Karlsruher Institut für Technologie

Bitte Logo über `\grouplogo{}` festlegen.

Corten: Refinement Types for Imperative Languages with Ownership

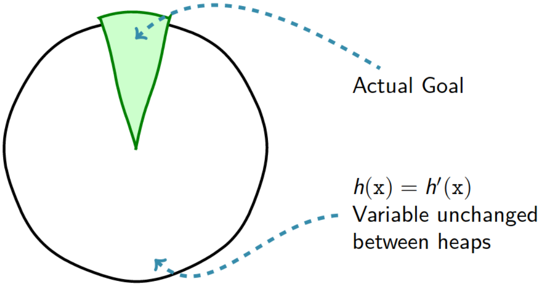
Abschlusspräsentation Masterarbeit

Carsten Csiky | 26th Oktober 2022

Inhaltsverzeichnis

Motivation

```
public IntList square(IntList list) {
  return list.map(x -> x*x);
}
```



Motivation

```rust
fn max(a: i32, b: i32) {
  if a > b { a } else { b }
}
```

Motivation

```rust
fn max(a: i32, b: i32) {
  if a > b { a } else { b }
}
```

- Return Value $(v) : v \geq a \land v \geq b$

Motivation

```rust
fn max(a: i32, b: i32) {
  if a > b { a } else { b }
}
```

- Return Value (v) : $v \geq a \wedge v \geq b$
- Rondon et al. [RKJ08]: Refinement Types for Functional Programming Languages

Motivation

```rust
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

Motivation

```
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

let $\Gamma = (a : \{v : \texttt{i32} \mid \text{true}\}, b : \{v : \texttt{i32} \mid \text{true}\})$ and $\tau = \{v : \texttt{i32} \mid v \geq a \land v \geq b\}$

$$\frac{}{\Gamma \vdash \texttt{if}\ a > b\ \{a\}\ \texttt{else}\ \{b\} : \tau}$$

Motivation

```
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

let $\Gamma = (a : \{v : \texttt{i32} \mid \text{true}\}, b : \{v : \texttt{i32} \mid \text{true}\})$ and $\tau = \{v : \texttt{i32} \mid v \geq a \land v \geq b\}$

$$\dfrac{\dfrac{}{\Gamma, a > b \vdash a : \tau} \qquad \dfrac{}{\Gamma, \neg(a > b) \vdash b : \tau}}{\Gamma \vdash \texttt{if } a > b \; \{a\} \texttt{ else } \{b\} : \tau}$$

Motivation

```
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

let $\Gamma = (a : \{v : \texttt{i32} \mid \text{true}\}, b : \{v : \texttt{i32} \mid \text{true}\})$ and $\tau = \{v : \texttt{i32} \mid v \geq a \land v \geq b\}$

$$\dfrac{\dfrac{\dfrac{}{\Gamma, a > b \vdash \{v : \texttt{i32} \mid v \doteq a\} \preceq \tau}}{\Gamma, a > b \vdash a : \tau} \qquad \dfrac{}{\Gamma, \neg(a > b) \vdash b : \tau}}{\Gamma \vdash \texttt{if}\ a > b\ \{a\}\ \texttt{else}\ \{b\} : \tau}$$

Motivation

```
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

let $\Gamma = (a : \{v : \texttt{i32} \mid \text{true}\}, b : \{v : \texttt{i32} \mid \text{true}\})$ and $\tau = \{v : \texttt{i32} \mid v \geq a \land v \geq b\}$

$$\dfrac{\dfrac{\dfrac{\star}{\Gamma, a > b \vdash a : \{v : \texttt{i32} \mid v \doteq a\}} \qquad \dfrac{}{\Gamma, a > b \vdash \{v : \texttt{i32} \mid v \doteq a\} \preceq \tau}}{\Gamma, a > b \vdash a : \tau} \qquad \dfrac{}{\Gamma, \neg(a > b) \vdash b : \tau}}{\Gamma \vdash \texttt{if}\ a > b\ \{a\}\ \texttt{else}\ \{b\} : \tau}$$

Department of Informatics – Institute of Information Security and Dependability (KASTEL)

Motivation

```
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

let $\Gamma = (a : \{v : \texttt{i32} \mid \text{true}\}, b : \{v : \texttt{i32} \mid \text{true}\})$ and $\tau = \{v : \texttt{i32} \mid v \geq a \wedge v \geq b\}$

$$\dfrac{\dfrac{\dfrac{\star}{\Gamma, a > b \vdash a : \{v : \texttt{i32} \mid v \doteq a\}} \qquad \dfrac{\text{SMT-VALID}\begin{pmatrix} \text{true} \wedge \text{true} \wedge a > b \\ \wedge\, v \doteq a \\ \Longrightarrow (v \geq a \wedge v \geq b) \end{pmatrix}}{\Gamma, a > b \vdash \{v : \texttt{i32} \mid v \doteq a\} \preceq \tau}}{\Gamma, a > b \vdash a : \tau} \qquad \dfrac{}{\Gamma, \neg(a > b) \vdash b : \tau}}{\Gamma \vdash \texttt{if } a > b\ \{a\}\ \texttt{else}\ \{b\} : \tau}$$

Motivation

```
//@ max(a: i32, b: i32) -> {v:i32 | v >= a && v >= b }
fn  max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

let $\Gamma = (a : \{v : \texttt{i32} \mid \text{true}\}, b : \{v : \texttt{i32} \mid \text{true}\})$ and $\tau = \{v : \texttt{i32} \mid v \geq a \land v \geq b\}$

$$
\dfrac{
\dfrac{
\dfrac{\star}{\Gamma, a > b \vdash a : \{v : \texttt{i32} \mid v \doteq a\}}
\qquad
\dfrac{\text{SMT-VALID}\left(\begin{array}{l}\text{true} \land \text{true} \land a > b \\ \land\, v \doteq a \\ \quad \Longrightarrow (v \geq a \land v \geq b)\end{array}\right)}{\Gamma, a > b \vdash \{v : \texttt{i32} \mid v \doteq a\} \preceq \tau}
}{\Gamma, a > b \vdash a : \tau}
\qquad
\dfrac{\vdots}{\Gamma, \neg(a > b) \vdash b : \tau}
}{\Gamma \vdash \texttt{if } a > b\ \{a\}\ \texttt{else}\ \{b\} : \tau}
$$

Motivation

```rust
fn  clamp(a: &mut i32, b: i32) {
  if *a > b { *a = b }
}
```

Motivation

Karlsruher Institut für Technologie

```
fn  clamp(a: &mut i32, b: i32) {
  if *a > b { *a = b }
}
fn  client(...) {
  ...
  clamp(&mut x, 5);
  clamp(&mut y, 6);
  print!(x);
  ...
}
```

Motivation

```
fn  clamp(a: &mut i32, b: i32) {
  if *a > b { *a = b }
}
fn  client(...) {
  ...
  clamp(&mut x, 5);
  clamp(&mut y, 6);
  print!(x);
  ...
}
```

What does this it `print(x)` output?

- In most imperative programming languages:
  - Could be: old `x` or `5`

Motivation

```rust
fn clamp(a: &mut i32, b: i32) {
  if *a > b { *a = b }
}
fn client(...) {
  ...
  clamp(&mut x, 5);
  clamp(&mut y, 6);
  print!(x);
  ...
}
```

What does this it `print(x)` output?

- In most imperative programming languages:
  - Could be: old `x` or `5`
  - But also `6` (if x aliases with y)!

Motivation

```
fn  clamp(a: &mut i32, b: i32) {
  if *a > b { *a = b }
}
fn  client(...) {
  ...
  clamp(&mut x, 5);
  clamp(&mut y, 6);
  print!(x);
  ...
}
```

What does this it `print(x)` output?

- In most imperative programming languages:
  - Could be: old `x` or `5`
  - But also `6` (if x aliases with y)!
- In Rust:
  - Just old `x` or `5`
  - And nothing else!

Motivation

```
fn  clamp(a: &mut i32, b: i32) {
  // borrows a
  // owns b
  if *a > b { *a = b }
  // "returns" the borrow of a
}
fn  client(...) { // owns x, y
  ...
  clamp(&mut x, 5); // lend x mutably
  clamp(&mut y, 6); // lend y mutably
  print!(x);
  ...
}
```

Ownership in Rust: Mutability XOR Aliasing

Each lexical scope tracks permissions for visible memory objects. Possible Permission Levels:

- Owner (e.g. `b`)
  - can: read, write
  - transfer ownership (if no outstanding borrows)
- Mutable Reference (e.g. `&mut x`)
  - can: read, write
  - guarantee: no aliasing
- Immutable Reference (e.g. `&v`)
  - can: read, alias
  - guarantee: no mutation

Motivation

Karlsruher Institut für Technologie

Consequences:

- unique data owner
- no global, mutable state
- no cycles in memory structure

Ownership in Rust: Mutability XOR Aliasing

Each lexical scope tracks permissions for visible memory objects. Possible Permission Levels:

- Owner (e.g. `b`)
  - can: read, write
  - transfer ownership (if no outstanding borrows)
- Mutable Reference (e.g. `&mut x`)
  - can: read, write
  - guarantee: no aliasing
- Immutable Reference (e.g. `&v`)
  - can: read, alias
  - guarantee: no mutation

Motivation

Consequences:

- unique data owner
- no global, mutable state
- no cycles in memory structure

Used for:

- safe non-gc memory management
- safe concurrency
- safe low-level hardware access
- ...

Ownership in Rust: Mutability XOR Aliasing

Each lexical scope tracks permissions for visible memory objects. Possible Permission Levels:

- Owner (e.g. `b`)
  - can: read, write
  - transfer ownership (if no outstanding borrows)
- Mutable Reference (e.g. `&mut x`)
  - can: read, write
  - guarantee: no aliasing
- Immutable Reference (e.g. `&v`)
  - can: read, alias
  - guarantee: no mutation

Motivation

Consequences:

- unique data owner
- no global, mutable state
- no cycles in memory structure

Used for:

- safe non-gc memory management
- safe concurrency
- safe low-level hardware access
- ...
- ⇒ show: program verification as well

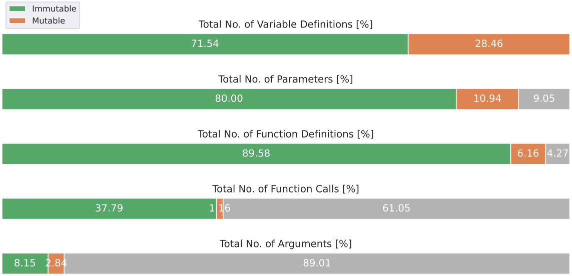
Ownership in Rust: Mutability XOR Aliasing

Each lexical scope tracks permissions for visible memory objects. Possible Permission Levels:

- Owner (e.g. `b`)
  - can: read, write
  - transfer ownership (if no outstanding borrows)
- Mutable Reference (e.g. `&mut x`)
  - can: read, write
  - guarantee: no aliasing
- Immutable Reference (e.g. `&v`)
  - can: read, alias
  - guarantee: no mutation

Empirical Use-Case Analysis

- public open-source code (crates.io)
- about 64 million lines of Rust code
- syntactical analysis



Refinement Types for Rust – Syntax

```
fn max(a: i32, b: i32) -> i32 {
  if a > b { a } else { b }
}
```

Addition of two macros

- `ty!`$\{l : b \mid \varphi\}$ in place of a type
- `relax_ctx!{ ... }` in place of a statement

Refinement Types for Rust – Syntax

```
fn max(
  a: ty!{ av: i32 | true },
  b: ty!{ bv : i32 | true }
) -> ty!{ v : i32 | v >= av && v >= bv } {
  if a > b { a } else { b }
}
```

Addition of two macros

- `ty!`$\{l : b \mid \varphi\}$ in place of a type
- `relax_ctx!{ ... }` in place of a statement

Refinement Types for Rust – Syntax

```
fn max(
  a: ty!{ av: i32 },
  b: ty!{ bv : i32 }
) -> ty!{ v : i32 | v >= av && v >= bv } {
  if a > b { a } else { b }
}
```

Addition of two macros

- `ty!`$\{l : b \mid \varphi\}$ in place of a type
- `relax_ctx!{ ... }` in place of a statement

Type Updates

```
fn decr() -> ty!{ w : i32 | w >= 0 } {
 let mut i = ... as ty!{ v: i32 | v > 0};
 i = i - 1;
 i
}
```

- Types need to change through execution
  - $\Rightarrow$ type updates
  - Separation of program-variables and logic-variables
  - Γ association of program- to logic-variables and predicate
  - $\Gamma \vdash s \Rightarrow \Gamma'$ (Statement Type Checking)
  - $\Gamma \vdash e : \tau$ (Expression Typing)
  - on assignment: *replace* association, *append* predicate
  - observation: assignment can not invalidate existing predicates

Type Updates

```
fn decr() -> ty!{ w : i32 | w >= 0 } {
  // Γ1 = ({}, true)
  let mut i = ... as ty!{ v: i32 | v > 0};
  // Γ2 = ({i ↦ v}, v > 0)
  i = i - 1;
  // Γ3 = ({i ↦ v2}, v > 0 ∧ v2 ≐ v − 1)
  i }
```

- Types need to change through execution
  - $\Rightarrow$ type updates
  - Separation of program-variables and logic-variables
  - Γ association of program- to logic-variables and predicate
  - $\Gamma \vdash s \Rightarrow \Gamma'$ (Statement Type Checking)
  - $\Gamma \vdash e : \tau$ (Expression Typing)
  - on assignment: *replace* association, *append* predicate
  - observation: assignment can not invalidate existing predicates

Type Updates

```
fn decr() -> ty!{ w : i32 | w >= 0 } {
  // Γ₁ = ({}, true)
  let mut i = ... as ty!{ v: i32 | v > 0};
  // Γ₂ = ({i ↦ v}, v > 0)
  i = i - 1;
  // Γ₃ = ({i ↦ v₂}, v > 0 ∧ v₂ ≐ v − 1)
  i }
```

$$\textsc{Intro-Sub}\ \frac{\Gamma \vdash e : \tau \qquad \Gamma \vdash \tau \preceq \tau'}{\Gamma \vdash e \ \texttt{as}\ \tau' : \tau'}$$

$$\textsc{Decl}\ \frac{\Gamma \vdash e : \{\beta : b \mid \varphi\}}{\Gamma \vdash \texttt{let}\ x = e \Rightarrow \Gamma[x \mapsto \beta], \varphi}$$

Type Updates

```
fn decr() -> ty!{ w : i32 | w >= 0 } {
  // Γ₁ = ({}, true)
  let mut i = ... as ty!{ v: i32 | v > 0};
  // Γ₂ = ({i ↦ v}, v > 0)
  i = i - 1;
  // Γ₃ = ({i ↦ v₂}, v > 0 ∧ v₂ ≐ v − 1)
  i }
```

$$\textsc{BinOp}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash x_1 \odot x_2 : \{\alpha : b \mid \alpha \simeq [\![x_1 \odot x_2]\!]\Gamma\}}$$

$$\textsc{Assign}\ \frac{\Gamma \vdash e : \{\beta : b \mid \varphi\}}{\Gamma \vdash x = e \Rightarrow \Gamma[x \mapsto \beta], \varphi}$$

Type Updates

```
fn decr() -> ty!{ w : i32 | w >= 0 } {
  // Γ1 = ({}, true)
  let mut i = ... as ty!{ v: i32 | v > 0};
  // Γ2 = ({i ↦ v}, v > 0)
  i = i - 1;
  // Γ3 = ({i ↦ v2}, v > 0 ∧ v2 ≐ v − 1)
  i }
```

$$\textsc{Seq}\ \frac{\Gamma \vdash s_1 \Rightarrow \Gamma' \qquad \Gamma' \vdash s_2 \Rightarrow \Gamma''}{\Gamma \vdash s_1; s_2 \Rightarrow \Gamma''}$$

References – Strong Updates

```
fn client() -> ty!{ v: i32 | v == 4 } {
  let a = 2;       // a : {v1 : i32 | v1 == 2}
  let b = &mut a;  // b : {v2 : &i32 | v2 == &a}
  *b = 0; // changes a's value and type
  let c = &mut b;  // c : {v3 : &i32 | v3 == &b}
  **c = 4; // changes a's value and type
  a
}
```

References – Strong Updates

```
fn client() -> ty!{ v: i32 | v == 4 } {
  let a = 2;        // a : {v₁ : i32 | v₁ == 2}
  let b = &mut a;   // b : {v₂ : &i32 | v₂ == &a}
  *b = 0; // changes a's value and type
  let c = &mut b;   // c : {v₃ : &i32 | v₃ == &b}
  **c = 4; // changes a's value and type
  a
}
```

$$\text{LIT}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash v : \{\alpha : b \mid \alpha \simeq [\![v]\!]\Gamma\}}$$

References – Strong Updates

```
fn client() -> ty!{ v: i32 | v == 4 } {
  let a = 2;        // a : {v₁ : i32 | v₁ == 2}
  let b = &mut a;   // b : {v₂ : &i32 | v₂ == &a}
  *b = 0; // changes a's value and type
  let c = &mut b;   // c : {v₃ : &i32 | v₃ == &b}
  **c = 4; // changes a's value and type
  a
}
```

$$\textsc{Ref}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash \&x : \{\alpha : \&b \mid \alpha \simeq [\![\&x]\!]\Gamma\}}$$

References – Strong Updates

```
fn client() -> ty!{ v: i32 | v == 4 } {
  let a = 2;        // a : {v1 : i32 | v1 == 2}
  let b = &mut a;   // b : {v2 : &i32 | v2 == &a}
  *b = 0; // changes a's value and type
  let c = &mut b;   // c : {v3 : &i32 | v3 == &b}
  **c = 4; // changes a's value and type
  a
}
```

$$\textsc{Assign-Strong}\ \frac{\Gamma(z) = \beta \quad \Gamma \vdash x \in \{\&y\} \qquad \Gamma \vdash \gamma \text{ fresh}}{\Gamma \vdash *x = z \Rightarrow \Gamma[y \mapsto \gamma], \gamma \doteq \beta}$$

(Also Assign-Weak)

References – Strong Updates

Karlsruher Institut für Technologie

```
fn client() -> ty!{ v: i32 | v == 4 } {
  let a = 2;        // a : {v₁ : i32 | v₁ == 2}
  let b = &mut a;   // b : {v₂ : &i32 | v₂ == &a}
  *b = 0; // changes a's value and type
  let c = &mut b;   // c : {v₃ : &i32 | v₃ == &b}
  **c = 4; // changes a's value and type
  a
}
```

$$\textsc{Var}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash x : \{\alpha : b \mid \alpha \simeq [\![x]\!]\Gamma\}}$$

Mutable Arguments


```
fn clamp(a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 }, b: ty!{ b1: i32 }) {
  if *a > b { *a = b }
}
fn client(...) {
  ...
  clamp(&mut x, 5);
  clamp(&mut y, 6);
  print!(x);
  ...
}
```

Mutable Arguments

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //                v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //          v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

Mutable Arguments

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //                 v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //          v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

- `ty!` $\{\alpha : \mathrm{b} \mid \varphi \Rightarrow \beta \mid \psi\}$
- Callee requires φ for reference destination α
- Callee ensures ψ for reference destination β
- Of course, multiple arguments possible

Mutable Arguments

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //                v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //           v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

Mutable Arguments

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //                 v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //             v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

Sub-Context

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ₁ = ({a ↦ v₁, arg₀ ↦ a₁, b ↦ b₁},
  //                v₁ ≐ &arg₀ ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ₂ = ({a ↦ v₁, arg₀ ↦ v₂, b ↦ b₁},
  //         v₂ ≤ b₁ ∧ v₁ ≐ &arg₀ ∧ true ∧ true)
}
```

- still left: proof obligation from signature $a_2 \leq b_1$
- i.e. is Γ_2 a valid end-state?

Sub-Context

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //                 v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //           v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

- still left: proof obligation from signature $a_2 \leq b_1$
- i.e. is Γ_2 a valid end-state?
- generalize notion of sub-types to context: sub-context

Sub-Context

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //              v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //         v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

- still left: proof obligation from signature $a_2 \leq b_1$
- i.e. is Γ_2 a valid end-state?
- generalize notion of sub-types to context: sub-context
- expected state: $\Gamma_e = (\{arg_0 \mapsto a_2, b \mapsto b_1\}, a_2 \leq b_1)$

Sub-Context

```
fn clamp(
  a: &mut ty!{ a1 : i32 | true => a2 | a2 <= b1 },
  b: ty!{ b1: i32 }
) {
  // Γ1 = ({a ↦ v1, arg0 ↦ a1, b ↦ b1},
  //            v1 ≐ &arg0 ∧ true ∧ true)
  if *a > b { *a = b }
  // Γ2 = ({a ↦ v1, arg0 ↦ v2, b ↦ b1},
  //       v2 ≤ b1 ∧ v1 ≐ &arg0 ∧ true ∧ true)
}
```

- still left: proof obligation from signature $a_2 \leq b_1$
- i.e. is Γ_2 a valid end-state?
- generalize notion of sub-types to context: sub-context
- expected state: $\Gamma_e = (\{arg_0 \mapsto a_2, b \mapsto b_1\}, a_2 \leq b_1)$
- show: $\Gamma_2 \preceq \Gamma_e$

Sub-Context

$$\preceq\text{-CTX}\ \frac{\vDash \Phi'[\mu'(x) \triangleright \mu(x) \ \mid x \in \mathrm{dom}(\mu')] \to \Phi \qquad \mathrm{dom}(\mu') \subseteq \mathrm{dom}(\mu)}{(\mu, \Phi) \preceq (\mu', \Phi')}$$

- still left: proof obligation from signature $a_2 \leq b_1$
- i.e. is Γ_2 a valid end-state?
- generalize notion of sub-types to context: sub-context
- expected state: $\Gamma_e = (\{arg_0 \mapsto a_2, b \mapsto b_1\}, a_2 \leq b_1)$
- show: $\Gamma_2 \preceq \Gamma_e$

Mutable Calls

```
fn client(...) {
  ...
  // Γ1 = ({x ↦ v1, y ↦ v2}, ...)
  clamp(&mut x, 5);
  // Γ2 = ({x ↦ v3, y ↦ v2}, ... ∧ v3 ≤ 5)
  clamp(&mut y, 6);
  // Γ3 = ({x ↦ v3, y ↦ v4}, ... ∧ v3 ≤ 5 ∧ v4 ≤ 6)
  print!(x);
  ...
}
```

- append predicates from callee to context
- update association of logic variables

SMT Request

...

Example Error Message

...

Ecosystem Integration

...

Soundness

Karlsruher Institut für Technologie

Progress

If $\Gamma \vdash s_1, \sigma : \Gamma \Rightarrow \Gamma_2$ and $s_1 \neq \texttt{unit}$, then there is a s_2 and σ_2 with $\langle s_1 \mid \sigma_1 \rangle \rightsquigarrow \langle s_2 \mid \sigma_2 \rangle$.

Corten strictly refines the base language, therefore progress depends on base type system.

Preservation

If $\Gamma \vdash s \Rightarrow \Gamma_2, \sigma : \Gamma$ and $\langle s \mid \sigma \rangle \rightsquigarrow \langle s_1 \mid \sigma_1 \rangle$, then there is a Γ_1 with $\Gamma_1 \vdash s_1 \Rightarrow \Gamma_2$ and $\sigma_2 : \Gamma_2$

Stronger property than base language preservation: Show that refined types are preserved

State Conformance

Karlsruher Institut für Technologie

State Conformance $\sigma : \Gamma$

A state σ is conformant with respect to a typing context $\Gamma = (\mu, \Phi)$ (written as $\sigma : \Gamma$), iff:

$$\Phi[\mu(x) \triangleright [\![\sigma(x)]\!] \mid x \in \mathrm{dom}(\mu)] \text{ is satisfiable}$$

I.e. a conformant type context does not contradict the execution state.

Examples:

- If $\sigma : (\emptyset, \Phi)$ then Φ is satisfiable
- If $\sigma : (\mu, \Phi_1 \wedge \Phi_2)$ then $\sigma : (\mu, \Phi_1)$ and $\sigma : (\mu, \Phi_1)$.
- If $\sigma : (\mu, \Phi)$ and $\mathrm{FV}(\Phi) \subseteq \mathrm{dom}(\mu)$, then $\vDash \Phi[\mu(x) \triangleright [\![\sigma(x)]\!] \mid x \in \mathrm{dom}(\mu)]$

Intermediate Steps

Karlsruher Institut für Technologie

Conformance of Symbolic Execution

If $\sigma : \Gamma, \Gamma \vdash \alpha$ fresh then $\sigma[x \mapsto [\![e]\!]\sigma] : \Gamma[x \mapsto \alpha], (\alpha \simeq [\![e]\!]\Gamma)$

where $(\alpha \simeq [\![e]\!]\Gamma)$ is the symbolic execution of e equated with α in context Γ

Reference Predicates are Conservative

If $\sigma : \Gamma$ and $\Gamma \vdash *x \in \{y_1, \ldots, y_n\}$ then $[\![\sigma(x)]\!] = \&y_i$ for some $i \in 1, \ldots, n$

Rare case where conservative typing requires

Sub-Context Relation is Conservative

If $\Gamma \preceq \Gamma'$ and $\sigma : \Gamma$ then $\sigma : \Gamma'$

Related Work

Refinement Types and Mutability

- Rondon et al. [RKJ10], Bakst and Jhala [BJ16]: Refinement Types for C subset. Lack of guarantees requires ad-hoc mechanisms to control aliasing
- Lanzinger [Lan21]: Property Types in Java (only immutable)
- Bachmeier [Bac22]: Extension using Ownership System

Rust verification

- Ullrich [Ull]: Translation to Lean; linear mutation chain. Denis et al [DJM21] similar, but to Why3
- Astrauskas et al. [Ast+19] (Prusti): heavy-weight verification, translation to separation logic (Viper)
- Matsushita et al. [MTK20]: constrained Horn clauses (RustHorn)

Flux – Refinement Types for Rust

- MIR vs. HIR
- specification in comments vs. embedding in types
- context inclusions vs. sub context
- distinction strong and weak references vs. dynamic choice by typ checking rules
- explicit introduction of logic variables vs. ad-hoc
- formalization based on RustBelt vs. formalization based on own language
- missing in Corten: records & inference
- otherwise: similar capabilities

Future Work

- Records & ADTs
- Predicate Generics (Abstract Predicates)
- Concurrency using Predicate Generics?

Department of Informatics – Institute of Information Security and Dependability (KASTEL)

Conclusion

- Working Refinement Type System for Rust with Mutability
- Minimal Interface
- Soundness Justification
- Evaluation

Department of Informatics – Institute of Information Security and Dependability (KASTEL)

Literatur

Karlsruher Institut für Technologie

[1] Vytautas Astrauskas u. a. „Leveraging rust types for modular specification and verification“. In: *Proceedings of the ACM on Programming Languages* 3 (OOPSLA 10. Okt. 2019), S. 1–30. ISSN: 2475-1421. DOI: `10.1145/3360573`. URL: `https://dl.acm.org/doi/10.1145/3360573` (besucht am 23. 02. 2022).

[2] Joshua Bachmeier. *Property Types for Mutable Data Structures in Java*. 2022. DOI: `10.5445/IR/1000150318`. URL: `https://publikationen.bibliothek.kit.edu/1000150318` (besucht am 03. 10. 2022).

[3] Alexander Bakst und Ranjit Jhala. „Predicate Abstraction for Linked Data Structures“. In: *Verification, Model Checking, and Abstract Interpretation*. Hrsg. von Barbara Jobstmann und K. Rustan M. Leino. Lecture Notes in Computer Science. Berlin, Heidelberg: Springer, 2016, S. 65–84. ISBN: 978-3-662-49122-5. DOI: `10.1007/978-3-662-49122-5_3`.

[4] Xavier Denis, Jacques-Henri Jourdan und Claude Marché. „The Creusot Environment for the Deductive Verification of Rust Programs“. Diss. Inria Saclay-Île de France, 2021.

[5] Florian Lanzinger. „Property Types in Java: Combining Type Systems and Deductive Verification“. Master Thesis. Karlsruher Institut für Technologie, Feb. 2021.

[6] Yusuke Matsushita, Takeshi Tsukada und Naoki Kobayashi. „RustHorn: CHC-based verification for Rust programs“. In: *European Symposium on Programming*. Springer, Cham, 2020, S. 484–514.

$\texttt{decr}$ Typing Tree

let $\Gamma_2 = \Gamma[i \mapsto v_1], v > 0$ and $\tau = \{v : \texttt{i32} \mid v > 0\}$

$$
\dfrac{
  \textsc{Decl}\;\dfrac{
    \textsc{Intro-Sub}\;\dfrac{\Gamma_1 \vdash ... : \tau' \qquad \Gamma_1 \vdash \tau' \preceq \tau}{\Gamma_1 \vdash \texttt{...\ as}\ \tau : \tau}
  }{\Gamma_1 \vdash \texttt{let i = ... as}\ \tau \Rightarrow \Gamma_2}
  \qquad
  \textsc{Ass}\;\dfrac{
    \textsc{BinOp}\;\dfrac{\Gamma_1 \vdash v_2 \text{ fresh}}{\Gamma_1 \vdash \texttt{i - 1} : \{v_2 : \text{i32} \mid v_2 \doteq v - 1\}}
  }{\Gamma_2 \vdash \texttt{i = i - 1} \Rightarrow \Gamma[i \mapsto v_2], v > 0, v_2 \doteq v - 1}
}{\Gamma_1 \vdash \texttt{let i = ... as}\ \tau\texttt{; i = i - 1} \Rightarrow \Gamma[i \mapsto v_2], v > 0, v_2 \doteq v - 1}\;\textsc{Seq}
$$

Expression Typing $\Gamma \vdash e : \tau$

$$\textsc{Lit}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash v : \{\alpha : b \mid \alpha \simeq [\![v]\!]\Gamma\}} \qquad \textsc{BinOp}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash x_1 \odot x_2 : \{\alpha : b \mid \alpha \simeq [\![x_1 \odot x_2]\!]\Gamma\}}$$

$$\textsc{Var}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash x : \{\alpha : b \mid \alpha \simeq [\![x]\!]\Gamma\}} \qquad \textsc{Intro-Sub}\ \frac{\Gamma \vdash e : \tau \qquad \Gamma \vdash \tau \preceq \tau'}{\Gamma \vdash e\ \texttt{as}\ \tau' : \tau'}$$

Statement Type Checking $\Gamma \vdash s \Rightarrow \Gamma'$

$$\textsc{If}\ \frac{\Gamma, \Gamma(x) \doteq \texttt{true} \vdash s_t \Rightarrow \Gamma' \qquad \Gamma, \Gamma(x) \doteq \texttt{false} \vdash s_e \Rightarrow \Gamma'}{\Gamma \vdash \text{if } x \text{ then } s_t \text{ else } s_e \Rightarrow \Gamma'}$$

$$\textsc{Seq}\ \frac{\Gamma \vdash s_1 \Rightarrow \Gamma' \qquad \Gamma' \vdash s_2 \Rightarrow \Gamma''}{\Gamma \vdash s_1 ; s_2 \Rightarrow \Gamma''}$$

$$\textsc{Decl}\ \frac{\Gamma \vdash e : \{\beta : b \mid \varphi\}}{\Gamma \vdash \texttt{let}\ x = e \Rightarrow \Gamma[x \mapsto \beta], \varphi} \qquad \textsc{Assign}\ \frac{\Gamma \vdash e : \{\beta : b \mid \varphi\}}{\Gamma \vdash x = e \Rightarrow \Gamma[x \mapsto \beta], \varphi}$$

Expression Typing $\Gamma \vdash e : \tau$

$$\textsc{Ref}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash \&x : \{\alpha : \&b \mid \alpha \simeq [\![\&x]\!]\Gamma\}}$$

$$\textsc{Var-Deref}\ \frac{\Gamma \vdash x \in \{\&y\} \qquad \Gamma \vdash y : \tau}{\Gamma \vdash *x : \tau}$$

Statement Type Checking $\Gamma \vdash s \Rightarrow \Gamma'$

$$\textsc{Assign-Strong}\ \frac{\Gamma(z) = \beta \qquad \Gamma \vdash x \in \{\&y\} \qquad \Gamma \vdash \gamma \text{ fresh}}{\Gamma \vdash *x = z \Rightarrow \Gamma[y \mapsto \gamma], \gamma \doteq \beta}$$

Expression Typing $\Gamma \vdash e : \tau$

$$\textsc{Ref}\ \frac{\Gamma \vdash \alpha \text{ fresh}}{\Gamma \vdash \&x : \{\alpha : \&b \mid \alpha \simeq [\![\&x]\!]\Gamma\}}$$

$$\textsc{Var-Deref}\ \frac{\Gamma \vdash x \in \{\&y\} \qquad \Gamma \vdash y : \tau}{\Gamma \vdash *x : \tau}$$

Statement Type Checking $\Gamma \vdash s \Rightarrow \Gamma'$

$$\textsc{Assign-Strong}\ \frac{\Gamma(z) = \beta \qquad \Gamma \vdash x \in \{\&y\} \qquad \Gamma \vdash \gamma \text{ fresh}}{\Gamma \vdash *x = z \Rightarrow \Gamma[y \mapsto \gamma], \gamma \doteq \beta}$$

$$\textsc{Assign-Weak}\ \frac{\dfrac{\Gamma \vdash e : \tau}{\Gamma \vdash y_i : \{\beta_i : b_i \mid \varphi_i\}} \qquad \dfrac{\Gamma \vdash x \in \{\&y_1, \dots, \&y_n\}}{\Gamma \vdash \tau \preceq \{\beta_i : b_i \mid \varphi_i\}}}{\Gamma \vdash *x = e \Rightarrow \Gamma}$$

Blöcke
in den KIT-Farben

Greenblock
Standard (`block`)

Blueblock
= `exampleblock`

Redblock
= `alertblock`

Brownblock

Purpleblock

Cyanblock

Yellowblock

Lightgreenblock

Orangeblock

Grayblock

Contentblock
(farblos)

Auflistungen

Text

- Auflistung
  Umbruch
- Auflistung
  - Auflistung
  - Auflistung

Bei Frames ohne Titel wird die Kopfzeile nicht angezeigt, und der freie Platz kann für Inhalte genutzt werden.

Bei Frames mit Option `[plain]` werden weder Kopf- noch Fußzeile angezeigt.

Beispielinhalt

Karlsruher Institut für Technologie

Bei Frames mit Option `[t]` werden die Inhalte nicht vertikal zentriert, sondern an der Oberkante begonnen.

Beispielinhalt: Literatur

Farbpalette

Karlsruher Institut für Technologie

kit-green100 kit-green90 kit-green80 kit-green70 kit-green60 kit-green50 kit-green40 kit-green30 kit-green25 kit-green20 kit-green15 kit-green10 kit-green5

kit-blue100 kit-blue90 kit-blue80 kit-blue70 kit-blue60 kit-blue50 kit-blue40 kit-blue30 kit-blue25 kit-blue20 kit-blue15 kit-blue10 kit-blue5

kit-red100 kit-red90 kit-red80 kit-red70 kit-red60 kit-red50 kit-red40 kit-red30 kit-red25 kit-red20 kit-red15 kit-red10 kit-red5

kit-gray100 kit-gray90 kit-gray80 kit-gray70 kit-gray60 kit-gray50 kit-gray40 kit-gray30 kit-gray25 kit-gray20 kit-gray15 kit-gray10 kit-gray5

kit-orange100 kit-orange90 kit-orange80 kit-orange70 kit-orange60 kit-orange50 kit-orange40 kit-orange30 kit-orange25 kit-orange20 kit-orange15 kit-orange10 kit-orange5

kit-lightgreen100 kit-lightgreen90 kit-lightgreen80 kit-lightgreen70 kit-lightgreen60 kit-lightgreen50 kit-lightgreen40 kit-lightgreen30 kit-lightgreen25 kit-lightgreen20 kit-lightgreen15 kit-lightgreen10 kit-lightgreen5

kit-brown100 kit-brown90 kit-brown80 kit-brown70 kit-brown60 kit-brown50 kit-brown40 kit-brown30 kit-brown25 kit-brown20 kit-brown15 kit-brown10 kit-brown5

kit-purple100 kit-purple90 kit-purple80 kit-purple70 kit-purple60 kit-purple50 kit-purple40 kit-purple30 kit-purple25 kit-purple20 kit-purple15 kit-purple10 kit-purple5

kit-cyan100 kit-cyan90 kit-cyan80 kit-cyan70 kit-cyan60 kit-cyan50 kit-cyan40 kit-cyan30 kit-cyan25 kit-cyan20 kit-cyan15 kit-cyan10 kit-cyan5