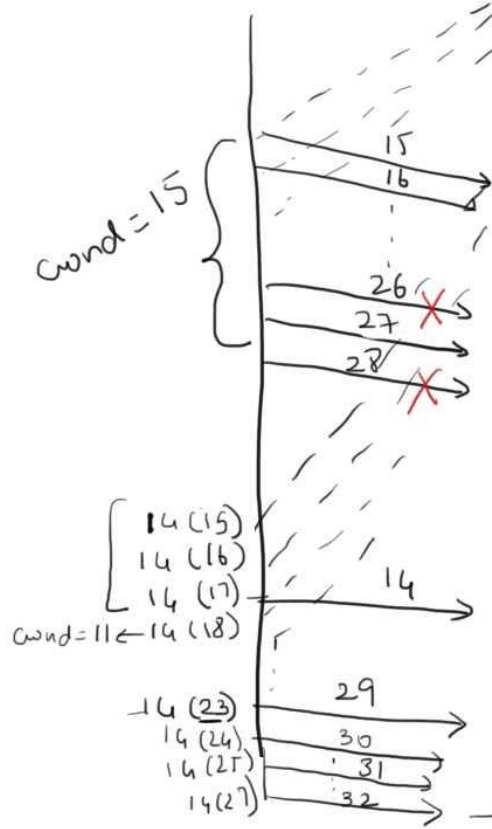
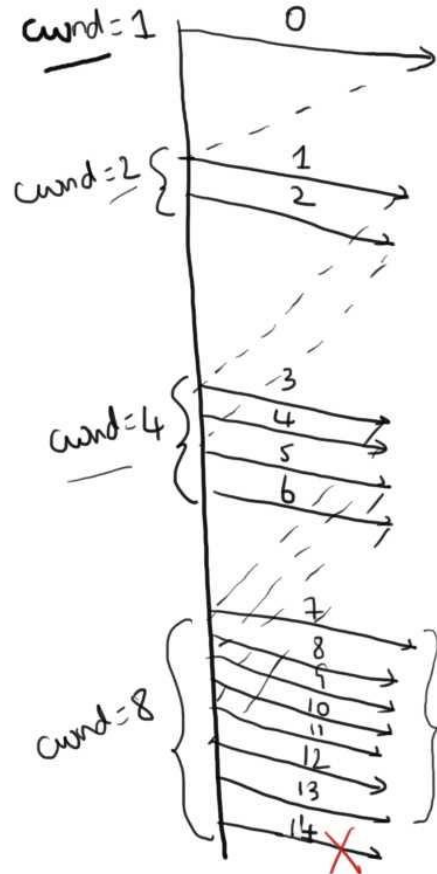


Comparison

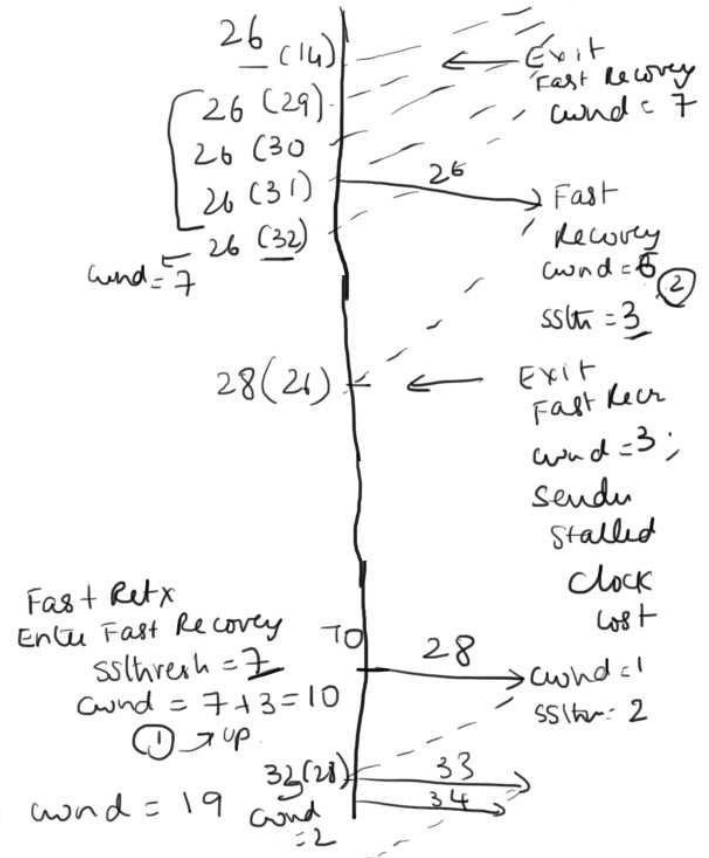


x (4) → who generated
L) asking for

TCP Reno



(Not drawn to scale) ① 14, 15-28
② 26 to 32
Pkt in Pipe 15



Other Versions

- TCP NewReno: Handles multiple losses per congestion window better (high loss rate scenarios)
- TCP Vegas: Uses packet delay to signal congestion than loss event
- TCP SACK: Employs selective acknowledgments

Summary

- TCP Tahoe: Go-back-N with slow start and congestion avoidance *Fast Retransmit*
 - Loss recovery slow; timeouts drain pipe
- TCP Reno: improves upon Tahoe
 - Better loss recover via duplicate acks (fast retransmit)
 - Prevents draining of pipe after fast retransmit (avoids slow-start)
- Ahead: Sliding window, flow control, and other miscellaneous things