

# Corten: Refinement Types for Imperative Languages with Ownership

**Abschlusspräsentation Masterarbeit**

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Motivation  
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Empirical Analysis

Solution

Soundness Justification

Related Work

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# Motivation

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Refinement Types **rondon\_liquid\_2008** in Functional Programming Languages

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let  $\Gamma = (a : \{v : i32 \mid \text{true}\}, b : \{v : i32 \mid \text{true}\})$  and  $\tau = \{v : i32 \mid v \geq a \wedge v \geq b\}$

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fn client(...) {  
    ...  
    clamp(&mut x, 5);  
    clamp(&mut y, 6);  
    print(x);  
    ...  
}
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What does this `print(x)` output?

- In most imperative programming languages:
  - Could be: old x or 5
  - But also 6 (if x aliases with y)!
- In Rust:
  - Just old x or 5
  - And nothing else!

# Motivation

```
fn clamp(a: &mut i32, b: i32) {
    // borrows a
    // owns b
    if *a > b { *a = b }
    // "returns" the borrow of a
}

fn client(...) { // owns x, y
    ...
    clamp(&mut x, 5); // lend x mutably
    clamp(&mut y, 6); // lend y mutably
    print(x);
    ...
}
```

## Ownership in Rust: Mutability XOR Aliasing

Each lexical scope tracks permissions for visible memory objects. Possible Permission Levels:

- Owner (e.g. b)
  - can: read, write
  - transfer ownership (if no outstanding borrows)
- Mutable Reference (e.g. &mut x)
  - can: read, write
  - guarantee: no aliasing
- Immutable Reference (e.g. &v)
  - can: read, alias
  - guarantee: no mutation

# Motivation

Consequences:

- unique data owner
- no global, mutable state
- no cycles in memory structure

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## Used for:

- safe non-gc memory management
- safe concurrency
- safe low-level hardware access
- ...

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- $\Rightarrow$  show: program verification as well

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# Extend Refinement Types to Rust

asd

Motivation  
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Empirical Analysis

Solution

Soundness Justification

Related Work

Conclusion / Future Work

## Backup-Teil

Folien, die nach `\beginbackup` eingefügt werden, zählen nicht in die Gesamtzahl der Folien.

# Blöcke

## in den KIT-Farben

Greenblock

Standard (block)

Blueblock

= exampleblock

Redblock

= alertblock

Brownblock

Purpleblock

Cyanblock

Yellowblock

Lightgreenblock

Orangeblock

Grayblock

**Contentblock**  
(farblos)



# Auflistungen

Text

- Auflistung  
Umbruch
- Auflistung
  - Auflistung
  - Auflistung

Zweiter Abschnitt  
○○○○

Farben  
○

Bei Frames ohne Titel wird die Kopfzeile nicht angezeigt, und der freie Platz kann für Inhalte genutzt werden.

Bei Frames mit Option `[plain]` werden weder Kopf- noch Fußzeile angezeigt.

# Beispielinhalt

Bei Frames mit Option [t] werden die Inhalte nicht vertikal zentriert, sondern an der Oberkante begonnen.

# Beispielinhalt: Literatur

# Farbpalette

