Problem Set 8

- 1. Write second order logic formulae to capture the following:
 - (a) There is a path from node s to node t in the graph. The signature is $\tau = \{E\}.$
 - (b) Every bounded non empty set has a least upper bound. The signature is $\tau = \{\leq\}$
 - (c) The graph is acyclic. The signature is $\tau = \{E\}$.
- 2. Let Σ be a finite alphabet. The atomic formulae in MSO defined over Σ^* are x = y, x < y, S(x, y), X(x) and $Q_a(x), a \in \Sigma$. Consider the following logic called MSO₀ having atomic formulae of the following forms:

$$Sing(X), X \subseteq Y, X < Y, S(X, Y), Q_a(X)$$

where

- -Sing(X) means that X is a SO variable of cardinality 1;
- $-X \subseteq Y$ means that every element of the SO variable X is contained in the SO variable Y;
- -X < Y means that SO variables X, Y have cardinality 1, and that the element in Y is greater than the element in X;
- -S(X,Y) means that SO variables X,Y have cardinality 1, and Y contains the successor of the element in X; and,
- $-Q_a(X)$ means that all positions in X are decorated by $a \in \Sigma$.

If φ is an atomic formula in MSO, then $\varphi \wedge \varphi, \neg \varphi, \varphi \vee \varphi, \forall x \varphi$ and $\forall X \varphi$ are formulae in MSO. Similarly, if φ is an atomic formula in MSO₀, then, $\varphi \wedge \varphi, \neg \varphi, \varphi \vee \varphi$ and $\forall X \varphi$ are formulae in MSO₀. Compare the expressiveness of MSO and MSO_0 .

- 3. For the formula $\exists x \forall y (x < y \rightarrow Q_a(y))$ give an equivalent MSO₀ formula.
- 4. Write an MSO formula for x < y without using <. You can use S in your MSO formula.
- 5. Consider the following NFA $N = (\{0, 1, 2, 3\}, \{a, b\}, \Delta, \{0\}, \{1\})$ with $\Delta(0, b) = \{1\}$, $\Delta(1, a) = \{2\}$, $\Delta(2, a) = \{2\}$, $\Delta(2, b) = \{3\}$ and $\Delta(3, b) = \{0\}$. Write an MSO formula with two SO variables that characterizes L(N).