# Interference-aware Scheduler for Serverless Computing

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### Background

Co-location between latency-sensitive services and throughput sensitive jobs is an appealing approach to improving memory utilization in multi-tenant datacenter systems.

Origin of problem: co-location(competition) and unpredictable workload(execution time)

Tradition OS scheduler: CFS not efficient (frequently context switch for short function)

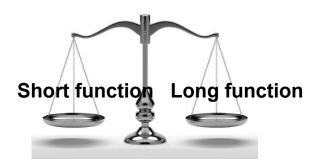
Our goal: develop scheduler to improve Service Level Objectives (SLOs)

#### Motivation

Solve problem-add another level

CFS - FIFO+CFS

Trade-off



# Compute SLO

SLO(service level objectives)

Definition: X% of function invocations are finished within a deadline

How to set: test the function in ideally isolated environment

E.X.: local PC, server

```
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.00020432472229003906
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.00023865699768066406
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002079010009765625
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.00022077560424804688
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.000209808349609375
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002155303955078125
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002135462829589844
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002079010009765625
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002079010009765625
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.00020647048950195312
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002065FSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.0002065FSRunxin:/home/sfswu/Desktop/project# python3 test1.py 0.00020647048950195312
```

```
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 3.0994415283203125e-06 root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 3.0994415283203125e-06 root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 3.0994415283203125e-06 root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 2.6226043701171875e-06 root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 2.86102294921875e-06 root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 2.6226043701171875e-06 root@SFSRunxin:/home/sfswu/Desktop/project# python3 test_auth.py 2.6226043701171875e-06
```



#### Influence of function execution time

Different OS: windows 10 vs ubuntu 20.04

```
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.00020432472229003906
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.00023865699768066406
root@SFSRunxin:/home/sfswu/Desktop/project# pvthon3 test1.pv
0.0002079010009765625
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.00022077560424804688
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.000209808349609375
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.0002155303955078125
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.00021338462829589844
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.0002079010009765625
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.00020647048950195312
root@SFSRunxin:/home/sfswu/Desktop/project# python3 test1.py
0.00020170211791992188
```

(base) C:\Users\xitian\Desktop>python test1.py 0.0019485950469970703

(base) C:\Users\xitian\Desktop>python test1.py 0.0009770393371582031

(base) C:\Users\xitian\Desktop>cd Desktop 系统找不到指定的路径。

(base) C:\Users\xitian\Desktop>python test1.py 0.0

(base) C:\Users\xitian\Desktop>python test1.py 0.001954793930053711

(base) C:\Users\xitian\Desktop>python test1.py 0.0019533634185791016



#### Influence of function execution time

Sever vs ubuntu

Execution time:

windows 10 < ubuntu 22.04

Ubuntu 22.04 < server

conclusion:

Which ideally isolated environment is the best

Keep the environment the same



# step

- 1. Extract the critical function
- 2. Run the function several time (same environment)
- 3. Compute the average value of the execution time.
- 4. Modify this value

```
generatePolicy(principalId, effect, resource):
start time=time.time()
authResponse = Empty()
authResponse.principalId = principalId
if effect and resource:
    policyDocument = Empty()
    policyDocument.Version = '2012-10-17'
    policyDocument.Statement = [None]
    statementOne = Empty()
    statementOne.Action = 'execute-api:Invoke'
    statementOne.Effect = effect
    statementOne.Resource = resource
    policyDocument.Statement[0] = statementOne
    authResponse.policyDocument = policyDocument
authResponse.context = {
  "stringKey": "stringval",
  "numberKey": 123,
  "booleanKey": True
end time=time.time()
execute time=end time-start time
print(execute time)
return authResponse
resource = "arn:aws:execute-api:{regionId}:{accountId}:{apiId}/{stage}/{httpVerb}/[{resource}/[{child-resources
generatePolicy('user', 'Allow', resource)
```

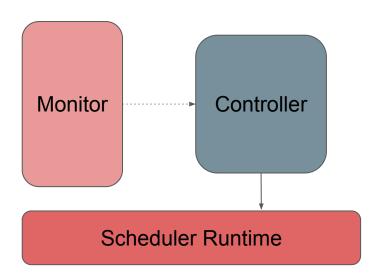
```
import os
import sys
import pyaes
import argparse
import time
parser = argparse.ArgumentParser()
parser.add argument("--default plaintext", default="defaultplaintext", help="Default plain text if plaintext messag
parser.add argument("-k", "--key", dest="KEY", default="6368616e676520746869732070617373", help="Secret key")
args = parser.parse args()
KEY = args.KEY.encode(encoding = 'UTF-8')
def AESModeCTR(plaintext):
    start time=time.time()
    counter = pyaes.Counter(initial value = 0)
    aes = pyaes.AESModeOfOperationCTR(KEY, counter = counter)
    ciphertext = aes.encrypt(plaintext)
    end_time=time.time()
    execute_time=end_time-start_time
    print(execute time)
    return ciphertext
if name == ' main ':
    plaintext = args.default plaintext
    AESModeCTR(plaintext)
                                                                        \Pi \Pi \Pi \Pi \Pi
```

# Framework Implementation - Overview

- Our Framework involves three main components

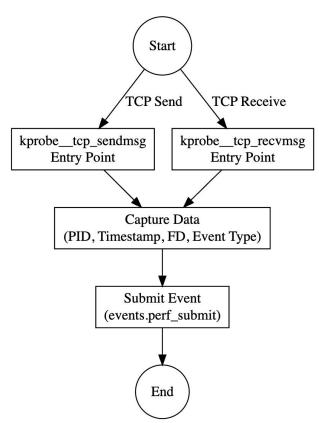
- Scheduler Runtime
- Monitor
- Controller

## Framework Implementation - Overview

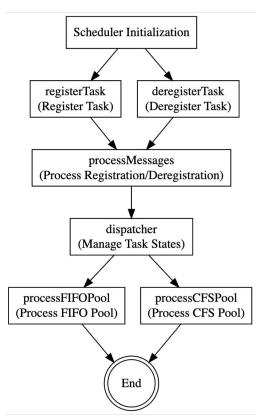


# Framework Implementation - Monitor

Specifically, we use kprobes on tcp\_sendmsg and tcp\_recvmsg functions, capturing data such as process ID (PID), timestamp, file descriptor (FD), and event type

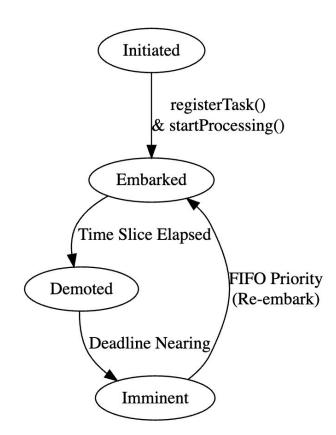


## Framework Implementation - Runtime

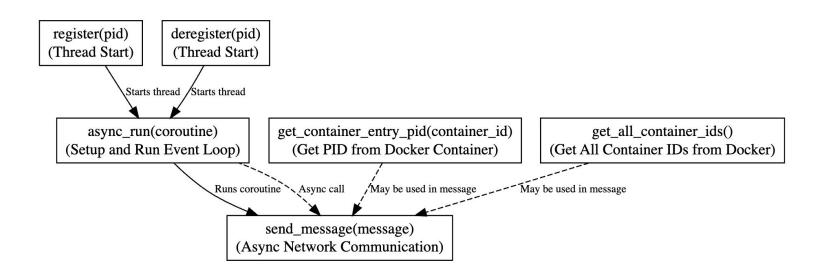


# Framework Implementation - Runtime

The tasks transition between various states (Initiated, Embarked, Demoted, Imminent).



# Framework Implementation - Controller



# Register

One big challenge in designing and implementing the register is that our platform uses TCP socket to trace the requests. So the register has to handle request status changes over multiple packets.

#### Example of the trace:

CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: recv,

Timestamp: 93642890.399628 milliseconds

CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: recv,

Timestamp: 93642890.407442 milliseconds

CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send,

Timestamp: 93647891.92718 milliseconds

CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send,

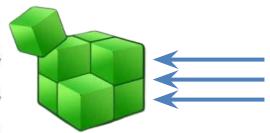
Timestamp: 93647891.940696 milliseconds

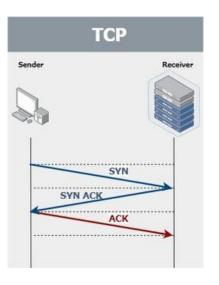
CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send,

Timestamp: 93647892.062655 milliseconds

CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send,

Timestamp: 93647892.064929 milliseconds





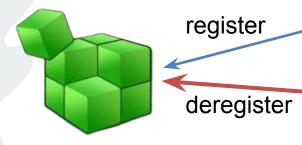
A dictionary variable was used to maintain a register, monitoring the current active requests.



# Register

#### When to register/deregister requests?

if FD not in dict and event == recv:
 dict[FD]=request status
if FD in dict and event == send:
 del dict[FD]



The register is maintained as a global variable in the monitor program.

CPU: 28, PID: 61883, FD: 3843769, Size: 36 Event: recv, Timestamp: 93642890.399628 milliseconds CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: recv, Timestamp: 93642890.407442 milliseconds CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send, Timestamp: 93647891.92718 milliseconds CPU: 28, PID: 61863, FD: 3843769, Size: 36, Event: send, Timestamp: 93647891.940696 milliseconds CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send, Timestamp: 93647892.062655 milliseconds CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send, Timestamp: 93647892.062659 milliseconds CPU: 28, PID: 61883, FD: 3843769, Size: 36, Event: send, Timestamp: 93647892.064929 milliseconds



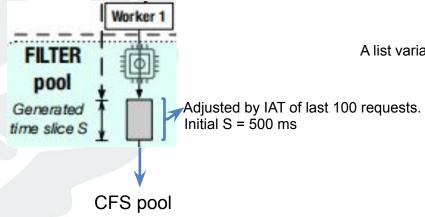
#### Time slice

The inter arrival rate (IAT) was used to dynamically adjust the time slice S=IAT \* c, c is the number of commuter cores.

Design of dynamical time slice:

inter arrival time = 1/arrival rate

- (1)Accommodate the skewed execution duration of the workload.
- (2)Strike a balance between queuing delay and execution time.



A list variable is maintained to record the last 100 requests in the monitor program.

```
history requests append((event.fd_event.timestamp / 1000000))

n = len(history_requests)

if n <= N and n>1:

IAT = (event.timestamp / 1000000 - history_requests[0][1]) / (n-1)

elif n>N:

IAT = (event.timestamp / 1000000 - history_requests[0][1]) / N

history_requests.pop(0)
```



#### Time slice

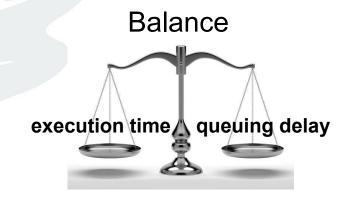
S=IAT \* c  $\rho=\lambda/(c\mu)$ , according to the M/G/c model with Kendall's notation in queuing theory.

ρ: traffic intensity per core
λ: arrival rate of the requests
μ: service rate of a single core
c: the number of cores used

$$S \xrightarrow{bound} \mu \xrightarrow{affect} \rho \xrightarrow{affect}$$
 queuing delay

#### **Design intuition**

All functions whose execution duration is shorter than S run to completion without being preempted.





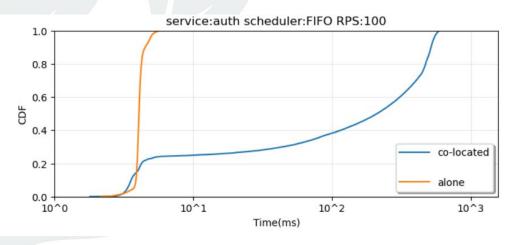


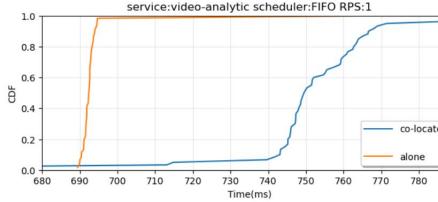
#### Evaluation:baseline

Baseline

Service: 1.auth 2.video-analytic

Scheduler: FIFO





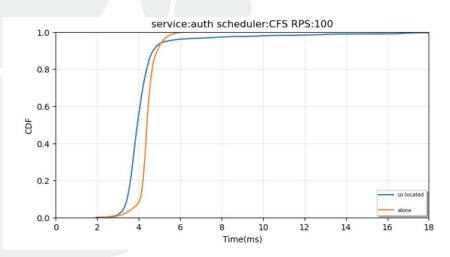


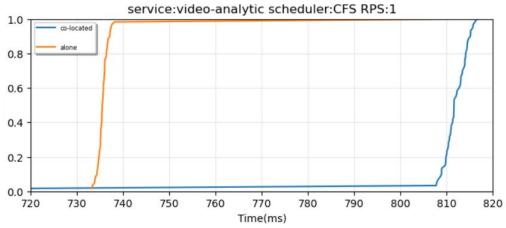
#### Evaluation:baseline

Baseline

Service: 1.auth 2.video-analytic

Scheduler: CFS







#### **Evaluation**

evaluation

Service: 1.auth 2.video-analytic

Scheduler: project scheduler

