

# Proving that $L_{a,b}$ is not a regular language

## Lemma

*There is no finite state automaton accepting  $L_{a,b}$ .*

## Proof.

By Pigeon Hole Principle.

Suppose  $\exists i, j \in \mathbb{N}$  such that  $i \neq j$ ,  
automaton reaches the same state after reading both  $a^i, a^j$ .

Then  $a^i \cdot b^j$  and  $a^j \cdot b^j$  are both accepted or both rejected,  
which is a contradiction.



# Pumping lemma

A recipe for proving that a given language is non-regular.

## Lemma (Pumping Lemma)

*If  $L$  is a regular language, then  $\exists p \in \mathbb{N}$  such that for any strings  $x, y, z$  with  $x \cdot y \cdot z \in L$  and  $|y| \geq p$ ,*

- ① *there exist strings  $u, v, w$ , s.t.  $y$  can be written as  $y = u \cdot v \cdot w$ ,*
- ②  *$\forall i \geq 0 \ x \cdot u \cdot v^i \cdot w \cdot z \in L$ ,*
- ③  *$|v| > 0$ .*

To prove that a given language  $L$  is not regular, the contrapositive of the above statement is useful.

# Contrapositive of the pumping lemma

## Lemma

We say that a language  $L$  satisfies **Property-NR** if the following conditions hold:

$$\forall p \geq 0,$$

$$\exists x, y, z \text{ such that } x \cdot y \cdot z \in L \text{ and } |y| \geq p,$$

$$\forall u, v, w \text{ such that } |v| > 0, y = u \cdot v \cdot w,$$

$$\exists i \ x \cdot u \cdot v^i \cdot w \cdot z \notin L.$$

If  $L$  satisfies Property-NR then  $L$  is not regular.

# Using the pumping lemma

We say that a language  $L$  satisfies **Property-NR** if the following conditions hold:

- ☹  $\forall p \geq 0,$
- ☺  $\exists x, y, z$  such that  $x \cdot y \cdot z \in L$  and  $|y| \geq p,$
- ☹  $\forall u, v, w$  such that  $|v| > 0, y = u \cdot v \cdot w,$
- ☺  $\exists i$   $x \cdot u \cdot v^i \cdot w \cdot z \notin L.$

If  $L$  satisfies Property-NR then  $L$  is not regular.

We will now use the lemma to prove that  $L_{a,b} = \{a^n b^n \mid n \geq 0\}$  is not regular.

For any chosen  $p \geq 0$ , let  $x := a^p$ ,  
 $y := b^p$ ,  $z = \epsilon$ .

For any split of  $y$  as  $u \cdot v \cdot w$ , if we  
take  $x \cdot u \cdot v^i \cdot w = a^p b^q$ , where  $q > p$  as  
long as  $i > 0$ .

In particular,  $x \cdot u \cdot v^2 \cdot w \cdot z \notin L$ .