

# Internet Technology

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<https://www.coursera.org/course/insidetheinternet>



# Copyright Thanks

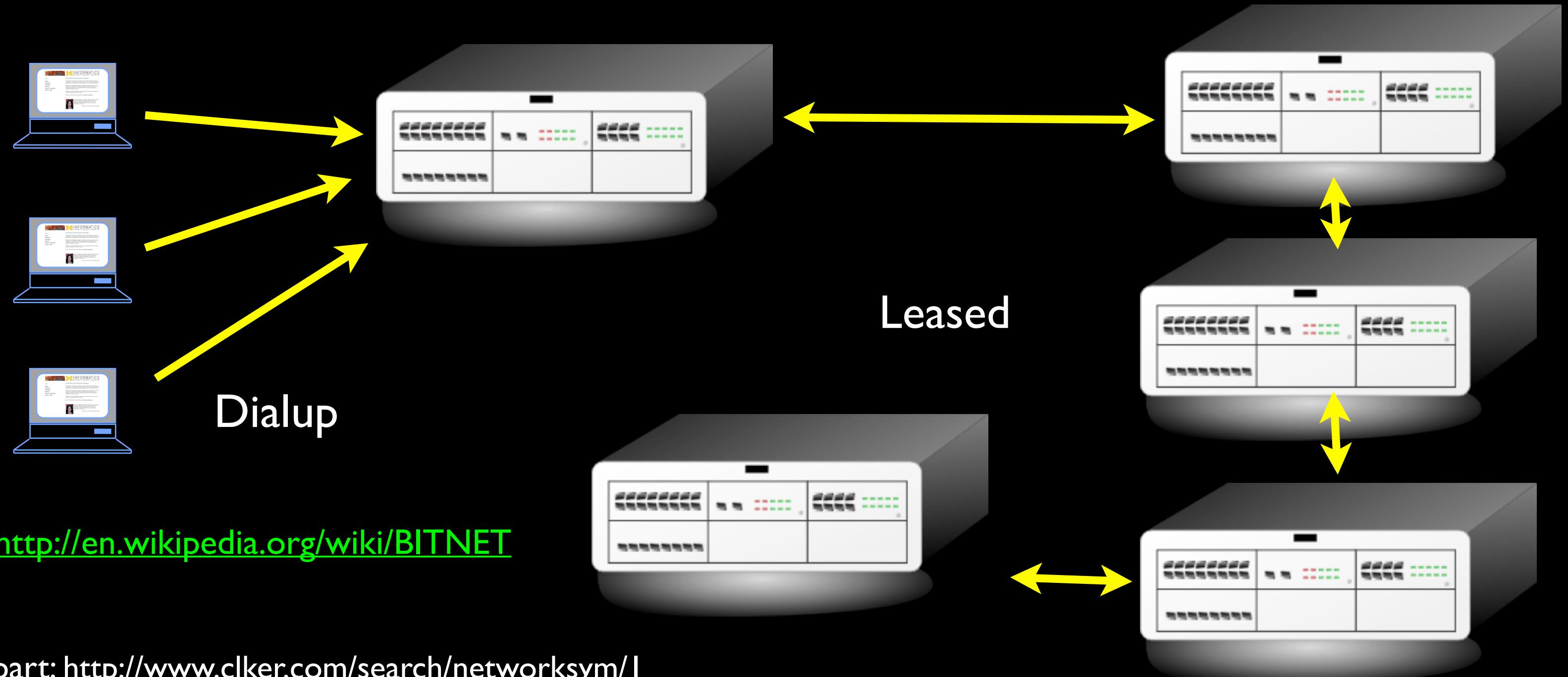
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BUT WHEN SHE TRACED THE  
KILLER'S IP ADDRESS... IT WAS  
IN THE 192.168/16 BLOCK!

GASP!

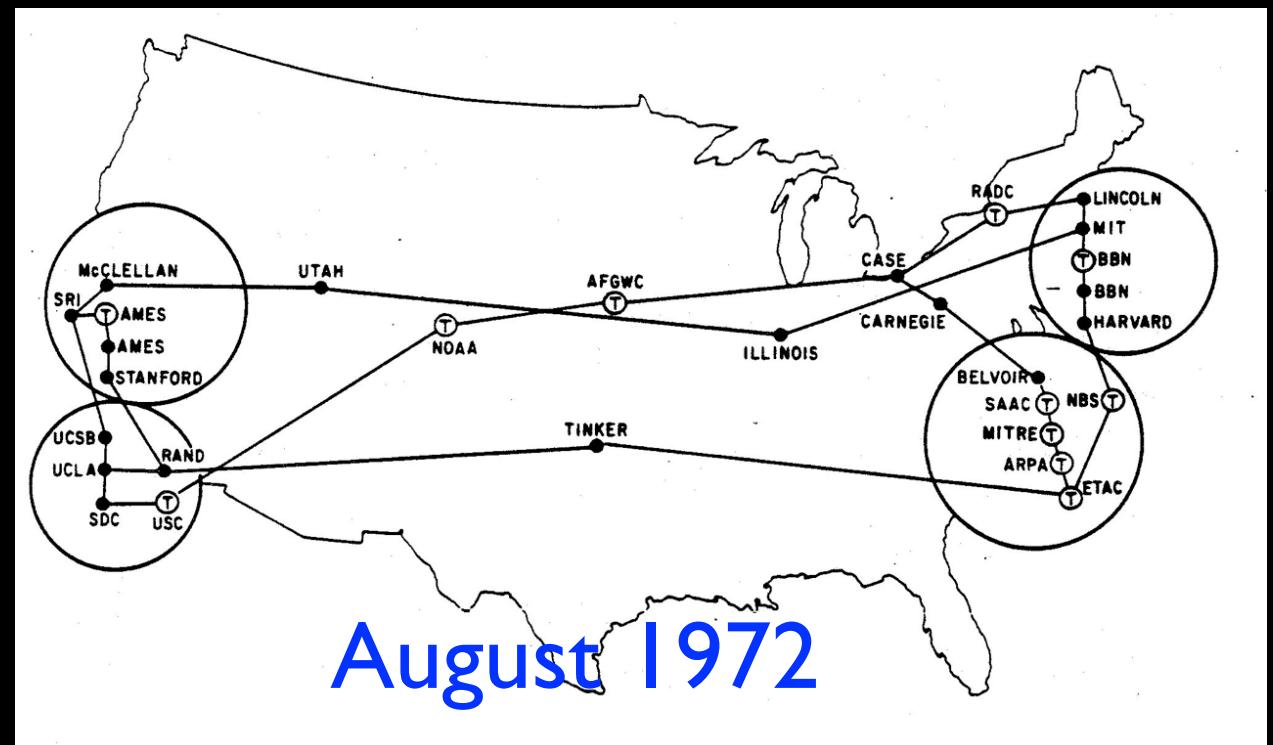
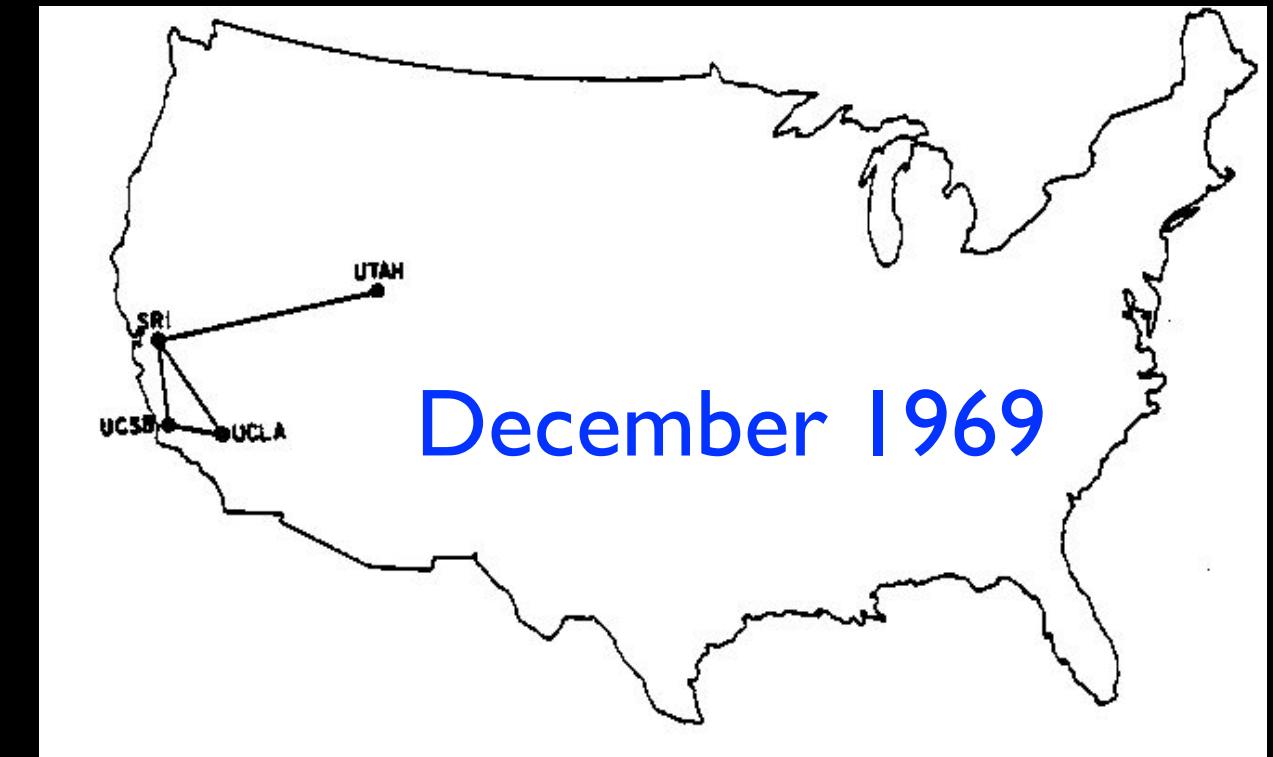
<http://xkcd.com/742/>

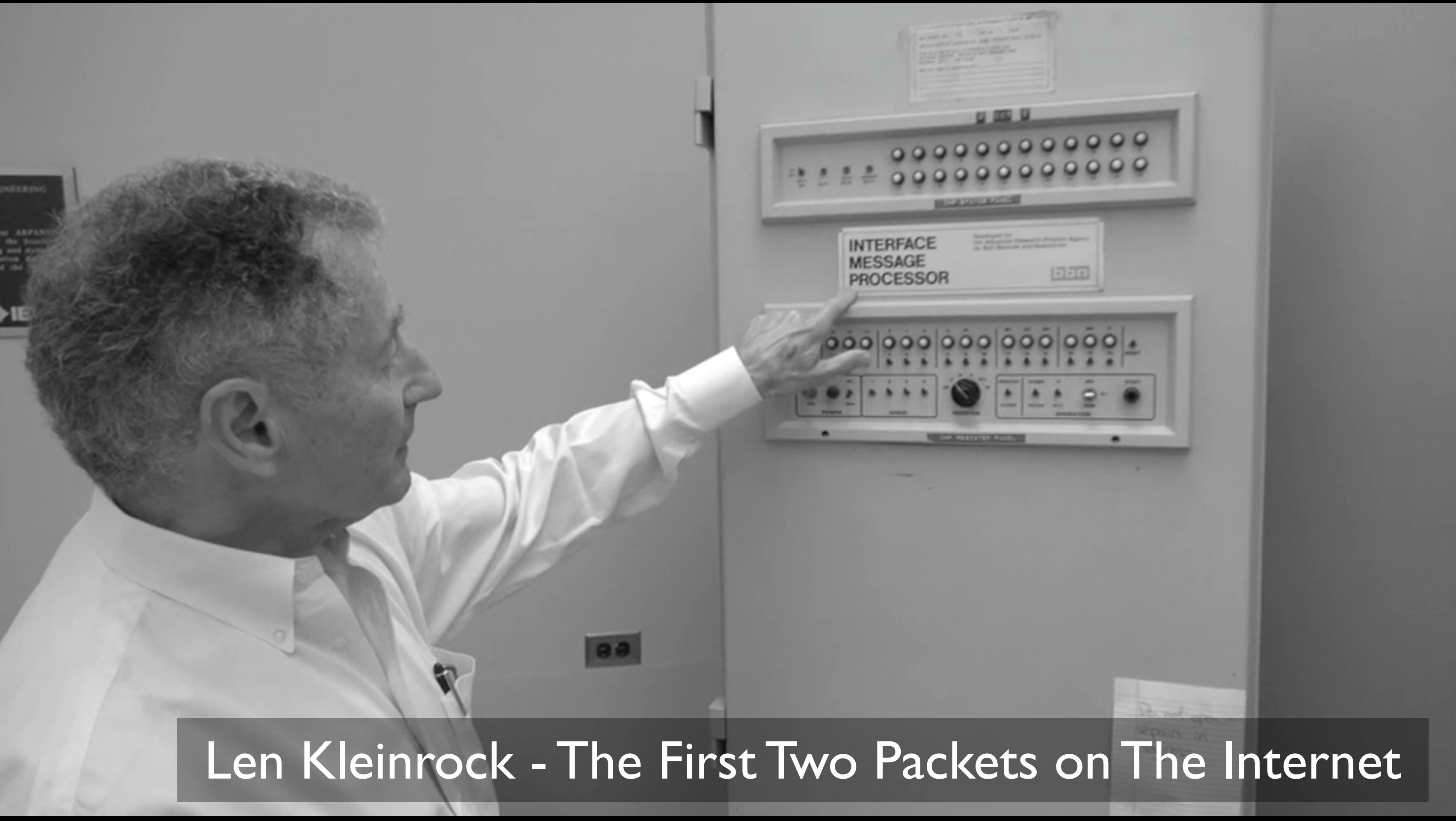
# Store and Forward Networking



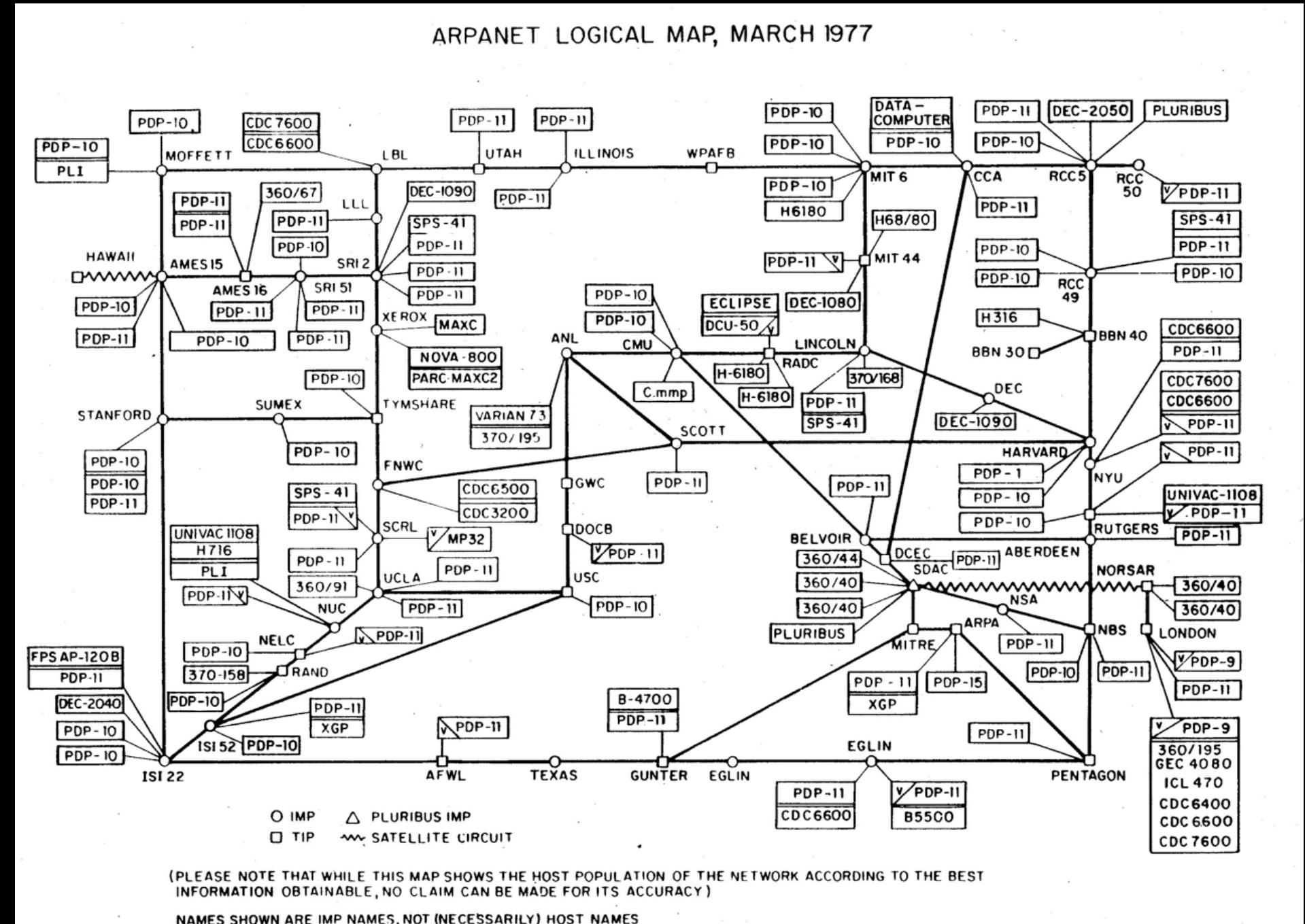
# Research Networks 1960-1980's

- How can we avoid having a direct connection between all pairs of computers?
- How to transport messages efficiently?
- How can we dynamically handle outages?





Len Kleinrock - The First Two Packets on The Internet



Heart, F., McKenzie, A., McQuillian, J., and Walden, D., ARPANET Completion Report,  
 Bolt, Beranek and Newman, Burlington, MA, January 4, 1978.  
<http://som.csudh.edu/fac/lpress/history/arpamaps/arpnetmar77.jpg>

# Efficient Message Transmission: Packet Switching

- Challenge: in a simple approach, like store-and-forward, large messages block small ones
- Break each message into **packets**
- Can allow the packets from a single message to travel over different paths, dynamically adjusting for use
- Use special-purpose computers, called **routers**, for the traffic control

# Packet Switching - Postcards

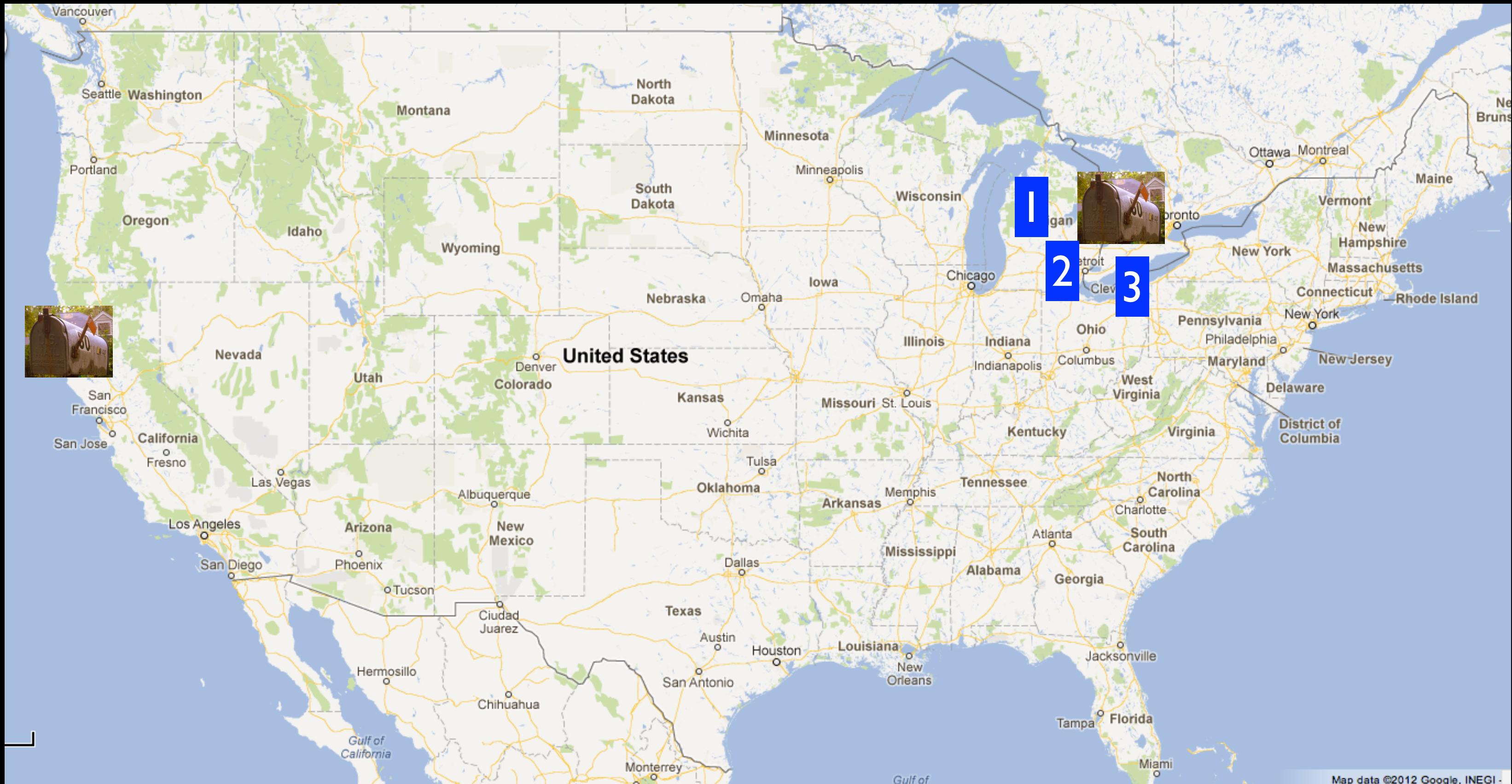
Hello there, have a nice day.

Hello ther (1, csev, daphne)

e, have a (2, csev, daphne)

nice day. (3, csev, daphne)



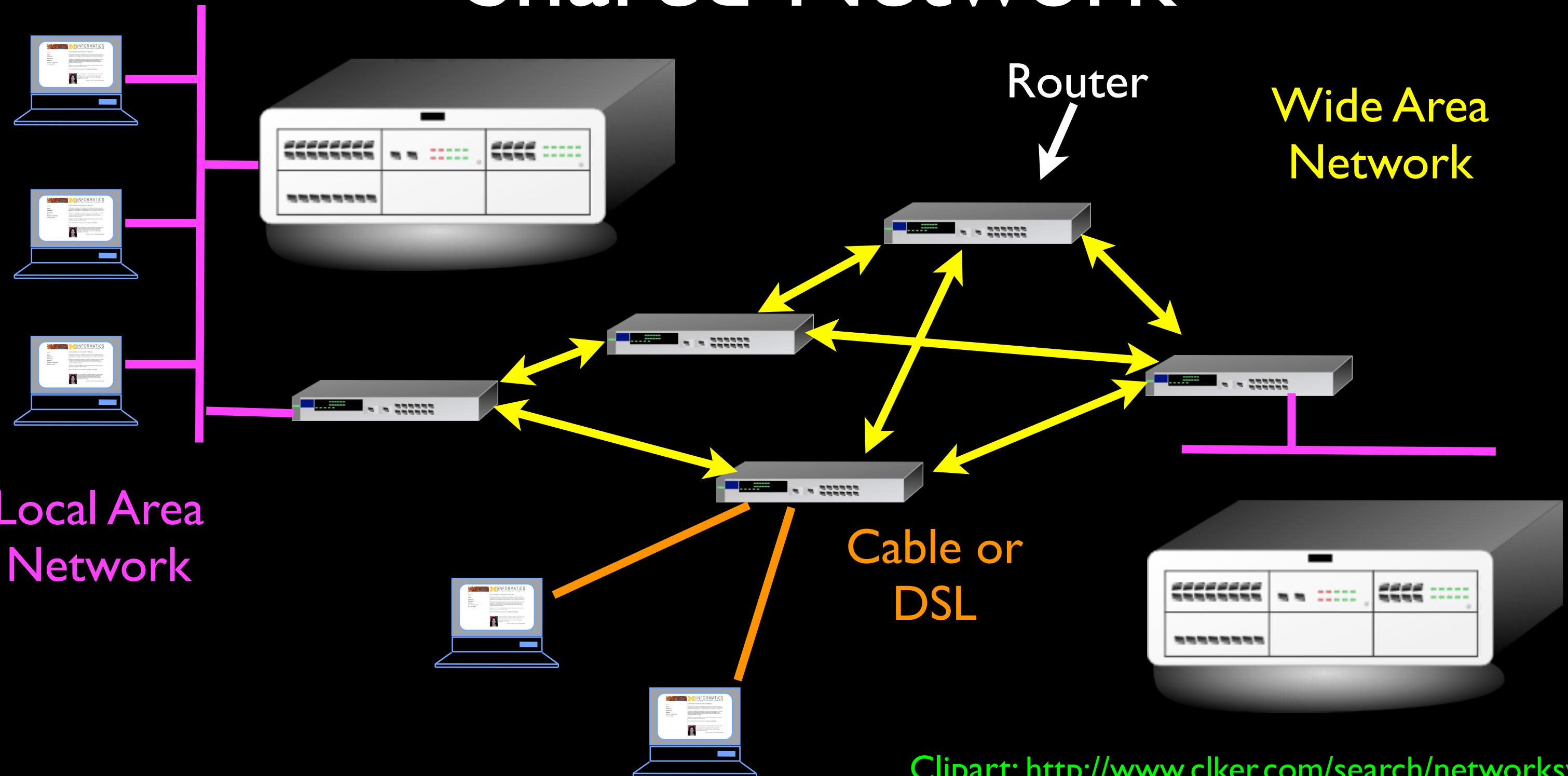


# Packet Switching - Postcards



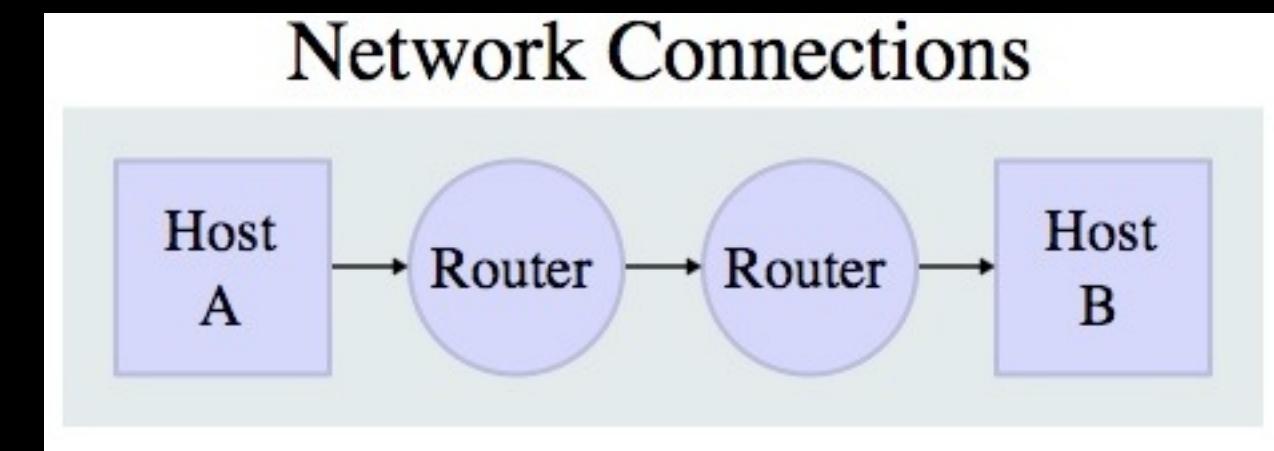
Hello there, have a nice day.

# Shared Network



# Shared Networks

- In order to keep cost low and the connections short geographically - data would be forwarded through several routers.
- Getting across the country usually takes about 10 “hops”
- Network designers continually add and remove links to “tune” their networks



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Layered Network Model

- A **layered** approach allows the problem of designing a network to be broken into more manageable sub problems
- Best-known model: **TCP/IP**—the “Internet Protocol Suite”
- There was also a 7 layer **OSI**: Open System Interconnection Model

Application Layer  
Web, E-Mail, File Transfer

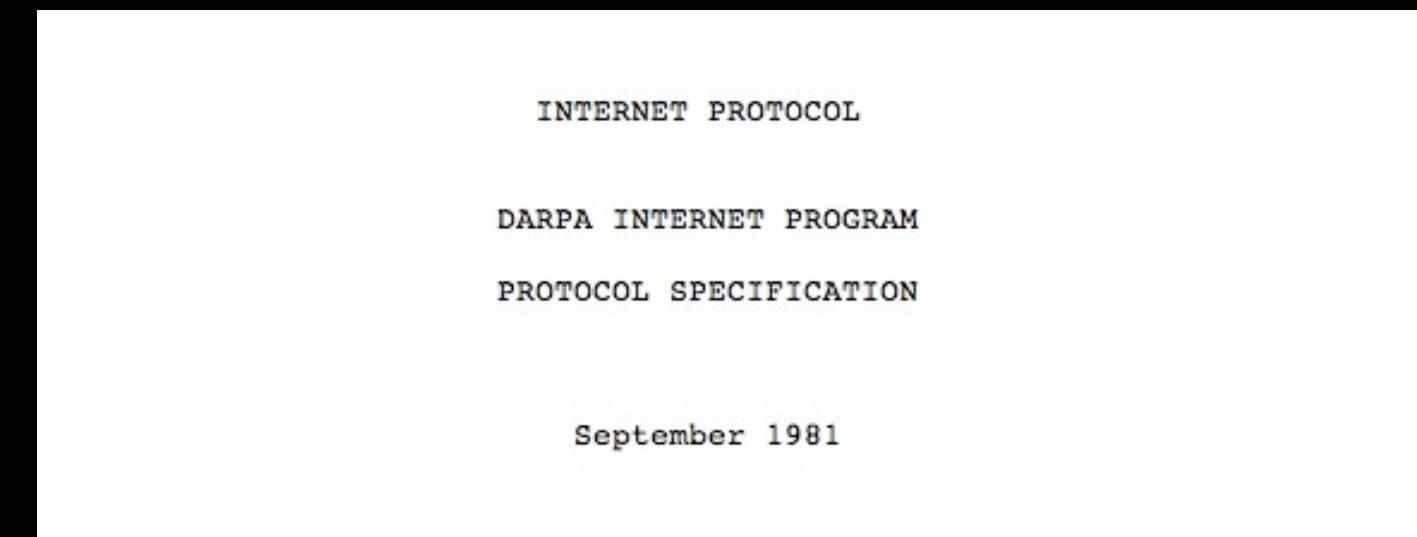
Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Internet Standards

- The standards for all of the Internet protocols (inner workings) are developed by an organization
- Internet Engineering Task Force (IETF)
- [www.ietf.org](http://www.ietf.org)
- Standards are called “RFCs” - “Request for Comments”



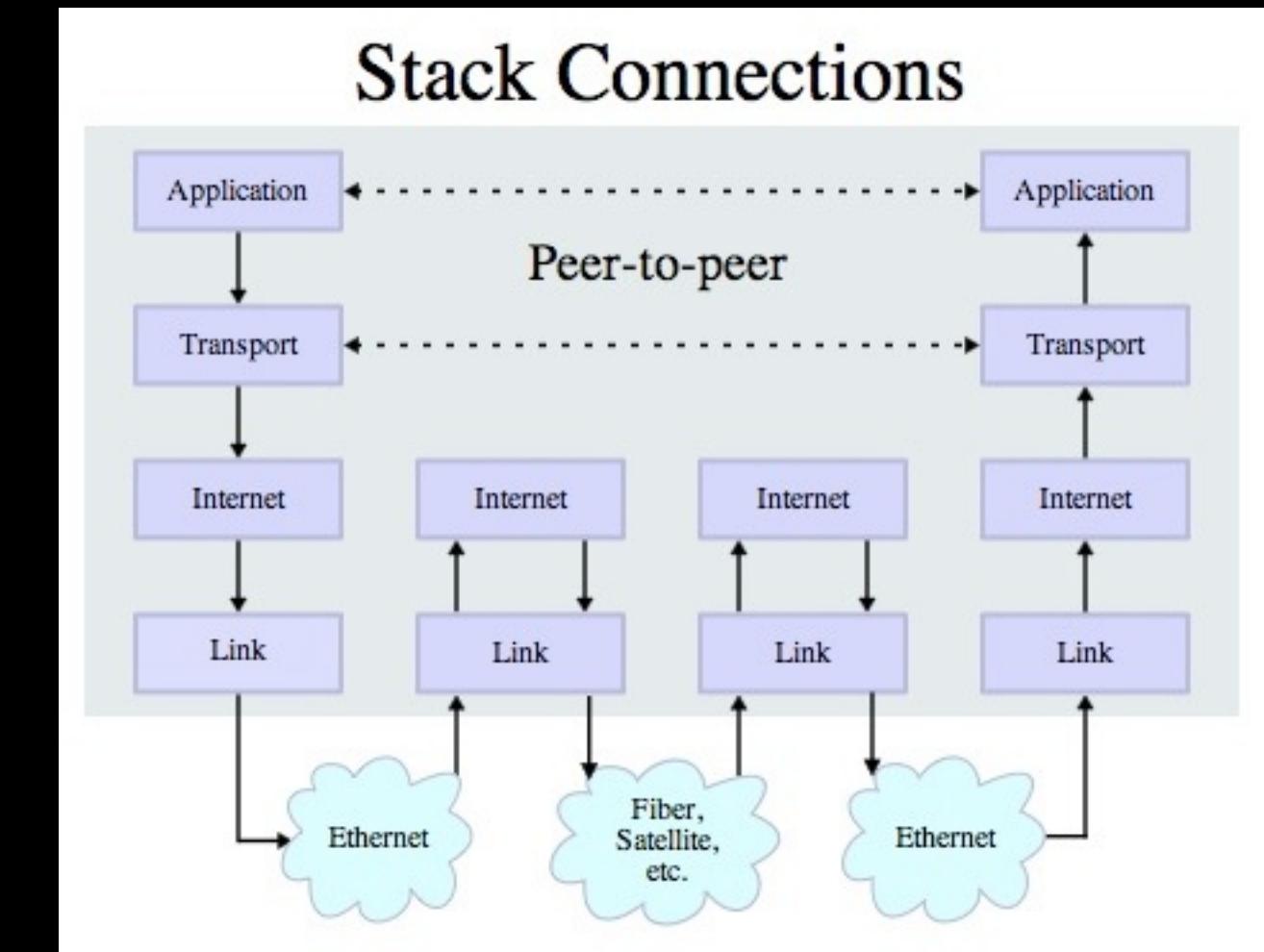
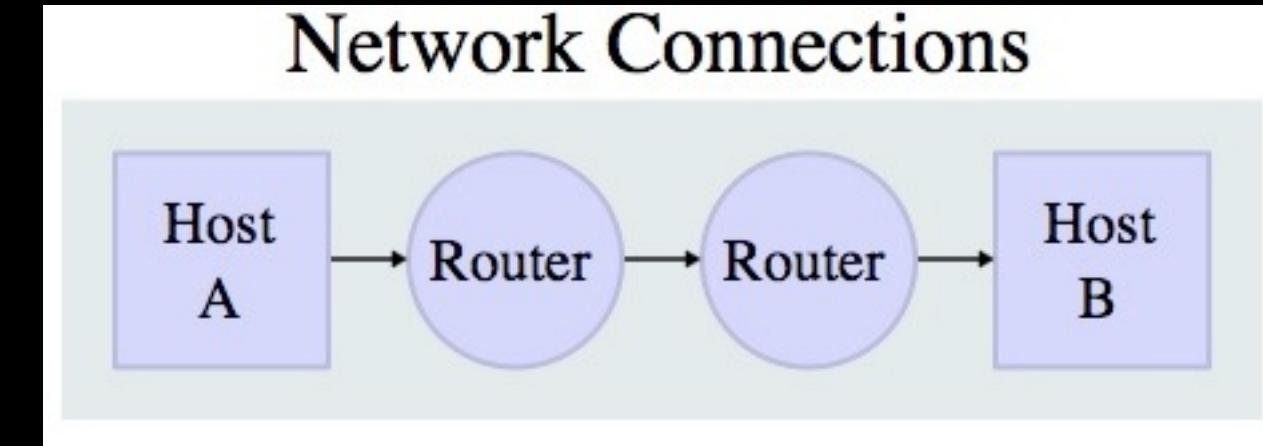
The internet protocol treats each internet datagram as an independent entity unrelated to any other internet datagram. There are no connections or logical circuits (virtual or otherwise).

The internet protocol uses four key mechanisms in providing its service: Type of Service, Time to Live, Options, and Header Checksum.

Source: <http://tools.ietf.org/html/rfc791>

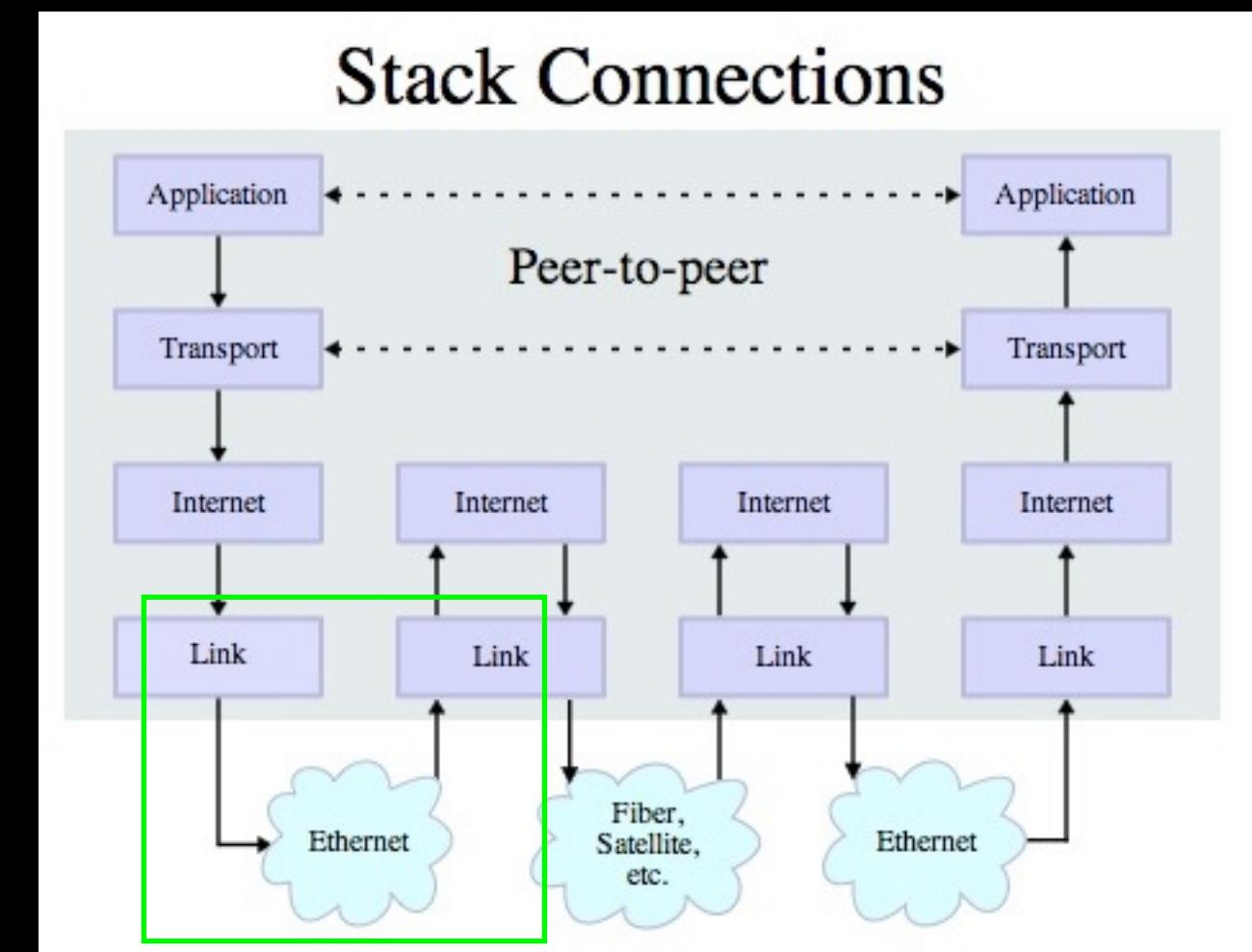
# Layered Architecture

- The Physical and Internet Layers are like trucks and trains - they haul stuff and get it to the right loading dock - it takes multiple steps
- The Transport layer checks to see if the trucks made it and send the stuff again if necessary



# Link Layer (aka Physical Layer)

- As your data crosses the country may use a different physical medium for each “hop”
- Wire, Wireless, Fiber Optic, etc.
- The link is “one hop” - Is it up or down? Connected or not?
- Very narrow focus - no view at all of the “whole Internet”

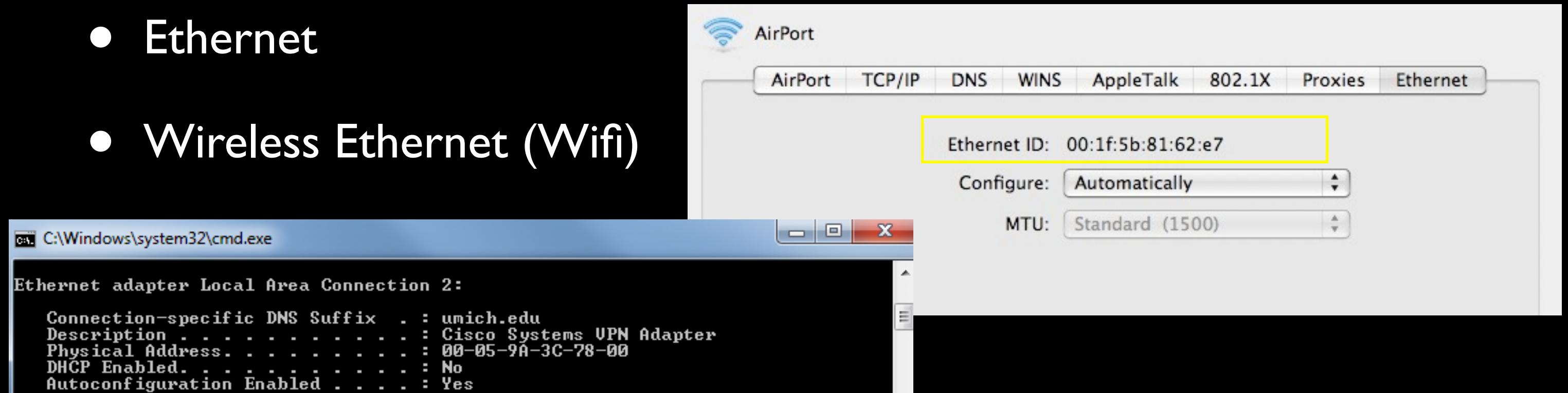


# Problems solved by the Link Layer

- How does data get pushed onto a link?
- How is the link shared?
- Common Link Technologies
  - Ethernet
  - WiFi
  - Cable modem
  - DSL
  - Satellite
  - Optical

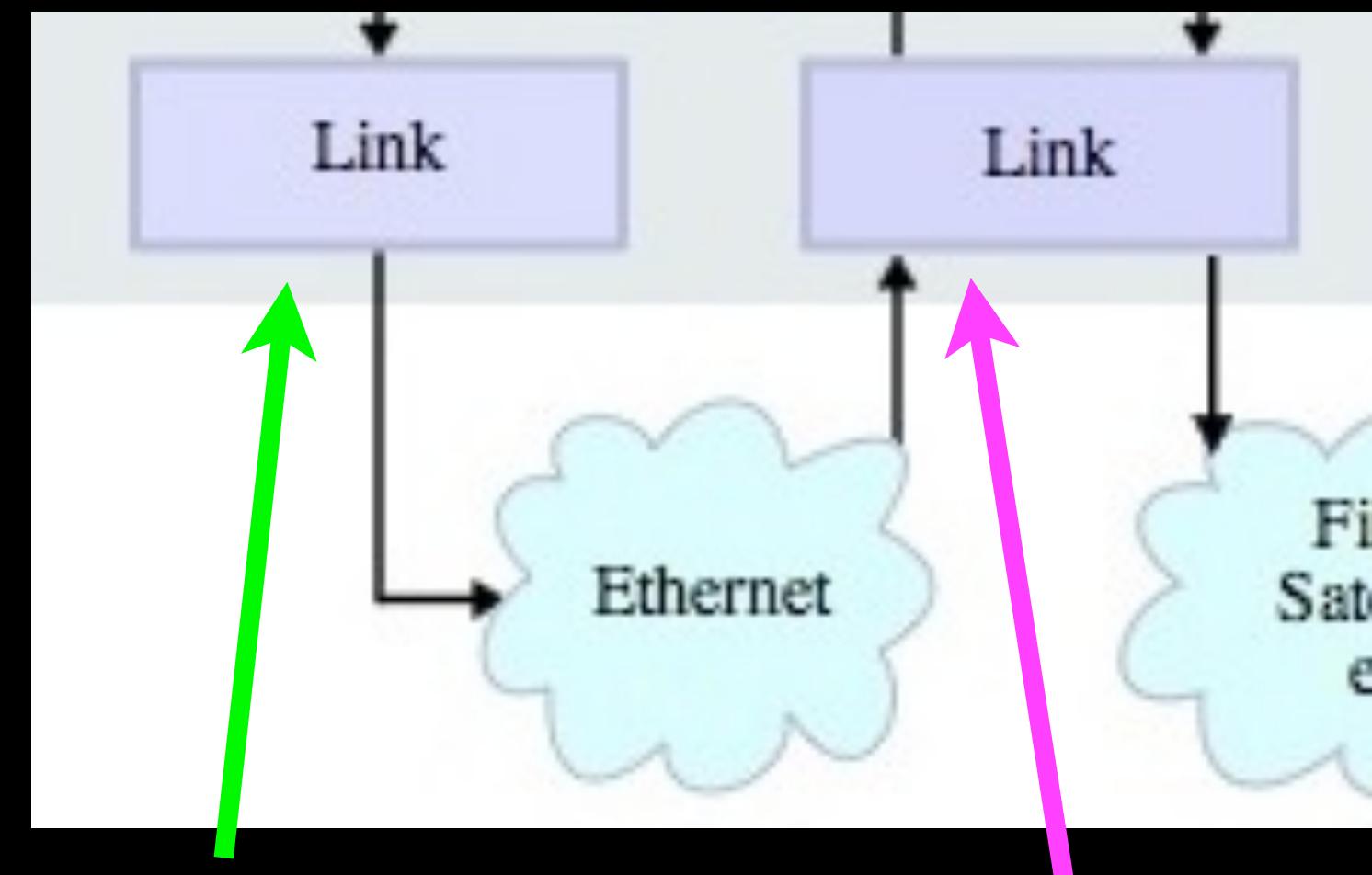
# Link Layer Addresses

- Many physical layer devices have addresses built in to them by the manufacturer
  - Ethernet
  - Wireless Ethernet (Wifi)



# Link Layer

- Physical addresses are to allow systems to identify themselves on the ends of a single link
- Physical addresses go no farther than one link
- Sometimes links like Wifi and Wired Ethernet are shared with multiple computers

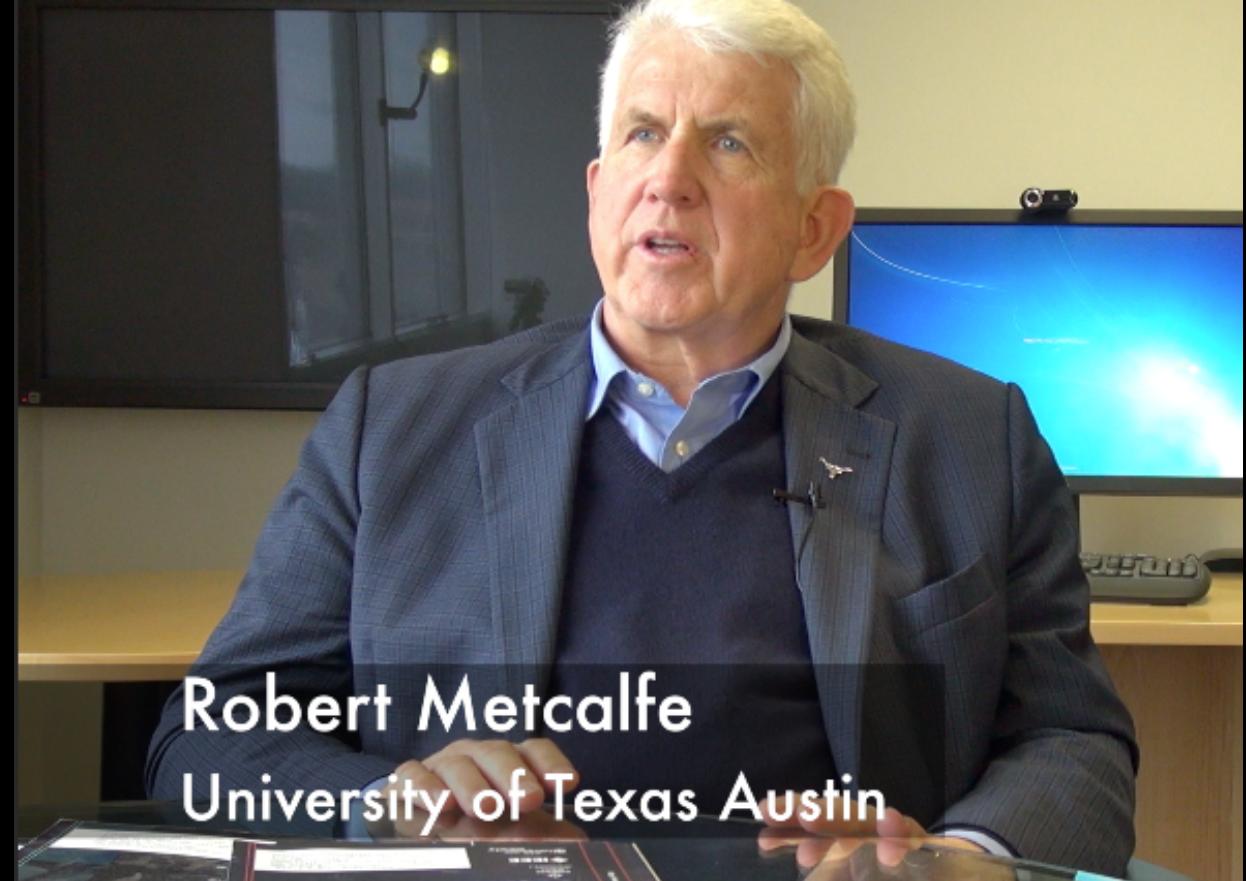


# Sharing Nicely - Avoiding Chaos

- **CSMA/CD Carrier Sense  
Multiple Access with  
Collision Detection**
- **To avoid garbled messages,  
systems must observe  
“rules” (Protocols)**
- **Ethernet rules are simple**
  - Wait for silence
  - Begin Transmitting data
  - Listen for your own data
  - If you cannot hear your own data clearly, assume a collision, stop and wait before trying again
  - Each system waits a different amount of time to avoid “too much politeness”

# Ethernet

- Invented at PARC (Xerox)
- The first Local-Area-Network
- Connected PC's to laser printers
- Inspired by an earlier wireless network called Aloha from the University of Hawaii



# **Internetwork Layer (IP)**

[http://en.wikipedia.org/wiki/Internet\\_Protocol](http://en.wikipedia.org/wiki/Internet_Protocol)

<http://en.wikipedia.org/wiki/Traceroute>

<http://en.wikipedia.org/wiki/Ping>

**Application Layer**  
Web, E-Mail, File Transfer

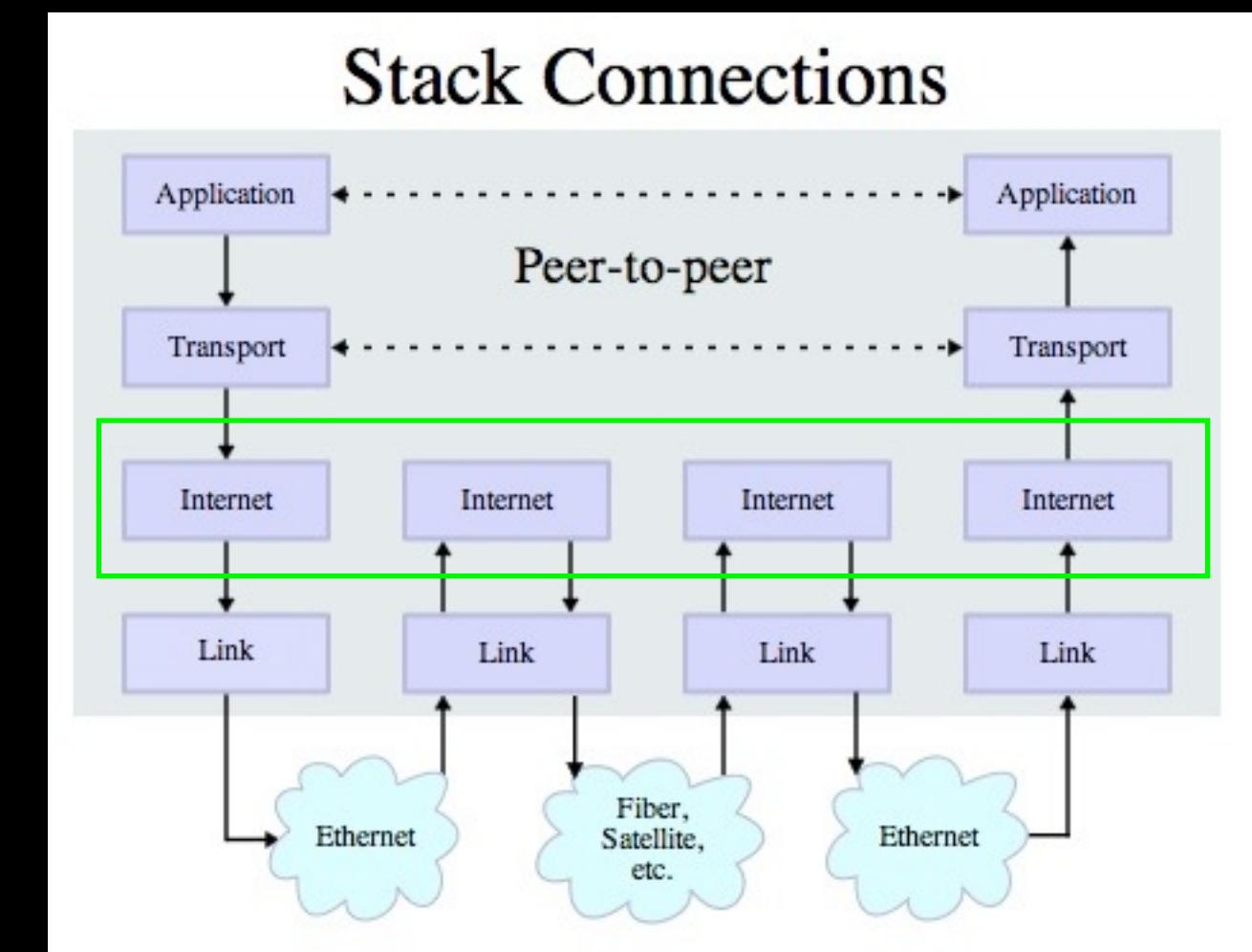
**Transport Layer (TCP)**  
Reliable Connections

**Internetwork Layer (IP)**  
Simple, Unreliable

**Link Layer (Ethernet, WiFi)**  
Physical Connections

# Internet Protocol Layer

- Goal: Gets your data from this computer to the other computer half way across the world
- Each router knows about nearby routers
- IP Is best effort - it is OK to drop data if things go bad...



# IP Addresses

- The IP address is the worldwide number which is associated with one particular workstation or server
- Every system which will send packets directly out across the Internet must have a unique IP address
- IP addresses are based on where station is connected
- IP addresses are not controlled by a single organization - address ranges are assigned
- They are like phone numbers – they get reorganized once in a great while

What Is My IP Address? - Lookup IP, Hide IP, Change IP, Trace IP and more... - MOZILLA FIREFOX

File Edit View History Bookmarks Tools Help

IP http://whatismyipaddress.com/ what's my IP address

Most Visited Getting Started Latest Headlines Readability

IP What Is My IP Address? - Lookup IP, ...

WhatIsMyIPAddress.com Dedicated to IP address discussion

My IP Change IP Hide IP IP Lookup Trace Email Blacklist Check IP Tools FAQs Resources

Ads by Google IP Address Domain 99 Where My IP IP Routing Hiding IP

BOOKMARK

Google™ Custom Search Search

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**Comcast Business Internet**  
Get Comcast Business Class Today Find Local Offer For Your Business  
[www.ComcastBusinessService.com](http://www.ComcastBusinessService.com)

**IP Scan**  
Scan your Network IP Free Network IP Scan.  
[IP.Scan.Qualys.com](http://IP.Scan.Qualys.com)

**Ip Address Providers**  
Voice, Data, T-1, Mobile, Web Host. Get price quote. Special packages.  
[grow.cbeyond.net](http://grow.cbeyond.net)

**Reverse Email Lookup**  
1) Type in Email Address. 2) Get Owner Name & Info.  
[EmailFinder.com](http://EmailFinder.com)

**Embarq is now CenturyLink**  
Try One Voice & Data

What is my IP address?

Your IP address is **68.42.65.147**  
(Now detects many proxy servers)

Free Trial.

Map Satellite Hybrid Terrain

Wisconsin Michigan Ontario Peterborough Barrie Oshawa

Madison Milwaukee Flint Kitchener London Hamilton

Rockford Elgin Racine Grand Rapids Davenport Peoria

Chicago Fort Wayne Indianapolis Dayton Lima

Cleveland Akron Pittsburgh Columbus Cincinnati

Ohio Pennsylvania Erie Harrisburg

West Virginia Arlington

Ann Arbor, Michigan United States

IP Address Location: Ann Arbor, Michigan United States  
Read about [GeoLocation accuracy](#).

Tools

IP Lookup Blacklist Check Trace Email Visual Traceroute Traceroute

68.42.65.147 Trace Now

IP Lookup now shows ISP, Organization, Proxy Status, and Connection Type!

CLICK TO READ MORE

# IP Address Format

- Four numbers with dots - each number 1-255 (32 bits)
- Kind of like phone numbers with an “area code”
- The prefix of the address is “which network”
- While the data is traversing the Internet - all that matters is the network number

(734) 764 1855

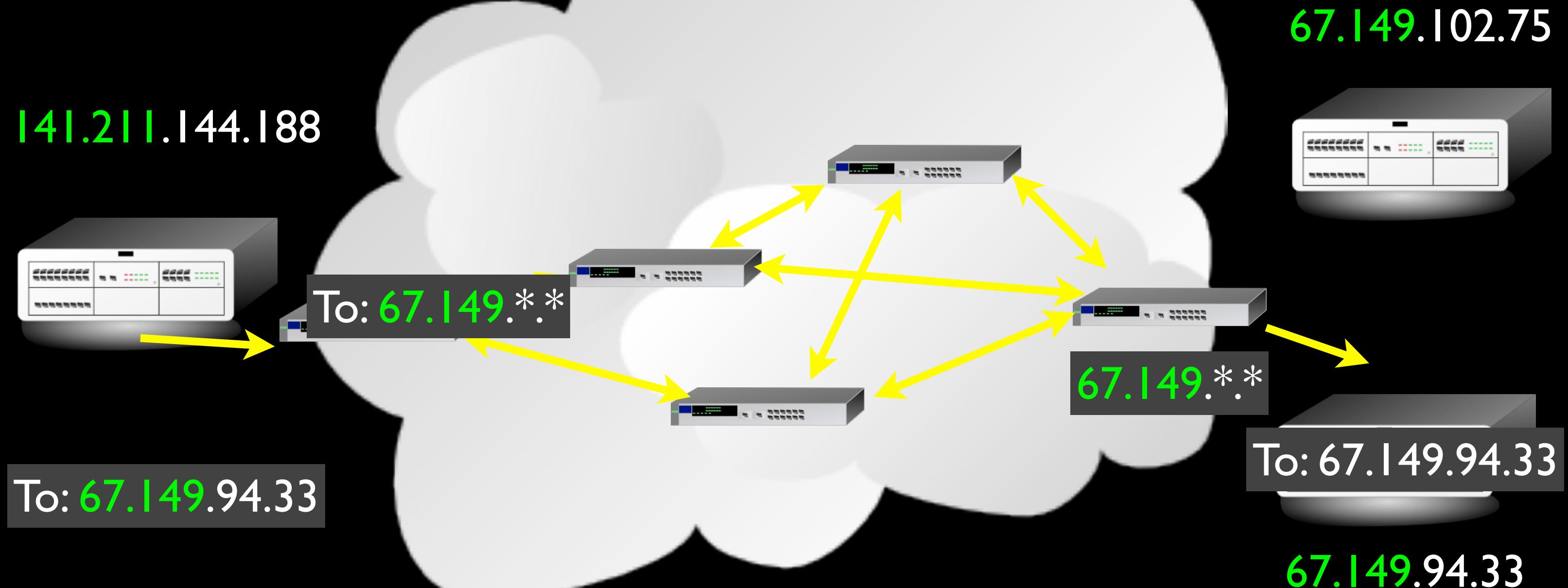
Area code

141.211.144.188

Network  
Number

141.211.\*.\*

While in the network, all  
that matters is the  
Network number.

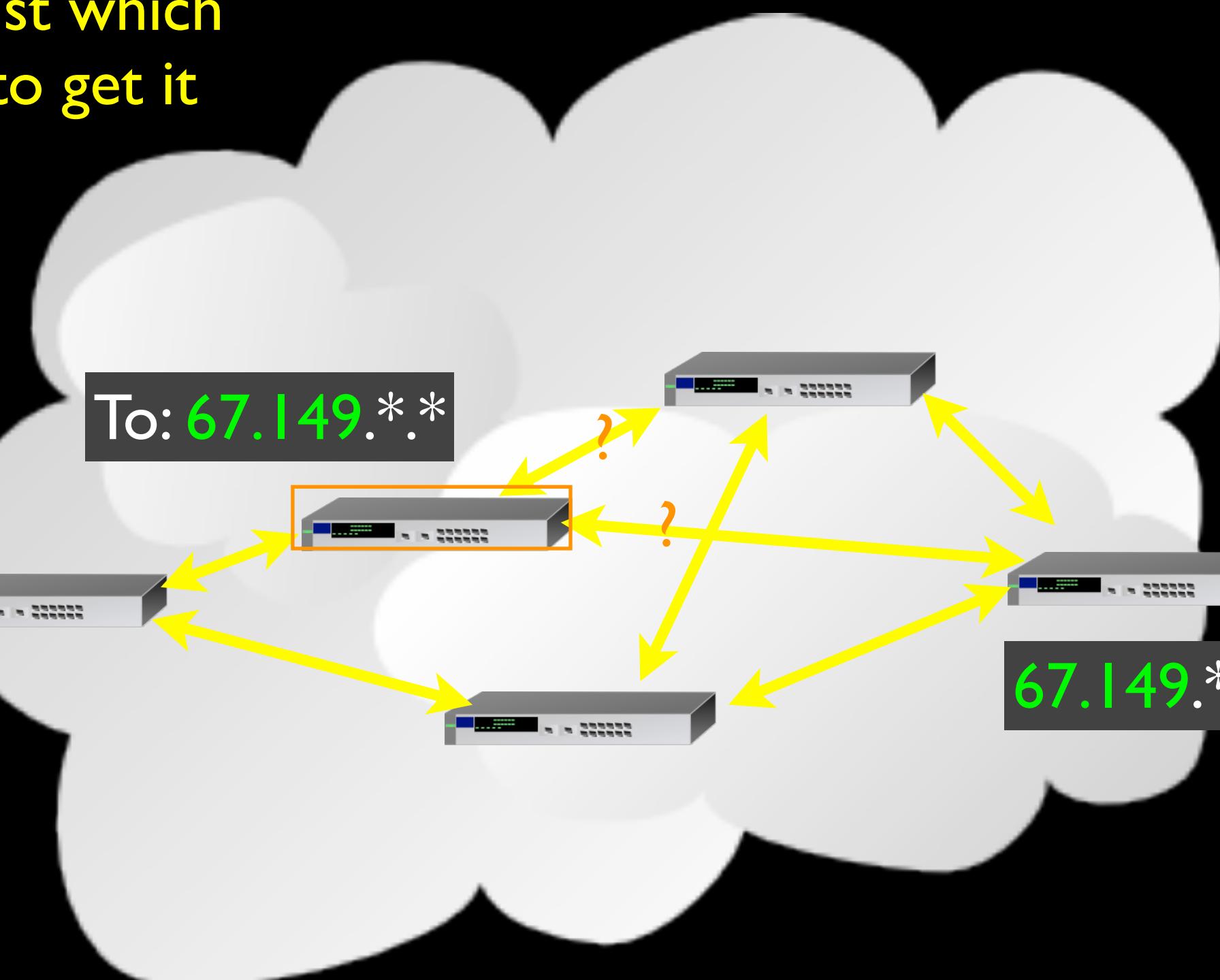


No single router knows the whole network - just which way to send data to get it “closer”

141.211.144.188



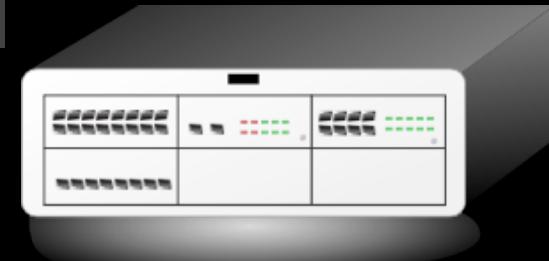
To: 67.149.\*.\*



67.149.102.75



67.149.\*.\*

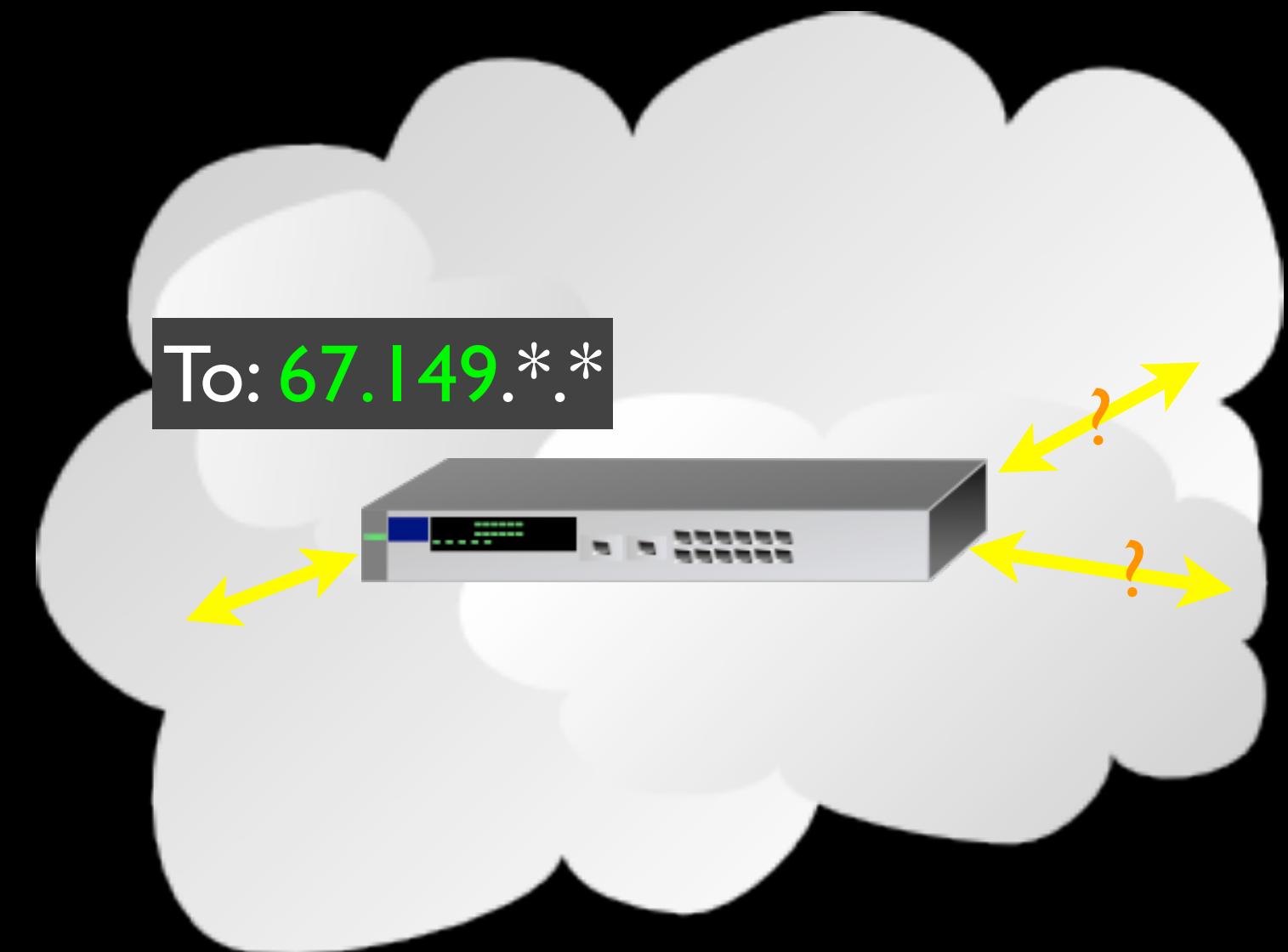


67.149.94.33

# Router Tables

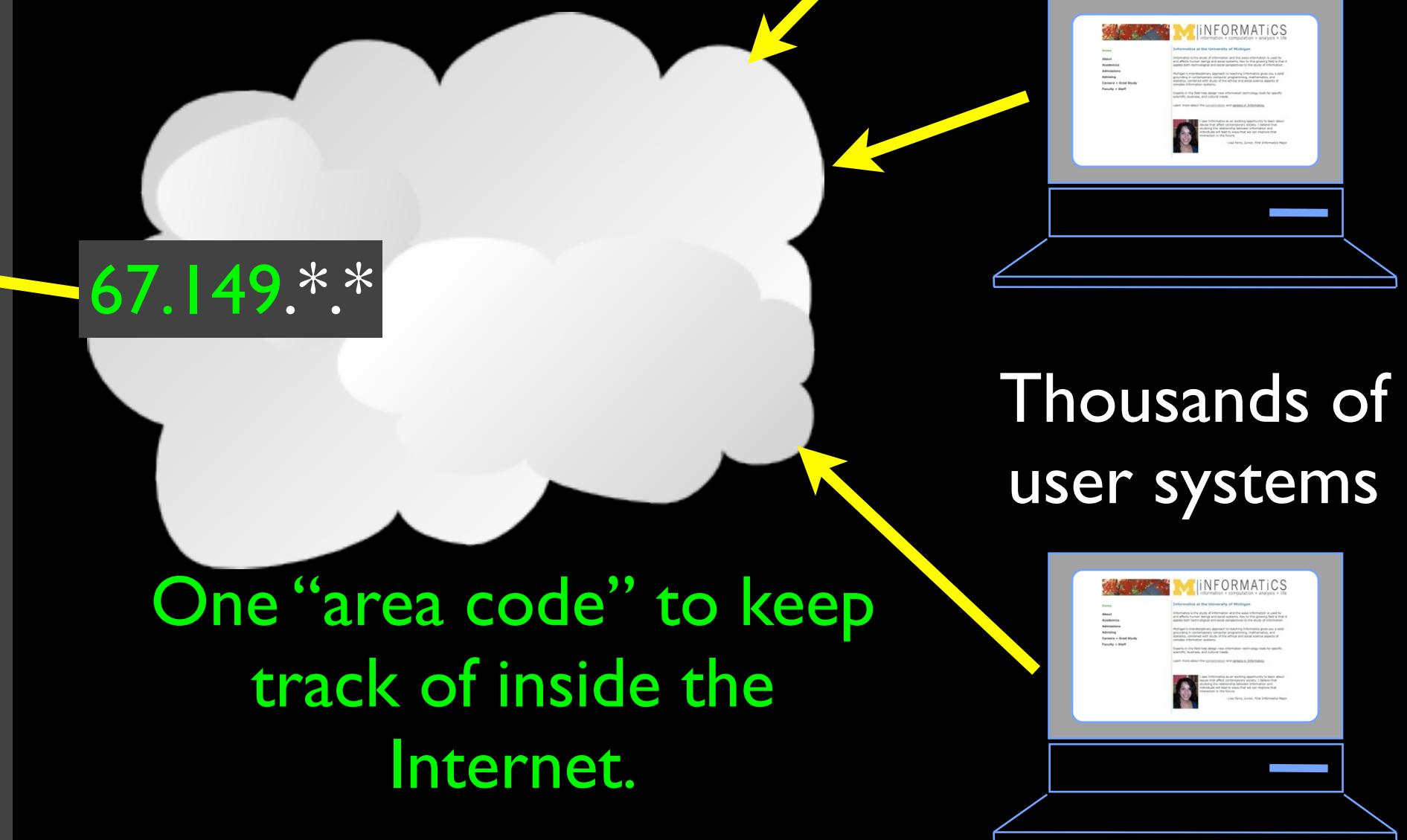
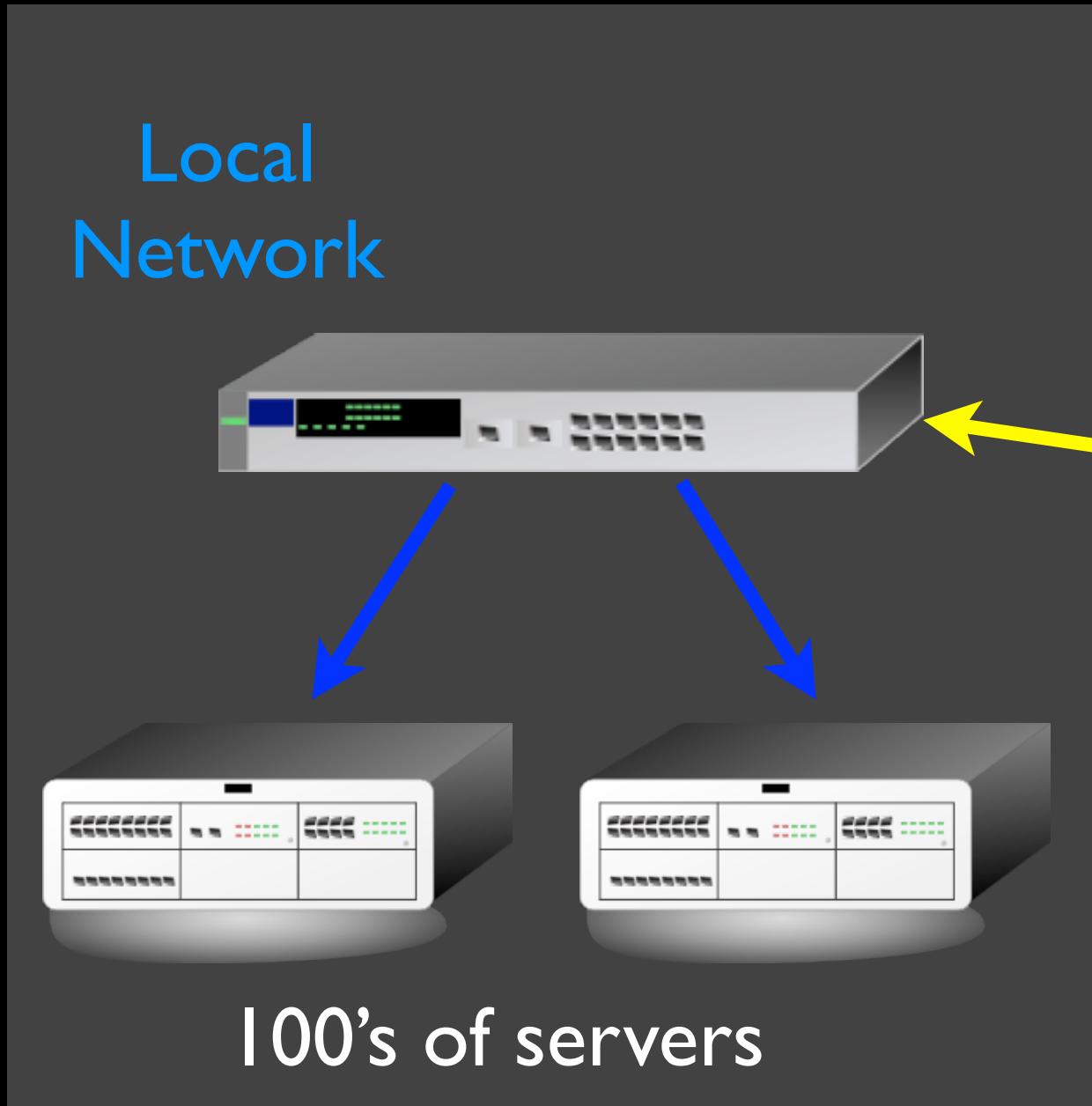
Lists of where to send packets, based on  
destination network address;  
bandwidth on adjacent links;  
traffic on adjacent links;  
state of neighbor nodes (up or not);  
...

Updated dynamically  
Routers “ask each other” for information

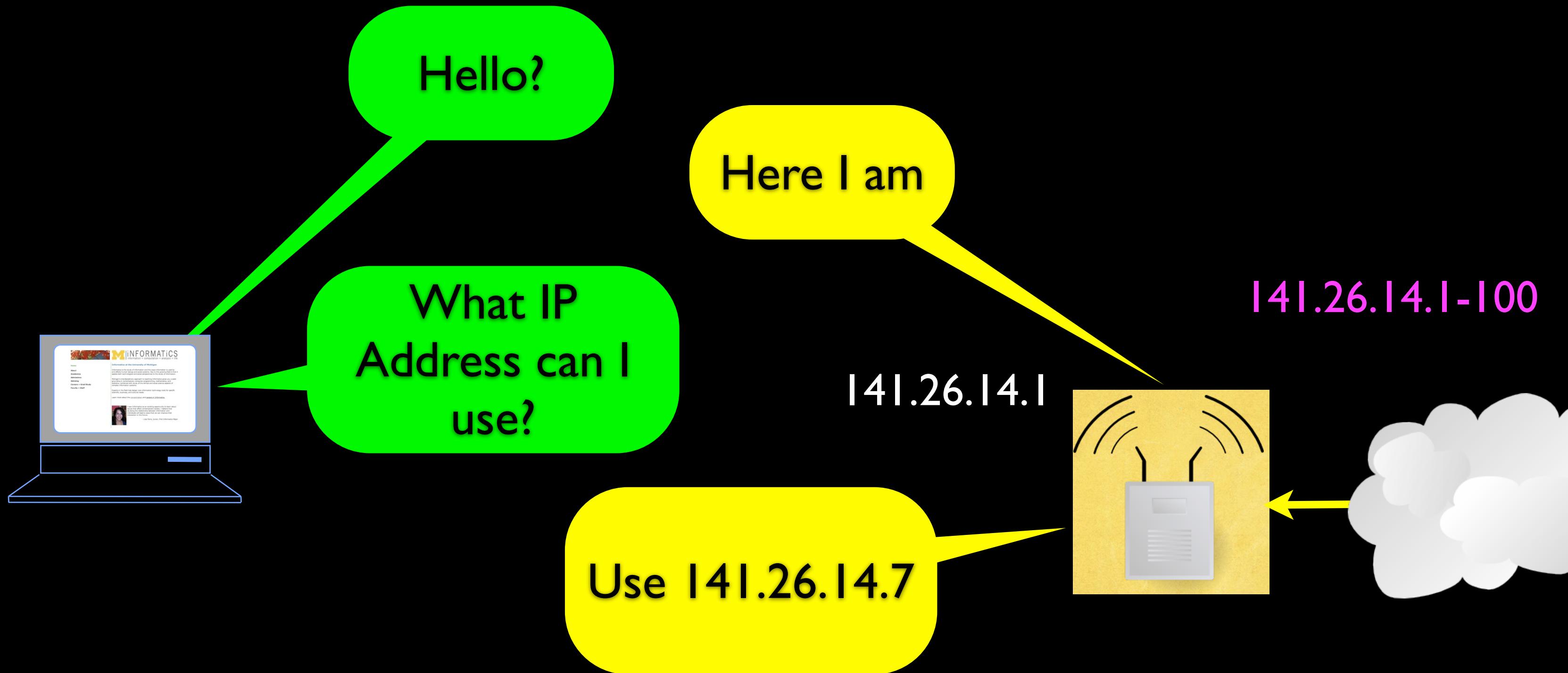


# IP Is Simple

Thousands of network connections.  
Billions of bytes of data per seconds.



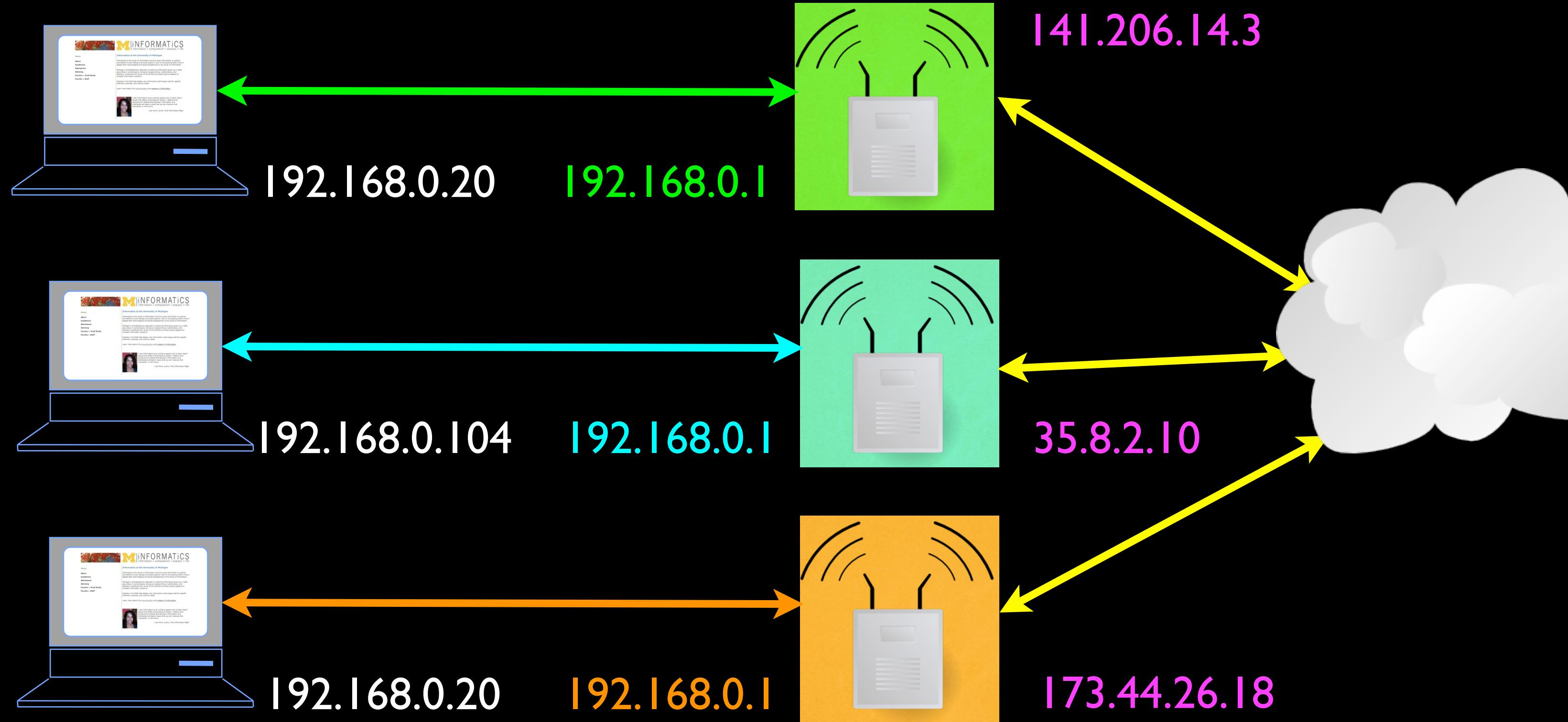
# DHCP = Dynamic Host Configuration Protocol



# Non-Routable Addresses

- A typical home router does Network Address Translation (NAT)
- Your ISP gives your home router a real global routable address
- Your router gives out local addresses in a special range (192.168.\*.\* )
- The router maps remote addresses for each connection you make from within your home network

[http://en.wikipedia.org/wiki/Network\\_address\\_translation](http://en.wikipedia.org/wiki/Network_address_translation)



NAT = Network Address Translation

Clipart: <http://www.clker.com/search/networksym/>

BUT WHEN SHE TRACED THE  
KILLER'S IP ADDRESS... IT WAS  
IN THE 192.168/16 BLOCK!

GASP!

<http://xkcd.com/742/>

# Peering into the Internet

- Most systems have a command that will reveal the route taken across the internet (traceroute on Mac and tracert on Windows)
- Each IP packet has a field called “Time to Live” - TTL
- The TTL is used to deal with loops in the network - normally if routers got confused and ended up with a loop - the network would clog up rapidly.



# How Traceroute Works

- Normal packets are sent with a Time to Live (TTL) of 255 hops
- Trace route sends a packet with TTL=1, TTL=2, ...
- So each packet gets part-way there and then gets dropped and traceroute gets a notification of where the drop happens
- This builds a map of the nodes that a packet visits when crossing the Internet.

# Traceroute

```
$ traceroute www.stanford.edu
```

```
traceroute to www5.stanford.edu (171.67.20.37), 64 hops max, 40 byte packets
```

1	141.211.203.252	(141.211.203.252)	1.390 ms	0.534 ms	0.490 ms
2	v-bin-seb.r-bin-seb.umnet.umich.edu	(192.122.183.61)	0.591 ms	0.558 ms	0.570 ms
3	v-bin-seb-i2-aa.merit-aa2.umnet.umich.edu	(192.12.80.33)	6.610 ms	6.545 ms	6.654 ms
4	192.122.183.30	(192.122.183.30)	7.919 ms	7.209 ms	7.122 ms
5	so-4-3-0.0.rtr.kans.net.internet2.edu	(64.57.28.36)	17.672 ms	17.836 ms	17.673 ms
6	so-0-1-0.0.rtr.hous.net.internet2.edu	(64.57.28.57)	31.800 ms	41.967 ms	31.787 ms
7	so-3-0-0.0.rtr.losa.net.internet2.edu	(64.57.28.44)	63.478 ms	63.704 ms	63.710 ms
8	hpr-lax-hpr--i2-newnet.cenic.net	(137.164.26.132)	63.093 ms	63.026 ms	63.384 ms
9	svl-hpr--lax-hpr-10ge.cenic.net	(137.164.25.13)	71.242 ms	71.542 ms	76.282 ms
10	oak-hpr--svl-hpr-10ge.cenic.net	(137.164.25.9)	72.744 ms	72.243 ms	72.556 ms
11	hpr-stan-ge--oak-hpr.cenic.net	(137.164.27.158)	73.763 ms	73.396 ms	73.665 ms
12	bbra-rtr.Stanford.EDU	(171.64.1.134)	73.577 ms	73.682 ms	73.492 ms
13	***				
14	www5.Stanford.EDU	(171.67.20.37)	77.317 ms	77.128 ms	77.648 ms

# Traceroute

```
$ traceroute www.msu.edu
```

```
traceroute to www.msu.edu (35.8.10.30), 64 hops max, 40 byte packets
```

```
1 141.211.203.252 (141.211.203.252) 2.644 ms 0.973 ms 14.162 ms
```

```
2 v-bin-seb.r-bin-seb.umnet.umich.edu (192.122.183.61) 1.847 ms 0.561 ms 0.496 ms
```

```
3 v-bin-seb-i2-aa.merit-aa2.umnet.umich.edu (192.12.80.33) 6.490 ms 6.499 ms 6.529 ms
```

```
4 lt-0-3-0x1.eq-chi2.mich.net (198.108.23.121) 8.096 ms 8.113 ms 8.103 ms
```

```
5 xe-0-0-0x23.msu6.mich.net (198.108.23.213) 7.831 ms 7.962 ms 7.965 ms
```

```
6 192.122.183.227 (192.122.183.227) 12.953 ms 12.339 ms 10.322 ms
```

```
7 cc-tl-gel-23.net.msu.edu (35.9.101.209) 9.522 ms 9.406 ms 9.817 ms
```

```
8 * * *
```

# Traceroute

```
$ traceroute www.pku.edu.cn
```

```
traceroute: Warning: www.pku.edu.cn has multiple addresses; using 162.105.129.104  
traceroute to www.pku.edu.cn (162.105.129.104), 64 hops max, 40 byte packets
```

```
1 141.211.203.252 (141.211.203.252) 1.228 ms 0.584 ms 0.592 ms  
2 v-bin-seb.r-bin-seb.umnet.umich.edu (192.122.183.61) 0.604 ms 0.565 ms 0.466 ms  
3 v-bin-seb-i2-aa.merit-aa2.umnet.umich.edu (192.12.80.33) 7.511 ms 6.641 ms 6.588 ms
```

```
4 192.122.183.30 (192.122.183.30) 12.078 ms 6.989 ms 7.619 ms
```

Michigan

```
5 192.31.99.133 (192.31.99.133) 7.666 ms 8.953 ms 17.861 ms
```

Tennessee

```
6 192.31.99.170 (192.31.99.170) 59.275 ms 59.273 ms 59.108 ms
```

```
7 134.75.108.209 (134.75.108.209) 173.614 ms 173.552 ms 173.333 ms
```

```
8 134.75.107.10 (134.75.107.10) 256.760 ms 134.75.107.18 (134.75.107.18) 256.574 ms 256.53
```

```
9 202.112.53.17 (202.112.53.17) 256.761 ms 256.801 ms 256.688 ms
```

Seoul

```
10 202.112.61.157 (202.112.61.157) 257.416 ms 257.960 ms 257.747 ms
```

```
11 202.112.53.194 (202.112.53.194) 256.827 ms 257.068 ms 256.962 ms
```

```
12 202.112.41.202 (202.112.41.202) 256.800 ms 257.053 ms 256.933 ms
```

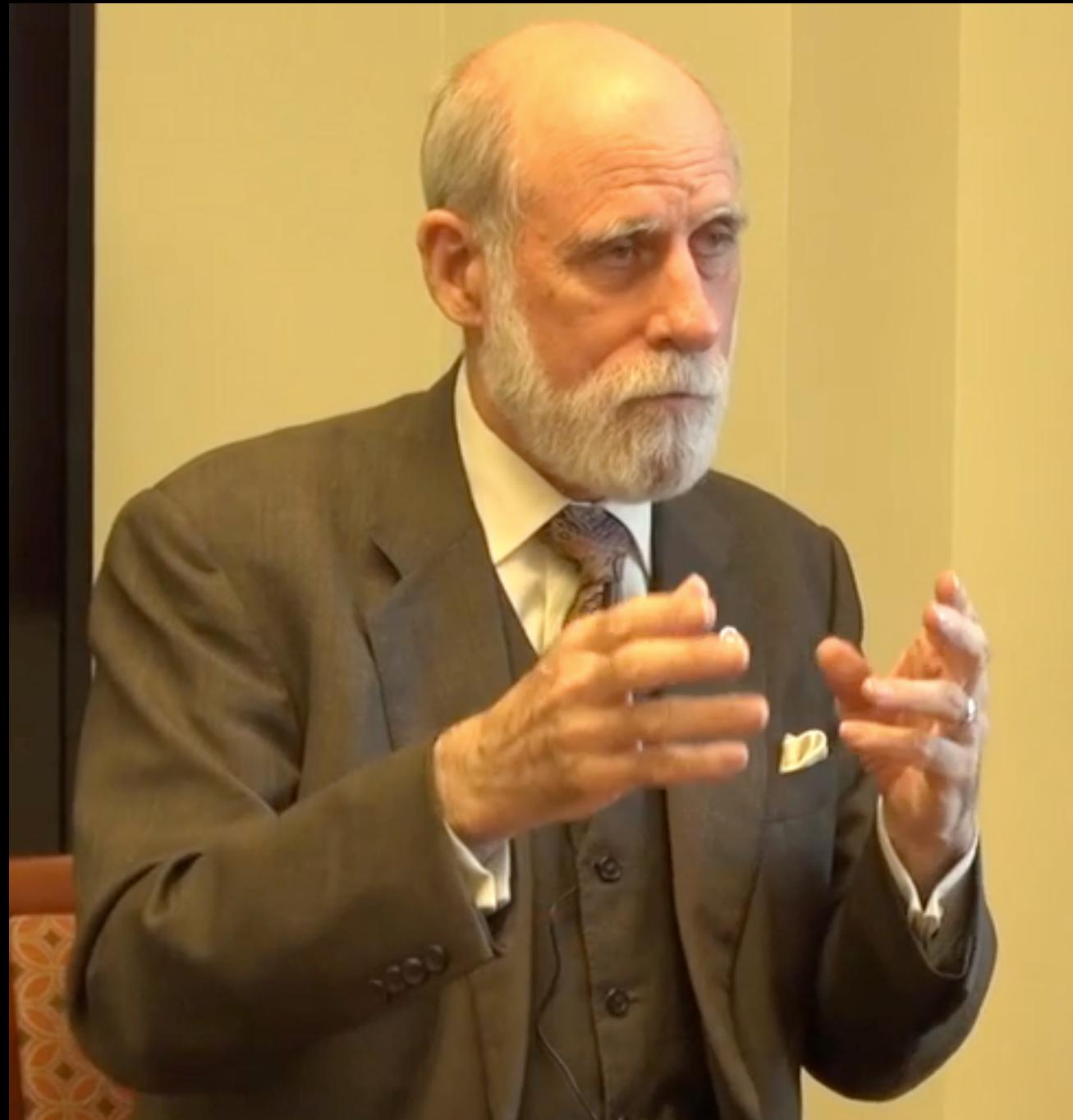
Beijing

# The perfect is the enemy of the good

*Le mieux est l'ennemi du bien. --Voltaire*

- **IP Does:** Best effort to get data across bunch of hops from one network to another network
- **IP Does Not:** Guarantee delivery - if things go bad - the data can vanish
- Best effort to keep track of the good and bad paths for traffic - tries to pick better paths when possible
- This makes it fast and scalable to very large networks - and ultimately “reliable” because it does not try to do too much

# Vint Cerf:A Brief History of Packets



- Instrumental in the design and development of the ARPANET
- Vint was a graduate student as the notions of packet-switching were emerging across academia

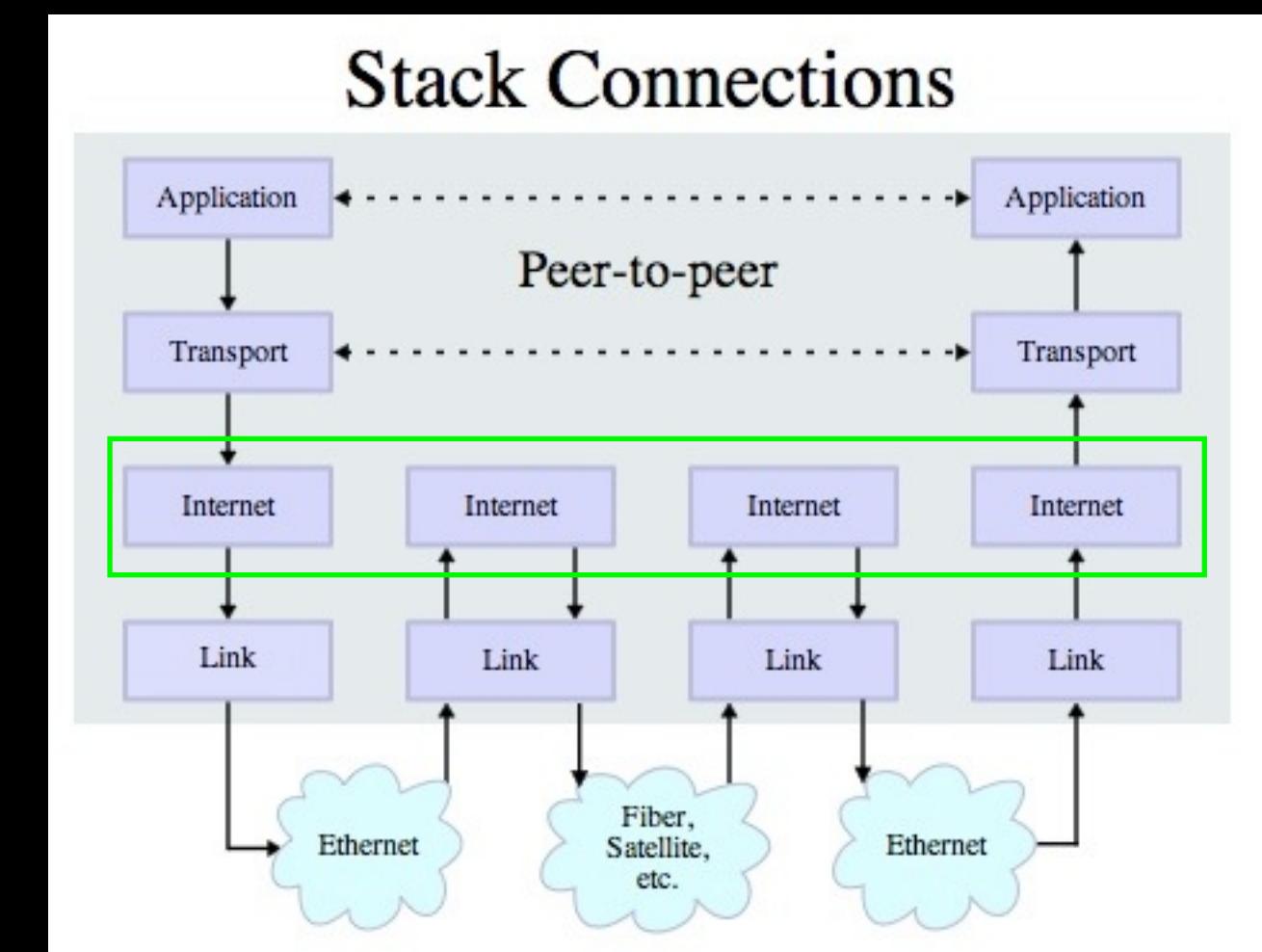
# Domain Name System

The Domain Name System  
convert user-friendly  
names, like

**www.umich.edu**

to network-friendly IP  
addresses, like

**141.211.32.166**



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

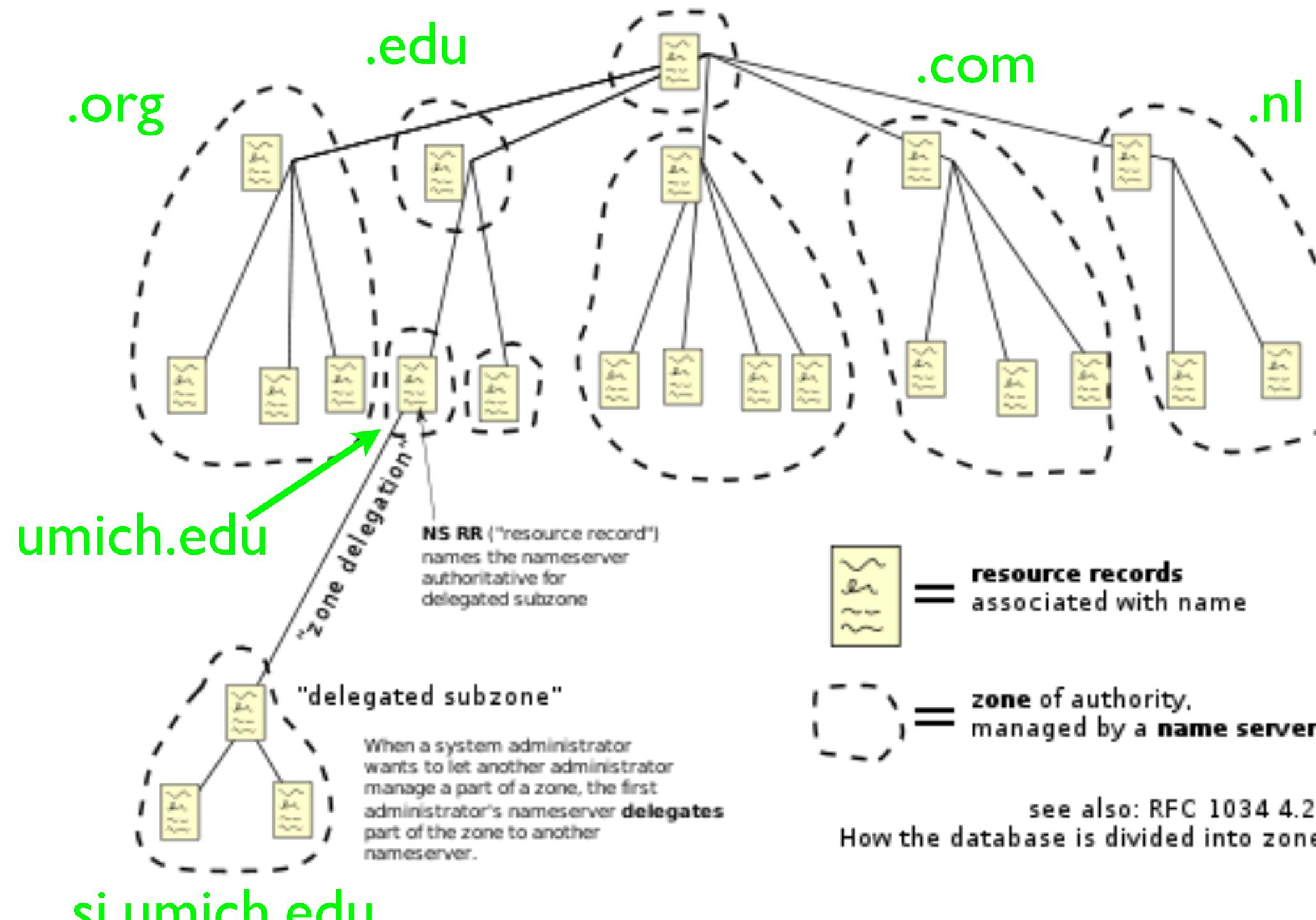
# Domain Name System

- Numeric addresses like 141.211.63.45 are great for Internet routers but lousy for people
- Each campus ends up with a lot of networks (141.211.\*.\* , 65.43.21.\* )
- Sometimes (rarely) the IP address numbers get reorganized
- When servers physically move they need new IP addresses

# DNS: Internet Address Book

- The Domain Name System is a big fast distributed database of Internet names to Internet “phone numbers”
- IP Addresses reflect technical “geography”
  - 141.211.63.44 - read left to right like a phone number
  - Domain names reflect organizational structure
    - www.si.umich.edu - read right to left like postal address
    - 2455 North Quad, Ann Arbor, MI, USA, Earth

## Domain Name Space



# Internet Layer (IP)

[http://en.wikipedia.org/wiki/Internet\\_Protocol](http://en.wikipedia.org/wiki/Internet_Protocol)

<http://en.wikipedia.org/wiki/Traceroute>

<http://en.wikipedia.org/wiki/Ping>

Application Layer  
Web, E-Mail, File Transfer

Transport Layer (TCP)  
Reliable Connections

Internet Layer (IP)  
Simple, Scalable, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Transport Layer

[http://en.wikipedia.org/wiki/  
Transmission\\_Control\\_Protocol](http://en.wikipedia.org/wiki/Transmission_Control_Protocol)

Application Layer  
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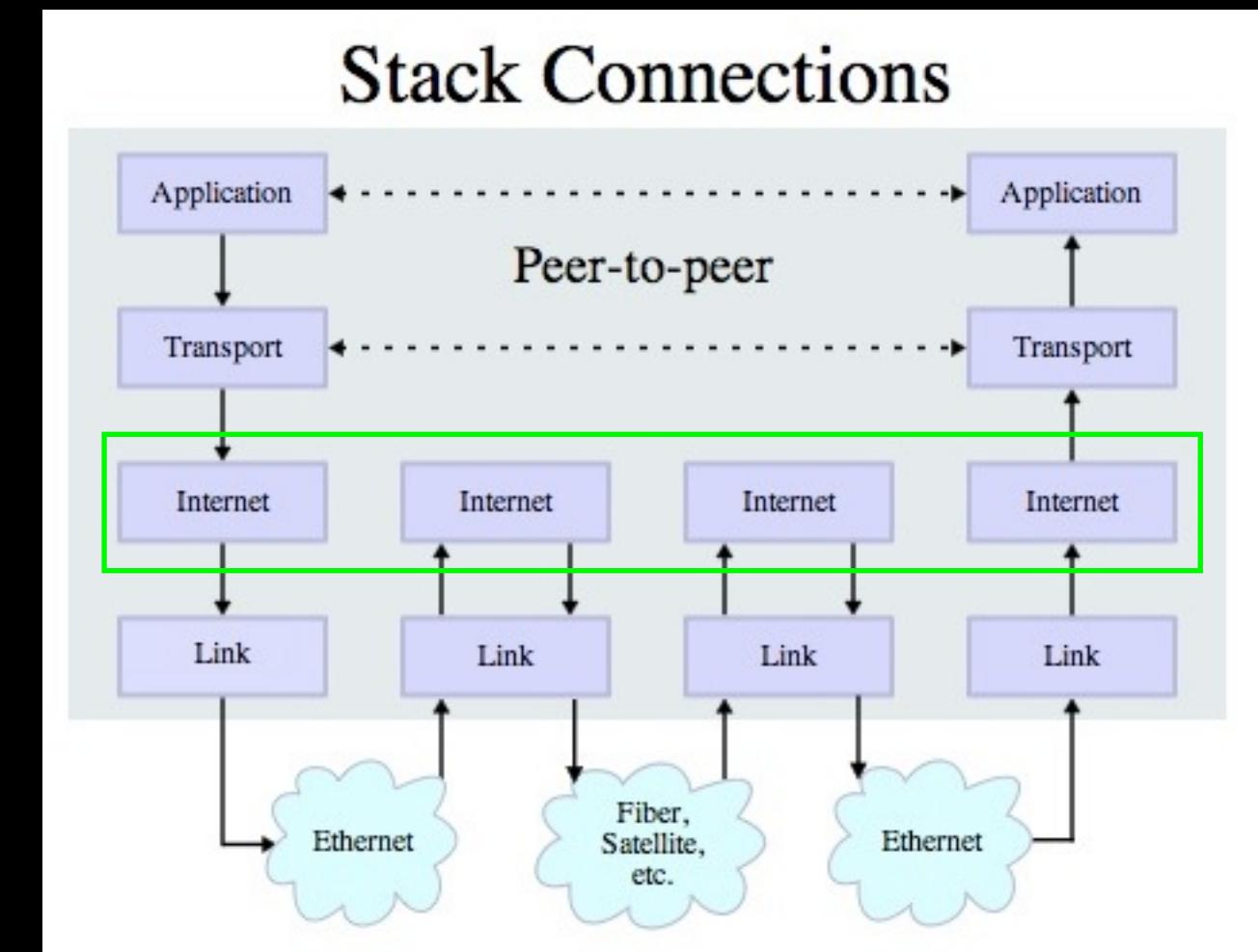
Link Layer (Ethernet, WiFi)  
Physical Connections

# Review: The Magic of IP

- What it does - Tries to get one packet across a 5-20 of hops from one network to another network
- Keeps track of the good and bad paths for traffic - tries to pick better paths when possible
- But no guarantee of delivery - if things go bad - the data vanishes
- This makes it fast and scalable - and ultimately “reliable” because it does not try to do too “everything”

# Internet Protocol

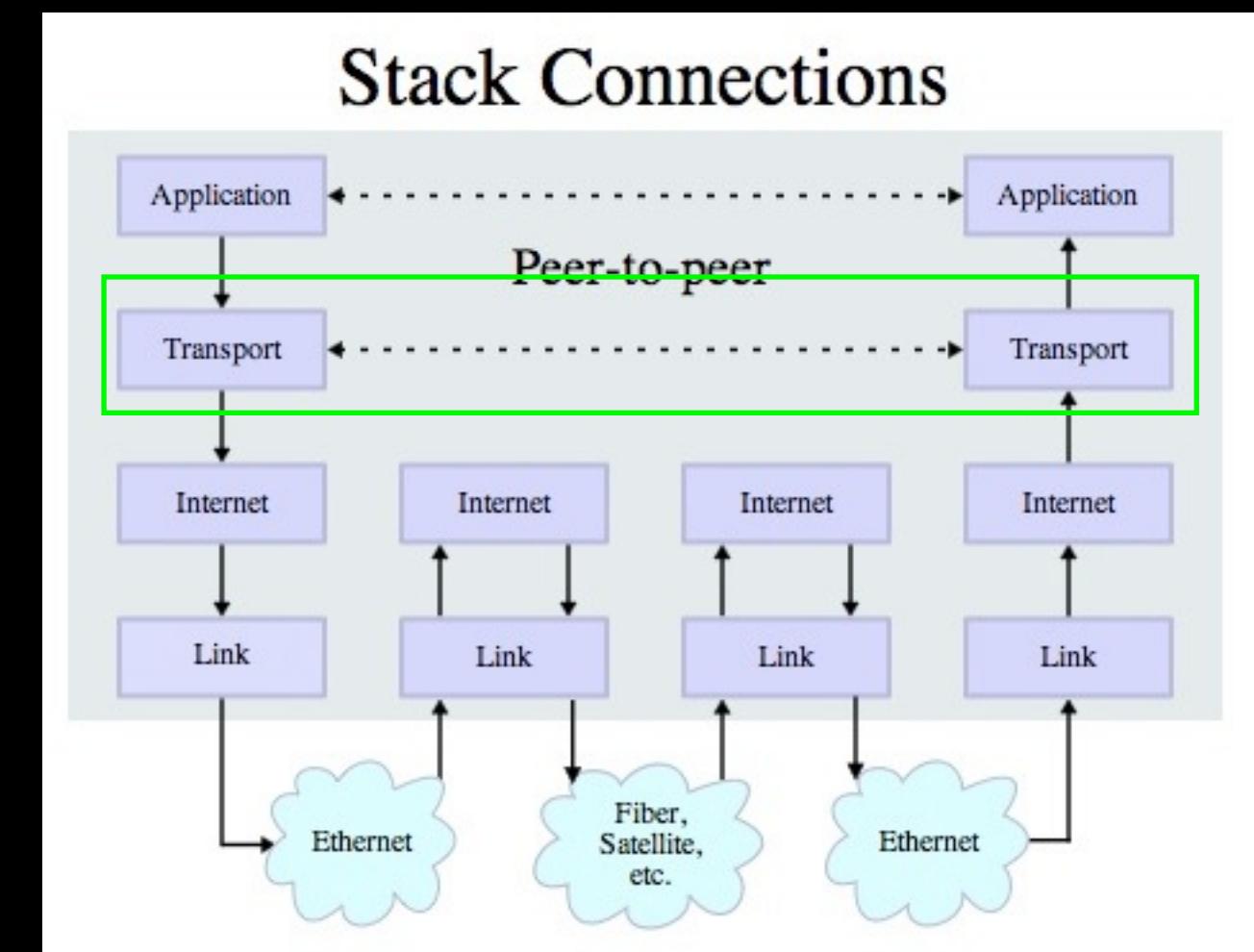
- So many links / hops
- So many routes
- Things can change dynamically and IP has to react (links up/down)
- IP can drop packets



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

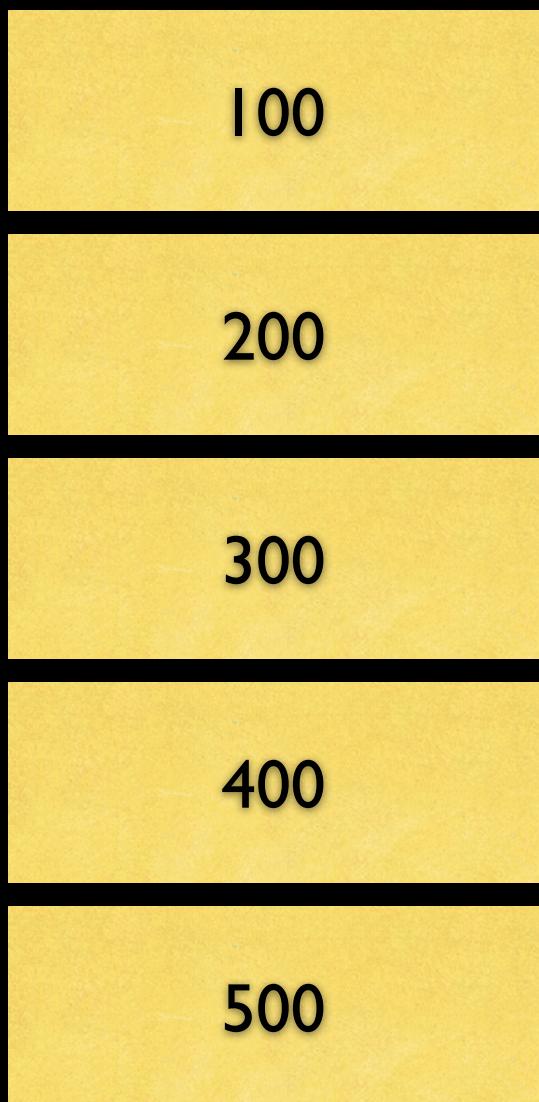
# Transmission Protocol (TCP)

- Built on top of IP
- Assumes IP might lose some data
- In case data gets lost - **we keep a copy of the data** until we get an acknowledgement
- If it takes “too long” - just send it again



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

Sender



Receiver

Break Messages  
into Pieces

Sender

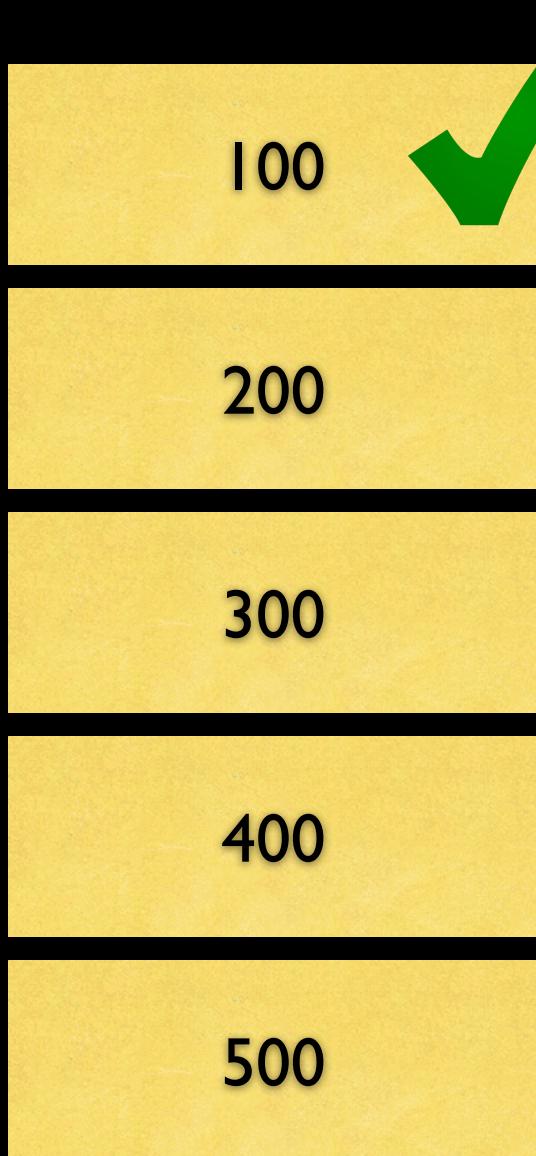


Receiver



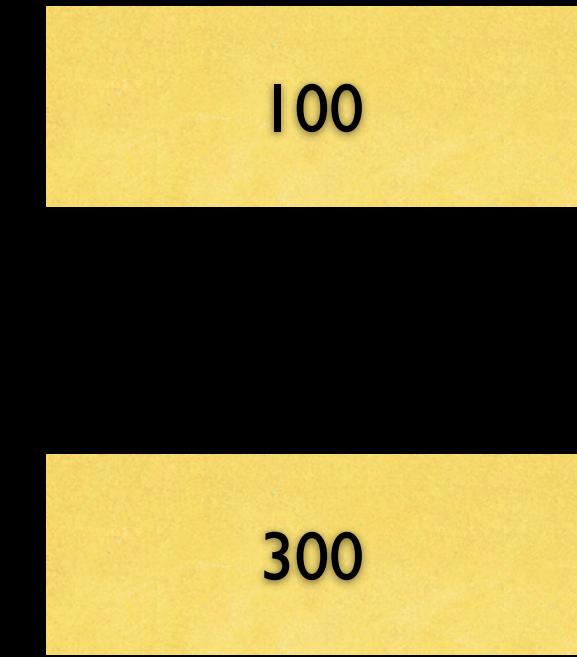
Break Messages  
into Pieces

Sender



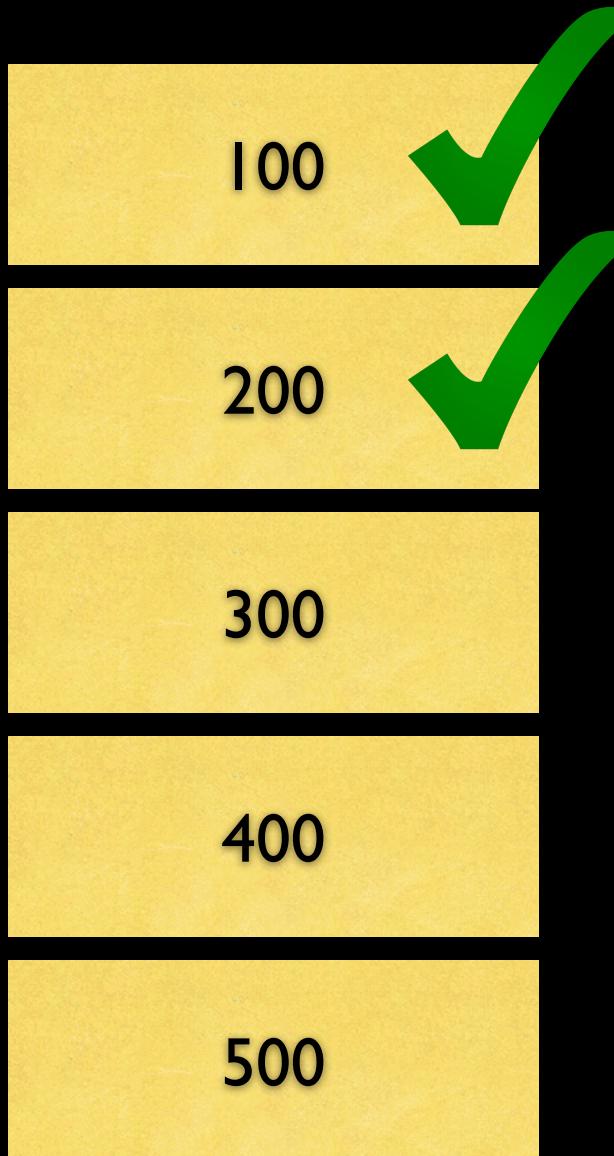
Got 100  
Where is 200

Receiver



Break Messages  
into Pieces

Sender



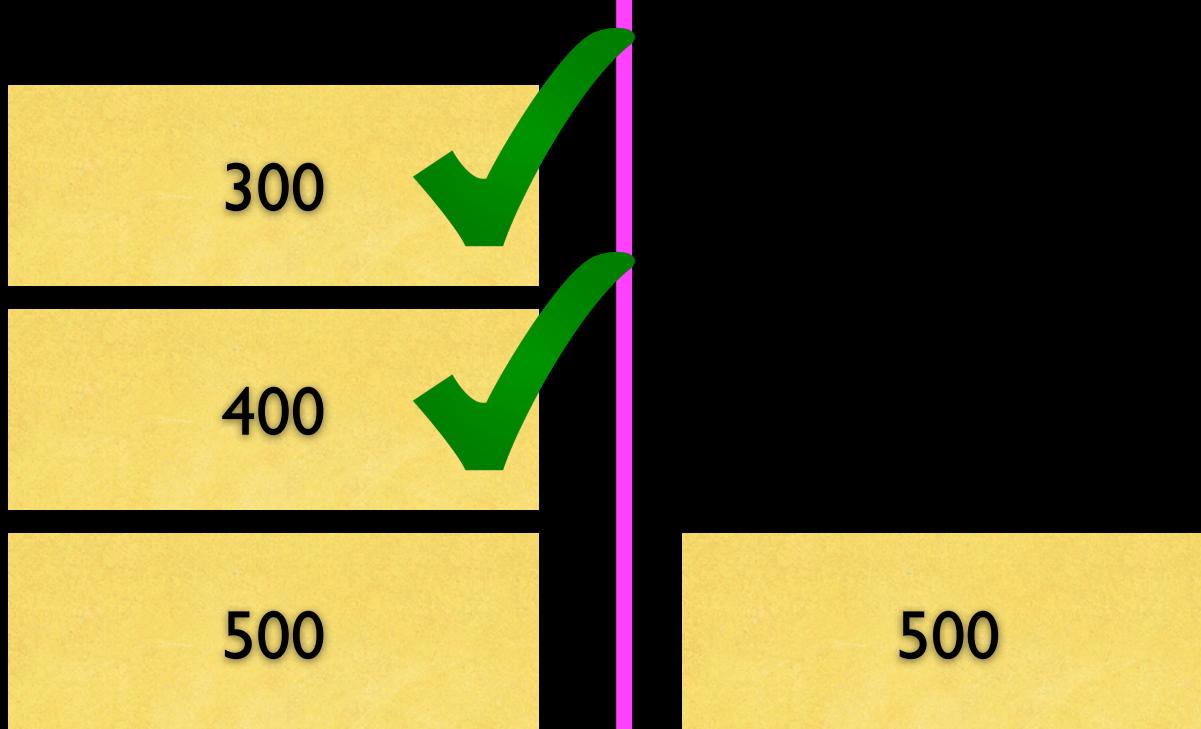
Got 200

Receiver



Break Messages  
into Pieces

Sender



Got 400

Receiver

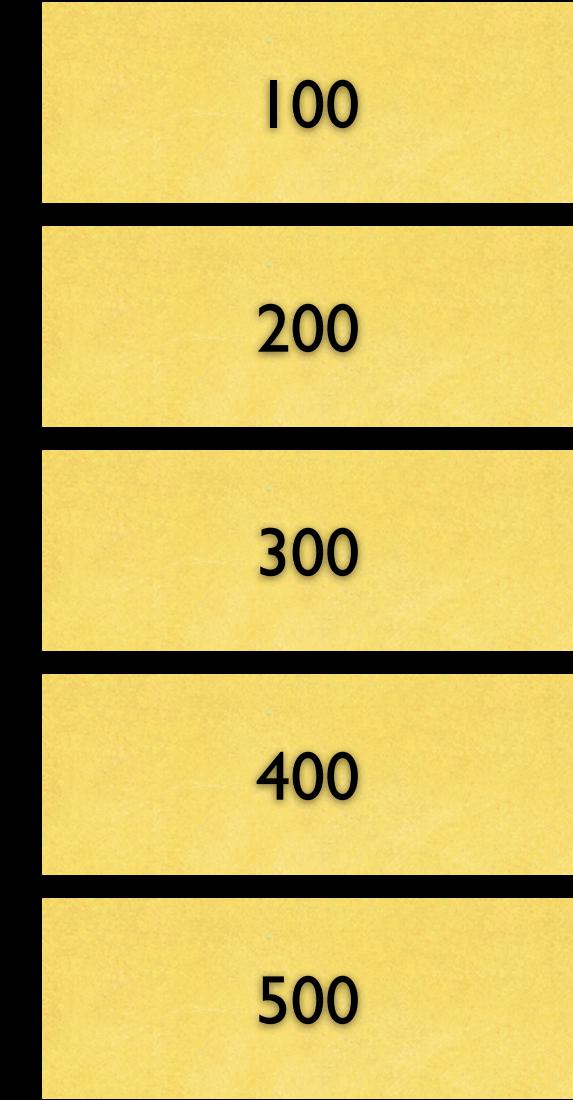
Break Messages  
into Pieces

Sender



Got 500

Receiver



Break Messages  
into Pieces

Sender

Receiver

100

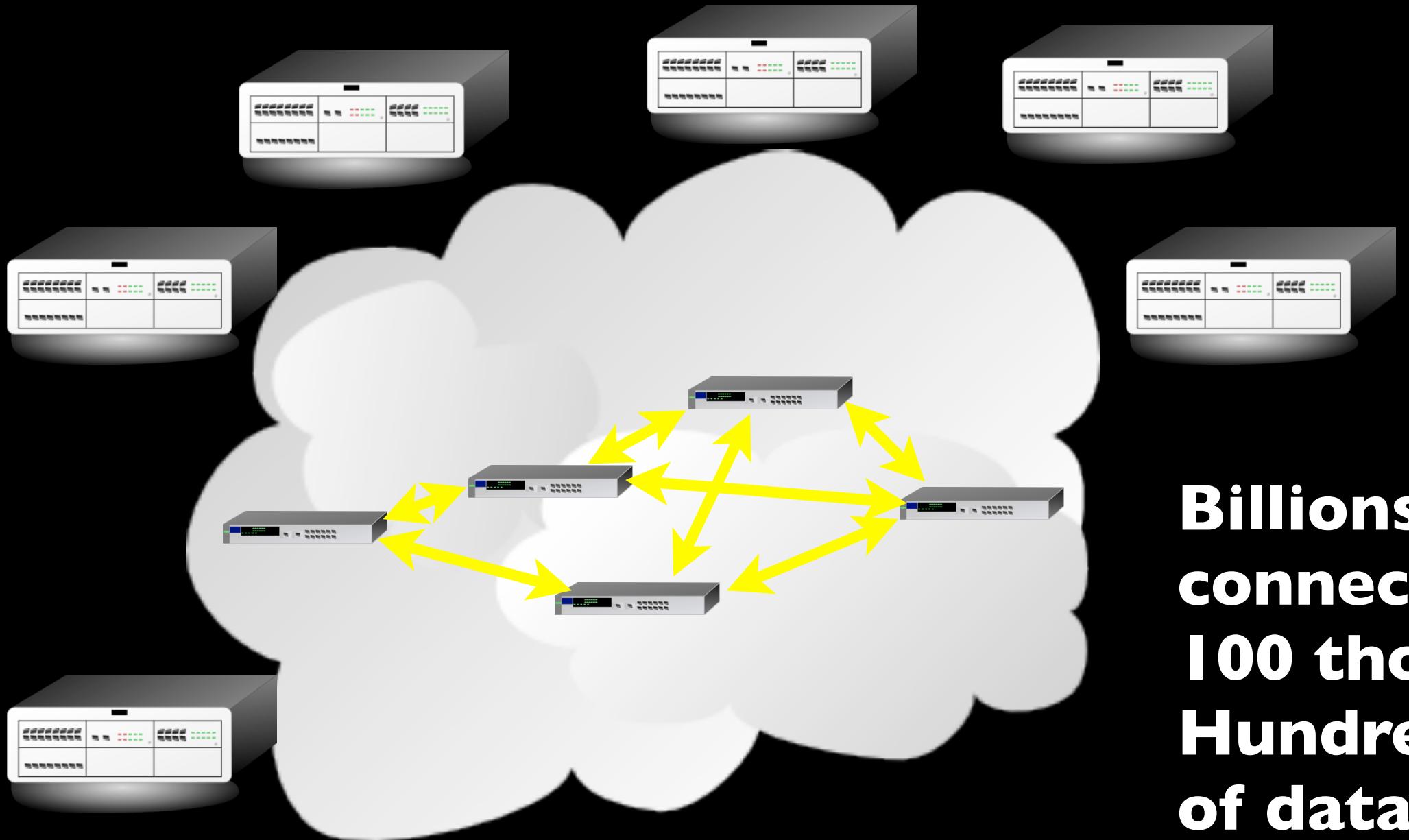
200

300

400

500

Break Messages  
into Pieces

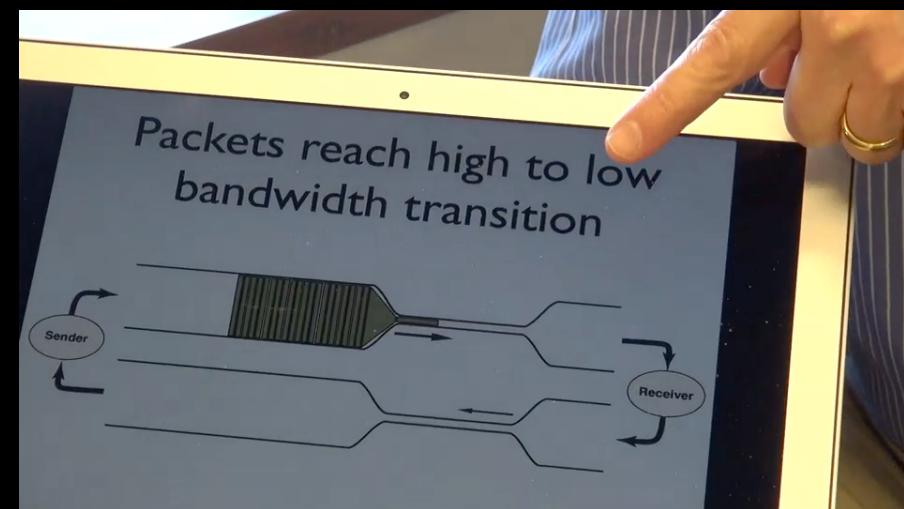
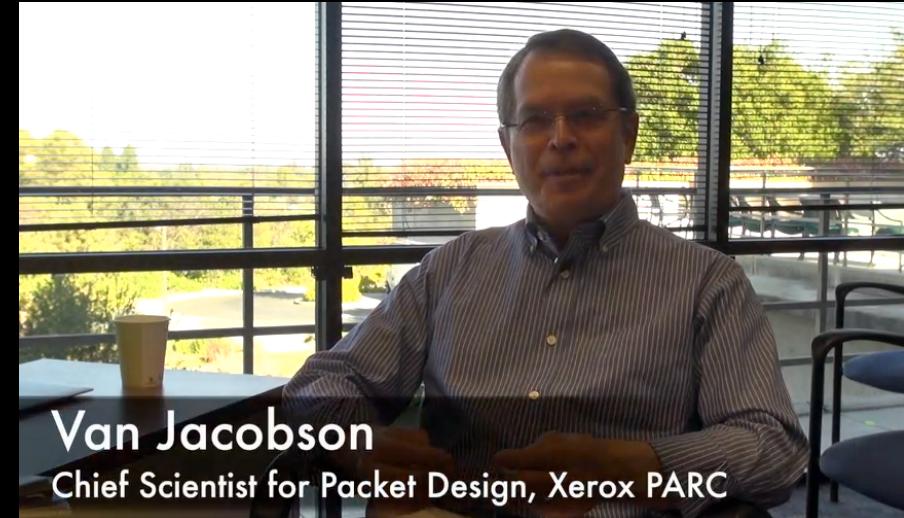


**Billions of computers  
connected to the internet;  
100 thousands of routers.  
Hundreds of billions bytes  
of data enroute at any  
moment.**

**Storage of enroute data  
done at the edges only!**

# One (of many) Scary Problem(s)

- In 1987 as local campuses with 10 MBit networks were connected together using 56Kbit leased lines, things kind of fell apart
- At some point, when there was a little too much traffic, it all fell apart...



<http://www.youtube.com/watch?v=IVgIMeRYmWI>

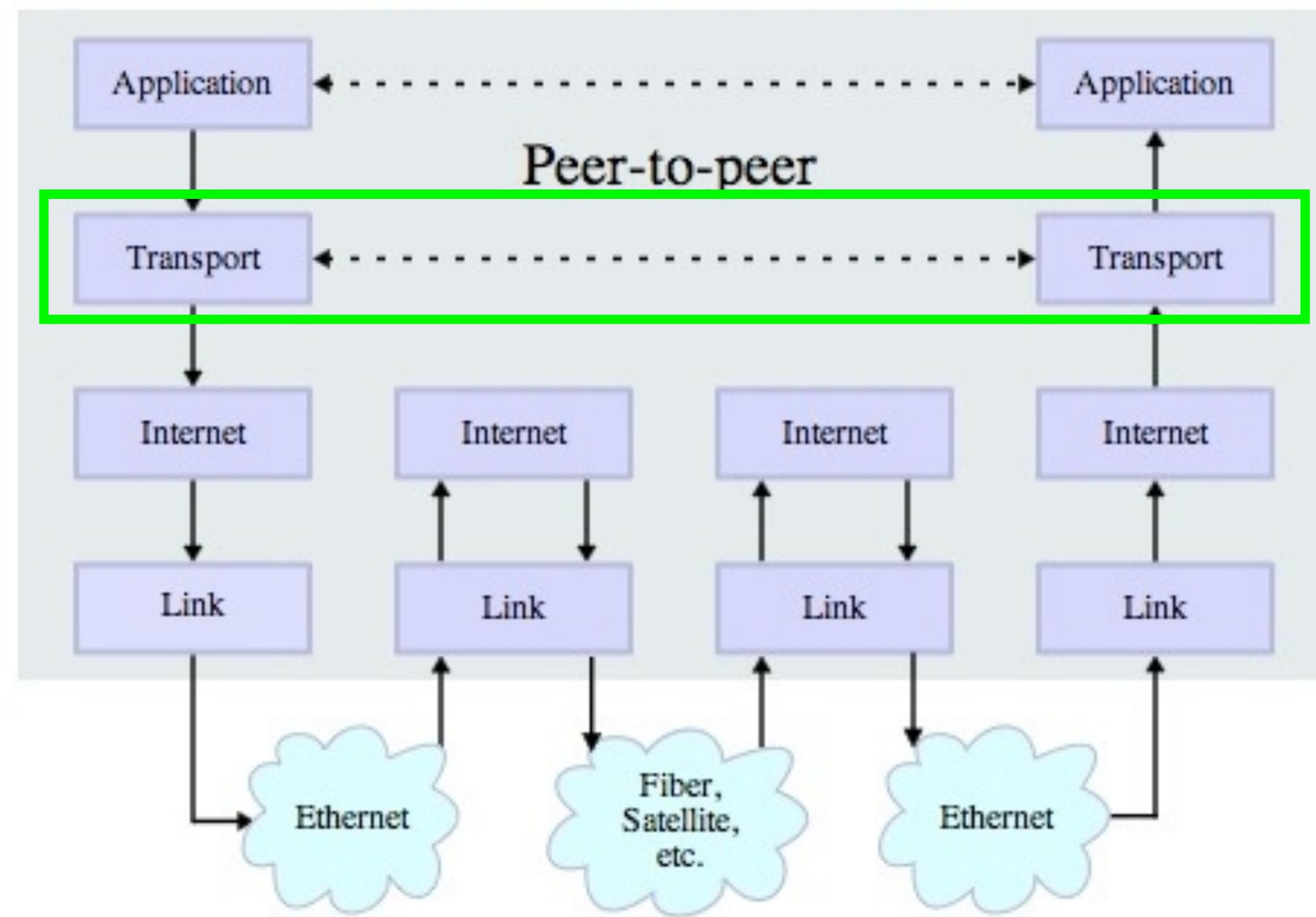
[http://en.wikipedia.org/wiki/Van\\_Jacobson](http://en.wikipedia.org/wiki/Van_Jacobson)

[http://en.wikipedia.org/wiki/TCP\\_congestion\\_avoidance\\_algorithm](http://en.wikipedia.org/wiki/TCP_congestion_avoidance_algorithm)

# Transmission Protocol (TCP)

- The responsibility of the transport layer is to present a reliable end-to-end pipe to the application
- Data either arrives in the proper order or the connection is closed
- TCP keeps buffers in the sending and destination system to keep data which has arrived out of order or to retransmit if necessary
- TCP provides individual connections between applications

# Stack Connections



# Application Layer

Application Layer  
Web, E-Mail, File Transfer

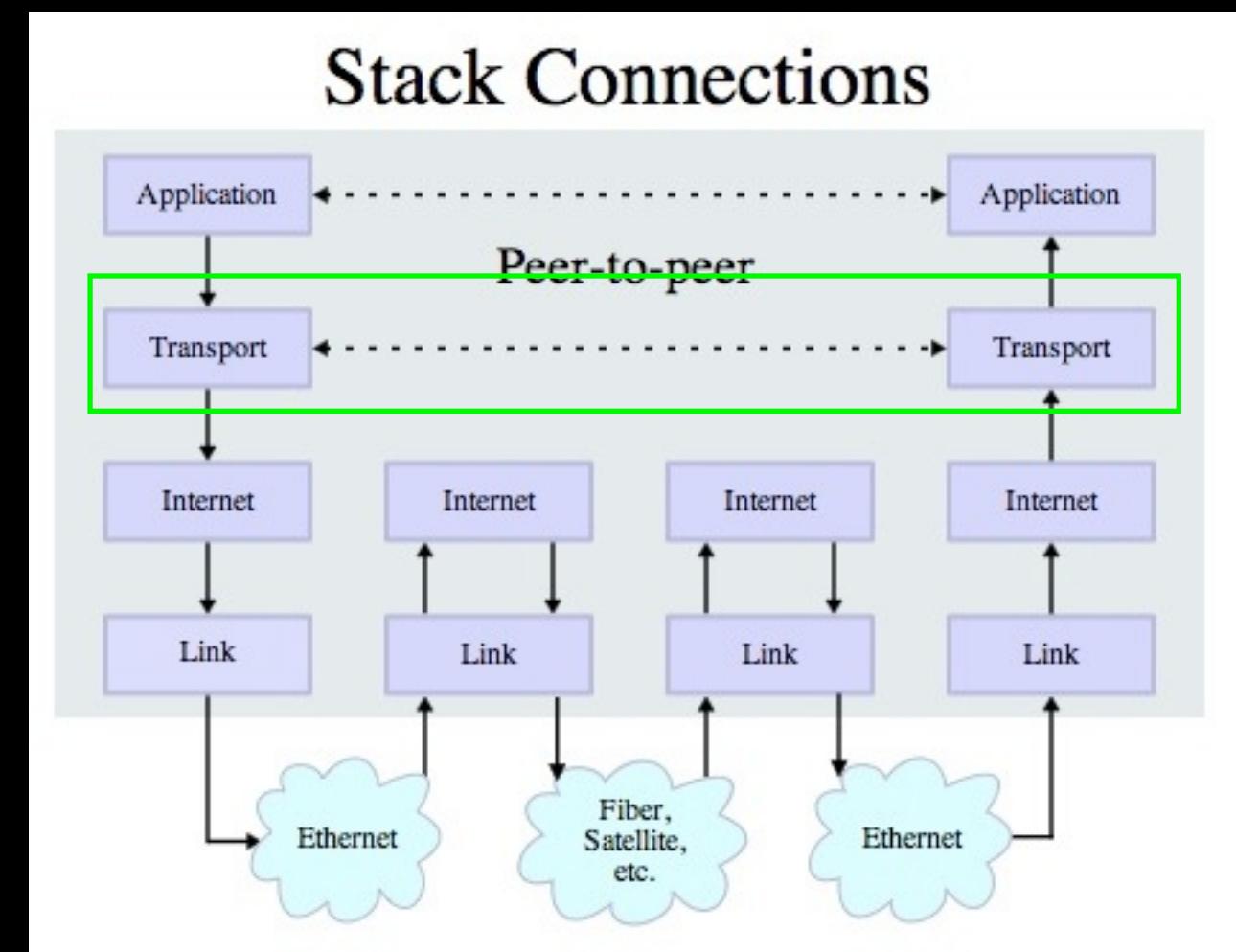
Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Quick Review

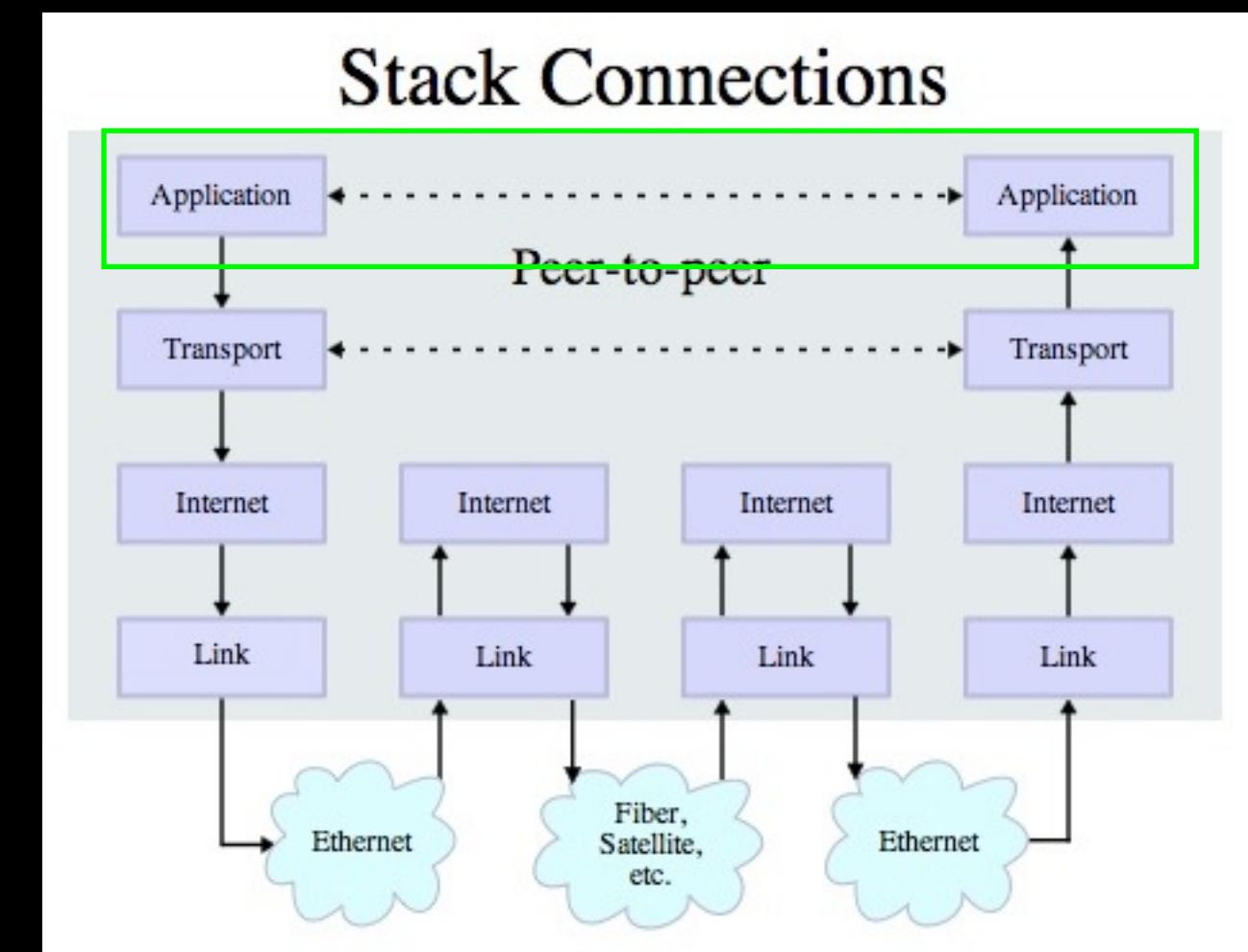
- Link layer: gets the data onto the link, and manages collisions on a single hop
- Internet layer: moves the data over one hop, trying to get it “closer” to its destination
- Transport layer: Assumes that the internet layer may lose data, so request retransmission when needed—provides a nice reliable pipe from source to destination



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Application Protocol

- Since TCP gives us a reliable pipe, what do we want to do with the pipe? What problem do we want to solve?
- Mail
- World Wide Web
- Stream kitty videos



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Two Questions for the Application Layer

- Which application gets the data?
  - Ports
- What are the rules for talking with that application?
  - Protocols

[http://en.wikipedia.org/wiki/TCP\\_and\\_UDP\\_port](http://en.wikipedia.org/wiki/TCP_and_UDP_port)

[http://en.wikipedia.org/wiki/List\\_of\\_TCP\\_and\\_UDP\\_port\\_numbers](http://en.wikipedia.org/wiki/List_of_TCP_and_UDP_port_numbers)

# Ports

- Like extensions in a phone number
- The IP address network number (the area code) gets to the LAN
- The IP address host number (the telephone number) gets you to the destination machine
- The port (the extension) gets you to a specific application

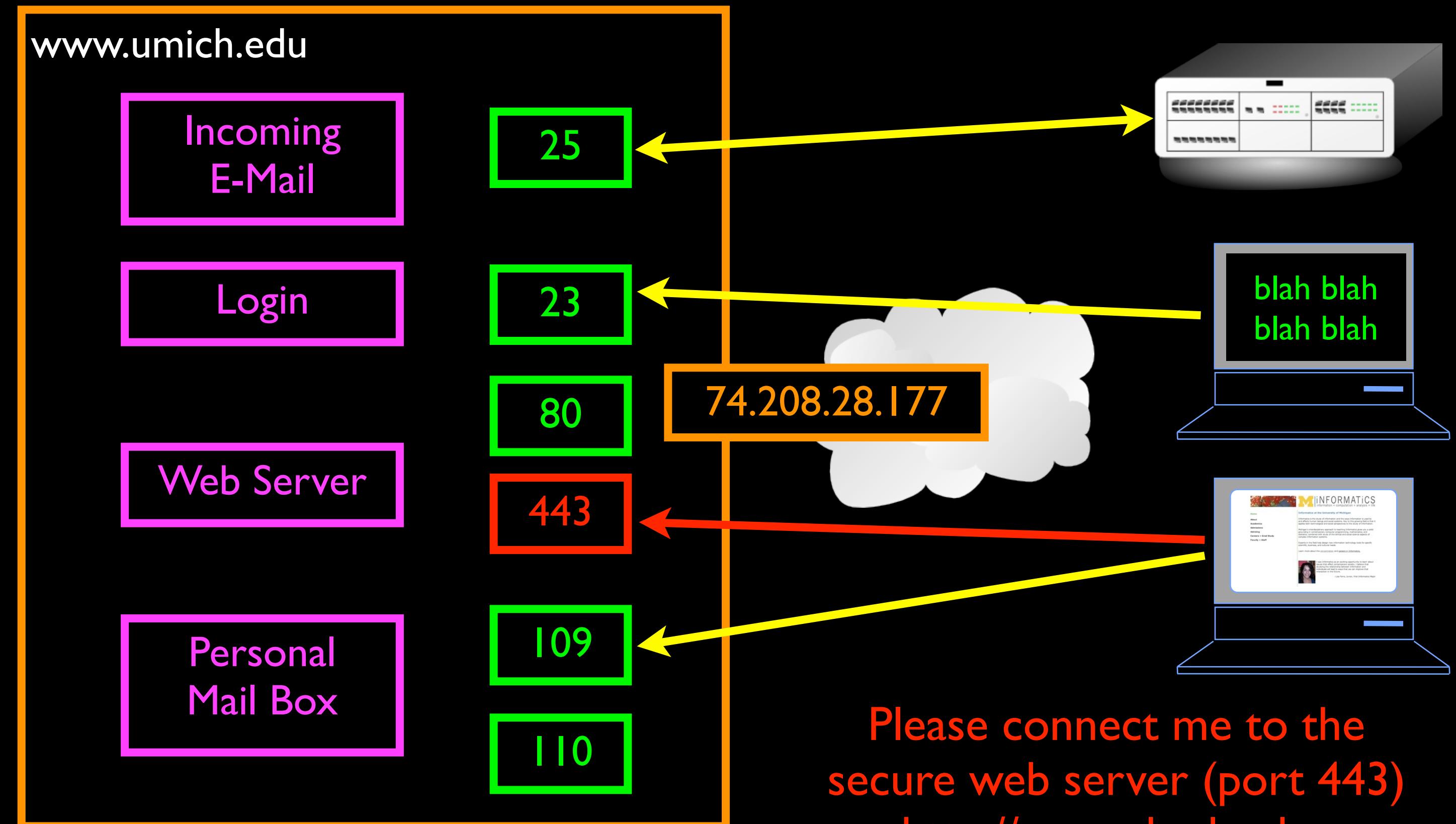
(734) 764 1855, ext. 27

141.211.144.188  
Port 25

# TCP, Ports, and Connections

[http://en.wikipedia.org/wiki/TCP\\_and\\_UDP\\_port](http://en.wikipedia.org/wiki/TCP_and_UDP_port)

[http://en.wikipedia.org/wiki/List\\_of\\_TCP\\_and\\_UDP\\_port\\_numbers](http://en.wikipedia.org/wiki/List_of_TCP_and_UDP_port_numbers)



# Common TCP Ports

- Telnet (23) - Login
- SSH (22) - Secure Login
- HTTP (80)
- HTTPS (443) - Secure
- SMTP (25) (Mail)
- IMAP (143/220/993) - Mail Retrieval
- POP (109/110) - Mail Retrieval
- DNS (53) - Domain Name
- FTP (21) - File Transfer

[http://en.wikipedia.org/wiki/List\\_of\\_TCP\\_and\\_UDP\\_port\\_numbers](http://en.wikipedia.org/wiki/List_of_TCP_and_UDP_port_numbers)

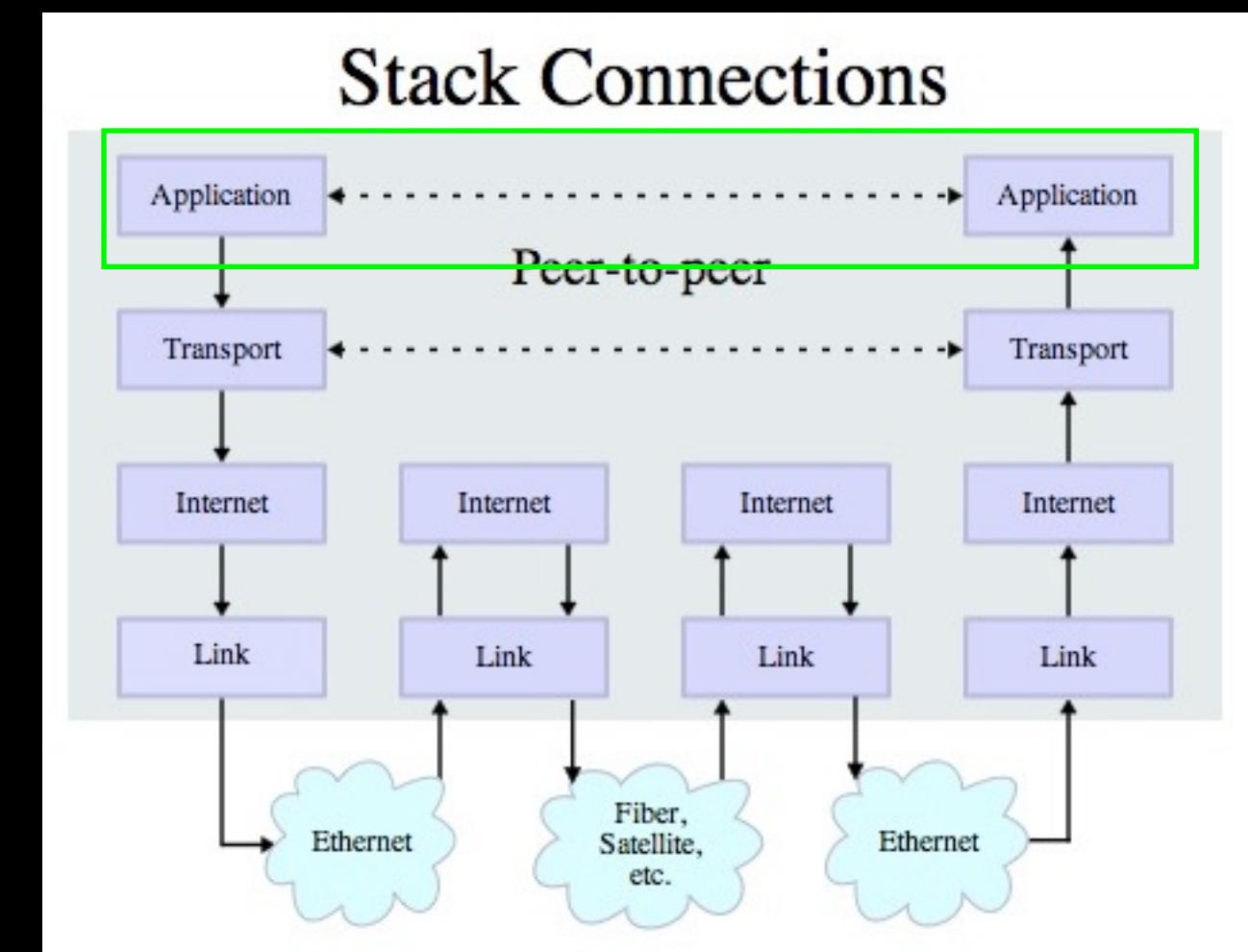
# Application Protocols

<http://en.wikipedia.org/wiki/Http>

<http://en.wikipedia.org/wiki/Pop3>

# Application Protocol

- Since TCP gives us a reliable pipe, what do we want to do with the pipe? What problem do we want to solve?
- Mail
- World Wide Web
- Stream kitty videos



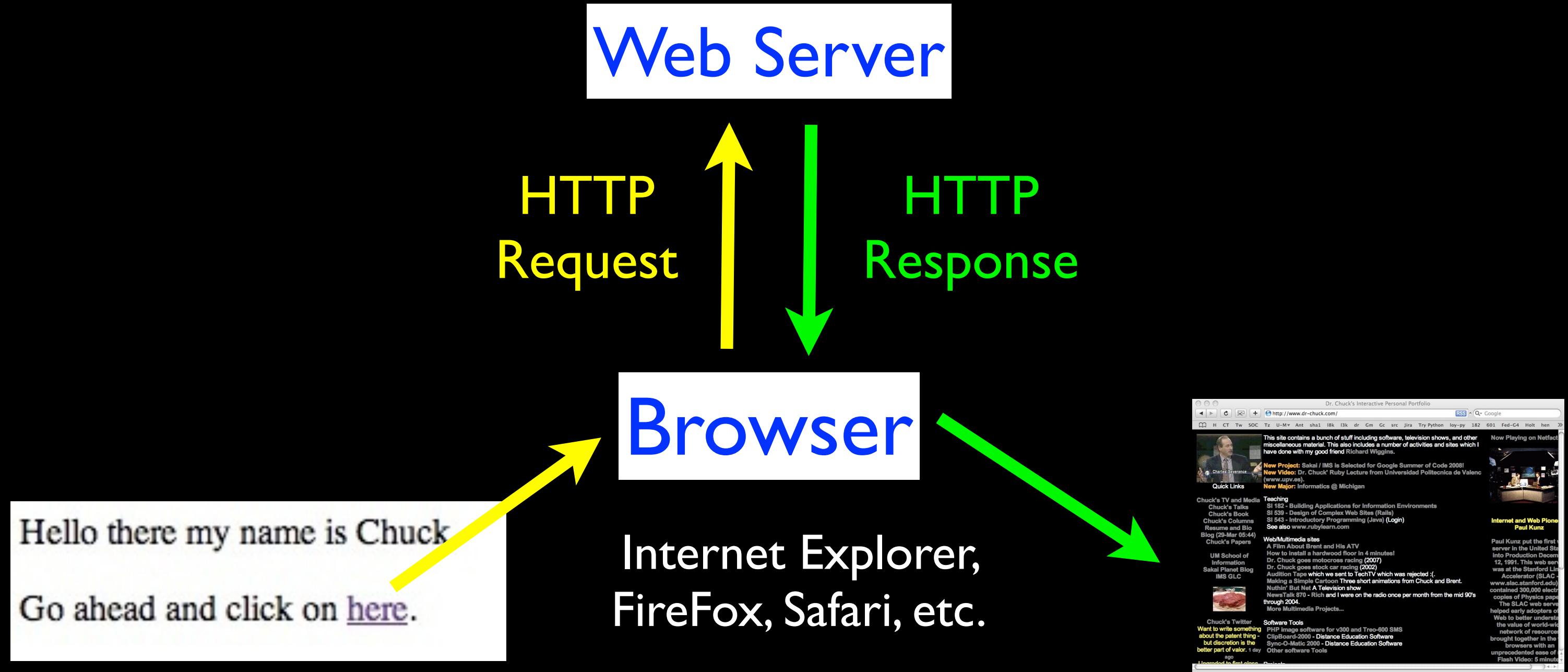
Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# HTTP - Hypertext Transport Protocol

- The dominant Application Layer Protocol on the Internet
- Invented for the Web - to Retrieve HTML, Images, Documents etc
- Extended to be data in addition to documents - RSS, Web Services, etc..
- Basic Concept - Make a Connection - Request a document - Retrieve the Document - Close the Connection

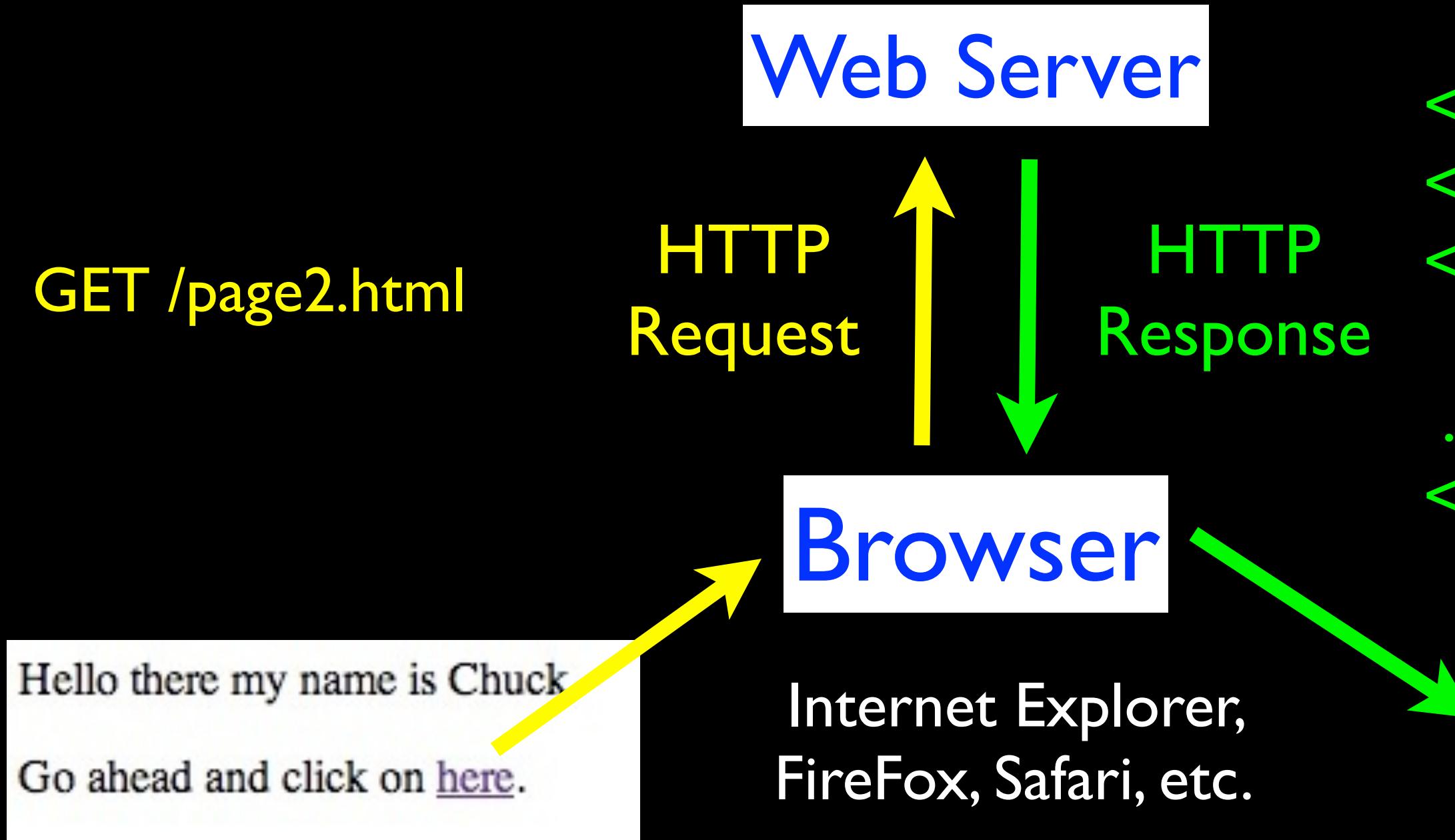
<http://en.wikipedia.org/wiki/Http>

# HTTP Request / Response Cycle



[http://www.oreilly.com/openbook/cgi/ch04\\_02.html](http://www.oreilly.com/openbook/cgi/ch04_02.html)

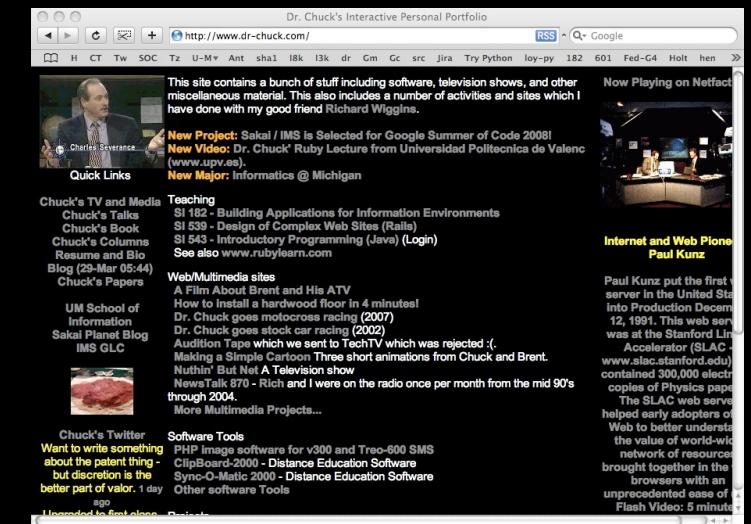
# HTTP Request / Response Cycle



[http://www.oreilly.com/openbook/cgi/ch04\\_02.html](http://www.oreilly.com/openbook/cgi/ch04_02.html)

<head> .. </head>  
<body>  
<h1>Welcome to my  
application</h1>

....  
</body>



Source: <http://www.dr-chuck.com/>

# Internet Standards

- The standards for all of the Internet protocols (inner workings) are developed by an organization
- Internet Engineering Task Force (IETF)
- [www.ietf.org](http://www.ietf.org)
- Standards are called “RFCs” - “Request for Comments”

Network Working Group  
Request for Comments: 1945  
Category: Informational

T. Berners-Lee  
MIT/LCS  
R. Fielding  
UC Irvine  
H. Frystyk  
MIT/LCS  
May 1996

## Hypertext Transfer Protocol -- HTTP/1.0

### Status of This Memo

This memo provides information for the Internet community. This memo does not specify an Internet standard of any kind. Distribution of this memo is unlimited.

### IESG Note:

The IESG has concerns about this protocol, and expects this document to be replaced relatively soon by a standards track document.

### Abstract

The Hypertext Transfer Protocol (HTTP) is an application-level protocol with the lightness and speed necessary for distributed, collaborative, hypermedia information systems. It is a generic, stateless, object-oriented protocol which can be used for many tasks, such as name servers and distributed object management systems, through extension of its request methods (commands). A feature of HTTP is the typing of data representation, allowing systems to be built independently of the data being transferred.

Source: <http://www.ietf.org/rfc/rfc1945.txt>

### 5.1.2 Request-URI

The Request-URI is a Uniform Resource Identifier (Section 3.2) and identifies the resource upon which to apply the request.

```
Request-URI      = absoluteURI | abs_path
```

The two options for Request-URI are dependent on the nature of the request.

The absoluteURI form is only allowed when the request is being made to a proxy. The proxy is requested to forward the request and return the response. If the request is GET or HEAD and a prior response is cached, the proxy may use the cached message if it passes any restrictions in the Expires header field. Note that the proxy may forward the request on to another proxy or directly to the server specified by the absoluteURI. In order to avoid request loops, a proxy must be able to recognize all of its server names, including any aliases, local variations, and the numeric IP address. An example Request-Line would be:

```
GET http://www.w3.org/pub/WWW/TheProject.html HTTP/1.0
```

The most common form of Request-URI is that used to identify a resource on an origin server or gateway. In this case, only the absolute path of the URI is transmitted (see Section 3.2.1, `abs_path`). For example, a client wishing to retrieve the resource above directly from the origin server would create a TCP connection to port 80 of the host "`www.w3.org`" and send the line:

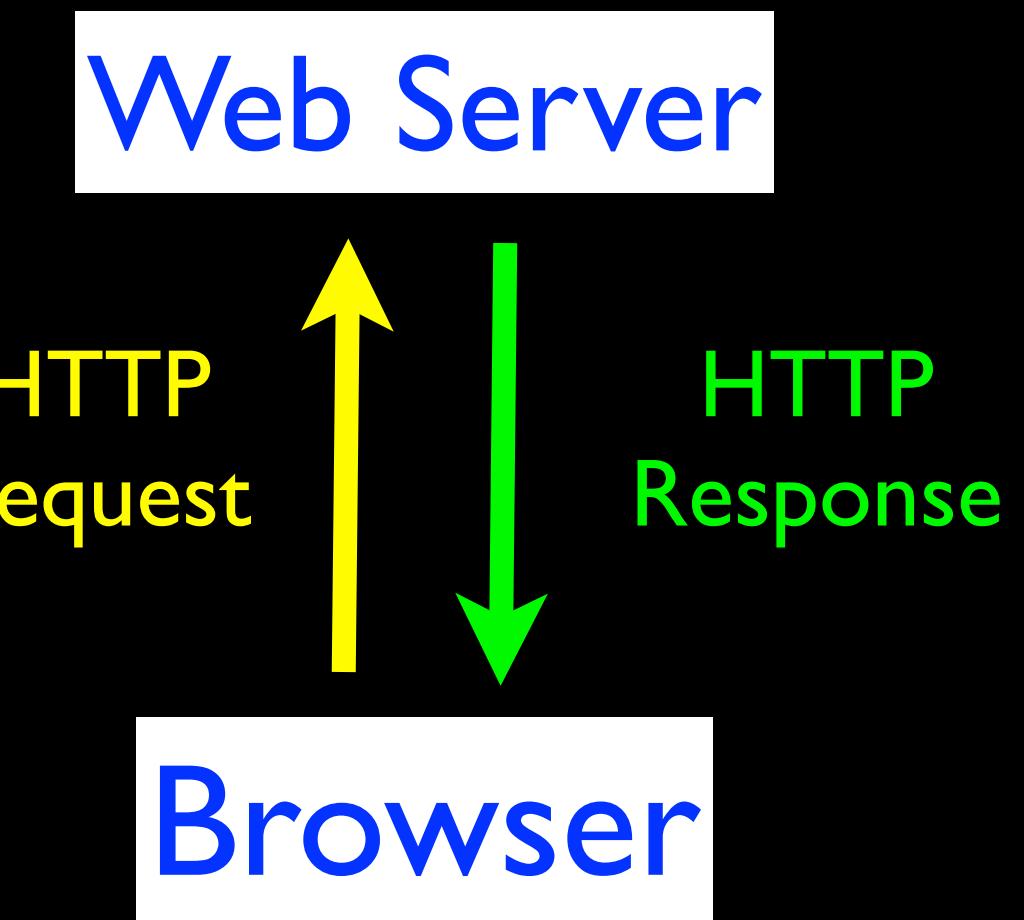
```
GET /pub/WWW/TheProject.html HTTP/1.0
```

followed by the remainder of the Full-Request. Note that the absolute path cannot be empty; if none is present in the original URI, it must be given as "/" (the server root).

The Request-URI is transmitted as an encoded string, where some characters may be escaped using the "% HEX HEX" encoding defined by RFC 1738 [4]. The origin server must decode the Request-URI in order to properly interpret the request.

# “Hacking” HTTP

```
Last login: Wed Oct 10 04:20:19 on ttys2  
si-csev-mbp:~ csev$ telnet www.dr-chuck.com 80  
Trying 74.208.28.177...  
Connected to www.dr-chuck.com.  
Escape character is '^]'.  
GET http://www.dr-chuck.com/page1.htm  
<h1>The First Page</h1>  
<p>  
If you like, you can switch to the  
<a href="http://www.dr-chuck.com/page2.htm">  
Second Page</a>.  
</p>
```



Port 80 is the non-encrypted HTTP port

# Accurate Hacking in the Movies

- Matrix Reloaded
- Bourne Ultimatum
- Die Hard 4
- ...



```
80/tcp open http  
81/tcp open  
10 [mobile]  
# nmap -v -SS -O 10.2.2.2  
Starting nmap 0.2.54BETA25  
Insufficient responses for TCP sequencing (3), OS detection  
accurate  
Interesting ports on 10.2.2.2:  
(The 1539 ports scanned but not shown below are in state: closed)  
Port      State    Service  
22/tcp    open     ssh  
No exact OS matches for host  
Nmap run completed -- 1 IP address (1 host up) scanned  
# sshnuke 10.2.2.2 -rootpw="Z10H0101"  
Connecting to 10.2.2.2:ssh ... successful.  
ReAttempting to exploit SSHv1 CRC32 ... successful.  
IP Resetting root password to "Z10H0101"... successful.  
System open: Access Level <9>  
# ssh 10.2.2.2 -l root  
root@10.2.2.2's password:   
DC RTF CONTROL  
ACCESS GRANTED
```

<http://nmap.org/movies.html> (scroll down for video)  
Or search YouTube for "Trinity hacking scene"



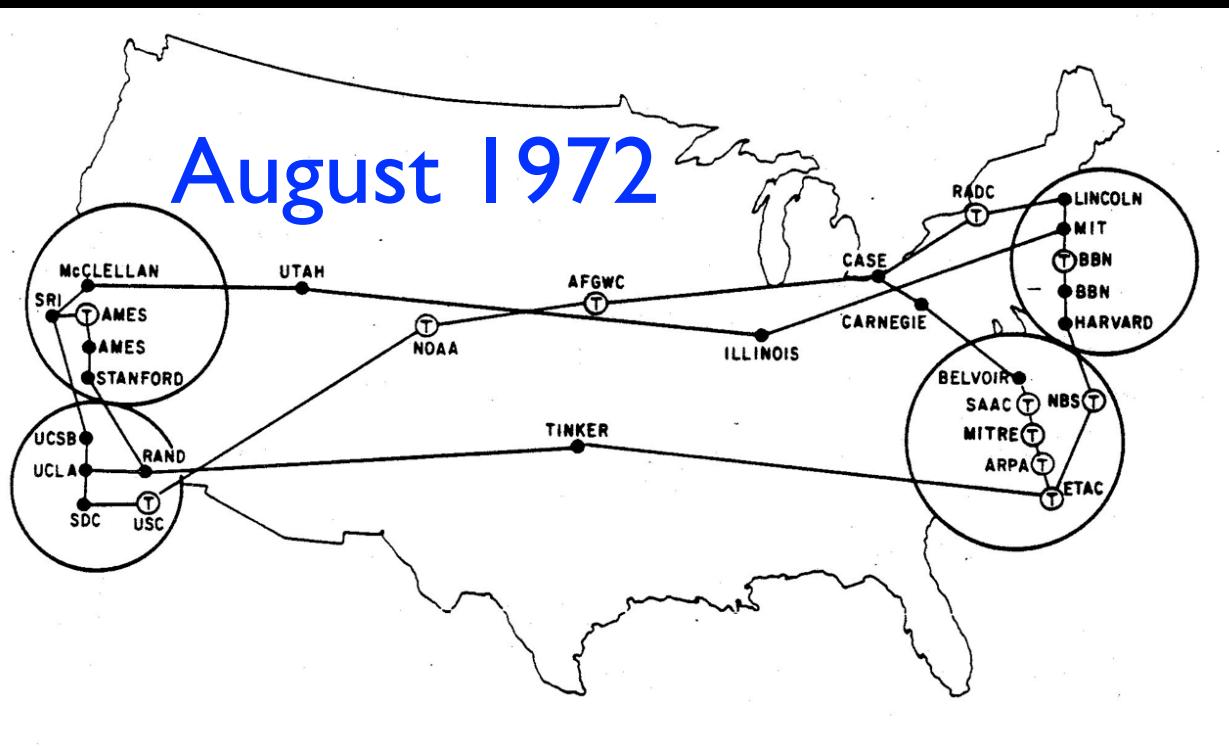
```
80/tcp      open     http  
81/tcp      open     hostc2.nc  
10/ssh     [ mobile]  
11  
12 nmap -v -sS -O 10.2.2.2  
13 Starting nmap 0.2.54BETA25  
13 Insufficient responses for TCP sequencing (3). OS detection  
13 accurate  
14 Interesting ports on 10.2.2.2:  
14 (The 1539 ports scanned but not shown below are in state: closed)  
51 Port      State    Service  
51 22/tcp    open     ssh  
58  
68 No exact OS matches for host  
68  
24 Nmap run completed -- 1 IP address (1 host up) scanned  
50 # sshnuke 10.2.2.2 -rootpw="Z10N0101"  
Connecting to 10.2.2.2:ssh ... successful.  
Re-Attempting to exploit SSHv1 CRC32 ... successful.  
IP Resetting root password to "Z10N0101".  
System open: Access Level <9>  
Hn # ssh 10.2.2.2 -l root  
root@10.2.2.2's password: ■  
RTF CONTROL  
ACCESS GRANTED
```

<http://nmap.org/movies.html> (scroll down for video)  
Or search YouTube for "Trinity hacking scene"

# Application Layer Summary

- We start with a “pipe” abstraction - we can send and receive data on the same “socket”
- We can optionally add a security layer to TCP using SSL - Secure Socket Layer (aka TLS - Transport Layer Security)
- We use well known “port numbers” so that applications can find a particular application \*within\* a server such as a mail server, web service, etc

# The Architecture of the Internet



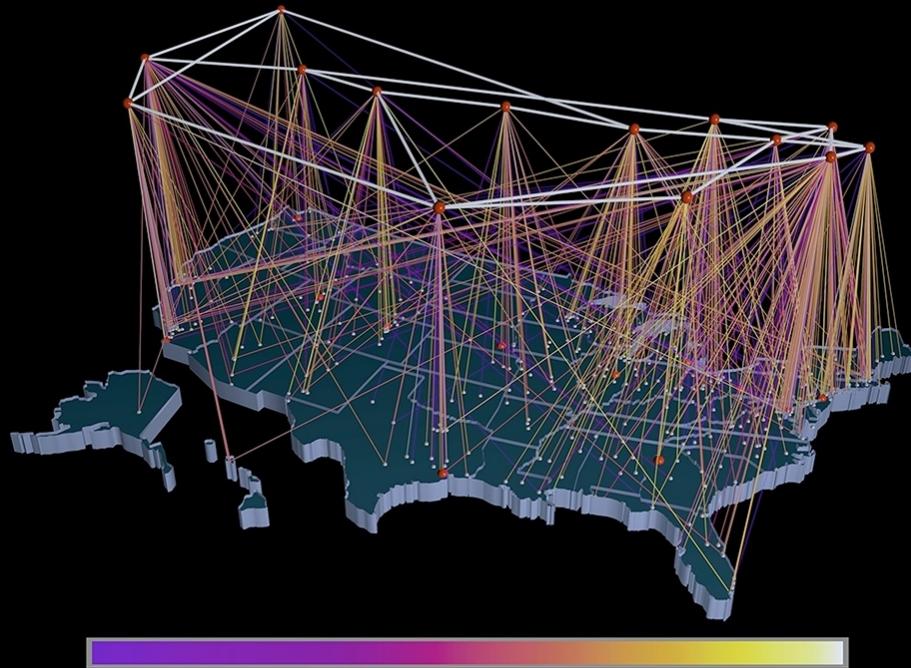
Application Layer  
Web, E-Mail, File Transfer

Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# The Architecture of the Internet

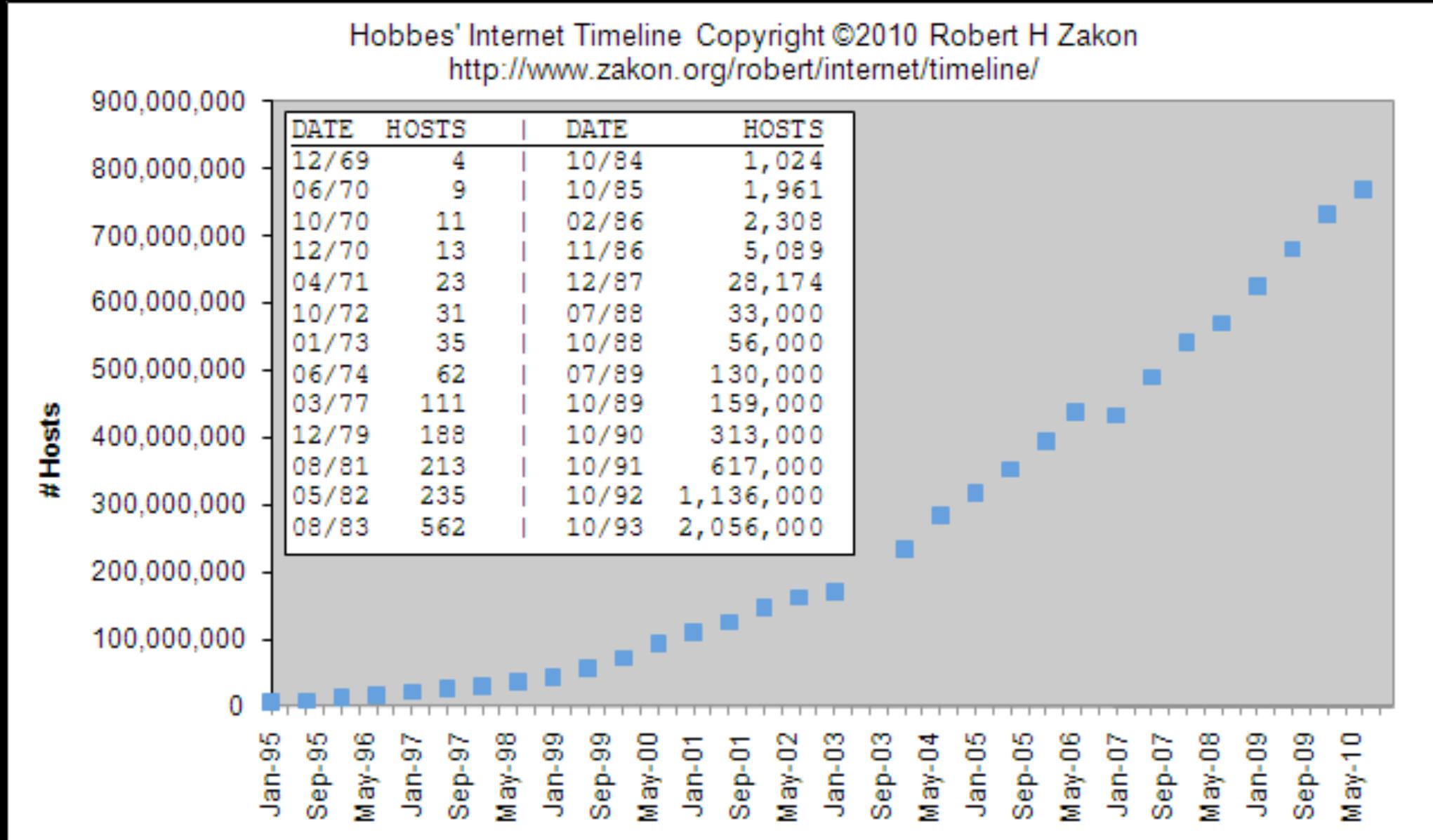


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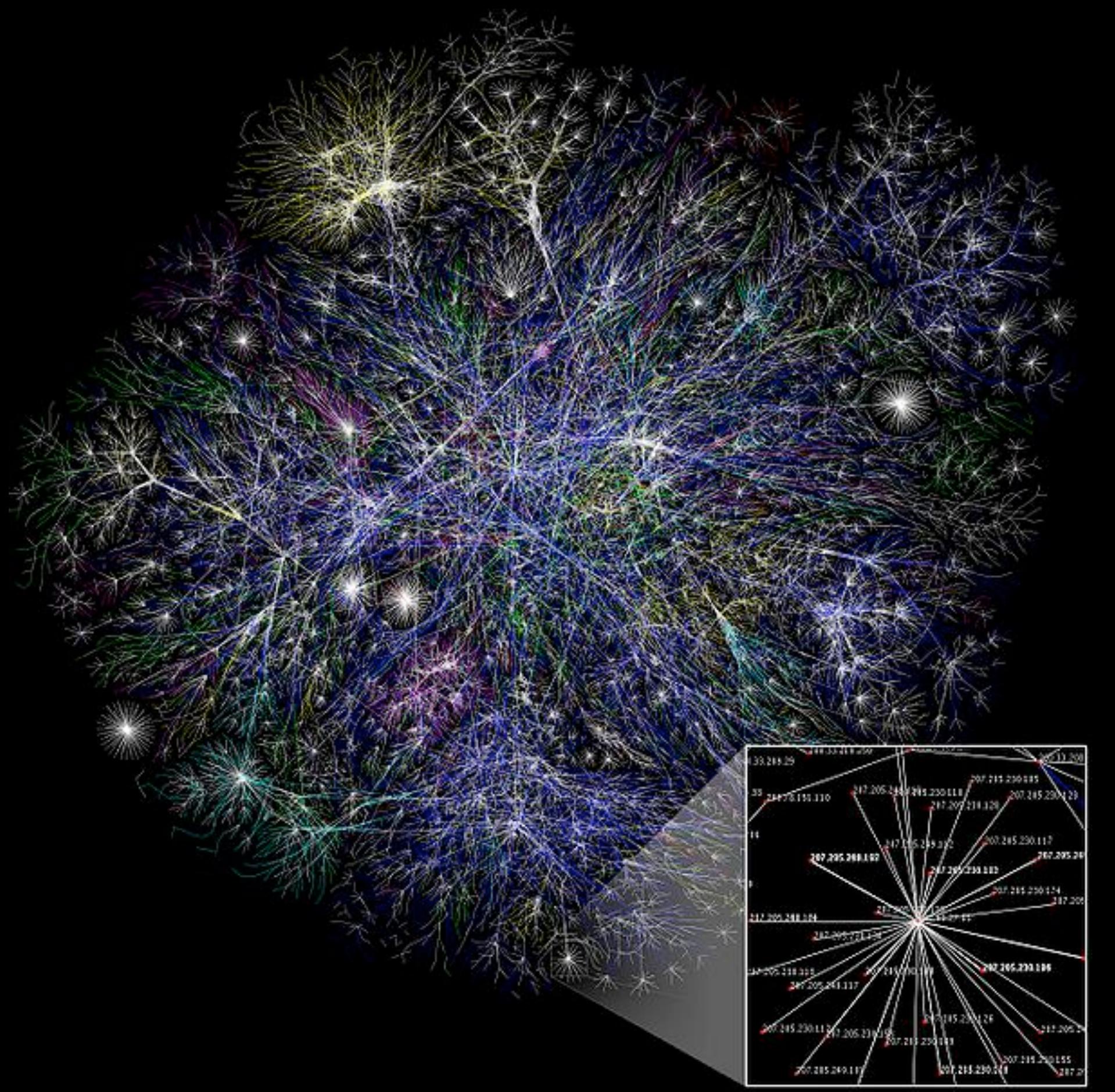


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**Application Layer**  
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**Reliable Connections**

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**Simple, Unreliable**

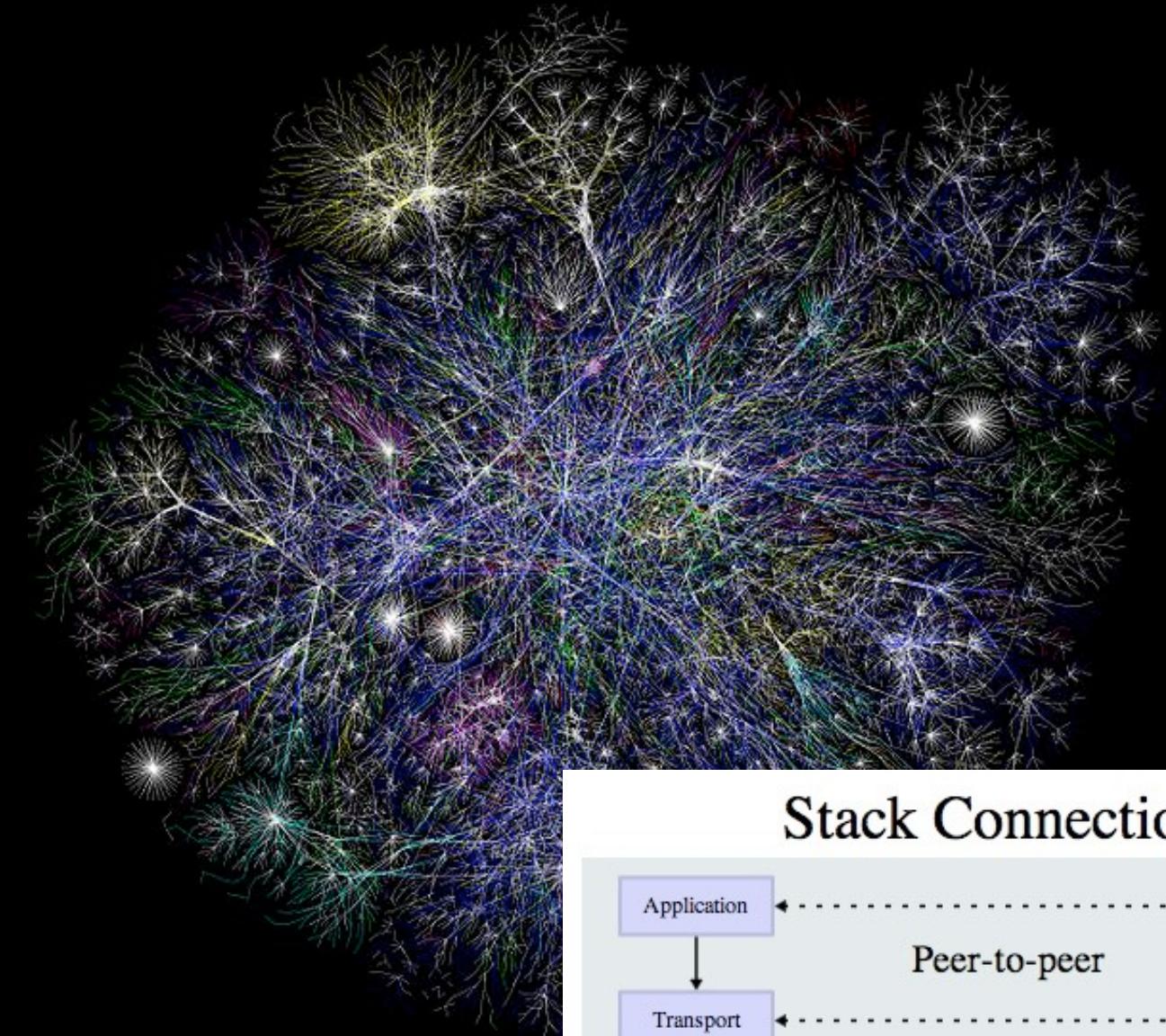
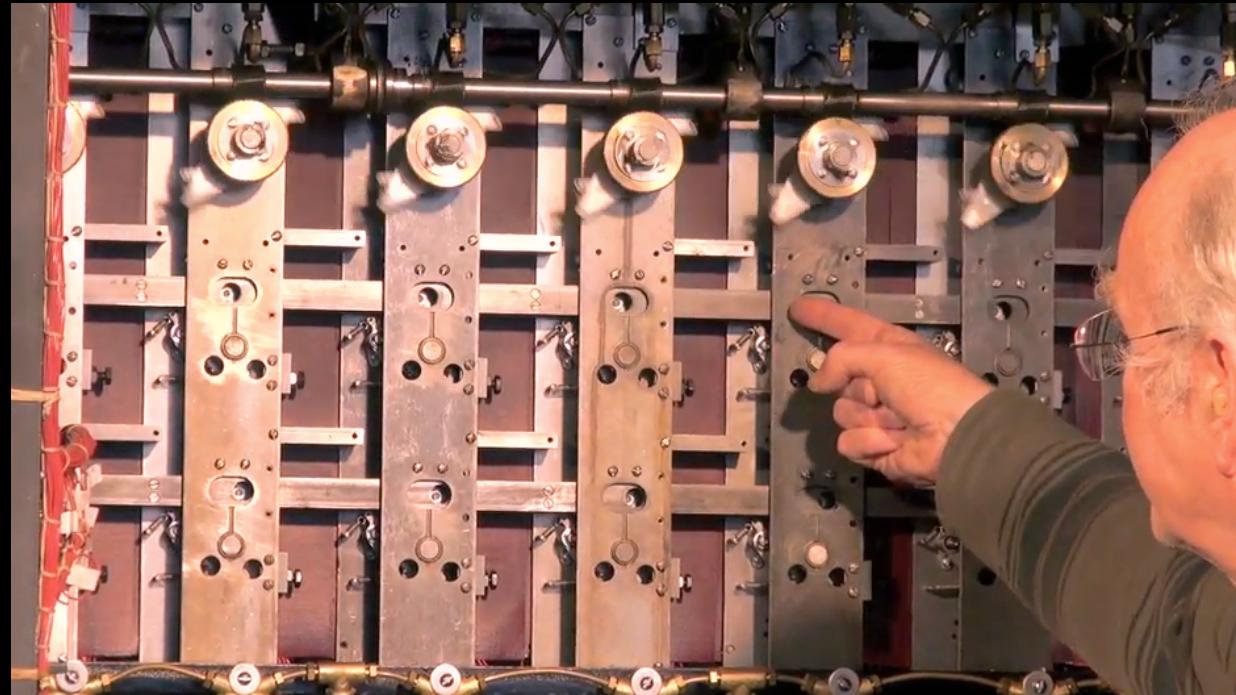
**Link Layer (Ethernet, WiFi)**  
**Physical Connections**

# The Internet: An Amazing Design

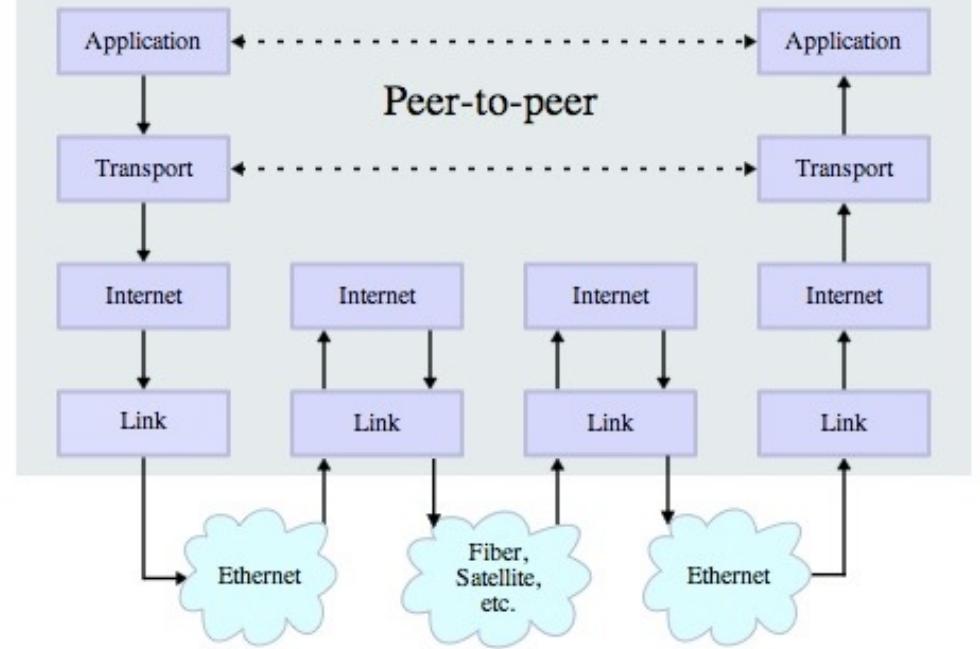
- Hundreds of millions of computers
- Thousands of routers inside the Internet
- Hundreds of millions of simultaneous connections
- Trillions of bytes of data moved per second around the world
- And it works

# The Internet

- It is said that “The Internet is the largest single engineering effort ever created by mankind”
- It was created to work in an organic way - to repair itself and automatically adjust when parts fail
- No one part of the Internet knows all of the Internet (like life)
- It is never 100% up - but it seems up all the time



## Stack Connections



# We are not done experimenting...

- There is still very active exploration on how network technology can be improved
- Content-Centric Networking is only one advanced idea
- Routers in the future can have \*lots\* of memory - lets try not to send the same piece of data more than once





# Additional Source Information

- xkcd, <http://xkcd.com/742/>, CC: BY-NC, <http://creativecommons.org/licenses/by-nc/2.5/>
- Internet Protocol Suite Diagrams: Kbrose, Wikimedia Commons, [http://upload.wikimedia.org/wikipedia/commons/c/c4/IP\\_stack\\_connections.svg](http://upload.wikimedia.org/wikipedia/commons/c/c4/IP_stack_connections.svg), CC:BY-SA, <http://creativecommons.org/licenses/by-sa/3.0/deed.en>
- All your bases are belong to me: Karin Dalziel, Flickr, <http://www.flickr.com/photos/nirak/270213335/>, CC:BY, <http://creativecommons.org/licenses/by/2.0/deed.en>
- Internet Map: The Opte Project, Wikimedia Commons, [http://upload.wikimedia.org/wikipedia/commons/d/d2/Internet\\_map\\_1024.jpg](http://upload.wikimedia.org/wikipedia/commons/d/d2/Internet_map_1024.jpg), CC:BY, <http://creativecommons.org/licenses/by/2.5/deed.en>

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- Please contact me if you are interested in reusing or remixing these materials in your own teaching or educational context