

# The Power of Security Policy Modeling

1. Message authentication code (MAC)-- data integrity mechanism that provides **integrity**, but **no confidentiality**.
2. Chosen plaintext secure encryption (CPA-secure encryption) provides **confidentiality against eavesdropping**, **but is not secure** against an active attacker who tampers with traffic.
3. Intuitively, combining the two primitives should provide both confidentiality and integrity against an active adversary.
  - How to do this integration?

# Integrating Confidentiality and Integrity

- $k_e$  and  $k_m$  denote the encryption and MAC keys, respectively

- $x || y$  denotes concatenation of  $x$  and  $y$

TLS: First, compute a checksum over the message, append it to the message, and encrypt the result. In symbols, compute

$$t := \text{MAC}(k_m, m) \text{ and output } c := E(k_e, m || t).$$

IPsec: First, encrypt the message and then output the resulting ciphertext followed by a checksum computed over the ciphertext. In symbols, compute

$c_0 := E(k_e, m)$  and output  $c := c_0 || \text{MAC}(k_m, c_0)$ .

SSH: Send the concatenation of the separately computed encryption and checksum. In symbols, compute and output

$$c := E(k_e, m) \parallel MAC(k_m, m).$$

During decryption, if the relevant integrity tag fails to verify, the decryption algorithm outputs a distinguished symbol (\$) to indicate error

Which method is Right and which is better?

# Threat Model

- Threat model associated with authenticated encryption:
  - the attacker is able to obtain the encryption of arbitrary messages of its choice
  - Attacker's goal :
    - Learn information about the decryption of a well-formed challenge ciphertext (thereby defeating confidentiality),
    - or generate a new well-formed ciphertext different from all ciphertexts previously given to the attacker (thereby defeating integrity).
- **If the attacker cannot do either then we say that the system provides authenticated encryption**

# Choice: TLS

- Not generically secure:
  - there are specific instances of encryption and MAC such that the TLS combination does not provide authenticated encryption.
  - However, for specific encryption systems, such as randomized counter mode encryption, TLS method provides authenticated encryption even if the MAC is only weakly secure (so called, one-time secure). The reason is that the MAC is protected by the encryption and therefore need not be a fully secure MAC; weak MAC security is sufficient.

# Choice: IPSEC

- The IPsec construction can be shown to provide authenticated encryption for any MAC and CPAsecure encryption.
- The basic reason is that the MAC locks the ciphertext so that any modification of the ciphertext **en-route will be detected by the decryptor.**

# Choice: SSH

- The SSH construction is known to be secure when a very specific MAC is used, but may not be secure for a general purpose MAC. To see why, recall that a MAC need not preserve confidentiality and therefore  $\text{MAC}(k, m)$  may leak information about the encrypted plaintext.



# Choices

- Based on these comparisons, a designer can choose the appropriate method for the application at hand.
  - When countermode encryption is used, the TLS construction is adequate even if a simple MAC is used.
  - Otherwise, one should use the IPsec construction.
- This clear understanding is only made possible thanks to the precise formulation of authenticated encryption

# What do we learn

- Using the definition of authenticated encryption, the National Institute of Standards and Technology (NIST) was able to publish precise encryption modes, called CCM and GCM, designed to meet the definition
- Once the goals of authenticated encryption were clearly spelled out, it turned out that authenticated encryption can be built far more efficiently than by combining encryption and MAC algorithms

Reference: Privacy and Cybersecurity: the next 100 years  
Carl Landwehr et al., Vol 100, Proceedings IEEE 2012, 13 May 2012